# NATIONAL INSTITUTE OF STANDARDS AND TECHNOLOGY

Intelligent Systems Division

## Knowledge Driven Planning and Modeling for Part Handling

Planning for a Kitting Workstation

Stephen Balakirsky Zeid Kootbally Thomas Kramer Anthony Pietromartire Craig Schlenoff stephen.balakirsky@nist.gov zeid.kootbally@nist.gov thomas.kramer@nist.gov pietromartire.anthony@nist.gov craig.schlenoff@nist.gov CONTENTS

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#### 1 The Kitting Domain

The foundation for the knowledge representation is domain specific information that is produced by an expert in the particular field of study. This includes information on items ranging from what actions and attributes are relevant, to what the necessary conditions are for an action to occur and what the likely results of the action are. We have chosen to encode this basic information in a formalism know as a state variable representation (SVR) [4]. This information will then flow up the abstraction and be transformed into the ontology, planning language, and robot language.

Before building a SVR, the domain for kitting needs to be specified. The domain for kitting contains some fixed equipment: a robot, a work table, end effectors, end effector holders, and an end effector changing station. Items that enter the workstation include kit trays, boxes in which to put kits, boxes that contain empty kit trays, and part supplies. Items that leave the workstation may be boxes with finished kits inside, empty part trays, empty boxes. An external agent is responsible of moving the items that leave the workstation. We assume that the workstation has only one work table, one changing station, and one robot.

#### 2 State-Variable Representation

In a SVR, each state is represented by a tuple of values of n state variables  $\{x_1, \ldots, x_n\}$ , and each action is represented by a partial function that maps this tuple into some other tuple of values of the n state variables.

To build the SVR, the group has taken a very systematic approach of identifying and modeling the concepts. Because the industrial robot field is so broad, the group decided to limit its efforts to a single type of operation, namely kitting. A scenario was developed that described, in detail, the types of operations that would be performed in kitting, the sequencing of steps, the parts and machines that were needed, constraints on the process such as pre- and post-conditions, etc. For this scenario, a set of concepts were extracted and defined. These concepts served as the initial requirements for the kitting SVR. The concepts were then modeling in our SVR, building off of the definitions and relationships that were identified in the scenario. A SVR relies on the elements of constant variable symbols, object variable symbols, state variable symbols, and planning operators. These are defined for the kitting domain in the rest of this section.

#### 2.1 Constant Variable Symbols

For the kitting domain, there is a finite set of constant variable symbols that must be represented. In the SVR, constant variable symbols are partitioned into disjoint classes corresponding to the objects of the domain. The finite set of all constant variable symbols in the kitting domain is partitioned into the following sets:

- A set of *Part*: A *Part* is the basic item that will be used to fill a kit.
- A set of *PartsTray*: *Parts* arrive at the workstation in *PartsTrays*. Each *Part* is at a known position in the *PartsTray*. Each *PartsTray* contains one type of *Part*.
- A set of KitTray: A KitTray can hold Parts in known positions.
- A set of *Kit*: A *Kit* consists of a *KitTray* and, possibly, some *Parts*. A *Kit* is empty when it does not contain any *Part* and finished when it contains all the *Parts* that constitute a kit.

- $\blacksquare$  A set of WorkTable: A WorkTable is an area in the kitting workstation where KitTrays are placed to build Kits.
- A set of LargeBoxWithKits: A LargeBoxWithKits contains only finished Kits.
- A set of LargeBoxWithEmptyKitTrays: A LargeBoxWithEmptyKitTrays is a box that contains only empty KitTrays.
- A set of  $Robot \{robot\_1, robot\_2, ...\}$ : A Robot in the kitting workstation is a robotic arm that can move objects in order to build Kits.
- A set of *EndEffector*: *EndEffectors* are used in a kitting workstation to manipulate *Parts*, *PartsTrays*, *KitTrays*, and *Kits*. An *EndEffector* is attached to a *Robot* in order to grasp objects.
- $\blacksquare$  A set of EndEffHolder: An EndEffHolder is a storage unit that holds one type of EndEffector.
- A set of EndEffChStation: An EndEffChStation is made up of End-EffHolders.

#### 2.2 Object Variable Symbols

Object variable symbols are typed variables which range over a class or the union of classes of constant variable symbols. Examples of object variable symbols are  $r \in Robots$ ,  $kt \in KitTrays$ , etc.

#### 2.3 State Variable Symbols

A state variable symbol is defined as follows:

 $x: A_1 \times \cdots \times A_i \times S \to B_1 \cup \cdots \cup B_j \cup bool \cup \{\} \cup numeric \ (i, j \geq 1)$  is a function from the set of states (S) and at least one set of constant variable symbols  $A_1 \times \cdots \times A_i$  into a set  $B_1 \cup \cdots \cup B_j \cup bool \cup \{\} \cup numeric$  where:

- $\blacksquare$  B<sub>1</sub>  $\cup \cdots \cup$  B<sub>i</sub> is a set of constant variable symbols
- $\blacksquare$  bool is a boolean
- $\blacksquare$   $\{\}$  is an empty set

#### *numeric* is a numerical value

The use of state variable symbols reduces the possibility of inconsistent states and generates a smaller state space. The following state variable symbols are used in the kitting domain.

#### ■ endeff-location

 $EndEffector \times S \rightarrow Robot \cup EndEffHolder$ : designates the location of an EndEffector in the workstation. An EndEffector is either attached to a Robot or placed in an EndEffHolder.

#### ■ robot-with-endeff

 $Robot \times S \rightarrow EndEffector \cup \{\}:$  designates the EndEffector attached to a Robot if there is one attached, otherwise nothing.

#### on-wtable

 $WorkTable \times S \rightarrow Kit \cup KitTray \cup \{\}$ : designates the object placed on the WorkTable, i.e., a Kit, a KitTray, or nothing.

#### ■ kit-location

 $Kit \times S \rightarrow LargeBoxWithKits \cup WorkTable \cup Robot$ : designates the different possible locations of a Kit in the workstation, i.e., in a Large-BoxWithKits, on the WorkTable, or being held by a Robot.

#### ■ kittray-location

 $KitTray \times S \rightarrow LargeBoxWithEmptyKitTrays \cup WorkTable \cup Robot:$  designates the different possible locations of a KitTray in the workstation, i.e., in a LargeBoxWithEmptyKitTrays, on a WorkTable or being held by a Robot.

#### ■ part-location

 $Part \times S \rightarrow PartsTray \cup Kit \cup Robot$ : designates the different possible locations of a Part in the workstation, i.e., in a PartsTray, in a Kit, or being held by a Robot.

#### ■ robot-holds

 $Robot \times S \rightarrow KitTray \cup Kit \cup Part \cup \{\}$ : designates the object being held by a Robot, i.e., a KitTray, a Kit, a Part, or nothing. It is assumed that the Robot is already equipped with the appropriate EndEffector.

#### ■ lbwk-full

 $LargeBoxWithKits \times S \rightarrow bool$ : designates if a LargeBoxWithKits is full or not.

#### ■ Ibwekt-empty

 $LargeBoxWithEmptyKitTrays \times S \rightarrow bool:$  designates if a LargeBoxWithEmptyKitTrays is empty or not.

#### partstray-empty

 $PartsTray \times S \rightarrow bool$ : designates if a PartsTray is empty or not.

#### ■ endffector-type

 $EndEffector \times S \rightarrow KitTray \cup Kit \cup Part$ : designates the type of object an EndEffector can hold, i.e., a KitTray, a Kit, or a Part.

#### ■ endeffholder-holds-endeff

 $EndEffHolder \times S \rightarrow EndEffector \cup \{\}:$  designates wether an EndEffHolder is holding an EndEffector or nothing.

#### ■ endeffholder-location

 $EndEffHolder \times S \rightarrow EndEffChStation$ : designates the EndEffChStation where the EndEffHolder is located.

#### ■ endeffchstation-has-endeffholder

 $EndEffChStation \times S \rightarrow EndEffHolder$ : designates the EndEffHolder the EndEffChStation contains.

#### ■ found-part

 $PartsTray \times S \rightarrow Part \cup \{\}$ : designates wether a Part is found in a PartsTray or not.

#### ■ origin-part

 $Part \times S \rightarrow PartsTray$ : designates the PartsTray where the Part is found.

#### ■ quantity-parts-in-partstray

 $PartsTray \times S \rightarrow numeric$ : designates the number of parts that PartsTray contains.

#### ■ quantity-parts-in-kit

 $Kit \times PartsTray \times S \rightarrow numeric$ : designates the number of parts from PartsTray that Kit contains.

