EXPERIMENT NO - 3

Implementation of Lexical Analyzer

AIM:

To implement lexical analyzer to check if a number is odd or even.

CODE:

```
#include<stdio.h>
#include<stdlib.h>
#include<string.h>
#include<ctype.h>
int isKeyword(char buffer[]){
  char keywords[32][10] =
{"auto","break","case","char","const","continue","default","do","double","else","enum","extern","fl
oat", "for", "goto", "if", "int", "long", "register", "return", "short", "signed", "sizeof", "static", "struct", "switc
h","typedef","union","unsigned","void","volatile","while"};
  int i, flag = 0;
  for(i = 0; i < 32; ++i){
    if(strcmp(keywords[i], buffer) == 0){
      flag = 1;
      break;
    }
  }
  return flag;
}
int main(){
  char ch, buffer[15], operators[] = "+-*/\%=";
  FILE *fp;
  int i,j=0;
  fp = fopen("code.txt","r");
  if(fp == NULL){
```

```
printf("error while opening the file\n");
    exit(0);
  }
  while((ch = fgetc(fp)) != EOF){
    for(i = 0; i < 6; ++i){
      if(ch == operators[i])
         printf("%c is operator\n", ch);
    }
    if(isalnum(ch)){
      buffer[j++] = ch;
    }
    else if((ch == ' ' | | ch == '\n') && (j != 0)){
      buffer[j] = '\0';
      j = 0;
      if(isKeyword(buffer) == 1)
         printf("%s is keyword\n", buffer);
      else
         printf("%s is indentifier\n", buffer);
    }
  }
  fclose(fp);
  return 0;
}
File:
#include<bits/stdc++.h>
using namespace std;
int main(){
  int a;
  if(a%2==0)
    cout<<"even";
  else cout<<"odd";
```

```
return 0;
```

```
/ is operator
+ is operator
+ is operator
includebitsstdch is indentifier
using is indentifier
namespace is indentifier
std is indentifier
int is keyword
main is indentifier
int is keyword
a is indentifier
% is operator
= is operator
= is operator
ifa20 is indentifier
couteven is indentifier
else is keyword
coutodd is indentifier
return is keyword
0 is indentifier
*** stack smashing detected ***: terminated
```

RESULT:

The above program was created and executed successfully.

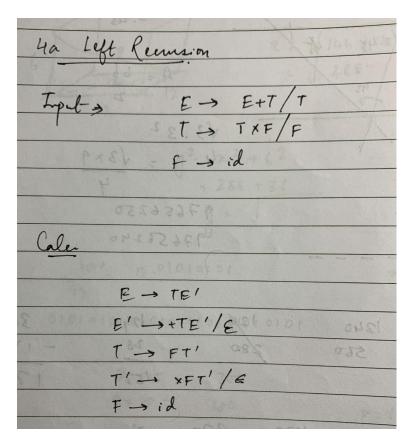
EXPERIMENT NO - 4a

Left Recursion Elimination

AIM:

To write a program to eliminate left recursion from the given grammar using manual calculation and using a C++ program.

MANUAL CALCULATION:



CODE:

```
#include<iostream>
#include<string>
#include<vector>
using namespace std;
int main()
{
   vector<string> ans;
   for(int tt=0;tt<3;tt++)</pre>
```

```
{
cout<<"re "<<tt+1<<"\n";
string ip,op1,op2,temp;
int sizes[10] = {};
char c;
int n,j,l;
cout<<"Enter the Parent Non-Terminal : ";</pre>
cin>>c;
ip.push_back(c);
op1 += ip + "\'->";
ip += "->";
op2+=ip;
cout<<"Enter the number of productions : ";</pre>
cin>>n;
for(int i=0;i<n;i++)
{ cout<<"Enter Production "<<i+1<<":";
  cin>>temp;
  sizes[i] = temp.size();
  ip+=temp;
  if(i!=n-1)
    ip += "|";
}
// cout<<"Production Rule : "<<ip<<endl;</pre>
for(int i=0,k=3;i<n;i++)
{
  if(ip[0] == ip[k])
  {
    // cout<<"Production "<<i+1<<" has left recursion."<<endl;
    if(ip[k] != '#')
    {
       for(l=k+1;l<k+sizes[i];l++)
```

```
op1.push_back(ip[l]);
       k=l+1;
       op1.push_back(ip[0]);
       op1 += "\'|";
    }
  }
  else
  {
    // cout<<"Production "<<i+1<<" does not have left recursion."<<endl;
    if(ip[k] != '#')
    {
      for(j=k;j<k+sizes[i];j++)</pre>
         op2.push_back(ip[j]);
       k=j+1;
      op2.push_back(ip[0]);
      op2 += "\'|";
    }
    else
    {
      op2.push_back(ip[0]);
      op2 += "\"";
    }}}
op1 += "#";
ans.push_back(op2);
ans.push_back(op1);
// cout<<op2<<endl;</pre>
// cout<<op1<<endl;
}
for(int i=0;i<ans.size();i++)</pre>
  cout<<ans[i]<<"\n";
return 0;
```

```
Enter the Parent Non-Terminal : E
Enter the number of productions : 2
Enter Production 1 : E+T
Enter Production 2 : T
re 2
Enter the Parent Non-Terminal : T
Enter the number of productions : 2
Enter Production 1 : TXF
Enter Production 2 : F
re 3
Enter the Parent Non-Terminal : F
Enter the number of productions : 1
Enter Production 1 : id
E->TE'
E'->+TE' | #
T->FT'
T'->XFT'|#
F->idF'
F'->#
```

RESULT:

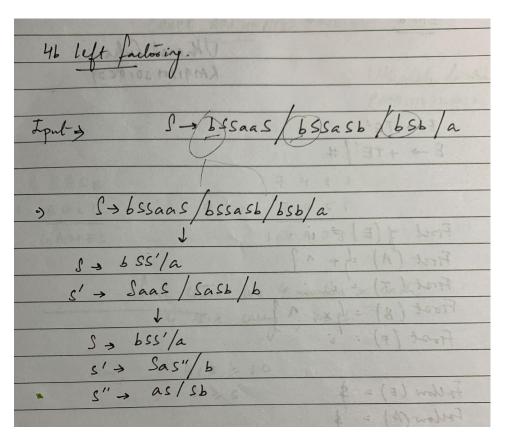
Left Recursion Elimination was successfully calculated manually and the program was successfully executed using an Online C++ compiler.

