

Lecture ÷ Tries - 1

Agenda

- Introduction
- Operations
 - search
 - Insert
 - Delete
- Shortest unique prefix

Trie:

Retrieval \rightarrow fetch, search.

DS that stores data from top to bottom.
characters
bits [next class]

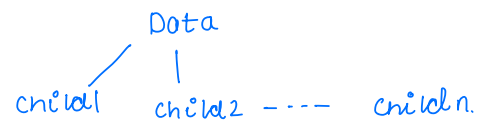
Structure

Trees [Binary]



```
class Node {  
    int data;  
    Node left;  
    Node right;  
}
```

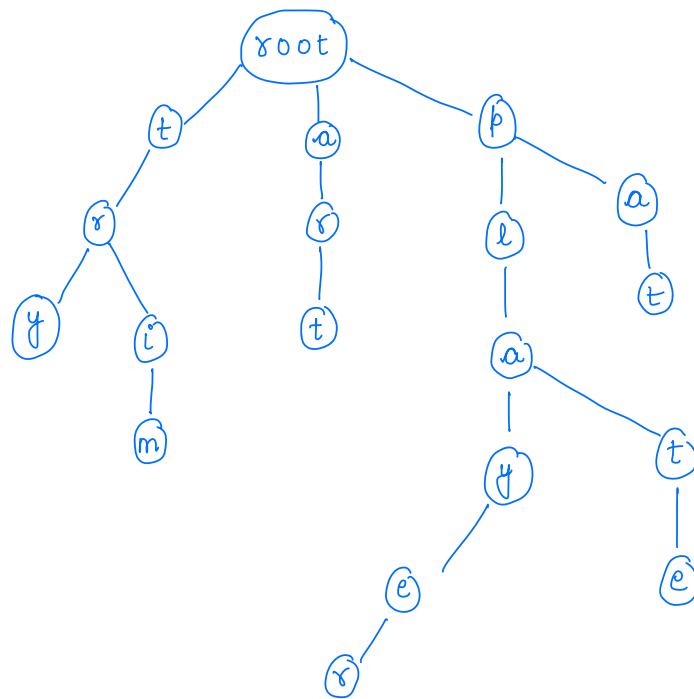
Tries



```
class Node {  
    char data;  
    Node[] children;  
    Node(x) {  
        data = x;  
    }  
}
```

11 words

try
art
play
trim
plate
pat
player.
trim.



Qu: Check valid words.

isValid(plate) = true

isValid(srk) = false

isValid(pla) = false

isValid(player) = true

isValid(aye) = false

Map: TC: $O(len)$
SC: $O(len)$
↑
1 word

n words:

TC: $O(n * len)$

SC: $O(n * len)$

try
art
play
trim
plate
pat
player.

Idea!

Travel complete word in trie and check for leaf.

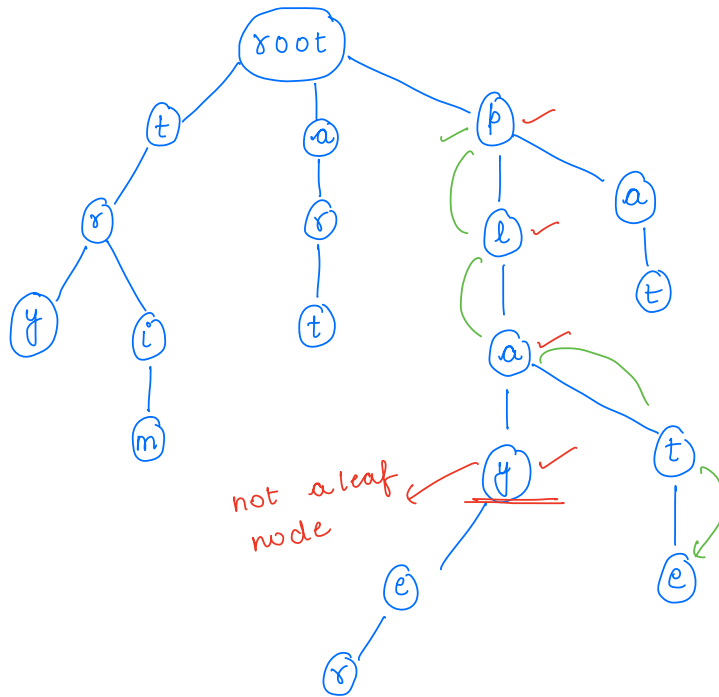


plate :- true
↑↑↑↑

play → Idea does not work here
↑↑↑

Idea 2: Use a flag to denote end of word.

