

Data structures III – Trees

Recursion Refresher - 1

- For some problems, it may be natural to define the problem in terms of the problem itself
- Recursion means that a function calls itself
- Useful for problems that can be represented by a simpler version of the same problem
- Example: the factorial function

$$6! = 6 * 5 * 4 * 3 * 2 * 1$$

- Can be written as:

$$6! = 6 * 5!$$

- Thus factorial function can be expressed as:

$$n! = n * (n-1)!$$

Recursion Refresher - 2

Iterative Solution

```
Public int fac(int numb){  
    int product=1;  
    while(numb>1){  
        product *= numb;  
        numb--;  
    }  
    return product;  
}
```

Recursive Solution

```
public int fac(int numb){  
    if(numb<=1)  
        return 1;  
    else  
        return numb*fac(numb-1);  
}
```

Recursion Refresher - 3

Exponent Function

```
int exp(int num, int power)
{
    if(power ==0)
        return 1;

    return num*exp(num, power-1);
}
```

General Form

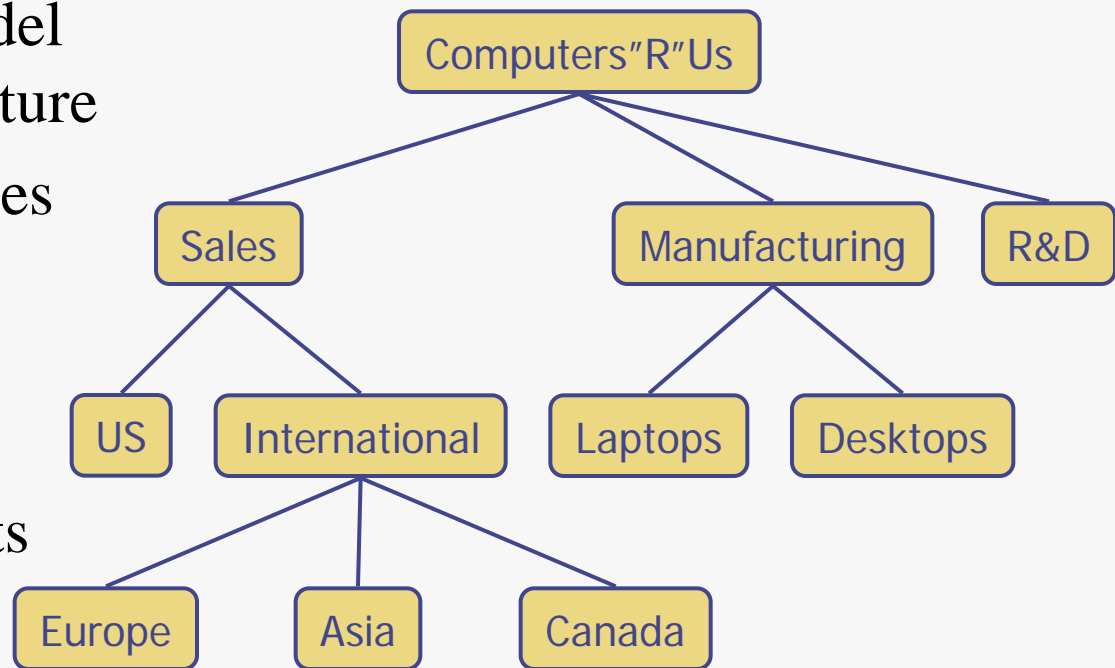
```
int recur_fn(param)
{
    if(stopping condition)
        return stopping value;

    return function of
    recur_fn(revised param);
}
```

TREES

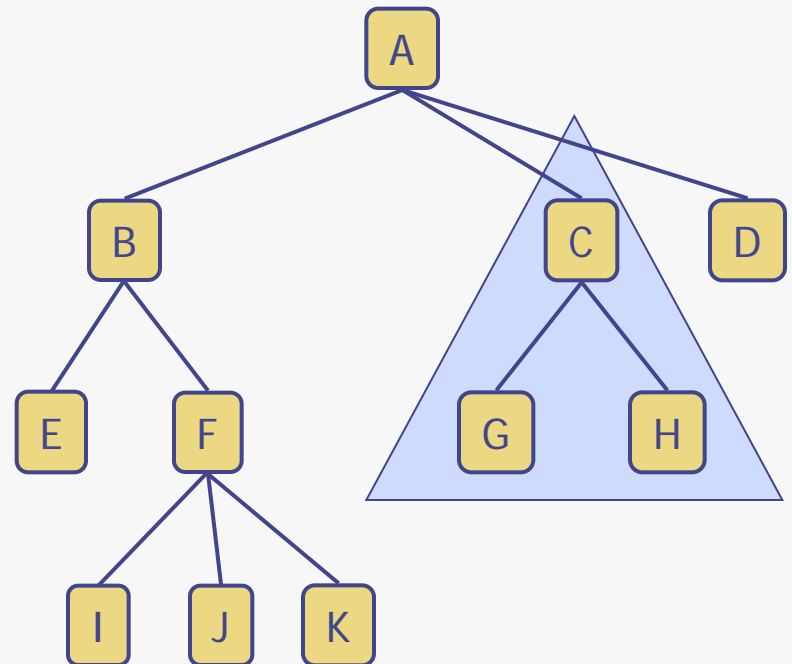
What is a Tree?