#### ■ capacity-parts-in-kit

 $Kit \times PartsTray \times S \rightarrow numeric$ : designates the number of parts from PartsTray that Kit can contain.

#### 2.4 Predicates and Functions

In PDDL, predicates are used to encode Boolean state variables, while functions are used to model updates of numerical values [2]. This section describes the predicates and functions derived from the state variables described in section 2.3. We recall the following definition of a state variable ((section 2.3))  $x: A_1 \times \cdots \times A_i \times S \to B_1 \cup \cdots \cup B_j \ (i, j \geq 1)$  that is used to convert state variables into predicates as follows:

$$\begin{array}{c} \blacksquare \ \, A_1 \times \cdots \times A_i \times S \to B_1 \cup \cdots \cup B_j \ (i,j \geq 1) \\ \\ \square \ \, \text{predicate\_1}(\mathcal{A},\mathcal{B}) \\ \\ \square \ \, \cdots \\ \\ \square \ \, \text{predicate\_n}(\mathcal{A},\mathcal{B}) \end{array}$$
 
$$\begin{array}{c} \square \ \, \text{where} \ \, \mathcal{A} \in \{A_1,\ldots,A_i\} \ \, \text{and} \ \, \mathcal{B} \in \{B_1,\ldots,B_i\} \ \, (i,j \geq 1) \end{array}$$

#### 2.4.1 Predicates

The state variables in our current kitting domain contains the following predicates.

#### ■ endeff-location

Ш	enderi-location-robot $(EnaEffector, Robot)$ ; IRUE in $EnaEffector$	)7
	is attached to <i>Robot</i>	
	$\verb endeff-location-endeffholder  (\textit{EndEffector}, \textit{EndEffHolder})  ; TRUE \\$	E
	iff EndEffector is in EndEffHolder	

#### ■ robot-with-endeff

any EndEffector

∐ r	${\sf robot\text{-}with\text{-}endeft}(Robot,\!End\!Ef\!fector)$	;TRUE	iff	Robot	is
$\epsilon$	equipped with EndEffector				
□ r	robot-with-no-endeff $(Robot)$ :TRUE iff	Robot is no	ot ear	iipped v	vith

#### ■ on-wtable

 $\square$  on-wtable-kit(WorkTable, Kit); TRUE iff Kit is on the WorkTable

	on-wtable-kittray ( $WorkTable, KitTray$ ); TRUE iff $KitTray$ is on the WorkTable
	$workTable\text{-}empty(\mathit{WorkTable})$ ; TRUE iff there is nothing on the $\mathit{WorkTable}$
■ kit-lo	ocation
	$\verb+kit-location-lbwk+(Kit, Large Box With Kits) ; TRUE iff \textit{Kit} is in the \textit{LargeBoxWithKits}$
	$kit\text{-}location\text{-wtable}(\mathit{Kit}, \mathit{WorkTable}) \; ; \\ TRUE \; iff \; \mathit{Kit} \; is \; on \; the \; \mathit{WorkTable} \\ \mathit{Table}$
	$kit\text{-}location\text{-}robot(\mathit{Kit}, Robot)$ ; TRUE iff $\mathit{Kit}$ is being held by the $Robot$
■ kittr	ay-location
	$\label{location-lowekt}  kittray-location-lbwekt (\textit{KitTray}, LargeBoxWithEmptyKitTrays) \\ ; TRUE iff \textit{KitTray} is in the \textit{LargeBoxWithEmptyKitTrays} \\$
	${\sf kittray\text{-}location\text{-}robot}(KitTray,Robot)$ ; TRUE iff $KitTray$ is being held by the $Robot$
	kittray-location-wtable ( $KitTray,WorkTable)$ ;TRUE iff $KitTray$ is on the WorkTable
■ part-	location
	$part\text{-}location\text{-}partsTray(\mathit{Part}, \mathit{PartsTray})$ ;TRUE iff $\mathit{Part}$ is in the $\mathit{PartsTray}$
	part-location-kit( $Part,Kit$ ); TRUE iff $Part$ is in the $Kit$
	part-location-robot(Part,Robot) ; TRUE iff $Part$ is being held by the $Robot$
■ robo	t-holds
	robot-holds-kittray ( $Robot, KitTray)$ ; TRUE iff $Robot$ is holding a $KitTray$
	robot-holds-kit(Robot,Kit) ; TRUE iff Robot is holding a $\mathit{Kit}$
	${\sf robot\text{-}holds\text{-}part}(Robot,Part) \ ; \\ {\sf TRUE} \ {\sf iff} \ Robot \ {\sf is} \ {\sf holding} \ {\sf a} \ Part$
	robot-empty(Robot); TRUE iff $Robot$ is not holding anything
■ Ibwk	-full

$\hfill \Box$ lbwk-not-full ( $LargeBoxWithKits)$ ;TRUE iff $LargeBoxWithKits$ is not full
■ Ibwekt-empty
$\begin{tabular}{ll} \Box & \begin{tabular}{ll} $
■ partstray-empty
$\square$ partstray-not-empty( $PartsTray$ ); TRUE iff $PartsTray$ is not empty
■ endeff-type
$\square$ endeff-type-kittray( $EndEffector, KitTray$ ) ;TRUE iff $EndEffector$ is capable of holding a $KitTray$
$\square$ endeff-type-kit( $EndEffector, Kit$ ); TRUE iff $EndEffector$ is capable of holding a $Kit$
$\Box$ endeff-type-part ( $EndEffector, Part)$ ;TRUE iff $EndEffector$ is capable of holding a Part
endeffholder-holds-endeff
$\square$ endeffholder-holds-endeff( $EndEffHolder,EndEffector$ ) ;TRUE iff $EndEffHolder$ is holding $EndEffector$
$\square$ endeffholder-empty( $EndEffHolder$ ) ;TRUE iff $EndEffHolder$ is empty (not holding an $EndEffector$ )
endeffholder-location
$\hfill\Box$ endeffholder-location ( $EndEffHolder, EndEffChStation)$ ; TRUE iff $EndEffHolder$ is in EndEffChStation
endeffchstation-has-endeffholder
$\begin{tabular}{ll} $\square$ endeffchstation-has-endeffholder ($EndEffChStation, EndEffHolder)$ \\ ; TRUE iff $EndEffChStation$ contains $EndEffHolder$ \\ \end{tabular}$
■ found-part
$\  \   \Box \   found\text{-part}(\mathit{Part}, \mathit{PartsTray}) \ ; TRUE \ iff \ \mathit{Part} \ is \ found \ in \ \mathit{PartsTray}$
■ origin-part
$\square$ origin-part( $Part, PartsTray$ ); TRUE iff $Part$ is from $PartsTray$ .

#### 2.4.2 Functions

In a planning model, numeric fluents represent function symbols that can take an infinite set of values. Introducing functions into planning not only makes it possible to deal with numerical values in a more general way than allowed for by a purely relational language but makes it possible to model operators in a more compact and sometimes also more natural way. The state variables in our current kitting domain contains the following functions.

■ quantity-parts-in-partstray
 □ quantity-parts-in-partstray(PartsTray); Quantity of parts in PartsTray
 ■ quantity-parts-in-kit
 □ quantity-parts-in-kit(Kit,PartsTray); Quantity of parts from PartsTray that is in Kit
 ■ capacity-parts-in-kit
 □ capacity-parts-in-kit(Kit,PartsTray); Quantity of parts from PartsTray that Kit can contain

#### 2.5 Planning Operators and Actions

The planning operators presented in this section are expressed in classical representation instead of state variable representation. In classical representation, states are represented as sets of logical atoms (predicates) that are true or false within some interpretation. Actions are represented by planning operators that change the truth values of these atoms.

#### 2.5.1 Planning Operators

In classical planning, a planning operator [4] is a triple o=(name(o), preconditions(o), effects(o)) whose elements are as follows:

■ name(o) is a syntactic expression of the form  $n(x_1, ..., x_k)$ , where n is a symbol called an operator symbol,  $x_1, ..., x_k$  are all of the object variable symbols that appear anywhere in o, and n is unique (i.e., no two operators can have the same operator symbol).

■ preconditions(o) and effects(o) are sets of literals (i.e., atoms and negations of atoms). Literals that are true in preconditions(o) but false in effects(o) are removed by using negations of the appropriate atoms.

Our kitting domain is composed of ten operators which are defined below.

