Dinic's Algorithm

Shusen Wang

Comparisons

- *m*: the number of edges.
- n: the number of vertices.
- Time complexity of Edmonds–Karp algorithm [1] is $O(m^2 \cdot n)$.

Reference

1. Jack Edmonds and Richard M Karp. Theoretical improvements in algorithmic efficiency for network flow problems. *Journal of the ACM*. 19 (2): 248–264, 1972.

Comparisons

- *m*: the number of edges.
- n: the number of vertices.
- Time complexity of Edmonds–Karp algorithm [1] is $O(m^2 \cdot n)$.
- Time complexity of Dinic's algorithm [2] is $O(m \cdot n^2)$. (Faster, because n is small than m.)

Reference

- 1. Jack Edmonds and Richard M Karp. Theoretical improvements in algorithmic efficiency for network flow problems. *Journal of the ACM*. 19 (2): 248–264, 1972.
- 2. Yefim Dinitz. Algorithm for solution of a problem of maximum flow in a network with power estimation. *Proceedings of the USSR Academy of Sciences*, 11: 1277–1280, 1970.

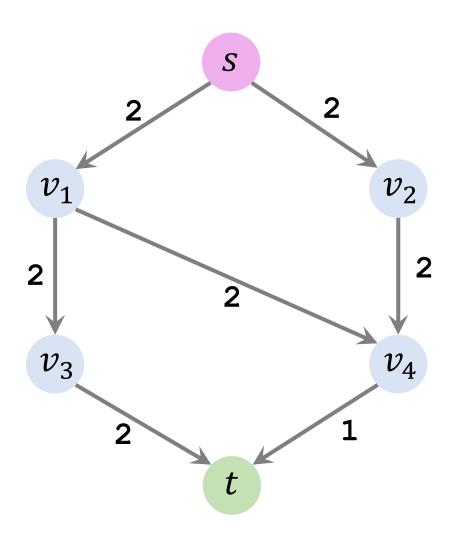
Dinic's Algorithm

- Yefim Dinitz published "Dinitz's Algorithm" in USSR, 1970 [1].
- "Dinitz's Algorithm" was introduced to the westerners by Even & Tarjan's 1975 paper [2].
- Even & Tarjan misspelled <u>Dinitz's</u> name as "<u>Dinic"</u> [2].

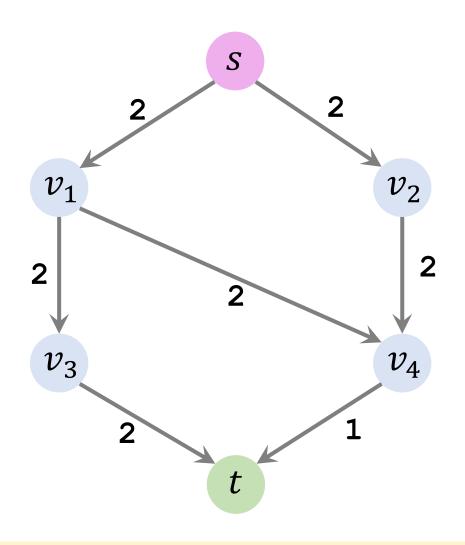
Reference

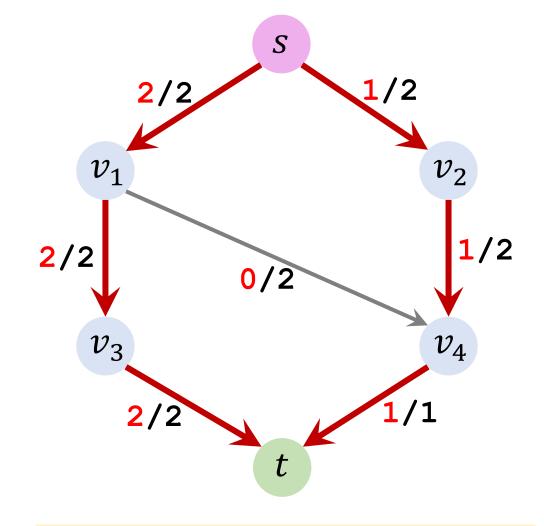
- 1. Yefim Dinitz. Algorithm for solution of a problem of maximum flow in a network with power estimation. *Proceedings of the USSR Academy of Sciences*, 11: 1277–1280, 1970.
- 2. Shimon Even and R. Endre Tarjan. Network Flow and Testing Graph Connectivity. *SIAM Journal on Computing*, 4 (4): 507–518, 1975.

Key Concept: Blocking Flow



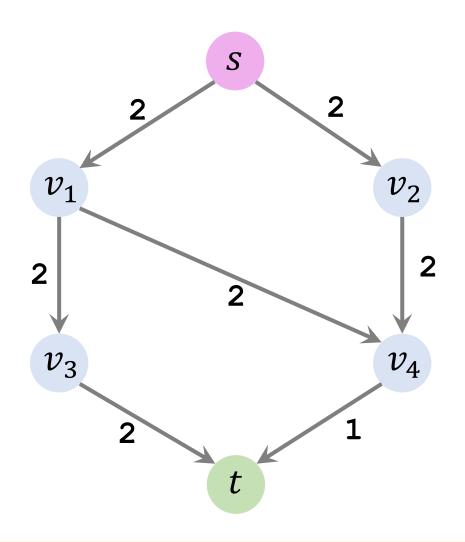
- A flow is blocking flow if no more flow from source to sink can be found.
- Max flow is blocking flow; blocking flow may not be max flow.
- Blocking flow can be found by the naïve algorithm.