- In computer science, a tree is an abstract model of a hierarchical structure
- A tree consists of nodes with a parent-child relation
- Applications:
 - Organization charts
 - File systems
 - Programming environments



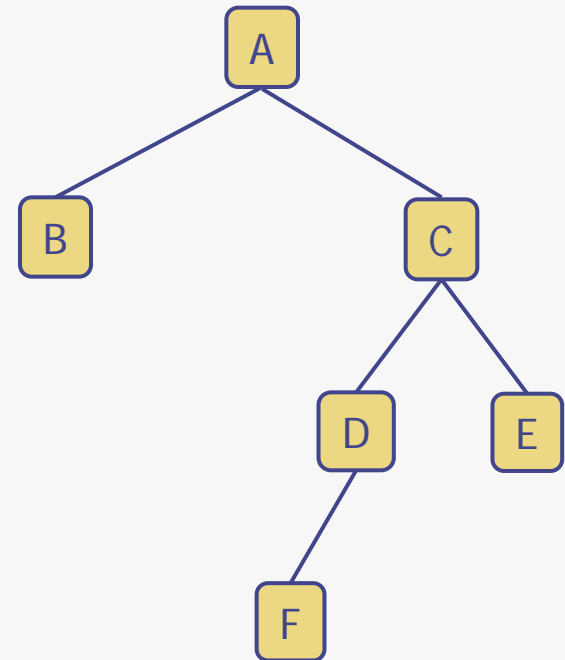
Tree Terminology

- **Root:** node without parent (A)
- **Internal node:** node with at least one child (A, B, C, F)
- **External node** (a.k.a. **leaf**): node without children (E, I, J, K, G, H, D)
- **Ancestors** of a node: parent, grandparent, grand-grandparent, etc. (B is grandparent of I)
- **Descendant** of a node: child, grandchild, grand-grandchild, etc.
- **Depth** of a node: number of ancestors (depth of E is 2)
- **Height** of a tree: maximum depth of any node (3)
- **Subtree:** tree consisting of a node and its descendants



Traversing a Tree

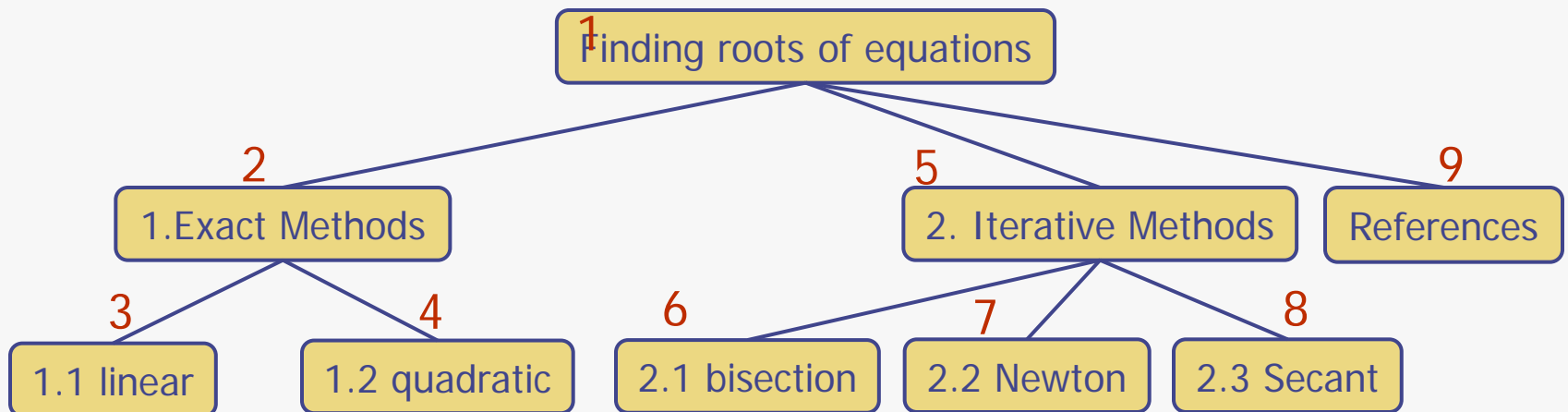
- Trees can be traversed in three orders
- Order chosen depends on
 - What the tree is being used for
 - What the traversal is supposed to accomplish
- Traversal is done **recursively!**
 - Treat L, R as trees in their own right
 - Recursively visit them
- **Pre-order:** root, L, R
 - A, B, C, D, F, E
- **In-order:** L, root, R
 - B, A, F, D, C, E
- **Post-order:** L, R, root
 - B, F, D, E, C, A