EXPERIMENT NO - 4b Left Factoring

AIM:

To write a program for performing left factoring to the given grammar using manual calculation and using a C++ program.

MANUAL CALCULATION:



CODE:

```
#include <iostream>
#include <string>
using namespace std;
int main()
{
    int n,j,l,i,m;
    int len[10] = {};
    string a, b1, b2, flag;
    char c;
```

```
cout << "Enter the Parent Non-Terminal : ";</pre>
cin >> c;
a.push_back(c);
b1 += a + "\'->";
b2 += a + "\'\'->";;
a += "->";
cout << "Enter total number of productions : ";</pre>
cin >> n;
for (i = 0; i < n; i++)
{
  cout << "Enter the Production " << i + 1 << " : ";
  cin >> flag;
  len[i] = flag.size();
  a += flag;
  if (i != n - 1)
  a += "|";
              }
}
cout << "The Production Rule is : " << a << endl;</pre>
char x = a[3];
for (i = 0, m = 3; i < n; i++)
{
  if (x != a[m])
  {
     while (a[m++] != '|');
  }
  else
  {
       if (a[m + 1] != '|')
    {
```

```
b1 += "|" + a.substr(m + 1, len[i] - 1);
       a.erase(m - 1, len[i] + 1);
    }
    else
    {
      b1 += "#";
       a.insert(m + 1, 1, a[0]);
       a.insert(m + 2, 1, '\");
       m += 4;
    }
  }
}
char y = b1[6];
for (i = 0, m = 6; i < n - 1; i++)
{
  if (y == b1[m])
  {
    if (b1[m + 1] != '|')
    {
       flag.clear();
       for (int s = m + 1; s < b1.length(); s++)
          flag.push_back(b1[s]);
                                }
       b2 += "|" + flag;
       b1.erase(m - 1, flag.length() + 2);
    }
    else
    {
      b1.insert(m + 1, 1, b1[0]);
       b1.insert(m + 2, 2, '\");
```

```
Enter the Parent Non-Terminal : S
Enter total number of productions : 4
Enter the Production 1 : bSSaaS
Enter the Production 2 : bSSaSb
Enter the Production 3 : bSb
Enter the Production 4 : a
The Production Rule is : S->bSSaaS|bSSaSb|bSb|a
After Left Factoring :
S-|a
S'->|
S''->|aaS|SSaSb|Sb|
```

RESULT:

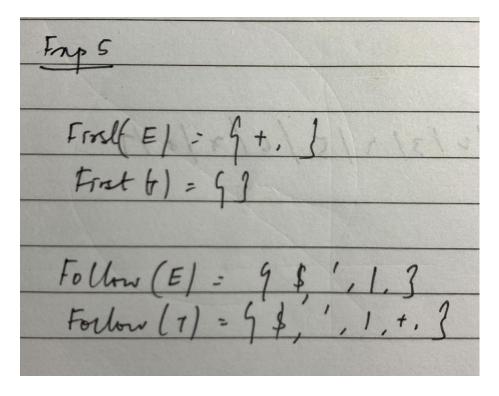
Left Factoring was successfully calculated manually and the program was successfully executed using an Online C++ compiler.

EXPERIMENT NO - 5 FIRST AND FOLLOW computation

AIM:

To write a program for finding first and follow for the given grammar using manual calculation and using a C++ program.

MANUAL CALCULATION:



CODE:

#include<stdio.h>
#include<ctype.h>
#include<string.h>

// Functions to calculate Follow
void followfirst(char, int, int);
void follow(char c);

// Function to calculate First
void findfirst(char, int, int);

```
int count, n = 0;
// Stores the final result
// of the First Sets
char calc_first[10][100];
// Stores the final result
// of the Follow Sets
char calc_follow[10][100];
int m = 0;
// Stores the production rules
char production[10][10];
char f[10], first[10];
int k;
char ck;
int e;
int main(int argc, char **argv)
{
        int jm = 0;
        int km = 0;
        int i, choice;
        char c, ch;
        count = 8;
        // The Input grammar
        strcpy(production[0], "S=01A");
        strcpy(production[1], "A=0S1SA");
        strcpy(production[2], "A=#");
```

```
int kay;
char done[count];
int ptr = -1;
// Initializing the calc_first array
for(k = 0; k < count; k++) {
        for(kay = 0; kay < 100; kay++) {
                 calc_first[k][kay] = '!';
        }
}
int point1 = 0, point2, xxx;
for(k = 0; k < count; k++)
{
        c = production[k][0];
        point2 = 0;
        xxx = 0;
        // Checking if First of c has
        // already been calculated
        for(kay = 0; kay <= ptr; kay++)
                 if(c == done[kay])
                         xxx = 1;
        if (xxx == 1)
                 continue;
        // Function call
        findfirst(c, 0, 0);
        ptr += 1;
```

```
// Adding c to the calculated list
        done[ptr] = c;
        printf("\n First(%c) = { ", c);
        calc_first[point1][point2++] = c;
        // Printing the First Sets of the grammar
        for(i = 0 + jm; i < n; i++) {
                int lark = 0, chk = 0;
                for(lark = 0; lark < point2; lark++) {</pre>
                        if (first[i] == calc_first[point1][lark])
                        {
                                 chk = 1;
                                 break;
                        }
                }
                if(chk == 0)
                {
                         printf("%c, ", first[i]);
                        calc_first[point1][point2++] = first[i];
                }
        }
        printf("}\n");
        jm = n;
        point1++;
printf("\n");
printf("-----\n\n");
char donee[count];
```

