

Lecture: Advanced Recursion

Qul Given no. n , find sum of natural no from 1 to n .

$$\text{sum}(5) = \underbrace{1 + 2 + 3 + 4 + 5}_{\text{sum}(4)}$$

$$\begin{array}{ccc} \text{sum}(5) & = & \text{sum}(4) + 5. \\ \uparrow & & \uparrow \\ \text{problem} & & \text{sub-problem.} \end{array}$$

```
int sum(int n) {  
    if (n==1) {  
        return 1;  
    }  
    sa = sum(n-1);  
    return sa + n;  
}
```

Tc:

$$\text{sum}(5) \text{ ans} = 5 + 10 = 15$$

$$\downarrow \uparrow 10$$

$$\text{sum}(4) \text{ ans} = 4 + 6 = 10$$

$$\downarrow \uparrow 6$$

$$\text{sum}(3) \text{ ans} = 3 + 3 = 6$$

$$\downarrow \uparrow 3$$

$$\text{sum}(2) \text{ ans} = 2 + 1 = 3$$

$$\downarrow \uparrow 1$$

$$\text{sum}(1) \rightarrow 1$$

Assumption:

Given n , find and return
sum from 1 to n .

main logic

$$\text{sum}(n)$$



$$\text{ans} = \text{sa} + n.$$

$$\text{sa} = \text{sum}(n-1)$$

Base case

Smallest subproblem that

we can solve

$$n==1 \text{ return } 1.$$

$$n==0 \text{ return } 0$$

SC: h/w

Q42 fibonacci number : Given n. find and return nth fibonacci number

Eg: 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, ...

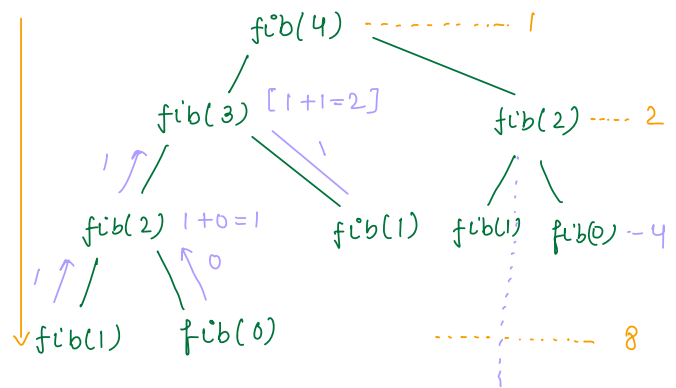
\uparrow \uparrow \uparrow \uparrow
 0th 1st 2nd 3rd -----

$$\underset{2}{\text{fib}(3)} = \underset{1}{\text{fib}(1)} + \underset{1}{\text{fib}(2)}$$

$$\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$$

```

int fib(int n) {
    if (n == 0 || n == 1) {
        return n;
    }
    int sa1 = fib(n-1);
    int sa2 = fib(n-2);
    return sa1 + sa2;
}
  
```



$$1 + 2 + 4 + 8 + \dots + 2^n$$

$$2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^n$$

GP

$$\simeq 2^n$$

TC: 2^n

SC: $O(1)$

Ques

```
void solve(int n) {
```

```
1. if (n==0) {
```

```
    return;
```

```
}
```

```
2. print(n);
```

```
3. solve(n-1);
```

```
}
```

Solve(-3) : a) -3 -2 -1

b) -1 -2 -3

c) stack overflow error

sol) -3 -2 -1 0

solve(-3)



solve(-4)



solve(-5)



solve(-6)



solve(-7)



solve(-8)



solve(-∞)

-3 -4 -5 -6

-7 -8

⇒ stored in stack
space

Qu4:

```
void solve(int n) {
```

```
1. if (n==0) {
```

```
    return;
```

```
}
```

```
2. print(n);
```

```
3. solve(n-1);
```

```
}
```