Preorder Traversal

- A traversal visits the nodes of a tree in a systematic manner
- In a preorder traversal, **a node is visited before its descendants**
- Application: print a structured document, e.g. Table of Contents

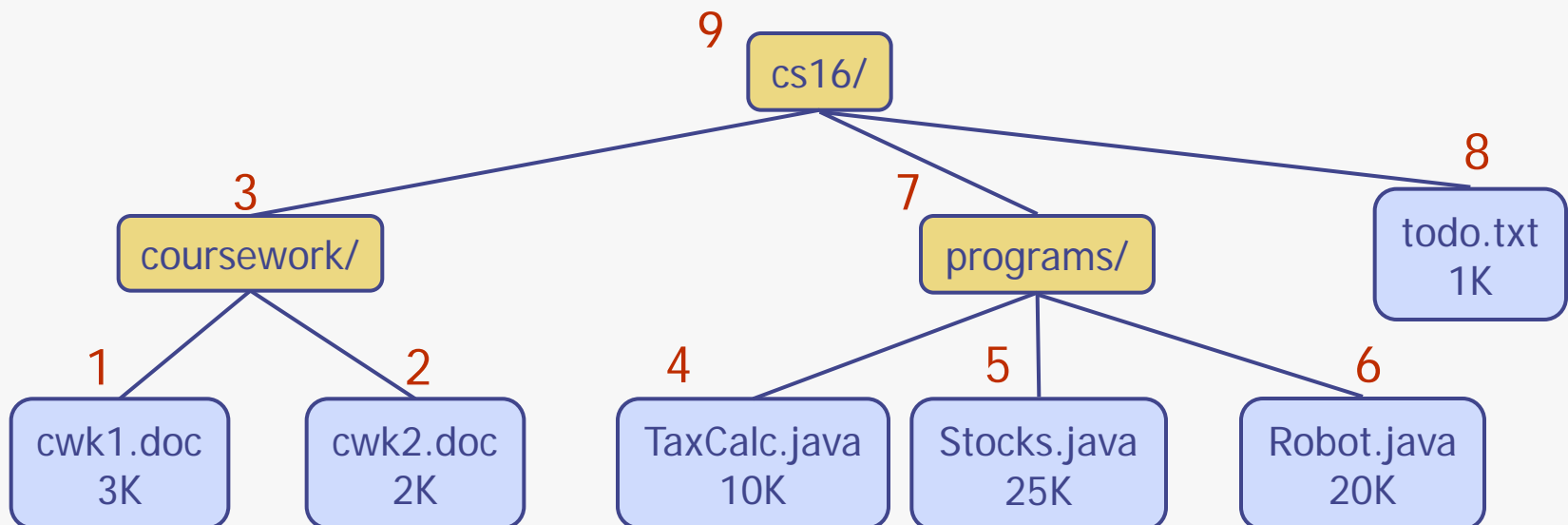
Algorithm *preOrder(v)*
visit(v)
for each child *w* of *v*
preOrder(w)



Postorder Traversal

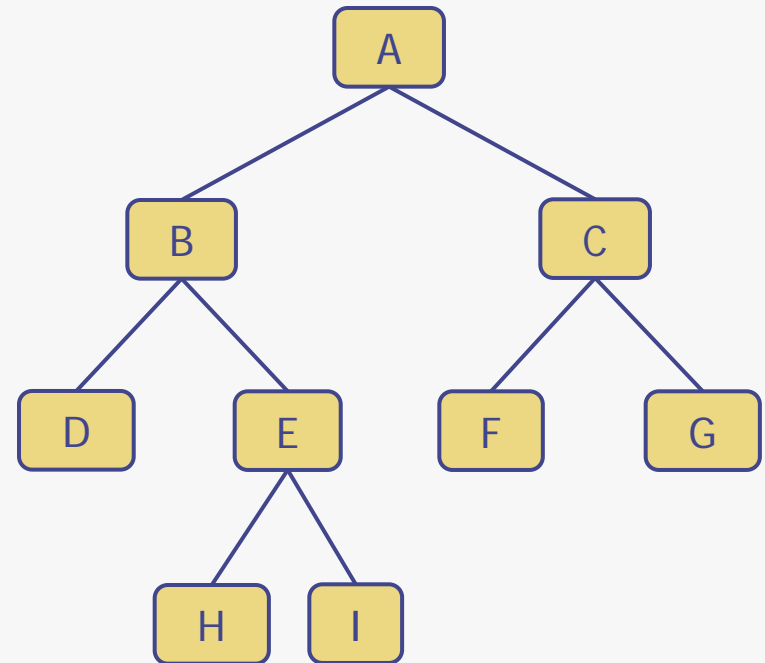
- Visits nodes in this order
- In a postorder traversal, **a node is visited *after* its descendants**
- Application: disk space used by files in a directory and its subdirectories

