

Lecture ÷ Graph I

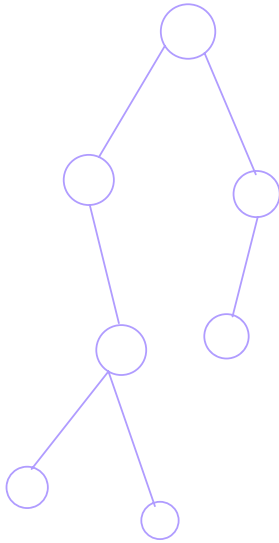
Agenda

- Introduction to graphs
- classification and inputs
- Store a graph
- Bfs and Dfs.

Graph:

Collection of nodes connected to each other using edges.

[Tree] \rightarrow Graph



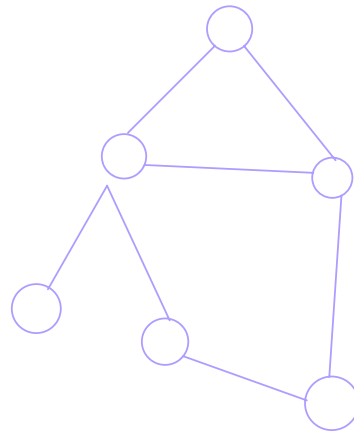
$$n = 7$$

$$\text{edges} = 6$$



$$\text{No of edges} = n - 1$$

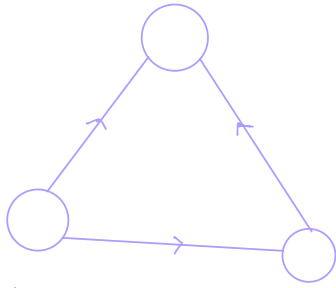
Graph



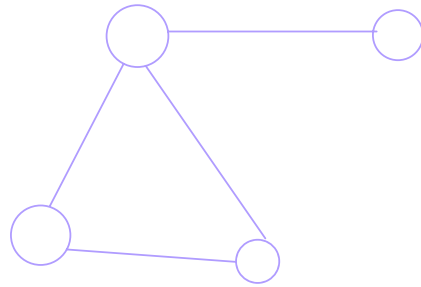
$$n = 6$$

$$\text{edges} = 7$$

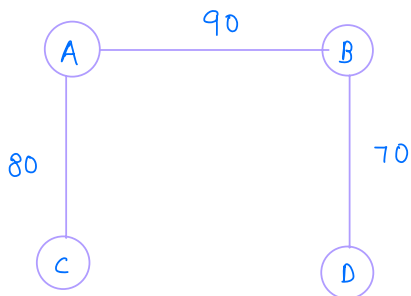
Classification of graphs.



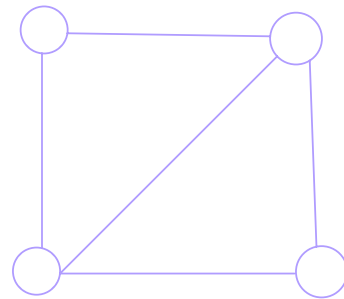
Directed graph
Instagram



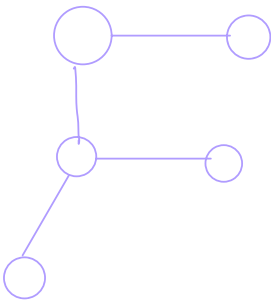
Undirected graph
facebook



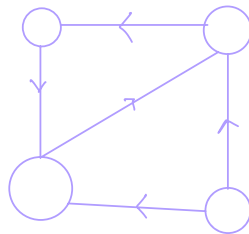
Weighted undirected.
Maps



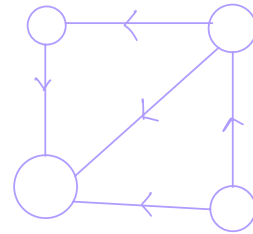
Undirected cyclic graph



Undirected **A**cylic graph



Directed cyclic



Directed acyclic

Graph input

Qn. Given an undirected graph with n nodes and m edges.