1. take-kittray (robot, kittray, lbwekt, endeff, worktable): The Robot robot equipped with the EndEffector endeff picks up the KitTray kittray from the LargeBoxWithEmptyKitTrays lbwekt.

preconditions	effects
robot-empty(robot)	$\neg robot\text{-empty}(robot)$
$\verb+kittray-location-lbwekt+(kittray, lbwekt)$	$\neg$ kittray-location-lbwekt( $kittray, lbwekt$ )
${\sf Ibwekt-not-empty}(\mathit{lbwekt})$	kittray-location-robot(kittray,robot)
$robot ext{-with-endeff}(robot,endeff)$	robot-holds-kittray $(robot, kittray)$
${\tt endeff-location-robot}(\mathit{endeff},\mathit{robot})$	
${\sf worktable\text{-}empty}(worktable)$	
endeff-type-kittray( $endeff,kittray$ )	

#### ■ preconditions

	robot-empty	(robot)	): '	robot	does	onumber not	hold	anything	•
--	-------------	---------	------	-------	------	-------------	------	----------	---

- $\square$  kittray-location-lbwekt(kittray, lbwekt): kittray is in lbwekt.
- $\square$  lbwekt-not-empty(lbwekt): lbwekt is not empty (contains at least one kit tray).
- $\square$  robot-with-endeff(robot,endeff): robot is equipped with endeff.
- $\square$  endeff-location-robot(endeff,robot): The end effector is on the robot's arm.
- worktable-empty(worktable): After picking up an empty kit tray from a large box of empty kit trays, the robot would normally place the kit tray on the work table. To put a kit tray on the work table, it is necessary that there is nothing on top of the work table. If the robot is allowed to pick up the kit tray while there is another object on the work table, the planning system may not be able to find a solution when it comes to put the kit tray on the work table. Therefore, it is necessary to check that the top of worktable is clear even before the robot picks up a kit tray from the large box of empty kit trays.
- $\square$  endeff-type-kittray(endeff, kittray): endeff in the robot's arm must be capable of handling kittray.

#### ■ effects

 $\Box$  ¬robot-empty(robot): robot's end effector is no longer empty since it contains kittray.

work table.

	$\square$ ¬kittray-location-lbwekt( $kitt$ $lbwekt$ since it is in the rob	tray,lbwekt): kittray is no longer in ot's end effector.
	☐ kit-tray-location(kittray.robe	ot): kittray is in robot's end effector.
	$\Box$ robot-holds-kittray( $robot$ , $kit$	,
2.	put-kittray (robot, kittray, worktable Tray kittray on the WorkTable we	e): The Robot robot puts the Kit- orktable.
	preconditions	effects
	$ \hline \textbf{kittray-location-robot}(kittray,robot) $	$\neg$ kittray-location-robot $(kittray, robot)$
	${\sf robot\text{-}holds\text{-}kittray}(robot, kittray)$	$\neg$ robot-holds-kittray $(robot, kittray)$
	${\sf worktable\text{-}empty}(worktable)$	$\neg$ worktable-empty $(worktable)$
		kittray-location-wtable(kittray,worktable)
		robot-empty(robot)
		${\sf on\text{-}wtable\text{-}kittray}(worktable\text{-}kittray)}$
	lacksquare $preconditions$	
	$\square$ kittray-location-robot( $kittray$ tor.	y,robot): kittray is in robot's end effec-
	$\square$ robot-holds-kittray( $robot$ , $kit$	tray): robot holds kittray.
	- 1	: There is nothing on <i>worktable</i> .
	lacktriangledown effects	
	$\square$ ¬kittray-location-robot( $kittray$ ) end effector since it is place	ray, robot): kittray is no longer it robot's ed on worktable.
	$\  \   \Box \  \   \neg robot-holds-kittray(robot,k) \  \   more.$	ittray): robot is not holding kittray any-
	$\square$ ¬worktable-empty( $worktable$ since there is something on	e): worktable is not empty anymore top of it.
	$\Box$ kittray-location-wtable( $kittrate$	ay, worktable): kittray is on worktable.
	$\square$ robot-empty( $robot$ ): $robot$ is	s not holding anything.
	$\Box$ on-wtable-kittray( $worktable$ , of it.	kittray): worktable has kittray on top

3. take-kit(robot,kit,worktable,endeff): The  $Robot\ robot\ equipped\ with$  the  $EndEffector\ endeff\ picks\ up\ the\ Kit\ kit\ from\ the\ WorkTable$ 

preconditions	$\mid$ effects		
$ \overline{ \text{kit-location-wtable}(kit,worktable) } $	$\neg$ kit-location-wtable $(kit, worktable)$		
robot-empty(robot)	$\neg robot\text{-empty}(robot)$		
$on\text{-wtable\text{-}kit}(worktable,\!kit)$	$\neg$ on-wtable-kit $(worktable, kit)$		
${\sf robot\text{-}with\text{-}endeff}(\textit{robot}, \textit{endeff})$	$kit ext{-location-robot}(kit,robot)$		
$endeff\text{-}type\text{-}kit(\mathit{endeff},\!\mathit{kit})$	$robot ext{-holds-kit}(robot,kit)$		
	$worktable\text{-}empty(\mathit{worktable})$		
lacksquare $preconditions$			
$\Box$ kit-location-wtable( $kit, u$	worktable): kit is located on worktable.		
	ot is not holding any object.		
	(kit): worktable has kit on top of it.		
•	endeff): robot is equipped with endeff.		
•	it): The type of endeff is capable of han-		
lacktriangledown effects			
	worktable): kit is not on worktable.		
•	obot is holding an object (kit).		
	le,kit): worktable does not have kit on top		
**·	bot): kit is being held by robot.		
$\Box$ robot-holds-kit( $robot,kit$	,		
	ble): worktable does not have any object on		
4. put-kit(robot,kit,lbwk): The Re LargeBoxWithKits lbwk.	obot robot puts down the Kit kit in the		
preconditions ef	fects		
kit-location-robot $(kit, robot)$ $\neg$ k	$\overline{\text{kit-location-robot}(kit,robot)}$		
$robot-holds-kit(robot,kit)$ $\neg r$	$robot ext{-holds-kit}(robot,kit)$		
$Ibwk ext{-not-full}(\mathit{lbwk})$ kit	s-location-lbwk $(kit, lbwk)$		
rol	$pot-empty(\mathit{robot})$		
lacksquare $preconditions$			
•	bot): kit is held by robot.		
$\Box$ robot-holds-kit(robot,kit			
` '	k should not be full so it can contain $kit$ .		
,	v should not be full so it can contain not.		
■ effects			
$\sqcup$ $\neg$ kit-location-robot( $kit, n$ )	robot): kit is not being held by robot.		

	$\Box$ ¬robot-holds-kit( $robot,kit$ ): $robot$	_
	$\square$ kit-location-lbwk $(kit, lbwk)$ : $kit$	
	$\square$ robot-empty(robot): robot is n anymore).	ot holding anything (not holding $kit$
5.	look-for-part(robot,part,partstray,kit looks for the Part part in the PartT	/
	preconditions	effects
	part-not-searched	¬part-not-searched
	$robot ext{-}empty(\mathit{robot})$	found-part(partstray)
	$robot ext{-with-endeff}(\mathit{robot},\!\mathit{endeff})$	
	on-wtable-kit $(worktable, kit)$	
	$endeff$ -location-robot $(\mathit{endeff}, robot)$	
	part-location-partstray(part,partstray)	
	kit-location-wtable( $kit, worktable$ )	
	endeff-type-part(endeff,part)	
	partstray-not-empty(partstray)	
	■ preconditions	
		set to true in the initial state in the and means that a part has not been
	want the operator look-for-paerator take-part to simulate a picked up by a robot. It is nector is empty to prepare for the robot-with-endeff (robot, endeff Again, since we want the optake-part, we want to make	ould not be holding anything. We art to be directly followed by the opsensor identifying a part before being essary to check that robot's end effect execution of the operator take-part.  The property is equipped with endeff. because that robot is already equipped.
	with endeff.	
	$\Box$ on-wtable-kit( $worktable, kit$ ):	_
	$\sqcup$ endeff-location-robot(endeff, refor the operator take-part.	obot): endeff is on robot so it is ready
	$\square$ part-location-partstray( $part,pa$	rtstray): part is in partstray.
	$\Box$ kit-location-wtable( $kit, worktab$	ole): kit is on worktable.
	$\Box$ endeff-type-part(endeff,part):	endeff can handle part.
	,	: partstray contains at least one part.
	■ effects	