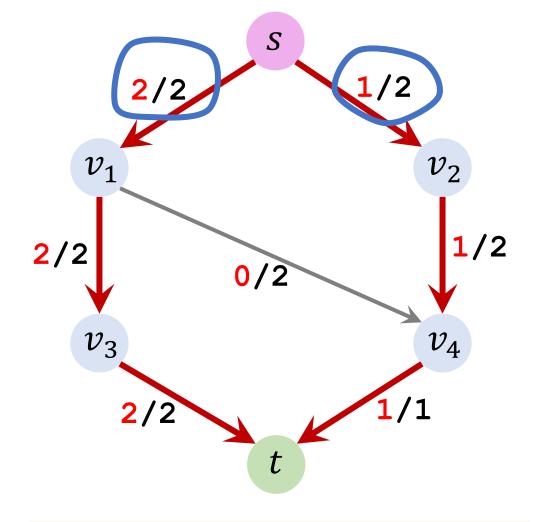




Original Graph

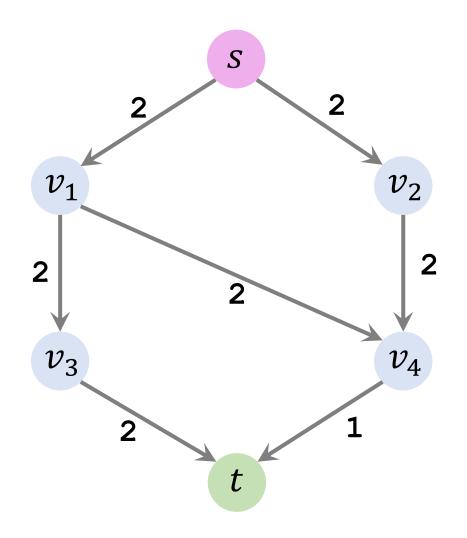
A Blocking Flow

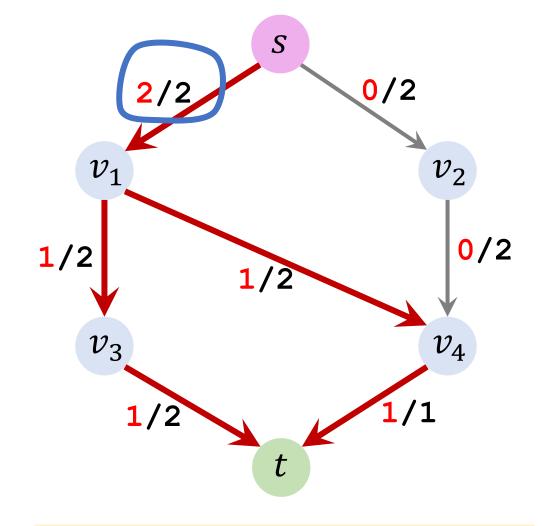




Original Graph

A Blocking Flow



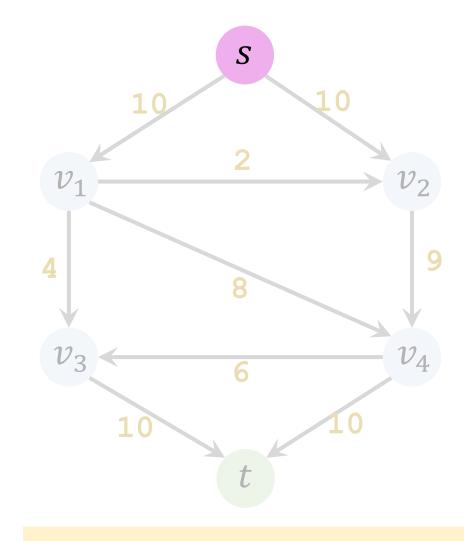


Original Graph

Another Blocking Flow

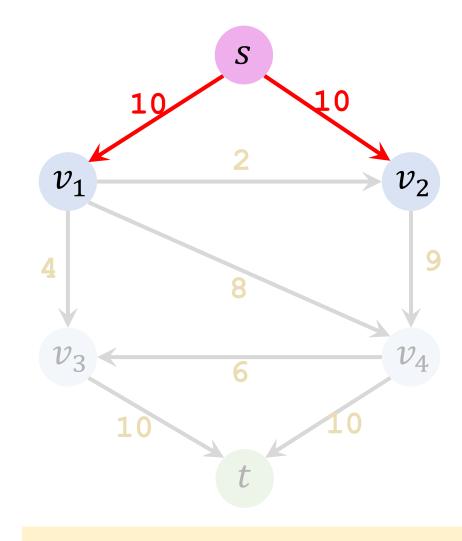
Key Concept: Level Graph

S

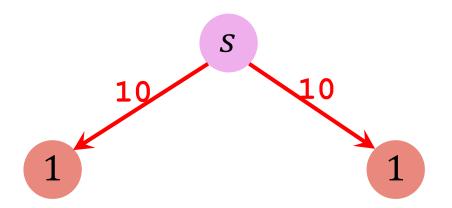


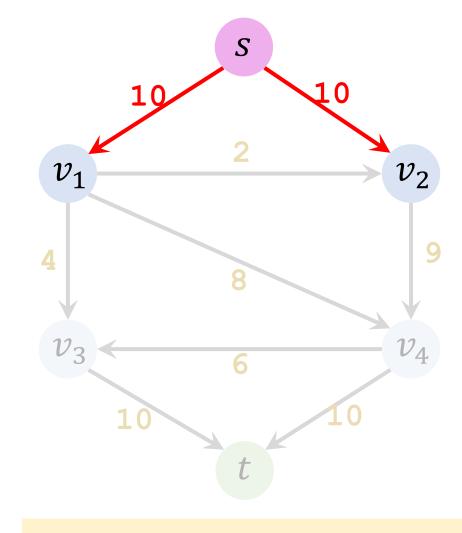
Level Graph

S

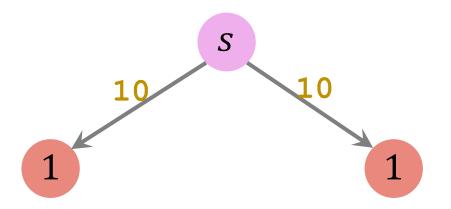


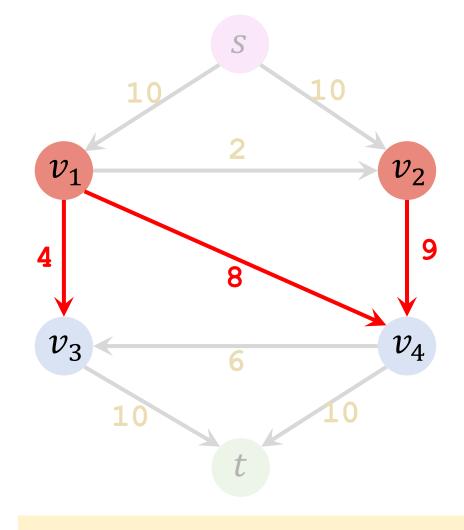
Level Graph



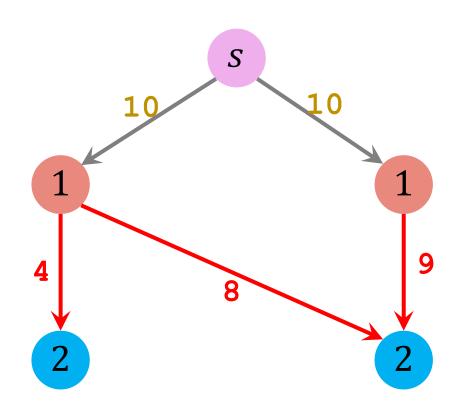


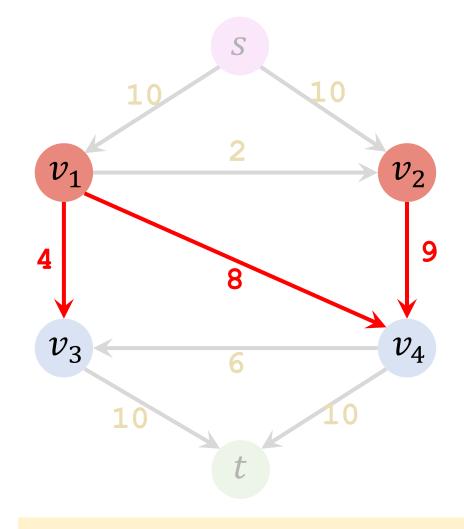
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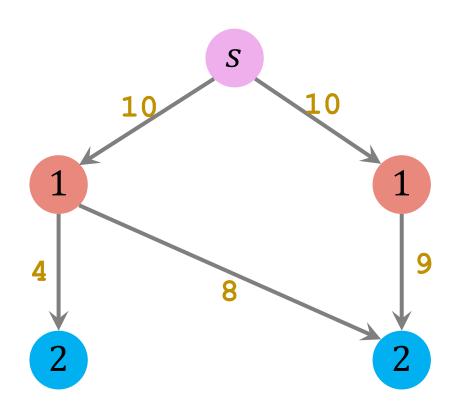


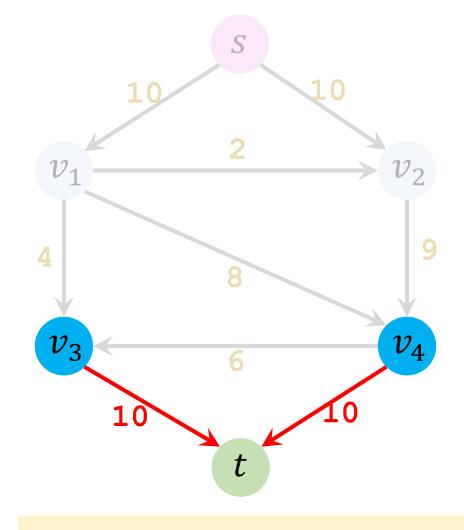
Level Graph



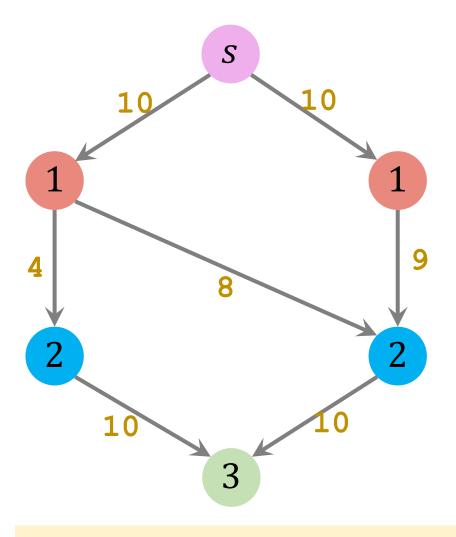


Level Graph

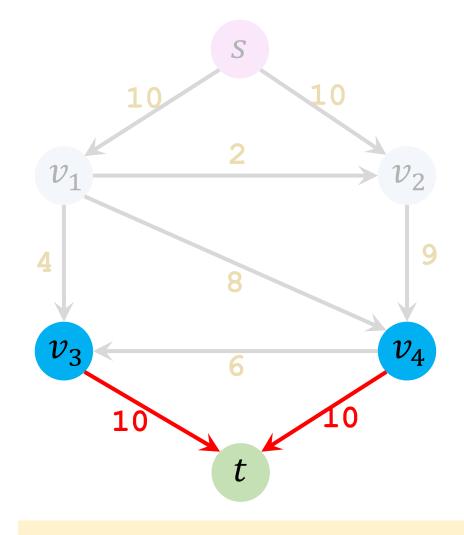


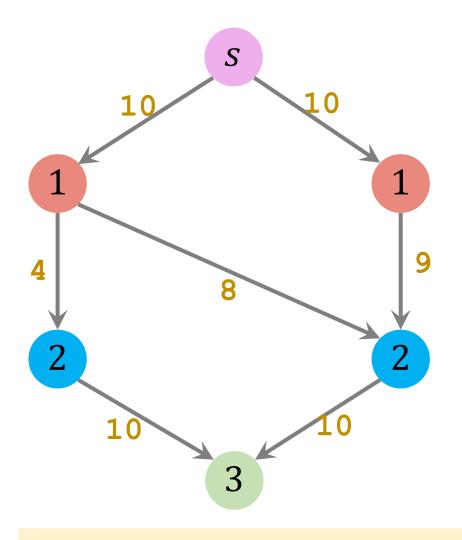


Level Graph

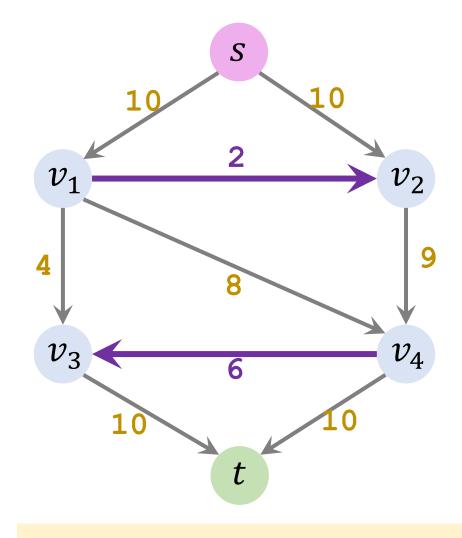


Level Graph

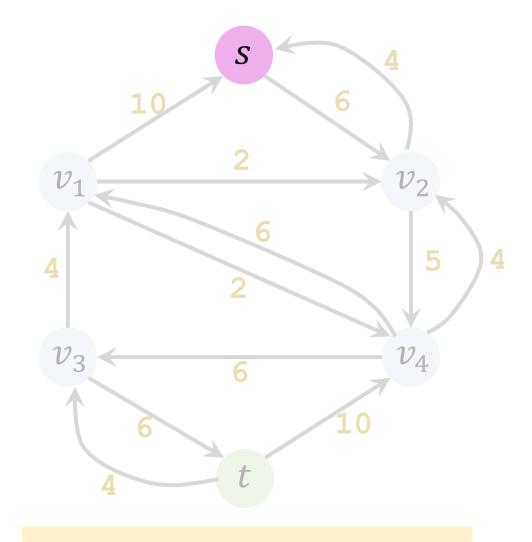




Level Graph

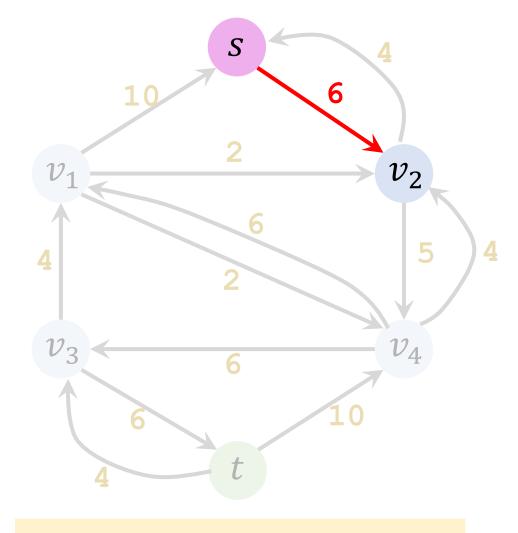


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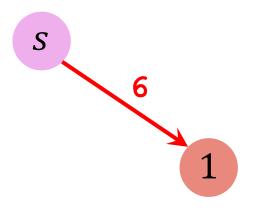


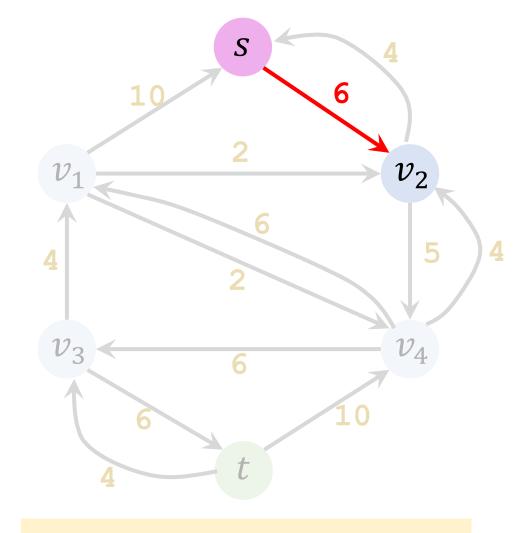
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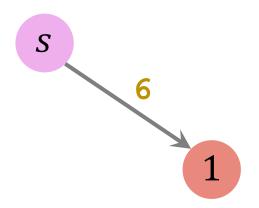