~~solve(3);~~

~~solve(2)~~

~~solve(1)~~

~~solve(0)~~

a) 1 2 3

b) 1 2 3 0

c) 3 2 1 0

d) 3 2 1

3 2 1

Qu5:

```
void solve(int n) {
```

```
1. if (n==0) {
```

```
    return;
```

```
}
```

```
2. solve(n-1);
```

```
3. print(n)
```

```
}
```

~~a) 1 2 3~~

b) 1 2 3 0

c) 3 2 1 0

d) 3 2 1

~~solve(3) | 2 3~~

~~solve(2) | 2 3~~

~~solve(1) | 2 3~~

~~solve(0)~~

1 2 3

Ques

Tower of Hanoi

There are n disks placed on tower A of different sizes.

Goal: Move all disks from tower A to tower B using tower C.

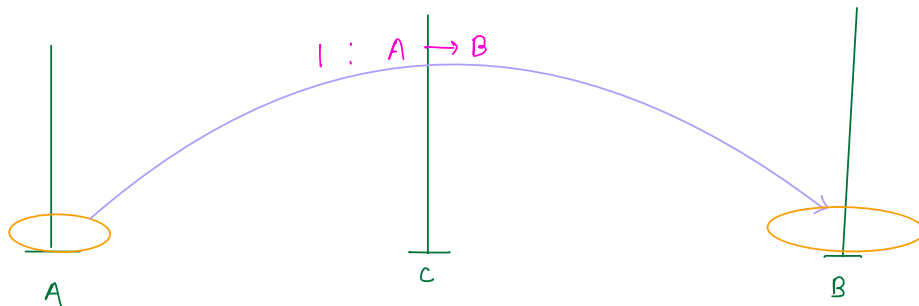
Constraint:

Only 1 disk can be moved at a time.

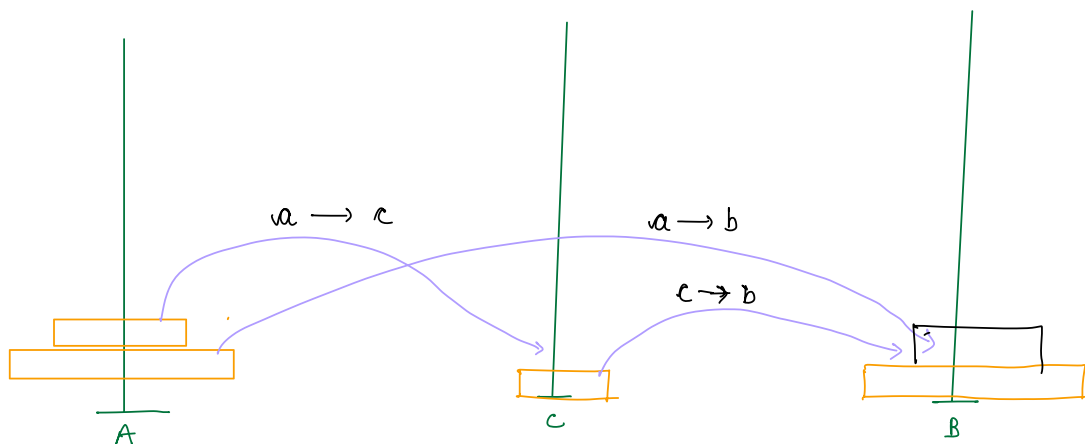
Large disk cannot be placed on smaller disk.

Problem: print the moves such that all disks move from tower A to B.