}

```
ptr = -1;
// Initializing the calc_follow array
for(k = 0; k < count; k++) {
        for(kay = 0; kay < 100; kay++) {
                 calc_follow[k][kay] = '!';
        }
}
point1 = 0;
int land = 0;
for(e = 0; e < count; e++)
{
        ck = production[e][0];
        point2 = 0;
        xxx = 0;
        // Checking if Follow of ck
        // has alredy been calculated
        for(kay = 0; kay <= ptr; kay++)</pre>
                 if(ck == donee[kay])
                         xxx = 1;
        if (xxx == 1)
                 continue;
        land += 1;
        // Function call
        follow(ck);
        ptr += 1;
        // Adding ck to the calculated list
```

```
donee[ptr] = ck;
                 printf(" Follow(%c) = { ", ck);
                 calc_follow[point1][point2++] = ck;
                 // Printing the Follow Sets of the grammar
                 for(i = 0 + km; i < m; i++) {
                         int lark = 0, chk = 0;
                         for(lark = 0; lark < point2; lark++)</pre>
                          {
                                  if (f[i] == calc_follow[point1][lark])
                                  {
                                           chk = 1;
                                           break;
                                  }
                          }
                         if(chk == 0)
                          {
                                  printf("%c, ", f[i]);
                                  calc_follow[point1][point2++] = f[i];
                          }
                 }
                 printf(" \n\n");
                 km = m;
                 point1++;
        }
}
void follow(char c)
{
        int i, j;
```

```
// set of the start symbol
        if(production[0][0] == c) {
                 f[m++] = '$';
        }
        for(i = 0; i < 10; i++)
        {
                 for(j = 2; j < 10; j++)
                 {
                          if(production[i][j] == c)
                          {
                                  if(production[i][j+1] != '\0')
                                  {
                                           // Calculate the first of the next
                                           // Non-Terminal in the production
                                           followfirst(production[i][j+1], i, (j+2));
                                  }
                                  if(production[i][j+1]=='\0' && c!=production[i][0])
                                  {
                                           // Calculate the follow of the Non-Terminal
                                           // in the L.H.S. of the production
                                           follow(production[i][0]);
                                  }
                         }
                 }
        }
}
void findfirst(char c, int q1, int q2)
{
```

// Adding "\$" to the follow

```
int j;
// The case where we
// encounter a Terminal
if(!(isupper(c))) {
        first[n++] = c;
}
for(j = 0; j < count; j++)
{
        if(production[j][0] == c)
        {
                if(production[j][2] == '#')
                {
                         if(production[q1][q2] == '\0')
                                 first[n++] = '#';
                         else if(production[q1][q2] != '\0'
                                          && (q1 != 0 || q2 != 0))
                         {
                                 // Recursion to calculate First of New
                                 // Non-Terminal we encounter after epsilon
                                 findfirst(production[q1][q2], q1, (q2+1));
                         }
                         else
                                 first[n++] = '#';
                }
                else if(!isupper(production[j][2]))
                {
                         first[n++] = production[j][2];
                }
                else
                {
```

```
// Recursion to calculate First of
                                  // New Non-Terminal we encounter
                                  // at the beginning
                                  findfirst(production[j][2], j, 3);
                         }
                 }
        }
}
void followfirst(char c, int c1, int c2)
{
        int k;
        // The case where we encounter
        // a Terminal
        if(!(isupper(c)))
                 f[m++] = c;
        else
        {
                 int i = 0, j = 1;
                 for(i = 0; i < count; i++)
                 {
                         if(calc_first[i][0] == c)
                                  break;
                 }
                 //Including the First set of the
                 // Non-Terminal in the Follow of
                 // the original query
                 while(calc_first[i][j] != '!')
                 {
```

```
if(calc_first[i][j] != '#')
                         {
                                 f[m++] = calc_first[i][j];
                         }
                         else
                         {
                                 if(production[c1][c2] == '\0')
                                 {
                                         // Case where we reach the
                                         // end of a production
                                         follow(production[c1][0]);
                                 }
                                 else
                                 {
                                         // Recursion to the next symbol
                                         // in case we encounter a "#"
                                         followfirst(production[c1][c2], c1, c2+1);
                                 }
                         }
                         j++;
                }
        }
}
```

```
First(E) = { +, }
First() = { }

Follow(E) = { $, ', |, }

Follow() = { $, ', |, +, }
```

RESULT:

First and follow fwas successfully calculated manually and the program was successfully executed using an Online C++ compiler

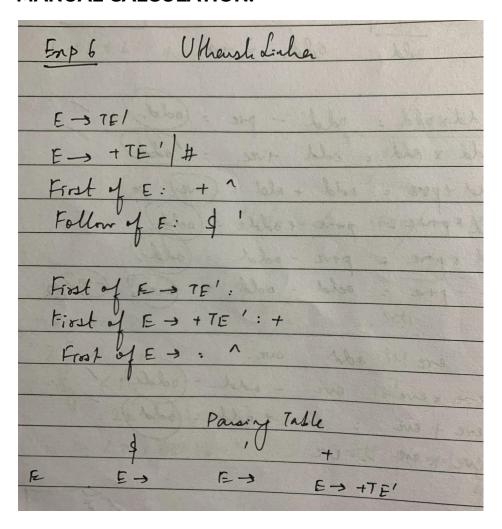
EXPERIMENT NO - 6

Predictive Parsing Table

AIM:

To write a program for making Predictive Parsing Table to the given grammar using manual calculation and using a C++ program.

MANUAL CALCULATION:



CODE:

#include<stdio.h>

#include<string.h>

#define TSIZE 128

int table[100][TSIZE];

char terminal[TSIZE];

```
char nonterminal[26];
struct product {
  char str[100];
  int len;
}pro[20];
int no_pro;
char first[26][TSIZE];
char follow[26][TSIZE];
char first_rhs[100][TSIZE];
int isNT(char c) {
  return c >= 'A' && c <= 'Z';
}
void readFromFile() {
  FILE* fptr;
  fptr = fopen("TEXT.txt", "r");
  char buffer[255];
  int i;
  int j;
  while (fgets(buffer, sizeof(buffer), fptr)) {
    printf("%s", buffer);
    j = 0;
    nonterminal[buffer[0] - 'A'] = 1;
    for (i = 0; i < strlen(buffer) - 1; ++i) {
       if (buffer[i] == '|') {
         ++no_pro;
         pro[no_pro - 1].str[j] = '\0';
         pro[no_pro - 1].len = j;
         pro[no_pro].str[0] = pro[no_pro - 1].str[0];
         pro[no_pro].str[1] = pro[no_pro - 1].str[1];
         pro[no_pro].str[2] = pro[no_pro - 1].str[2];
         j = 3;
```

```
}
       else {
         pro[no_pro].str[j] = buffer[i];
         ++j;
         if (!isNT(buffer[i]) && buffer[i] != '-' && buffer[i] != '>') {
            terminal[buffer[i]] = 1;
         }
       }
    }
    pro[no_pro].len = j;
    ++no_pro;
  }
}
void add_FIRST_A_to_FOLLOW_B(char A, char B) {
  int i;
  for (i = 0; i < TSIZE; ++i) {
    if (i != '^')
       follow[B - 'A'][i] = follow[B - 'A'][i] | | first[A - 'A'][i];
  }
}
void add_FOLLOW_A_to_FOLLOW_B(char A, char B) {
  int i;
  for (i = 0; i < TSIZE; ++i) {
    if (i != '^')
       follow[B - 'A'][i] = follow[B - 'A'][i] |  | follow[A - 'A'][i];
 }
}
void FOLLOW() {
  int t = 0;
  int i, j, k, x;
  while (t++ < no_pro) {
```

```
for (k = 0; k < 26; ++k) {
       if (!nonterminal[k]) continue;
       char nt = k + 'A';
       for (i = 0; i < no_pro; ++i) {
         for (j = 3; j < pro[i].len; ++j) {
            if (nt == pro[i].str[j]) {
              for (x = j + 1; x < pro[i].len; ++x) {
                 char sc = pro[i].str[x];
                 if (isNT(sc)) {
                   add_FIRST_A_to_FOLLOW_B(sc, nt);
                   if (first[sc - 'A']['^'])
                      continue;
                 }
                 else {
                   follow[nt - 'A'][sc] = 1;
                 }
                 break;
              }
              if (x == pro[i].len)
                 add_FOLLOW_A_to_FOLLOW_B(pro[i].str[0], nt);
            }
         }
       }
    }
  }
void add_FIRST_A_to_FIRST_B(char A, char B) {
  int i;
  for (i = 0; i < TSIZE; ++i) {
    if (i != '^') {
       first[B - 'A'][i] = first[A - 'A'][i] || first[B - 'A'][i];
```

