

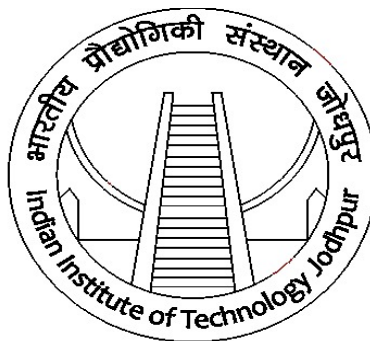
INDIAN INSTITUTE OF TECHNOLOGY- JODHPUR

GRAPH THEORY AND APPLICATIONS(GTA)
COURSE CODE: CSL7410

Lecture Scribing Assignment: Week 7

Done by:
Sachin Kumar
M20MA064

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Week 7

Hall's Theorem*

7.1 Some Theorems and Lemma

In this Section we will learn about the Properties of bipartite graph and some related lemma

7.1.1 Hall's Theorem

Statement : Let G be a Bipartite graph with bipartition (X, Y) . Then G contains a matching that saturates every vertex in X if and only if $|N(S)| \geq |S|$ for all $S \subseteq X$.

Proof: Let us assume that G contains a matching M which saturates every vertex in x .
Let s be a subset of x .

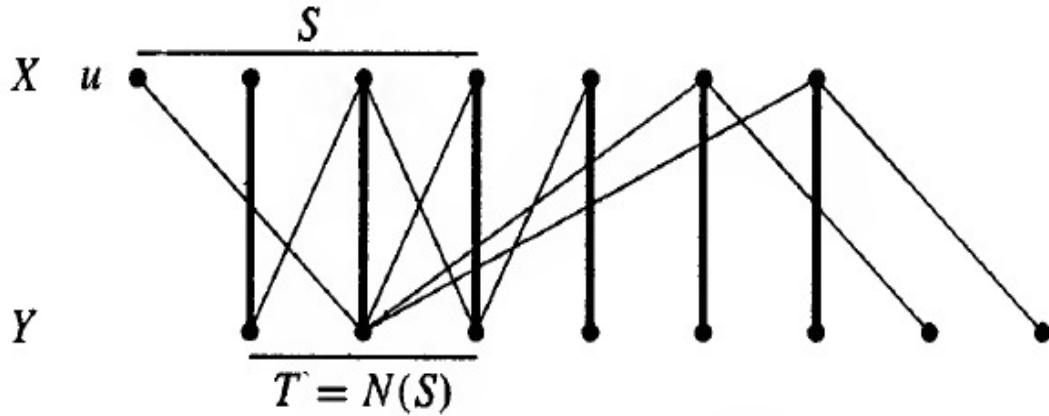
\implies The vertices of S are matched under M with distinct vertices in $N(S)$
 $|N(S)| \geq |S|$ for all $S \subseteq X$

Converse:-

Let us consider a bipartite graph G let us assume that $|N(S)| \geq |S| \quad \forall S \subseteq X$

To prove:- G contains a matching that saturates every vertex in X .

*Lecturer: Anand Mishra. Scribe: Sachin Kumar (M20MA064).



Let us assume the contradiction that G contains no matching saturating all vertices in X .

Let us assume that M^* is a maximum matching in G .

$\implies M^*$ does not saturates all vertices in X .

Let u be an M^* unsaturated vertex in X let Z denotes the set of vertices connected to u by M^* - alternating Path.

By Berge's theorem it follows that u is the only M^* - unsaturated vertex in Z
let the set $S = Z \cap X$ and $T = Z \cap Y$

clearly, the vertices of S -set(u) are matched under M^* with the vertices.

Since, $|T| = |S| - 1$

$N(S) \subseteq T \quad N(S) = T$

$|N(S)| = |T| = |S| - 1$

$\implies |N(S)| < |S|$

$|N(S)| \geq |S|$

Some definitions

Vertex cover : A vertex cover of graph G is a set $\theta \subseteq V(G)$ that contains at least one end point of every edge.

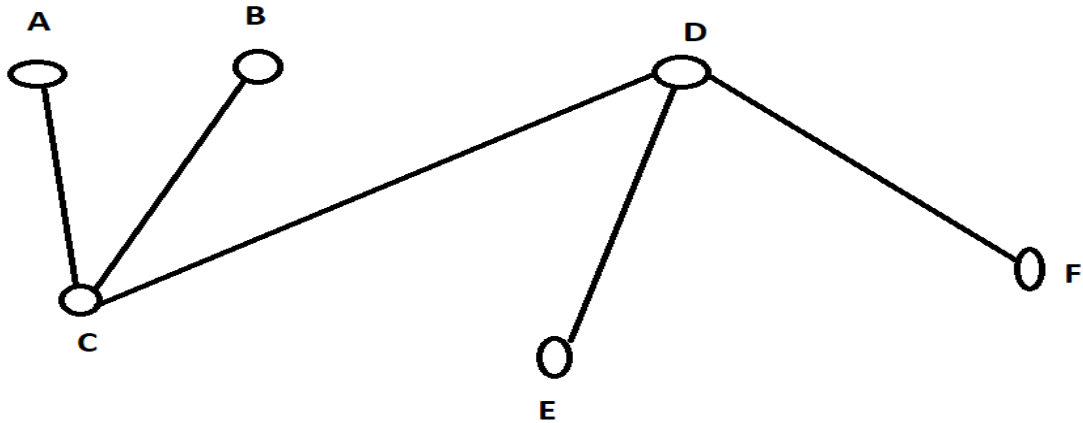
example:

$$\theta_1 = \{C, D\}$$

$$\theta_2 = \{A, B, E, F\}$$

$$\theta_3 = \{D, A, B\}$$

$$\theta_4 = \{C, E, F\}$$



Edge Cover: An edge cover of a graph is set of edge such that every vertex of the graph is incident to the at least one edge of the set.