Algorithm *postOrder(v)*
for each child *w* of *v*
 postOrder(w)
visit(*v*)



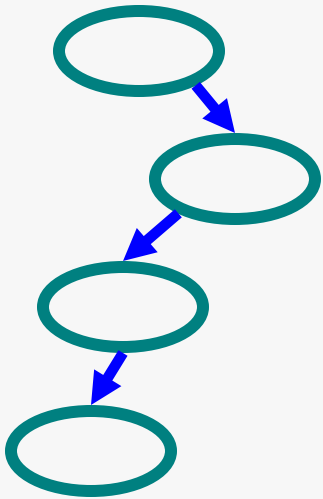
Binary Trees

- A binary tree is a tree with the following properties:
 - Each internal node has at most two children (exactly two for *proper* binary trees)
 - The children of a node are an ordered pair
- We call the children of an internal node *left child* and *right child*

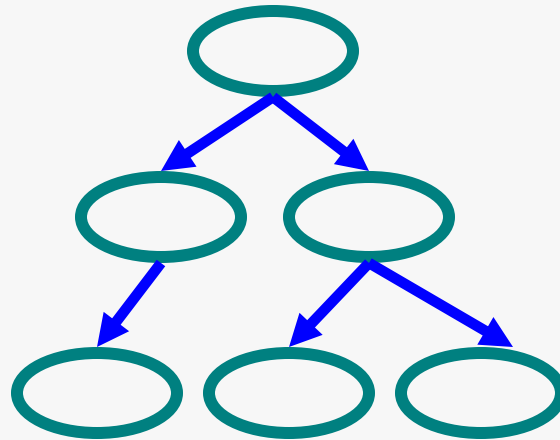


Types of Binary Trees

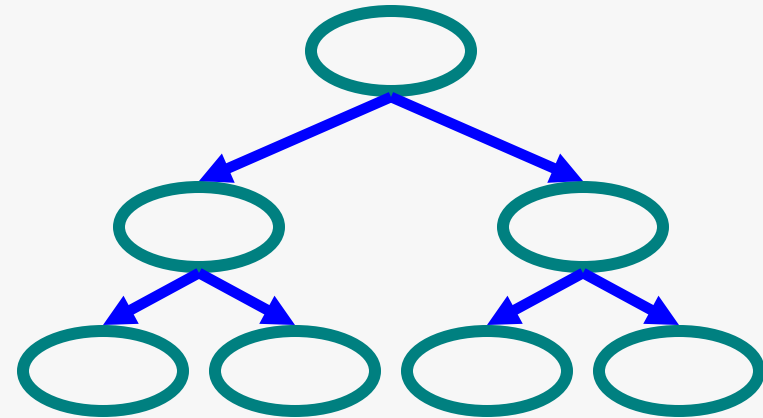
- Degenerate – only one child
- Balanced – mostly two children
- Proper – always two children



Degenerate
binary tree



Balanced binary
tree

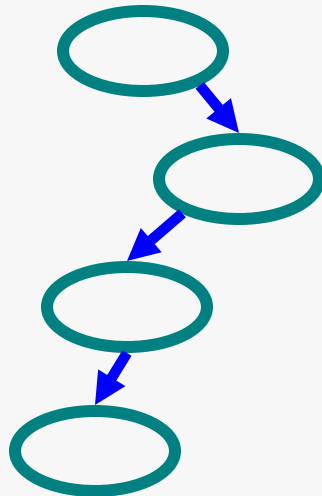


Proper binary
tree

Binary Trees Properties

- *Degenerate*

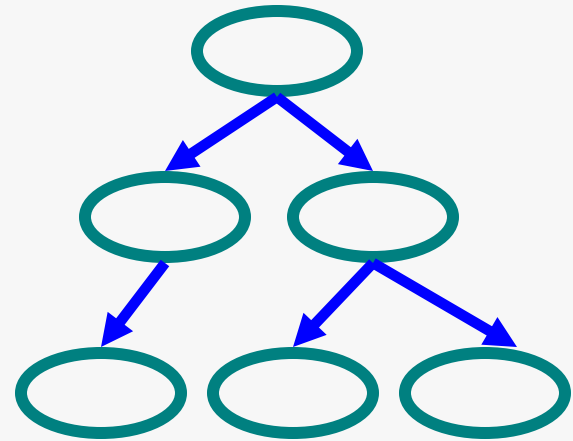
- Height = $O(n)$ for n nodes
- Similar to linear list



