

NATIONAL INSTITUTE OF STANDARDS AND
TECHNOLOGY

Intelligent Systems Division

Knowledge Driven Planning and
Modeling for Part Handling

Planning for a Kitting Workstation

Stephen Balakirsky

stephen.balakirsky@nist.gov

Zeid Kootbally

zeid.kootbally@nist.gov

Thomas Kramer

thomas.kramer@nist.gov

Anthony Pietromartire

pietromartire.anthony@nist.gov

Craig Schlenoff

craig.schlenoff@nist.gov

Contents

1	Introduction	1
2	Domain Specific Information	4
2.1	Constant Variable Symbols	5
2.2	Object Variable Symbols	6
2.3	State Variable Symbols	6
2.4	Predicates and Functions	8
2.4.1	Predicates	9
2.4.2	Functions	11
2.5	Planning Operators and Actions	12
2.5.1	Planning Operators	12
2.5.2	Actions	21
3	Ontology	22
4	Planning Language	23
4.1	The PDDL Domain File	23
4.2	PDDL Problem File	26
4.2.1	Initial State	29
4.2.2	Goal State	30
4.3	Plan	31
5	Robot Language	34
6	Planner	35
6.1	Requirements	35
6.2	Download and Install CBC	35

6.3	Compile the Planner	36
6.4	Run the Planner	37
7	Tools	38
7.1	The Generator	38
7.1.1	Prequisites	38
7.1.2	How to Run the Generator Tool	39
7.1.3	Functionalities	39
7.2	XML to OWL	42
7.2.1	owlPrinter	42
7.2.2	kittingParser	43
7.2.3	compactOwl and compareOwl	44

1 Introduction

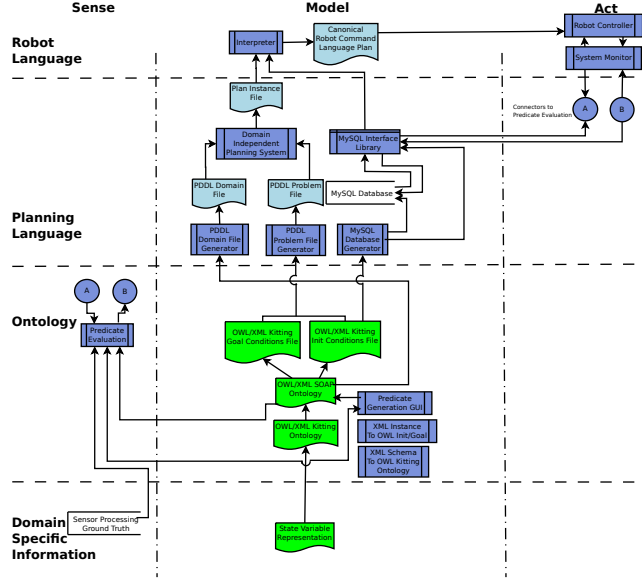


Figure 1: Knowledge Driven Design extensions – In this figure, green shaded boxes with curved bottoms represent hand generated files while light blue shaded boxes with curved bottoms represent automatically created boxes. Rectangular boxes represent processes and libraries.

The knowledge driven methodology presented in this section is not purposed to act as a stand-alone system architecture. Rather it is intended to be an extension to well developed hierarchical, deliberative architectures such as 4D/RCS [1]. The overall knowledge driven methodology of the system is depicted in Figure 1. The figure is organized vertically by the representation that is used for the knowledge and horizontally by the classical sense-model-act paradigm of intelligent systems. The remainder of this section gives a brief description of each level of the hierarchy to help the reader understand the basic concepts implemented within the system architecture. The reader may find a more detailed description of each component and each level of the architecture in other documented publications [2, 3].

- On the vertical axis, knowledge begins with Domain Specific Information (DSI) that captures operational knowledge that is necessary to be successful in the particular domain in which the system is designed to operate. This includes information on items ranging from what actions

and attributes are relevant, to what the necessary conditions are for an action to occur and what the likely results of the action are. The authors have chosen to encode this basic information in a formalism known as a state variable representation [7].

- At the next level up, the information encoded in the DSI is then organized into a domain independent representation. A base ontology (OWL/XML Kitting) contains all of the basic information that was determined to be needed during the evaluation of the use cases and scenarios. The knowledge is represented in as compact a form as possible with knowledge classes inheriting common attributes from parent classes. The OWL/XML SOAP ontology describes not only aspects of actions and predicates but also the individual actions and predicates that are necessary for the domain under study. The instance files describe the initial and goal states for the system through the **Kitting Init Conditions File** and the **Kitting Goal Conditions File**, respectively. The initial state file must contain a description of the environment that is complete enough for a planning system to be able to create a valid sequence of actions that will achieve the given goal state. The goal state file only needs to contain information that is relevant to the end goal of the system.

Since both the OWL and XML implementations of the knowledge representation are file based, real time information proved to be problematic. In order to solve this problem, an automatically generated MySQL database has been introduced as part of the knowledge representation.

- At the next level up, aspects of this knowledge are automatically extracted and encoded in a form that is optimized for a planning system to utilize (the Planning Language). The planning language used in the knowledge driven system is expressed with the Planning Domain Definition Language (PDDL) [6] (version 3.0). The PDDL input format consists of two files that specify the domain and the problem. As shown in Figure 1, these files are automatically generated from the ontology. The domain file represents actions along with required preconditions and expected results. The problem file represents the initial state of the system and the desired goal. From these two files, a domain independent planning system [4] was used to produce a static Plan Instance File.
- Once a plan has been formulated, the knowledge is transformed into a

representation that is optimized for use by a robotic system (the Robot Language). The interpreter combines knowledge from the plan with knowledge from the MySQL database to form a sequence of sequential actions that the robot controller is able to execute. The authors devised a canonical robot command language (CRCL) in which such lists can be written. The purpose of the CRCL is to provide generic commands that implement the functionality of typical industrial robots without being specific either to the language of the planning system that makes a plan or to the language used by a robot controller that executes a plan.

This document describes each level of the architecture from the Domain Specific Information level to the Robot Language level. The goal of this document is to help a user understand the different tools and techniques to generate the appropriate files in order to build a kit.

Each of the following section corresponds to each level of the architecture described in Figure 1 starting from the bottom level to the top level of the architecture.

2 Domain Specific Information

The foundation for the knowledge representation is domain specific information that is produced by an expert in the particular field of study. This includes information on items ranging from what actions and attributes are relevant, to what the necessary conditions are for an action to occur and what the likely results of the action are. We have chosen to encode this basic information in a formalism known as a state variable representation (SVR) [7]. This information will then flow up the abstraction and be transformed into the ontology, planning language, and robot language.

Before building a SVR, the domain for kitting needs to be specified. The domain for kitting contains some fixed equipment: a robot, a work table, end effectors, end effector holders, and an end effector changing station. Items that enter the workstation include kit trays, boxes in which to put kits, boxes that contain empty kit trays, and part supplies. Items that leave the workstation may be boxes with finished kits inside, empty part trays, empty boxes. An external agent is responsible of moving the items that leave the workstation. We assume that the workstation has only one work table, one changing station, and one robot.

In a State Variable Representation (SVR), each state is represented by a tuple of values of n state variables $\{x_1, \dots, x_n\}$, and each action is represented by a partial function that maps this tuple into some other tuple of values of the n state variables.

To build the SVR, the group has taken a very systematic approach of identifying and modeling the concepts. Because the industrial robot field is so broad, the group decided to limit its efforts to a single type of operation, namely kitting. A scenario was developed that described, in detail, the types of operations that would be performed in kitting, the sequencing of steps, the parts and machines that were needed, constraints on the process such as pre- and post-conditions, etc. For this scenario, a set of concepts were extracted and defined. These concepts served as the initial requirements for the kitting SVR. The concepts were then modeling in our SVR, building off of the definitions and relationships that were identified in the scenario. A SVR relies on the elements of constant variable symbols, object variable symbols, state variable symbols, and planning operators. These are defined for the kitting domain in the rest of this section.

2.1 Constant Variable Symbols

For the kitting domain, there is a finite set of constant variable symbols that must be represented. In the SVR, constant variable symbols are partitioned into disjoint classes corresponding to the objects of the domain. The finite set of all constant variable symbols in the kitting domain is partitioned into the following sets:

- A set of *Part*: A *Part* is the basic item that will be used to fill a kit.
- A set of *PartsTray*: *Parts* arrive at the workstation in *PartsTrays*. Each *Part* is at a known position in the *PartsTray*. Each *PartsTray* contains one type of *Part*.
- A set of *KitTray*: A *KitTray* can hold *Parts* in known positions.
- A set of *Kit*: A *Kit* consists of a *KitTray* and, possibly, some *Parts*. A *Kit* is empty when it does not contain any *Part* and finished when it contains all the *Parts* that constitute a kit.
- A set of *WorkTable*: A *WorkTable* is an area in the kitting workstation where *KitTrays* are placed to build *Kits*.
- A set of *LargeBoxWithKits*: A *LargeBoxWithKits* contains only finished *Kits*.
- A set of *LargeBoxWithEmptyKitTrays*: A *LargeBoxWithEmptyKitTrays* is a box that contains only empty *KitTrays*.
- A set of *Robot* $\{robot_1, robot_2, \dots\}$: A *Robot* in the kitting workstation is a robotic arm that can move objects in order to build *Kits*.
- A set of *EndEffector*: *EndEffectors* are used in a kitting workstation to manipulate *Parts*, *PartsTrays*, *KitTrays*, and *Kits*. An *EndEffector* is attached to a *Robot* in order to grasp objects.
- A set of *EndEffHolder*: An *EndEffHolder* is a storage unit that holds one type of *EndEffector*.
- A set of *EndEffChStation*: An *EndEffChStation* is made up of *EndEffHolders*.

2.2 Object Variable Symbols

Object variable symbols are typed variables which range over a class or the union of classes of constant variable symbols. Examples of object variable symbols are $r \in Robots$, $kt \in KitTrays$, etc.

2.3 State Variable Symbols

A state variable symbol is defined as follows:

$x : A_1 \times \dots \times A_i \times S \rightarrow B_1 \cup \dots \cup B_j \cup bool \cup \{\} \cup numeric$ ($i, j \geq 1$) is a function from the set of states (S) and at least one set of constant variable symbols $A_1 \times \dots \times A_i$ into a set $B_1 \cup \dots \cup B_j \cup bool \cup \{\} \cup numeric$ where:

- $B_1 \cup \dots \cup B_j$ is a set of constant variable symbols
- *bool* is a boolean
- $\{\}$ is an empty set
- *numeric* is a numerical value

The use of state variable symbols reduces the possibility of inconsistent states and generates a smaller state space. The following state variable symbols are used in the kitting domain.