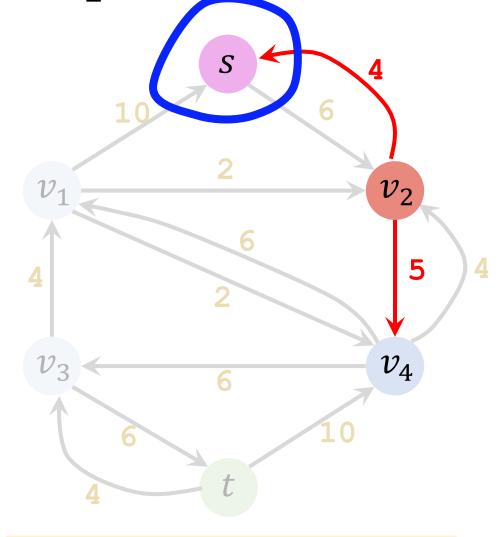
Level Graph



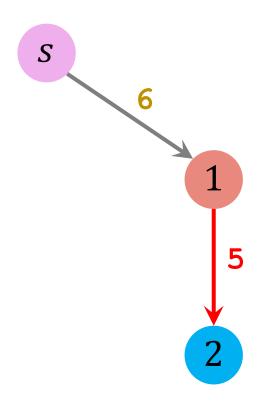


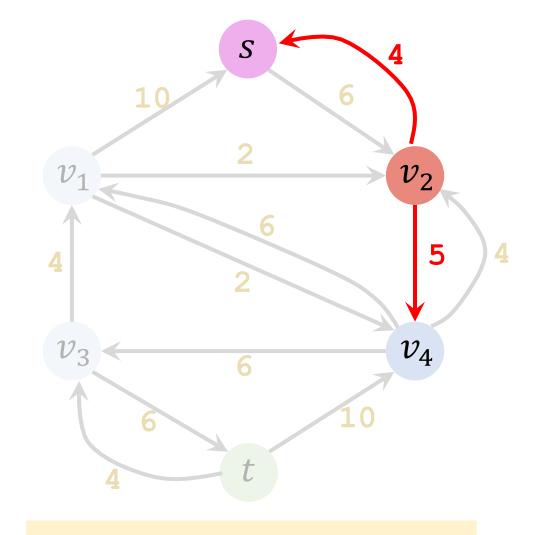
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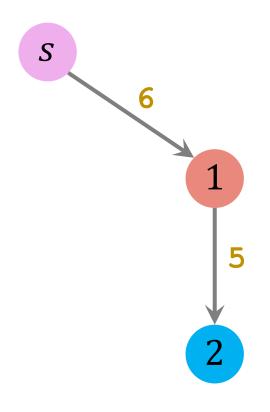


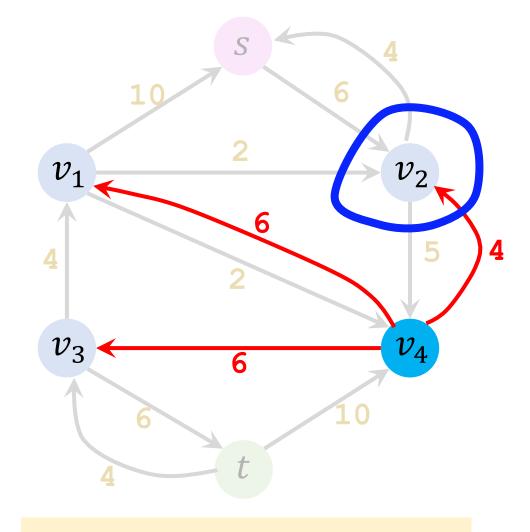
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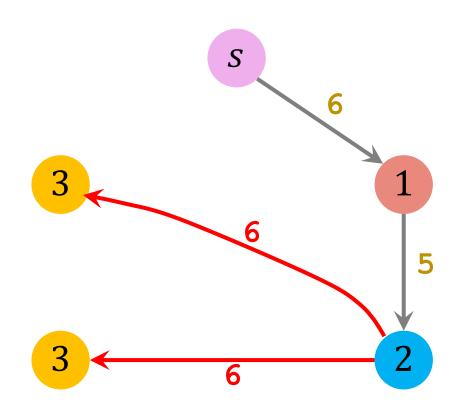


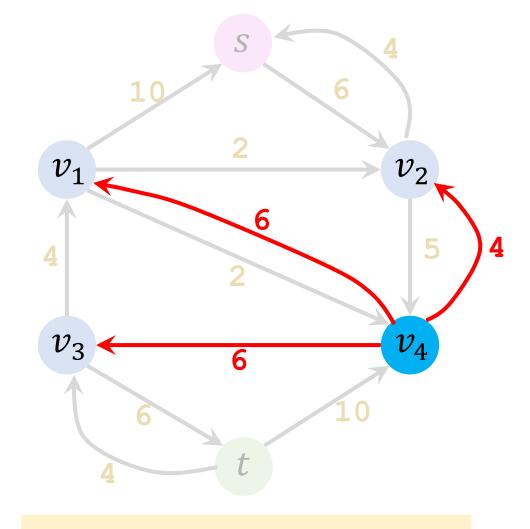
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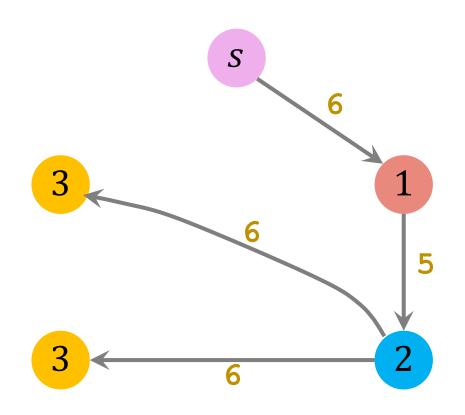


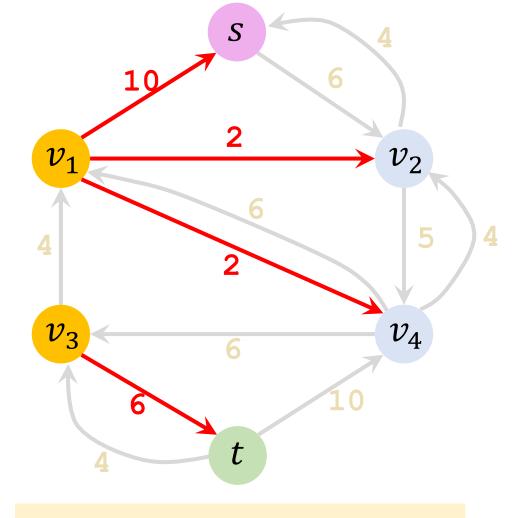
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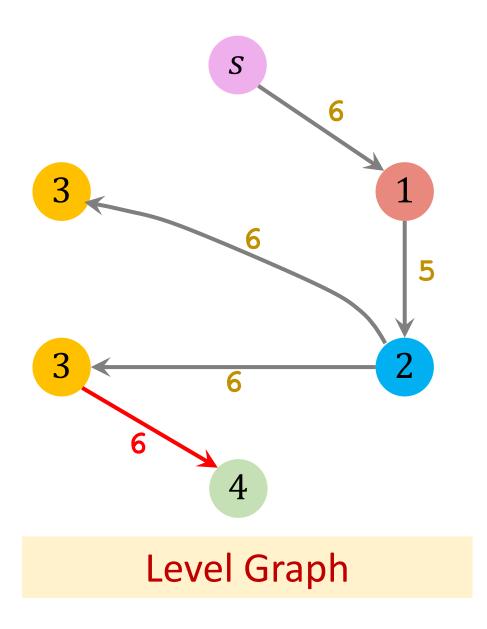


