

PIB File Specification PIB File Specification

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1. Overview **1. Overview**

The Platform-Independent Binary or PIB file format was developed to provide a lightweight, machine-
independent format for the plotting and analysis of data from the NRC Reactor Safety Data Dank. The NRC independent format for the plotting and analysis of data from the NRC Reactor Safety Data Dank. The NRC Data Bank was originally developed at the Idaho National Engineering Laboratory, to collect, store, and make available, data from many domestic and foreign light water reactor safety research programs. The NRC Data available, data from many domestic and foreign light water reactor safety research programs. The NRC Data Bank is currently being modified and relocated to the NRC offices in Rockville, MD. The XMGR5 plotting and Bank is currently being modified and relocated to the NRC offices in Rockville, XMGRS plotting and data analysis software is used to read this experimental data as well as calculated data generated from several data analysis software is used to read this experimental data as well as calculated data generated from several reactor safety analysis programs including RELAP, MELCOR and FRAPCON. reactor safety analysis programs including RELAP, MELCOR and FRAPCON.

The data files maintained by the NRC Data Bank consist of large time-dependent data files containing the results of experimental calculations. A single experiment can generate several hundred megabytes of data. ASCII files of experimental calculations. A single experiment can generate several hundred megabytes of data. files provide a portable data format for exchange of data between different machines, however, it is generally very inefficient and time consuming to store and extract data from large ASCII files. Native binary files can be used to store and retrieve data in a very efficient manner using direct access pointers, however, they cannot be read store and retrieve data in a very efficient manner using direct access pointers, however, they cannot be read directly by other machine types. The PIB format uses the External Data Representation (XDR) protocol to provide a portable, machine-independent format that can be accessed with an efficiency approaching that of the native binary format. binary format.

Two utility programs, Convert and Merge were developed to provide a means of generating data files from the Two utility programs, Convert and Merge were developed to provide a means of generating data files the NRC Data Bank for XMGR5. Figure 1 illustrates the usage these programs as well as the role of the NRC Data NRC Data Bank for XMGRS. Figure 1 illustrates the usage these programs as well the role of the NRC Data Bank in the creation of PIB data files. Data is obtained from experimental facilities on magnetic tape. Each facility sends data in its own format, and a single experiment may encompass several tapes. These tapes are facility sends data in its own format, and a. single experiment may encompass several tapes. These tapes are processed by the NRC Data bank and stored in a common binary format. Data stored in the NRC Data Bank processed by the NRC Databank and stored in a common binary format. Data stored in the NRC Data Bank can be dumped into an ASCII file using a standard format known as TWX. The Convert utility reads a single can be dumped into an ASCII file using a standard format known as TWX .. The Convert utility reads a single TWX file and converts it to a binary file that can be read by XMGR5. This BIN format was initially used to TWX file and converts it to a binary fIle that can be read by XMGRS. This BIN format was initially used to provide experimental data to XMGR5. Although it is possible for XMGR5 to read from several files provide experimental data to XMGRS. Although it is possible for XMGRS to read from several files simultaneously, it became clear that it would be desirable to merge all the data files from a single experiment into a single file that could then be read by XMGR5. Due to the wide range of Unix platforms used to analyze the a single file that could then be read by XMGRS. Due to the wide range of Unix platforms used to analyze the data, there is also a benefit to switching to an efficient platform independent format. The Merge program can data, there is also a benefIt to switching to an efficient platform independent format. The Merge program can

read multiple BIN or PIB files and merge them into a single PIB file. In addition, the Merge program is used to verify data integrity, perform statistical analyses, apply digital filters, compress the channel data, and perform basic data management functions. Much of the example code shown in this report was derived from the Merge data management functions. Much of the example code shown in this report derived from the Merge program. program.

This document is meant to serve as a guide for creating PIB files. Several C-language code segments are provided to illustrate programming approaches. As such, a general knowledge of C-language programming is assumed. to illustrate programming approaches. As such, a general knowledge of Clanguage programming is assumed.