- **endeff-location**
 $EndEffector \times S \rightarrow Robot \cup EndEffHolder$: designates the location of an *EndEffector* in the workstation. An *EndEffector* is either attached to a *Robot* or placed in an *EndEffHolder*.
- **robot-with-endeff**
 $Robot \times S \rightarrow EndEffector \cup \{\}$: designates the *EndEffector* attached to a *Robot* if there is one attached, otherwise nothing.
- **on-wtable**
 $WorkTable \times S \rightarrow Kit \cup KitTray \cup \{\}$: designates the object placed on the *WorkTable*, i.e., a *Kit*, a *KitTray*, or nothing.
- **kit-location**
 $Kit \times S \rightarrow LargeBoxWithKits \cup WorkTable \cup Robot$: designates the different possible locations of a *Kit* in the workstation, i.e., in a *LargeBoxWithKits*, on the *WorkTable*, or being held by a *Robot*.

■ **kittray-location**

$KitTray \times S \rightarrow LargeBoxWithEmptyKitTrays \cup WorkTable \cup Robot$: designates the different possible locations of a *KitTray* in the workstation, i.e., in a *LargeBoxWithEmptyKitTrays*, on a *WorkTable* or being held by a *Robot*.

■ **part-location**

$Part \times S \rightarrow PartsTray \cup Kit \cup Robot$: designates the different possible locations of a *Part* in the workstation, i.e., in a *PartsTray*, in a *Kit*, or being held by a *Robot*.

■ **robot-holds**

$Robot \times S \rightarrow KitTray \cup Kit \cup Part \cup \{\}$: designates the object being held by a *Robot*, i.e., a *KitTray*, a *Kit*, a *Part*, or nothing. It is assumed that the *Robot* is already equipped with the appropriate *EndEffector*.

■ **lbwk-full**

$LargeBoxWithKits \times S \rightarrow bool$: designates if a *LargeBoxWithKits* is full or not.

■ **lbwekt-empty**

$LargeBoxWithEmptyKitTrays \times S \rightarrow bool$: designates if a *LargeBoxWithEmptyKitTrays* is empty or not.

■ **partstray-empty**

$PartsTray \times S \rightarrow bool$: designates if a *PartsTray* is empty or not.

■ **endeffector-type**

$EndEffector \times S \rightarrow KitTray \cup Kit \cup Part$: designates the type of object an *EndEffector* can hold, i.e., a *KitTray*, a *Kit*, or a *Part*.

■ **endeffholder-holds-endeff**

$EndEffHolder \times S \rightarrow EndEffector \cup \{\}$: designates whether an *EndEffHolder* is holding an *EndEffector* or nothing.

■ **endeffholder-location**

$EndEffHolder \times S \rightarrow EndEffChStation$: designates the *EndEffChStation* where the *EndEffHolder* is located.

■ **endeffchstation-has-endeffholder**

$EndEffChStation \times S \rightarrow EndEffHolder$: designates the *EndEffHolder* the *EndEffChStation* contains.

- **found-part**
 $PartsTray \times S \rightarrow Part \cup \{\}$: designates whether a *Part* is found in a *PartsTray* or not.
- **origin-part**
 $Part \times S \rightarrow PartsTray$: designates the *PartsTray* where the *Part* is found.
- **quantity-parts-in-partstray**
 $PartsTray \times S \rightarrow numeric$: designates the number of parts that *PartsTray* contains.
- **quantity-parts-in-kit**
 $Kit \times PartsTray \times S \rightarrow numeric$: designates the number of parts from *PartsTray* that *Kit* contains.
- **capacity-parts-in-kit**
 $Kit \times PartsTray \times S \rightarrow numeric$: designates the number of parts from *PartsTray* that *Kit* **can** contain.

2.4 Predicates and Functions

In PDDL, predicates are used to encode Boolean state variables, while functions are used to model updates of numerical values [5]. This section describes the predicates and functions derived from the state variables described in section 2.3. We recall the following definition of a state variable ((section 2.3)) $x : A_1 \times \dots \times A_i \times S \rightarrow B_1 \cup \dots \cup B_j$ ($i, j \geq 1$) that is used to convert state variables into predicates as follows:

- $A_1 \times \dots \times A_i \times S \rightarrow B_1 \cup \dots \cup B_j$ ($i, j \geq 1$)
 - `predicate_1(\mathcal{A}, \mathcal{B})`
 - ...
 - `predicate_n(\mathcal{A}, \mathcal{B})`

Where $\mathcal{A} \in \{A_1, \dots, A_i\}$ and $\mathcal{B} \in \{B_1, \dots, B_j\}$ ($i, j \geq 1$)

2.4.1 Predicates

The state variables in our current kitting domain contains the following predicates.

■ endeff-location

- ☐ `endeff-location-robot(EndEffector,Robot) ;TRUE` iff *EndEffector* is attached to *Robot*
- ☐ `endeff-location-endeffholder(EndEffector,EndEffHolder) ;TRUE` iff *EndEffector* is in *EndEffHolder*

■ robot-with-endeff

- ☐ `robot-with-endeff(Robot,EndEffector) ;TRUE` iff *Robot* is equipped with *EndEffector*
- ☐ `robot-with-no-endeff(Robot) ;TRUE` iff *Robot* is not equipped with any *EndEffector*

■ on-wtable

- ☐ `on-wtable-kit(WorkTable,Kit) ;TRUE` iff *Kit* is on the *WorkTable*
- ☐ `on-wtable-kittray(WorkTable,KitTray) ;TRUE` iff *KitTray* is on the *WorkTable*
- ☐ `worktable-empty(WorkTable) ;TRUE` iff there is nothing on the *WorkTable*

■ kit-location

- ☐ `kit-location-lbwk(Kit,LargeBoxWithKits) ;TRUE` iff *Kit* is in the *LargeBoxWithKits*
- ☐ `kit-location-wtable(Kit,WorkTable) ;TRUE` iff *Kit* is on the *WorkTable*
- ☐ `kit-location-robot(Kit,Robot) ;TRUE` iff *Kit* is being held by the *Robot*

■ kittray-location

- ☐ `kittray-location-lbwekt(KitTray,LargeBoxWithEmptyKitTrays) ;TRUE` iff *KitTray* is in the *LargeBoxWithEmptyKitTrays*

- ☐ $\text{kittray-location-robot}(KitTray, Robot)$;TRUE iff $KitTray$ is being held by the $Robot$
- ☐ $\text{kittray-location-wtable}(KitTray, WorkTable)$;TRUE iff $KitTray$ is on the $WorkTable$

■ part-location

- ☐ $\text{part-location-partstray}(Part, PartsTray)$;TRUE iff $Part$ is in the $PartsTray$
- ☐ $\text{part-location-kit}(Part, Kit)$;TRUE iff $Part$ is in the Kit
- ☐ $\text{part-location-robot}(Part, Robot)$;TRUE iff $Part$ is being held by the $Robot$

■ robot-holds

- ☐ $\text{robot-holds-kittray}(Robot, KitTray)$;TRUE iff $Robot$ is holding a $KitTray$
- ☐ $\text{robot-holds-kit}(Robot, Kit)$;TRUE iff $Robot$ is holding a Kit
- ☐ $\text{robot-holds-part}(Robot, Part)$;TRUE iff $Robot$ is holding a $Part$
- ☐ $\text{robot-empty}(Robot)$;TRUE iff $Robot$ is not holding anything

■ lbwk-full

- ☐ $\text{lbwk-not-full}(LargeBoxWithKits)$;TRUE iff $LargeBoxWithKits$ is not full

■ lbwekt-empty

- ☐ $\text{lbwekt-not-empty}(LargeBoxWithEmptyKitTrays)$;TRUE iff $LargeBoxWithEmptyKitTrays$ is not empty

■ partstray-empty

- ☐ $\text{partstray-not-empty}(PartsTray)$;TRUE iff $PartsTray$ is not empty

■ endeff-type

- ☐ $\text{endeff-type-kittray}(EndEffector, KitTray)$;TRUE iff $EndEffector$ is capable of holding a $KitTray$
- ☐ $\text{endeff-type-kit}(EndEffector, Kit)$;TRUE iff $EndEffector$ is capable of holding a Kit

□ `endeff-type-part(EndEffector,Part)` ;TRUE iff *EndEffector* is capable of holding a *Part*

■ `endeffholder-holds-endeff`

□ `endeffholder-holds-endeff(EndEffHolder,EndEffector)` ;TRUE iff *EndEffHolder* is holding *EndEffector*

□ `endeffholder-empty(EndEffHolder)` ;TRUE iff *EndEffHolder* is empty (not holding an *EndEffector*)

■ `endeffholder-location`

□ `endeffholder-location(EndEffHolder,EndEffChStation)` ;TRUE iff *EndEffHolder* is in *EndEffChStation*

■ `endeffchstation-has-endeffholder`

□ `endeffchstation-has-endeffholder(EndEffChStation,EndEffHolder)` ;TRUE iff *EndEffChStation* contains *EndEffHolder*

■ `found-part`

□ `found-part(Part,PartsTray)` ;TRUE iff *Part* is found in *PartsTray*

■ `origin-part`

□ `origin-part(Part,PartsTray)` ;TRUE iff *Part* is from *PartsTray*.

2.4.2 Functions

In a planning model, numeric fluents represent function symbols that can take an infinite set of values. Introducing functions into planning not only makes it possible to deal with numerical values in a more general way than allowed for by a purely relational language but makes it possible to model operators in a more compact and sometimes also more natural way. The state variables in our current kitting domain contains the following functions.

■ `quantity-parts-in-partstray`

□ `quantity-parts-in-partstray(PartsTray)` ;Quantity of parts in *PartsTray*

■ `quantity-parts-in-kit`

□ `quantity-parts-in-kit(Kit,PartsTray)` ;Quantity of parts from *PartsTray* that is in *Kit*

■ `capacity-parts-in-kit`

□ `capacity-parts-in-kit(Kit,PartsTray)` ;Quantity of parts from *PartsTray* that *Kit* **can** contain

2.5 Planning Operators and Actions

The planning operators presented in this section are expressed in classical representation instead of state variable representation. In classical representation, states are represented as sets of logical atoms (predicates) that are true or false within some interpretation. Actions are represented by planning operators that change the truth values of these atoms.

2.5.1 Planning Operators

In classical planning, a planning operator [7] is a triple $o = (\text{name}(o), \text{preconditions}(o), \text{effects}(o))$ whose elements are as follows:

- `name(o)` is a syntactic expression of the form $n(x_1, \dots, x_k)$, where n is a symbol called an operator symbol, x_1, \dots, x_k are all of the object variable symbols that appear anywhere in o , and n is unique (i.e., no two operators can have the same operator symbol).
- `preconditions(o)` and `effects(o)` are sets of literals (i.e., atoms and negations of atoms). Literals that are true in `preconditions(o)` but false in `effects(o)` are removed by using negations of the appropriate atoms.