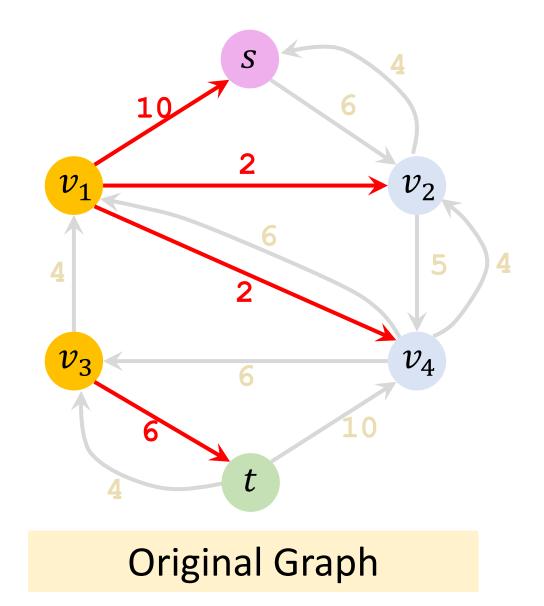
Level Graph





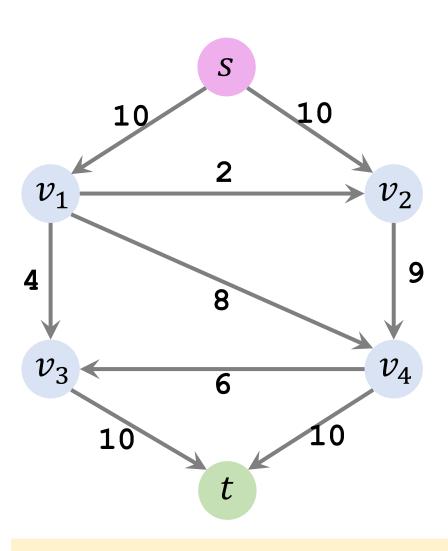
Level Graph



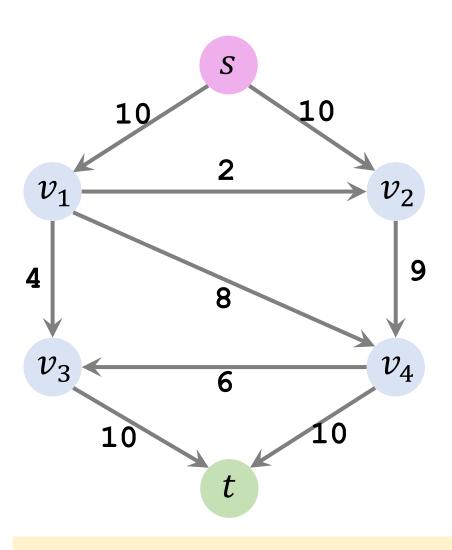


Dinic's Algorithm

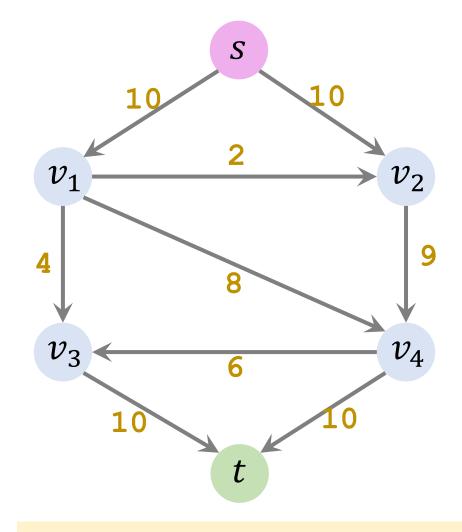
Initialization



Initialization

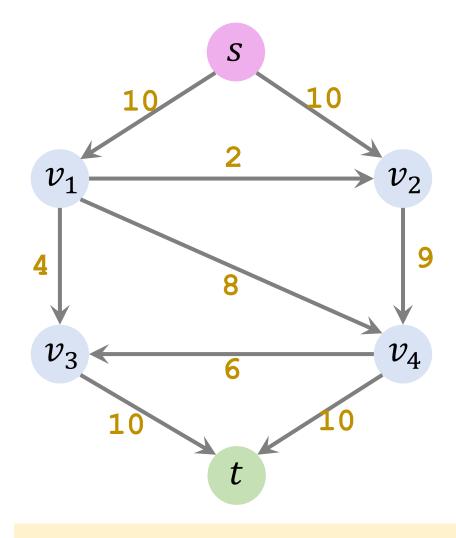


Original Graph



Residual Graph

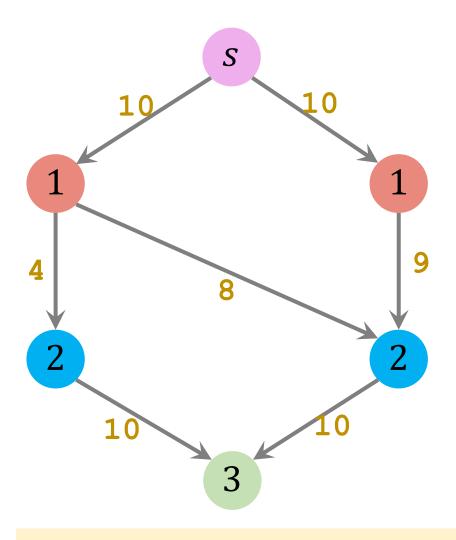
Iteration 1: Construct level graph



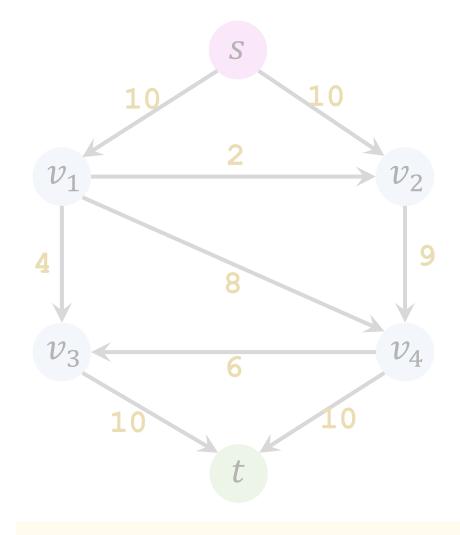
Level Graph

Residual Graph

Iteration 1: Construct level graph

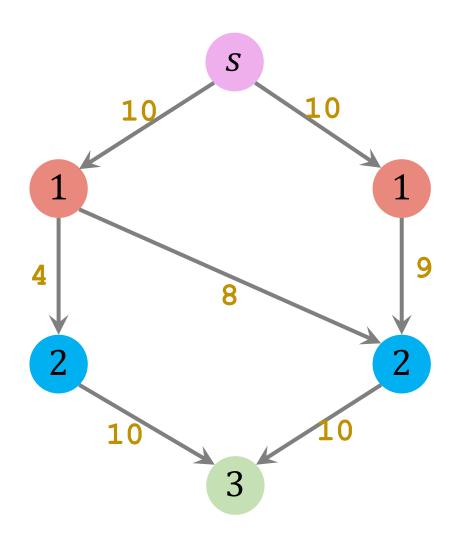


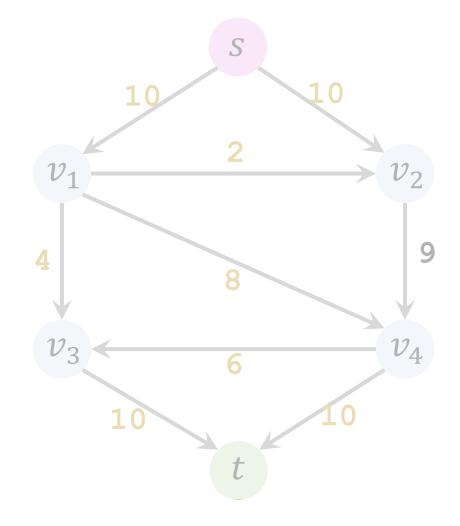
Level Graph



Residual Graph

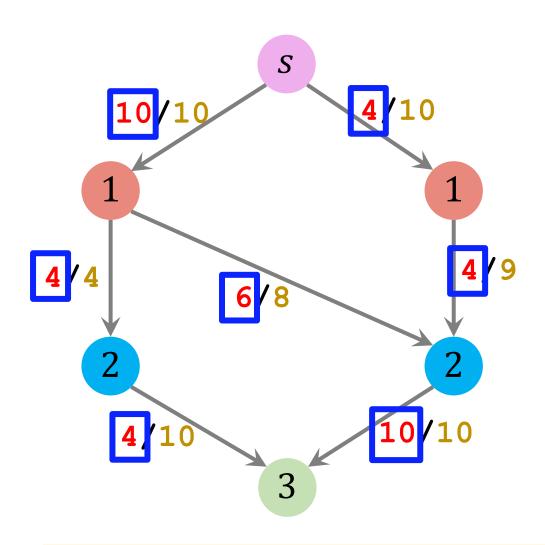
Iteration 1: Find blocking flow in level graph

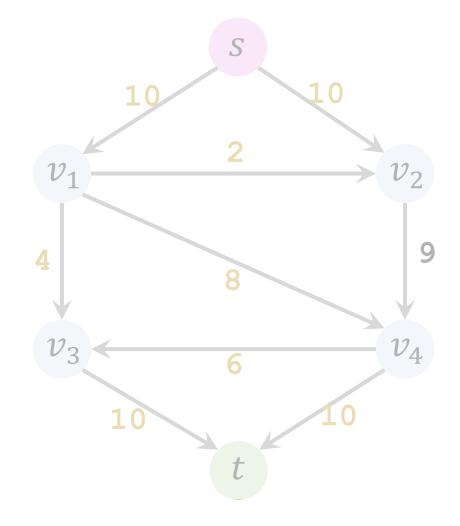




A flow is **blocking flow** if no more flow from source to sink can be found.

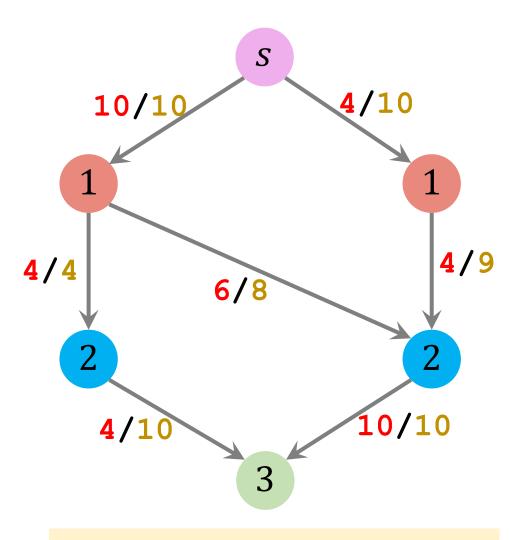
Iteration 1: Find blocking flow in level graph



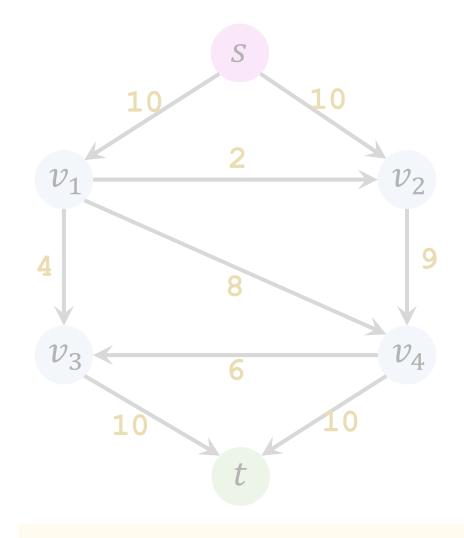


Blocking flow can be found using the naïve algorithm.

