

A

$$\begin{aligned}\sum_{i=1}^n \sum_{j=1}^m \gcd(i, j) &= \sum_{i=1}^n \sum_{j=1}^m \sum_{d=1}^n [\gcd(i, j) = d] d \\&= \sum_{d=1}^n d \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{m}{d} \rfloor} [\gcd(i, j) = 1] \\&= \sum_{d=1}^n d \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{m}{d} \rfloor} \sum_{k|i, k|j} \mu(k) \\&= \sum_{d=1}^n d \sum_{k=1}^{\lfloor \frac{n}{d} \rfloor} \mu(k) \sum_{k|i}^{\lfloor \frac{n}{d} \rfloor} \sum_{k|j}^{\lfloor \frac{m}{d} \rfloor} 1 \\&= \sum_{d=1}^n d \sum_{k=1}^{\lfloor \frac{n}{d} \rfloor} \mu(k) \lfloor \frac{n}{kd} \rfloor \lfloor \frac{m}{kd} \rfloor\end{aligned}$$

B

有个神奇的柿子 $D(ab) = \sum_{i|a} \sum_{j|b} [\gcd(i, j) = 1]$

$$\begin{aligned}Ans &= \sum_{i=1}^n \sum_{j=1}^m \sum_{a|i} \sum_{b|j} [\gcd(a, b) = 1] \\&= \sum_{a=1}^n \sum_{b=1}^m \lfloor \frac{n}{a} \rfloor \lfloor \frac{m}{b} \rfloor [\gcd(a, b) = 1] \\&= \sum_{a=1}^n \sum_{b=1}^m \lfloor \frac{n}{a} \rfloor \lfloor \frac{m}{b} \rfloor \sum_{k|a, k|b} \mu(k) \\&= \sum_{k=1}^n \mu(k) \sum_{a=1}^{\lfloor \frac{n}{k} \rfloor} \sum_{b=1}^{\lfloor \frac{m}{k} \rfloor} \lfloor \frac{n}{ak} \rfloor \lfloor \frac{m}{bk} \rfloor\end{aligned}$$

可以发现右边那两个东西实际上只和 n/k 和 m/k 有关

所以说再整除分块预处理出来 $\sum_{i=1}^a \lfloor \frac{n}{i} \rfloor$

对每个询问都整除分块，发现再处理一下 μ 的前缀和即可

C

$$f(k) = \sum_{d=1}^a \mu(d) \lfloor \frac{a}{dk} \rfloor \lfloor \frac{b}{dk} \rfloor$$

显然可以整除分块

D

C 加点容斥即可

E

$$\begin{aligned}
& \sum_{i=1}^n \sum_{j=1}^m \gcd(i, j)^k \\
&= \sum_{i=1}^n \sum_{j=1}^m \sum_{d=1}^n [\gcd(i, j) = d] d^k \\
&= \sum_{d=1}^n d^k \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{m}{d} \rfloor} [\gcd(i, j) = 1] \\
&= \sum_{d=1}^n d^k \sum_{i=1}^{\lfloor \frac{n}{d} \rfloor} \sum_{j=1}^{\lfloor \frac{m}{d} \rfloor} \sum_{l|i, l|j} \mu(l) \\
&= \sum_{d=1}^n d^k \sum_{l=1}^{\lfloor \frac{n}{d} \rfloor} \mu(l) \lfloor \frac{n}{dl} \rfloor \lfloor \frac{m}{dl} \rfloor
\end{aligned}$$

设 $D = dl$

$$= \sum_{d=1}^n \lfloor \frac{n}{D} \rfloor \lfloor \frac{m}{D} \rfloor \sum_{d|D} d^k \mu(\frac{D}{d})$$

发现右边 $f(D) = \sum_{d|D} d^k \mu(\frac{D}{d})$, 是积性函数 ($d^k, \mu(\frac{D}{d})$ 均为积性函数, 他们的狄利克雷卷积也为积性函数)

$$f(p) = p^k - 1$$

$$f(p^{x_p}) = p^{x_p k} - p^{(x_p-1)k} = p^{(x_p-1)k} * (p - 1)$$

令 p 为 a 的最小质因数

$$f(a) = \begin{cases} f(a/(p^{x_p})) * f(p^{x_p}) = f(a/p) \times p^k & (x_p > 1) \\ f(a/p) \times (p^k - 1) & (x_p = 1) \end{cases}$$

然后就可以线性递推了

左边的 $\sum_{d=1}^n \lfloor \frac{n}{D} \rfloor \lfloor \frac{m}{D} \rfloor$ 可以整除分块

线性递推代码

```

void check()
{
    sum[1]=1;
    for (int i=2;i<=5000000;i++)
    {
        if (!is_prime[i])
        {
            prime[++cnt]=i;
            sum[i]=ksm(i,k)-1;
            if (sum[i]<0) sum[i]+=mod;
        }
        for (int j=1;j<=cnt;j++)
        {
            if (prime[j]*i>5000000) break;
            is_prime[i*prime[j]]=1;
            if (i%prime[j]==0)
            {
                sum[i*prime[j]]=sum[i]*(sum[prime[j]]+1)%mod;
                break;
            }
            else
            {
                sum[i*prime[j]]=sum[i]*sum[prime[j]]%mod;
            }
        }
    }
    for (int i=2;i<=5000000;i++)
        sum[i]=(sum[i-1]+sum[i])%mod;
}

```