Idea 2: Use a flag to denote end of word.

```
class Node {
```

```
char data;
```

```
Node[] children;
```

```
boolean isEnd;
```

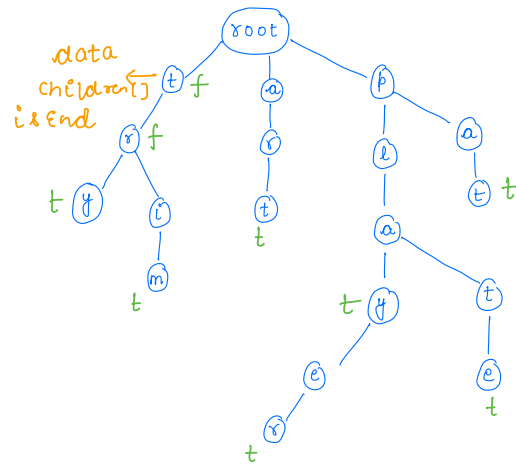
```
Node(x) {
```

```
data = x;
```

```
children = new Node(26);
```

```
i & end = false;
```

1



Modified Idea:

Travel complete word in trie and check whether the last character is end of word or not?

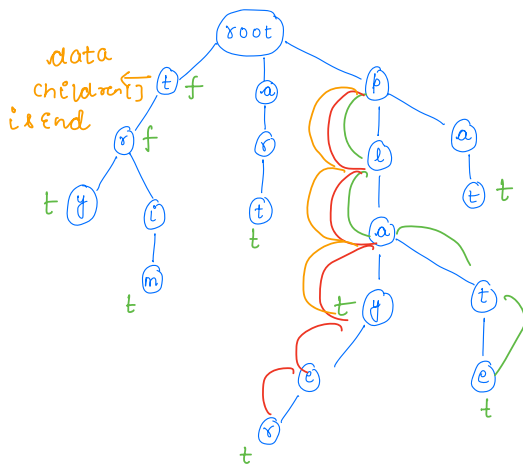


plate :- true



player: true

A green arrow points down from the top of the diagram to a pair of red arrows pointing up.

play : true

↑ ↑ ↑ ↑

pla :-

##

Insertion [Insert a word in a trie]

word = trim

```
void insert(root, word) {
```

```
    curr = root;
```

```
    n = word.length;
```

```
    for(i=0; i<n; i++) {
```

```
        ch = word.charAt(i);
```

```
        idx = ch - 'a';
```

```
        if(curr.children[idx] == null) {
```

```
            Node nn = new Node(ch);
```

```
            curr.children[idx] = nn;
```

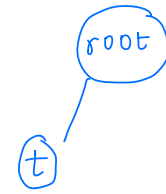
```
        }
```

```
        curr = curr.children[idx];
```

```
    }
```

```
    curr.isEnd = true;
```

```
}
```



search for a word whether it is valid or not?

```
boolean search (root, word) {
```

$$\text{curr} = \text{root};$$

$n = \text{word.length};$

```
for (i=0; i < n; i++) {
```

```
ch = word.charAt(i);
```

$$dx = ch - 'a';$$

```
if (curr.children[idx] == null) {
```

```
return false;
```