}

```
}
  }
}
void FIRST() {
  int i, j;
  int t = 0;
  while (t < no_pro) {
    for (i = 0; i < no_pro; ++i) {
       for (j = 3; j < pro[i].len; ++j) {
          char sc = pro[i].str[j];
          if (isNT(sc)) {
            add_FIRST_A_to_FIRST_B(sc, pro[i].str[0]);
            if (first[sc - 'A']['^'])
               continue;
         }
          else {
            first[pro[i].str[0] - 'A'][sc] = 1;
         }
          break;
       if (j == pro[i].len)
         first[pro[i].str[0] - 'A']['^'] = 1;
    }
    ++t;
  }
}
void add_FIRST_A_to_FIRST_RHS__B(char A, int B) {
  int i;
  for (i = 0; i < TSIZE; ++i) {
    if (i != '^')
       first_rhs[B][i] = first[A - 'A'][i] || first_rhs[B][i];
```

```
}
}
// Calculates FIRST(ß) for each A->ß
void FIRST_RHS() {
  int i, j;
  int t = 0;
  while (t < no_pro) {
    for (i = 0; i < no_pro; ++i) {
       for (j = 3; j < pro[i].len; ++j) {
         char sc = pro[i].str[j];
         if (isNT(sc)) {
            add_FIRST_A_to_FIRST_RHS__B(sc, i);
            if (first[sc - 'A']['^'])
              continue;
         }
         else {
            first_rhs[i][sc] = 1;
         }
         break;
       if (j == pro[i].len)
         first_rhs[i]['^'] = 1;
    }
    ++t;
  }
}
int main() {
  readFromFile();
  follow[pro[0].str[0] - 'A']['$'] = 1;
  FIRST();
  FOLLOW();
```

```
FIRST_RHS();
int i, j, k;
// display first of each variable
printf("\n");
for (i = 0; i < no_pro; ++i) {
  if (i == 0 || (pro[i - 1].str[0] != pro[i].str[0])) {
     char c = pro[i].str[0];
     printf("FIRST OF %c: ", c);
     for (j = 0; j < TSIZE; ++j) {
       if (first[c - 'A'][j]) {
          printf("%c ", j);
       }
     }
     printf("\n");
  }
}
// display follow of each variable
printf("\n");
for (i = 0; i < no_pro; ++i) {
  if (i == 0 | | (pro[i - 1].str[0] != pro[i].str[0])) {
     char c = pro[i].str[0];
     printf("FOLLOW OF %c: ", c);
     for (j = 0; j < TSIZE; ++j) {
       if (follow[c - 'A'][j]) {
          printf("%c ", j);
       }
     }
     printf("\n");
  }
```

```
}
// display first of each variable \mbox{\ensuremath{\ensuremath{\wp}}}
// in form A->ß
printf("\n");
for (i = 0; i < no_pro; ++i) {
  printf("FIRST OF %s: ", pro[i].str);
  for (j = 0; j < TSIZE; ++j) {
    if (first_rhs[i][j]) {
       printf("%c ", j);
    }
  }
  printf("\n");
}
// the parse table contains '$'
// set terminal['$'] = 1
// to include '$' in the parse table
terminal['$'] = 1;
// the parse table do not read '^'
// as input
// so we set terminal['^{\prime}] = 0
// to remove '^' from terminals
terminal['^{'}] = 0;
// printing parse table
printf("\n");
printf("\n\t************** LL(1) PARSING TABLE **************\n");
printf("\t----\n");
printf("%-10s", "");
for (i = 0; i < TSIZE; ++i) {
```

```
if (terminal[i]) printf("%-10c", i);
}
printf("\n");
int p = 0;
for (i = 0; i < no_pro; ++i) {
  if (i != 0 && (pro[i].str[0] != pro[i - 1].str[0]))
     p = p + 1;
  for (j = 0; j < TSIZE; ++j) {
     if (first_rhs[i][j] && j != '^') {
       table[p][j] = i + 1;
     }
     else if (first_rhs[i]['^']) {
       for (k = 0; k < TSIZE; ++k) {
          if (follow[pro[i].str[0] - 'A'][k]) {
            table[p][k] = i + 1;
          }
       }
     }
  }
}
k = 0;
for (i = 0; i < no_pro; ++i) {
  if (i == 0 | | (pro[i - 1].str[0] != pro[i].str[0])) {
     printf("%-10c", pro[i].str[0]);
     for (j = 0; j < TSIZE; ++j) {
       if (table[k][j]) {
          printf("%-10s", pro[table[k][j] - 1].str);
       }
       else if (terminal[j]) {
          printf("%-10s", "");
       }
```

```
}
    ++k;
    printf("\n");
}
}
```

TEXT.txt FILE:

E->TE'

E->+TE'|#

OUTPUT:

RESULT:

Predictive Parsing Table was successfully calculated manually and the program was successfully executed using an Online C++ compiler.

EXPERIMENT NO - 7

Shift Reduce Parsing

AIM:

To write a program for performing Shift Reduce Parsing to the given grammar using manual calculation and using a C++ program.

