

Programming Test

Secure Multiparty Computation

- Yao's millionaire problem
 - Alice knows a secret x , Bob knows a secret y
 - Both together want to know whether $x > y$, but none wish to reveal x or y
- Let, $1 \leq x, y \leq 100$, and Bob has public-private keys (PU_b and PR_b)

Alice

1. Compute $C = E(PU_b, M1)$, $M1$ -- a large random number
2. Compute $C1 = C - x$ and sends $C1$ to Bob

8. Check if x -th number in the sequence is congruent to $M1 \bmod p$, if so, $x < y$, otherwise $x > y$
9. Alice tells the conclusion to Bob

Bob

3. Bob computes $M2_i = D(PR_b, C1+i)$, for $1 \leq i \leq 100$
4. Choose a large prime $p (< M1)$; Bob can know the size of $M1$
5. Compute $Z_i = M2_i \bmod p$, $1 \leq i < 100$
6. Verify if $|Z_i - Z_j| \geq 2$ for all (i,j) and $0 < Z_i < p-1$, for all i , otherwise try another prime and repeat from step-4
7. Send to Alice the sequence: $Z1, Z2, \dots, Z_y, Z_{y+1} + 1, Z_{y+2} + 1, \dots, Z_{100} + 1, p$

10. Receive conclusion

Secure Multiparty Computation: example

Global: $1 \leq x, y \leq 4$

Alice ($x = 4$)

1. Compute $C = E(\text{PUB}, M1) = 19, M1 = 39$
2. Compute $C1 = C - x = 15$

8. 4-th number (9) is not congruent to 39 (mod 31),
hence $x > y$
9. Send result of $x > y$ to Bob

Bob ($y = 2$): RSA with parameters $\text{PUB} = 7$ and $\text{PRb} = 23, n = 11$

3. $M2_1 = D(\text{PRb}, 15+1) = 16; M2_1 = D(\text{PRb}, 15+2) = 18; M2_1 = D(\text{PRb}, 15+3) = 2; M2_1 = D(\text{PRb}, 15+4) = 39$
4. $P=31;$
5. $Z_1 = 26 \bmod 31 = 26$; similarly $Z_2 = 18, Z_3 = 2, Z_4 = 8$
6. Verify all the conditions
7. Send: (26, 18, 2+1, 8+1, 31

10. Receive result of $x > y$