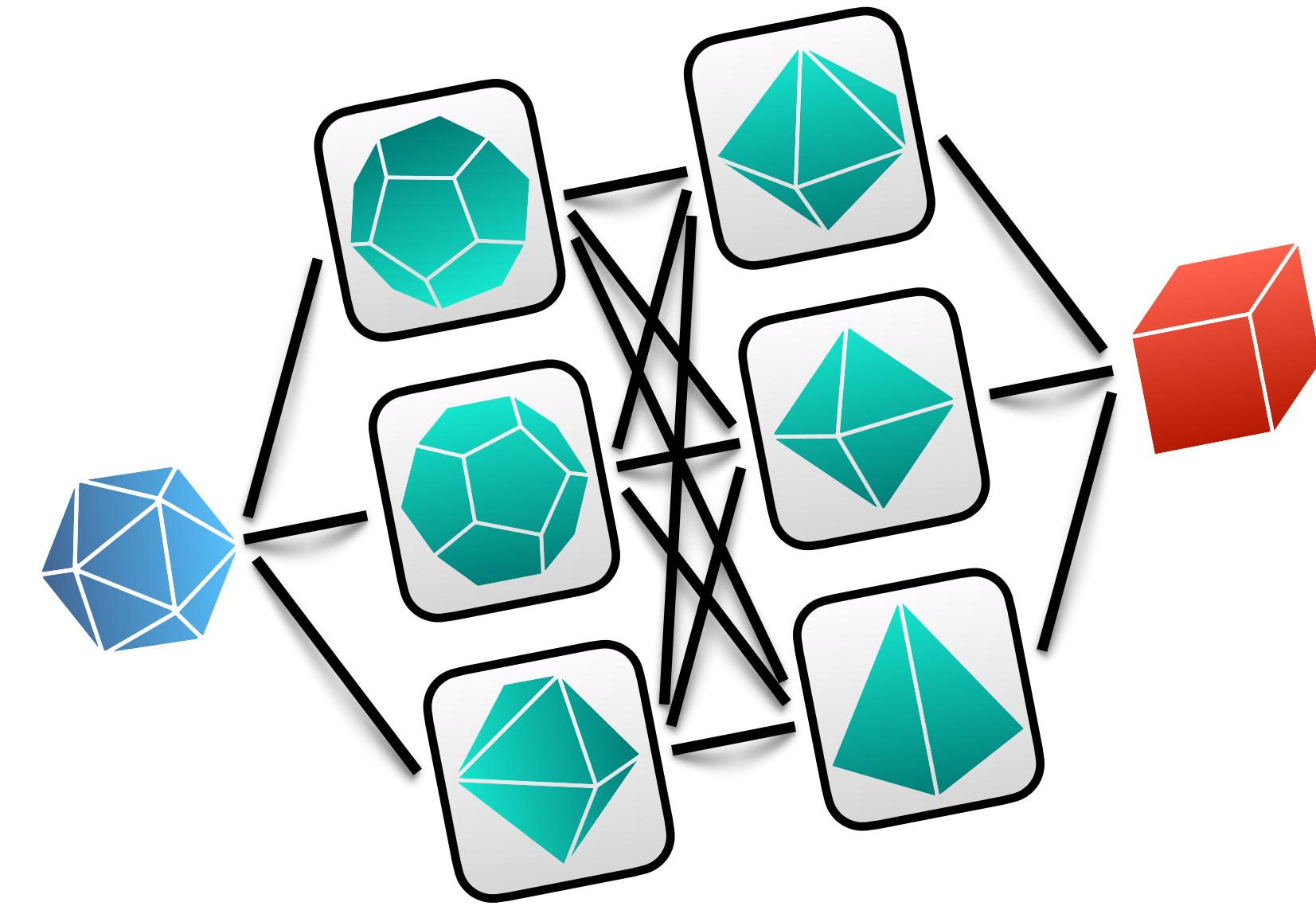




Did I forget to hit  
record? Please  
remind me!



# Algorithm Configuration and friends

**MIE1666: Machine Learning for Mathematical Optimization**

Largely based on Hoos, Holger H. "Automated algorithm configuration and parameter tuning." Autonomous search. Springer, Berlin, Heidelberg, 2011. 37-71.

Some examples from Chapter 13 of Integer Programming by Wolsey

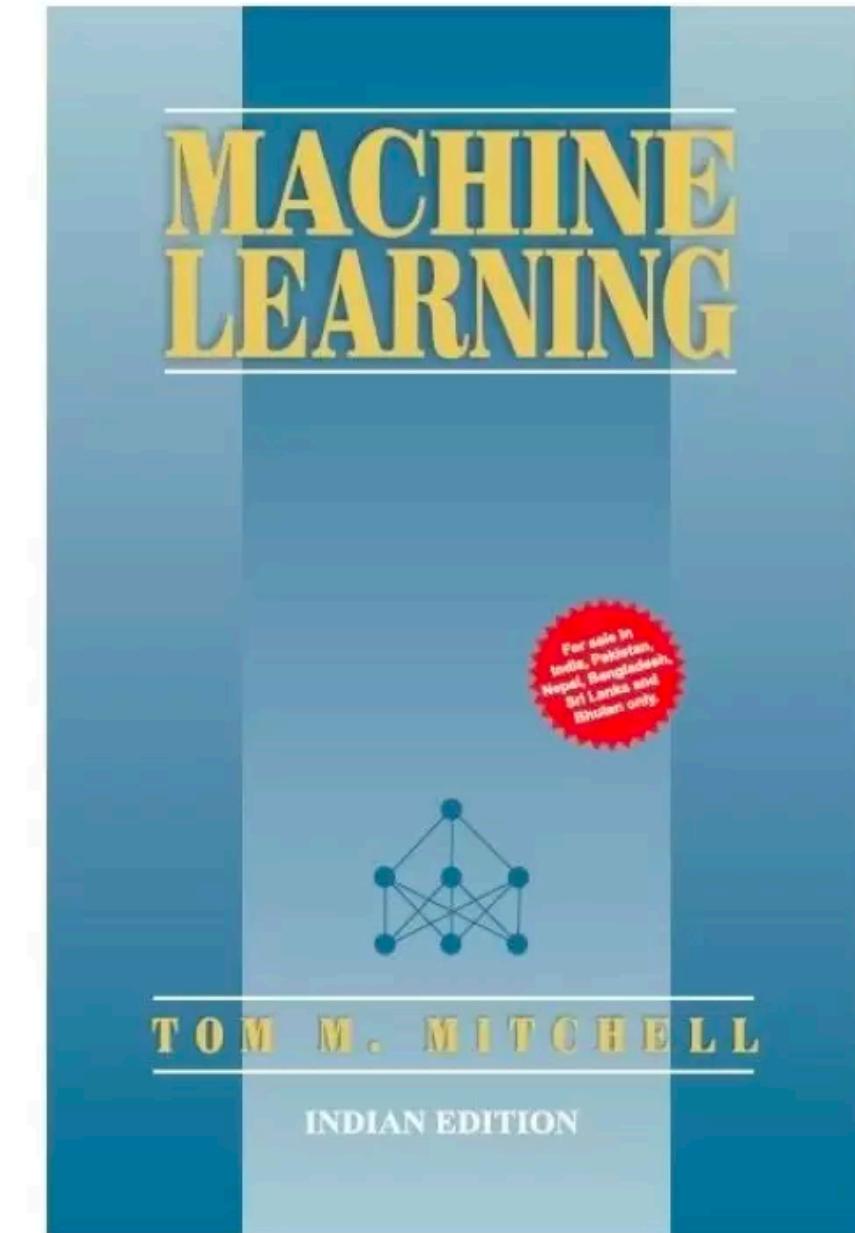
Elias B. Khalil – 04/10/21

# Machine Learning:

Study of algorithms that

- improve their performance P
- at some task T
- with experience E

well-defined learning task:  $\langle P, T, E \rangle$



# Why should this work at all?

The main theoretical basis of supervised learning:

With a **sufficient amount of “similar” data**

+

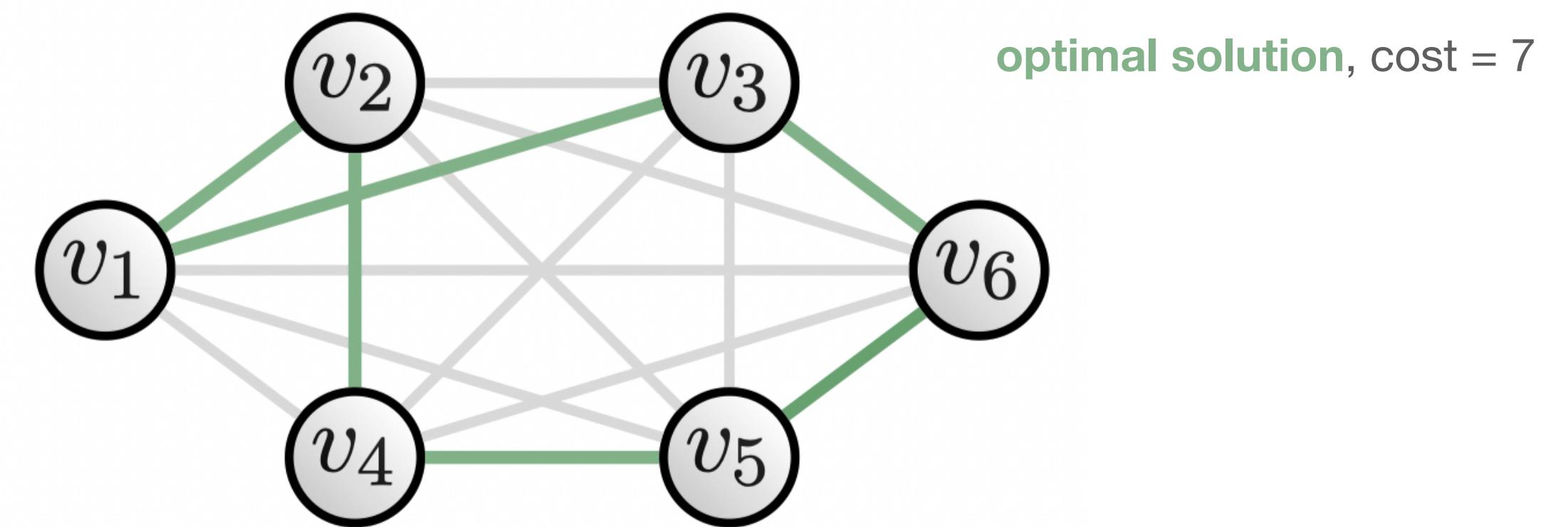
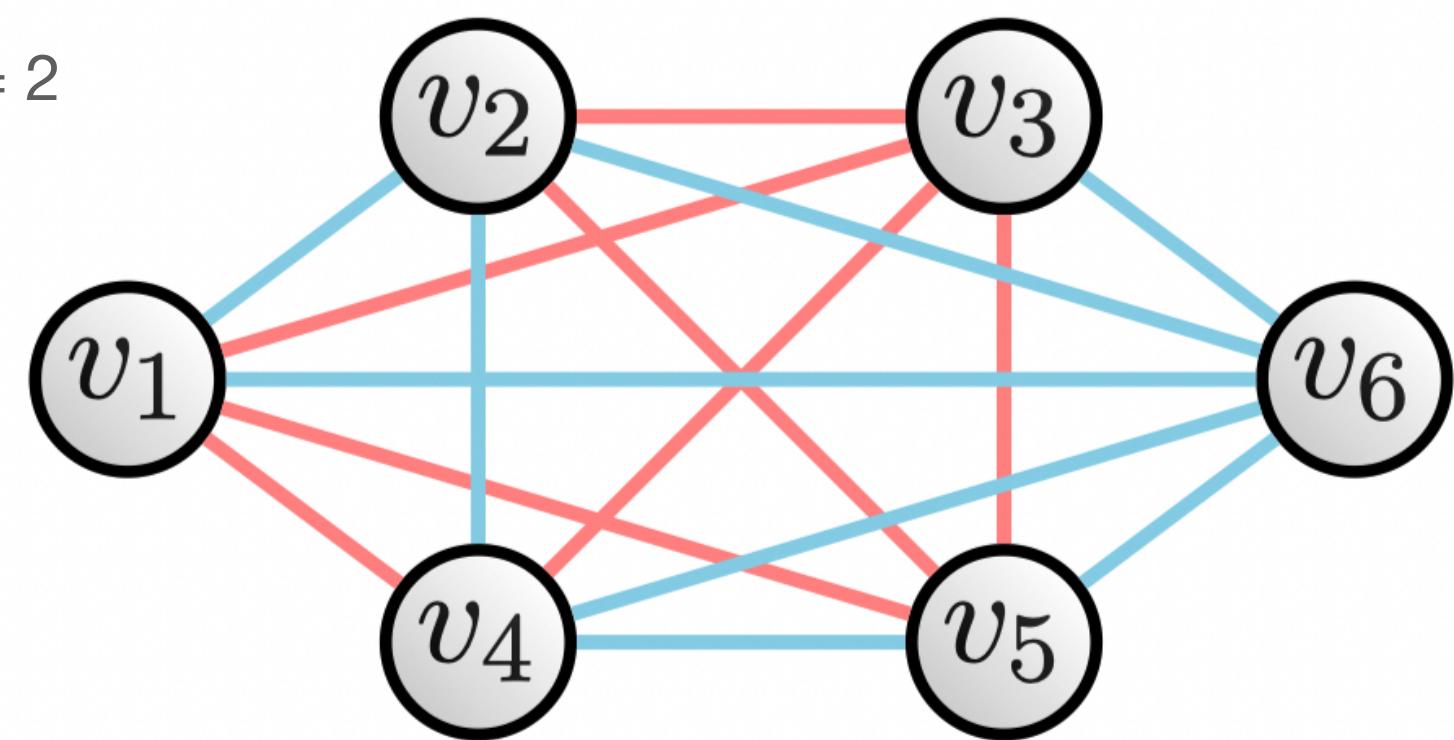
an **expressive model class**:

Minimizing the loss function on the training data yields a highly **accurate model on unseen test data**, with high probability

1. **Data:**  $S = \{(x_i, y_i)\}_{i=1,\dots,n}$ 
  - $x_i$ : data example with  $d$  attributes
  - $y_i$ : label of example (what you care about)
2. **Classification model:** a function  $f_{(a,b,c,\dots)}$ 
  - Maps from  $X$  to  $Y$
  - $(a,b,c,\dots)$  are the **parameters**
3. **Loss function:**  $L(y, f(x))$ 
  - Penalizes the model's mistakes

# Graph Optimization

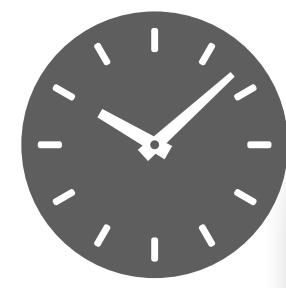
cost = 1  
cost = 2



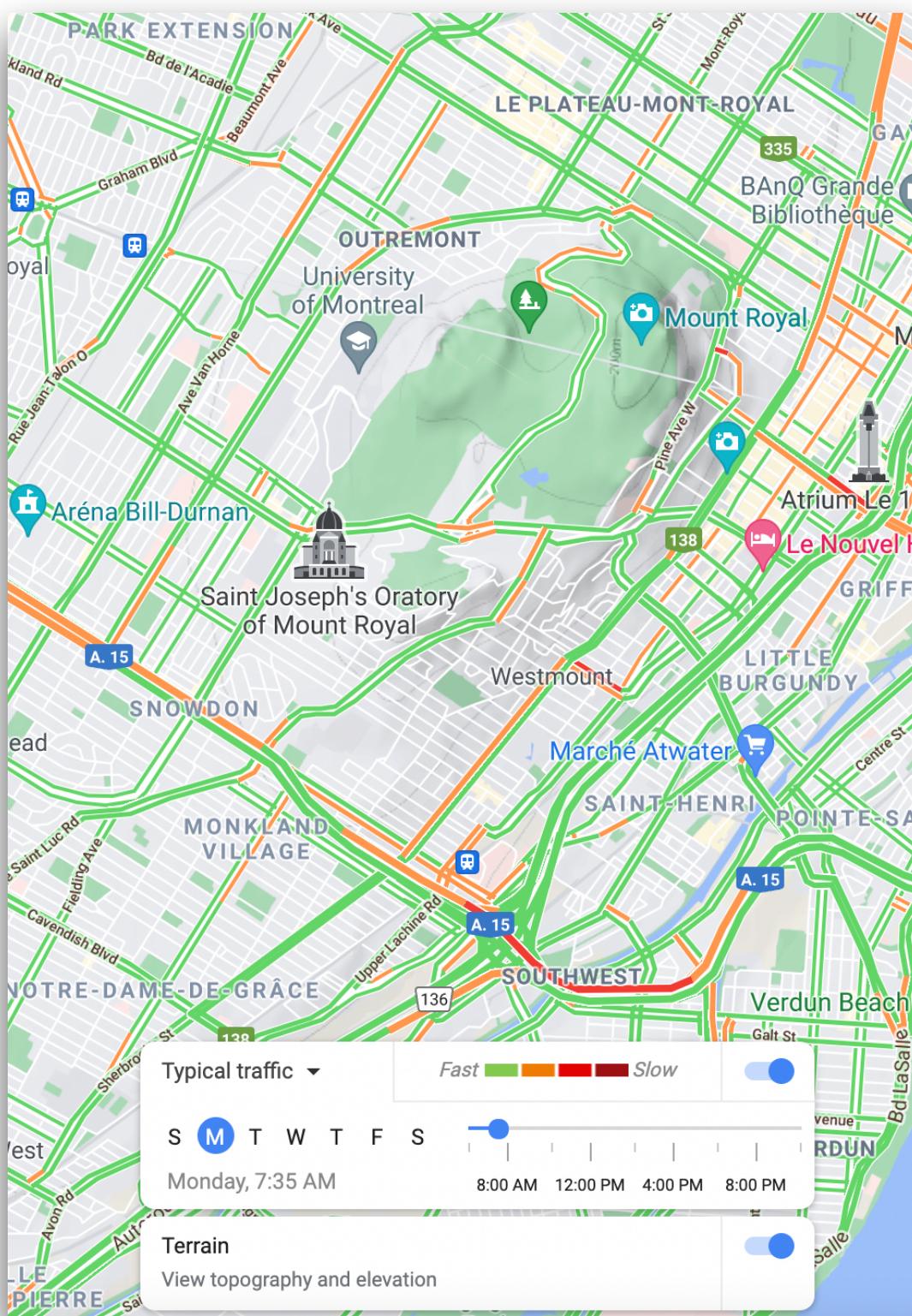
Travelling Salesperson Problem (TSP)

# TSPs in Montréal

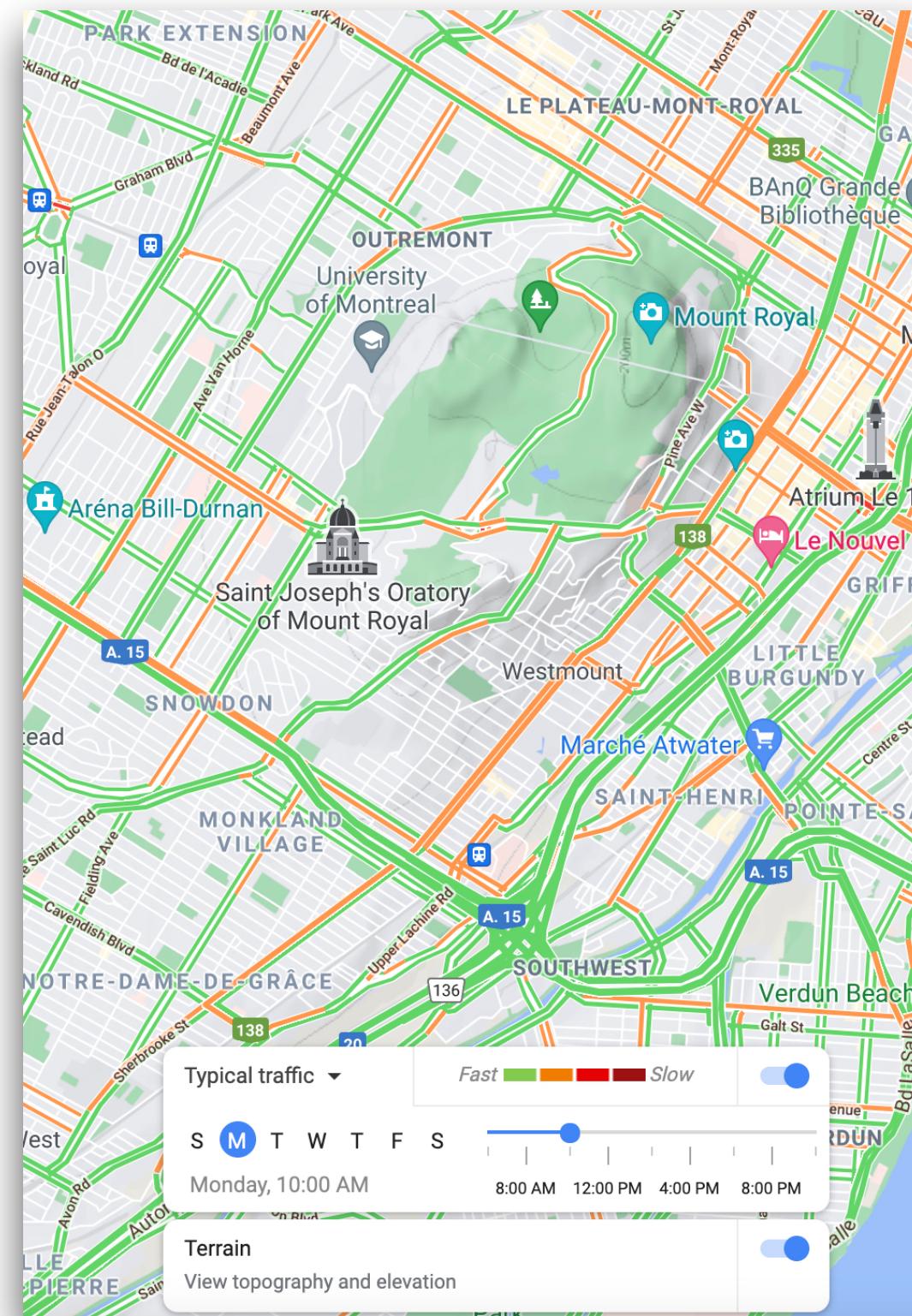
Every few hours, optimally **route** the salesperson through a set of locations