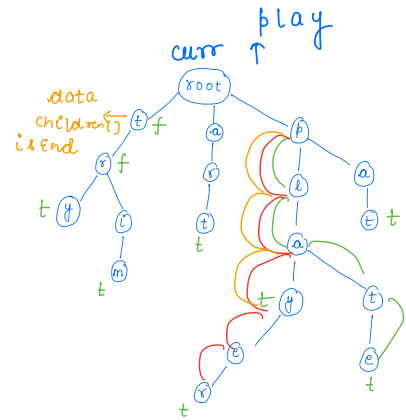
}

```
curr = curr.children[idx];
```

}

```
return curr.isEnd;
```

1



Qn

Frequency of words

trim

try

art

play

trim

trimmer

plate

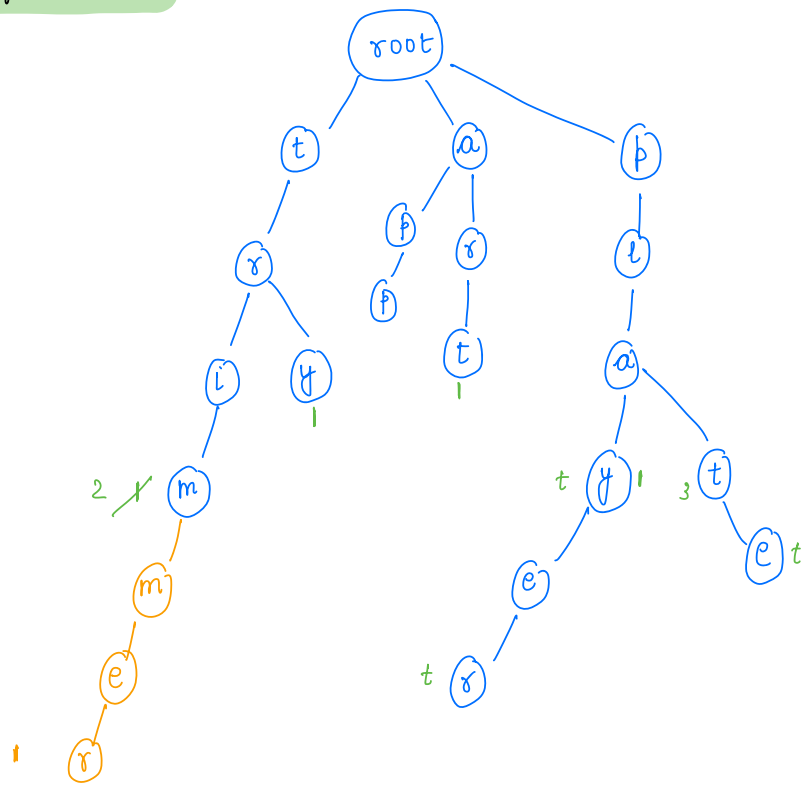
player

play

art

play

app



frequency(trim) = 2

freq(art) = 2

freq(play) =

freq(try) =

Structure:

```
class Node {  
    char data;  
    Node[] children;  
    boolean isEnd;  
    int freq;  
    Node(x) {  
        data = x;  
        children = new Node[26];  
        isEnd = false;  
        freq = 0;  
    }  
}
```

Insert

```
void insert(root, word) {  
    curr = root;  
    n = word.length;  
    for(i=0; i < n; i++) {  
        ch = word.charAt(i);  
        idx = ch - 'a';  
        if(curr.children[idx] == null) {  
            Node nn = new Node(ch);  
            curr.children[idx] = nn;  
        }  
        curr = curr.children[idx];  
    }  
    curr.isEnd = true;  
    curr.freq += 1;  
}
```

find frequency

```
int frequency (root, word) {  
    curr = root;  
    n = word.length;  
    for (i=0; i < n; i++) {  
        ch = word.charAt(i);  
        idx = ch - 'a';  
        if (curr.children[idx] == null) {  
            return 0;  
        }  
        curr = curr.children[idx];  
    }  
    if (curr.isEnd) {  
        return curr.freq;  $\implies$  return curr.freq  
    }  
    return 0;  
}
```

TC: $O(\text{len}) = 1 \text{ word}$.