Our kitting domain is composed of ten operators which are defined below.

1. `take-kittray(robot,kittray,lbwekt,endeffect,worktable)`: The *Robot robot* equipped with the *EndEffector endeffect* picks up the *KitTray kittray* from the *LargeBoxWithEmptyKitTrays lbwekt*.

<i>preconditions</i>	<i>effects</i>
<code>robot-empty(robot)</code>	\neg <code>robot-empty(robot)</code>
<code>kittray-location-lbwekt(kittray,lbwekt)</code>	\neg <code>kittray-location-lbwekt(kittray,lbwekt)</code>
<code>lbwekt-not-empty(lbwekt)</code>	<code>kittray-location-robot(kittray,robot)</code>
<code>robot-with-endeff(robot,endeff)</code>	<code>robot-holds-kittray(robot,kittray)</code>
<code>endeff-location-robot(endeff,robot)</code>	
<code>worktable-empty(worktable)</code>	
<code>endeff-type-kittray(endeff,kittray)</code>	

■ *preconditions*

- ☐ `robot-empty(robot)`: *robot* does not hold anything.
- ☐ `kittray-location-lbwekt(kittray,lbwekt)`: *kittray* is in *lbwekt*.
- ☐ `lbwekt-not-empty(lbwekt)`: *lbwekt* is not empty (contains at least one kit tray).
- ☐ `robot-with-endeff(robot,endeff)`: *robot* is equipped with *endeff*.
- ☐ `endeff-location-robot(endeff,robot)`: The end effector is on the robot's arm.
- ☐ `worktable-empty(worktable)`: After picking up an empty kit tray from a large box of empty kit trays, the robot would normally place the kit tray on the work table. To put a kit tray on the work table, it is necessary that there is nothing on top of the work table. If the robot is allowed to pick up the kit tray while there is another object on the work table, the planning system may not be able to find a solution when it comes to put the kit tray on the work table. Therefore, it is necessary to check that the top of *worktable* is clear even before the robot picks up a kit tray from the large box of empty kit trays.
- ☐ `endeff-type-kittray(endeff,kittray)`: *endeff* in the robot's arm must be capable of handling *kittray*.

■ *effects*

- ☐ \neg `robot-empty(robot)`: *robot*'s end effector is no longer empty since it contains *kittray*.
- ☐ \neg `kittray-location-lbwekt(kittray,lbwekt)`: *kittray* is no longer in *lbwekt* since it is in the robot's end effector.
- ☐ `kit-tray-location(kittray,robot)`: *kittray* is in *robot*'s end effector.
- ☐ `robot-holds-kittray(robot,kittray)`: *robot* is holding *kittray*.

2. *put-kittray(robot,kittray,worktable)*: The Robot *robot* puts the Kit-Tray *kittray* on the WorkTable *worktable*.

<i>preconditions</i>	<i>effects</i>
kittray-location-robot(<i>kittray</i> , <i>robot</i>)	\neg kittray-location-robot(<i>kittray</i> , <i>robot</i>)
robot-holds-kittray(<i>robot</i> , <i>kittray</i>)	\neg robot-holds-kittray(<i>robot</i> , <i>kittray</i>)
worktable-empty(<i>worktable</i>)	\neg worktable-empty(<i>worktable</i>)
	kittray-location-wtable(<i>kittray</i> , <i>worktable</i>)
	robot-empty(<i>robot</i>)
	on-wtable-kittray(<i>worktable</i> , <i>kittray</i>)

■ *preconditions*

- kittray-location-robot(*kittray*,*robot*): *kittray* is in *robot*'s end effector.
- robot-holds-kittray(*robot*,*kittray*): *robot* holds *kittray*.
- worktable-empty(*worktable*): There is nothing on *worktable*.

■ *effects*

- \neg kittray-location-robot(*kittray*,*robot*): *kittray* is no longer it *robot*'s end effector since it is placed on *worktable*.
- \neg robot-holds-kittray(*robot*,*kittray*): *robot* is not holding *kittray* anymore.
- \neg worktable-empty(*worktable*): *worktable* is not empty anymore since there is something on top of it.
- kittray-location-wtable(*kittray*,*worktable*): *kittray* is on *worktable*.
- robot-empty(*robot*): *robot* is not holding anything.
- on-wtable-kittray(*worktable*,*kittray*): *worktable* has *kittray* on top of it.

3. *take-kit*(*robot*,*kit*,*worktable*,*endeff*): The *Robot robot* equipped with the *EndEffector endeff* picks up the *Kit kit* from the *WorkTable worktable*.

<i>preconditions</i>	<i>effects</i>
kit-location-wtable(<i>kit</i> , <i>worktable</i>)	\neg kit-location-wtable(<i>kit</i> , <i>worktable</i>)
robot-empty(<i>robot</i>)	\neg robot-empty(<i>robot</i>)
on-wtable-kit(<i>worktable</i> , <i>kit</i>)	\neg on-wtable-kit(<i>worktable</i> , <i>kit</i>)
robot-with-endeff(<i>robot</i> , <i>endeff</i>)	kit-location-robot(<i>kit</i> , <i>robot</i>)
endeff-type-kit(<i>endeff</i> , <i>kit</i>)	robot-holds-kit(<i>robot</i> , <i>kit</i>)
	worktable-empty(<i>worktable</i>)

■ *preconditions*

- kit-location-wtable(*kit*,*worktable*): *kit* is located on *worktable*.
- robot-empty(*robot*): *robot* is not holding any object.
- on-wtable-kit(*worktable*,*kit*): *worktable* has *kit* on top of it.

- $\text{robot-with-endeff}(\text{robot}, \text{endeff})$: *robot* is equipped with *endeff*.
- $\text{endeff-type-kit}(\text{endeff}, \text{kit})$: The type of *endeff* is capable of handling *kit*.

■ *effects*

- $\neg \text{kit-location-wtable}(\text{kit}, \text{worktable})$: *kit* is not on *worktable*.
- $\neg \text{robot-empty}(\text{robot})$: *robot* is holding an object (*kit*).
- $\neg \text{on-wtable-kit}(\text{worktable}, \text{kit})$: *worktable* does not have *kit* on top of it.
- $\text{kit-location-robot}(\text{kit}, \text{robot})$: *kit* is being held by *robot*.
- $\text{robot-holds-kit}(\text{robot}, \text{kit})$: *robot* is holding *kit*.
- $\text{worktable-empty}(\text{worktable})$: *worktable* does not have any object on top of it.

4. $\text{put-kit}(\text{robot}, \text{kit}, \text{lbwk})$: The *Robot robot* puts down the *Kit kit* in the *LargeBoxWithKits lbwk*.

<i>preconditions</i>	<i>effects</i>
$\text{kit-location-robot}(\text{kit}, \text{robot})$	$\neg \text{kit-location-robot}(\text{kit}, \text{robot})$
$\text{robot-holds-kit}(\text{robot}, \text{kit})$	$\neg \text{robot-holds-kit}(\text{robot}, \text{kit})$
$\text{lbwk-not-full}(\text{lbwk})$	$\text{kit-location-lbwk}(\text{kit}, \text{lbwk})$
	$\text{robot-empty}(\text{robot})$

■ *preconditions*

- $\text{kit-location-robot}(\text{kit}, \text{robot})$: *kit* is held by *robot*.
- $\text{robot-holds-kit}(\text{robot}, \text{kit})$: *robot* is holding *kit*.
- $\text{lbwk-not-full}(\text{lbwk})$: *lbwk* should not be full so it can contain *kit*.

■ *effects*

- $\neg \text{kit-location-robot}(\text{kit}, \text{robot})$: *kit* is not being held by *robot*.
- $\neg \text{robot-holds-kit}(\text{robot}, \text{kit})$: *robot* is not holding *kit*.
- $\text{kit-location-lbwk}(\text{kit}, \text{lbwk})$: *kit* has been placed in *lbwk*.
- $\text{robot-empty}(\text{robot})$: *robot* is not holding anything (not holding *kit* anymore).

5. $\text{look-for-part}(\text{robot}, \text{part}, \text{partstray}, \text{kit}, \text{worktable}, \text{endeff})$: A sensor looks for the *Part part* in the *PartTray partstray*.

<i>preconditions</i>	<i>effects</i>
part-not-searched	\neg part-not-searched
robot-empty(<i>robot</i>)	found-part(<i>partstray</i>)
robot-with-endeff(<i>robot</i> , <i>endeff</i>)	
on-wtable-kit(<i>worktable</i> , <i>kit</i>)	
endeff-location-robot(<i>endeff</i> , <i>robot</i>)	
part-location-partstray(<i>part</i> , <i>partstray</i>)	
kit-location-wtable(<i>kit</i> , <i>worktable</i>)	
endeff-type-part(<i>endeff</i> , <i>part</i>)	
partstray-not-empty(<i>partstray</i>)	

■ *preconditions*

- ☐ **part-not-searched**: This flag is set to true in the initial state in the problem file (see section 4.2) and means that a part has not been searched yet.
- ☐ **robot-empty(*robot*)**: *robot* should not be holding anything. We want the operator *look-for-part* to be directly followed by the operator *take-part* to simulate a sensor identifying a part before being picked up by a robot. It is necessary to check that *robot*'s end effector is empty to prepare for the execution of the operator *take-part*.
- ☐ **robot-with-endeff(*robot*,*endeff*)**: *robot* is equipped with *endeff*. Again, since we want the operator following *look-for-part* to be *take-part*, we want to make sure that *robot* is already equipped with *endeff*.
- ☐ **on-wtable-kit(*worktable*,*kit*)**: *worktable* has *kit* on top of it.
- ☐ **endeff-location-robot(*endeff*,*robot*)**: *endeff* is on *robot* so it is ready for the operator *take-part*.
- ☐ **part-location-partstray(*part*,*partstray*)**: *part* is in *partstray*.
- ☐ **kit-location-wtable(*kit*,*worktable*)**: *kit* is on *worktable*.
- ☐ **endeff-type-part(*endeff*,*part*)**: *endeff* can handle *part*.
- ☐ **partstray-not-empty(*partstray*)**: *partstray* contains at least one part.

■ *effects*

- ☐ **\neg part-not-searched**: This flag is set to true so that the operator *look-for-part* can be called again to look for another part in the workstation.
- ☐ **found-part(*partstray*)**: A part from *partstray* has been found.