MANUAL CALCULATION:

	1300	at the bod observed and the
Maruel	Cale.	The set of second beautiful of the set
Enp. 7.		
Λ Λ-	1 - 1 4	
$A \rightarrow aAB$	aBe aAc	306 6 60
$B \rightarrow C$		
A -> a AB B -> c 3 Input It	ing aaccc	
	U	
N,-1	Front	0.7
Stach	Imput	Palion
a \$	aacci \$	Shift
	acus	shift
aaş	ccc \$	shift
nac g	ccs	Reduce B = C
aa B g	CCF	Lhift
aabes	c \$	Reduce 4-raBC
a A &	c \$	chift.
aAc 4 A\$	\$.	reduce A + aAc
4 \$	- 4	Je cup liel
And the second s		

CODE:

#include<iostream>

```
#include<string.h>
using namespace std;
struct prodn
{
       char p1[10];
       char p2[10];
};
int main()
{
       char input[20],stack[50],temp[50],ch[2],*t1,*t2,*t;
       int i,j,s1,s2,s,count=0;
       struct prodn p[10];
       FILE *fp=fopen("input.txt","r");
       stack[0]='\0';
       cout<<"Enter the Input String:\n";</pre>
       cin>>input;
       while(!feof(fp))
        {
               fscanf(fp,"%s\n",temp);
               t1=strtok(temp,"->");
               t2=strtok(NULL,"->");
               strcpy(p[count].p1,t1);
               strcpy(p[count].p2,t2);
               count++;
       }
       i=0;
       while(1)
        {
               if(i<strlen(input))</pre>
                {
```

```
ch[0]=input[i];
       ch[1]='\0';
       i++;
       strcat(stack,ch);
       cout<<"\n"<<stack;
}
for(j=0;j<count;j++)
{
       t=strstr(stack,p[j].p2);
       if(t!=NULL)
        {
               s1=strlen(stack);
               s2=strlen(t);
               s=s1-s2;
               stack[s]='\0';
               strcat(stack,p[j].p1);
               cout<<"\n"<<stack;
               j=-1;
        }
}
if(strcmp(stack,"E")==0&&i==strlen(input))
{
       cout<<"\n\nAccepted";</pre>
       break;
}
if(i==strlen(input))
{
       cout<<"\n\nNot Accepted";</pre>
       break;
}
```

```
return 0;

TEXT FILE:
A=>aAB|aBc|aAc
B->C
```

```
Parking Action
                 Limport Boltzen
                                     SMITT.
                                    Reduce 5-3202
52300
                                    5251.6
                 (25
                                    Reduce 5-x3E3
529.2
                 18
                                    SECTION AND ADDRESS.
                  Ħ
                                    Seduce 5->282
                 B
                                    Accepted
... Program finished with exit code #
Press ENTER to exit console.
```

RESULT:

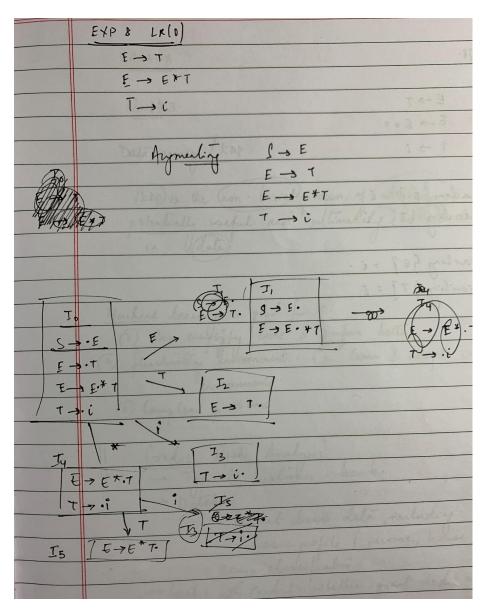
Shift Reduce Parsing was successfully calculated manually and the program was successfully executed using an Online C++ compiler.

Computation of LR(0) items

AIM:

To write a program for finding LR(0) items for the given grammar using manual calculation and using a C++ program.

MANUAL CALCULATION:



CODE:

```
#include<iostream>
#include<conio.h>
#include<string.h>
using namespace std;
char prod[20][20],listofvar[26]="ABCDEFGHIJKLMNOPQR";
int novar=1,i=0,j=0,k=0,n=0,m=0,arr[30];
int noitem=0;
struct Grammar
{
       char lhs;
       char rhs[8];
}g[20],item[20],clos[20][10];
int isvariable(char variable)
{
       for(int i=0;i<novar;i++)</pre>
              if(g[i].lhs==variable)
                      return i+1;
       return 0;
}
void findclosure(int z, char a)
{
       int n=0,i=0,j=0,k=0,l=0;
       for(i=0;i < arr[z];i++)
       {
              for(j=0;j<strlen(clos[z][i].rhs);j++)
```

```
{
                      if(clos[z][i].rhs[j]=='.' && clos[z][i].rhs[j+1]==a)
                       {
                              clos[noitem][n].lhs=clos[z][i].lhs;
                              strcpy(clos[noitem][n].rhs,clos[z][i].rhs);
                              char temp=clos[noitem][n].rhs[j];
                              clos[noitem][n].rhs[j]=clos[noitem][n].rhs[j+1];
                              clos[noitem][n].rhs[j+1]=temp;
                              n=n+1;
                       }
               }
       }
       for(i=0;i< n;i++)
        {
               for(j=0;j<strlen(clos[noitem][i].rhs);j++)
               {
                      if(clos[noitem][i].rhs[j]=='.' && isvariable(clos[noitem][i].rhs[j+1])>0)
                       {
                              for(k=0;k<novar;k++)
                               {
                                      if(clos[noitem][i].rhs[j+1] == clos[0][k].lhs)
                                      {
                                              for(l=0;l< n;l++)
                                                      if(clos[noitem][1].lhs==clos[0][k].lhs
&& strcmp(clos[noitem][l].rhs,clos[0][k].rhs)==0)
                                                             break;
                                              if(l==n)
                                                      clos[noitem][n].lhs=clos[0][k].lhs;
                                              strcpy(clos[noitem][n].rhs,clos[0][k].rhs);
                                                     n=n+1;
```

```
}
                                      }
                              }
                       }
               }
       }
       arr[noitem]=n;
       int flag=0;
       for(i=0;i<noitem;i++)
               if(arr[i]==n)
               {
                       for(j=0;j<\!arr[i];j++)
                              int c=0;
                              for(k=0;k<arr[i];k++)
                                      if(clos[noitem][k].lhs == clos[i][k].lhs \ \&\&
strcmp(clos[noitem][k].rhs,clos[i][k].rhs)==0)
                                              c=c+1;
                              if(c==arr[i])
                               {
                                      flag=1;
                                      goto exit;
                               }
               }
       }
       exit:;
       if(flag==0)
               arr[noitem++]=n;
}
```

```
int main()
{
       cout << "ENTER THE PRODUCTIONS OF THE GRAMMAR(0 TO END) :\n";
       do
              cin>>prod[i++];
       \} while (strcmp (prod[i-1],"0")!=0);\\
       for(n=0;n< i-1;n++)
              m=0;
              j=novar;
              g[novar++].lhs=prod[n][0];
              for(k=3;k<strlen(prod[n]);k++)
               {
                     if(prod[n][k] != '|')
                     g[j].rhs[m++]=prod[n][k];
                     if(prod[n][k]=='|')
                             g[j].rhs[m]='\0';
                             m=0;
                             j=novar;
                             g[novar++].lhs=prod[n][0];
                     }
              }
       for(i=0;i<26;i++)
              if(!isvariable(listofvar[i]))
                     break;
       g[0].lhs=listofvar[i];
```

```
char temp[2]=\{g[1].lhs,'\0'\};
strcat(g[0].rhs,temp);
cout<<"\n\n augumented grammar \n";</pre>
for(i=0;i<novar;i++)
        cout<<endl<<g[i].lhs<<"->"<<g[i].rhs<<" ";
for(i=0;i<novar;i++)
        clos[noitem][i].lhs=g[i].lhs;
        strcpy(clos[noitem][i].rhs,g[i].rhs);
        if(strcmp(clos[noitem][i].rhs,"ε")==0)
               strcpy(clos[noitem][i].rhs,".");
        else
        {
               for(int j=strlen(clos[noitem][i].rhs)+1;j>=0;j--)
                       clos[noitem][i].rhs[j]=clos[noitem][i].rhs[j-1];
               clos[noitem][i].rhs[0]='.';
        }
arr[noitem++]=novar;
for(int z=0;z<noitem;z++)</pre>
{
        char list[10];
       int 1=0;
        for(j=0;j<arr[z];j++)
        {
               for(k=0;k\leq trlen(clos[z][j].rhs)-1;k++)
                {
                       if(clos[z][j].rhs[k]=='.')
                        {
```

```
for(m=0;m< l;m++)
                                             if(list[m]==clos[z][j].rhs[k+1])
                                                     break;
                                     if(m==1)
                                             list[l++]=clos[z][j].rhs[k+1];
                              }
                      }
               }
              for(int x=0;x<1;x++)
                      findclosure(z,list[x]);
       cout << "\n The SET OF ITEMS ARE \n\n";
       for(int z=0; z<noitem; z++)
       {
              cout << "\n I" << z << "\n\n";
              for(j=0;j<arr[z];j++)
                      cout << clos[z][j].lhs << "-> "<< clos[z][j].rhs << "\n";
       }
}
```

```
ENTER THE PRODUCTIONS OF THE GRAMMAR(0 TO END) :
E->T
E->E*T
T->i
 augumented grammar
A->E
E->T
E->E*T
T->i
THE SET OF ITEMS ARE
 I0
A->.E
E->.T
E->.E*T
T->.i
 11
A->E.
E->E.*T
 12
E->T.
13
T->i.
 14
E->E*.T
T->.i
 15
E->E*T.
```

RESULT:

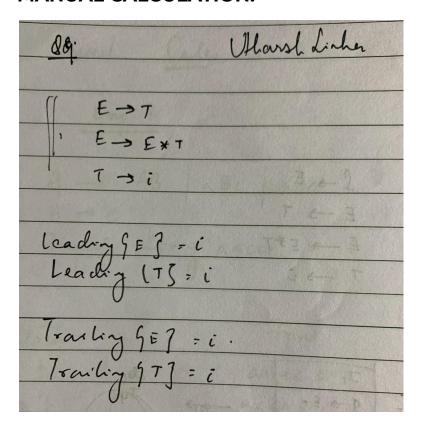
LR(0) items was successfully calculated manually and the program was successfully executed using an Online C++ compiler.

Computation of LEADING AND TRAILING

AIM:

To write a program for finding Leading and Trailing to the given grammar using manual calculation and using a C++ program.

MANUAL CALCULATION:



CODE:

#include<iostream>

#include<conio.h>

#include<stdio.h>

#include<string.h>

#include<stdlib.h>

using namespace std;

int vars,terms,i,j,k,m,rep,count,temp=-1;

```
char var[10],term[10],lead[10][10],trail[10][10];
struct grammar
{
       int prodno;
       char lhs,rhs[20][20];
}gram[50];
void get()
{
       cout<<"\nLEADING AND TRAILING\n";</pre>
       cout<<"\nEnter the no. of variables : ";</pre>
       cin>>vars;
       cout<<"\nEnter the variables : \n";</pre>
       for(i=0;i < vars;i++)
        {
               cin>>gram[i].lhs;
               var[i]=gram[i].lhs;
       cout<<"\nEnter the no. of terminals : ";</pre>
       cin>>terms;
       cout<<"\nEnter the terminals : ";</pre>
       for(j=0;j < terms;j++)
               cin>>term[j];
       cout<<"\nPRODUCTION DETAILS\n";</pre>
       for(i=0;i\leq vars;i++)
        {
               cout << "\nEnter the no. of production of "<< gram[i].lhs << ":";
               cin>>gram[i].prodno;
               for(j=0;j<gram[i].prodno;j++)
                {
                       cout<<gram[i].lhs<<"->";
```

```
cin>>gram[i].rhs[j];
               }
       }
}
void leading()
{
       for(i=0;i<vars;i++)
               for(j=0;j<gram[i].prodno;j++)</pre>
               {
                      for(k=0;k<terms;k++)
                              if(gram[i].rhs[j][0]==term[k])
                                     lead[i][k]=1;
                              else
                              {
                                     if(gram[i].rhs[j][1] == term[k])
                                             lead[i][k]=1;
                              }
                      }
               }
       for(rep=0;rep<vars;rep++)
       {
              for(i=0;i<vars;i++)
               {
                      for(j=0;j<gram[i].prodno;j++)
                      {
                              for(m=1;m<vars;m++)
                              {
```

```
if(gram[i].rhs[j][0]==var[m])
                                             temp=m;
                                             goto out;
                                      }
                              }
                              out:
                              for(k=0;k<terms;k++)
                              {
                                      if(lead[temp][k]==1)
                                             lead[i][k]=1;
                              }
                      }
               }
       }
}
void trailing()
{
       for(i=0;i<vars;i++)
       {
               for(j=0;j \leq gram[i].prodno;j++)
               {
                      count=0;
                      while(gram[i].rhs[j][count]!='\x0')
                              count++;
                      for(k=0;k<terms;k++)
                       {
                              if(gram[i].rhs[j][count-1] == term[k]) \\
                                      trail[i][k]=1;
                              else
```

```
{
                             if(gram[i].rhs[j][count-2]==term[k])
                                    trail[i][k]=1;
                      }
              }
       }
}
for(rep=0;rep<vars;rep++)
       for(i=0;i<vars;i++)
       {
              for(j=0;j \le gram[i].prodno;j++)
                     count=0;
                     while(gram[i].rhs[j][count]!='\x0')
                             count++;
                     for(m=1;m<vars;m++)
                      {
                             if(gram[i].rhs[j][count-1]==var[m])
                                    temp=m;
                     for(k=0;k<terms;k++)
                      {
                             if(trail[temp][k]==1)
                                    trail[i][k]=1;
                      }
              }
}
```