Break: 8:19 - 8:30 AM

* Deletion in a tree

trim

try

art

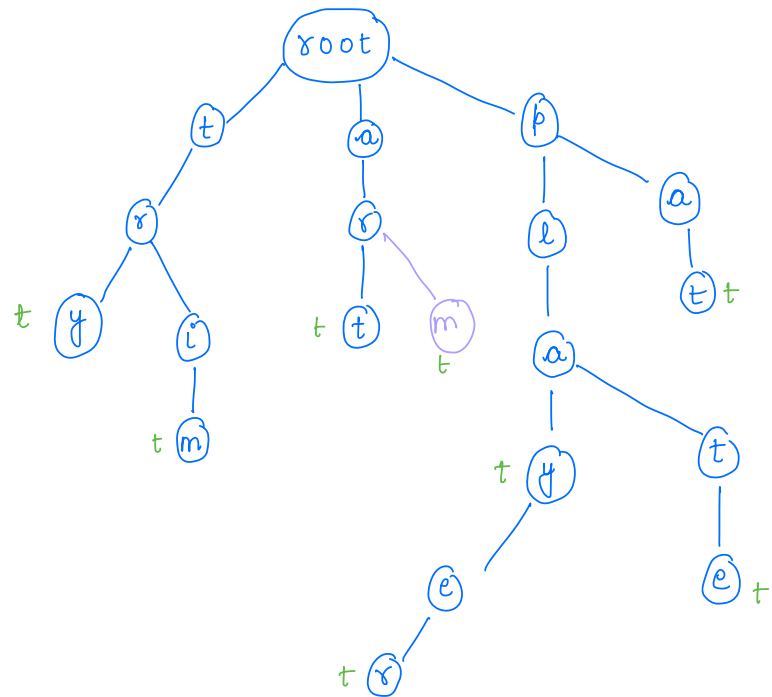
play

plate

player

arm

pot



delete(trim) =

delete(play) =

Ideal:

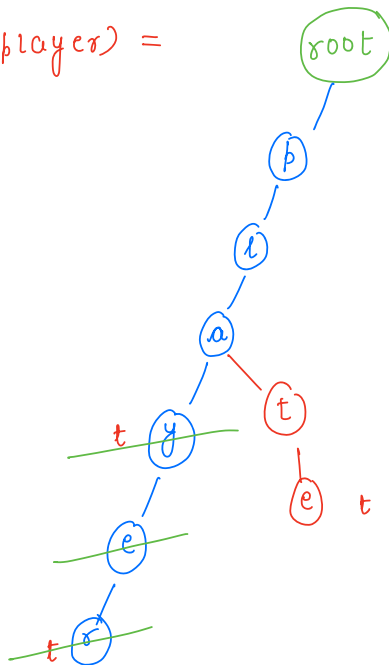
Traverse complete word and mark end of word = false.
 ↳ space is wasted

Idea 2

Reach end of word .
while coming back,
if leaf node - delete
else - stop.