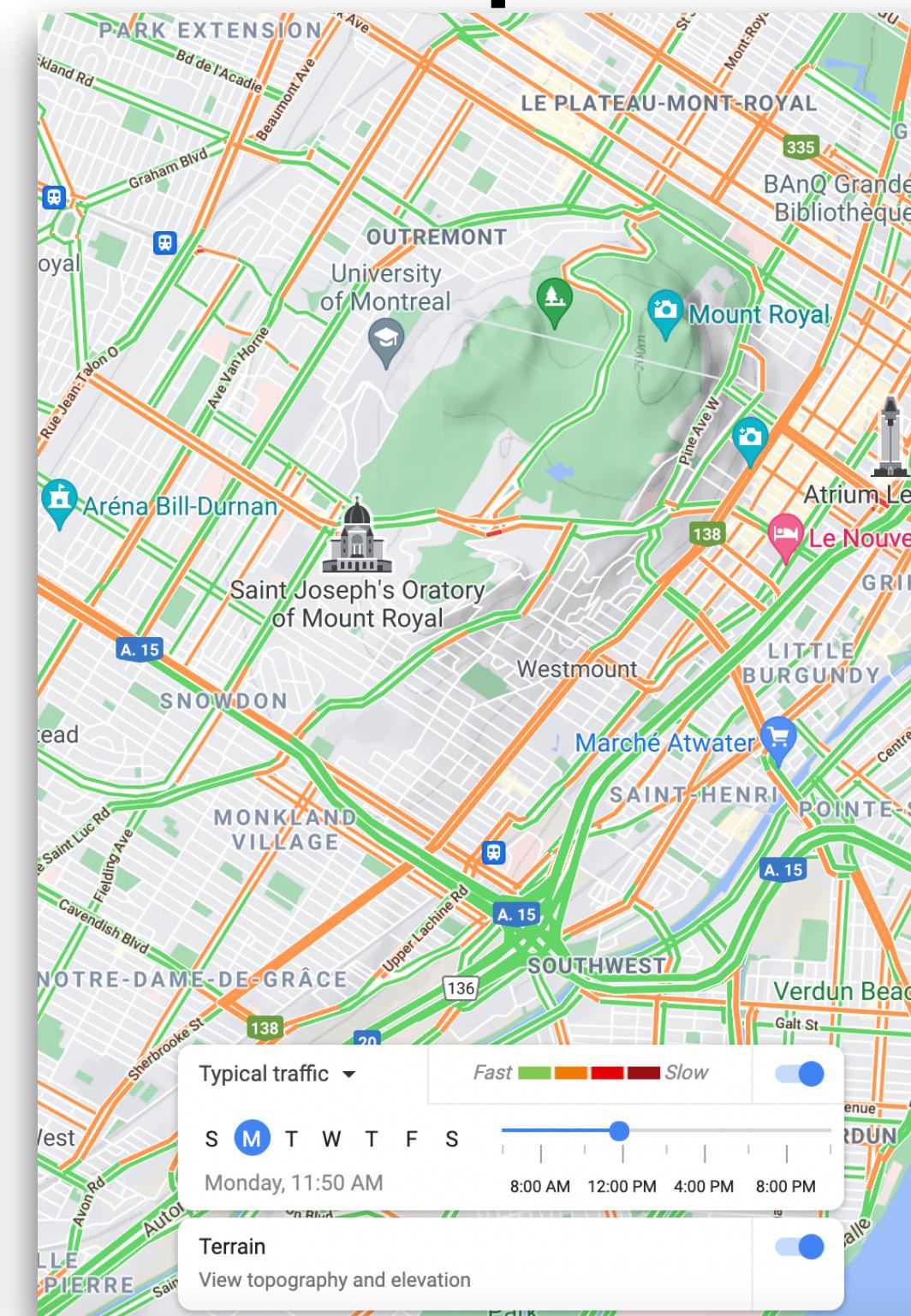
8 am



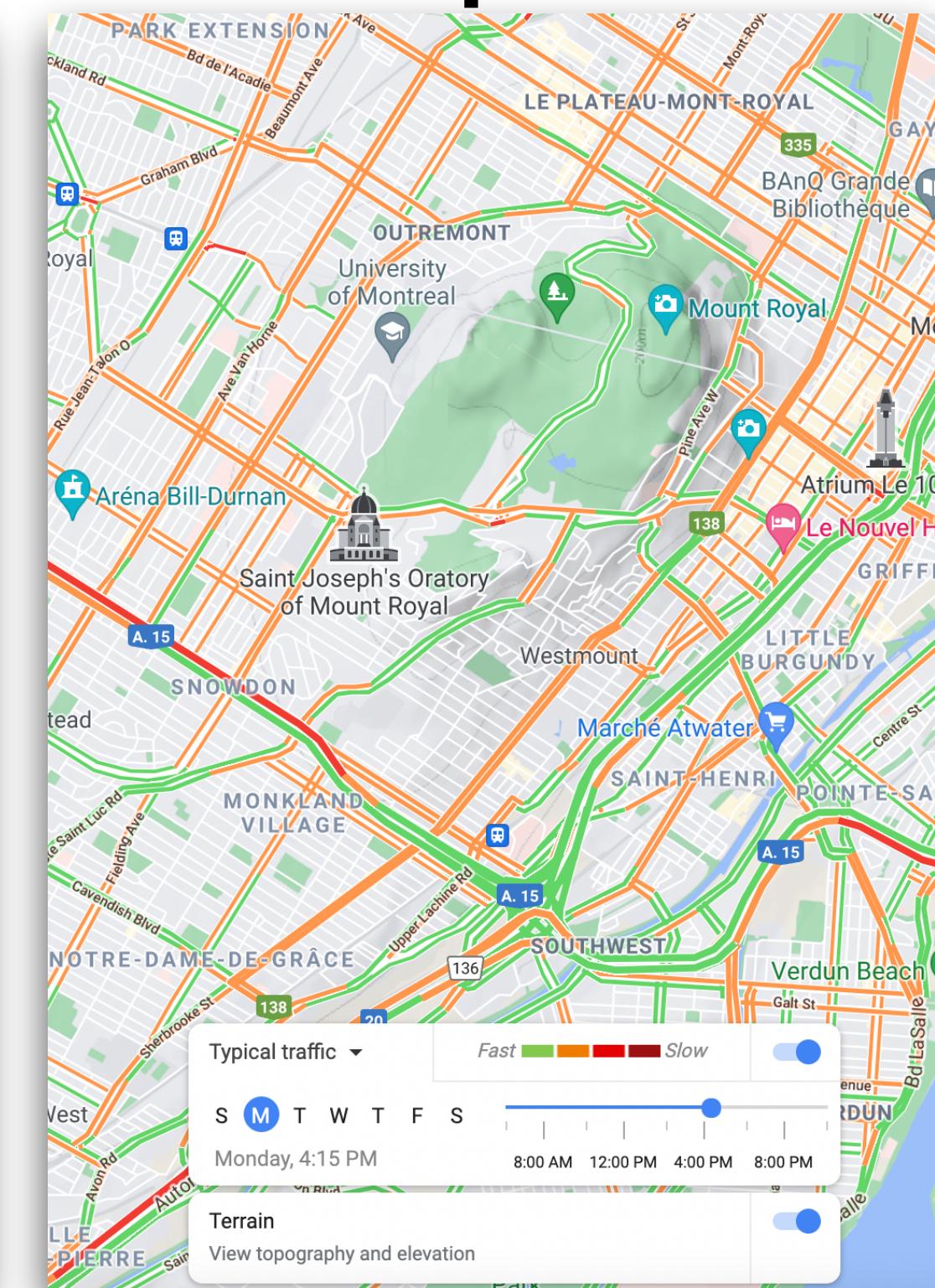
10 am



12 pm



4 pm



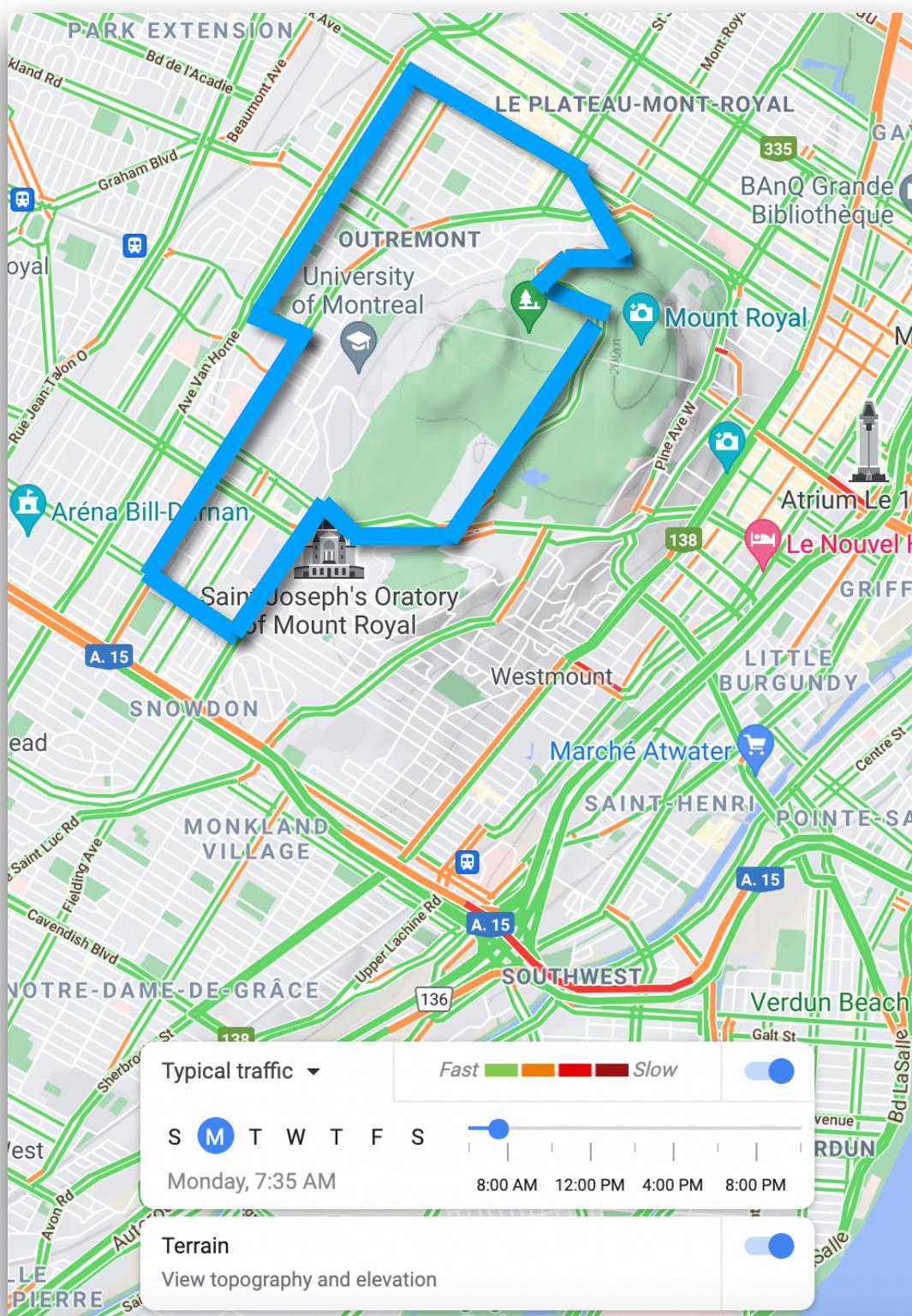
Google Maps, Montréal, Québec, Canada

# TSPs in Montréal

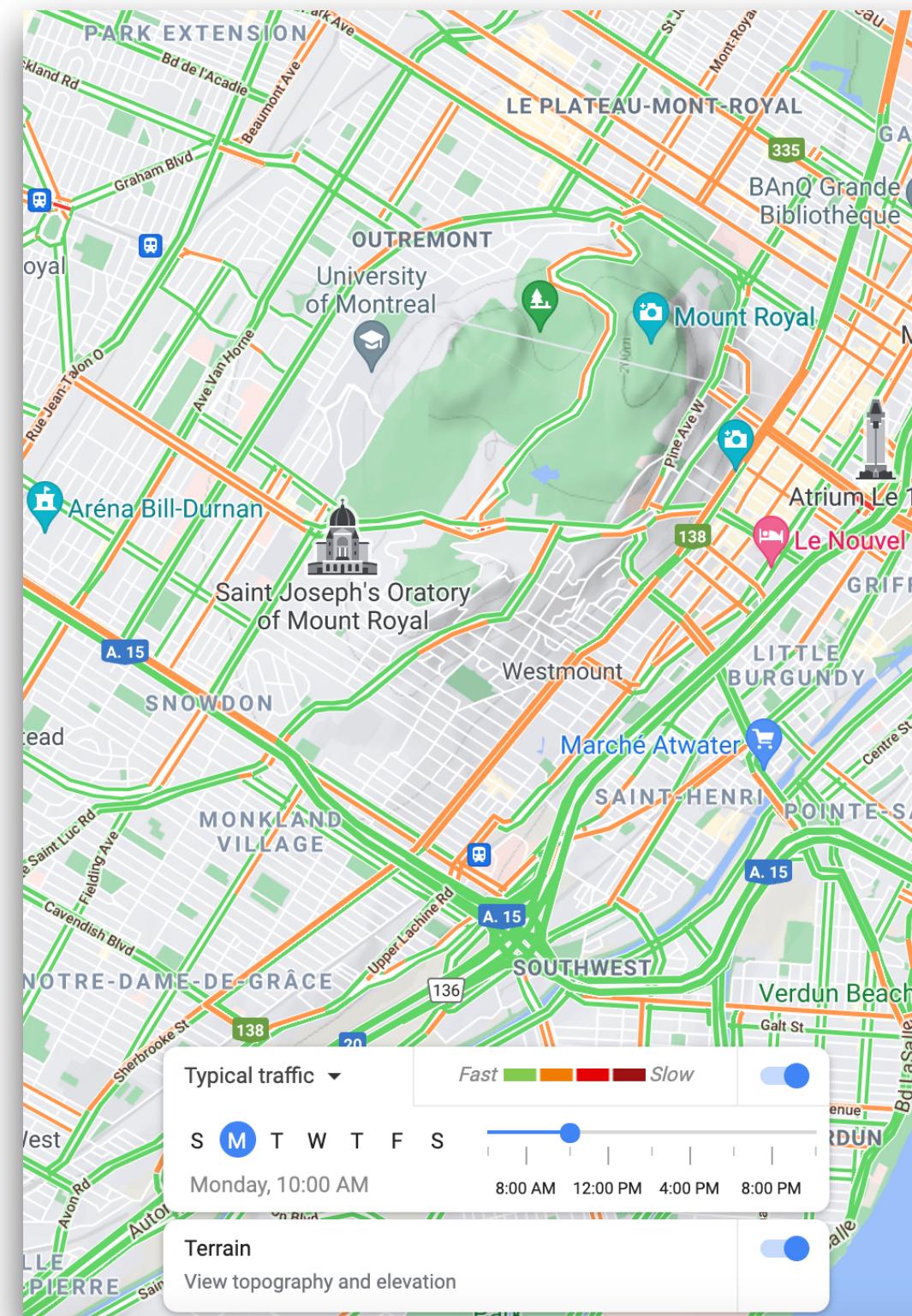
Every few hours, optimally **route** the salesperson through a set of locations