Iteration 1: Update the residual graph

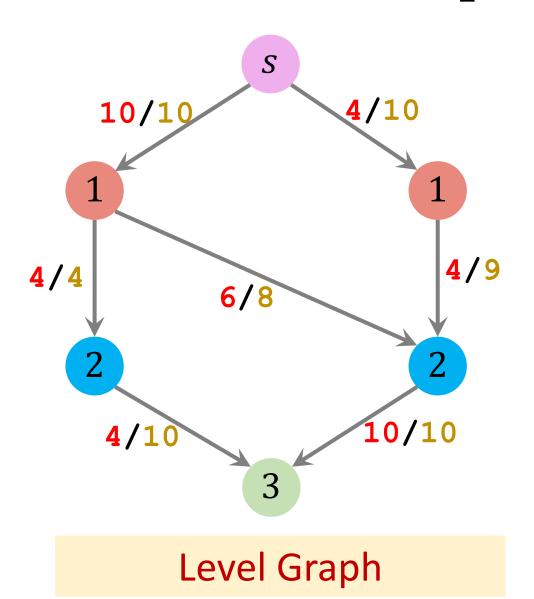


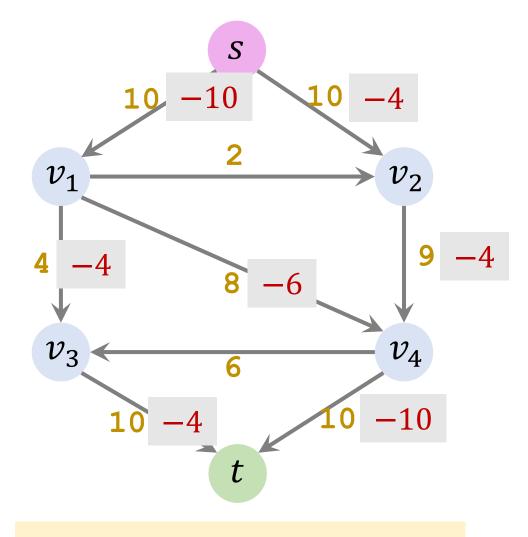
Level Graph



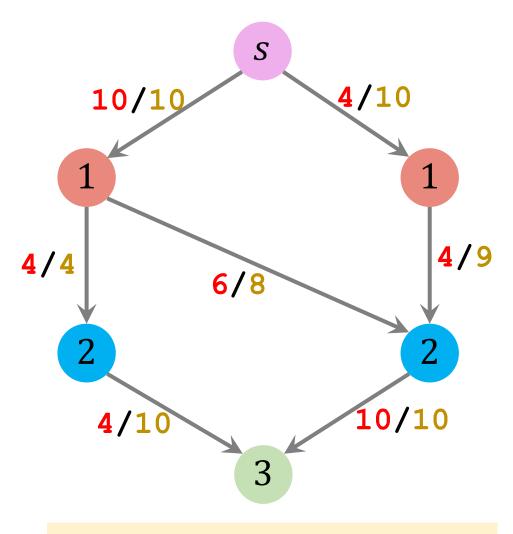
Old Residual Graph

Iteration 1: Update the residual graph

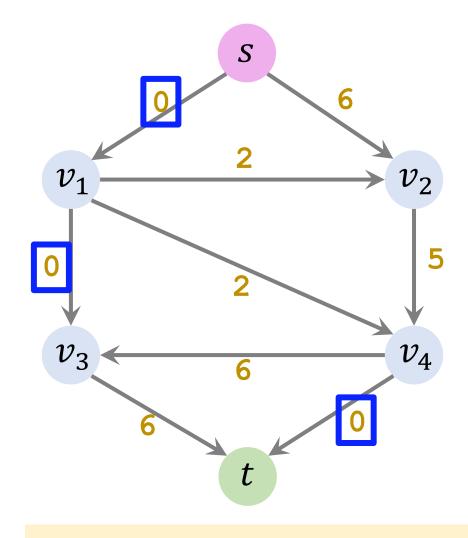




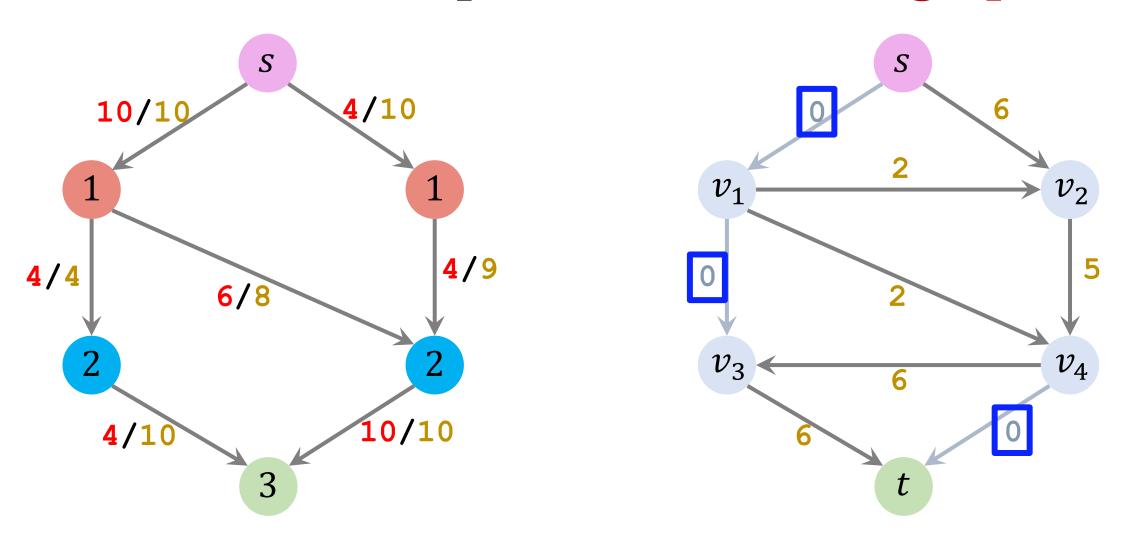
Old Residual Graph



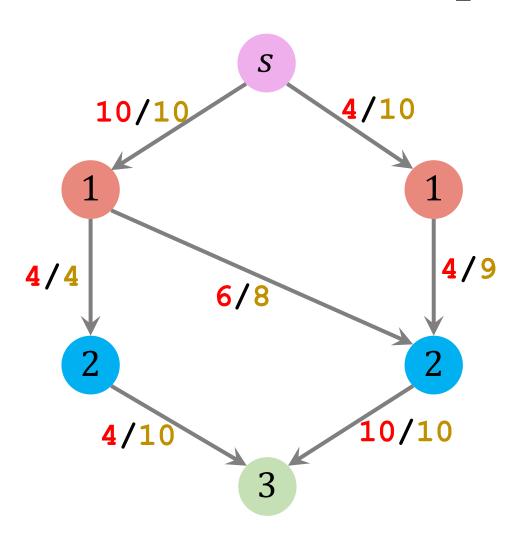
Level Graph

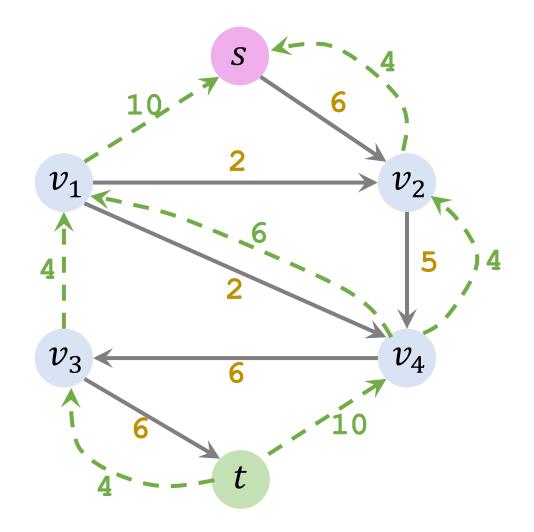


New Residual Graph

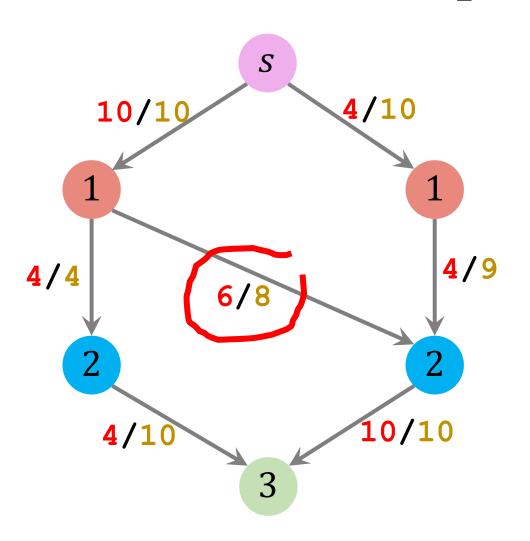


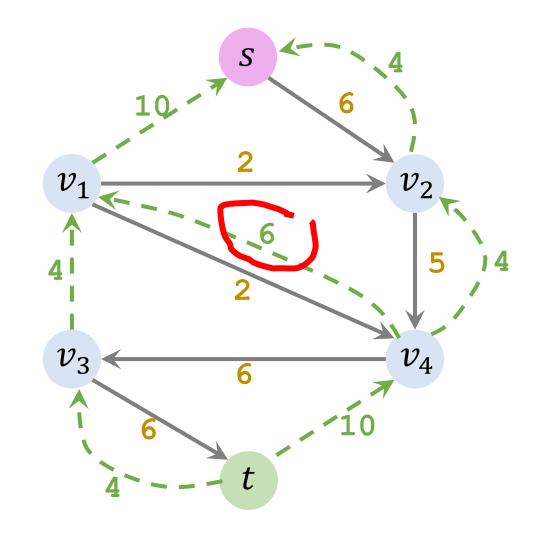
Removed saturated edges from residual graph.





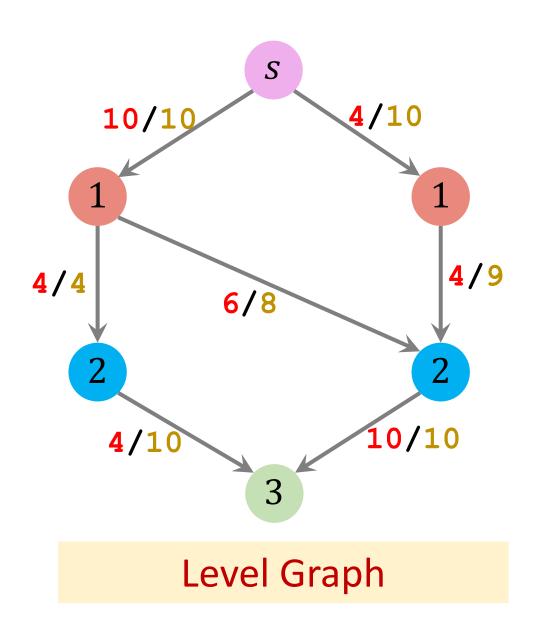
Add flows to the residual graph as backward paths.

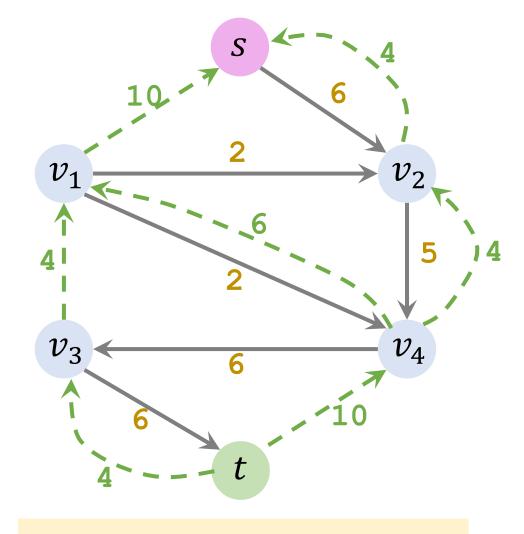




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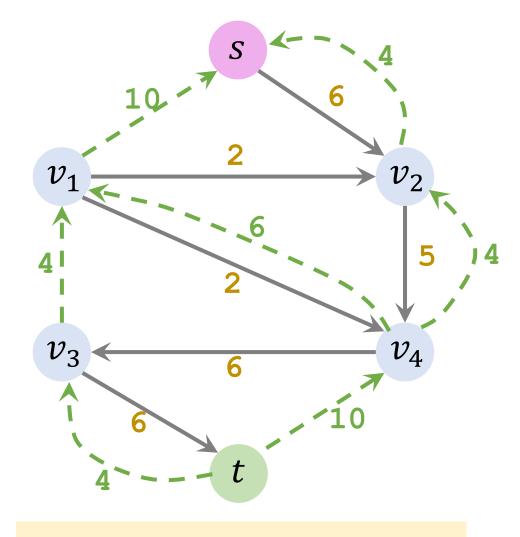
Iteration 1: Finished





