# **Binary Search**

**Shusen Wang** 

$$\operatorname{arr} = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \begin{bmatrix} 5 \\ 1 \end{bmatrix} \begin{bmatrix} 12 \\ 2 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \end{bmatrix} \begin{bmatrix} 17 \\ 4 \end{bmatrix} \begin{bmatrix} 26 \\ 5 \end{bmatrix} \begin{bmatrix} 32 \\ 6 \end{bmatrix} \begin{bmatrix} 51 \\ 7 \end{bmatrix} \begin{bmatrix} 53 \\ 8 \end{bmatrix} \begin{bmatrix} 64 \\ 9 \end{bmatrix}$$

$$arr = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \begin{bmatrix} 5 \\ 1 \end{bmatrix} \begin{bmatrix} 12 \\ 2 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \end{bmatrix} \begin{bmatrix} 17 \\ 4 \end{bmatrix} \begin{bmatrix} 26 \\ 5 \end{bmatrix} \begin{bmatrix} 32 \\ 6 \end{bmatrix} \begin{bmatrix} 51 \\ 7 \end{bmatrix} \begin{bmatrix} 53 \\ 8 \end{bmatrix} \begin{bmatrix} 64 \\ 9 \end{bmatrix}$$

- Inputs: (i) an array whose elements are in the ascending order and (ii) a key.
- Goal: Search for the key in the array. If found, return its index; if not found, return −1.

$$arr = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \begin{bmatrix} 5 \\ 1 \end{bmatrix} \begin{bmatrix} 12 \\ 2 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \end{bmatrix} \begin{bmatrix} 17 \\ 4 \end{bmatrix} \begin{bmatrix} 26 \\ 5 \end{bmatrix} \begin{bmatrix} 32 \\ 6 \end{bmatrix} \begin{bmatrix} 51 \\ 7 \end{bmatrix} \begin{bmatrix} 53 \\ 8 \end{bmatrix} \begin{bmatrix} 64 \\ 9 \end{bmatrix}$$

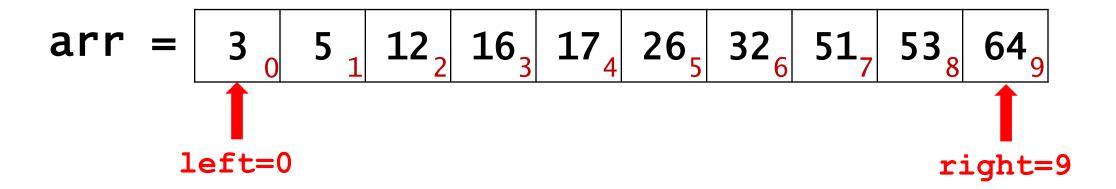
#### **Example 1:**

- Search for the element 53.
- Return 8.

$$arr = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \begin{bmatrix} 5 \\ 1 \end{bmatrix} \begin{bmatrix} 12 \\ 2 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \end{bmatrix} \begin{bmatrix} 17 \\ 4 \end{bmatrix} \begin{bmatrix} 26 \\ 5 \end{bmatrix} \begin{bmatrix} 32 \\ 6 \end{bmatrix} \begin{bmatrix} 51 \\ 7 \end{bmatrix} \begin{bmatrix} 53 \\ 8 \end{bmatrix} \begin{bmatrix} 64 \\ 9 \end{bmatrix}$$

#### **Example 2:**

- Search for the element 9.
- Return -1.



mid ← [(left + right)/2].
 If arr[mid] == key ==> return mid.
 If arr[mid] > key ==> right ← mid-1.
 If arr[mid] < key ==> left ← mid+1.

mid ← [(left + right)/2].
If arr[mid]==key ==> return mid.
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If arr[mid] == key ==> return mid.
If arr[mid] > key ==> right ← mid-1.
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```

```
arr = 3 0 5 1 12 16 17 26 32 51 53 64 9

left=0 mid=4 right=9
```

```
mid ← [(left + right)/2].
If arr[mid] == key ==> return mid.
If arr[mid] > key ==> right ← mid-1.
If arr[mid] < key ==> left ← mid+1.
```

```
arr = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \begin{bmatrix} 5 \\ 1 \end{bmatrix} \begin{bmatrix} 12 \\ 2 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \end{bmatrix} \begin{bmatrix} 17 \\ 4 \end{bmatrix} \begin{bmatrix} 26 \\ 5 \end{bmatrix} \begin{bmatrix} 32 \\ 6 \end{bmatrix} \begin{bmatrix} 51 \\ 7 \end{bmatrix} \begin{bmatrix} 53 \\ 8 \end{bmatrix} \begin{bmatrix} 64 \\ 9 \end{bmatrix}
left=5
```

```
mid ← [(left + right)/2].
If arr[mid] == key ==> return mid.
If arr[mid] > key ==> right ← mid-1.
If arr[mid] < key ==> left ← mid+1.
```