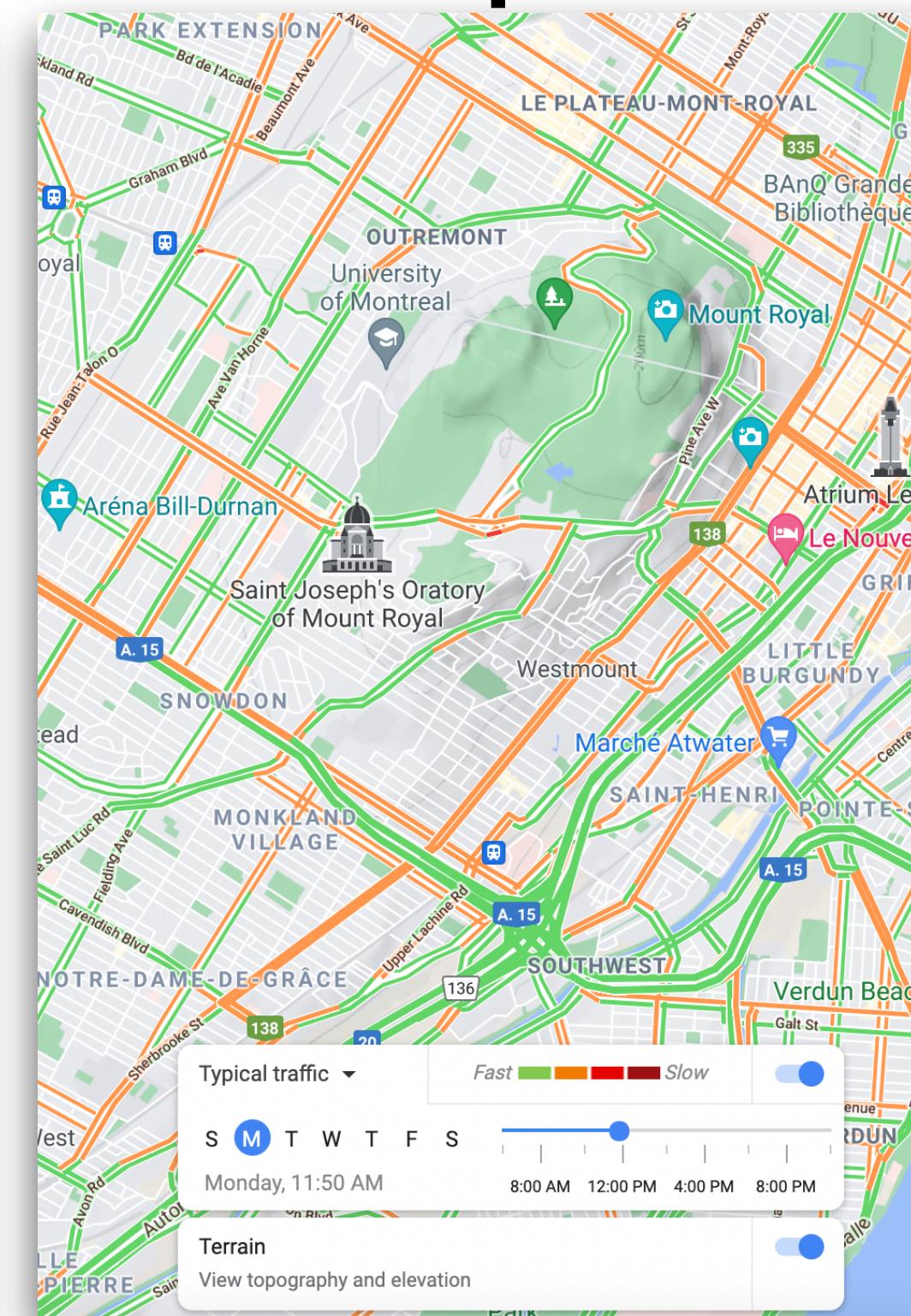
8 am



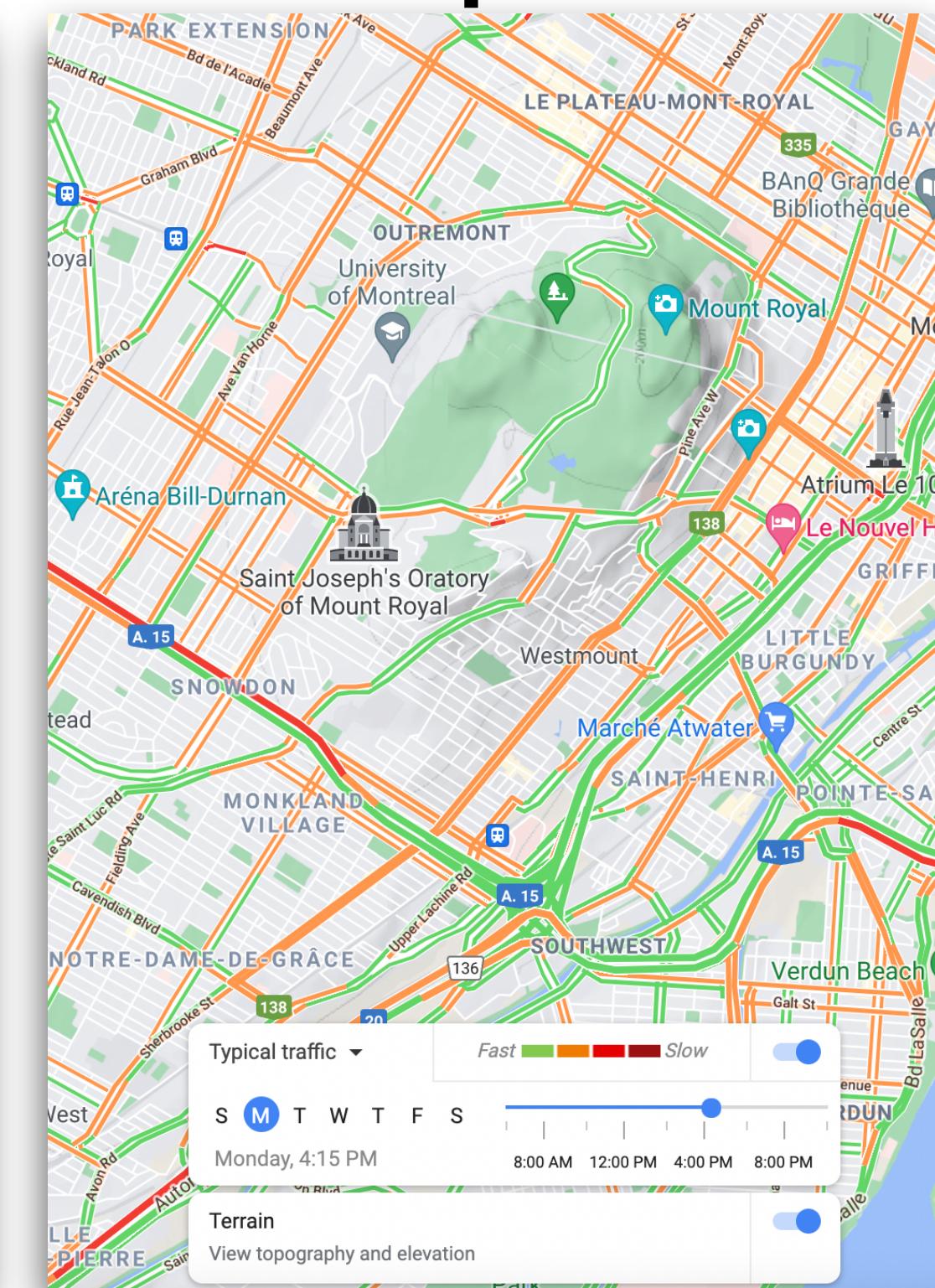
10 am



12 pm



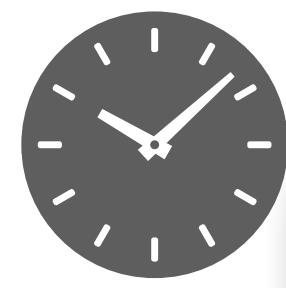
4 pm



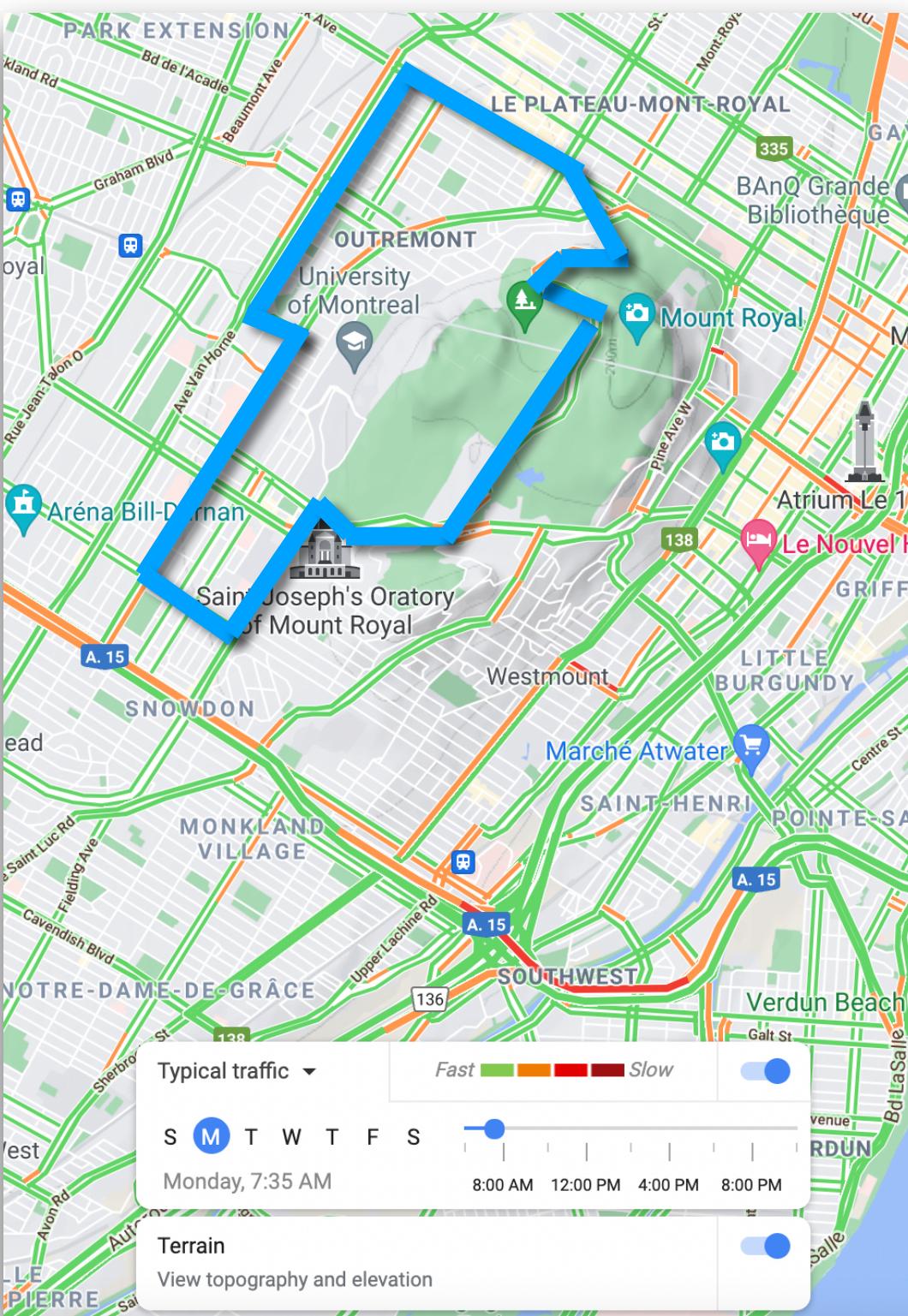