Residual Graph

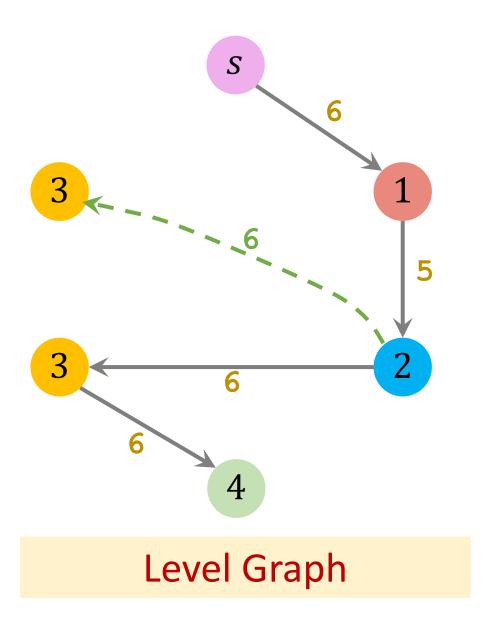
Iteration 2: Construct level graph

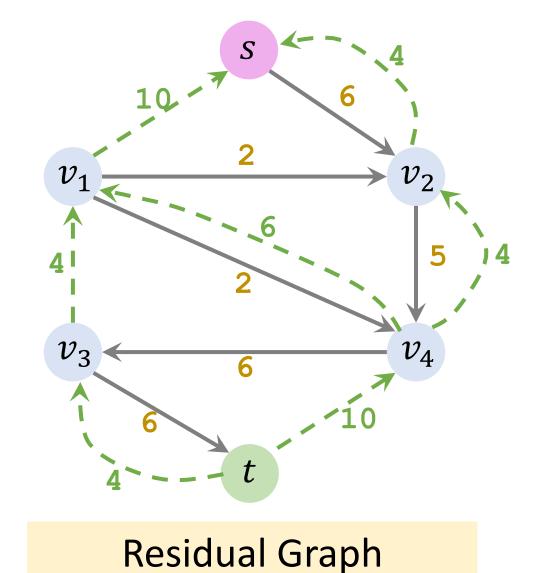


Level Graph

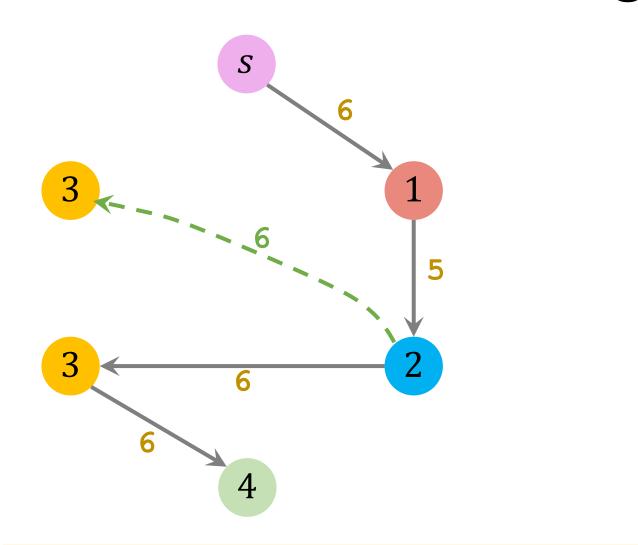
Residual Graph

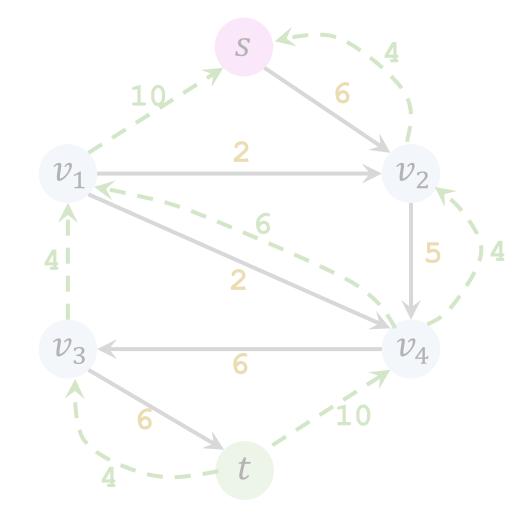
Iteration 2: Construct level graph





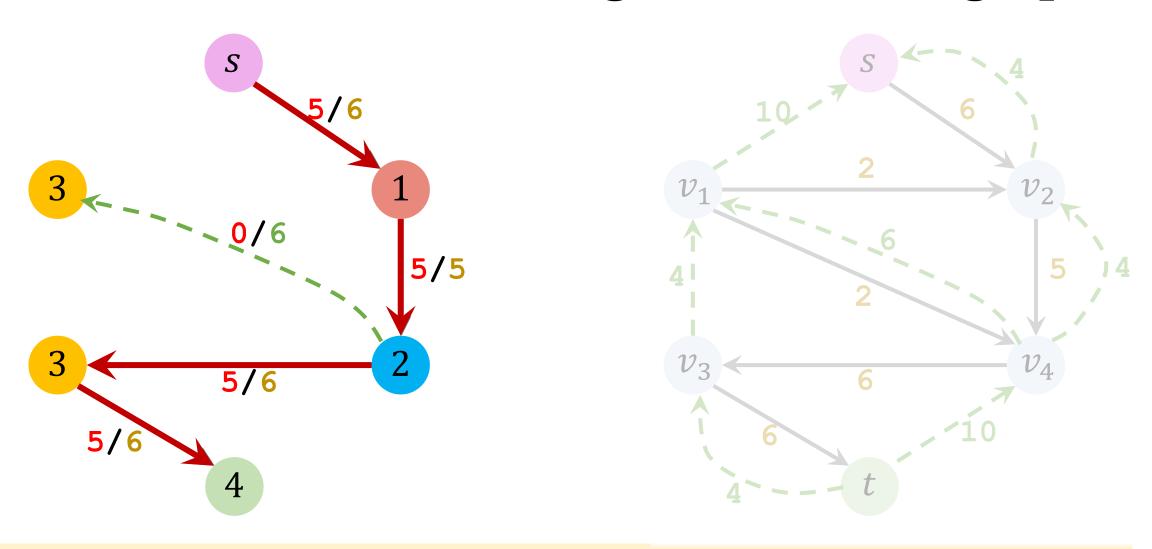
Iteration 2: Find blocking flow in level graph



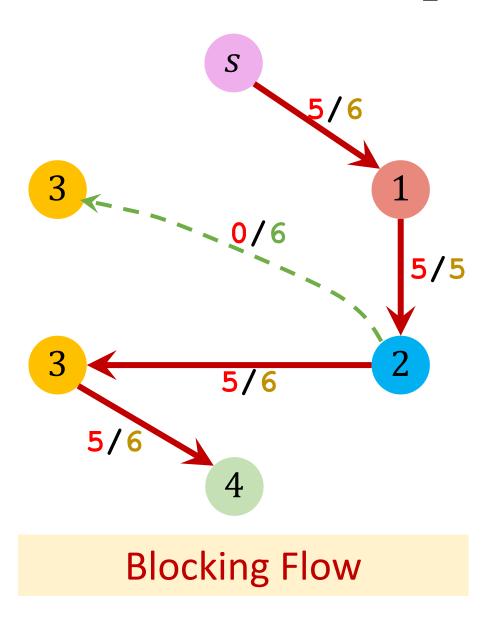


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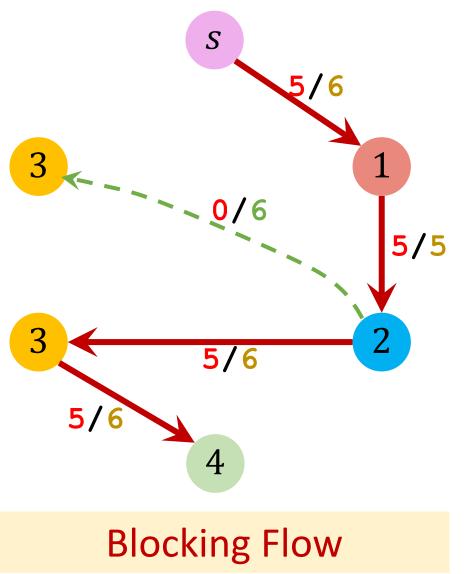
Iteration 2: Find blocking flow in level graph



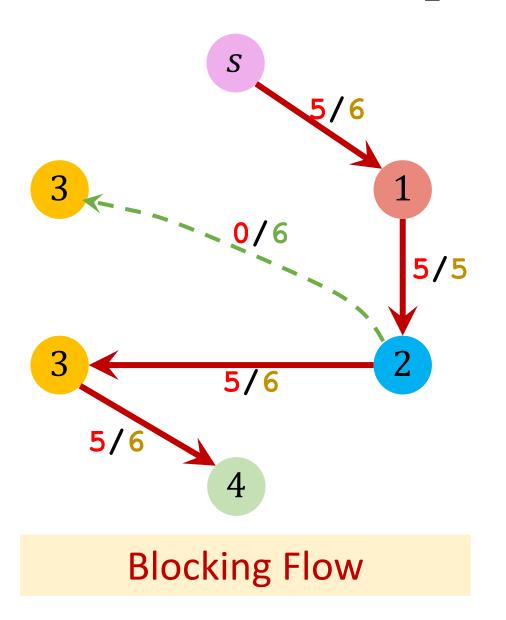
Blocking flow can be found using the naïve algorithm.