Input: 1st line: n and m .
 ↓ ↓
 no of nodes no of edges

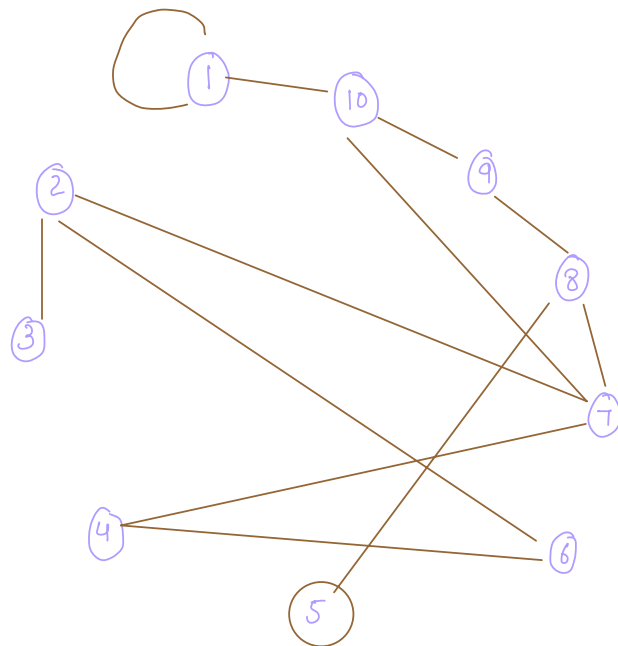
followed by m lines:

Each line contains u and v indicating an edge from u to v .

Eg: n m
 10 12

u	v
2	3
4	7
8	9
2	7
7	8
10	1
4	6
5	8
2	6
10	9
7	10
1	1

Visualisation

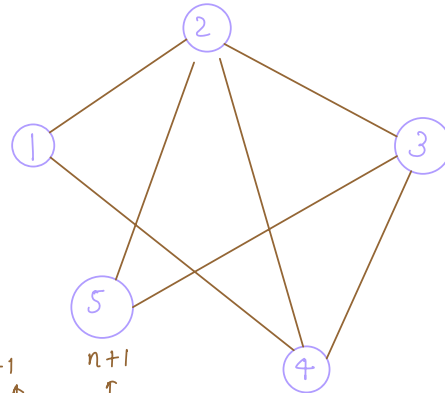


Approach 1 Adjacency matrix

ip $n = 5$, $m = 7$

Edges:

u	v
1	2
1	4
2	3
2	4
2	5
3	4
3	5



$\overset{n+1}{\uparrow}$ $\overset{n+1}{\uparrow}$
 $\text{mat}[6][6]$

u	v
1	2
1	4
2	3
2	4
2	5
3	4
3	5

	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	0	0 1	0	0 1	0
2	0	0 1	0	0 1	0 1	0 1
3	0	0	0 1	0	0 1	0 1
4	0	0 1	0 1	0 1	0	0
5	0	0	0 1	0 1	0	0

$\hookrightarrow \text{mat}[i][j] = 1 \{ \text{edge b/w } i \text{ and } j \}$

	Unweighted	Weighted
Undirected	$\text{mat}[u][v] = 1$ $\text{mat}[v][u] = 1$	$\text{mat}[u][v] = wt$ $\text{mat}[v][u] = wt$
Directed	$\text{mat}[u][v] = 1$	$\text{mat}[u][v] = wt$

Code:

```
int[][] buildGraph(n, m, edges[][]) {  
    mat[n+1][n+1];  
    for(i=0; i< m ; i++){  
        u = edges[i][0];  
        v = edges[i][1];  
        mat[u][v] = 1;  
        mat[v][u] = 1;  
    }  
    return mat;  
}
```

TC: $O(\text{edges})$

SC: $O(n*n) \approx O(n^2)$

└ Space wasted

Confusion:

$\text{mat}[u][v] = \text{wt}$

└ u & v have an edge

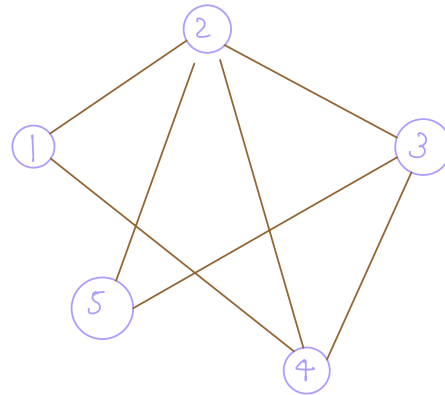
└ wt b/w u and v is wt.