Google Maps, Montréal, Québec, Canada

# TSPs in Montréal

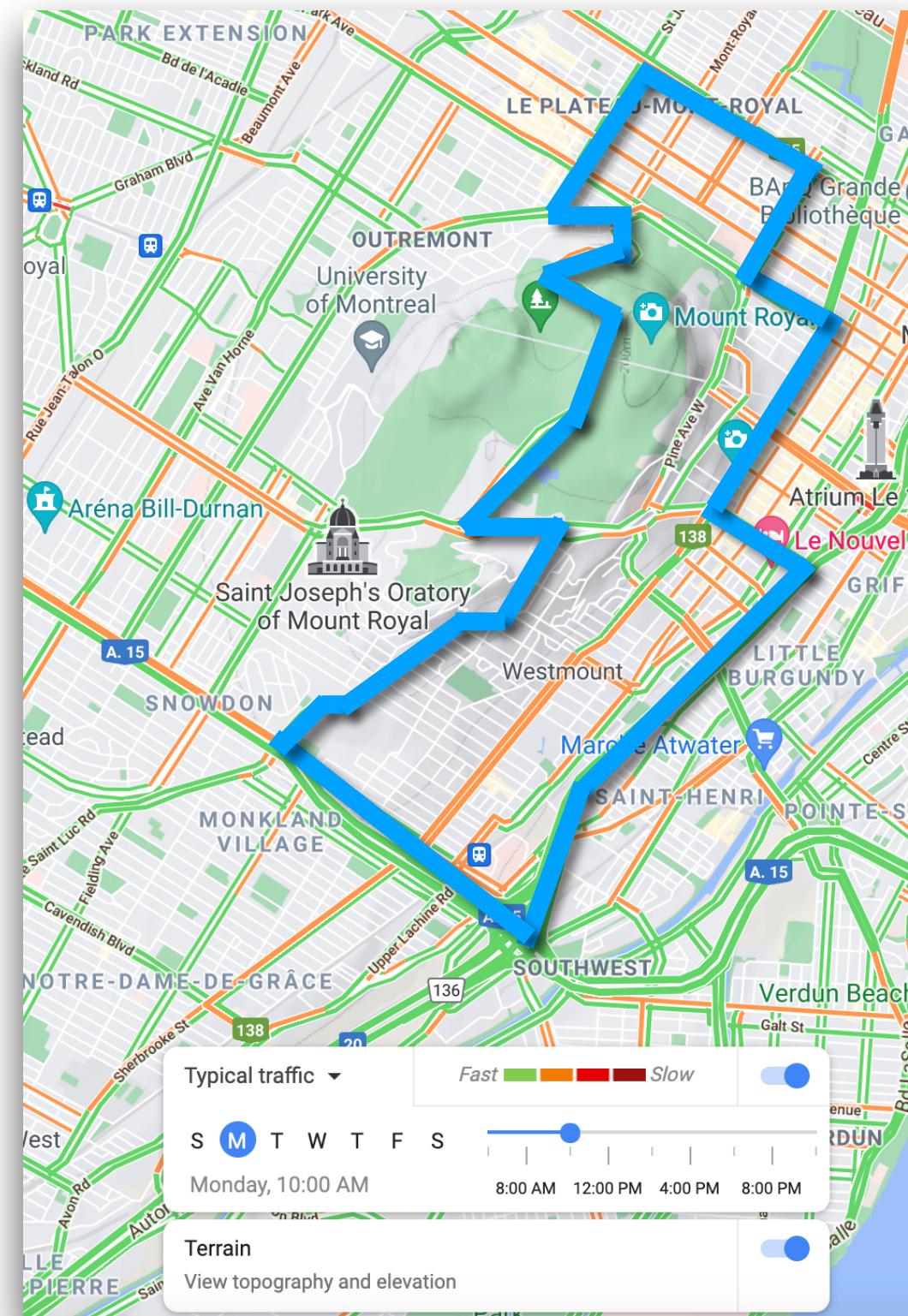
Every few hours, optimally **route** the salesperson through a set of locations