Eg: $n=1$

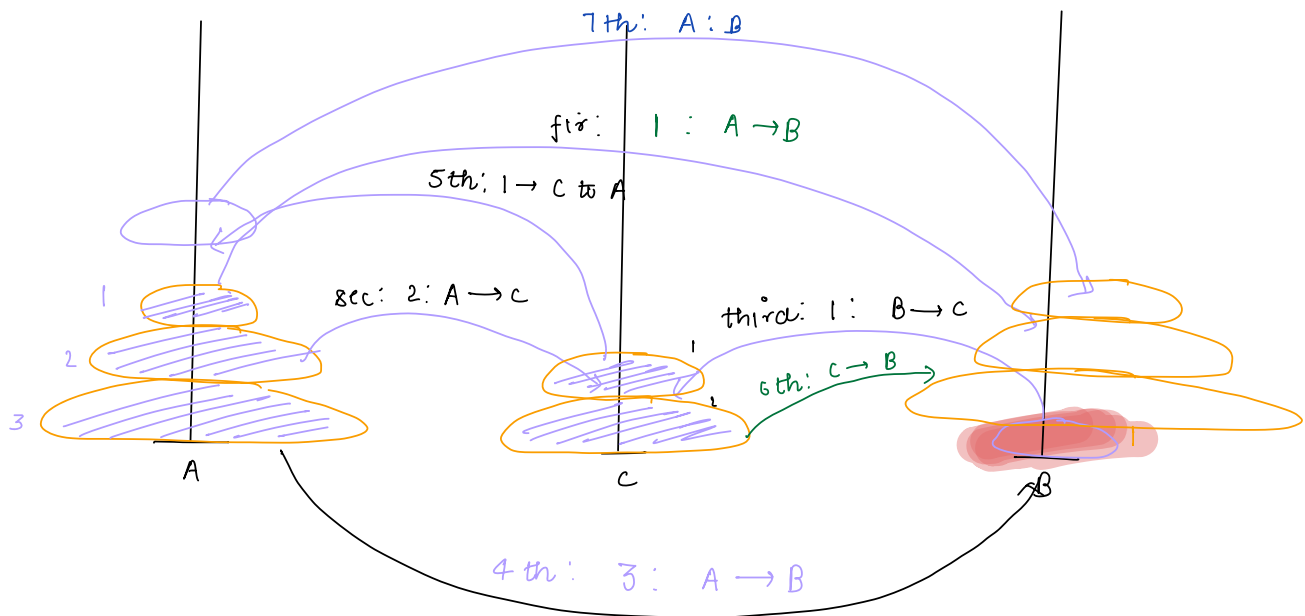


$n=2$



$n=3$

$$\left[\begin{array}{l} n=2: \text{source: } A \quad \text{dest: } C \quad \text{aux: } B \\ n=2: \text{source: } C \quad \text{dest: } B \quad \text{aux: } A \end{array} \right]$$