Degenerate
binary tree

- *Balanced*

- Height = $O(\log_2(n))$ for n nodes
- Useful for searches



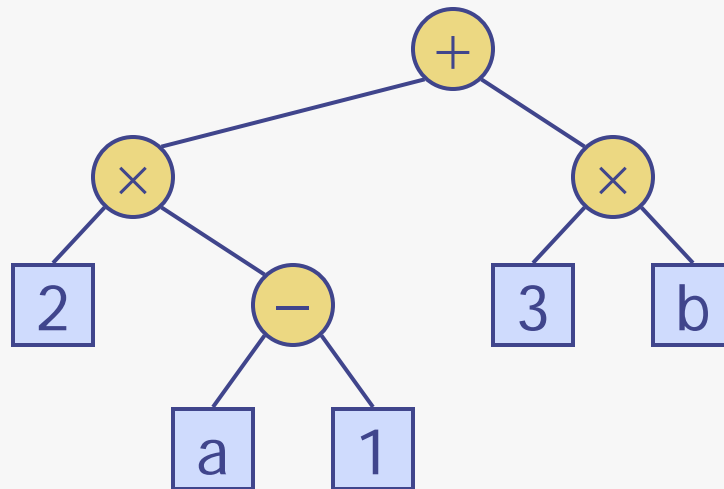
Balanced binary
tree

Simple example applications

- arithmetic expressions
- decision processes
- Searching (later lecture)

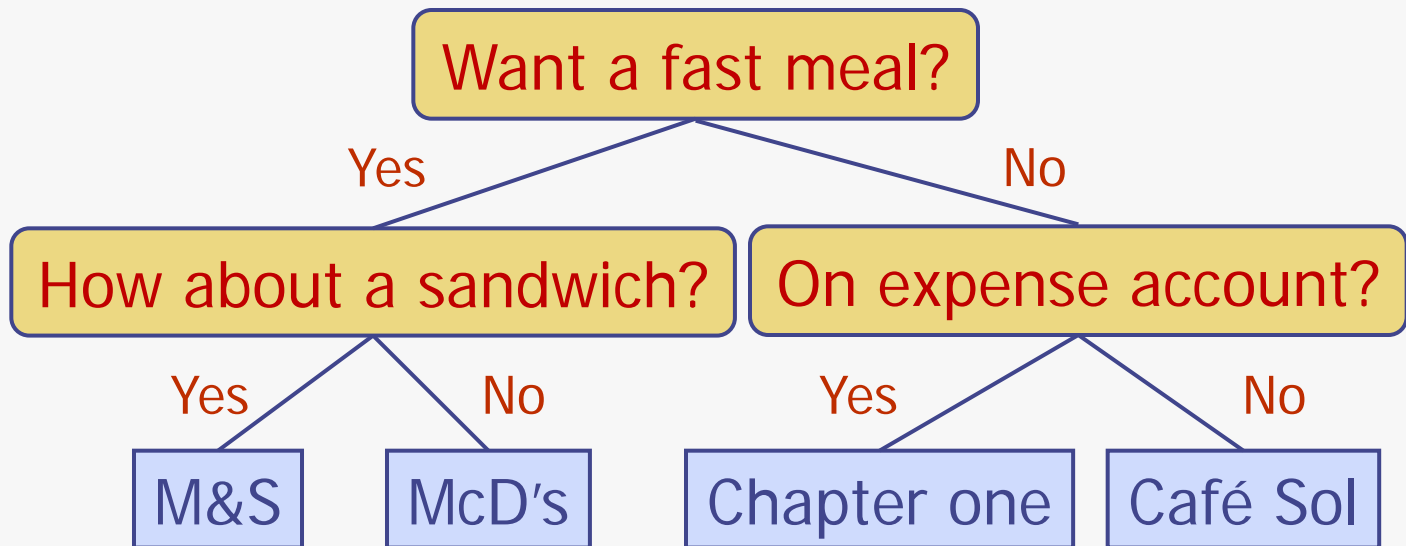
Arithmetic Expression Tree

- Binary tree associated with an arithmetic expression
 - internal nodes: operators
 - external nodes: operands
- Example: arithmetic expression tree for the expression $(2 \times (a - 1) + (3 \times b))$