6. *take-part*(*robot*,*part*,*partstray*,*endeff*,*worktable*,*kit*): The *Robot robot* uses the *EndEffector endeff* to pick up the *Part part* from the *PartTray partstray*.

<i>preconditions</i>	<i>effects</i>
part-location-partstray(<i>part</i> , <i>partstray</i>)	\neg part-location-partstray(<i>part</i> , <i>partstray</i>)
robot-empty(<i>robot</i>)	\neg robot-empty(<i>robot</i>)
endeff-location-robot(<i>endeff</i> , <i>robot</i>)	part-location-robot(<i>part</i> , <i>robot</i>)
robot-with-endeff(<i>robot</i> , <i>endeff</i>)	robot-holds-part(<i>robot</i> , <i>part</i>)
on-wtable-kit(<i>worktable</i> , <i>kit</i>)	decrease quantity-parts-in-partstray(<i>partstray</i>)
kit-location-wtable(<i>kit</i> , <i>worktable</i>)	
endeff-type-part(<i>endeff</i> , <i>part</i>)	
partstray-not-empty(<i>partstray</i>)	
found-part(<i>part</i> , <i>partstray</i>)	

■ *preconditions*

- ☐ part-location-partstray(*part*,*partstray*): *part* to be picked up is in *partstray*.
- ☐ robot-empty(*robot*): *robot* is not holding any object.
- ☐ endeff-location-robot(*endeff*,*robot*): *endeff* is on *robot*.
- ☐ robot-with-endeff(*robot*,*endeff*): *robot* is equipped with *endeff*.
- ☐ on-wtable-kit(*worktable*,*kit*): *worktable* has *kit* on top of it.
- ☐ kit-location-wtable(*kit*,*worktable*): *kit* is on *worktable*. Once a part is picked up by the robot, the next logical action would be to put the part in the kit. For this to happen, the kit needs to be already on the work table so it can hold the part.
- ☐ endeff-type-part(*endeff*,*part*): *endeff* is the type for *part* handling.
- ☐ partstray-not-empty(*partstray*): *partstray* is not empty and contains at least one part.
- ☐ found-part(*part*,*partstray*): *part* has been found in *partstray*. found-part is set to true in the *effects* of the operator *look-for-part*.

■ *effects*

- ☐ \neg part-location-partstray(*part*,*partstray*): *part* is not in *partstray* anymore since it was picked up by *robot*.
- ☐ \neg robot-empty(*robot*): *robot* is now holding *part* and is not empty anymore.
- ☐ part-location-robot(*part*,*robot*): *part* is held by *robot*.
- ☐ robot-holds-part(*robot*,*part*): *robot* is holding *part*.
- ☐ decrease quantity-parts-in-partstray(*partstray*): After picking up a part from *partstray* the number of parts in *partstray* is decreased by one. This is expressed with the *decrease* function.

7. *put-part*(*robot*,*part*,*kit*,*worktable*,*partstray*): The Robot *robot* puts the Part *part* in the Kit *kit*.

<i>preconditions</i>	<i>effects</i>
$\text{part-location-robot}(\text{part}, \text{robot})$	$\neg \text{part-location-robot}(\text{part}, \text{robot})$
$\text{robot-holds-part}(\text{robot}, \text{part})$	$\neg \text{robot-holds-part}(\text{robot}, \text{part})$
$\text{on-wtable-kit}(\text{worktable}, \text{kit})$	$\neg \text{found-part}(\text{part}, \text{partstray})$
$\text{kit-location-wtable}(\text{kit}, \text{worktable})$	$\text{robot-empty}(\text{robot})$
$\text{origin-part}(\text{part}, \text{partstray})$	$\text{part-location-kit}(\text{part}, \text{kit})$
$(< (\text{quantity-parts-in-kit}(\text{kit}, \text{partstray}))$	$(\text{increase } (\text{quantity-kit}(\text{kit}, \text{partstray})))$
$(\text{capacity-kit}(\text{kit}, \text{partstray})))$	part-not-searched

■ *preconditions*

- ☐ $\text{part-location-robot}(\text{part}, \text{robot})$: *part* is held by *robot*.
- ☐ $\text{robot-holds-part}(\text{robot}, \text{part})$: *robot* is holding *part*.
- ☐ $\text{on-wtable-kit}(\text{worktable}, \text{kit})$: *worktable* has *kit* on top of it.
- ☐ $\text{kit-location-wtable}(\text{kit}, \text{worktable})$: *kit* is on *worktable*.
- ☐ $\text{origin-part}(\text{part}, \text{partstray})$: *part* is from *partstray*. This is used to tell the type of *part*.
- ☐ $(< (\text{quantity-kit}(\text{kit}, \text{partstray})) (\text{capacity-kit}(\text{kit}, \text{partstray})))$: The quantity of parts of type *partstray* in *kit* should be lesser than the capacity *kit* can hold for this type of part.

■ *effects*

- ☐ $\neg \text{part-location-robot}(\text{part}, \text{robot})$: *part* is not held by *robot*.
- ☐ $\neg \text{robot-holds-part}(\text{robot}, \text{part})$: *robot* is not holding *part*.
- ☐ $\text{robot-empty}(\text{robot})$: *robot* is not holding any object.
- ☐ $\text{part-location}(\text{part}, \text{kit})$: *part* is located in *kit*.
- ☐ $(\text{increase } (\text{quantity-kit}(\text{kit}, \text{partstray})))$: Once *part* is placed in *kit*, the quantity of *parts* in *kit* is increased by one.
- ☐ part-not-searched : This flag is set to true so another part search (through the operator *look-for-part*) is made after *part* is placed in *kit*.

8. *attach-endeff*(*robot*, *endeff*, *endeffholder*, *endeffchstation*): The *Robot robot* attaches the *EndEffector endeff* which is situated in the *EndEffHolder endeffholder*.

preconditions

$\text{endeff-location-endeffholder}(\text{endeff}, \text{endeffholder})$
 $\text{robot-with-no-endeff}(\text{robot})$
 $\text{endeffholder-holds-endeff}(\text{endeffholder}, \text{endeff})$
 $\text{endeffholder-location}(\text{endeffholder}, \text{endeffchstation})$
 $\text{endeffchstation-contains-endeffholder}(\text{endeffchstation}, \text{endeffholder})$

effects

$\neg \text{endeff-location-endeffholder}(\text{endeff}, \text{endeffholder})$
 $\neg \text{endeffholder-holds-endeff}(\text{endeffholder}, \text{endeff})$
 $\neg \text{robot-with-no-endeff}(\text{robot})$
 $\text{robot-empty}(\text{robot})$
 $\text{endeff-location-robot}(\text{endeff}, \text{robot})$
 $\text{robot-with-endeff}(\text{robot}, \text{endeff})$
 $\text{endeffholder-empty}(\text{endeffholder})$

■ *preconditions*

- $\text{endeff-location-endeffholder}(\text{endeff}, \text{endeffholder})$: *endeff* is located in *endeffholder*.
- $\text{robot-with-no-endeff}(\text{robot})$: *robot* is not equipped with any *endeff*.
- $\text{endeffholder-holds-endeff}(\text{endeffholder}, \text{endeff})$: *endeffholder* is holding *endeff*.
- $\text{endeffholder-location}(\text{endeffholder}, \text{endeffchstation})$: *endeffholder* is in *endeffchstation*.
- $\text{endeffchstation-contains-endeffholder}(\text{endeffchstation}, \text{endeffholder})$: *endeffchstation* contains *endeffholder*.

■ *effects*

- $\neg \text{endeff-location-endeffholder}(\text{endeff}, \text{endeffholder})$: *endeff* is not in *endeffholder* anymore since it has been attached to *robot*.
- $\neg \text{endeffholder-holds-endeff}(\text{endeffholder}, \text{endeff})$: *endeffholder* is not holding *endeff* anymore.
- $\neg \text{robot-with-no-endeff}(\text{robot})$: *robot* is now equipped with *endeff*.
- $\text{robot-empty}(\text{robot})$: *robot* is not holding any object.
- $\text{endeff-location-robot}(\text{endeff}, \text{robot})$: *endeff* is on *robot*.
- $\text{robot-with-endeff}(\text{robot}, \text{endeff})$: *robot* is equipped with *endeff*.
- $\text{endeffholder-empty}(\text{endeffholder})$: *endeffholder* is not holding any *endeff*.

9. **remove-endeff**(*robot*, *endeff*, *endeffholder*, *endeffchstation*): The *Robot* *robot* removes the *EndEffector* *endeff* and puts it in the *EndEffHolder* *endeffholder*.

preconditions

$\text{endeff-location-robot}(\text{endeff}, \text{robot})$
 $\text{robot-with-endeff}(\text{robot}, \text{endeff})$
 $\text{robot-empty}(\text{robot})$
 $\text{endeffholder-location}(\text{endeffholder}, \text{endeffchstation})$
 $\text{endeffchstation-contains-endeffholder}(\text{endeffchstation}, \text{endeffholder})$
 $\text{endeffholder-empty}(\text{endeffholder})$

effects

$\neg \text{endeff-location-robot}(\text{endeff}, \text{robot})$
 $\neg \text{robot-with-endeff}(\text{robot}, \text{endeff})$
 $\neg \text{endeffholder-empty}(\text{endeffholder})$
 $\text{endeff-location-endeffholder}(\text{endeff}, \text{endeffholder})$
 $\text{endeffholder-holds-endeff}(\text{endeffholder}, \text{endeff})$
 $\text{robot-with-no-endeff}(\text{robot})$

■ *preconditions*

- ☐ $\text{endeff-location-robot}(\text{endeff}, \text{robot})$: *endeff* is on *robot*.
- ☐ $\text{robot-with-endeff}(\text{robot}, \text{endeff})$: *robot* is holding *endeff*.
- ☐ $\text{robot-empty}(\text{robot})$: *robot* is not holding anything.
- ☐ $\text{endeffholder-location}(\text{endeffholder}, \text{endeffchstation})$: *endeffholder* is in *endeffchstation*.
- ☐ $\text{endeffchstation-contains-endeffholder}(\text{endeffchstation}, \text{endeffholder})$
- ☐ $\text{endeffholder-empty}(\text{endeffholder})$: *endeffholder* does not contain *endeff*.

■ *effects*

- ☐ $\neg \text{endeff-location-robot}(\text{endeff}, \text{robot})$: *endeff* is not on *robot* anymore.
- ☐ $\neg \text{robot-with-endeff}(\text{robot}, \text{endeff})$: *robot* does not have *endeff* anymore.
- ☐ $\neg \text{endeffholder-empty}(\text{endeffholder})$: *endeffholder* does not contain any *endeff*.
- ☐ $\text{endeff-location-endeffholder}(\text{endeff}, \text{endeffholder})$: *endeff* is situated in *endeffholder*.
- ☐ $\text{endeffholder-holds-endeff}(\text{endeffholder}, \text{endeff})$: *endeffholder* holds *endeff*.
- ☐ $\text{robot-with-no-endeff}(\text{robot})$: *robot* is not equipped with *endeff*.