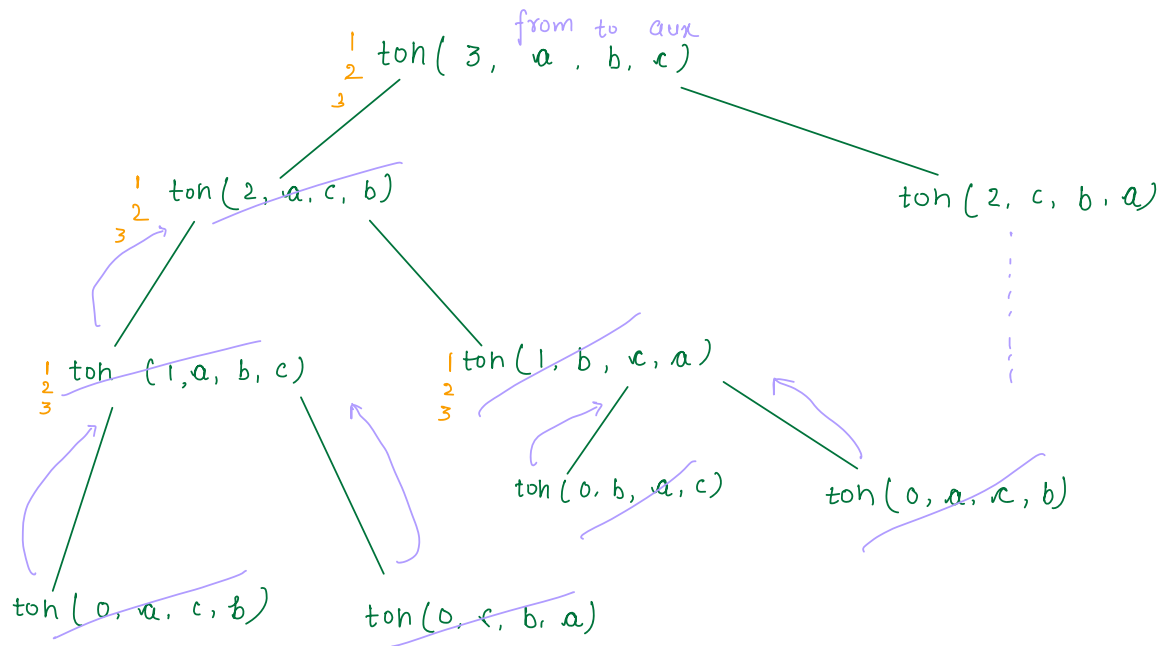
Pseudocode:

```

void towerofHanoi(int n, char from, char to, char aux) {
    if (n == 0) {
        return;
    }
    1 towerofHanoi(n-1, from, aux, to);
    2 print(n + " : " + from + "  $\rightarrow$  " + to);
    3 towerofHanoi(n-1, aux, to, from);
}
  
```

Dry run:

```
void towerofHanoi(int n, char from, char to, char aux) {
    if (n == 0) {
        return;
    }
    1 towerofHanoi(n-1, from, aux, to);
    2 print(n + ":" + from + " → " + to);
    3 towerofHanoi(n-1, aux, to, from);
}
```



prints: 1 : a → b

2 : a → c

1 : b → c

3 : a → b

TC: $n=1 [1] \rightarrow 2^1 - 1$: $O(2^n)$

$n=2 [3] \rightarrow 2^2 - 1$

$n=3 [7] \rightarrow 2^3 - 1$

\vdots

$n [2^n - 1]$

SC: h/w

Break: 8:33 AM

Q: Generate parenthesis

Given no. n , Generate all valid parenthesis of sequence len: $2 * n$

Eg: $n=1$ () (()) ()
Invalid

$n=2$ (())
() ()