Decision Binary Tree

- Binary tree associated with a decision process
 - internal nodes: questions with yes/no answer
 - external nodes: decisions
- Example: dining decision



Binary Tree ADT

- Some necessary methods that would be part of the Binary Tree Class:
 - TreeNode **getLeft**()
 - TreeNode **getRight**()
 - boolean **hasLeft**(p)
 - boolean **hasRight**(p)
- Update methods may be defined by data structures implementing the Binary Tree Class

Inorder Traversal

- In an inorder traversal a node is visited after its left subtree and before its right subtree
- Visit nodes in this order
 - Left subtree
 - Root node
 - Right subtree
- Application: to produce a drawing of a binary tree

Algorithm *inOrder*(*v*)

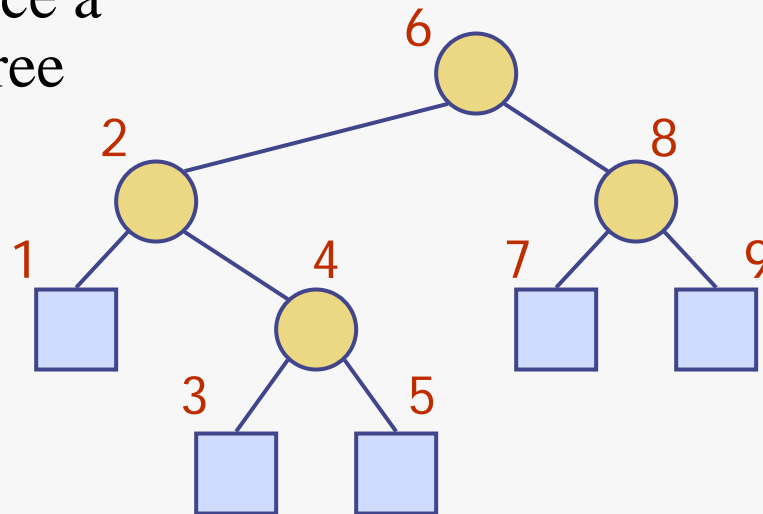
if *hasLeft* (*v*)

inOrder (*getLeft* (*v*))

visit(*v*)

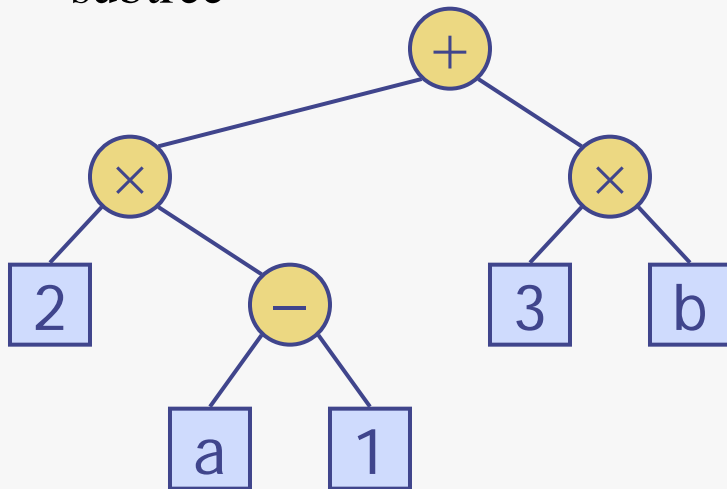
if *hasRight* (*v*)

inOrder (*getRight* (*v*))



Print Arithmetic Expressions

- Specialization of an inorder traversal
 - print operand or operator when visiting node
 - print “(“ before traversing left subtree
 - print “)” after traversing right subtree



Algorithm *printExpr* (*v*)

if *hasLeft* (*v*)

print("(")

printExpr(*left*(*v*))

print(*v.element* ())

if *hasRight* (*v*)

