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## Crowdsensing Data Trading for Unknown Market: Privacy, Stability, and Conflicts

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### Key Question

How to design a Crowdsensing Data Trading Framework considering **privacy** and **stability** for **unknown market** in **centralized**<sup>[1]</sup> and **decentralized**<sup>[2]</sup> settings?

### Introduction

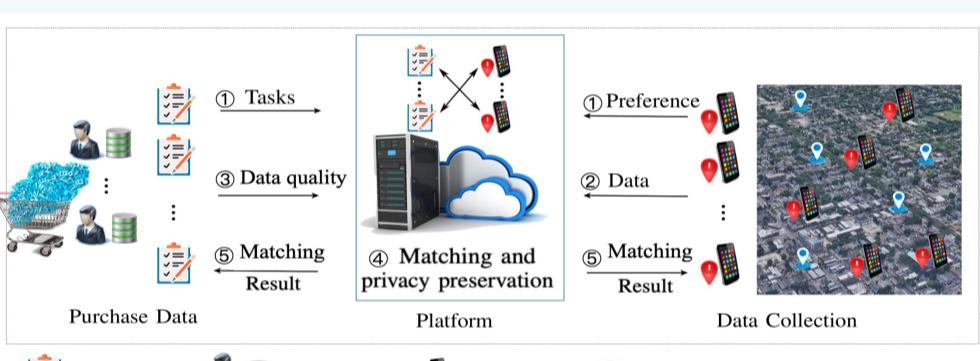
#### ➤ Crowdsensing Data Trading (CDT)

A new data trading paradigm where the Mobile CrowdSensing (MCS) technique is adopted to provide data sources, e.g., Thingful, ThingSpeak.

#### ➤ Concept of Matching Markets

- ✓ Both sides of the markets can't just choose what you want even if you can afford it.
- ✓ One of them also have to be chosen.
- ✓ They choose each other according to the preferences of each other.

#### ➤ Components of CDT systems



PS-CDT platform

**Platform:** As a broker, it provides credible data trading services for sellers and buyers.

**Buyers:** Propose and publish their data requirements to the platform to collect data.

**Sellers:** A crowd of mobile users to provide data collection service to buyers.

#### ➤ Existing Problems

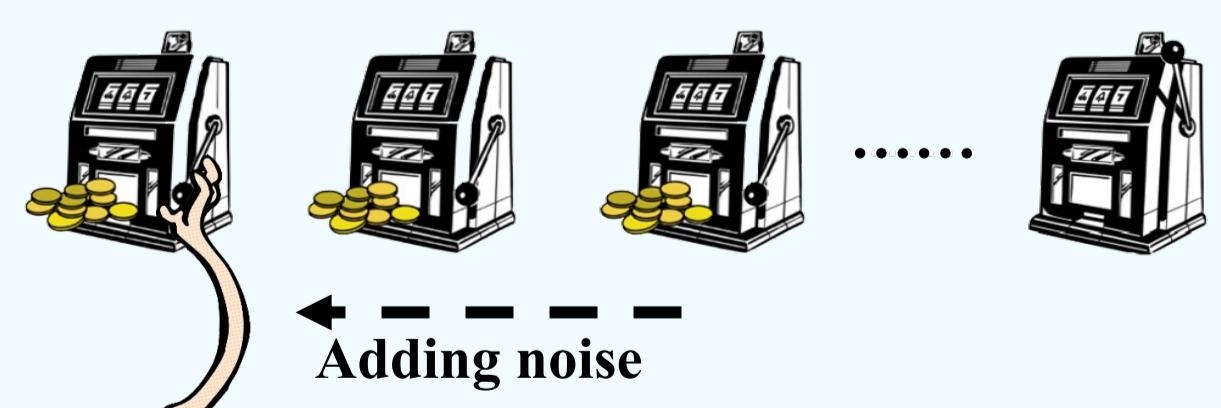
- A few existing CDTs consider the stability of the Data Trading Market.
- The Data Trading Market is unknown in practice, i.e., the preference sequences over sellers are unknown by buyers.
- The private information of sellers needs to be preserved.
- Decentralized CDT has more practical significance.
- Two matching requests for the same seller would create a competitive matching conflict.

