

Data Management Assignment 5

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- A.**
1. (a) Transaction 4 (T4) is needs to be redone, it committed after the last checkpoint and therefore this transaction must be reflected in the final state.
 - (b) Transactions 2 and 3 started before the checkpoint but did not commit, they have to be undone. As well as transactions T5 and T6 who did not commit and therefore need to be undone. T6 has yet made a operation as in our system the log is ahead of anything, so theoretically you can ignore 6 as nothing can be undone.
 - (c) T1 and T2 have committed before the last checkpoint, thus they are not affected by the crash.
 2. (a) The possible values of a are 1 or 2, before the last checkpoint it was changed to 1 and after it it was again changed to 2.
The possible values for b are 2 or 3. The checkpoint insures that after it the disk contains 2 and then b is altered to 3, which we do not know if it actually gets written to the disk.
Finally, c can be 0 or 1 as it is changed only after the checkpoint and again we cannot know whether the value 1 was written on the disk before the crash.
 - (b) T1 is ignored, T2 needs to be redone and transaction 3 and 4 undone. Therefore the values of a, b, c are: a=2, b=2, c=0.
- B.**
1. (a) A database history is recoverable if every transaction commits only after all transactions from whom values are read, are already committed. More precisely, the commit of a transaction has to be after all commits of the transaction from whom values are read during this transaction.
 - (b) A database history is cascadeless if every transaction only reads values from other transaction who have already committed. If a transaction reads a value who was written by a transaction that has not committed then it is not cascadeless.

- (c) The schedule is strict if a transaction only changes(writes) values who have already committed in all transaction who write to this value.

- 2. (a)
- (b)

- C.** (a)
- (b)