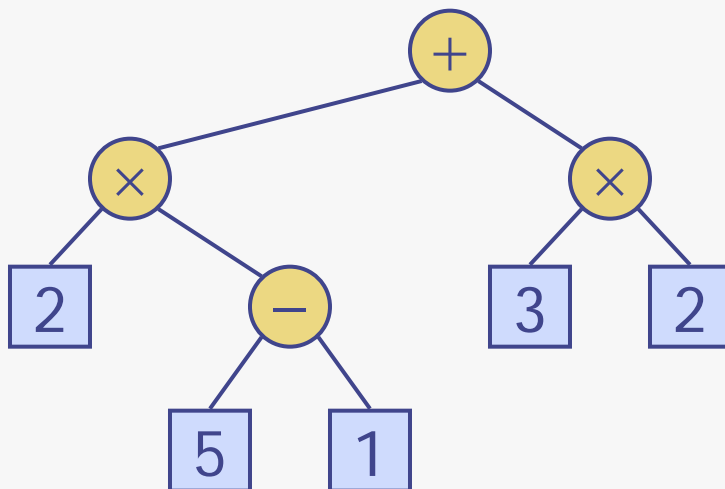
printExpr(*right*(*v*))

print (")")

$((2 \times (a - 1)) + (3 \times b))$

Evaluate Arithmetic Expressions

- Specialization of a postorder traversal
 - recursive method returning the value of a subtree
 - when visiting an internal node, combine the values of the subtrees



Algorithm *evalExpr(v)*

if *!hasLeft(v) && !hasRight(v)*

return *v.element ()*

else

x \leftarrow *evalExpr(left (v))*

y \leftarrow *evalExpr(right (v))*

$\diamond \leftarrow$ operator stored at *v*

return *x* \diamond *y*

x \diamond *y* \leftarrow 5-1

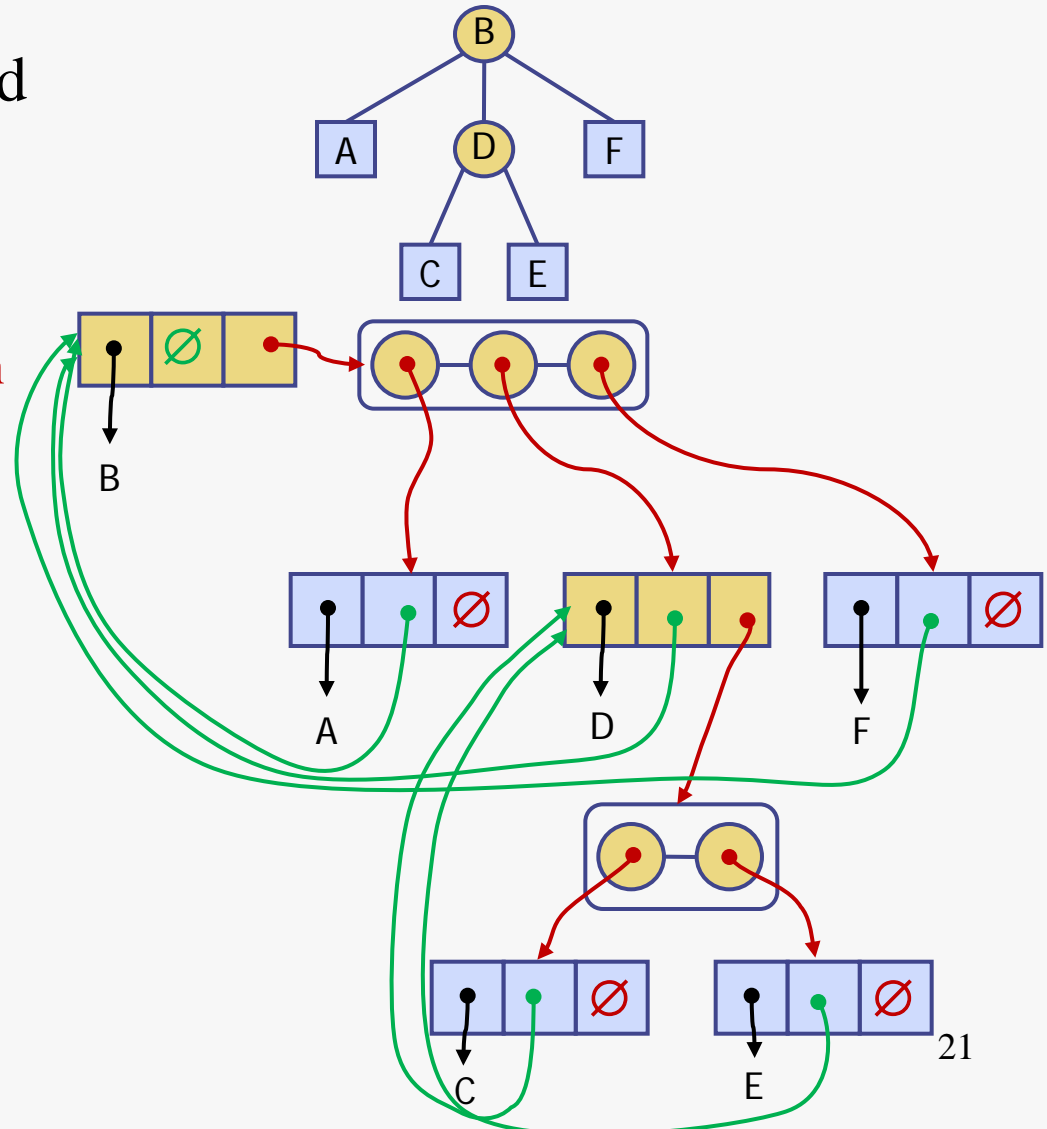
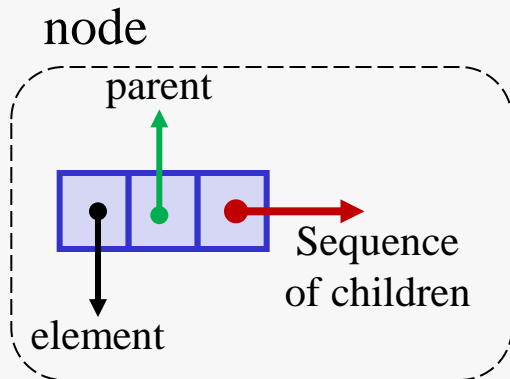
x \diamond *y* \leftarrow 2 x (5-1)