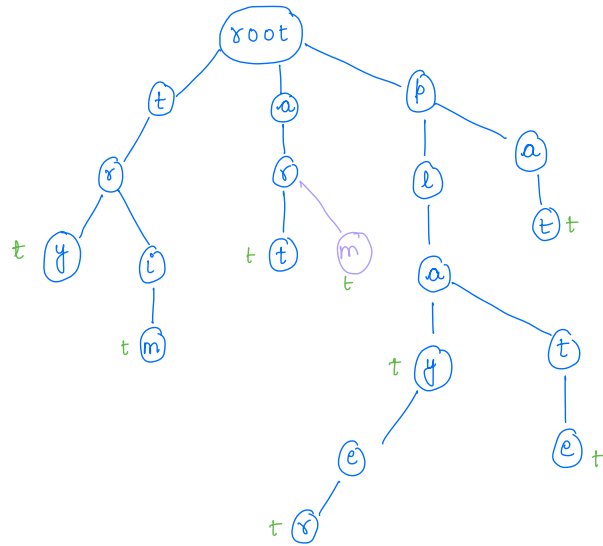
delete(trim) = works fine

delete(player) =



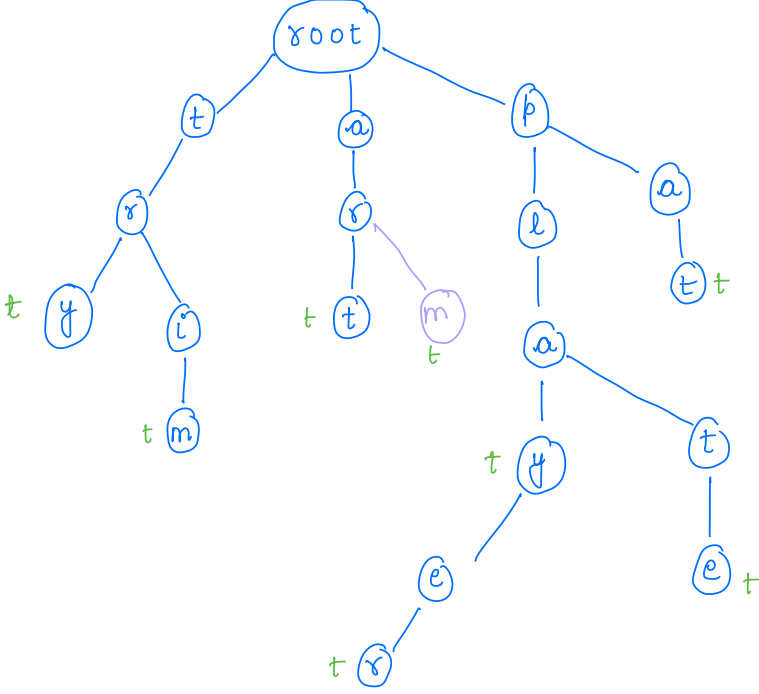
Modified idea

Reach end of word . —
while coming back —
if leaf node && isend = false.
delete it
else
stop



Idea3:

find last / latest node that can't be deleted.



delete(player)

Code:

```
void deleteWord(root, word) {
```

$$Cu\gamma\gamma = \sigma_{00}t;$$

```
lastNode = null;
```

deleteChar = ' ';

$n = \text{word.length}()$

```
for (i=0; i<n; i++) {
```

```
cnt = 0;
```

```
for(j=0; j<26; j++) {
```

```
if (curr.children[j] != null) {
```

cnt++;

```
if (cnt > 1 || curr.isEnd == true) {
```

```
lastNode = curr;
```

deleteChar = word.charAt(i)

}

```
idx = word.charAt(i) - 'a';
```

```
curr = curr.children[idx];
```

```
curr = i & end = false;
```

// find whether last character is a leaf

node or not?

```
cnt = 0;
```

```
for(j=0; j<26; j++) {
```

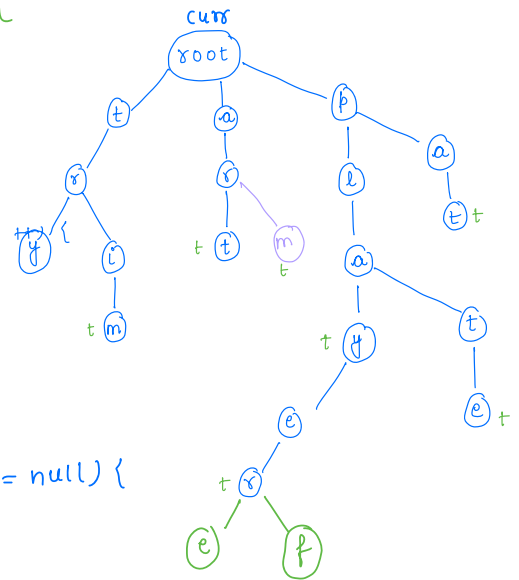
```
if (curr.children[j] != null) {
```

cnt++;

```
if(count > 0) { return; };
```

```
lastNode->children[deletechar - 'a'] = null;
```

SC: $O(1)$



Dry run:

delete (player) \rightarrow

$$\text{cur } \gamma = \text{root}$$

```
deletechar = ' ';
```

$i = 0$, $p \rightarrow$ curr = root
last = null
dc = ''

$$\text{last} = \text{root}$$
$$dc = p$$
$$\text{Cu}^{+} = \textcircled{p}$$
$$i^0 = 1, \quad \ell$$
$$\text{curl} \gamma = p$$
$$100t = p$$
$$dc = L$$
$$cu\gamma\gamma = 2$$