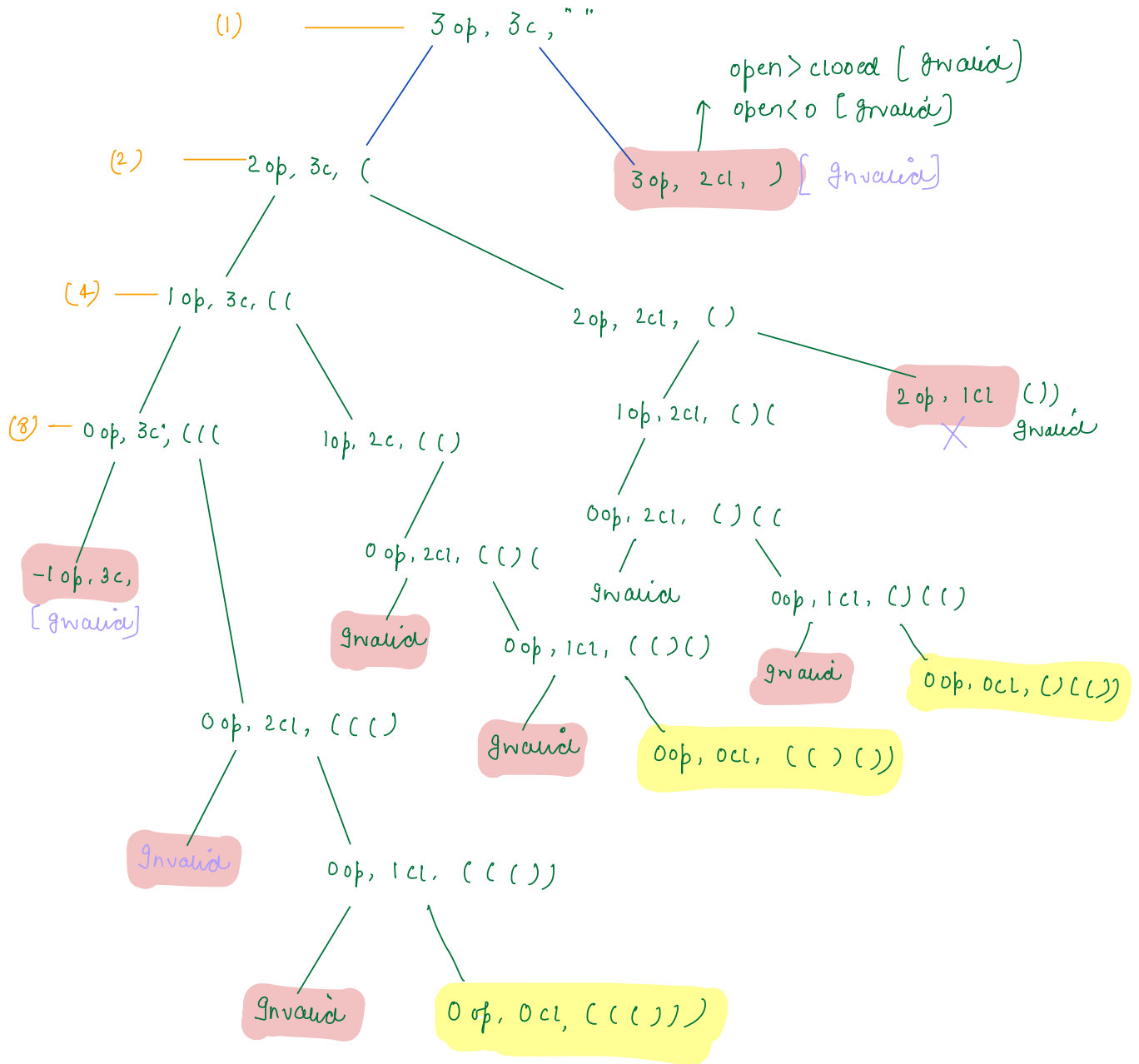
$n=3$ ((()))
(()) ()
() () ()
() (())

Brute force: Generate all parenthesis.

$n=2$. [(())]

↑
Generate all subsets [Backtracking]

↓
Check for valid ones [stack].

$$n=3 \quad [3 \text{ open, } 3 \text{ closed}]$$


```
void generate (int n) {
    generate (n, n, "");
}
```

```
void generate (int open, int closed, string seq) {
    if (open == 0 && closed == 0)
        print (seq);
        return;
    }
    if (open > 0) { // left part: deal with (
        generate (open-1, closed, seq + '(');
    }
    if (open < closed) { // right part: )
        generate (open, closed-1, seq + ')');
    }
}
```

TC: $O(2^n)$

SC: n/w

Thankyou 😊

Doubt:

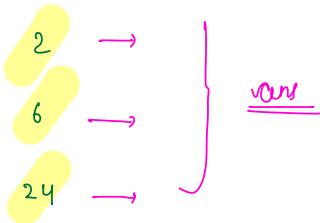
$$A = [2, 3, 4]$$

$$B[] = [2, 6, 24]$$

↑
factorial

↑
total no of ^{non-empty} subsequences such that
each el in subsequence has same prime factors.

$$B[] = [2, 6, 24]$$



✗ 2 6 → 2: 2 ✗
6: 2, 3

6 24

✗ 2 24 → 6: 2 3 ✓
24: 2 3

✗ 2 6 24

a b c

L c

b

a

a b

a c

b c

a b c