1. File Structure **1. File Structure**

The PIB file structure is composed of three distinct sections, a file header block, a channel header block, and the channel data block. The file header block is used to identify the file type and it contains the information channel data block. The file header block is used to identify the file type and it contains the information necessary to read the channel header block. The channel header block contains one channel header record for necessary to read the channel header block. The channel header block contains one channel header record for each channel in the file. The channel header records contain the information required to read the channel data each channel in the file. The channel header records contain the information required to read the channel data block. A data channel can represent either a set of time values or a set of dependent data values. Each dependent block. A data channel can represent either a set of time values or a dependent data values. Each dependent data channel includes a reference to a time data channel that provides the independent data values. Many data channel includes a reference to a time data channel that provides the independent data values. Many dependent data channels generally reference the same time channel. dependent data channels generally reference the same time channel.

In order to achieve platform independence, all data written to or retrieved from the PIB file is accessed by means of the External Data Representation (XDR) routines. In the UNIX operating system, the interface for the XDR of the External Data Representation (XDR) routines. In the UNIX operating system, the interface for the XDR primitives are defined in the include file $<$ rpc/xdr.h $>$, which is automatically included by $<$ rpc/rpc.h $>$. Refer to a programmers guide to remote procedure calls (RPC) for a description of the XDR protocol. One such reference is Power *Programming with RPC* by John Bloomer, O'Reilly & Associates, 1992. reference is *Power ProgrammingwithRPC* by John Bloomer, O'Reilly & Associates, 1992.

The process of creating a PIB file is illustrated in Figure 2. As will be discussed further in subsequent sections, the PIB file is written in two passes. During the first pass, the file header block, the channel header block, and the channel data block are written to the file. However, the file pointer information required for the channel header

block is not determined until the channel data block is written to file. For this reason, a second pass is required to update the channel header block. update the channel header block. . .

The first two process steps involve creating the file header block structure and writing it to an XDR stream. The first two process steps involve creating the file header structure and writing to an XDR stream. These steps are discussed in Section 2.A, File Header Block. Steps 3, 4, 9 and 10, which involve writing the channel header block, are covered in Section 2.B. Steps 5, 6 and 7 involve writing the channel data block and channel header block, are covered in Section 2.B. Steps 5, 6 and 7 involve writing the channel data block and updating the channel header records. These steps are discussed in Section 2.C. updating the channel header records. These steps are discussed in Section 2.C.

A. File Header Block **A. File Header Block**

The PIB file begins with a header block that is used to identify the file type and to provide the information necessary to read the subsequent channel header records. It may also include an optional list of filenames that necessary to read the subsequent channel header records. **It** may also include an optional list of that were used to create the PIB file along with an integer value indicating their file type. Currently, integer values of 1000 and 2000 are reserved to identify binary and PIB file types respectively. This file list can prove useful in tracking down data sources, especially if several files are combined using the merge utility. Table 1 illustrates the layout of the file header block. layout of the file header block.

Table 1. File Header Block Layout **Table 1.** File Header Block Layout

The following xdr_fileHead routine can be used to read or write the file header block to an existing XDR stream. The header data is retrieved from or placed into a pibHeader structure as defined below: The header data is retrieved from or placed into a pibHeader structure as defined below:

Source Listing 1. File Header XDR Access Routine **Source Listing 1.** File Header XDR Access Routine

```
struct pibHeader {
struct pibHeader { 
  char fileTyp[80];
  int size;
  int numOfChnls;
   int numOfFiles;
   char fromfile[80] [256]
a list of up to 80 filenames */
char fromfile[80] [256]i/* a list of up to 80 filenames */ 
  int fromfiletype[80];
   char tofile[256];
   1*
string indicating file type */
char fileTyp[80]i /* string indicating file type */ 
   int size; \frac{1}{2} \frac{1}{2} \frac{1}{2} \frac{1}{2} and used enter.0 */
   int size;<br>int numOfChnls; \frac{1}{2} /* the number of data channels contained in this file */
   1*
the number of data files in the file list */
int numOfFilesi /* the number of data files in the file list */ 
   chaft fromfiletype[80]; /* a fist of up to so filehames<br>int fromfiletype[80]; /* corresponding file types */
   1*
the name of the file being created */
char tofile[256]i /* the name of the file being created */ 
} i
```
 $\overline{}$

```
bool_t xdr_fileHead(XDR *xdrs, struct pibHeader *hdr)<br>{
     int i;
int i; 
     bool_t rc;
     u int slen;
int slen; 
     char *cptr;
char *cptr; 
     cptr = hdr->fileType;slen = 80;if(! (rc = xdr string(xdrs, &cptr, slen))) return rc;
if(! (rc xdr string(xdrs, &cptr, slen») return rc; 
     if(!(IC = xdr_string(xdrs, &cptr, sien))) return<br>if(!(rc = xdr_int(xdrs, &hdr->size))) return rc;
     if(! (rc = xdr_int(xdrs, &hdr->size))) return rc;<br>if(! (rc = xdr_int(xdrs, &hdr->size))) return rc;<br>if(! (rc = xdr_int(xdrs, &hdr->numOfChnls))) return rc;
     if(! (rc = xdr_{int}(xdr_s, \theta dr->numOfChnls))) return rc;
     if(! (rc = xdr int(xdrs, &hdr->numOfFiles))) return rc;
if(! (rc = xdr-int(xdrs, &hdr->numOfFiles») return rc; 
     for(i=0; i<hatr->numOfFiles; i++) {
          cptr = &hdr->fromfile(i] [0];
cptr &hdr->fromfile[i] [0]; 
          slen = 256;
          sien - 250,<br>if(!(rc = xdr_string(xdrs, &cptr, slen))) return rc;
     } 
     for(i=0; i<hdr->numOfFiles; i++)
for(i=O; i<hdr->numOfFiles; i++) { 
          if(!(rc = xdr int(xdrs, &hdr->fromfiletype[i]))) return rc;
if(! (rc = xdr_int(xdrs, &hdr->fromfiletype[i]») return rc; 
     } 
     .<br>cptr = &hdr->tofile[0];
     slen = 256;sien = 250,<br>if(!(rc = xdr_string(xdrs, &cptr, slen))) return rc;
     return rc;
return rc; 
\mathbf{)}
```
The following code segment demonstrates how to create the XDR stream and write the header using the xdr_fileHead routine. It assumes that the filename and numChan variables, corresponding to the name of the file to be created, and the number of data channels, respectively, have been previously defined. to be created, and the number of data channels, respectively, have been previously defined.

Source Listing 2. Example of xdr fileHead Usage Source Listing 2. Example of xdr _fileHead Usage

```
/* declarations */
/* declarations */ 
XDR xdrs;
xdrs; 
FILE fileHndl;
FILE fileHndl; 
struct pibHeader hdr;
struct pibHeader hdr; 
int rc;
int rc; 
/* copy data to the header structure */
/* copy data to the header structure */ 
strcpy(hdr->filetype, "NRCDB V2.0, K. R. Jones");
strcpy(hdr->filetype, "NRCDB V2.0, K. R. Jones"); 
bdr->size = 0;hdr->numOfChnls = numChan;
hdr->numOfChnls = numChan; 
hot->numOfFiles = 0;
strcpy(hdr->toFile, filename);
strcpy(hdr->toFile, filename); 
/* open the output file */
/* open the output file */ 
fileHndl=fopen(filename,"w");
fileHndl=fopen(filename,"w"); 
if (fileHndl == NULL) I
if (fileHndl == NULL) { 
      fprintf(stderr, "Unable to open file %s", filename);
fprintf(stderr, "Unable to open file %s", filename); 
      return;
return; 
/* create the XDR object */
/* create the XDR object */ 
xdrstdio_create(&xdrs, fileHndl, XDR_ENCODE);
/* write the header structure to the XDR stream */
/* write the header structure to the XDR stream */ 
rc = xdr fileHead(&xdrs, &hdr);<br>if(!rc) \overline{t}if(!rc) \overline{t}:LC/ (<br>fprintf(stderr,"Error writting PIB file...");
      return;
return; 
} ;
```
B. Channel Header Block **B. Channel Header Block**

The channel header block immediately follows the file header block. The channel header block contains one The channel header block immediately the file header block. The channel header contains one channel header record for each data channel. The layout of the channel header record is illustrated in Table 2 channel header record for each data channel. The layout of the channel header record is illustrated in Table 2 along with a description of each field. The channel name is stored in ASCII characters and must be padded to with NULL characters to a length of 24 bytes. This ensures a fixed size for the channel header record, allowing the channel header block to be read very efficiently. Each channel is assigned a unique zero based index to identify that is used to identify the channel.