Approach 2 Adjacency list

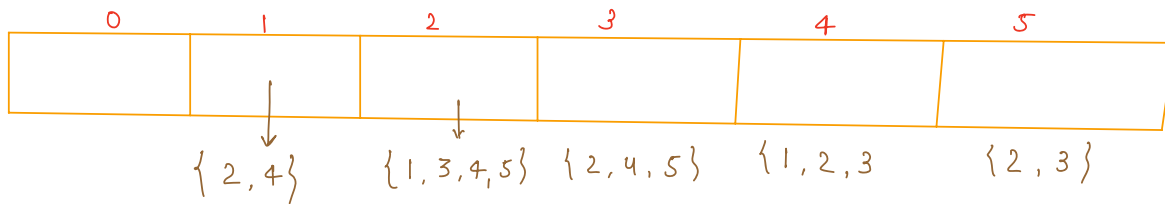
ip $n = 5$, $m = 7$

Edges:

u	v
1	2
1	4
2	3
2	4
2	5
3	4
3	5



List<Integer> $l[n+1]$



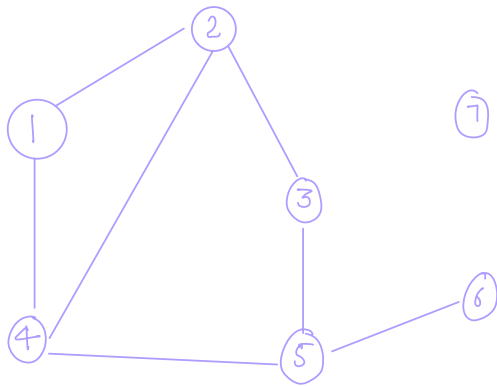
	Unweighted	Weighted
Undirected	$l[u].add(v)$ $l[v].add(u)$	$l[u].add(pair(v, wt))$ $l[v].add(pair(u, wt))$
Directed	$l[u].add(v)$	$l[u].add(pair(v, wt))$

code:

```
List<Integer>[] buildGraph(n, m, edges[][]) {  
    List<Integer>[] l;  
    for (i = 0; i < m; i++) {  
        u = edges[i][0];  
        v = edges[i][1];  
        l[u].add(v);  
        l[v].add(u);  
    }  
    return l;  
}
```

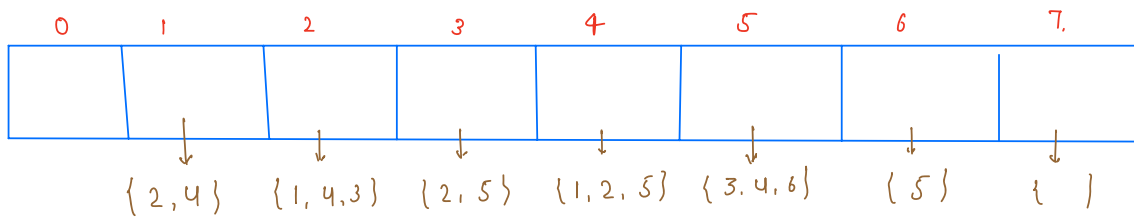

Bfs

Q: Given an undirected graph and source node and destination node. Check if destination node can be visited from source node.



s	d	
1	6	true
1	7	false

Graph:



Approach:

source = 1, destination = 6

Queue:

| 1 2 3 4 5 6 |

visited[] =

0	1	2	3	4	5	6	7
f	f +	f +	f +	f +	f +	f +	f

0	1	2	3	4	5	6	7
	↓	↓	↓	↓	↓	↓	↓
	{ 2, 4 }	{ 1, 4, 3 }	{ 2, 5 }	{ 1, 2, 5 }	{ 3, 4, 6 }	{ 5 }	{ }

Code:

```
boolean bfs(n, m, edges[][], src, dest) {  
    List<Integer>[] graph = buildGraph(n, m, edges);  
    boolean[] vis = new boolean[n+1];  
    Queue<Integer> q;  
    q.add(src);  
    vis[src] = true;  
    while( ! q.isEmpty() ) {  
        curr = q.poll();  
        List<Integer> nbrs = graph[curr];  
        for(int el: nbrs) {  
            if(vis[el] == false) {  
                q.add(el);  
                vis[el] = true;  
            }  
        }  
    }  
    return vis[dest];  
}
```