#### ➤ Contributions

- ✓ To the best of our knowledge, this is the first CDT work that takes the unknown market, privacy preservation, and the stability of the data trading into consideration simultaneously in centralized and decentralized settings.
- ✓ We define a novel metric, i.e.,  $\delta$ -stability to measure the stability of the markets.
- ✓ We propose the DPS-CB and CDPS-CB mechanisms to solve the privacy, stability, and conflicts-avoiding problems.

### System, Modeling, and Problem

#### ➤ $\epsilon$ -Differentially private bandit model



Platform	Game players
Sellers	Arms
Select a seller	Pull an arm
Data quality	Reward
Protected Data quality	Perturbed reward

$$\checkmark \mathbb{P}\{\Phi_i(q_i^{1:l-1}) \in \mathcal{X}\} \leq e^\epsilon \cdot \mathbb{P}\{\Phi_i(q_i^{1:l-1'}) \in \mathcal{X}\} \quad (1)$$

where  $\epsilon > 0$  is a small constant that the policy provides, indicating the privacy-preserving level.

A bandit policy  $\Phi_i$  of play  $i$  is a sequence of arm-pulling decisions.

$q_i^{1:l} = \{q_i^1, \dots, q_i^l\}$ ,  $q_i^{1:l-1'}$  is its adjacent sequence.

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#### ➤ $\delta$ -Stable Matching Model

##### Definition of preference

*Unknown* preference sequences of the buyer

✓ Denoted by  $\pi_k^l = \{\dots, \pi_i^l, \dots\}$ ,  $\pi_i^l = \{\dots, j, j', \dots\}$ .

✓  $\pi_i^l(j)$  denotes the rank of seller  $j$  in  $\pi_i^l$ .

✓  $v_i = \{\dots, v_i^l(j), \dots\}$  denotes the value. (*Unknown*)

Preference sequence of the seller

✓ Denoted by  $\pi_j = \{\dots, i, i', \dots\}$ .

✓  $\pi_j(i)$  denotes the rank of task  $i$  in  $\pi_j$ .

Adding noise → Matching is **not truly stable**

**Definition of  $\delta$ -stable:** We say a market outcome  $M^l$  is  $\delta$ -stable with a probability less equal than  $1 - \delta$  that a preference sequence is invalid, i.e., there exists two matching pairs  $\langle i, j \rangle$  and  $\langle i, j^* \rangle$ ,  $\forall i \in T, \forall j, j^* \in S$ , satisfies  $\pi_i^l(j) <_i \pi_i^l(j^*)$ ,  $\hat{\pi}_i^l(j) <_i \hat{\pi}_i^l(j^*)$  and  $\hat{v}_i^l(j^*) > \xi_0^l$ , denoted by  $\widehat{M}^k$ .  $\xi_0^l$  is a perturbed care bound and  $\delta$  is a constant less than but close to 1.  $\pi_i^l(j) <_i \pi_i^l(j^*)$ : task  $i$  prefers seller  $j$  to  $j^*$  in  $l^{\text{th}}$  round.

#### ➤ Problem formulation

Our goal is to make the optimal matching in each round according to the built perturbed preference sequences, i.e., to maximize the expected accumulative reward for each task, assuring the  $\epsilon$ -differential privacy and  $\delta$ -stable of market outcomes in each rounds.

$$\text{Maximize: } \sum_l q_i^l(m^l(i))$$

Subject to: Eq. (1) holds

$M^l$  is  $\delta$ -stable

### DPS-CB and CDPS-CB mechanisms

#### Algorithm 1: DPS-CB mechanism

```

Input: the total rounds  $N$ , the preference sequences set  $\{\bar{\pi}_j | \forall j \in \mathcal{S}\}$  of sellers.
Output:  $\{M^l | l = 1, 2, \dots\}$ 
1 for  $l = 1, \dots, N$  do
2   if  $l \leq T$  then
3      $m^l(l) \leftarrow j, \forall j \in \mathcal{S}$ ;
4     Get  $\hat{q}_i^l(j)$  as the corresponding reward according to Eqs. (6-8) while using  $\epsilon$  as the privacy budget under the hybrid differentially private mechanism;
5   else if  $l = T + 1$  then
6     Compute the DP-UCB indexes  $I_i^l(j), \forall i \in \mathcal{T}, \forall j \in \mathcal{S}$  according to Eq. (9);
7     Sort the sellers by the DP-UCB index to build the initial perturbed preference sequence  $\hat{\pi}_i^l$  of each task over sellers;
8     Compute stable matching to get the market outcome  $M^l$  according to  $\{\bar{\pi}_j | \forall j \in \mathcal{S}\}$  and  $\{\hat{\pi}_i^l | \forall i \in \mathcal{T}\}$  using the Gale and Shapley algorithm;
9   else
10    Update  $I_i^l(j), \forall i \in \mathcal{T}, \forall j \in \mathcal{S}$  and  $\{\hat{\pi}_i^l | \forall i \in \mathcal{T}\}$  according to Eqs. (6-9).
11    Compute stable matching to get the market outcome  $M^l$  in the way of Step 8.
12  end
13 end