10. **create-kit**(*kit*, *kittray*, *worktable*): The *KitTray* *kittray* is converted into the *Kit* *kit* once the *KitTray* *kittray* is on the *WorkTable* *worktable*.

<i>preconditions</i>	<i>effects</i>
$\text{on-wtable-kittray}(\text{worktable}, \text{kittray})$	$\neg \text{on-wtable-kittray}(\text{worktable}, \text{kittray})$ $\text{on-wtable-kit}(\text{worktable}, \text{kit})$ $\text{kit-location-wtable}(\text{kit}, \text{worktable})$

■ *preconditions*

- ☐ $\text{on-wtable-kittray}(\text{worktable}, \text{kittray})$: *worktable* has *kittray* on top of it.

■ *effects*

- $\neg \text{on-wtable-kittray}(\text{worktable}, \text{kittray})$: The object *kittray* is destroyed and is thus not on *worktable* anymore.
- $\text{on-wtable-kit}(\text{worktable}, \text{kit})$: *worktable* now has *kit* on top of it.
- $\text{kit-location-wtable}(\text{kit}, \text{worktable})$: *kit* is on *worktable*.

2.5.2 Actions

An action *a* can be obtained by substituting the object variable symbols that appear anywhere in the operator with constant variable symbols. For instance, the operator *take-part*(*robot*, *part*, *partstray*, *endeff*) in the kitting domain can be translated into the action *take-part*(*robot_1*, *part_1*, *partstray_1*, *endeff_2*) where *robot_1*, *part_1*, *partstray_1*, and *endeff_2* are constant variable symbols in the classes *Robot*, *Part*, *PartsTray*, and *EndEffector*, respectively.

3 Ontology

4 Planning Language

The Planning Domain Definition Language (PDDL) [6] is an attempt by the domain independent planning community to formulate a standard language for planning. A community of planning researchers has been producing planning systems that comply with this formalism since the first International Planning Competition held in 1998. This competition series continues today, with the seventh competition being held in 2011. PDDL is constantly adding extensions to the base language in order to represent more expressive problem domains. Our work is based on PDDL Version 3.

By placing our knowledge in a PDDL representation, we enable the use of an entire family of open source planning systems. Each PDDL file-set consists of two files that specify the domain and the problem.

4.1 The PDDL Domain File

The PDDL domain file is composed of four sections that include requirements, types, predicates and functions, and actions. An excerpt of the PDDL domain file is depicted in Figure 2.

- line 1: The keyword **domain** signals a planner that this file contains information on the domain. **kitting-domain** is the name given to the domain.
- line 2: The **:requirements** field specifies which section the domain relies on. The planning system can examine this statement to determine if it is capable of solving problems in this domain. A keyword (symbol starting with a colon) used in a **:requirements** field is called a requirement flag; the domain is said to declare a requirement for that flag. The requirement flags present in the kitting domain are:
 - **:strips**: The most basic subset of PDDL, consisting of STRIPS only.
 - **:typing**: PDDL has a special syntax for declaring parameter and object types. **:typing** allows types names in declaration of variables.
 - **:fluents**: A domain's set of requirements allow a planner to quickly tell if it is likely to be able to handle the domain. For example, this version of the kitting world requires fluents numeric,

```

1. (define (domain kitting-domain)
2.   (:requirements :strips :typing :fluents)
3.   (:types
4.     EndEffector
5.     EndEffectorHolder
6.     Kit
7.     KitTray
8.     LargeBoxWithEmptyKitTrays
9.     LargeBoxWithKits
10.    Part
11.    PartsTray
12.    EndEffectorChangingStation
13.    Robot
14.    WorkTable
15.  )
16.
17.  (:predicates
18.    (endeffector-location-robot ?endeffector - EndEffector ?robot - Robot)
19.    (on-worktable-kit ?worktable - WorkTable ?kit - Kit)
20.  )
21.
22.  (:functions
23.    (quantity-partstray ?partstray - PartsTray)
24.    (quantity-kit ?kit - Kit ?partstray - PartsTray)
25.    (capacity-kit ?kit - Kit ?partstray - PartsTray)
26.  )
27.
28.  (:action take-kittray
29.    :parameters(
30.      ?robot - Robot
31.      ?kittray - KitTray
32.      ?largeboxwithemptykittrays - LargeBoxWithEmptyKitTrays
33.      ?endeffector - EndEffector
34.      ?worktable - WorkTable)
35.    :precondition(and
36.      (robot-empty ?robot)
37.      (lbwekt-not-empty ?largeboxwithemptykittrays)
38.      (robot-with-endeffector ?robot ?endeffector)
39.      (kittray-location-lbwekt ?kittray ?largeboxwithemptykittrays)
40.      (endeffector-location-robot ?endeffector ?robot)
41.      (worktable-empty ?worktable)
42.      (endeffector-type-kittray ?endeffector ?kittray))
43.    :effect(and
44.      (robot-holds-kittray ?robot ?kittray)
45.      (kittray-location-robot ?kittray ?robot)
46.      (not (robot-empty ?robot))
47.      (not (kittray-location-lbwekt ?kittray ?largeboxwithemptykittrays)))
48.  )
49. )
50.

```

Figure 2: Excerpt of the PDDL domain file for kitting.

so a straight STRIPS-representation planner would not be able to handle it. A fluent is a term (`:functions`) with time-varying

value (i.e., a value that can change as a result of performing an action).

- line 3–15: Type names have to be declared before they are used (before `:predicates` and `:functions`). This is done with the declaration `(:types name1 ... namen)`.
- line 17–20: The `:predicates` part of a domain definition specify only what are the predicate names used in the domain, and their number of arguments (and argument types, if the domain uses `:typing`). The “meaning” of a predicate, in the sense of for what combinations of arguments it can be true and its relationship to other predicates, is determined by the effects that actions in the domain can have on the predicate, and by what instances of the predicate are listed as true in the initial state of the problem definition.

It is common to make a distinction between static and dynamic predicates. A *static* predicate is not changed by any action. Thus in a problem, the true and false instances of a *static* predicate will always be precisely those listed in the initial state specification of the problem definition. Note that there is no syntactic difference between *static* and *dynamic* predicates in PDDL, they look exactly the same in the `:predicates` declaration part of the domain.

A predicate is build using the structure `(predicate_name ?X - type_of_X)`. A list of parameters of the same type in a predicate can be abbreviated to `(predicate_name ?X ?Y ?Z - type_of_XYZ)`. Note that the hyphen between parameter and type name is surrounded by whitespace.

- line 22–26: A fluent is similar to a state variable/predicate except that its value is a number instead of true or false. The initial value of a function is set in the initial state of the problem file and changes when an action is executed. The declaration of functions is similar to predicates.
- line 28–48: The domain definition contains operators (called *actions* in PDDL). An action statement specifies a way that a planner affects the state of the world. The statement includes parameters, preconditions, and effects. All parts of an action definition except the name are, according to the PDDL specification, optional (although, of course, an action without effects is pretty useless). However, for an action that has no preconditions some planners may require an “empty” precon-

dition, on the form `:precondition ()` or `:precondition (and)`, and some planners may also require an empty `:parameter` list for actions without parameters).

- line 29–34: The `:parameters` section declare all the parameters used by predicates and functions in `preconditions` and `effects`.
- line 35–42: The `:preconditions` section is a conjunction of predicates and functions that need to be true in the world in order for the action to be invoked.
- line 43–47: The `:effects` equation dictates the changes in the world that will occur due to the execution of the action.

4.2 PDDL Problem File

The second file of the PDDL file-set is a problem file. The problem file specifies information about the specific instance of the given problem. This file contains the initial conditions and definition of the world (in the `init` section) and the final state that the world must be brought to (in the `goal` section). Using an example of kit to build, this section only describes the initial and goal states explicitly. The operators detailed in Section 2.5 are used by a planner to generate the other states as needed.

In this example, the *Robot* has to build a kit that contains two *Parts* of type A, two *Part* of type B and one *Part* of type C. The kitting process is completed once the *Kit* is placed in the *LargeBoxWithKits*. The PDDL problem file for the kitting domain is presented below.

```

1. (define (problem kitting-problem)
2.   (:domain kitting-domain)
3.   (:objects
4.     robot_1 - Robot
5.     changing_station_1 - EndEffectorChangingStation
6.     kit_tray_1 - KitTray
7.     kit_a2b2c1 - Kit
8.     empty_kit_tray_supply - LargeBoxWithEmptyKitTrays
9.     finished_kit_receiver - LargeBoxWithKits
10.    work_table_1 - WorkTable
11.    part_a_tray part_b_tray part_c_tray - PartsTray
12.    part_a_1 part_a_2 part_a_3 part_a_4 - Part
13.    part_b_1 part_b_2 part_b_3 part_b_4 - Part
14.    part_c_1 part_c_2 part_c_3 part_c_4 - Part
15.    part_gripper tray_gripper - EndEffector
16.    part_gripper_holder tray_gripper_holder - EndEffectorHolder
17.  )
18.)
19. (:init
20.  (robot-with-no-endeffector robot_1)
21.  (part-not-searched)
22.  (lbwekt-not-empty empty_kit_tray_supply)
23.  (lbwk-not-full finished_kit_receiver)
24.  (partstray-not-empty part_a_tray)
25.  (partstray-not-empty part_b_tray)
26.  (partstray-not-empty part_c_tray)
27.  (endeffector-location-endeffectorholder part_gripper part_gripper_holder)
28.  (endeffector-location-endeffectorholder tray_gripper tray_gripper_holder)
29.  (endeffectorholder-holds-endeffector part_gripper_holder part_gripper)
30.  (endeffectorholder-holds-endeffector tray_gripper_holder tray_gripper)
31.  (endeffectorholder-location tray_gripper_holder changing_station_1)
32.  (endeffectorholder-location part_gripper_holder changing_station_1)
33.  (endeffectorchangingstation-contains-endeffectorholder changing_station_1 tray_gripper_holder)
34.  (endeffectorchangingstation-contains-endeffectorholder changing_station_1 part_gripper_holder)
35.  (worktable-empty work_table_1)
36.  (kittray-location-lbwekt kit_tray_1 empty_kit_tray_supply)
37.
38.  (part-location-partstray part_a_1 part_a_tray)
39.  (part-location-partstray part_a_2 part_a_tray)
40.  (part-location-partstray part_a_3 part_a_tray)
41.  (part-location-partstray part_a_4 part_a_tray)
42.  (part-location-partstray part_b_1 part_b_tray)
43.  (part-location-partstray part_b_2 part_b_tray)
44.  (part-location-partstray part_b_3 part_b_tray)
45.  (part-location-partstray part_b_4 part_b_tray)
46.  (part-location-partstray part_c_1 part_c_tray)
47.  (part-location-partstray part_c_2 part_c_tray)
48.  (part-location-partstray part_c_3 part_c_tray)
49.  (part-location-partstray part_c_4 part_c_tray)
50.
51.  (endeffector-type-part part_gripper part_a_1)
52.  (endeffector-type-part part_gripper part_a_2)
53.  (endeffector-type-part part_gripper part_a_3)
54.  (endeffector-type-part part_gripper part_a_4)