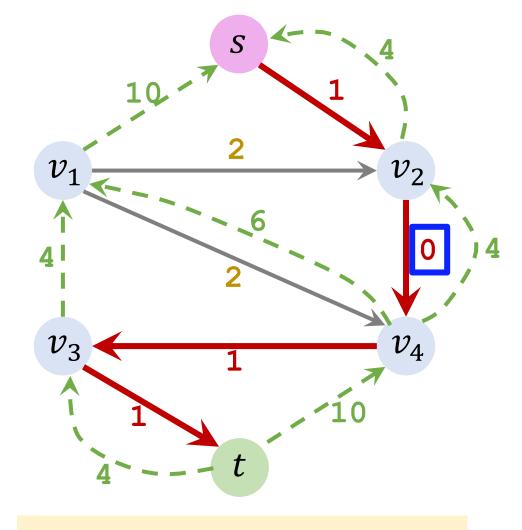


Old Residual Graph

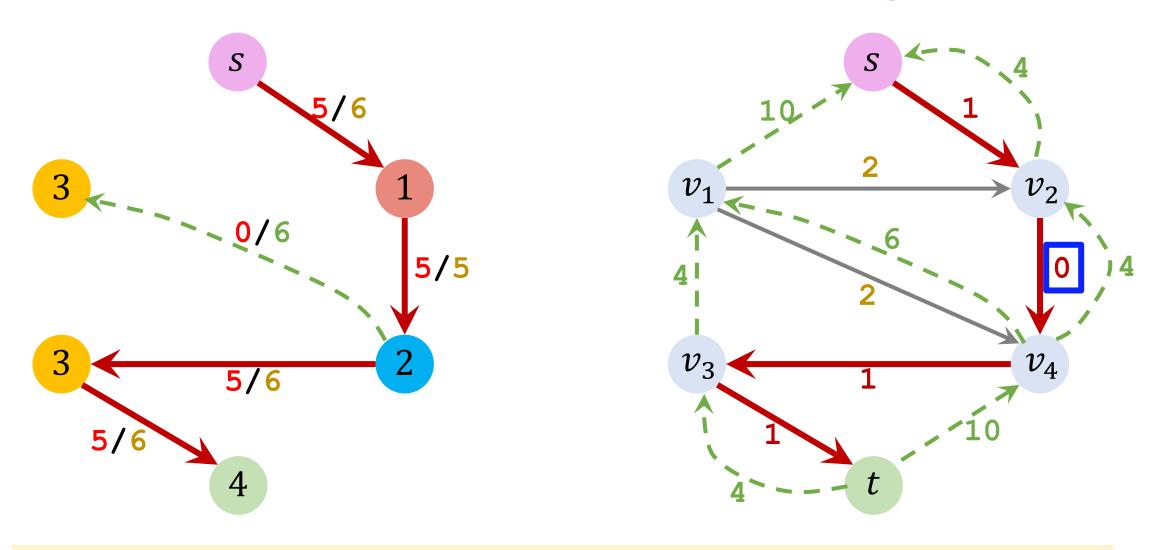


w Old Residual Graph

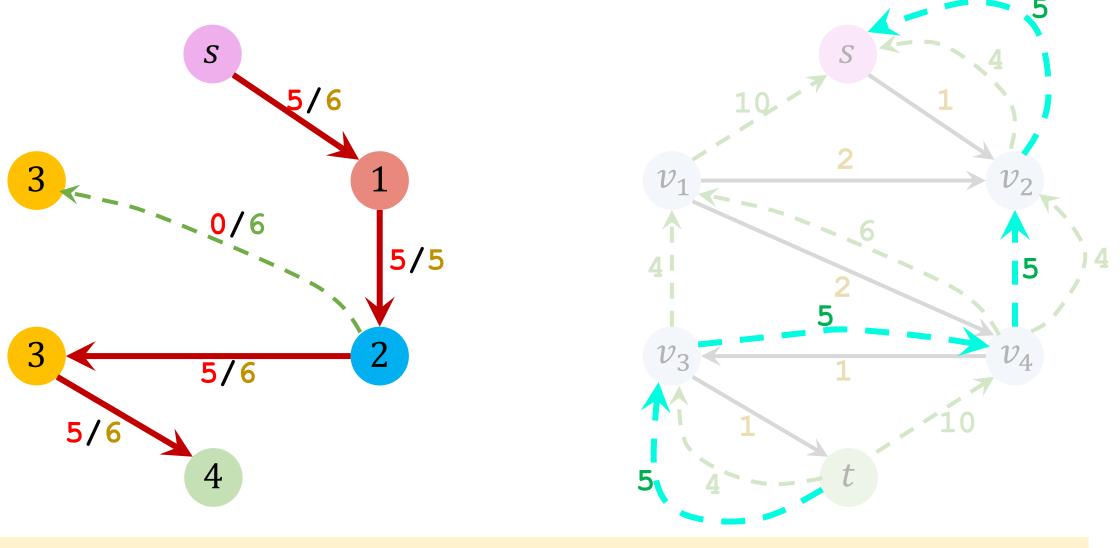




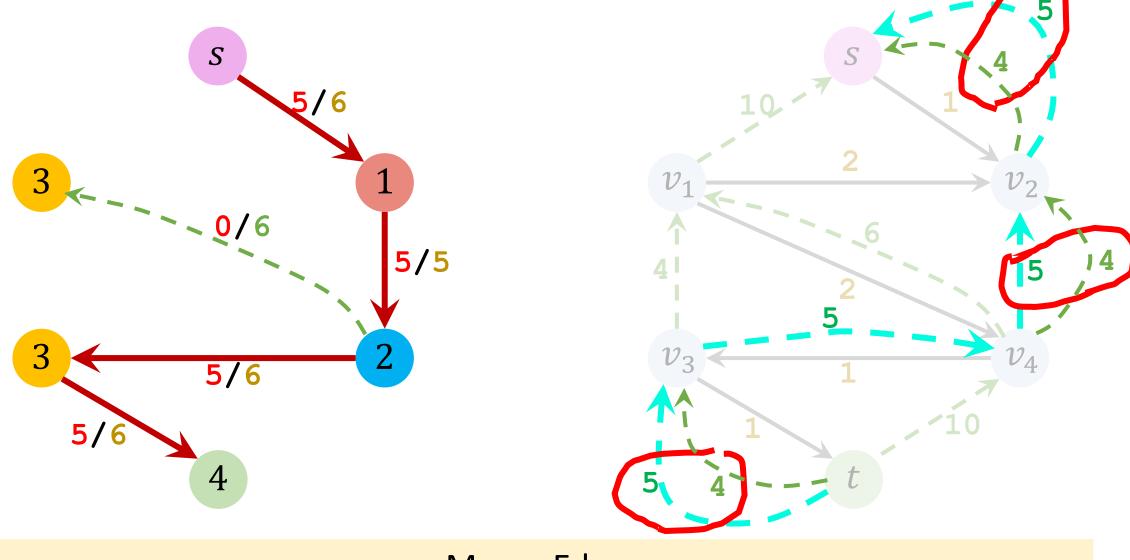
New Residual Graph



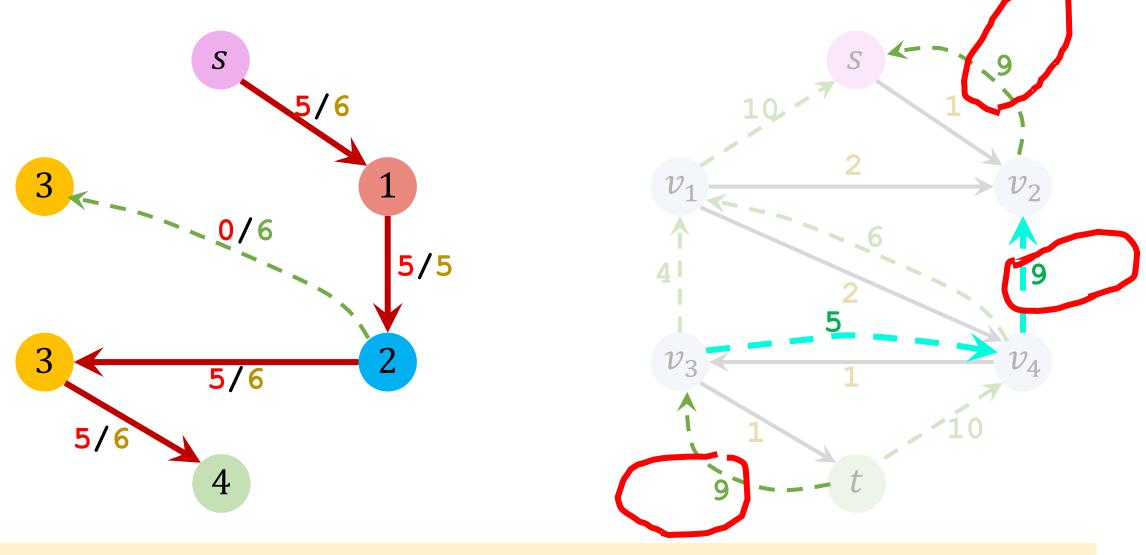
Removed saturated edges from residual graph.



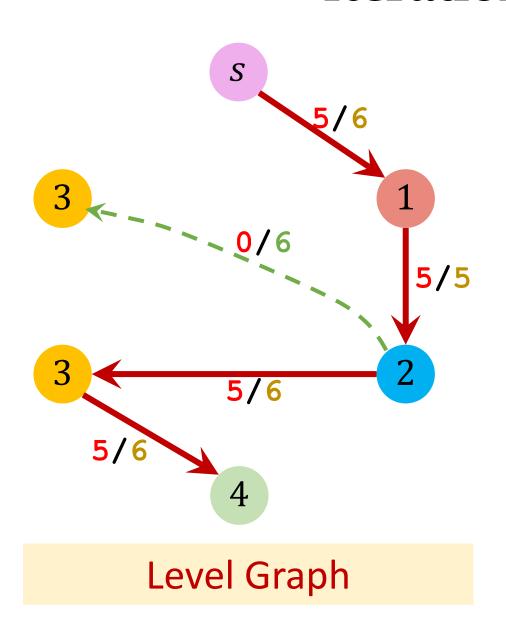
Add flows to the residual graph as backward paths.