	□ ¬part-not-searched: This flag is set to true so that the operate look-for-part can be called again to look for another part in the workstation.	
	$\Box$ found-part(partstray): A part from partstray has been found.	
6.	6. take-part(robot,part,partstray,endeff,worktable,kit): The Robot robouses the EndEffector endeff to pick up the Part part from the PartTrapartstray.	
	$preconditions \hspace{1.5cm} ig \hspace{0.5cm} effects$	
		stray)
	on-wtable-kit $(worktable, kit)$ decrease quantity-parts-in-partstray kit-location-wtable $(kit, worktable)$ endeff-type-part $(endeff, part)$ partstray-not-empty $(partstray)$	y(partstray)
	$found ext{-}part(part,partstray)$	
	lacksquare $preconditions$	
	$\square$ part-location-partstray( $part, partstray$ ): $part$ to be picked up is $partstray$ .	in
	$\square$ robot-empty(robot): robot is not holding any object.	
	$\square$ endeff-location-robot(endeff,robot): endeff is on robot.	
	$\square$ robot-with-endeff(robot,endeff): robot is equipped with endeff.	
	$\Box$ on-wtable-kit(worktable,kit): worktable has kit on top of it.	
	□ kit-location-wtable(kit,worktable): kit is on worktable. Once a part is picked up by the robot, the next logical action would be to protect the part in the kit. For this to happen, the kit needs to be alread on the work table so it can hold the part.	ut ly
	$\square$ endeff-type-part(endeff,part): endeff is the type for part handling	g.
	$\Box$ partstray-not-empty(partstray): partstray is not empty and contain at least one part.	ns
	☐ found-part(part,partstray): part has been found in partstrate found-part is set to true in the effects of the operator look-for-partstray.	
	lacktriangledown $effects$	
	$\square$ ¬part-location-partstray( $part, partstray$ ): $part$ is not in $partstrapers$ anymore since it was picked up by $robot$ .	iy
	$\Box$ ¬robot-empty(robot): robot is now holding part and is not emp anymore.	ty

$\square$ robot-holds-part(robot,part): robot is holding part.			
	ray(partstray): After picking up a r of parts in partstray is decreased the decrease function.		
$7. \ \ \underline{put\text{-}part}(robot,part,kit,worktable,part,kit,wort,kit,wort,kit,wort,kit,wort,kit,wort,kit,wort,kit,wort,kit,wort,kit,wort,kit,wort,ki$	stray): The Robot robot puts the		
Part part in the Kit kit.			
preconditions	effects		
part-location-robot(part,robot)	$\neg part ext{-location-robot}(part,robot)$		
${\sf robot-holds-part}(robot,part)$	$\neg robot-holds-part(robot,part)$		
${\sf on\text{-}wtable\text{-}kit}(worktable,kit)$	$\neg found\text{-}part(part,partstray)$		
kit-location-wtable(kit,worktable)	robot-empty(robot)		
origin-part(part,partstray)	$part ext{-location-kit}(part,kit)$		
(<(quantity-parts-in-kit(kit,partstray))	(increase (quantity-kit(kit,partstray)))		
(capacity-kit(kit,partstray)))	part-not-searched		
<ul> <li>■ preconditions</li> <li>□ part-location-robot(part,robot): part is held by robot.</li> <li>□ robot-holds-part(robot,part): robot is holding part.</li> <li>□ on-wtable-kit(worktable,kit): worktable has kit on top of it.</li> <li>□ kit-location-wtable(kit,worktable): kit is on worktable.</li> </ul>			
<ul> <li>□ origin-part(part,partstray): part is from partstray. This is used to tell the type of part.</li> <li>□ (&lt; (quantity-kit(kit,partstray)) (capacity-kit(kit,partstray))): The quantity of parts of type partstray in kit should be lesser than the capacity kit can hold for this type of part.</li> </ul>			
lacksquare effects			
<ul> <li>□ ¬part-location-robot(part,robot): part is not held by robot.</li> <li>□ ¬robot-holds-part(robot,part): robot is not holding part.</li> <li>□ robot-empty(robot): robot is not holding any object.</li> <li>□ part-location(part,kit): part is located in kit.</li> <li>□ (increase (quantity-kit(kit,partstray))): Once part is placed in kit, the quantity of parts in kit is increased by one.</li> <li>□ part-not-searched: This flag is set to true so another part search (through the operator look-for-part) is made after part is placed in kit.</li> </ul>			

 $\square$  part-location-robot(part, robot): part is held by robot.

 ${\it EndEffHolder\ endeffholder}.$ 

8. attach-endeff(robot, endeff, endeffholder, endeffchstation): The Robot robot attaches the EndEffector endeff which is situated in the

```
preconditions
endeff-location-endeffholder (endeff, endeffholder)
robot-with-no-endeff(robot)
endeffholder-holds-endeff(endeffholder,endeff)
endeffholder-location (endeffholder, endeffchstation)
endef f chstation-contains-endef f holder (endef f chstation, endef f holder)
effects
\negendeff-location-endeffholder(endeff,endeffholder)
\negendeffholder-holds-endeff(endeffholder, endeff)
\negrobot-with-no-endeff(robot)
robot-empty(robot)
endeff-location-robot(endeff,robot)
robot-with-endeff(robot,endeff)
endeffholder-empty(endeffholder)
 ■ preconditions
     \square endeff-location-endeffholder (endeff, endeffholder): endeff is located
         in endeffholder.
      \square robot-with-no-endeff(robot): robot is not equipped with any endeff.
      \square endeffholder-holds-endeff(endeffholder,endeff):
                                                               endeffholder is
         holding endeff.
      \square endeffholder-location(endeffholder,endeffchstation): endeffholder is
         in endeffchstation.
      \square endeffchstation-contains-endeffholder (endeffchstation, endeffholder):
         endeffchstation contains endeffholder.
 \blacksquare effects
     \square ¬endeff-location-endeffholder(endeff,endeffholder): endeff is not in
         endeffholder anymore since it has been attached to robot.
      \Box ¬endeffholder-holds-endeff(endeffholder,endeff):
                                                                endeffholder is
         not holding endeff anymore.
      \square ¬robot-with-no-endeff(robot): robot is now equipped with endeff.
      \square robot-empty(robot): robot is not holding any object.
      \square endeff-location-robot(endeff,robot): endeff is on robot.
      \square robot-with-endeff(robot, endeff): robot is equipped with endeff.
      \square endeffholder-empty(endeffholder): endeffholder is not holding any
         endeff.
```

9. remove-endeff (robot, endeff, endeffholder, endeffchstation): The Robot robot removes the EndEffector endeff and puts it in the EndEffHolder endeffholder.

```
preconditions
endeff-location-robot(endeff,robot)
robot-with-endeff(robot,endeff)
robot-empty(robot)
endeffholder-location (endeffholder, endeffchstation)
endef f chstation-contains-endef f holder (endef f chstation, endef f holder)
endeffholder-empty(endeffholder)
effects
\negendeff-location-robot(endeff, robot)
\negrobot-with-endeff(robot, endeff)
\negendeffholder-empty(endeffholder)
endeff-location-endeffholder(endeff, endeffholder)
endeffholder-holds-endeff(endeffholder,endeff)
robot-with-no-endeff(robot)
 ■ preconditions
      \square endeff-location-robot(endeff,robot): endeff is on robot.
      \square robot-with-endeff(robot,endeff): robot is holding endeff.
      \square robot-empty(robot): robot is not holding anything.
      \square endeffholder-location(endeffholder,endeffchstation): endeffholder is
         in endeffchstation.
      \square endeffchstation-contains-endeffholder (endeffchstation, endeffholder)
      \square endeffholder-empty(endeffholder): endeffholder does not contain
         endeff.
 ■ effects
     \square ¬endeff-location-robot(endeff,robot): endeff is not on robot any-
      \square ¬robot-with-endeff(robot, endeff): robot does not have endeff any-
         more.
      □ ¬endeffholder-empty(endeffholder): endeffholder does not contain
         any endeff.
      \square endeff-location-endeffholder (endeff, endeffholder):
                                                                    endeff
                                                                               is
         situated in endeffholder.
      \square endeffholder-holds-endeff(endeffholder,endeff): endeffholder holds
      \square robot-with-no-endeff(robot): robot is not equipped with endeff.
```

10. create-kit(kit,kittray,worktable): The KitTray kittray is converted into the Kit kit once the KitTray kittray is on the WorkTable worktable.

preconditions	effects	
$- \\ \\ \text{on-wtable-kittray}(worktable, kittray) \\$	$\neg$ on-wtable-kittray( $worktable, kittray$ )	
	on-wtable-kit $(worktable, kit)$	
	$kit ext{-location-wtable}(kit, worktable)$	
<ul><li>■ preconditions</li><li>□ on-wtable-kittray(worktable, of it.</li></ul>	kittray): worktable has kittray on top	
lacksquare $effects$		
$\square$ ¬on-wtable-kittray( $worktable$ , $kittray$ ): The object $kittray$ is destroyed and is thus not on $worktable$ anymore.		
$\square$ on-wtable-kit( $worktable, kit$ ): $worktable$ now has $kit$ on top of it.		
$\square$ kit-location-wtable( $kit, worktable$ ): $kit$ is on $worktable$ .		