```

```

55. (endeffector-type-part part_gripper part_b_1)
56. (endeffector-type-part part_gripper part_b_2)
57. (endeffector-type-part part_gripper part_b_3)
58. (endeffector-type-part part_gripper part_b_4)
59. (endeffector-type-part part_gripper part_c_1)
60. (endeffector-type-part part_gripper part_c_2)
61. (endeffector-type-part part_gripper part_c_3)
62. (endeffector-type-part part_gripper part_c_4)
63. (endeffector-type-kittray tray_gripper kit_tray_1)
64. (endeffector-type-kit tray_gripper kit_a2b2c1)
65.
66. (= (capacity-kit kit_a2b2c1 part_a_tray) 2)
67. (= (capacity-kit kit_a2b2c1 part_b_tray) 2)
68. (= (capacity-kit kit_a2b2c1 part_c_tray) 1)
69. (= (quantity-kit kit_a2b2c1 part_a_tray) 0)
70. (= (quantity-kit kit_a2b2c1 part_b_tray) 0)
71. (= (quantity-kit kit_a2b2c1 part_c_tray) 0)
72. (= (quantity-partstray part_a_tray) 4)
73. (= (quantity-partstray part_b_tray) 4)
74. (= (quantity-partstray part_c_tray) 4)
75.
76. (origin-part part_a_1 part_a_tray)
77. (origin-part part_a_2 part_a_tray)
78. (origin-part part_a_3 part_a_tray)
79. (origin-part part_a_4 part_a_tray)
80. (origin-part part_b_1 part_b_tray)
81. (origin-part part_b_2 part_b_tray)
82. (origin-part part_b_3 part_b_tray)
83. (origin-part part_b_4 part_b_tray)
84. (origin-part part_c_1 part_c_tray)
85. (origin-part part_c_2 part_c_tray)
86. (origin-part part_c_3 part_c_tray)
87. (origin-part part_c_4 part_c_tray)
88. )
89.
90. (:goal
91.   (and
92.     (= (quantity-kit kit_a2b2c1 part_a_tray)
93.       (capacity-kit kit_a2b2c1 part_a_tray))
94.     (= (quantity-kit kit_a2b2c1 part_b_tray)
95.       (capacity-kit kit_a2b2c1 part_b_tray))
96.     (= (quantity-kit kit_a2b2c1 part_c_tray)
97.       (capacity-kit kit_a2b2c1 part_c_tray))
98.     (kit-location-lbwk kit_a2b2c1 finished_kit_receiver)
99.   )
100. )

```

- line 1: Signal a planner that the file contains all the element part of a problem. `kitting-problem` is the name given to this problem.
- line 2: `:domain` refers to the domain that the current problem is associated to. In this case, the problem refers to the domain `kitting-domain`. Note that `kitting-domain` is the name given to the kitting domain as presented in section 4.1.

- line 3–17: `:objects` declare objects present in the problem instance. The syntax for `:objects` is `object1 - Type ... objectn - Type`.

4.2.1 Initial State

The initial state S_0 (Figure 3) defines the environment in its initial condition. The initial state of the kitting problem in PDDL format is described below.

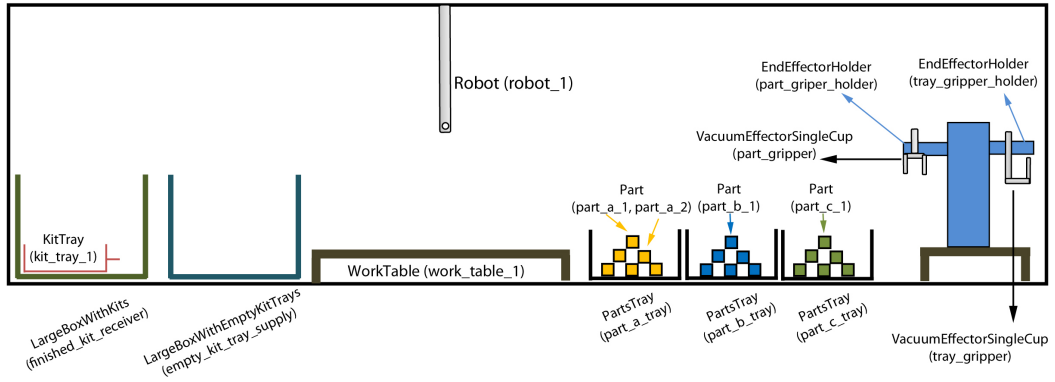


Figure 3: Initial state S_0 .

- line 19: `:init` signals a planner that the predicates and functions in this section are true in the initial state.
- line 20–87: Predicates true in the initial state of the environment. Since PDDL uses a close world assumption, predicates that are not present in the initial state are automatically set to false. This section also set the initial values for functions. Some relevant sections are presented:
 - line 21: The predicate `part-not-searched` is set to true so that the operator *look-for-part* can be activated during a plan search.
 - line 65–67: Functions describing the quantity of parts of a type that `kit_a2b2c1` can contain. In this example, `kit_a2b2c1` can have 2 parts of type A (`part_a_tray`), 2 parts of type B (`part_b_tray`), and 1 part of type C (`part_c_tray`).
 - line 68–70: Functions that represent the quantity of parts of a specific type that are already in `kit_a2b2c1`. `kit_a2b2c1` has no parts of type A, B, and C.

- line 71–73: Functions that describe the quantity of parts available in their respective parts tray. This also can be read as: *In the workstation, there are 4 parts of type A available, 4 parts of type B available, and 4 parts of type C available.*
- line 75–86: Predicates that describe the type of each specific part in the workstation.

4.2.2 Goal State

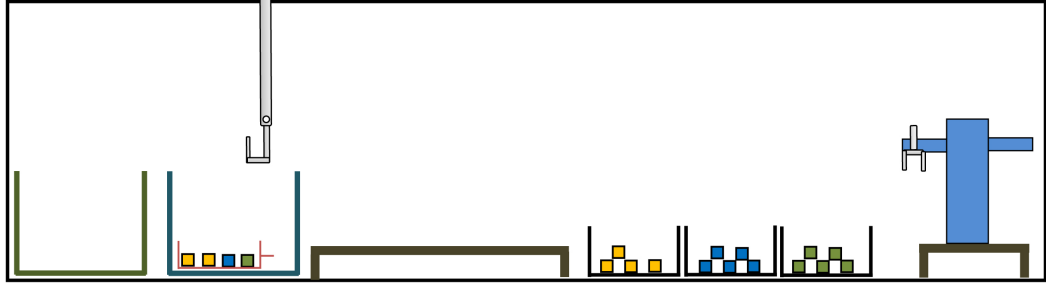


Figure 4: Goal state S_G .

Figure 4 depicts the goal state S_G for the kitting workstation. The expression of the goal state in PDDL is described below.

- line 89: `:goal` is a keyword used to signal a planner about the goal state to reach. All the predicates and functions in the goal state must be true.
- line 91–93: The quantity of parts of a specific type in `kit_a2b2c1` should match the capacity of parts of a specific type for `kit_a2b2c1`. The quantity of parts in `kit_a2b2c1` is increased in the operator `put-part`. The initial quantity of parts in `kit_a2b2c1` and its capacity are set in the initial state. Note that we are not specifying which instance of `Part` should go in `kit_a2b2c1` but rather the number of `Parts` of a specific type that `kit_a2b2c1` must have.
- line 94: `kit_a2b2c1` should be placed in the large box with kits `finished_kit_receiver`.

4.3 Plan

This section shows an example of a plan generated by a planner (see section 6) for the PDDL domain and problem files discussed previously in this document.

Figure 5 displays the different states and actions used by the planner to generate a plan starting from the initial state S_0 to the goal state S_G . The actions $A_1 \dots A_{17}$ are described below.

- $A1: (\text{attach-endeffectector} \quad robot_1 \quad tray_gripper \quad tray_gripper_holder \quad changing_station_1)$
- $A2: (\text{take-kittray} \quad robot_1 \quad kit_tray_1 \quad empty_kit_tray_supply \quad tray_gripper \quad work_table_1)$
- $A3: (\text{put-kittray} \quad robot_1 \quad kit_tray_1 \quad work_table_1)$
- $A4: (\text{create-kit} \quad kit_a2b2c1 \quad kit_tray_1 \quad work_table_1)$
- $A5: (\text{remove-endeffectector} \quad robot_1 \quad tray_gripper \quad tray_gripper_holder \quad changing_station_1)$
- $A6: (\text{attach-endeffectector} \quad robot_1 \quad part_gripper \quad part_gripper_holder \quad changing_station_1)$
- $A7: (\text{look-for-part} \quad robot_1 \quad part_c_1 \quad part_c_tray \quad kit_a2b2c1 \quad work_table_1 \quad part_gripper)$
- $A8: (\text{take-part} \quad robot_1 \quad part_c_1 \quad part_c_tray \quad part_gripper \quad work_table_1 \quad kit_a2b2c1)$
- $A9: (\text{put-part} \quad robot_1 \quad part_c_1 \quad kit_a2b2c1 \quad work_table_1 \quad part_c_tray)$
- $A10: (\text{look-for-part} \quad robot_1 \quad part_b_2 \quad part_b_tray \quad kit_a2b2c1 \quad work_table_1 \quad part_gripper)$
- $A11: (\text{take-part} \quad robot_1 \quad part_b_2 \quad part_b_tray \quad part_gripper \quad work_table_1 \quad kit_a2b2c1)$
- $A12: (\text{put-part} \quad robot_1 \quad part_b_2 \quad kit_a2b2c1 \quad work_table_1 \quad part_b_tray)$
- $A13: (\text{look-for-part} \quad robot_1 \quad part_b_1 \quad part_b_tray \quad kit_a2b2c1 \quad work_table_1 \quad part_gripper)$
- $A14: (\text{take-part} \quad robot_1 \quad part_b_1 \quad part_b_tray \quad part_gripper \quad work_table_1 \quad kit_a2b2c1)$