x \diamond *y* \leftarrow 3 x 2

x \diamond *y* \leftarrow 2 x (5-1) + (3 x 2)

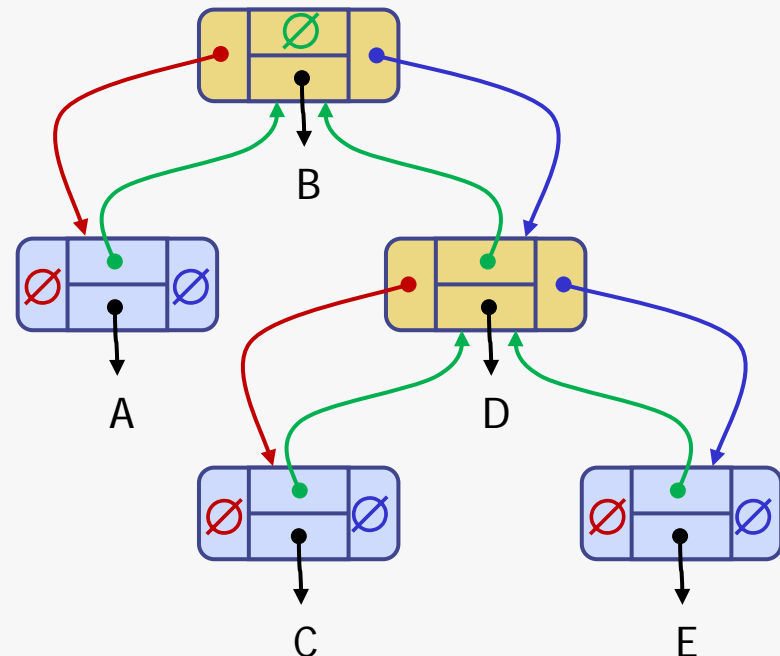
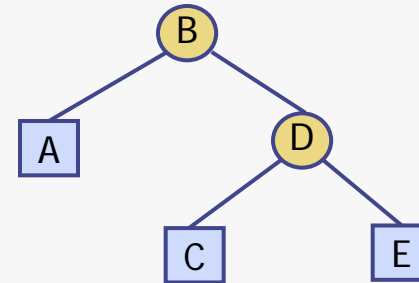
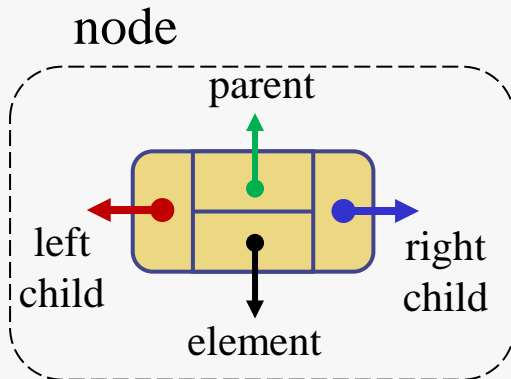
Linked Structure for Trees

- A node is represented by an object storing
 - Element
 - Parent node
 - Sequence of children nodes



Linked Structure for Binary Trees

- A node is represented by an object storing
 - Element
 - Parent node
 - Left child node
 - Right child node



Acknowledgements

- Some of this material has been taken from a variety of sources
- the most notable of which is from
*Goodrich and Tamassia “Data Structures and Algorithms in Java”
John Wiley & Sons.*