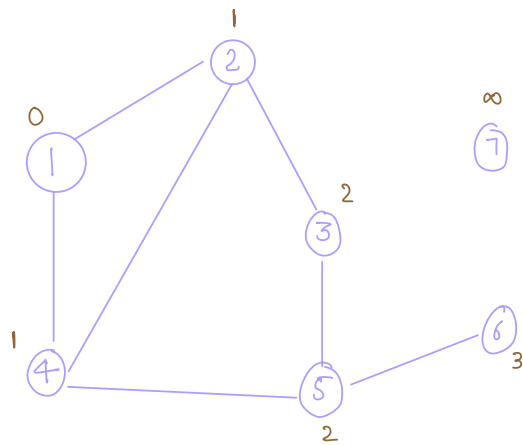
TC: $O(n + e)$ \approx Linear

Sc: $O(n + \underbrace{n + e})$

↑
Queue
size

↑
Adjacency list

Qu: find distance from source node to destination node



s	d	distance
1	6	3
1	7	∞

Queue:



	0	1	2	3	4	5	6	7
visited[] =	f	f t	f t	f t	f t	f t	f t	f

	0	1	2	3	4	5	6	7
dist[] =	∞	∞ 0	∞ 1	∞ 2	∞ 1	∞ 2	∞ 3	∞

0	1	2	3	4	5	6	7
	↓	↓	↓	↓	↓	↓	↓
	{ 2, 4 }	{ 1, 4, 3 }	{ 2, 5 }	{ 1, 2, 5 }	{ 3, 4, 6 }	{ 5 }	{ }

Code:

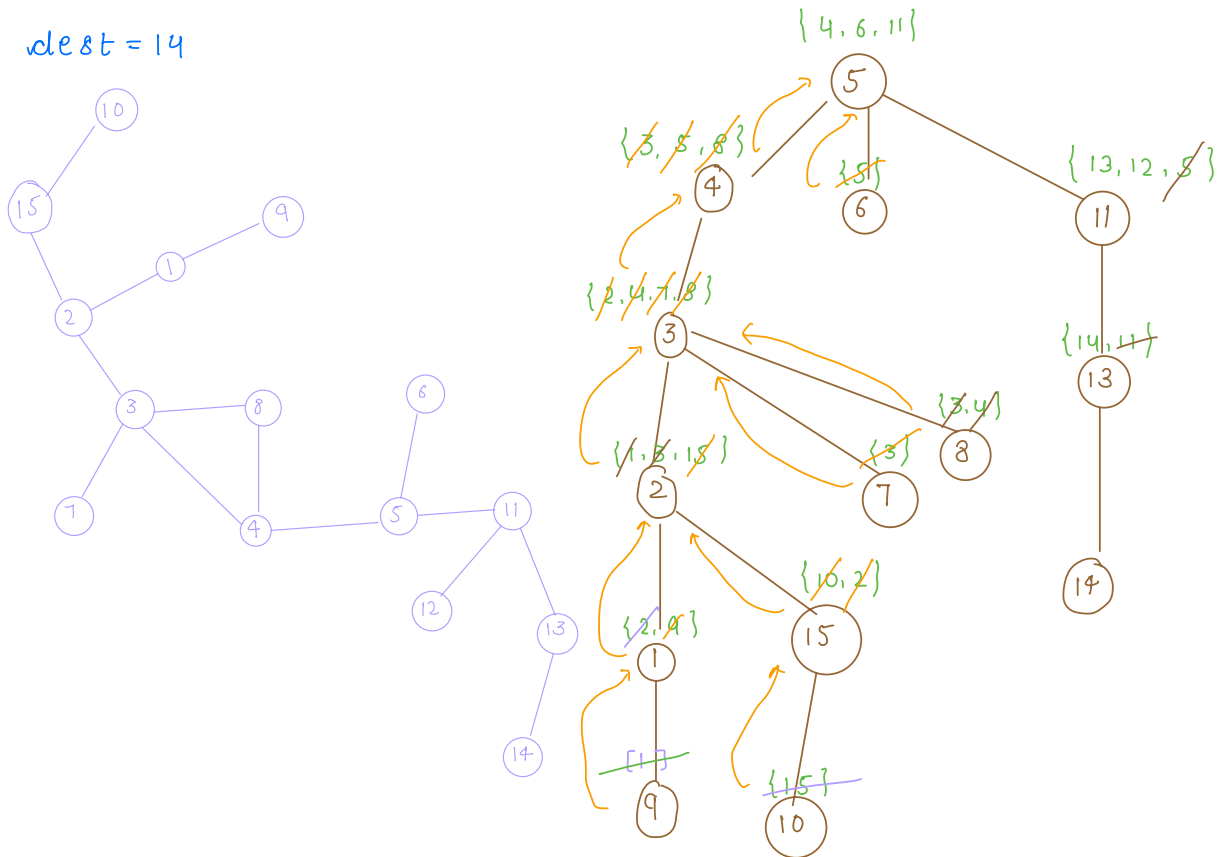
```
int getDistance(n, m, edges[][] , src, dest) {  
  
    List<Integer>[] graph = buildGraph(n, m, edges);  
  
    boolean[] vis = new boolean[n+1];  
    int[] dis = new int[n+1]; // Initialise as ∞.  
    Queue<Integer> q;  
  
    q.add(src);  
  
    vis[src] = true;  
    dis[src] = 0;  
    while( ! q.isEmpty() ) {  
  
        curr = q.poll();  
  
        List<Integer> neighbours = graph[curr];  
  
        for(int el: neighbours) {  
            if(vis[el] == false) {  
  
                q.add(el);  
  
                vis[el] = true;  
                dis[el] = dis[curr] + 1;  
            }  
        }  
    }  
  
    return dis[dest];  
}
```