Table 2. Channel Header Record Layout **Table** 2. Channel Header Record Layout

The most difficult fields to populate are the pointers to the data and time data, ptrToData and ptrToTime, The most difficult fields to populate the pointers to the data and time data, ptrToData and ptrToTime, respectively. These pointers are most easily determined using the xdr_getpos function prior to writing the channel data. Unfortunately, these values are not known at the time the channel header block is written. For this reason, the channel header block is written in two passes. During the first pass, write the file header block, the channel header block and then the data block. As the data block is written, use the xdr_getpos function to determine the file pointer prior to writing each data channel. After the data block has been written, rewind the file to the start of the channel header block and rewrite the channel header block. file to the start of the channel header block and rewrite the channel header block.

The following xdr_chanHead routine can be used to read or write a single channel header record to an existing XDR stream. The channel data is retrieved from or placed into a pibChnlRec structure as defined below: XDR stream. The channel data is retrieved from or placed into a pibChnlRec structure as defined below:


```
Source Listing 3. Channel Header XDR Access Routine
Source Listing 3. Channel Header XDR Access Routine
```

```
rc = (xdr bytes(xdrs, &cptr, &slen, slen)
        Auf_bytes(Aufs, &Cprf, &Sien,<br>&& xdr_int(xdrs, &Chan->Index)
        %& xdr_int(xdrs, &Chan->size)<br>&& xdr_int(xdrs, &Chan->totalSize)
        aa xdr_int(xdrs, achan->tocarsize)<br>&& xdr_int(xdrs, &Chan->timeIndex)
        %& xdr_int(xdrs, &Chan->ptrToData)
        %& xdr_int(xdrs, &Chan->ptrrobata)
        %& xdr_int(xdrs, &Chan->ptriori<br>&& xdr_int(xdrs, &Chan->eucode)
        ad Adr_int(Adrs, achan > cacode
        ad nar inc (nars) domain recens)
        && xdr int(xdrs, &Chan->orgFile)
        ad Adr_int(Adrs, comm Porgrine
        && xdr int(xdrs, &Chan->cmpMode)
&& xdr-int(xdrs, &Chan->cmpMode) 
        && xdr int(xdrs, &Chan->cmpSize)
&& xdr-int(xdrs, &Chan->cmpSize) 
        %& xdr_int(xdrs, &Chan->sparel)
        && xdr int(xdrs, &Chan->spare2)
&& xdr-int(xdrs, &Chan->sparel) 
        ad Adr_int(Adrs, achan->sparez)<br>&& xdr_int(xdrs, &Chan->spare3));
return rc;
return rc; 
cptr = Chan->name;<br>rc = (xdr_bytes(xdrs, &cptr, &slen, slen)
        && xdr int(xdrs, &Chan->Index) 
                                                  Size)<br>ndex)
        && xdr-int(xdrs, &Chan->timeIndex) 
                                                  Data)
                                                  Time)<br>e)
        && xdr-int(xdrs, &Chan->eucode) 
        && xdr-int(xdrs, &Chan->recNo) 
        && xdr-int(xdrs, &Chan->orgIndex) 
        && xdr-int(xdrs, &Chan->orgFile) 
        && xdr-int(xdrs, &Chan->status) 
        && xdr-int(xdrs, &Chan->spare2)
```
 $\,$ }

Typically, an array of pibChnlRec structures will be allocated and filled with the channel information, then the Typically, an array of pibChnlRec structures allocated and filled with·the channel infonnation, then the xdr chanHead routine will be called once for each element of the array.