arr = 
$$\begin{bmatrix} 3 \\ 0 \end{bmatrix} \begin{bmatrix} 5 \\ 1 \end{bmatrix} \begin{bmatrix} 12 \\ 2 \end{bmatrix} \begin{bmatrix} 16 \\ 3 \end{bmatrix} \begin{bmatrix} 17 \\ 4 \end{bmatrix} \begin{bmatrix} 26 \\ 5 \end{bmatrix} \begin{bmatrix} 32 \\ 6 \end{bmatrix} \begin{bmatrix} 51 \\ 7 \end{bmatrix} \begin{bmatrix} 53 \\ 8 \end{bmatrix} \begin{bmatrix} 64 \\ 9 \end{bmatrix}$$
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- mid  $\leftarrow \lfloor (\text{left} + \text{right})/2 \rfloor$ .
- If arr[mid] == key ==> return mid.
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```
arr = 3 0 5 1 12 16 17 26 32 51 53 64 9 left=5 right=6
```

- mid  $\leftarrow \lfloor (\text{left} + \text{right})/2 \rfloor$ .
- If arr[mid] == key ==> return mid.
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arr = 3 0 5 1 12 16 17 26 32 51 53 64 9 left=5 right=6
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```
arr = 3 5 12 16 17 26 32 51 53 64 9

left=5

right=6
```

- mid  $\leftarrow \lfloor (\text{left} + \text{right})/2 \rfloor$ .
- If arr[mid] == key ==> return mid.
- If arr[mid]>key ==> right ← mid-1.
- If arr[mid] < key ==> left ← mid+1.

```
arr = \begin{bmatrix} 3 & 5 & 12 & 16 & 17 & 26 & 32 & 51 & 53 & 64 & 9 \end{bmatrix}
```

- mid  $\leftarrow \lfloor (\text{left} + \text{right})/2 \rfloor$ .
- If arr[mid] == key ==> return mid.
- If arr[mid]>key ==> right ← mid-1.
- If arr[mid] < key ==> left ← mid+1.

#### **Time Complexity**

- Each iteration reduces the size of array by half.
- Let *n* be the size of the array.
- The total number of iterations is at most  $\log_2 n$ .
- Per-iteration time complexity: O(1).

Overall time complexity:  $O(\log n)$ .

```
int search(int arr[], int left, int right, int key) {
    while (left <= right) {
          int mid = (left + right) / 2;
          if (key == arr[mid])
               return mid; // key is found
          if (key > arr[mid])
               left = mid + 1; // search in the right half
               right = mid - 1; // search in the left half
     return -1; // key is not found
```

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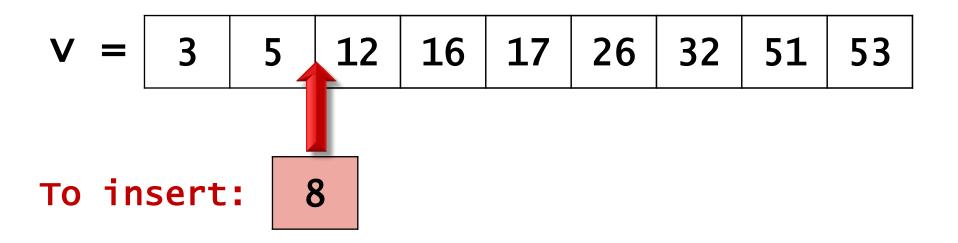
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#### **Vector**

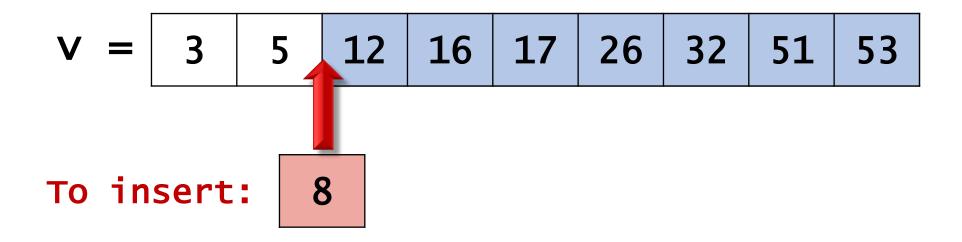
- The ascending order must be kept; otherwise, search would take O(n) time.
- Inserting an item into the middle has O(n) time complexity (on average).

#### **Vector**



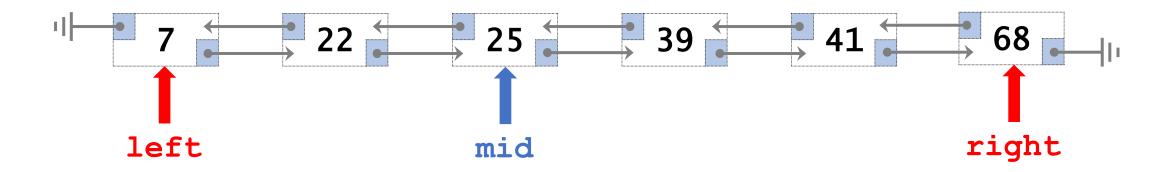
• Inserting an item into the middle has O(n) time complexity (on average).

#### **Vector**



• Inserting an item into the middle has O(n) time complexity (on average).

#### List



- Can we perform binary search in the list?
- No. Given left and right, we cannot get mid efficiently.

## **Comparisons**

**Vector** 

List

SearchInsertion $O(\log n)$ O(n)O(n)O(1)

# **Comparisons**

	Search	Insertion
Vector	$O(\log n)$	O(n)
List	O(n)	0(1)
Skip List	$O(\log n)$	$O(\log n)$

#### Thank You!