- A15:(*put-part* robot_1 part_b_1 kit_a2b2c1 work_table_1 part_b_tray)
- A16:(*look-for-part* robot_1 part_a_2 part_a_tray kit_a2b2c1 work_table_1 part_gripper)
- A17:(*take-part* robot_1 part_a_2 part_a_tray part_gripper work_table_1 kit_a2b2c1)
- A18:(*put-part* robot_1 part_a_2 kit_a2b2c1 work_table_1 part_a_tray)
- A19:(*look-for-part* robot_1 part_a_1 part_a_tray kit_a2b2c1 work_table_1 part_gripper)
- A20:(*take-part* robot_1 part_a_1 part_a_tray part_gripper work_table_1 kit_a2b2c1)
- A21:(*put-part* robot_1 part_a_1 kit_a2b2c1 work_table_1 part_a_tray)
- A22:(*remove-endeffector* robot_1 part_gripper part_gripper_holder changing_station_1)
- A23:(*attach-endeffector* robot_1 tray_gripper tray_gripper_holder changing_station_1)
- A24:(*take-kit* robot_1 kit_a2b2c1 work_table_1 tray_gripper)
- A25:(*put-kit* robot_1 kit_a2b2c1 finished_kit_receiver)

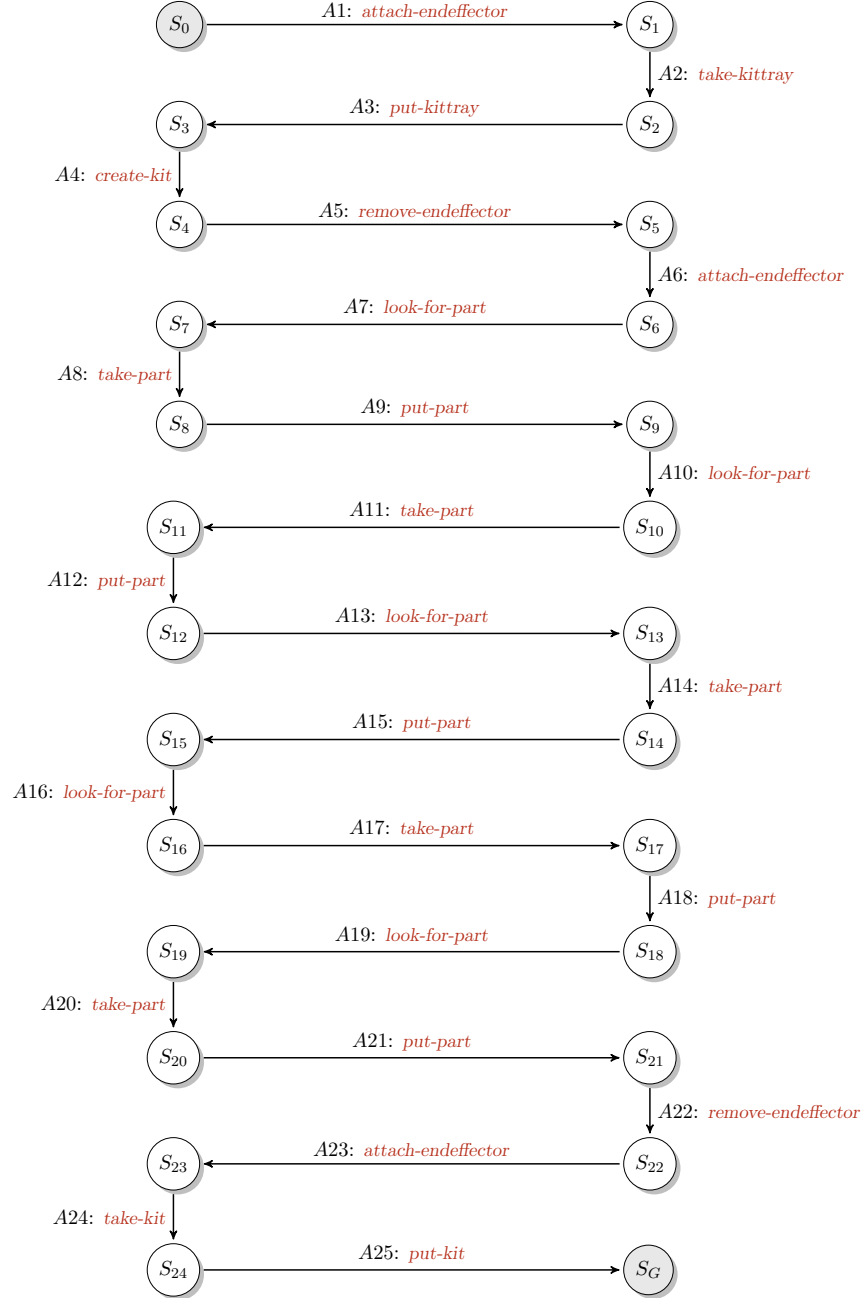


Figure 5: Example of plan generated with the kitting domain and problem files.

5 Robot Language

6 Planner

This section describes the steps to install and run a planner on the PDDL domain and problem files in order to generate a plan. The planner uses a forward-chaining partial-order planning [4].

6.1 Requirements

The planner requires:

- `cmake`
- The CBC mixed integer programming solver (<https://projects.coin-or.org/Cbc/>)
- `perl`, `bison` and `flex` to build the parser

These are packaged with most Linux distributions - on Ubuntu/Debian, the following should suffice:

```
sudo apt-get install cmake coinor-libcbc-dev coinor-libclp-dev \
coinor-libcoinutils-dev bison flex
```

6.2 Download and Install CBC

The CBC source code can be retrieved in two ways:

1. Local copy: `ipmas/planner/coin-Cbc.tar.gz`
2. SVN: `svn co https://projects.coin-or.org/svn/Cbc/stable/2.8 coin-Cbc`

If CBC is retrieved using option 1, unzip `coin-Cbc.tar.gz` to generate the `coin-Cbc` directory.

Perform the following steps.

- `cd coin-Cbc`

- `./configure -C`: Runs a configure script that generates the make file.
- `make`: Builds the Cbc library and executable program.
- `make test`: Builds and runs the Cbc unit test program.
- `make install`: Installs libraries, executables and header files in directories `coin-Cbc/lib`, `coin-Cbc/bin` and `coin-Cbc/include`.

6.3 Compile the Planner

The planner can be found at `ipmas/planner/popf2-11jun2011.tar.bz2`. Unzip `popf2-11jun2011.tar.bz2` to get the `tempo-sat-popf2` directory, then:

- `cd tempo-sat-popf2`
- `./build`

New files and directories are created in the `tempo-sat-popf2/compile/` directory. However, the executable is not generated at this point and errors should be displayed. To fix this, open `tempo-sat-popf2/compile/CMakeCache.txt` in a text editor and edit the following lines. Note that `<path>` is the absolute path that leads to the `coin-Cbc` directory.

- line 24: `CBC_INCLUDES:PATH=<path>/include`
- line 33: `CGL_INCLUDES:PATH=<path>/include`
- line 36: `CGL_LIBRARIES:FILEPATH=/usr/lib/libCgl.so.0`
- line 39: `CLP_INCLUDES:PATH=<path>/include`
- line 186: `COINUTILS_INCLUDES:PATH=<path>/include`
- line 230: `OSICLP_LIBRARIES:FILEPATH=/usr/lib/lib0siClp.so.0`
- line 233: `OSI_INCLUDES:PATH=<path>/include`
- line 236: `OSI_LIBRARIES:FILEPATH=/usr/lib/lib0si.so.0`

Recompile the planner:
`./build`

6.4 Run the Planner

To run the planner, the path to the PDDL domain and problem files should be identified. The format of the PDDL files must be `.pddl`. The following command run the planner on the PDDL files.

■ `./plan <domain> <problem> <solution>`

Where `<domain>` and `<problem>` are the paths that point to the PDDL domain and problem files, respectively.
`<solution>` is the output file that will contain the plan.

7 Tools

7.1 The Generator

The Generator tool is a graphical user interface developed in Java, allowing the user to store data from OWL files into a MySQL database. This tool also permits the user to query the database using the C++ function calls. The tool Generator is composed of the following functionalities:

1. Convert OWL documents into SQL syntaxes (OWL to SQL).
2. Translate SQL syntaxes to OWL language in order to modify an OWL document (SQL to OWL).
3. Convert the OWL language into C++ classes (OWL to C++).

To date, only steps 1. and 3. have been implemented and will be covered in this document.

7.1.1 Prerequisites

The description of the Generator tool is given for a Ubuntu Linux system. To run and use the Generator tool, different applications must be installed on the system.

Java Runtime Environment The Generator tool comes as a jar file. As such, the Java Runtime Environment should be installed on your system. This application can be found at www.oracle.com.

MySQL Server and Client The MySQL server and client should be installed and running on your system.

- *sudo apt-get update* (Update the package management tools)
- *sudo apt-get dist-upgrade* (Install the latest software)
- *sudo apt-get install mysql-server mysql-client* (Install the MySQL server and client packages). You will be asked to enter a password.

When done, you have a MySQL database ready to run. The following command will allow you to run MySQL.

- `mysql -u root -p`

- Enter the same password you used when you installed MySQL.

Finally, we need the plugin `libmysqlcppconn-dev` which allows C++ to connect to MySQL databases. It can be installed as follows:

- `sudo apt-get install libmysqlcppconn-dev`

7.1.2 How to Run the Generator Tool

The Generator tool can be launched using either one of these two following methods:

1. `java -jar Generator.jar`
2. Right-click on `Generator.jar` and select the option “Open With Open-JDK Java 6 Runtime”. Note that this message will be different for future releases of the Java Runtime Environment.

7.1.3 Functionalities

As mentioned in the Introduction, we are covering only steps 1. and 3. in the rest of this document, i.e., *OWL to SQL* and *OWL to C++*, respectively.

OWL to SQL To convert OWL classes and instances to SQL, the `Owl to SQL` tab should be selected (see Figure 6). The different fields are:

Generate SQL Files

- **Ontology Path:** This field requires the file `kittingInstances.owl`. Before doing so, you need to modify one line in this file. Open it with a text editor and find the line `Import(<file:kittingClasses.owl>)`.