C. Channel Data Block **C. Channel Data Block**

The channel data block follows immediately after the channel header block. It consists entirely of arrays of double precision floating point data output with the xdr routines. The channel data is written using the xdr_array primitive with a double precision element type. The following code segment illustrates writing the channel data block to an open XDR stream: block to an open stream:

Source Listing 4. Channel Data Block Write **Source Listing 4. Channel Data Block Write**

```
/* declarations */
/* declarations */ 
u int filePtr;
u int filePtri 
bool_t rc;
double *dptr;<br>u_int csize;
u int csize;
int i;
int ii 
int totalNumChan = [number of channels];
int totalNumChan = [number of channels]i 
for(i=0; i<totalNumChan; i++) {
      /* get the current file position
/* get the current file position 
             note: save this to the channel header record array so you can write the
note: save this to the channel header record array so you can write the 
      file pointers in the second pass */
file pointers in the second pass 
     */ 
      ,<br>filePtr = xdr_getpos(&xdrs);
      /* set up dptr and csize variables */
/* set up dptr and csize variables */ 
     dptr = [pointer to data];
dptr = [pointer to data]; 
     csize = [length of data array];
csize = [length of data array]; 
      /* compress the data */
/* compress the data */ 
     cmpMode =cmpres(dptr, &csize);
crnpMode =cmpres(dptr, &csize); 
      /* output channel data */
/* output channel data */ 
     rc = xdr-array(&xdrs, /* XDR stream */
      rc = xdr_array(&xdrs, /* XDR stream */<br>(char **)&dptr, /* pointer to the data to be written */
                                 (u int *)&csize,/* number of elements */
(u int *)&csize,/* number of elements */ 
                                csize, /* maximum number of elements */
csIze, /* maximum number of elements */ 
                               ,sizeof(double), /* element size */
.sizeof(double) , /* element size */ 
                                xdrdouble); /* primative used to encode element */
xdr_double); /* primative used to encode element */ 
     if(!rc) I
if (! rc) ( 
         (:rc) {<br>fprintf(stderr,"Error writting output file. exiting");
         exit(-1);
exit (-1) ; 
     \mathbf{I}/* now save the compression mode and data pointer to the
/* now save the compression mode and data pointer to the 
           channel header record array */
channel header record array */ 
     chanHeader[i].ptrToData = filePtr;
     chanHeader[i].cmpMode = cmpMode;
chanHeader[ij .cmpMode = cmpMode; 
\lambda
```
After the data block has been written, and all of the channel data pointers determined, the pointers to the time data can be determined. This is accomplished by setting the time data pointer value, ptrTotime equal to the data pointer value, ptrToData for the corresponding time channel. For time channels, ptrToTime should be set equal to its own data pointer. This process is illustrated in the following code segment: **to its own data pointer. This process is illustrated in the following code segment:**

a 8

Source Listing 5. Setting Time Pointers Source Listing 5. Setting Time Pointers

```
/* set pointers to time values */
/* set pointers to time values */ 
for(i-0; i<totalNumChan; i++)
for(i~O; i<totalNumChan; i++) { 
    chanHeader[i].ptrToTime = -1;
    if(chanHeader[i].timeIndex == 0) 1
if (chanHeader[il . time Index == 0) { 
        chanHeader[i].ptrToTime = chanHeader[i].ptrToData;
chanHeader[il.ptrToTime chanHeader[il.ptrToData; 
    lelsef
}else{ 
        chanHeader[i].ptrToTime = chanHeader(chanHeader[i].timeIndex].ptrToData;
chanHeader[il.ptrToTime = chanHeader[chanHeader[il.timeIndexl .ptrToData; 
    } 
    if(chanHeader[i].ptrToTime < 0)
if(chanHeader[il.ptrToTime < 0) 
        fprintf(stderr, "No time data found for channel:%s\n",
fprintf(stderr, "No time found for channel:%s\n", 
                        chanHeader[i].name);
chanHeader[il.name); 
        exit (-1);
   \mathbf{I}\overline{\phantom{a}}
```
After the data block has been written, and the channel header record structures have been updated to include the compression and file pointer information, the file is rewound to the start of the channel header block, and the compression and file pointer information, the file is rewound to the start of the channel. header block, and the channel header block is overwritten. channel header block is overwritten.