}

```
void display()
{
        for(i=0;i<vars;i++)
                cout<<"\nLEADING("<<gram[i].lhs<<") = ";
                for(j=0;j < terms;j++)
                 {
                         if(lead[i][j]==1)
                                 cout \!\!<\!\! term[j] \!\!<\!\! ",\!";
                }
        }
        cout << endl;
        for(i=0;i<vars;i++)
        {
                cout<<"\nTRAILING("<<gram[i].lhs<<") = ";
                for(j=0;j < terms;j++)
                 {
                         if(trail[i][j]==1)
                                 cout \!\!<\!\! term[j] \!\!<\!\! ",\!";
                }
        }
int main()
{
        get();
        leading();
        trailing();
        display();
```

```
LEADING AND TRAILING

Enter the no. of variables : 2

Enter the variables : E

T

Enter the no. of terminals : 2

Enter the terminals : *

i

PRODUCTION DETAILS

Enter the no. of production of E:2
E->T
E->E+T

Enter the no. of production of T:1
T->i

LEADING(E) = *,i,
LEADING(T) = i,

TRAILING(E) = *,i,
TRAILING(T) = i,
```

RESULT:

Leading and Trailing was successfully calculated manually and the program was successfully executed using an Online C++ compiler.

Intermediate code generation – Postfix, Prefix

AIM:

To generate the prefix and postfix for the given expression.

MANUAL CALCULATION:

Eng 10.	4 4/314	3
37 1000	Les 4 Les / Grand	3
Infin to just	fin	0 T
0 /	1	7
	30/(3)	4
Input Uning	Oulgut Stack	Operator Stack
((A+B) +(C/b))		(
(A+B) + (C/A))		(
A*R) + (LID)		U
*B) + (c/b))	A	(1
*B) + (c/o))	A	(1
B) + (UD))	A	((+
B) + (c/b))	A	((*
) + (c/b))	AB	((+
t(c/o))	AB*	
+ (c/o))	AB +	
(C/D))	AB +	1+
(4/1))	AB *	(+
C/D))	AB+	(+(
(75g)/b))	AB*C	41
(A))	AB + C	1+0
<u>'b))</u>	4B + C	(+(/

₽ ▷))	AB+C	The to profes	(+(/
))	AB*CD	9-1-0	(+(/
)	AB+CD/		(+
)	AB + Cb/	Ha:	(+
	AB + CD/+		
((AY8)+		115年 1	
120 11.1 A	m. 43 * CD/+	12 2 · Re	

CODE:

OPERATORS = set(['+', '-', '*', '/', '(', ')'])

PRI = {'+': 1, '-': 1, '*': 2, '/': 2}

INFIX ===> POSTFIX

def infix_to_postfix(formula):

stack = [] # only pop when the coming op has priority

```
output = "
for ch in formula:
  if ch not in OPERATORS:
     output += ch
  elif ch == '(':
     stack.append('(')
  elif ch == ')':
     while stack and stack[-1] != '(':
       output += stack.pop()
     stack.pop() # pop '('
  else:
     while stack and stack[-1] != '(' and PRI[ch] <= PRI[stack[-1]]:
       output += stack.pop()
     stack.append(ch)
     # leftover
while stack:
```

```
output += stack.pop()
  print(f'POSTFIX: {output}')
  return output
### INFIX ===> PREFIX ###
def infix_to_prefix(formula):
  op_stack = []
  exp_stack = []
  for ch in formula:
    if not ch in OPERATORS:
       exp_stack.append(ch)
    elif ch == '(':
       op_stack.append(ch)
    elif ch == ')':
       while op_stack[-1] != '(':
         op = op_stack.pop()
         a = exp_stack.pop()
```

```
b = \exp_{\text{stack.pop}}()
        exp_stack.append(op + b + a)
     op_stack.pop() # pop '('
  else:
     while op_stack and op_stack[-1] != '(' and PRI[ch] <= PRI[op_stack[-1]]:
        op = op_stack.pop()
        a = \exp_{\text{stack.pop}}()
        b = \exp_{\text{stack.pop}}()
        exp_stack.append(op + b + a)
     op_stack.append(ch)
     # leftover
while op_stack:
  op = op\_stack.pop()
  a = exp_stack.pop()
  b = \exp_{\text{stack.pop}}()
  exp_stack.append(op + b + a)
```

```
print(fPREFIX: {exp_stack[-1]}')

return exp_stack[-1]

expres = input("INPUT THE EXPRESSION: ")

pre = infix_to_prefix(expres)

pos = infix_to_postfix(expres)
```

```
INPUT THE EXPRESSION: ((A*B)+(C/D))
PREFIX: +*AB/CD
POSTFIX: AB*CD/+
```

RESULT:

The program was successfully executed and output was verified.

Intermediate code generation – Quadruple, Triple, Indirect triple

AIM:

To find the intermediate code generation - Quadruple, triple, indirect.