#### 2.5.2 Actions

An action a can be obtained by substituting the object variable symbols that appear anywhere in the operator with constant variable symbols. For instance, the operator take-part(robot,part,partstray,endeff) in the kitting domain can be translated into the action  $take-part(robot\_1,part\_1,partstray\_1,endeff\_2)$  where  $robot\_1$ ,  $part\_1$ ,  $partstray\_1$ , and  $endeff\_2$  are constant variable symbols in the classes Robot, Part, PartsTray, and EndEffector, respectively.

#### 3 Planning Language

The Planning Domain Definition Language (PDDL) [3] is an attempt by the domain independent planning community to formulate a standard language for planning. A community of planning researchers has been producing planning systems that comply with this formalism since the first International Planning Competition held in 1998. This competition series continues today, with the seventh competition being held in 2011. PDDL is constantly adding extensions to the base language in order to represent more expressive problem domains. Our work is based on PDDL Version 3.

By placing our knowledge in a PDDL representation, we enable the use of an entire family of open source planning systems. Each PDDL file-set consists of two files that specify the domain and the problem.

#### 3.1 The PDDL Domain File

The PDDL domain file is composed of four sections that include requirements, types, predicates and functions, and actions. An excerpt of the PDDL domain file is depicted in Figure 1.

- line 1: The keyword domain signals a planner that this file contains information on the domain. kitting-domain is the name given to the domain.
- line 2: The :requirements field specifies which section the domain relies on. The planning system can examine this statement to determine if it is capable of solving problems in this domain. A keyword (symbol starting with a colon) used in a :requirements field is called a requirement flag; the domain is said to declare a requirement for that flag. The requirement flags present in the kitting domain are:

$\tt:strips:$ The most basic subset of PDDL, consisting of STRIPS only.
:typing: PDDL has a special syntax for declaring parameter and object types. :typing allows types names in declaration of variables.
:fluents: A domain's set of requirements allow a planner to quickly tell if it is likely to be able to handle the domain. For

example, this version of the kitting world requires fluents numeric,

```
1. (define (domain kitting-domain)
 2.
        (:requirements :strips :typing :fluents)
 3.
        (:types
 4.
            {\tt EndEffector}
 5.
            EndEffectorHolder
 6.
 7.
           KitTray
 8.
            LargeBoxWithEmptyKitTrays
9.
            LargeBoxWithKits
10.
            Part
11.
            PartsTray
            {\tt EndEffectorChangingStation}
12.
13.
            Robot
            WorkTable
14.
15.
16.
        (:predicates
17.
18.
               (endeffector-location-robot ?endeffector - EndEffector ?robot - Robot)
               (on-worktable-kit ?worktable - WorkTable ?kit - Kit)
19.
20.
21.
22.
        (:functions
23.
               (quantity-partstray ?partstray - PartsTray)
24.
               (quantity-kit ?kit - Kit ?partstray - PartsTray)
25.
               (capacity-kit ?kit - Kit ?partstray - PartsTray)
       )
26.
27.
28.
        (:action take-kittray
29.
            :parameters(
                ?robot - Robot
30.
31.
                ?kittray - KitTray
32.
                ?largeboxwithemptykittrays - LargeBoxWithEmptyKitTrays
33.
                ?endeffector - EndEffector
                ?worktable - WorkTable)
34.
35.
            :precondition(and
36.
                (robot-empty ?robot)
37.
                (lbwekt-not-empty ?largeboxwithemptykittrays)
38.
                (robot-with-endeffector ?robot ?endeffector)
39.
                (kittray-location-lbwekt ?kittray ?largeboxwithemptykittrays)
40.
                (endeffector-location-robot ?endeffector ?robot)
41.
                (worktable-empty ?worktable)
42.
                (endeffector-type-kittray ?endeffector ?kittray))
43.
            :effect(and
44.
                (robot-holds-kittray ?robot ?kittray)
45.
                (kittray-location-robot ?kittray ?robot)
46.
                (not (robot-empty ?robot))
47.
                (not (kittray-location-lbwekt ?kittray ?largeboxwithemptykittrays)))
       )
48.
49.)
50.
```

Figure 1: Excerpt of the PDDL domain file for kitting.

so a straight STRIPS-representation planner would not be able to handle it. A fluent is a term (:functions) with time-varying value (i.e., a value that can change as a result of performing an action).

- line 3-15: Type names have to be declared before they are used (before :predicates and :functions). This is done with the declaration (:types  $name_1 \ldots name_n$ ).
- line 17–20: The :predicates part of a domain definition specify only what are the predicate names used in the domain, and their number of arguments (and argument types, if the domain uses :typing). The "meaning" of a predicate, in the sense of for what combinations of arguments it can be true and its relationship to other predicates, is determined by the effects that actions in the domain can have on the predicate, and by what instances of the predicate are listed as true in the initial state of the problem definition.

It is common to make a distinction between static and dynamic predicates. A *static* predicate is not changed by any action. Thus in a problem, the true and false instances of a *static* predicate will always be precisely those listed in the initial state specification of the problem definition. Note that there is no syntactic difference between *static* and *dynamic* predicates in PDDL, they look exactly the same in the :predicates declaration part of the domain.

A predicate is build using the structure (predicate\_name ?X - type\_of\_X). A list of parameters of the same type in a predicate can be abbreviated to (predicate\_name ?X ?Y ?Z - type\_of\_XYZ). Note that the hyphen between parameter and type name is surrounded by whitespace.

- line 22–26: A fluent is similar to a state variable/predicate except that its value is a number instead of true or false. The initial value of a function is set in the initial state of the problem file and changes when an action is executed. The declaration of functions is similar to predicates.
- line 28–48: The domain definition contains operators (called *actions* in PDDL). An action statement specifies a way that a planner affects the state of the world. The statement includes parameters, preconditions, and effects. All parts of an action definition except the name are, according to the PDDL specification, optional (although, of course, an action without effects is pretty useless). However, for an action that has no preconditions some planners may require an "empty" precon-

dition, on the form :precondition () or :precondition (and), and some planners may also require an empty :parameter list for actions without parameters).

- □ line 29–34: The :parameters section declare all the parameters used by predicates and functions in preconditions and effects.
- □ line 35–42: The :preconditions section is a conjunction of predicates and functions that need to be true in the world in order for the action to be invoked.
- □ line 43–47: The :effects equation dictates the changes in the world that will occur due to the execution of the action.

#### 3.2 PDDL Problem File

The second file of the PDDL file-set is a problem file. The problem file specifies information about the specific instance of the given problem. This file contains the initial conditions and definition of the world (in the init section) and the final state that the world must be brought to (in the goal section). Using an example of kit to build, this section only describes the initial and goal states explicitly. The operators detailed in Section 2.5 are used by a planner to generate the other states as needed.