1. Engineering Unit Codes **1. Engineering Unit Codes**

Each data channel must have an engineering unit code (eucode) defined in its channel header record Each data channel must have an engineering unit code (eucode) defined in its channel header record corresponding to the physical units in which the data channel was calibrated. Table 3 contains a list of all available engineering unit codes. The description and units columns shown in Table 3 are used to provide axis available engineering unit codes. The description and units columns shown in Table 3 are used to provide axis labels in the plotting software. These codes were originally developed for the NRC Data Bank to provide a labels in the plotting software. These codes were originally developed for the NRC Data Bank to provide a common set of units to handle data from a wide range of domestic and foreign research programs. common set of units to handle data from a wide range of domestic and foreign research programs.

In the event that none of the available unit codes are suitable for a given channel and an engineering unit code In the event that none of the available unit codes are suitable for a given channel and an engineering code must be added, contact the NRC Project Manager for the NRC Data Bank to reserve a unit code. This will must be added, contact the NRC Project Manager for the NRC Data Bank to reserve a unit code. This will prevent conflicting use of a code value and ensure that the appropriate modifications are made to the plotting and analysis software. analysis software.

Eng. Unit	${\tt Description}$	\sim Units
Code		
1	Core Heater Temperature	F
$\overline{\mathbf{c}}$	Fluid Temperature	F
3	Pressure	psiq
4	Strain	
5	Volumetric Flow	gpm
6	Fluid Velocity	ft/s
$\overline{7}$	Force	1b
8	Length	in
9	Voltage	
10	Material Temperature	F
11	Current	Amp
12	Specific Volume	ft^3/lbm
13	Decibels	dB
14	Pressure	psi
15	Pressure	psia
16	Differential Pressure	psid
17	Density	$1bm/ft^3$
18	Power.	kW
19	Heat Flux	Btu/s*ft^2
20	H. T. Coeff.	Btu/s*ft^2*F
21	Surface Temperature	\mathbf{F}_{\perp}
22	Saturation Temperature	F
23	Enthalpy	Btu/lbm
24	Mass Flux	$1bm/s*ft^2$
25	Mass Flow	lbm/s
26	Integrated Mass Flow	lbm
27	Momentum Flux	lbm/ft*s^2
28	Fluid Velocity	ft/s
29 [°]	Pump Speed	rpm
30	Elevation	ft
31	Quality	
32	Normalized Power	
33	Mass Flux	10e6 1bm/hr*ft^2
34	Temperature	F
35 ⁷	Time After Rupture	s
36	Time	s
37	Total Energy \mathcal{P}	Btu
38	Reactivity	\$
39	Stored Energy	Btu
40	Energy	Btu
41	Mass Balance	lbm
42	Power	MW

Table 3. Engineering Unit Codes **Table** 3. Engineering Unit Codes

--

 $\bar{\epsilon}$

1. Compression 1. Compression

The PIB data is compressed on an individual channel basis, as opposed to full file compression. This approach provides an efficient method of data retrieval, permitting direct access to the channel data and eliminating the need to uncompress the entire file prior to extracting the data for an individual channel. The compression method used is defined by the compression mode or cmpMode entry of the channel header record. Currently, method is defined by the compression mode or cmpMode entry of the channel header record. Currently, three compression modes are available as indicated in Table 4. compression modes are available as indicated in Table 4.

Table 4. Channel Compression Modes **Table** 4. Channel Compression Modes

Typically, compression is turned off for a channel (cmpMode **=** 0) if the achievable compression falls below 5%. Typically, compression is turned off for a channel (cmpMode = 0) if the achievable compression below 5%. This arbitrary threshold balances the potential savings in storage requirements against the overhead associated with uncompressing the data.

If the data channel does not vary over the entire range of time, a compression mode of cmpMode=1 is used. A single value is then stored in the channel data block to represent the entire range of data.