MANUAL CALCULATION:

Para um 2010521				
Enp 11 RA19/100301052/				
Fort Engrusson: ((A+B)+(C/D))				
Prelia : +*AB/CD				
Prefix: +*AB/CD Abelfix: AB *CD/+				
ti				
three address code generalion:				
Postfin: AB*CD/+ three address code generation: $ \begin{array}{cccccccccccccccccccccccccccccccccc$				
t2:= 0/D				
$t_2 := C/D$ $t_3 : t_1 + t_2$ $t_3 : t_1 + t_2$				
Quadrule:				
Guaduple:				
OPEration ARGument/ Argument 2. Result				
A 8 (t(1)				
/ C D t(2) + t(1) t(2) t(3)				
Triple				
Bart Bart Age				
Operation fogument/ Argument 2 pessent				
Operation fogument/ Argument 2 Pessent				
3				
+ (o) (1)				

CODE:

```
OPERATORS = set(['+', '-', '*', '/', '(', ')'])
PRI = {'+':1, '-':1, '*':2, '/':2}
### INFIX ===> POSTFIX ###
def infix_to_postfix(formula):
  stack = [] # only pop when the coming op has priority
  output = "
  for ch in formula:
    if ch not in OPERATORS:
       output += ch
    elif ch == '(':
       stack.append('(')
    elif ch == ')':
       while stack and stack[-1]!='(':
         output += stack.pop()
       stack.pop() # pop '('
    else:
       while stack and stack[-1] != '(' and PRI[ch] <= PRI[stack[-1]]:
          output += stack.pop()
       stack.append(ch)
  # leftover
  while stack:
       output += stack.pop()
  print(f'POSTFIX: {output}')
  return output
### INFIX ===> PREFIX ###
def infix_to_prefix(formula):
  op_stack = []
  exp_stack = []
```

```
for ch in formula:
     if not ch in OPERATORS:
       exp_stack.append(ch)
     elif ch == '(':
       op stack.append(ch)
     elif ch == ')':
       while op stack[-1] != '(':
          op = op stack.pop()
          a = exp_stack.pop()
          b = \exp \operatorname{stack.pop}()
          exp stack.append(op+b+a)
       op_stack.pop() # pop '('
     else:
       while op stack and op stack[-1] != '(' and PRI[ch] <= PRI[op stack[-1]]:
          op = op stack.pop()
          a = \exp \operatorname{stack.pop}()
          b = \exp \operatorname{stack.pop}()
          exp_stack.append( op+b+a )
       op stack.append(ch)
  # leftover
  while op stack:
     op = op_stack.pop()
     a = \exp \operatorname{stack.pop}()
     b = \exp_{\text{stack.pop}}()
     exp stack.append(op+b+a)
  print(f'PREFIX: {exp stack[-1]}')
  return exp stack[-1]
### THREE ADDRESS CODE GENERATION ###
def generate3AC(pos):
       print("### THREE ADDRESS CODE GENERATION ###")
```

```
\exp_{\text{stack}} = []
       t = 1
       for i in pos:
               if i not in OPERATORS:
                       exp_stack.append(i)
               else:
                      print(f't\{t\} := \{exp\_stack[-2]\} \{i\} \{exp\_stack[-1]\}')
                       exp_stack=exp_stack[:-2]
                      exp\_stack.append(f't\{t\}')
                      t+=1
expres = input("INPUT THE EXPRESSION: ")
pre = infix_to_prefix(expres)
pos = infix_to_postfix(expres)
generate3AC(pos)
def Quadruple(pos):
 stack = []
 op = []
 x = 1
 for i in pos:
  if i not in OPERATORS:
    stack.append(i)
  elif i == '-':
    op1 = stack.pop()
    stack.append("t(%s)" %x)
    print("{0:^4s} | {1:^4s} | {2:^4s}|{3:4s}".format(i,op1,"(-)"," t(%s)" %x))
    x = x+1
    if stack != []:
      op2 = stack.pop()
      op1 = stack.pop()
      print("{0:^4s} | {1:^4s} | {2:^4s}|{3:4s}".format("+",op1,op2," t(%s)" %x))
```

```
stack.append("t(%s)" %x)
                                 x = x+1
             elif i == '=':
                   op2 = stack.pop()
                   op1 = stack.pop()
                   print("\{0:^4s\} \mid \{1:^4s\} \mid \{2:^4s\} \mid \{3:4s\} \mid \{3:4s\} \mid \{0:^4s\} \mid
             else:
                   op1 = stack.pop()
                   op2 = stack.pop()
                   print("{0:^4s} | {1:^4s} | {2:^4s}||{3:4s}".format(i,op2,op1," t(%s)" %x))
                   stack.append("t(%s)" %x)
                   x = x+1
print("The quadruple for the expression ")
print(" OP | ARG 1 | ARG 2 | RESULT ")
Quadruple(pos)
def Triple(pos):
                          stack = []
                          op = []
                          x = 0
                          for i in pos:
                                 if i not in OPERATORS:
                                        stack.append(i)
                                  elif i == '-':
                                        op1 = stack.pop()
                                        stack.append("(%s)" %x)
                                        print("{0:^4s} | {1:^4s} | {2:^4s}".format(i,op1,"(-)"))
                                        x = x+1
                                         if stack != []:
                                               op2 = stack.pop()
                                               op1 = stack.pop()
```

```
print("{0:^4s} | {1:^4s} | {2:^4s}".format("+",op1,op2))
         stack.append("(%s)" %x)
        x = x+1
      elif i == '=':
       op2 = stack.pop()
       op1 = stack.pop()
       print("\{0:^4s\} \mid \{1:^4s\} \mid \{2:^4s\}".format(i,op1,op2))
      else:
       op1 = stack.pop()
       if stack != []:
        op2 = stack.pop()
         print("\{0:^4s\} | \{1:^4s\} | \{2:^4s\}".format(i,op2,op1))
         stack.append("(%s)" %x)
        x = x+1
print("The triple for given expression")
print(" OP | ARG 1 | ARG 2 ")
Triple(pos)
```

```
INPUT THE EXPRESSION: ((A*B)+(C/D))
PREFIX: +*AB/CD
POSTFIX: AB*CD/+
   THREE ADDRESS CODE GENERATION ###
The quadruple for the expression
            ARG 2
                    RESULT
      ARG 1
               D
        С
                    t(2)
              t(2) t(3)
       t(1)
The triple for given expression
       ARG 1 ARG 2
               В
        Α
        C
               D
       (0)
              (1)
```

RESULT: The program was successfully executed and output was verified.