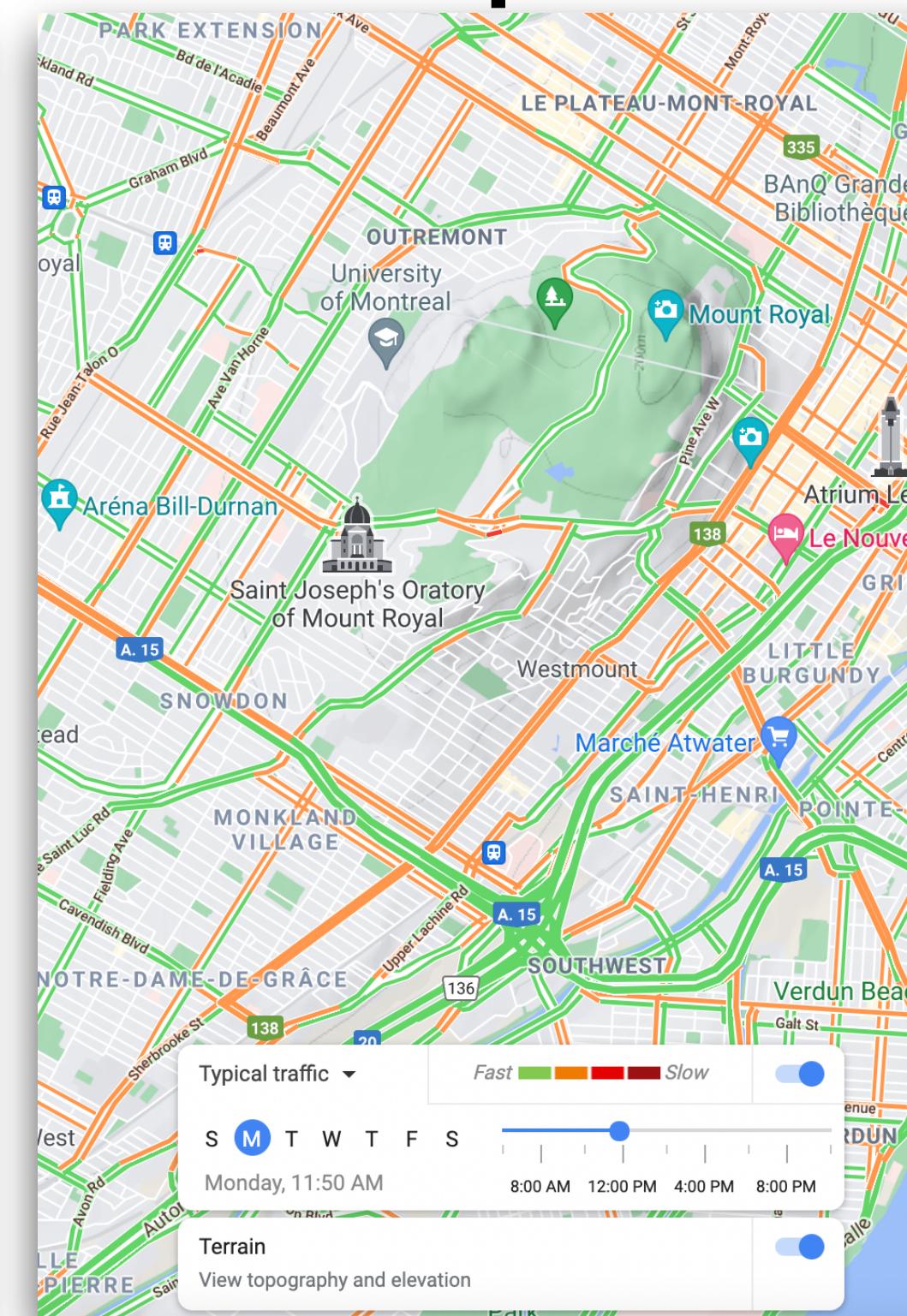
8 am



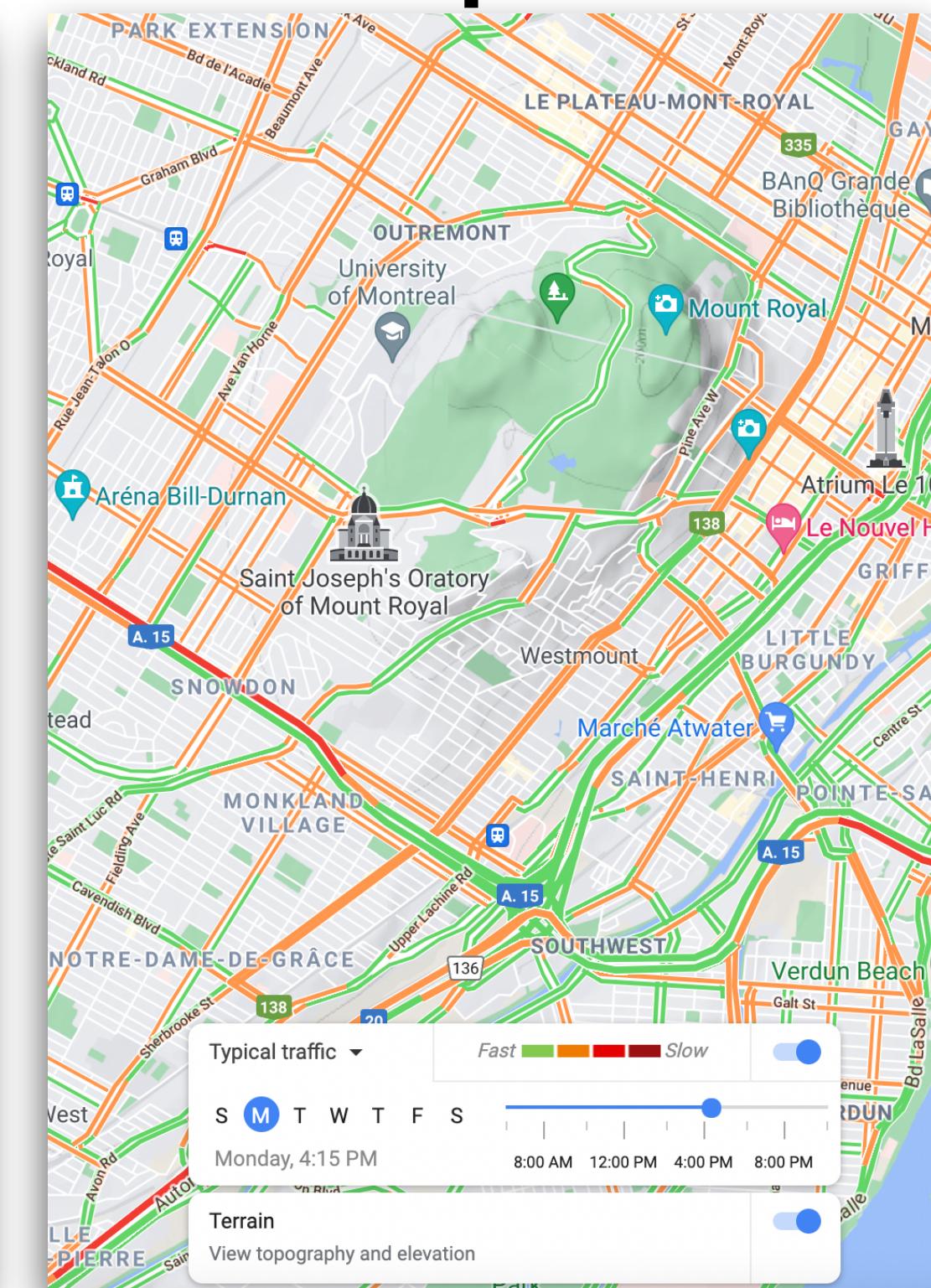
10 am



12 pm



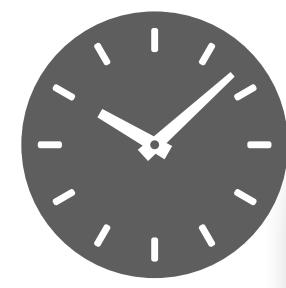
4 pm



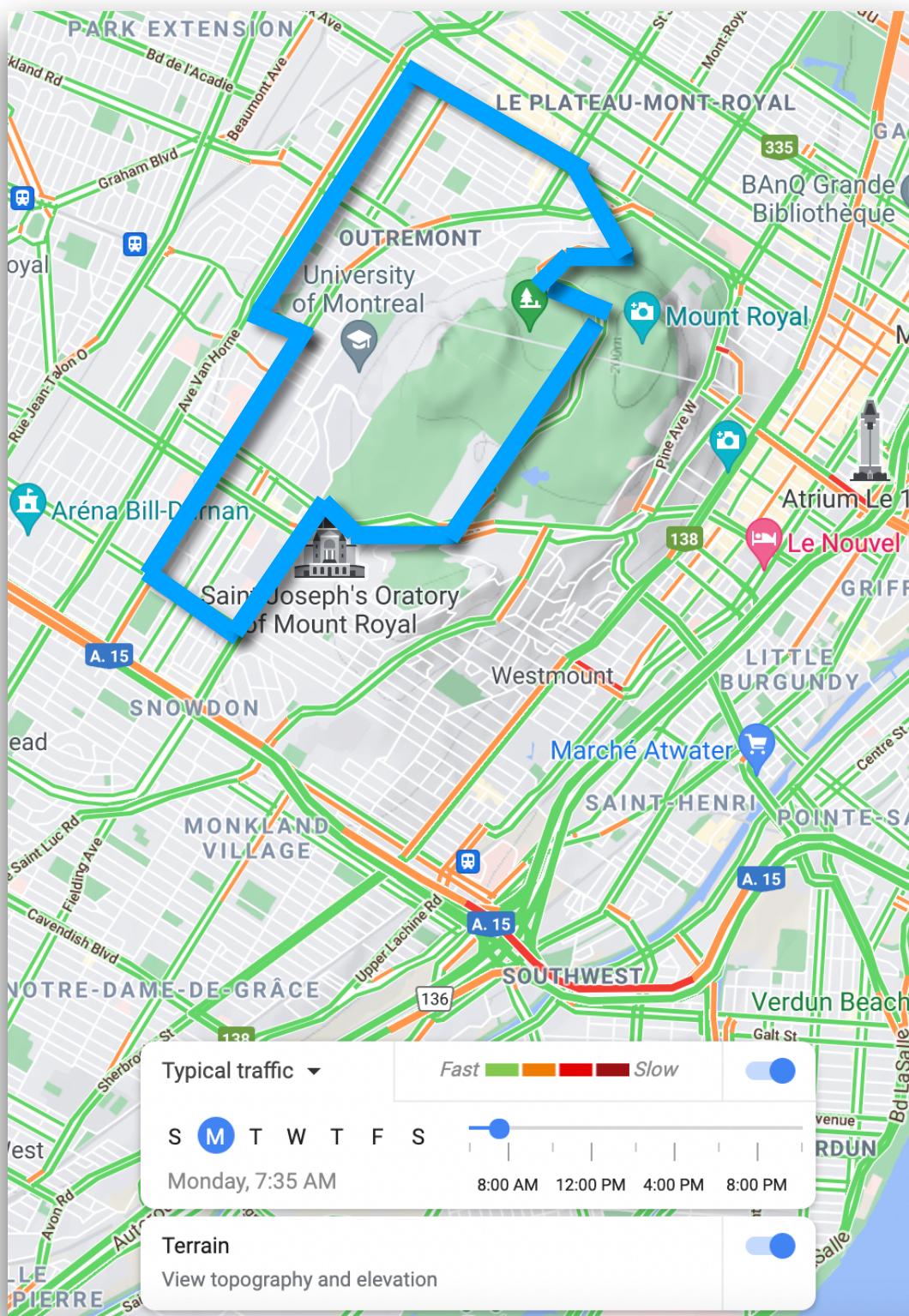
Google Maps, Montréal, Québec, Canada

# TSPs in Montréal

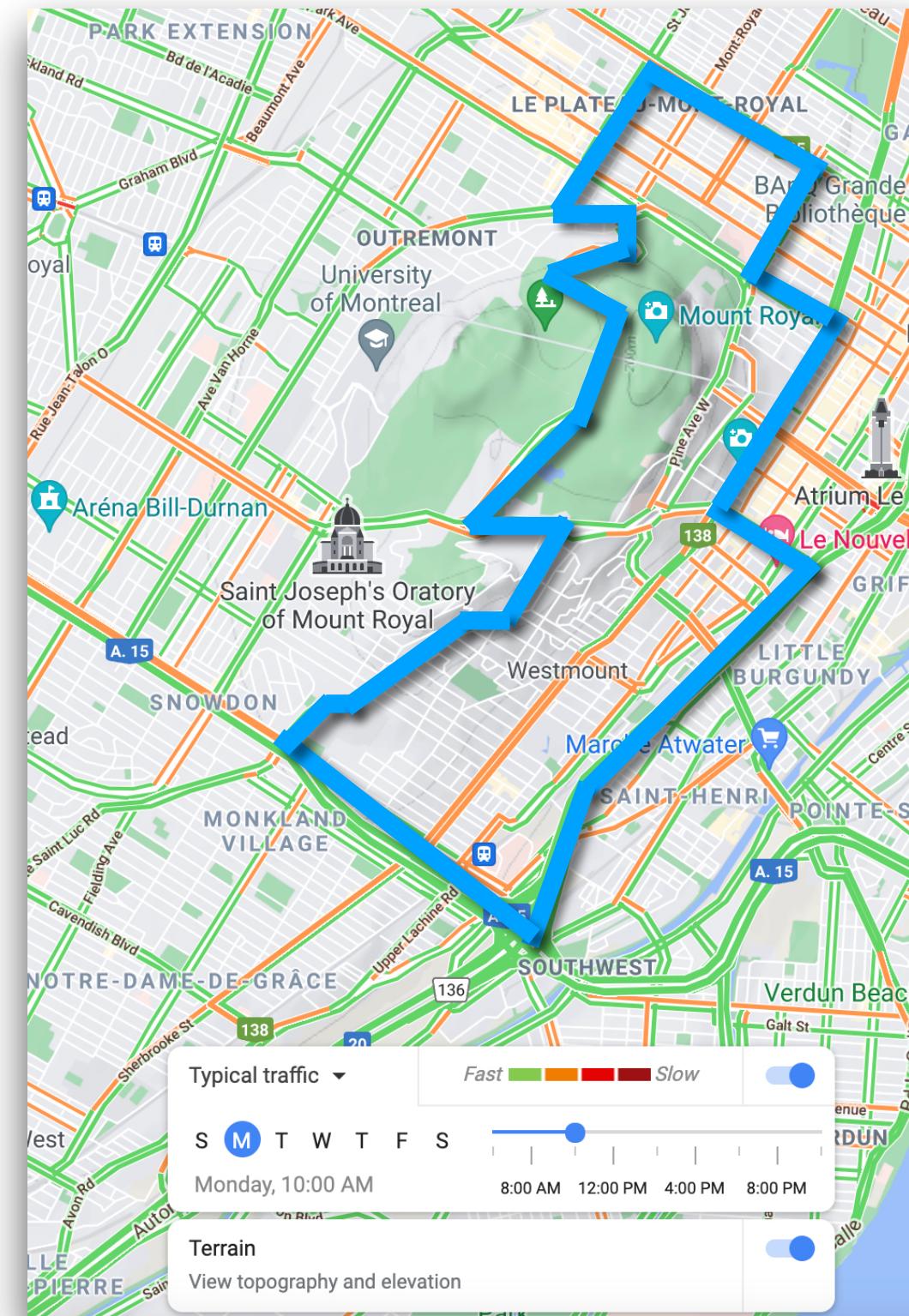
Every few hours, optimally **route** the salesperson through a set of locations