The double precision run length encoding compression method is typically used if the achievable compression exceeds 5%. Although it is a very simple algorithm, significant compression is achievable with very little impact on performance. In this method, each set of identical, consecutive values is replaced by two values; the first value being the number of values replaced and the second being the actual value. Regions that are not compressed are being the number of values replaced and the second being the actual value. Regions that are not compressed are preceded by a negative value that indicates the length of the uncompressed region. preceded by a negative value that indicates the length of the uncompressed region.

Table 5 illustrates application of the algorithm to a set of raw values. The initial two values 518.3 and 518.4 are replaced by the set of values -2.0, 518.4 and 518.4. The -2.0 value indicates that the next two values are not compressed. This is followed by the value 518.5 repeated twelve times which is replaced by the pair 12.0, 518.5. Similarly, the next four values have no repeats and are preceded by a value of 4.0, while the last eight values repeat and are replaced by the pair 8.0, 518.9. repeat and are replaced by the pair 8.0, 518.9.

Table 5. Example of Run Length Compression

The following two routines can be used to compress and uncompress the data, respectively. The following two routines can be used to compress and uncompress the data, respectively.

Source Listing 6. Compression Routines Source Listing 6. Compression Routines

```
/* define a temporary storage array used for compression */
/* define a temporary storage array used for compression */ 
#define DBUFSIZE 60000
#define DBUFSIZE 60000 
static double flbuf[DBUFSIZE];
static double flbuf[DBUFSIZE]; 
I**
  *
  *
  *
  *
  *
  *
  *
  *
  *
  *
Function: cmpres
/* Function: cmpres 
  Purpose: compress an array of double precision numbers
* Purpose: compress an array of double precision numbers 
  Arguments:
* Arguments: 
  data on entry contains raw channel data
* data on entry contains raw channel data 
  on exit contains compressed channel data
* on exit contains compressed channel data 
  size on entry contains length of raw channel data
* size on entry contains length of raw channel data 
  on exit contains length of compressed channel data
* on exit contains length of compressed channel data 
  /
*/ 
  Returns:
* Returns: 
  compression mode
* compression mode 
  \star 0 = no compression
  \begin{array}{r} 0 = 10 \text{ complex} \end{array}<br>
\begin{array}{r} 1 = \text{flat} \end{array} channel
  2 = run length compression
* 2 run length compression 
int cmpres(double **data, int *size) f
int cmpres(double **data, int *size) 
    int i, j;
    int cm;
int cm; 
    int reps,difs,dpos;
int reps,difs,dpos; 
    if(*size > DBUFSIZE)
if(*size > DBUFSIZE) 
        fprintf(stderr,"Channel length, %d, exceeds DBUFSIZE",*size);
fprintf(stderr,"Channel length, %d, exceeds DBUFSIZE",*size); 
{
```

```
exit (-1);
exit (-1); 
\mathbf{I}cm = 0;j = 0;
j= 0; 
dpos = -1;
dpos = -1; 
reps = 0;
reps = 0; 
\text{difs} = 0;for(i=l; i<*size; i++)
for(i=l; i<*size; i++) { 
    if(data[0][i] == data[0][i-1])
        if(difs) f
if (difs) { 
           {\text{rlbuf}}[{\text{dpos}}] = ({\text{double}}) (-1 \text{*difs});difs = 0;
           dpos = -1;\mathbf{F}reps++;
reps++; 
        else I
else { 
        if(reps)
if (reps) 
           reps++;
reps++; 
           flbuf[j++] = (double) \text{reps};\text{flow} = \text{down} = \text{temp}<br>\text{flow} = \text{data}[0][i-1];reps = 0;
reps = 0; 
           else
else { 
           if(dpos == -1)
if (dpos == -1) 
               dpos = j; j++;
dpos = j; 
              j++; 
            I
} 
           flbuf[j++] = data[0] [i-1];
           diffs++;\mathbf{)}\mathcal Y} 
if (reps)
if(reps) 
    reps++;
reps++; 
    fllbuf[j++] = (double)reps;flbuf[j++] = data[O] [i-l];
flbuf[j++] = data[O] [i-1]; 
} else if(difs) I
else if (difs) { 
    flbuf[j++] = data[O] [i-l];
flbuf[j++] = data[O] [i-1]; 
    difs++;
difs++; 
    flbuf[dpos] = (double) (-1*diffs);else f
else { 
    f_{\text{lbuf}}(j_{++}] = 1.0;f1[1] + f = 1.0;<br>f1[1] + f = 1.0;<br>f1[1] + f = 1.0;} 
if(j >= 0.95* (*size))
if(j >= 0.95* (*size)) { em = 0; 
   \overline{cm} = 0;else if (reps == *size)
else if (reps == *size) 
    *size = 1;
*size = 1; 
    cm = 1;else {<br>cm = 2;
else
    \frac{1}{2}*size = j;
    *size = j;
    for(i=O; i<*size; i++)
    for(i=O; i<*size; i++) { 
       data[O][i] = flbuf[i];
data[O] [i] = flbuf[i]; 
   \mathbf{1}return cm;
return em;
```
I