```

#### Algorithm 2: CDPS-CB mechanism

```

Input: the preference sequences set  $\{\bar{\pi}_j | \forall j \in \mathcal{S}\}$  of sellers, the Bernoulli mean  $p$ 
Output:  $\{M^l | l = 1, 2, \dots\}$ 
1 Initialization:
2  $I_i^l(j) = +\infty, \forall i \in \mathcal{T}, j \in \mathcal{S}$ ;
3 Find a matching that one-to-one maps from  $i$  to  $j$  randomly,  $\forall i \in \mathcal{T}, j \in \mathcal{S}$ ;
4 for  $i = 0, \dots, N$  do
5   for  $j = 1, \dots, T$  do
6     Sample an random value  $B^l(i)$  from  $Ber(p)$ ;
7     if  $B^l(i) = 0$  then
8       Update the set of feasible sellers according to Eqs (10);
9       Task  $i$  selects the seller with the maximum DP-UCB index to match:  $m^l(i) = \max\{I_i^l(j) | \forall j \in F^l(i)\}$ ;
10      end
11    else
12      Task  $i$  matches the same seller as the last round:  $m^l(i) = m^{l-1}(i)$ ;
13    end
14  end
15  if  $i$  wins the conflicts then
16     $M^l \leftarrow (i, m^l(i))$ ;
17    Update  $I_i^l(m^l(i))$  and  $\{\hat{\pi}_i^l | \forall i \in \mathcal{T}\}$ ;
18  end
19 end
20 end
21 end

```

#### ➤ Theoretical Analysis

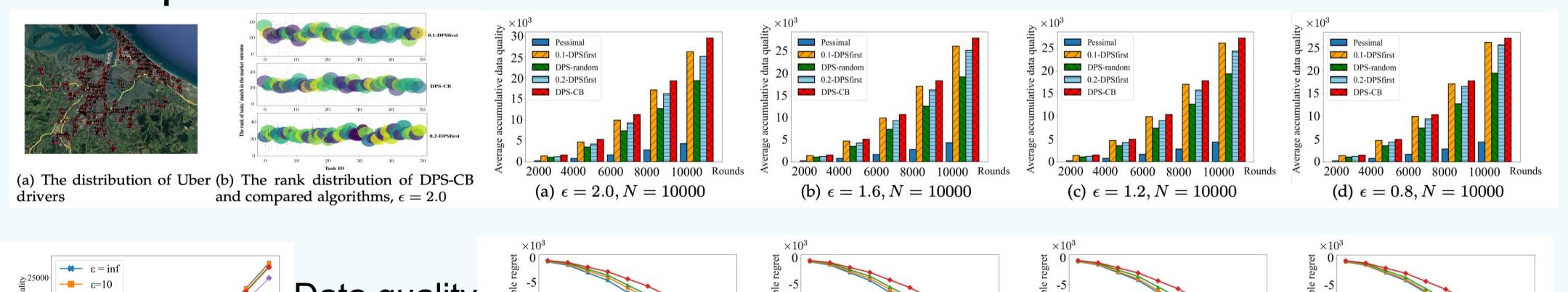
✓ The two mechanisms satisfies  $\epsilon$ -differential privacy.

✓ The market outcome computed by two mechanisms are  $\delta$ -stable.

✓ The two mechanisms can achieve *sublinear* pessimal stable regret

### Performance Analysis

#### ➤ The performance of DPS-CB mechanism.



#### ➤ The performance of CDPS-CB mechanism.