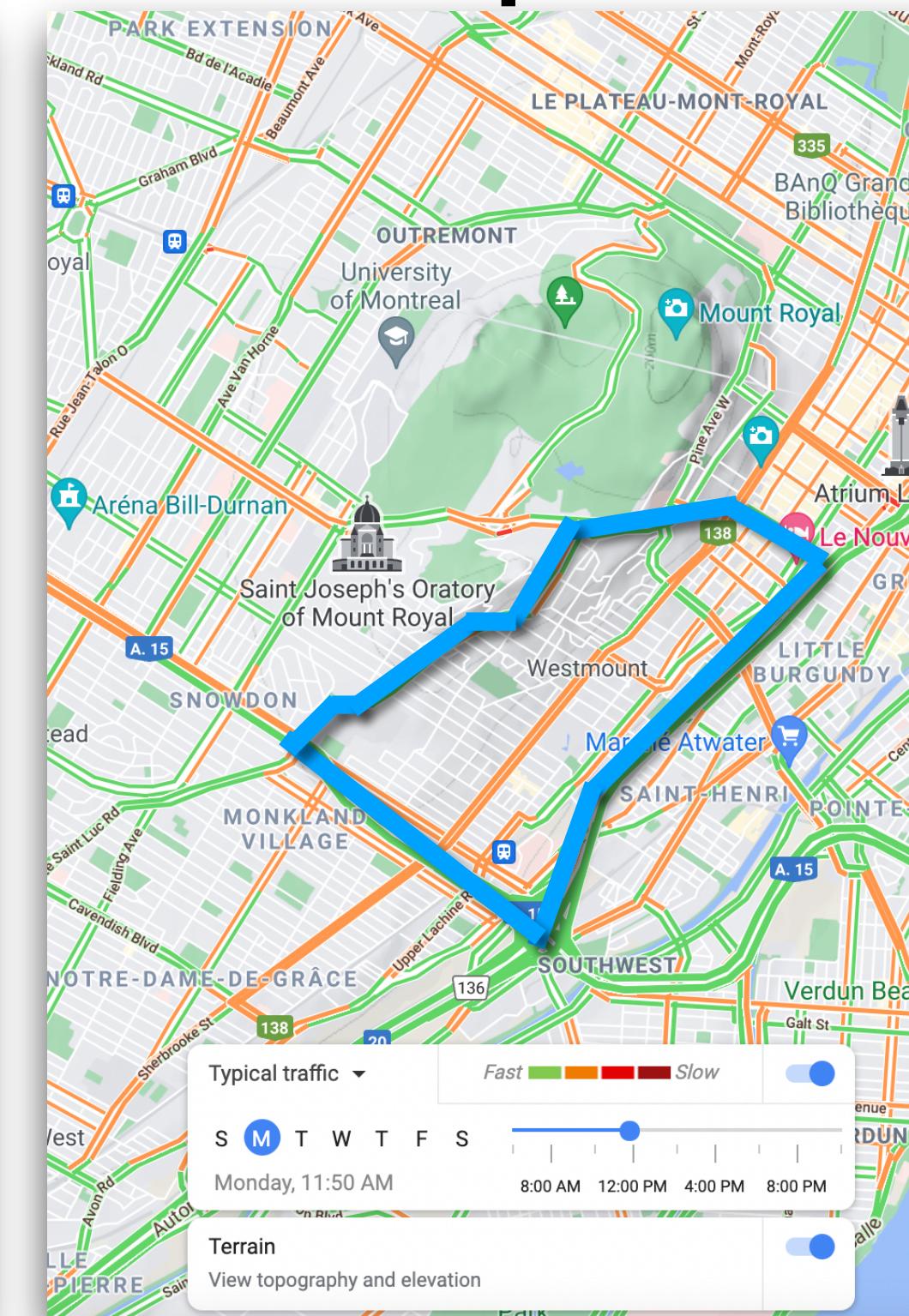
8 am



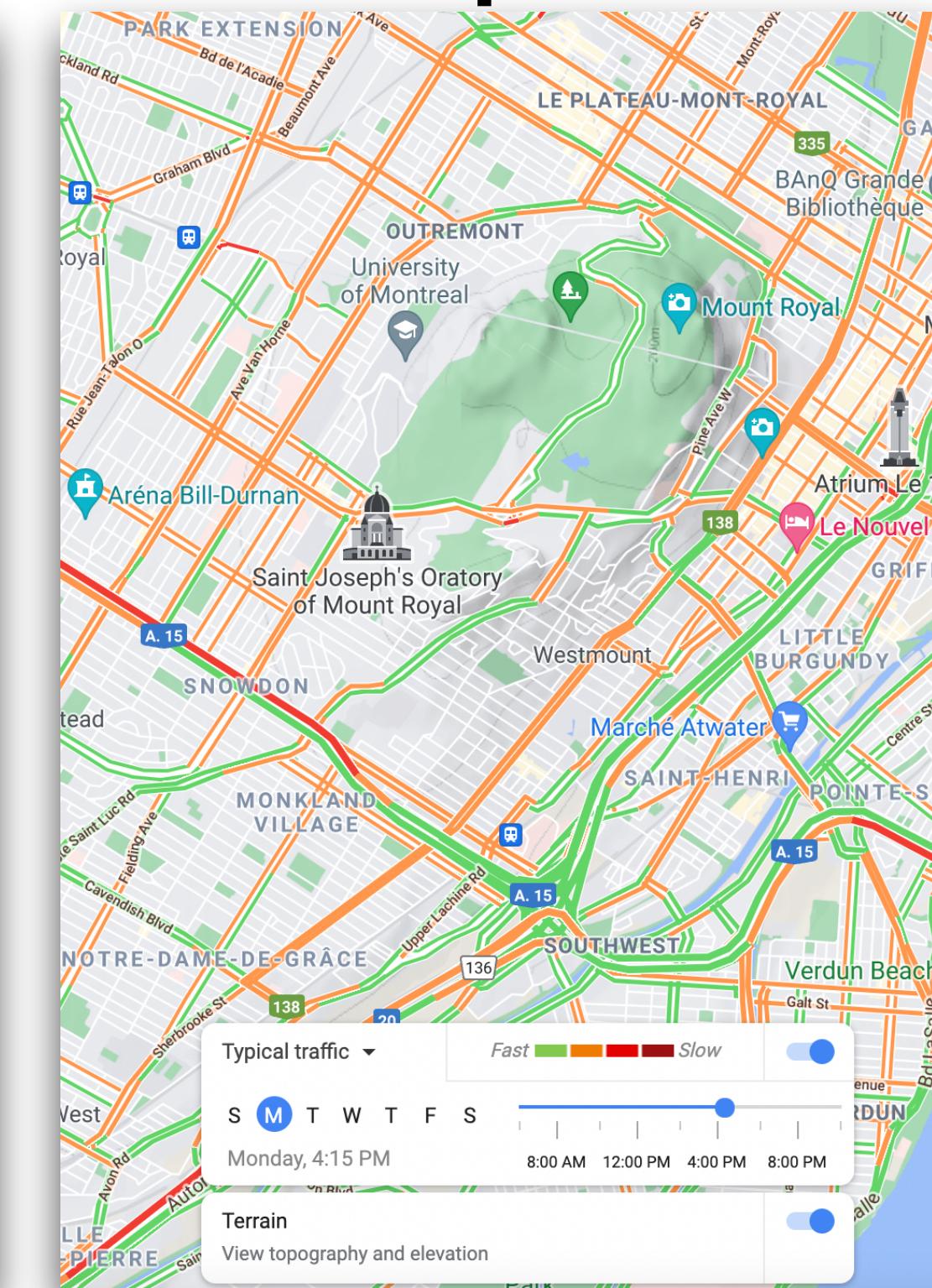
10 am



12 pm



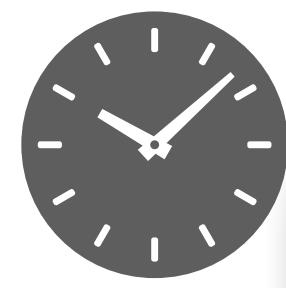
4 pm



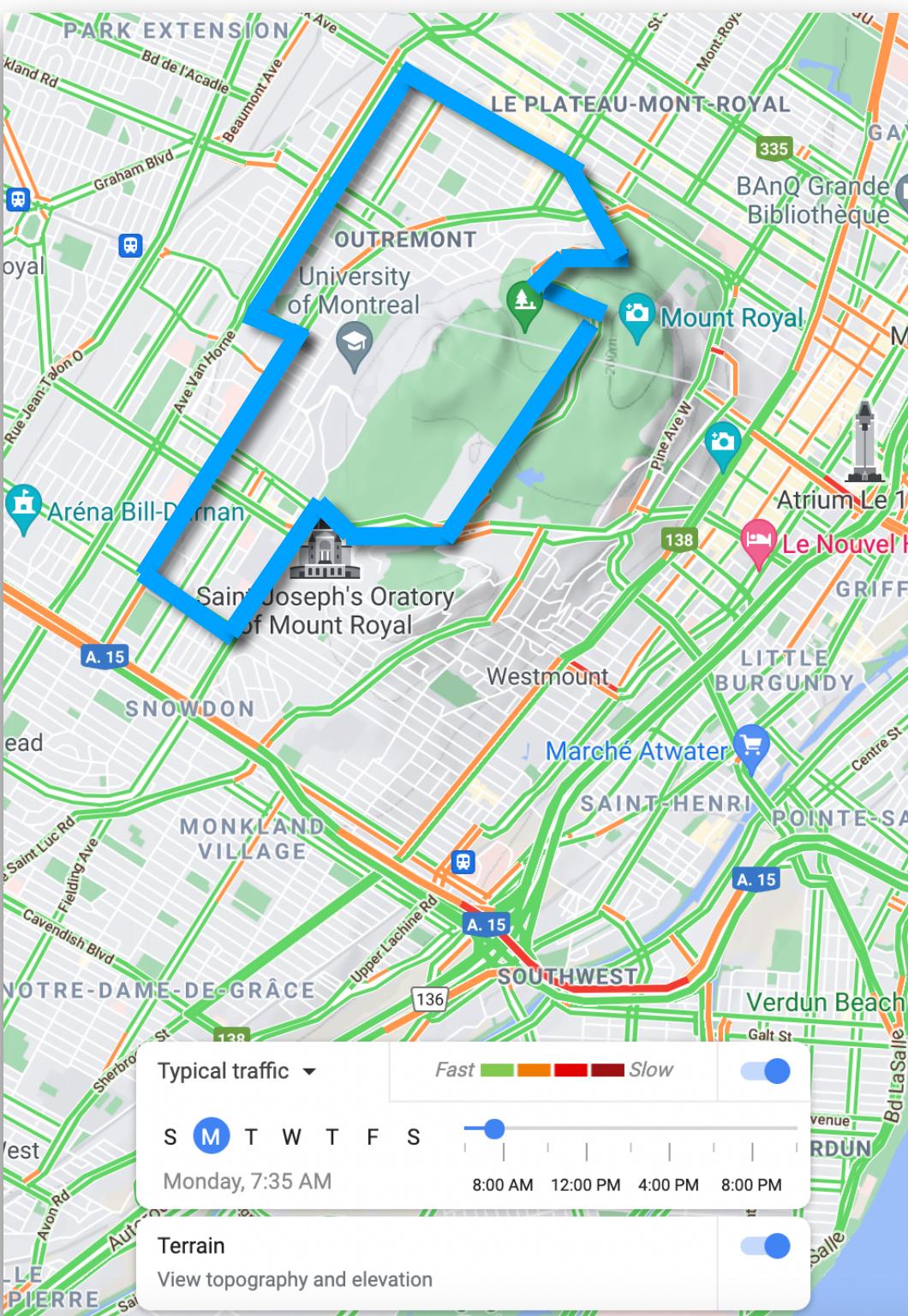
Google Maps, Montréal, Québec, Canada

# TSPs in Montréal

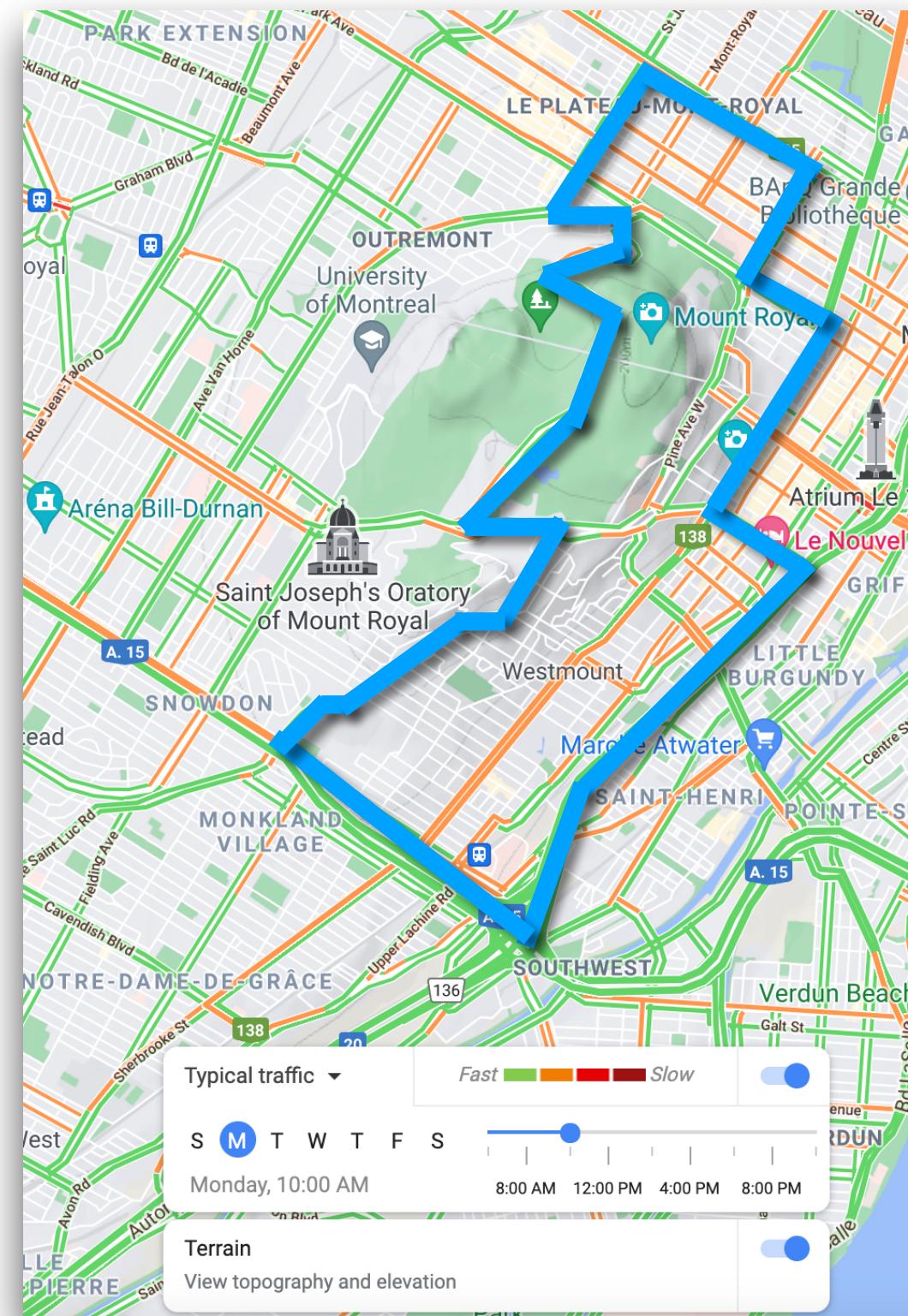
Every few hours, optimally **route** the salesperson through a set of locations