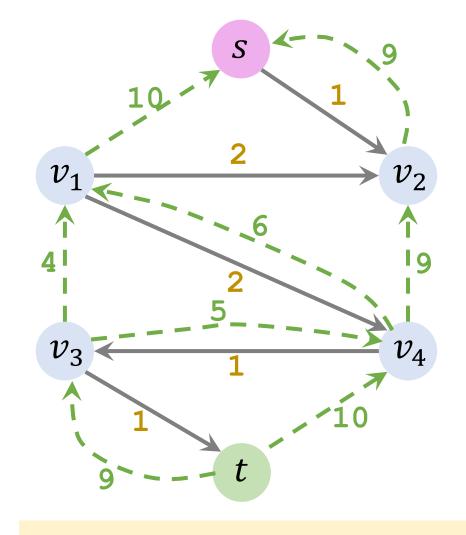


Merge Edges



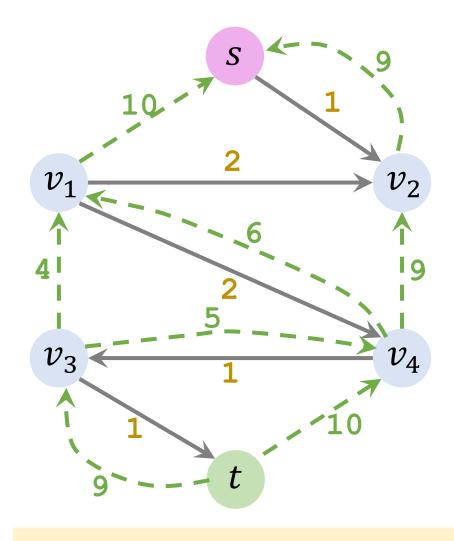
Iteration 2: Finished





Residual Graph

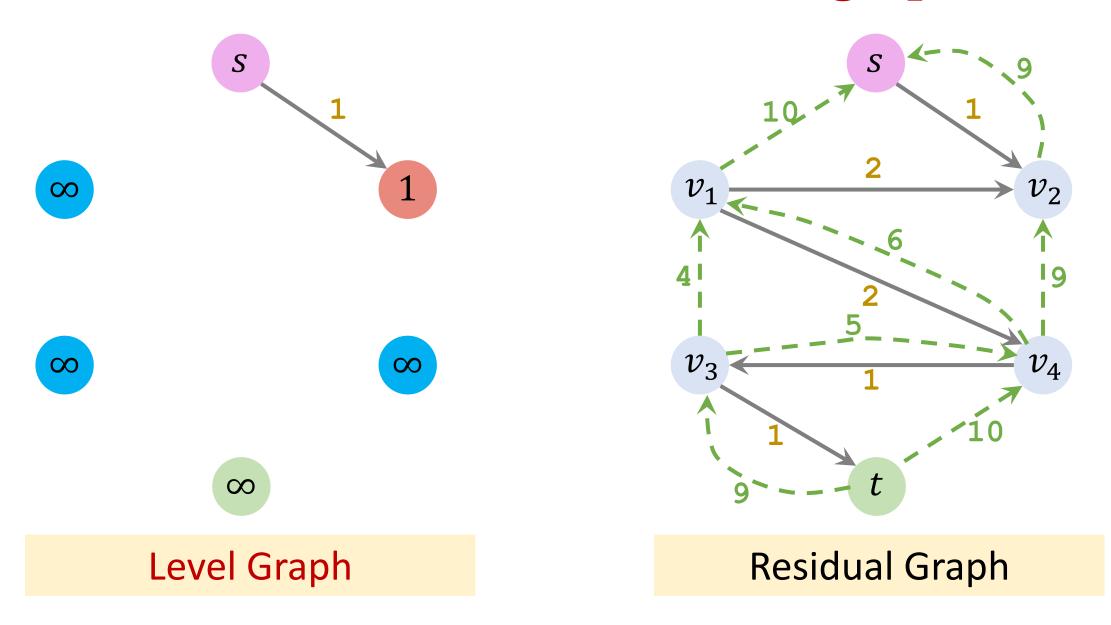
Iteration 3: Construct level graph



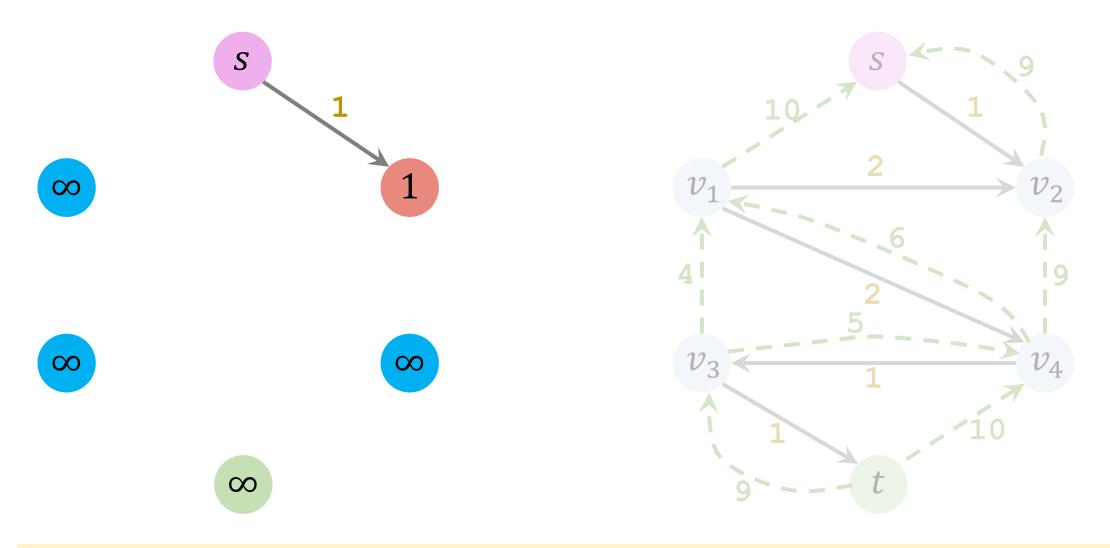
Level Graph

Residual Graph

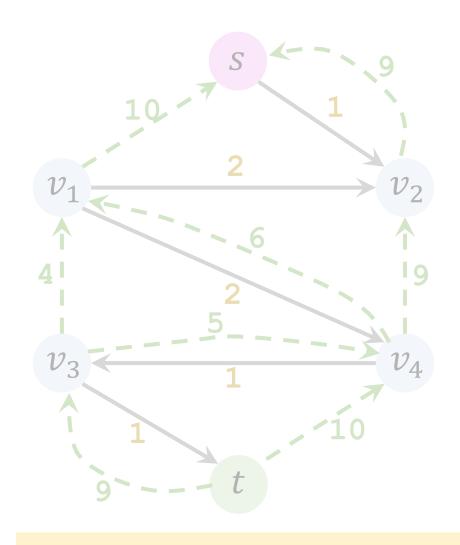
Iteration 3: Construct level graph



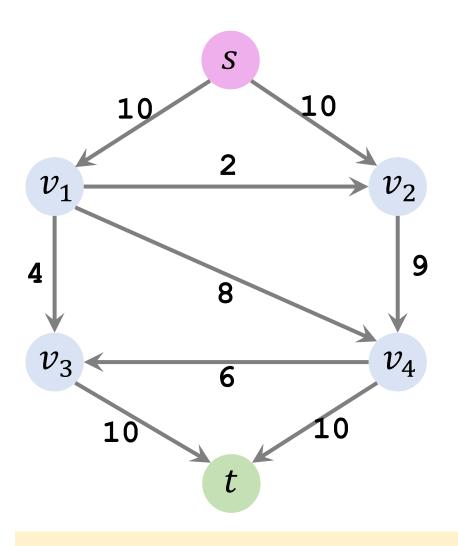
Iteration 3: Find blocking flow in level graph



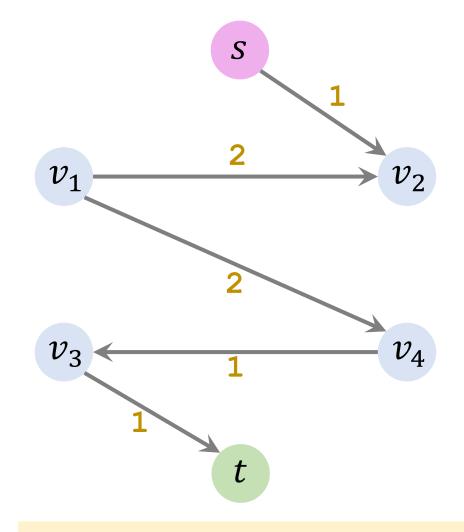
On the level graph, no flow can be found!



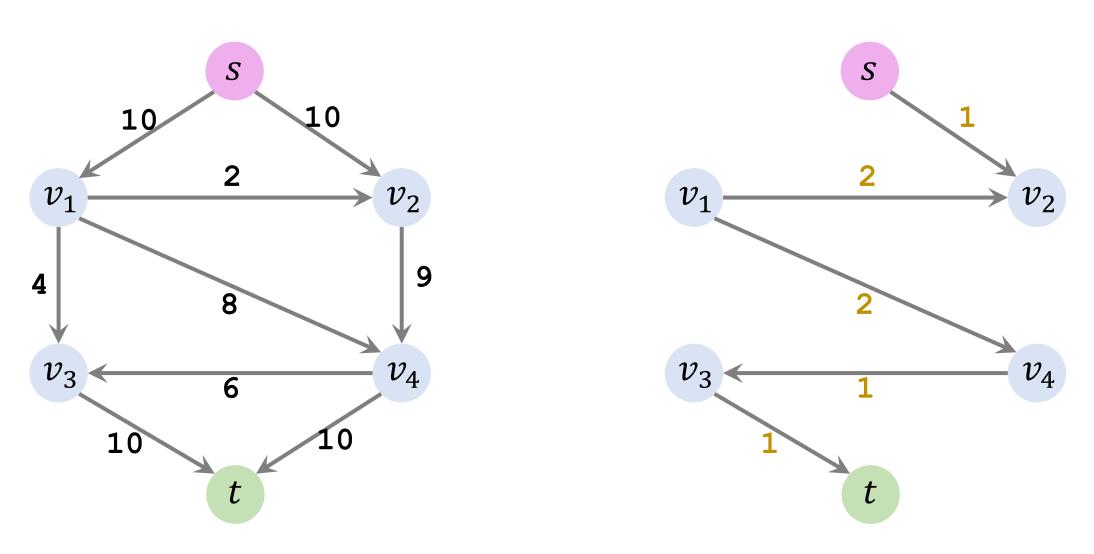
Residual Graph



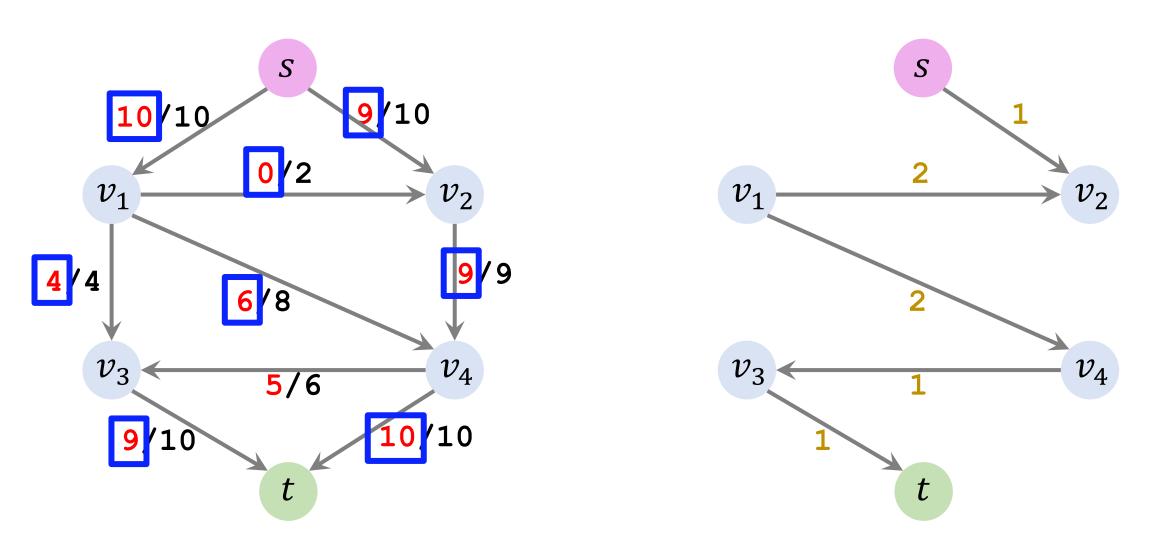
Original Graph



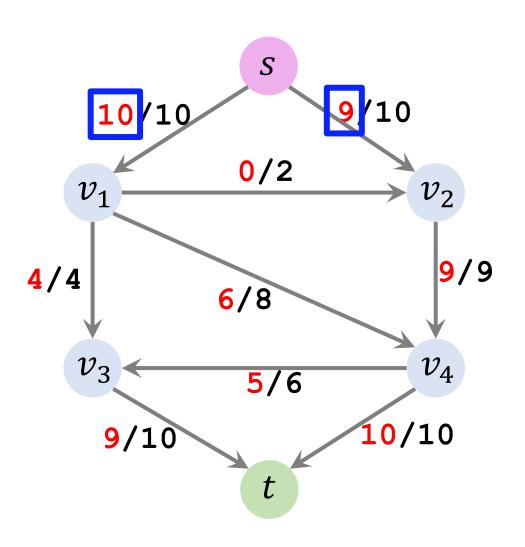
Residual Graph

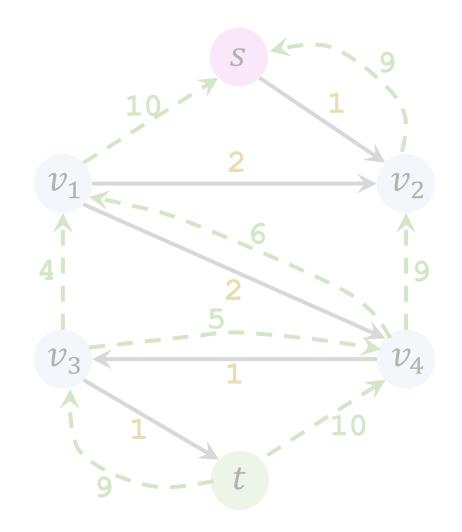


Flow = Capacity - Residual.



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Amount of Max Flow = 19.

Summary

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1. Initially, the residual graph is a copy of the original graph.

2. Repeat:

- a. Construct the level graph of the residual graph.
- b. Find a blocking flow on the level graph.
 - c. Update the residual graph (update the weights, remove saturated edges, and add backward edges.)

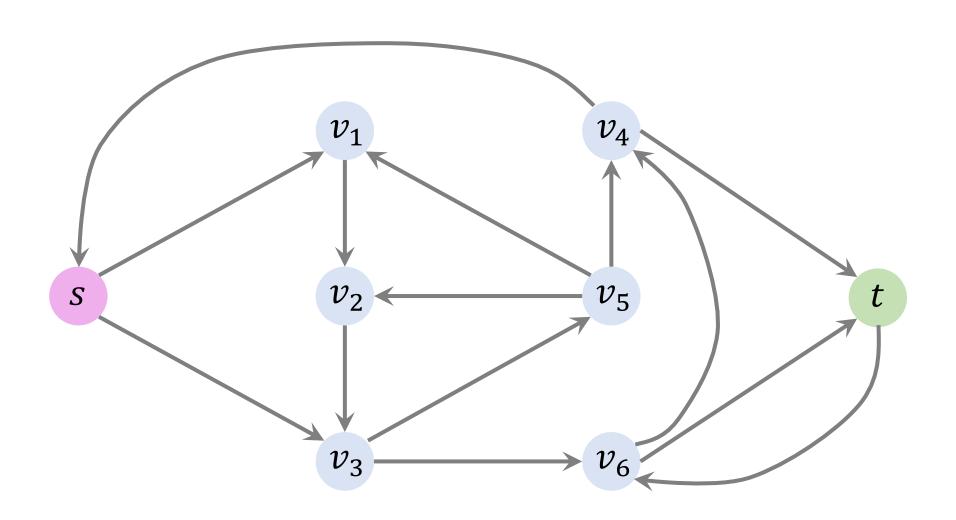
Time Complexity

Time complexity: $O(m \cdot n^2)$. (m is #edges; n is #vertices.)

- Dinic's algorithm has at most n-1 iterations.
- Per-iteration time complexity is O(mn).

Questions

Q1: What is the level graph?



Thank You!