In this example, the *Robot* has to build a kit that contains two *Parts* of type A, two *Part* of type B and one *Part* of type C. The kitting process is completed once the *Kit* is placed in the *LargeBoxWithKits*. The PDDL problem file for the kitting domain is presented below.

```
    (define (problem kitting-problem)

2.
       (:domain kitting-domain)
3.
       (:objects
 4.
           robot 1 - Robot
 5.
            changing_station_1 - EndEffectorChangingStation
 6.
           kit_tray_1 - KitTray
           kit_a2b2c1 - Kit
 7.
8.
           empty_kit_tray_supply - LargeBoxWithEmptyKitTrays
9.
           finished_kit_receiver - LargeBoxWithKits
10.
           work_table_1 - WorkTable
           part_a_tray part_b_tray part_c_tray - PartsTray
11.
           part_a_1 part_a_2 part_a_3 part_a_4 - Part
12.
13.
           part_b_1 part_b_2 part_b_3 part_b_4 - Part
14.
           part_c_1 part_c_2 part_c_3 part_c_4 - Part
15.
           part_gripper tray_gripper - EndEffector
16.
           part_gripper_holder tray_gripper_holder - EndEffectorHolder
17.
18.)
19. (:init
20.
       (robot-with-no-endeffector robot_1)
21.
       (part-not-searched)
22.
       (lbwekt-not-empty empty_kit_tray_supply)
       (lbwk-not-full finished_kit_receiver)
23.
24.
        (partstray-not-empty part_a_tray)
25.
        (partstray-not-empty part_b_tray)
26.
        (partstray-not-empty part_c_tray)
27.
        (endeffector-location-endeffectorholder part_gripper part_gripper_holder)
28.
        (endeffector\hbox{-location--endeffector} holder \hbox{ tray\_gripper tray\_gripper\_holder})
29.
        (endeffectorholder-holds-endeffector part_gripper_holder part_gripper)
30.
       (endeffectorholder-holds-endeffector tray_gripper_holder tray_gripper)
31.
        (endeffectorholder-location tray_gripper_holder changing_station_1)
32.
        (endeffectorholder-location part_gripper_holder changing_station_1)
33.
        (endeffectorchangingstation-contains-endeffectorholder changing_station_1 tray_gripper_holder)
34.
        (endeffectorchangingstation-contains-endeffectorholder changing_station_1 part_gripper_holder)
35.
        (worktable-empty work_table_1)
36.
        (kittray-location-lbwekt kit_tray_1 empty_kit_tray_supply)
37.
38.
        (part-location-partstray part_a_1 part_a_tray)
39.
        (part-location-partstray part_a_2 part_a_tray)
40.
        (part-location-partstray part_a_3 part_a_tray)
41.
        (part-location-partstray part_a_4 part_a_tray)
42.
        (part-location-partstray part_b_1 part_b_tray)
43.
        (part-location-partstray part_b_2 part_b_tray)
44.
        (part-location-partstray part_b_3 part_b_tray)
45.
        (part-location-partstray part_b_4 part_b_tray)
46.
        (part-location-partstray part_c_1 part_c_tray)
47
        (part-location-partstray part_c_2 part_c_tray)
48.
        (part-location-partstray part_c_3 part_c_tray)
49.
       (part-location-partstray part_c_4 part_c_tray)
50.
51.
       (endeffector-type-part part_gripper part_a_1)
52.
       (endeffector-type-part part_gripper part_a_2)
53.
        (endeffector-type-part part_gripper part_a_3)
54.
```

(endeffector-type-part part\_gripper part\_a\_4)

```
55.
        (endeffector-type-part part_gripper part_b_1)
56.
        (endeffector-type-part part_gripper part_b_2)
57.
        (endeffector-type-part part_gripper part_b_3)
58.
        (endeffector-type-part part_gripper part_b_4)
        (endeffector-type-part part_gripper part_c_1)
59.
        (endeffector-type-part part_gripper part_c_2)
60.
        (endeffector-type-part part_gripper part_c_3)
61.
        (endeffector-type-part part_gripper part_c_4)
62.
63.
        (endeffector-type-kittray tray_gripper kit_tray_1)
64.
        (endeffector-type-kit tray_gripper kit_a2b2c1)
65.
66.
        (= (capacity-kit kit_a2b2c1 part_a_tray) 2)
67.
        (= (capacity-kit kit_a2b2c1 part_b_tray) 2)
        (= (capacity-kit kit_a2b2c1 part_c_tray) 1)
68.
69.
        (= (quantity-kit kit_a2b2c1 part_a_tray) 0)
70.
        (= (quantity-kit kit_a2b2c1 part_b_tray) 0)
        (= (quantity-kit kit_a2b2c1 part_c_tray) 0)
71.
72.
        (= (quantity-partstray part_a_tray) 4)
        (= (quantity-partstray part_b_tray) 4)
73.
74.
        (= (quantity-partstray part_c_tray) 4)
75.
76.
        (origin-part part_a_1 part_a_tray)
77.
        (origin-part part_a_2 part_a_tray)
78.
        (origin-part part_a_3 part_a_tray)
79.
        (origin-part part_a_4 part_a_tray)
80.
        (origin-part part_b_1 part_b_tray)
        (origin-part part_b_2 part_b_tray)
81.
82.
        (origin-part part_b_3 part_b_tray)
        (origin-part part_b_4 part_b_tray)
83.
        (origin-part part_c_1 part_c_tray)
84.
85.
        (origin-part part_c_2 part_c_tray)
86.
        (origin-part part_c_3 part_c_tray)
87.
        (origin-part part_c_4 part_c_tray)
88.)
89.
90. (:goal
91.
        (and
92.
            (= (quantity-kit kit_a2b2c1 part_a_tray)
            (capacity-kit kit_a2b2c1 part_a_tray))
93.
94.
            (= (quantity-kit kit_a2b2c1 part_b_tray)
95.
            (capacity-kit kit_a2b2c1 part_b_tray))
96.
            (= (quantity-kit kit_a2b2c1 part_c_tray)
97.
            (capacity-kit kit_a2b2c1 part_c_tray))
98.
            (kit-location-lbwk kit_a2b2c1 finished_kit_receiver)
99.
        )
100.)
```

- line 1: Signal a planner that the file contains all the element part of a problem. kitting-problem is the name given to this problem.
- line 2: :domain refers to the domain that the current problem is associated to. In this case, the problem refers to the domain kitting-domain. Note that kitting-domain is the name given to the kitting domain as presented in section 3.1.

■ line 3-17: :objects declare objects present in the problem instance. The syntax for :objects is  $object_1$  - Type . . .  $object_n$  - Type.

#### 3.2.1 Initial State

The initial state  $S_0$  (Figure 2) defines the environment in its initial condition. The initial state of the kitting problem in PDDL format is described below.

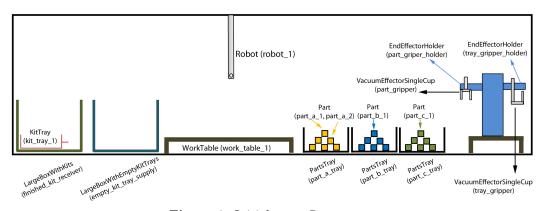


Figure 2: Initial state  $S_0$ .

- line 19: :init signals a planner that the predicates and functions in this section are true in the initial state.
- line 20–87: Predicates true in the initial state of the environment. Since PDDL uses a close world assumption, predicates that are not present in the initial state are automatically set to false. This section also set the initial values for functions. Some relevant sections are presented:
  - □ line 21: The predicate part-not-searched is set to true so that the operator *look-for-part* can be activated during a plan search.
  - □ line 65–67: Functions describing the quantity of parts of a type that  $kit_{-}a2b2c1$  can contain. In this example,  $kit_{-}a2b2c1$  can have 2 parts of type A  $(part_{-}a_{-}tray)$ , 2 parts of type B  $(part_{-}b_{-}tray)$ , and 1 part of type C  $(part_{-}c_{-}tray)$ .
  - □ line 68–70: Functions that represent the quantity of parts of a specific type that are already in  $kit_{-}a2b2c1$ .  $kit_{-}a2b2c1$  has no parts of type A, B, and C.

- □ line 71–73: Functions that describe the quantity of parts available in their respective parts tray. This also can be read as: In the workstation, there are 4 parts of type A available, 4 parts of type B available, and 4 parts of type C available.
- $\Box$  line 75–86: Predicates that describe the type of each specific part in the workstation.

#### 3.2.2 Goal State

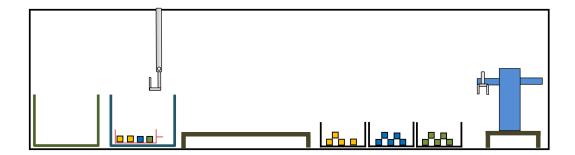


Figure 3: Goal state  $S_G$ .

Figure 3 depicts the goal state  $S_G$  for the kitting workstation. The expression of the goal state in PDDL is described below.

- line 89: :goal is a keyword used to signal a planner about the goal state to reach. All the predicates and functions in the goal state must be true.
- line 91–93: The quantity of parts of a specific type in  $kit\_a2b2c1$  should match the capacity of parts of a specific type for  $kit\_a2b2c1$ . The quantity of parts in  $kit\_a2b2c1$  is increased in the operator  $put\_part$ . The initial quantity of parts in  $kit\_a2b2c1$  and its capacity are set in the initial state. Note that we are not specifying which instance of Part should go in  $kit\_a2b2c1$  but rather the number of Parts of a specific type that  $kit\_a2b2c1$  must have.
- line 94:  $kit\_a2b2c1$  should be placed in the large box with kits  $finished\_kit\_receiver$ .

#### 3.3 Plan

This section shows an example of a plan generated by a planner (see section 4) for the PDDL domain and problem files discussed previously in this document.