PIB FILE SPECIFICATION

```
/* Function: uncmpres
/* Function: uncmpres 
  * Purpose: uncompress an array of double precision numbers
* Purpose: uncompress an array of double precision numbers 
  * Arguments:
* Arguments: 
  * data on entry contains compressed channel data
* data on entry contains compressed channel data 
  * on exit contains raw channel data
* on exit contains raw channel data 
  * dblbuf pointer to temporary array space
* dblbuf pointer to temporary array space 
  * csize on entry contains length of compressed channel data
* csize on entry contains length of .compressed channel data 
  * fsize on entry contains length of raw channel data
* fsize on entry contains length of raw channel data 
  * cm on entry contains compression mode
* cm on entry contains compression mode 
  * 0 = no compression
  * 1 = flat channel
  * 2 = run length compression
* 2 run length compression 
  * Returns:
* Returns: 
  * 0 */
 \star 0
 */ 
int uncmpres(double *data, double *dblbuf, int* csize, int fsize, int cm) f
int uncmpres(double *data, double *dblbuf, int* csize, int fsize, int cm) 
{ 
   int i,j,k;int difs, reps;
int difs, reps; 
   switch (cm)
switch (cm) 
   case 0:
case 0: 
       break; /* no compression */
/* no compression */ 
      break; 
   case 1:
case 1: 
       \text{for}(i = 0; i\leq \text{fsize}; i++)data[i] = data[0];} 
       break; /* flat */
break; /* flat */ 
   case 2:
case 2: 
       k = i = 0;while(i<*csize)
while (i<*csize) ( 
           if(data[i] < 0.0)
if(data[ij < 0.0) { 
              \diff = (int) ( -1*data[i] + 0.1 );<br>diff = (int) ( -1*data[i] + 0.1 );
                i++; 
                for(j = 0; j<difs; j++)
for (j = 0; j<difs; j++) { 
                    dblbuf[k++] = data[i++];
                } 
              else if(data[i] > 0.0)
else if(data[ij > 0.0) { 
               r = \frac{r}{1} (\frac{r}{1} \frac{r}{1} \frac{r}{1} \frac{r}{1} \frac{r}{1} \frac{r}{1} \frac{r}{1} \frac{r}{1} \frac{r}{1} \frac{r}{1}f = \frac{1}{10} \frac{1}{100} \frac{1}{100} \frac{1}{100} \frac{1}{100} \frac{1}{100}\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} = data[i+1];
              i+=2;
i+=2; 
              else
else { 
               return 1; /* Error uncompressing data cm=2 */
          \rightarrowif(k != fsize) f
if (k != fsize) { 
       } 
           return 2; /* Error uncompressing data cm=2 */
return 2; /* Error uncompressing data cm=2 */ 
       ,<br>for(i = 0; i<fsize; i++)
         data[i] = dblbuf[i];} 
       break; /* run length compression */
break; /* run length compression */ 
   return 0;
return 0; 
}
}
```