Break: 8:30-8:40

Dfs [Depth first search]

src = 5

dest = 14



code:

```
boolean isVisited(n, m, edges[][] , src, dest) {  
    List<Integer>[] graph = buildGraph(n, m, edges);  
    boolean[] vis = new boolean[n+1];  
    dfs(src, graph, vis);  
    return vis[dest];  
}
```

```
void dfs(s, graph, vis) {  
    if (vis[s] == true) {  
        return;  
    }  
    vis[s] = true;  
    List<Integer> neighbours = graph[s];  
    for (int el: neighbours) {  
        if (vis[el] == false) {  
            dfs(el, graph, vis);  
        }  
    }  
}
```

TC: $O(n+e)$

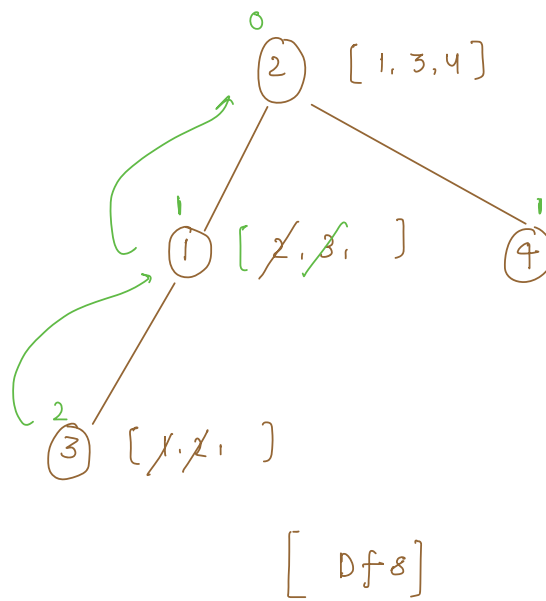
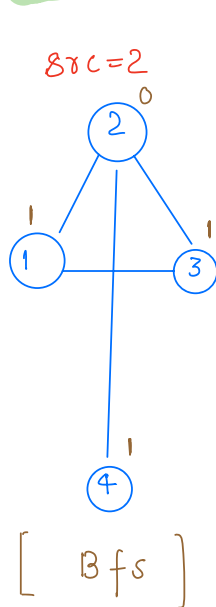
SC: $O(n+e) +$

↑
Adjacency
list

↑
stack space

$+ O(n)$
↑
vis[]

Shortest path from src to destination



Observation

1. Shortest path \longrightarrow Bfs
2. " " \longrightarrow cannot be calculated using Dfs

Thankyou 😊