Figure 4 displays the different states and actions used by the planner to generate a plan starting from the initial state  $S_0$  to the goal state  $S_G$ . The actions  $A_1 \ldots A_{17}$  are described below.

- $A1:(attach-endeffector robot\_1 tray\_gripper tray\_gripper\_holder changing\_station\_1)$
- A2:(take-kittray robot\_1 kit\_tray\_1 empty\_kit\_tray\_supply tray\_gripper work\_table\_1)
- A3:(put-kittray robot\_1 kit\_tray\_1 work\_table\_1)
- $\blacksquare$  A4:(create-kit kit\_a2b2c1 kit\_tray\_1 work\_table\_1)
- $A5:(remove-endeffector robot\_1 tray\_gripper tray\_gripper\_holder changing\_station\_1)$
- A6:(attach-endeffector robot\_1 part\_gripper part\_gripper\_holder changing\_station\_1)
- A7:(look-for-part robot\_1 part\_c\_1 part\_c\_tray kit\_a2b2c1 work\_table\_1 part\_gripper)
- A8:(take-part robot\_1 part\_c\_1 part\_c\_tray part\_gripper work\_table\_1 kit\_a2b2c1)
- $\blacksquare$  A9:(put-part robot\_1 part\_c\_1 kit\_a2b2c1 work\_table\_1 part\_c\_tray)
- A10:(look-for-part robot\_1 part\_b\_2 part\_b\_tray kit\_a2b2c1 work\_table\_1 part\_gripper)
- A11:(take-part robot\_1 part\_b\_2 part\_b\_tray part\_gripper work\_table\_1 kit\_a2b2c1)
- A12:(put-part robot\_1 part\_b\_2 kit\_a2b2c1 work\_table\_1 part\_b\_tray)
- A13:(look-for-part robot\_1 part\_b\_1 part\_b\_tray kit\_a2b2c1 work\_table\_1 part\_gripper)
- A14:(take-part robot\_1 part\_b\_1 part\_b\_tray part\_gripper work\_table\_1 kit\_a2b2c1)

- $\blacksquare$  A15:(put-part robot\_1 part\_b\_1 kit\_a2b2c1 work\_table\_1 part\_b\_tray)
- A16:(look-for-part robot\_1 part\_a\_2 part\_a\_tray kit\_a2b2c1 work\_table\_1 part\_gripper)
- A17:(take-part robot\_1 part\_a\_2 part\_a\_tray part\_gripper work\_table\_1 kit\_a2b2c1)
- A18:(put-part robot\_1 part\_a\_2 kit\_a2b2c1 work\_table\_1 part\_a\_tray)
- A19:(look-for-part robot\_1 part\_a\_1 part\_a\_tray kit\_a2b2c1 work\_table\_1 part\_gripper)
- A20:(take-part robot\_1 part\_a\_1 part\_a\_tray part\_gripper work\_table\_1 kit\_a2b2c1)
- A21:(put-part robot\_1 part\_a\_1 kit\_a2b2c1 work\_table\_1 part\_a\_tray)
- A22:(remove-endeffector robot\_1 part\_gripper part\_gripper\_holder changing\_station\_1)
- A23:(attach-endeffector robot\_1 tray\_gripper tray\_gripper\_holder changing\_station\_1)
- $\blacksquare$  A24:(take-kit robot\_1 kit\_a2b2c1 work\_table\_1 tray\_gripper)
- $\blacksquare$  A25:(put-kit robot\_1 kit\_a2b2c1 finished\_kit\_receiver)

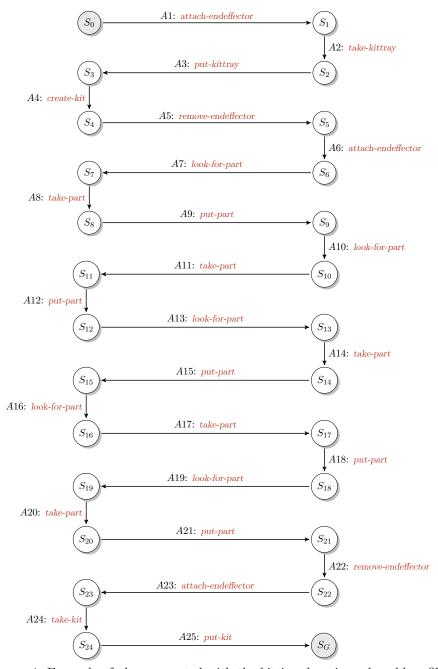


Figure 4: Example of plan generated with the kitting domain and problem files.

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#### 4 Planner

This section describes the steps to install and run a planner on the PDDL domain and problem files in order to generate a plan. The planner uses a forward-chaining partial-order planning [1].

#### 4.1 Requirements

The planner requires:

- cmake
- The CBC mixed integer programming solver (https://projects.coin-or.org/Cbc/)
- perl, bison and flex to build the parser

These are packaged with most Linux distributions - on Ubuntu/Debian, the following should suffice:

sudo apt-get install cmake coinor-libcbc-dev coinor-libclp-dev
\
coinor-libcoinutils-dev bison flex

#### 4.2 Download and Install CBC

The CBC source code can be obtained using subversion:

- svn co https://projects.coin-or.org/svn/Cbc/stable/2.7 coin-Cbc: Issues the subversion command to obtain the source code.
- cd coin-Cbc
- ./configure -C: Runs a configure script that generates the make file.
- make: Builds the Cbc library and executable program.
- make test: Builds and runs the Cbc unit test program.
- make install: Installs libraries, executables and header files in directories coin-Cbc/lib, coin-Cbc/bin and coin-Cbc/include.

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#### 4.3 Compile the Planner

Before compiling the planner, the file *compile/CMakeCache.txt* should be edited as follows. Note that <path> is the absolute path that leads to the coin-Cbc directory.

- CBC\_INCLUDES:PATH = <path>/coin-Cbc/build/include
- CGL\_INCLUDES:PATH = <path>/coin-Cbc/build/include
- CLP\_INCLUDES:PATH = <path>/coin-Cbc/build/include
- COINUTILS\_INCLUDES:PATH = <path>/coin-Cbc/build/include/coin
- OSI\_INCLUDES:PATH = <path>/coin-Cbc/build/include

To compile the planner, one should use:

./build

#### 4.4 Run the Planner

To run the planner, the path to the PDDL domain and problem files should be identified. The format of the PDDL files must be .pddl. The following command run the planner on the PDDL files.

./plan <domain> <problem> <solution>

Where <domain> and <problem> are the PDDL domain and problem files, respectively.<solution> is the output file containing the plan.

#### 5 Tools

#### 5.1 The Generator

The Generator tool is a graphical user interface developed in Java, allowing the user to store data from OWL files into a MySQL database. This tool also permits the user to query the database using the C++ function calls. The tool Generator is composed of the following functionalities:

- 1. Convert OWL documents into SQL syntaxes (OWL to SQL).
- 2. Translate SQL syntaxes to OWL language in order to modify an OWL document (SQL to OWL).
- 3. Convert the OWL language into C++ classes (OWL to C++).

To date, only steps 1. and 3. have been implemented and will be covered in this document.

#### 5.1.1 Prequisites

The description of the Generator tool is given for a Ubuntu Linux system. To run and use the Generator tool, different applications must be installed on the system.

Java Runtime Environment The Generator tool comes as a jar file. As such, the Java Runtime Environment should be installed on your system. This application can be found at www.oracle.com.

MySQL Server and Client The MySQL server and client should be installed and running on your system.

- sudo apt-get update (Update the package management tools)
- sudo apt-get dist-upgrade (Install the latest software)
- sudo apt-get install mysql-server mysql-client (Install the MySQL server and client packages). You will be asked to enter a password.

When done, you have a MySQL database ready to run. The following command will allow you to run MySQL.

- $\blacksquare$  mysql -u root -p
- Enter the same password you used when you installed MySQL.

Finally, we need the plugin libmysqlcppconn-dev which allows C++ to connect to MySQL databases. It can be installed as follows:

■ sudo apt-qet install libmysqlcppconn-dev

#### 5.1.2 How to Run the Generator Tool

The Generator tool can be launched using either one of these two following methods:

- 1. java -jar Generator.jar
- 2. Right-click on Generator.jar and select the option "Open With Open-JDK Java 6 Runtime". Note that this message will be different for future releases of the Java Runtime Environment.

#### 5.1.3 Functionalities

As mentioned in the Introduction, we are covering only steps 1. and 3. in the rest of this document, i.e., OWL to SQL and OWL to C++, respectively.