example:

$$EG = \{AC, BC, DE, DF\}$$

Independent set: set of vertices which. is not adjacent.

example: independent set = $\{A, B\}$ or

$$\text{or } \{E, F\} \text{ of } \{A, B, E, F\}$$

$\alpha(G)$ = maximum size of independent set

$\alpha'(G)$ = maximum size of matching

$\beta(G)$ = minimum size of vertices cover

$\beta'(G)$ = minimum size of edge

$n(G)$ = total number of vertices in graph G

$$\alpha(G) = \{A, B, E, F\}$$

$$\alpha'(G) = \{AC, DE\} \text{ or, } \{BC, DF\}$$

$$\beta(G) = \{C, D\}$$

$$\beta'(G) = \{AB, AC, BC, DE, DF\}$$

7.1.2 Theorem

Statement: if G is a simple graph then $\text{diam}(G) \geq 3$ then $\text{diam}(G^c) \leq 3$

proof: if $\text{diam}(G) \geq 3$ that implies

- (1) $\exists u \ \& \ v$ such that $uv \notin E(G)$
- (2) $u \ \& \ v$ does not have common neighbour.

hence $\text{diam}(G^c) \leq 3$

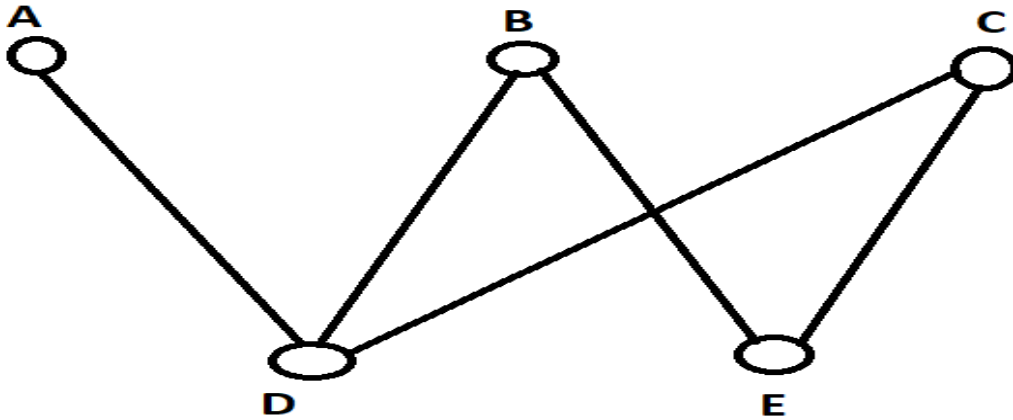
7.1.3 Lemma:

Statement: sum of maximum size of independent set and minimum size of vertex cover is equal to total number of vertices

$$\text{i.e.} \quad \alpha(G) + \beta(G) = n(G)$$

Proof: let s be an independent set of maximum size then every edge is incident to at least one vertex of s

$$S = \{A, B, C\}, \bar{S} = \{D, E\}$$



$$S \cup \bar{S} = V(G)$$

\bar{S} covers all the edges

\bar{S} is minimum size vertex cover $\Rightarrow \beta(G) = |\bar{S}|$ S is maximum size of independent set

$$\Rightarrow S = \alpha(G)$$

$$\therefore \alpha(G) + \beta(G) = |S| + |\bar{S}| = n(G)$$

hence on $\alpha(G) + \beta(G) = n(G)$

7.1.4 theorem

Statement : if G is a graph without isolated vertices, then

$$\alpha'(G) + \beta'(G) = n(G)$$

7.1.5 Theorem

Statement: if G is a bipartite graph with no isolated vertices then $\alpha(G) = \beta'(G)$

Proof: we have already know

$$\alpha(G) + \beta(G) = n(G) \dots\dots\dots(i)$$

$$\alpha'(G) + \beta'(G) = n(G) \dots\dots\dots(ii)$$

$$\alpha'(G) = \beta(G) \dots\dots\dots(iii)$$

from (i) and (ii)

$$\alpha(G) + \beta(G) = \alpha'(G) + \beta'(G)$$

$$\implies \alpha(G) + \beta(G) = \beta(G) + \beta'(G)$$

$$\implies \alpha(G) = \beta'(G)$$

7.1.6 theorem

Statement : G be a bipartite graph then $\alpha(G) = \frac{n(G)}{2}$ iff G has perfect matching .

Proof: Since $\alpha(G) + \beta(G) = n(G)$

$$\begin{aligned} \implies \alpha(G) &= n(G) - \beta(G) \\ &= n(G) - \alpha'(G) \end{aligned}$$

if G has perfect matching then maximum size of matching

$$\alpha'(G) = \frac{n(G)}{2}$$

$$\text{So } \alpha(G) = n(G) - \frac{n(G)}{2}$$

$$\alpha(G) = \frac{n(G)}{2}$$