$i=2$, a

$$\text{curr} = l$$
$$\text{curr} = a$$

$i=3, y:$

$$\text{curr} = a$$

la stnode = (a)

$$dc = y.$$

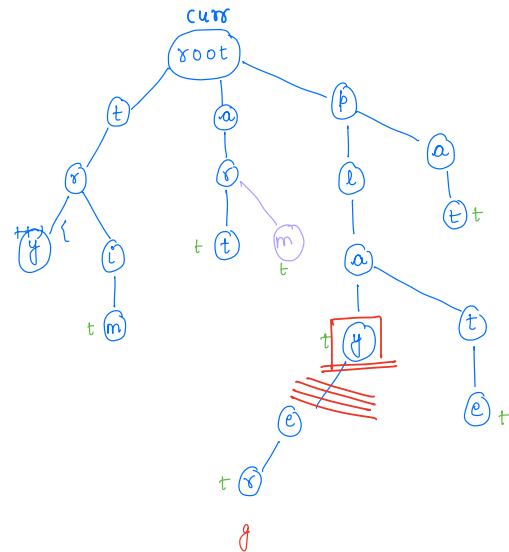
curr = y

 $i=4$ e:
$$\text{curr} = y$$

la stnode = (y)

$$dc = \underline{\underline{e}}$$
$$\text{curr} = (e)$$

$i = 5$ $\gamma :$

$$\text{curr} = (e)$$
$$\text{curr} = (\delta)$$


```
for (i=0; i<n; i++) {
```

```
cnt = 0;
```

```
for(j=0; j<26; j++){
```

```
if (curr.children[j] != null)
```

```
cnt++;
```

```
if (cnt > 1 || curr.isEnd == true) {
```

```
lastNode = curr;
```

```
deleteChar = word.charAt(i)
```

```
idx = word.charAt(i) - 'a';
```

```
curr = curr.children[idx];
```

0

```
lastNode->children[deletechar - 'a'] = null;
```

$$\textcircled{y} \cdot \begin{bmatrix} e \end{bmatrix} = \text{null}$$

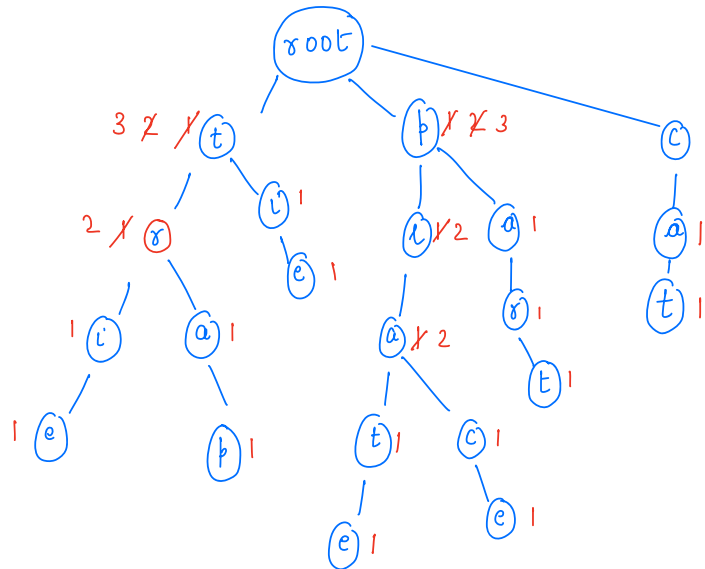
Qu

Shortest unique prefix

[tri[✓]e, trap, plat[✓]e, cat, part, plac[✓]e, tie]

↑ ↓ ↓ ↓ ↓ ↓ ↓

tri tra plat c pa plac ti



```

class Node {
    char data;
    Node[] children;
    boolean isEnd;
    int pc;
    Node(x) {
        data = x;
        children = new Node[26];
        isEnd = false;
        pc = 0;
    }
}

```

Insert

```
void insert(root, word) {  
    curr = root;  
    n = word.length;  
    for(i=0; i < n; i++) {  
        ch = word.charAt(i);  
        idx = ch - 'a';  
        if(curr.children[idx] == null) {  
            Node nn = new Node(ch);  
            curr.children[idx] = nn;  
        }  
        curr = curr.children[idx];  
        curr.pf += 1;  
    }  
    curr.isEnd = true;  
}
```

search shortest unique prefix

string shortestUniquePrefix
(root, word) {

curr = root;

n = word.length;

ans = "";

for (i = 0; i < n; i++) {

ch = word.charAt(i);

idx = ch - 'a';

if (curr.children[idx] == null) {

return "" / null; → depends on question

if (curr.children[idx].pf > 1) {

ans += ch;

} else { → curr.children[idx].pf == 1

ans += ch;

} return ans;

curr = curr.children[idx];

}

return ans;

}

Thankyou 😊