**OWL to SQL** To convert OWL classes and instances to SQL, the Owl to SQL tab should be selected (see Figure 5). The different fields are:

#### Generate SQL Files

■ Ontology Path: This field requires the file kittingInstances.owl. Before doing so, you need to modify one line in this file. Open it with a text editor and find the line Import(<file:kittingClasses.owl>).

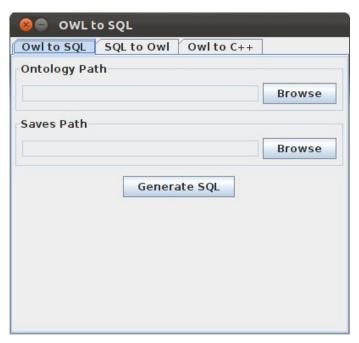


Figure 5: Owl to SQL tab.

Modify this line by giving the absolute path to the file kitting-Classes.owl. You should should have something that looks like Import(<file:/home/username/NIST/ipmas/Generator/kittingClasses.owl>). When this is done, save the file, and browse to kittingInstances.owl using the "Browse" button.

■ Browse to the directory where you want to save the SQL files.

Once the two previous steps are done, click on "Generate SQL". You should receive a message confirming the generation of the SQL files: kittingInstances.owlCreateTable.sql and kittingInstances.owlInsertInto.sql. The former is used to create tables, the latter is used to populate these tables;

**SQL Tables and Insertions** The next step is to create a database and to populate it.

■ Connect to mysql using  $mysql - u \ root - p$ , then enter your password. You should be in the mysql shell if this succeeded (mysql>).

■ Delete a previous database (if you already used this tool and you want to replace the existing database with this new one): mysql> DROP DATABASE OWL; (OWL is the name of the old database).

- Create a database:
  - □ mysql> CREATE DATABASE OWL;. Here, OWL is the name of the database (you can use a name of your choice).
  - $\square$  Before performing the following commands, we need to tell MySQL which database we are planning to work with (*OWL* in our case). This is done using:

- Populate the database with tables using kittingInstances.owlCreateTable.sql.
  - $\square$  mysql> source <path>/kittingInstances.owlCreateTable.sql;
- Populate the tables with data using kittingInstances.owlInsertInto.sql:
  - □ mysql> source <path>/kittingInstances.owlInsertInto.sql;

<path> designs the absolute path to the appropriate file.

**OWL to C++** The "Owl to C++" tab (see Figure 6) is used to generate C++ classes and scripts allowing the connection between C++ and MySQL. The different fields are explained below:

- Ontology Path: This is the path to the ontology (kittingClasses.owl in our example).
- Saves Path: Directory where the C++ files and scripts will be generated.
- Url: This is the url of the database. It's usually the IP address of the machine hosting the database (127.0.0.1 if it is local).
- User name: User name used to connect to the MySQL database.
- **Password**: Password associated to the user name to connect to the MySQL database.

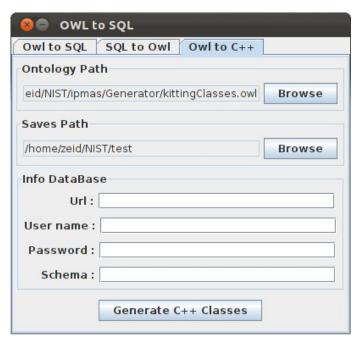


Figure 6: Owl to C++ tab.

■ Schema: This is the name of the database (*OWL* in our example).

When all the fields are completed, click the "Generate C++ Classes" button to generate C++ and script files.

#### 5.2 XML to OWL

This section describes tools that are used to generate OWL files from XML files.

#### 5.2.1 owlPrinter

owlPrinter reads an XML kitting data file corresponding to an XML schema for kitting (kitting.xsd) and writes an OWL instance file corresponding to the OWL class file kittingClasses.owl. The kitting.xsd file contains the same conceptual model as the kittingClasses.owl file, but in a different language.

The owlPrinter is useful because there is no OWL tool that will help generate

an OWL instance file and check the file adequately against an OWL class file. That is because OWL uses an open world model in which anything not explicitly or implicitly illegal is allowed. Hence many things that are errors to the writer of the instance file are not OWL errors. For example, if the name of an instance is misspelled, OWL will assume that there is a new instance that has not been explicitly declared as such, which is OK in OWL. If a reference to an instance name is misspelled in an XML data file corresponding to the kitting.xsd schema, that will be caught automatically by the owlPrinter (and other readily available XML tools). Several other types of error will not be caught by OWL tools but will not be made or will be detected if the OWL printer is used.

Another OWL problem that disappears in XML is that in OWL, there is no distinction between an instance file and a class file. An instance file can modify classes, intentionally, or accidentally. In XML there is no way a data file can modify a model.

To use the owlPrinter, use a text editor such as emacs or an XML tool such as XMLSpy to write an XML data file corresponding to the kitting.xsd schema and then run it through the owlPrinter with a command of the form:

bin/owlPrinter [XML file in] [OWL file out]

For example, the command

bin/owlPrinter data/kittingInstances.xml junk

will print the file junk, which will be identical to the kittingInstances.owl file in the owl directory (except for a couple comments).

#### 5.2.2 kittingParser

The kittingParser may be used to check an XML data file against the kitting.xsd schema. The schema is hard-coded into the kittingParser. If there is any error in the XML data file, the kittingParser prints a message and quits. If there is no error, the input file is echoed by printing an output file whose name is the same as that of the input file with "echo" appended. To run the kittingParser, give a command of the form:

#### bin/kittingParser [XML kitting data file in]

For example, the command

#### bin/kittingParser data/kittingInstances.xml

will read kittingInstances.xml and write kittingInstances.xmlecho. The two files will be identical except for comments. If the format of an input file differs from the format used for printing the output file, the two files will differ, but only in format.

The owlPrinter makes the same checks as the kittingParser, so there is no need to use the kittingParser.

All of the source code for the kittingParser was generated automatically by the GenXMiller generator.

#### 5.2.3 compactOwl and compareOwl

There is no need to read this section unless you are interested in how the *owlPrinter* was debugged.

The *compactOwl* and *compareOwl* utilities are used for checking that two different OWL instance files have the same statements. They have been used as follows to debug the *owlPrinter* (and the kitting.xsd file and the kittingInstances.xml file and the kittingInstances.owl file).

- 1. Write kitting.xsd to model the same information as kittingClasses.owl.
- 2. Write kittingInstances.xml to correspond to kitting.xsd and contain the same information as kittingInstances.owl.
- 3. Build the owlPrinter.
- 4. Run kittingInstances.xml through the *owlPrinter* to produce kittingInstances.owl.
- 5. Run kittingInstances.owl through *compactOwl* to produce one version of kittingInstancesCompact.owl.
- 6. Run kittingInstances.owl through *compactOwl* to produce a second version of kittingInstancesCompact.owl.

7. Run the two versions of kittingInstancesCompact.owl through compare-Owl. If compareOwl reports that the two files have the same statements, that means that steps 1, 2, and 3 have been done correctly, so debugging is finished. If compareOwl reports a pair of statements that differ, figure out why, go back to step 1, 2, or 3 (or edit kittingInstances.owl), fix the problem, and repeat the subsequent steps.

These utilities assume that the format of the input files is the same as either the format used by the *owlPrinter* or the format followed by the kittingInstances.owl file. If the format used by an input file is different from both of those, the utilities may fail.

The *compactOwl* utility compacts an OWL instance file by:

- 1. removing all occurrences of one or two blank lines. The blank lines must not contain spaces or tabs.
- 2. removing comments. The comments must have "//" as the first two characters on the line.
- 3. combining each OWL statement written on two or more lines so it is all on one line. The first non-space character on the second line must be a colon (:) or a double quote (").
- 4. rewriting numbers with decimal points so there are exactly six decimal places. Such numbers must have at least one digit on each side of the decimal point in the input file.
- 5. putting the **DifferentIndividuals** inside **DifferentIndividuals** statements into alphabetical order.

To run compactOwl use a command of the form:

bin/compactOwl < [owl file in] > [owl file out]

where [owl file in] and [owl file out] are replaced by file names.

The *compareOwl* utility compares two files that are expected to have the same lines, but in a different order, such as an automatically generated OWL file and a hand-written OWL file. It reads the two files and saves the lines of each one in two sets in alphabetical order (set::insert puts strings in

alphabetical order by default). Then it compares the lines of the two sets in order. If it finds two lines that do not match, it prints the line from the first file followed by the line from the second file. If all lines match, that is reported.

To run *compareOwl*, use a command of the form:

bin/compareOwl [first owl file in] [second owl file in]
where [first owl file in] and [second owl file in] are replaced by
the names of compacted OWL files.

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