DATA STRUCTURES & ALGORITHMS

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Lecture 02 Searching Algorithms

Searching

Finding elements in large amounts of data

Determine whether array contains value matching key value

☐ Linear search

☐Binary search

Searching an Array

Linear searchsmall arraysunsorted arrays

* Binary searchlarge arrayssorted arrays

Linear Search Algorithm

- Brute force algorithm, that starts searching the key from one end of the data set and proceed searching to the other end of the data set – Exhaustive Search.
 - Start at first element of the array.
 - Compare the value with the key
 - Continue with next element until a match is found or reach end of the array.

Note: On the average we have to compare the search key with half the elements in the array.

Linear Search - Example

* Array numlist contains:

| 17 | 23 | 5 | 11 | 2 | 29 | 3 |
|----|----|---|----|---|----|---|
| | | | | | | |

- * Searching for the the value 11: linear search examines 17, 23, 5, and 11
- * Searching for the the value 7: linear search examines 17, 23, 5, 11, 2, 29, and 3

Linear Search

* Algorithm:

```
set found to false; set position to -1; set index to 0
while index < number of elements. and found is false
   if list[index] is equal to search value
             found = true
            position = index
    end if
             add 1 to index
end while
return position
```

A Linear Search Function

```
int searchList(int list[], int numElems, int value)
   int index = 0; // Used as a subscript to search array
   int position = -1; // To record position of search value
  boolean found = false; // Flag to indicate if value was
                                                     found
  while (index < numElems && !found)
      if (list[index] == value) // If the value is found
         found = true; // Set the flag
        position = index; // Record the value's subscript
      index++; // Go to the next element
return position; // Return the position, or -1
```

Linear Search

* Benefits:

- Easy to understand and implement
- No assumption on data can be in any order

* Disadvantages:

- * Inefficient (slow) \rightarrow O(n)
- * Best case scenario?
- * Average complexity?
- * Worst case scenario?

Binary Search Algorithm

Need the dataset to be sorted (say ascending order)

```
At each step check the middle item

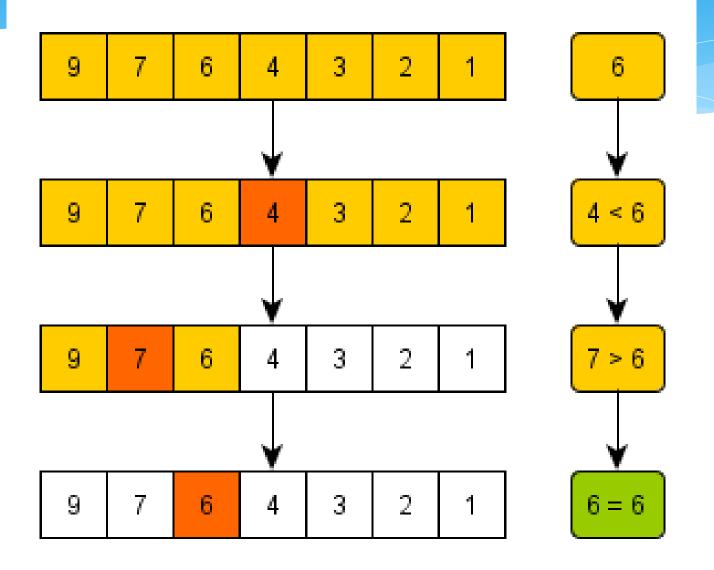
If key == middle item → item found

If key < middle item → binary search the first half

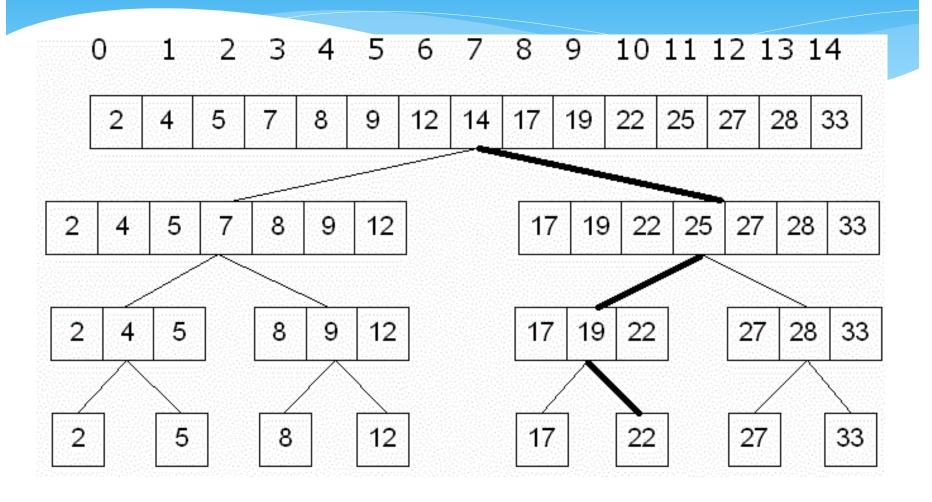
If key > middle item → binary search the second half
```

Search ends either finding the item or after checking an array of size 1.

Binary Search Example



Binary Search Tree



Efficiency

Searching and array of 1024 elements will take at most 10 passes to find a match or determine that the element does not exist in the array.

512, 256, 128, 64, 32, 16, 8, 4, 2, 1

- ➤ An array of one billion elements takes a maximum of 30 comparisons.
- The bigger the array the better a binary search is as compared to a linear search
- \triangleright O(log n)

Binary Search

Set first index to 0. Set last index to the last subscript in the array. Set found to false. Set position to -1.

While found is not true and first is less than or equal to last

Set middle to the subscript half-way between array[first] and array[last].

```
If array[middle] equals the desired value
Set found to true.
Set position to middle.
```

Else If array[middle] is greater than the desired value Set last to middle - 1.

Else

Set first to middle + 1.

End If.

End While.
Return position.

A Binary Search Function

```
int binarySearch(int array[], int size, int value)
  int first = 0, // First array element
    middle,
                   // Mid point of search
    position = -1; // Position of search value
  while (!found && first <= last)</pre>
    middle = (first + last) / 2; // Calculate mid point
    found = true;
      position = middle;
    else if (array[middle] > value) // If value is in lower half
      last = middle - 1;
    else
      first = middle + 1;
                           // If value is in upper half
  return position;
```

Carryout a Desk-Check for the above algorithm for the following data set:

| Var | 1st call | 2nd call | 3rd call | 4th call |
|-----------|----------|----------|----------|----------|
| | | | | |
| key | | | | |
| data | | | | |
| size | | | | |
| mid | | | | |
| data[mid] | | | | |
| return | | | | |

Answer:

| Var | 1st call | 2nd call | 3rd call | 4th call |
|-----------|----------|----------|----------|----------|
| key | 54 | 54 | 54 | 54 |
| data | &data[0] | &data[7] | &data[7] | &data[7] |
| size | 12 | 5 | 2 | 1 |
| mid | 6 | 2 | 1 | 0 |
| data[mid] | 47 | 65 | 64 | 55 |
| return | ? | ? | , | -1 |

Thank You