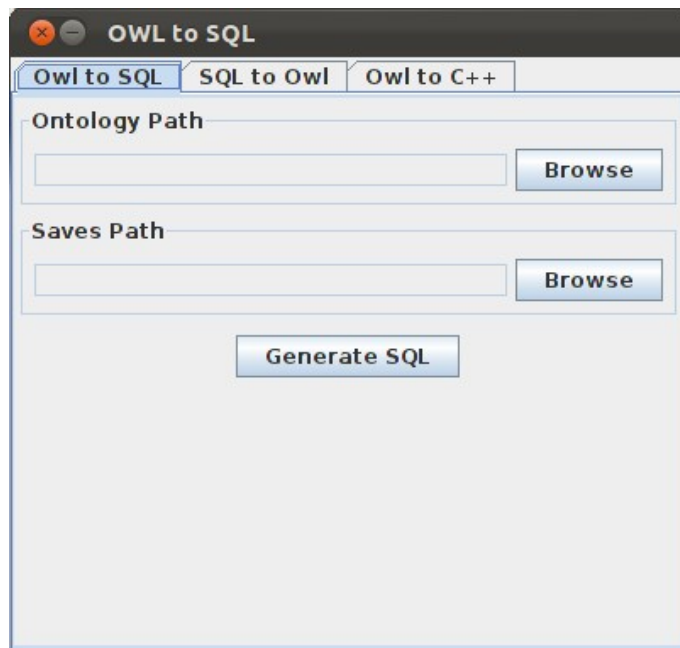


Figure 6: Owl to SQL tab.

Modify this line by giving the absolute path to the file `kittingClasses.owl`. You should have something that looks like `Import(<file:/home/username/NIST/ipmas/Generator/kittingClasses.owl>)`. When this is done, save the file, and browse to `kittingInstances.owl` using the “Browse” button.

- Browse to the directory where you want to save the SQL files.

Once the two previous steps are done, click on “Generate SQL”. You should receive a message confirming the generation of the SQL files: `kittingInstances.owlCreateTable.sql` and `kittingInstances.owlInsertInto.sql`. The former is used to create tables, the latter is used to populate these tables;

SQL Tables and Insertions The next step is to create a database and to populate it.

- Connect to mysql using `mysql -u root -p`, then enter your password. You should be in the mysql shell if this succeeded (`mysql>`).

- Delete a previous database (if you already used this tool and you want to replace the existing database with this new one) : `mysql> DROP DATABASE OWL;` (*OWL* is the name of the old database).

- Create a database:

- `mysql> CREATE DATABASE OWL;` Here, *OWL* is the name of the database (you can use a name of your choice).
- Before performing the following commands, we need to tell MySQL which database we are planning to work with (*OWL* in our case). This is done using:

`mysql> USE OWL`

- Populate the database with tables using `kittingInstances.owlCreateTable.sql`.

- `mysql> source <path>/kittingInstances.owlCreateTable.sql;`

- Populate the tables with data using `kittingInstances.owlInsertInto.sql`:

- `mysql> source <path>/kittingInstances.owlInsertInto.sql;`

`<path>` designs the absolute path to the appropriate file.

OWL to C++ The “Owl to C++” tab (see Figure 7) is used to generate C++ classes and scripts allowing the connection between C++ and MySQL. The different fields are explained below:

- **Ontology Path:** This is the path to the ontology (`kittingClasses.owl` in our example).
- **Saves Path:** Directory where the C++ files and scripts will be generated.
- **Url:** This is the url of the database. It’s usually the IP address of the machine hosting the database (127.0.0.1 if it is local).
- **User name:** User name used to connect to the MySQL database.
- **Password:** Password associated to the user name to connect to the MySQL database.

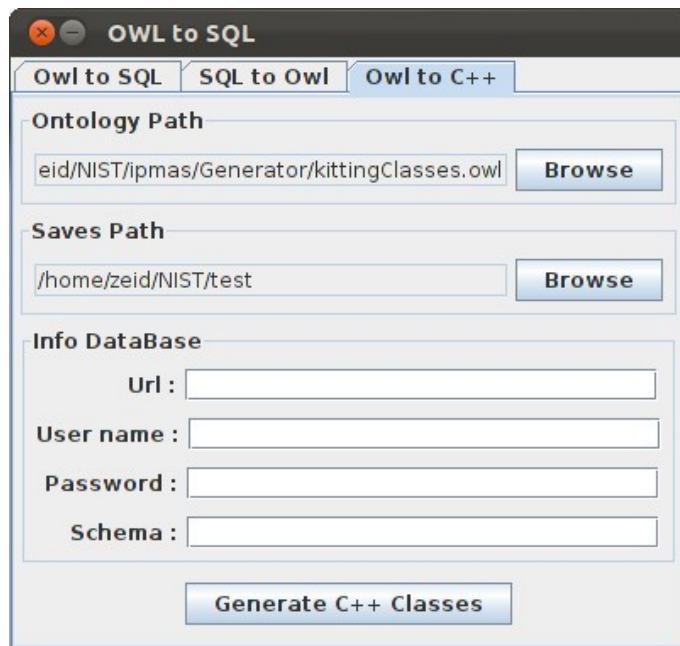


Figure 7: Owl to C++ tab.

■ **Schema:** This is the name of the database (*OWL* in our example).

When all the fields are completed, click the “Generate C++ Classes” button to generate C++ and script files.

7.2 XML to OWL

This section describes tools that are used to generate OWL files from XML files.

7.2.1 owlPrinter

owlPrinter reads an XML kitting data file corresponding to an XML schema for kitting (kitting.xsd) and writes an OWL instance file corresponding to the OWL class file kittingClasses.owl. The kitting.xsd file contains the same conceptual model as the kittingClasses.owl file, but in a different language.

The owlPrinter is useful because there is no OWL tool that will help generate

an OWL instance file and check the file adequately against an OWL class file. That is because OWL uses an open world model in which anything not explicitly or implicitly illegal is allowed. Hence many things that are errors to the writer of the instance file are not OWL errors. For example, if the name of an instance is misspelled, OWL will assume that there is a new instance that has not been explicitly declared as such, which is OK in OWL. If a reference to an instance name is misspelled in an XML data file corresponding to the `kitting.xsd` schema, that will be caught automatically by the `owlPrinter` (and other readily available XML tools). Several other types of error will not be caught by OWL tools but will not be made or will be detected if the OWL printer is used.

Another OWL problem that disappears in XML is that in OWL, there is no distinction between an instance file and a class file. An instance file can modify classes, intentionally, or accidentally. In XML there is no way a data file can modify a model.

To use the `owlPrinter`, use a text editor such as `emacs` or an XML tool such as `XMLSpy` to write an XML data file corresponding to the `kitting.xsd` schema and then run it through the `owlPrinter` with a command of the form:

```
bin/owlPrinter [XML file in] [OWL file out]
```

For example, the command

```
bin/owlPrinter data/kittingInstances.xml junk
```

will print the file `junk`, which will be identical to the `kittingInstances.owl` file in the `owl` directory (except for a couple comments).

7.2.2 `kittingParser`

The `kittingParser` may be used to check an XML data file against the `kitting.xsd` schema. The schema is hard-coded into the `kittingParser`. If there is any error in the XML data file, the `kittingParser` prints a message and quits. If there is no error, the input file is echoed by printing an output file whose name is the same as that of the input file with "echo" appended. To run the `kittingParser`, give a command of the form:

```
bin/kittingParser [XML kitting data file in]
```

For example, the command

```
bin/kittingParser data/kittingInstances.xml
```

will read `kittingInstances.xml` and write `kittingInstances.xmllecho`. The two files will be identical except for comments. If the format of an input file differs from the format used for printing the output file, the two files will differ, but only in format.

The `owlPrinter` makes the same checks as the `kittingParser`, so there is no need to use the `kittingParser`.

All of the source code for the `kittingParser` was generated automatically by the GenXMiller generator.

7.2.3 compactOwl and compareOwl

There is no need to read this section unless you are interested in how the *owlPrinter* was debugged.

The *compactOwl* and *compareOwl* utilities are used for checking that two different OWL instance files have the same statements. They have been used as follows to debug the *owlPrinter* (and the `kitting.xsd` file and the `kittingInstances.xml` file and the `kittingInstances.owl` file).

1. Write `kitting.xsd` to model the same information as `kittingClasses.owl`.
2. Write `kittingInstances.xml` to correspond to `kitting.xsd` and contain the same information as `kittingInstances.owl`.
3. Build the *owlPrinter*.
4. Run `kittingInstances.xml` through the *owlPrinter* to produce `kittingInstances.owl`.
5. Run `kittingInstances.owl` through *compactOwl* to produce one version of `kittingInstancesCompact.owl`.
6. Run `kittingInstances.owl` through *compactOwl* to produce a second version of `kittingInstancesCompact.owl`.

7. Run the two versions of `kittingInstancesCompact.owl` through *compareOwl*. If *compareOwl* reports that the two files have the same statements, that means that steps 1, 2, and 3 have been done correctly, so debugging is finished. If *compareOwl* reports a pair of statements that differ, figure out why, go back to step 1, 2, or 3 (or edit `kittingInstances.owl`), fix the problem, and repeat the subsequent steps.

These utilities assume that the format of the input files is the same as either the format used by the *owlPrinter* or the format followed by the `kittingInstances.owl` file. If the format used by an input file is different from both of those, the utilities may fail.

The *compactOwl* utility compacts an OWL instance file by:

1. removing all occurrences of one or two blank lines. The blank lines must not contain spaces or tabs.
2. removing comments. The comments must have “/” as the first two characters on the line.
3. combining each OWL statement written on two or more lines so it is all on one line. The first non-space character on the second line must be a colon (:) or a double quote (”).
4. rewriting numbers with decimal points so there are exactly six decimal places. Such numbers must have at least one digit on each side of the decimal point in the input file.
5. putting the **DifferentIndividuals** inside **DifferentIndividuals** statements into alphabetical order.

To run *compactOwl* use a command of the form:

```
bin/compactOwl < [owl file in] > [owl file out]
```

where `[owl file in]` and `[owl file out]` are replaced by file names.

The *compareOwl* utility compares two files that are expected to have the same lines, but in a different order, such as an automatically generated OWL file and a hand-written OWL file. It reads the two files and saves the lines of each one in two sets in alphabetical order (`set::insert` puts strings in

alphabetical order by default). Then it compares the lines of the two sets in order. If it finds two lines that do not match, it prints the line from the first file followed by the line from the second file. If all lines match, that is reported.

To run *compareOwl*, use a command of the form:

```
bin/compareOwl [first owl file in] [second owl file in]
```

where [first owl file in] and [second owl file in] are replaced by the names of compacted OWL files.

References

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List of Figures

1	Knowledge Driven Design extensions – In this figure, green shaded boxes with curved bottoms represent hand generated files while light blue shaded boxes with curved bottoms represent automatically created boxes. Rectangular boxes represent processes and libraries. . .	1
2	Excerpt of the PDDL domain file for kitting.	24
3	Initial state S_0	29
4	Goal state S_G	30
5	Example of plan generated with the kitting domain and problem files.	33
6	Owl to SQL tab.	40
7	Owl to C++ tab.	42