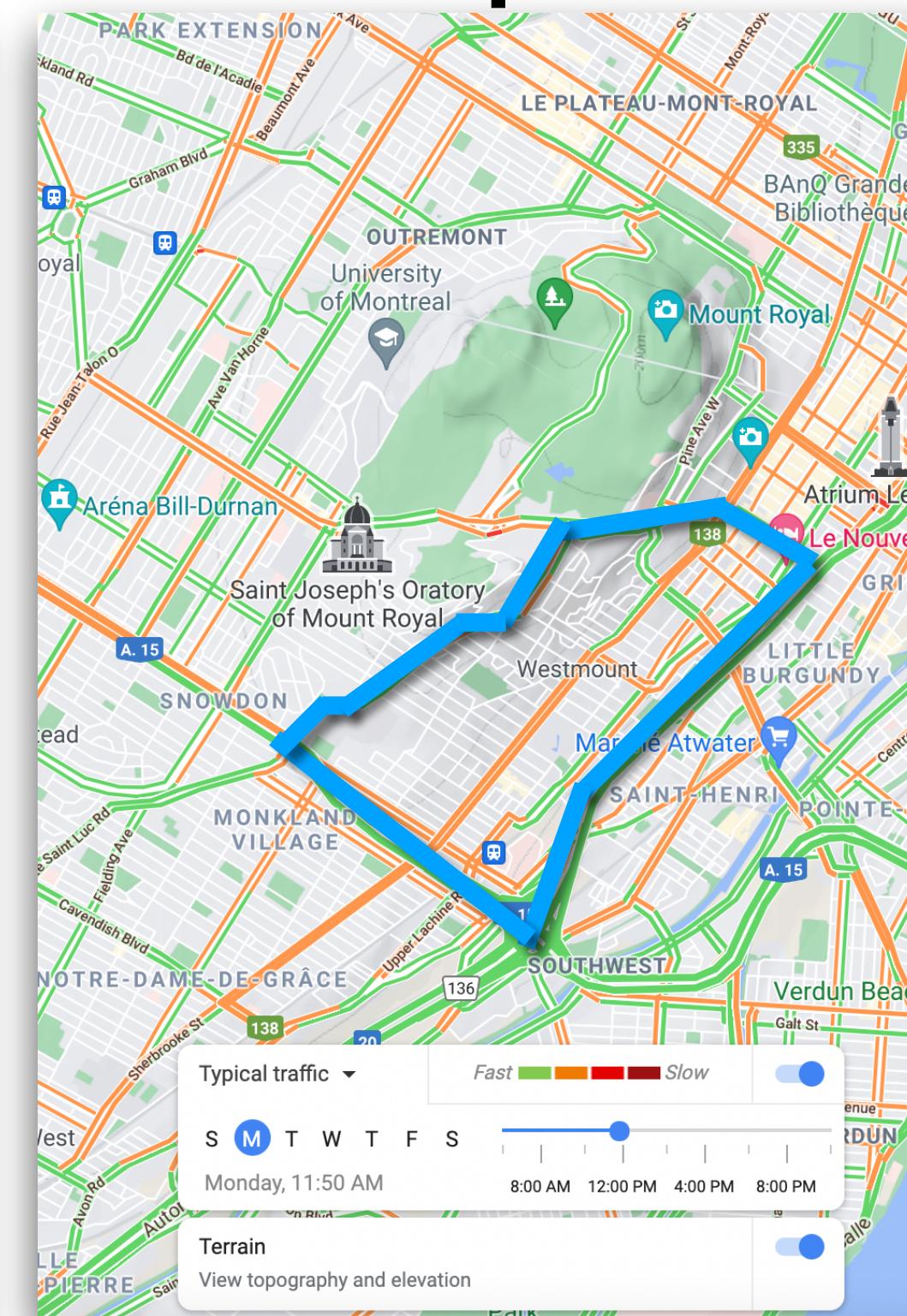
8 am



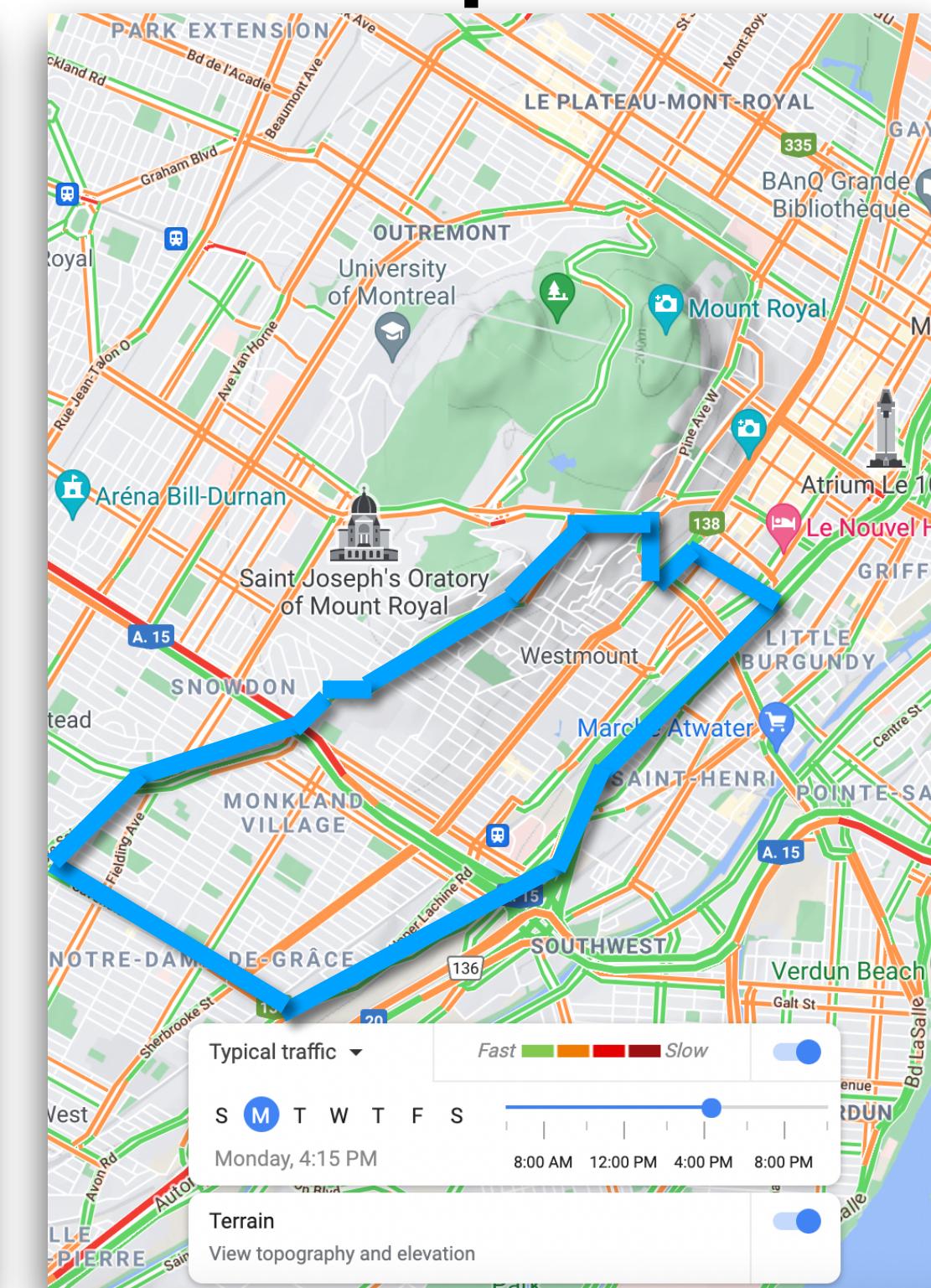
10 am



12 pm



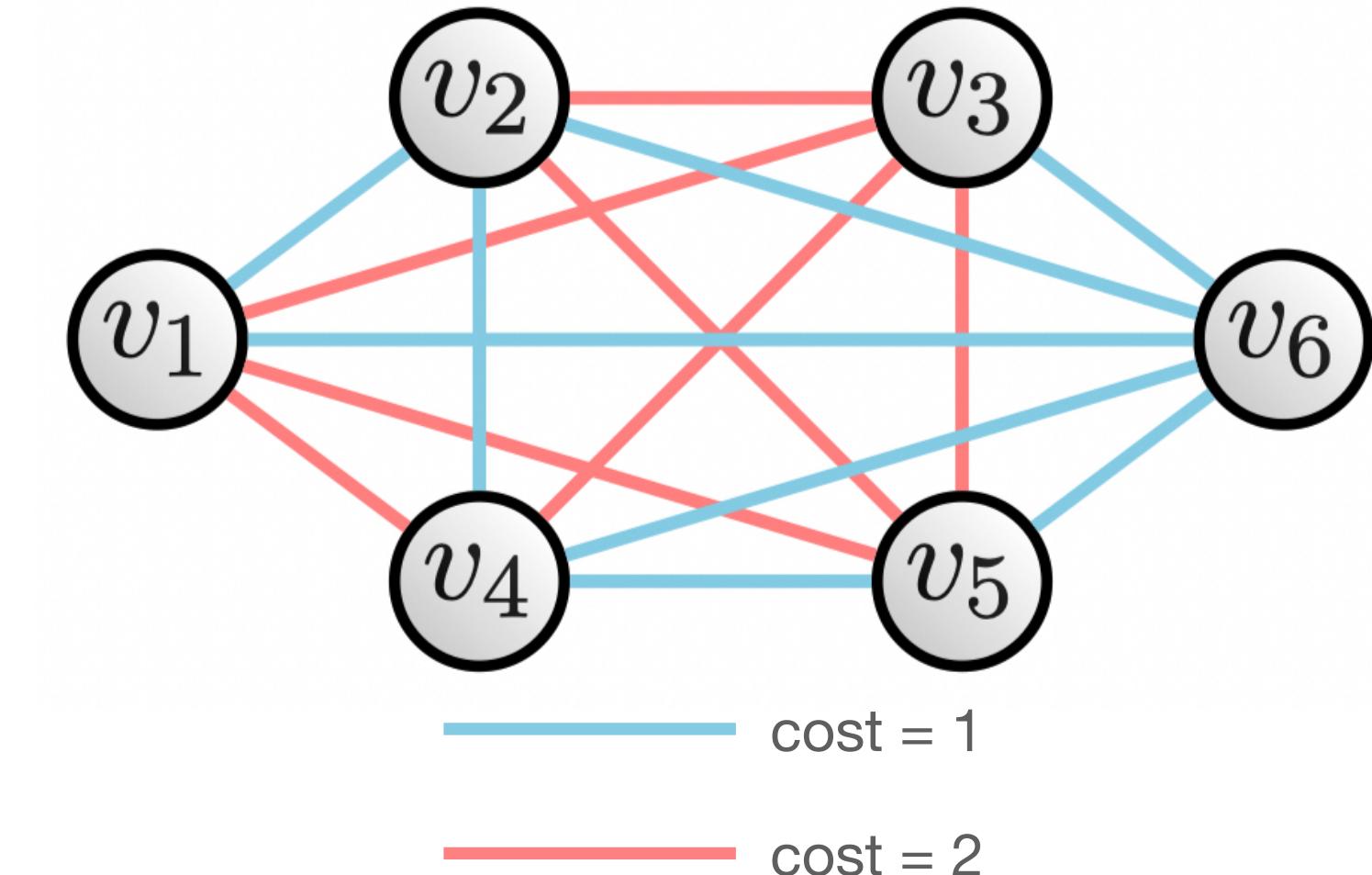
4 pm



Google Maps, Montréal, Québec, Canada

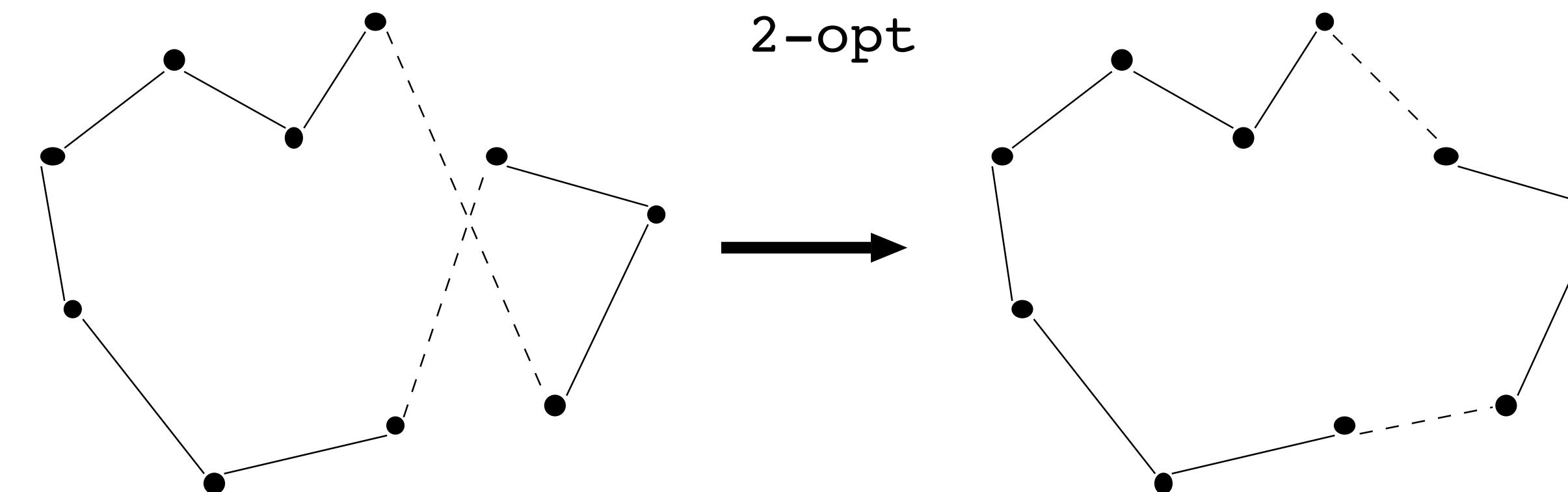
## Nearest Neighbour heuristic for the TSP:

- always choose at the current city the closest unvisited city
  - choose an arbitrary initial city  $\pi(1)$
  - at the  $i$ th step choose city  $\pi(i + 1)$  to be the city  $j$  that minimises  $\{d(\pi(i), j)\}; j \neq \pi(k), 1 \leq k \leq i$



## Iterative Improvement for the TSP

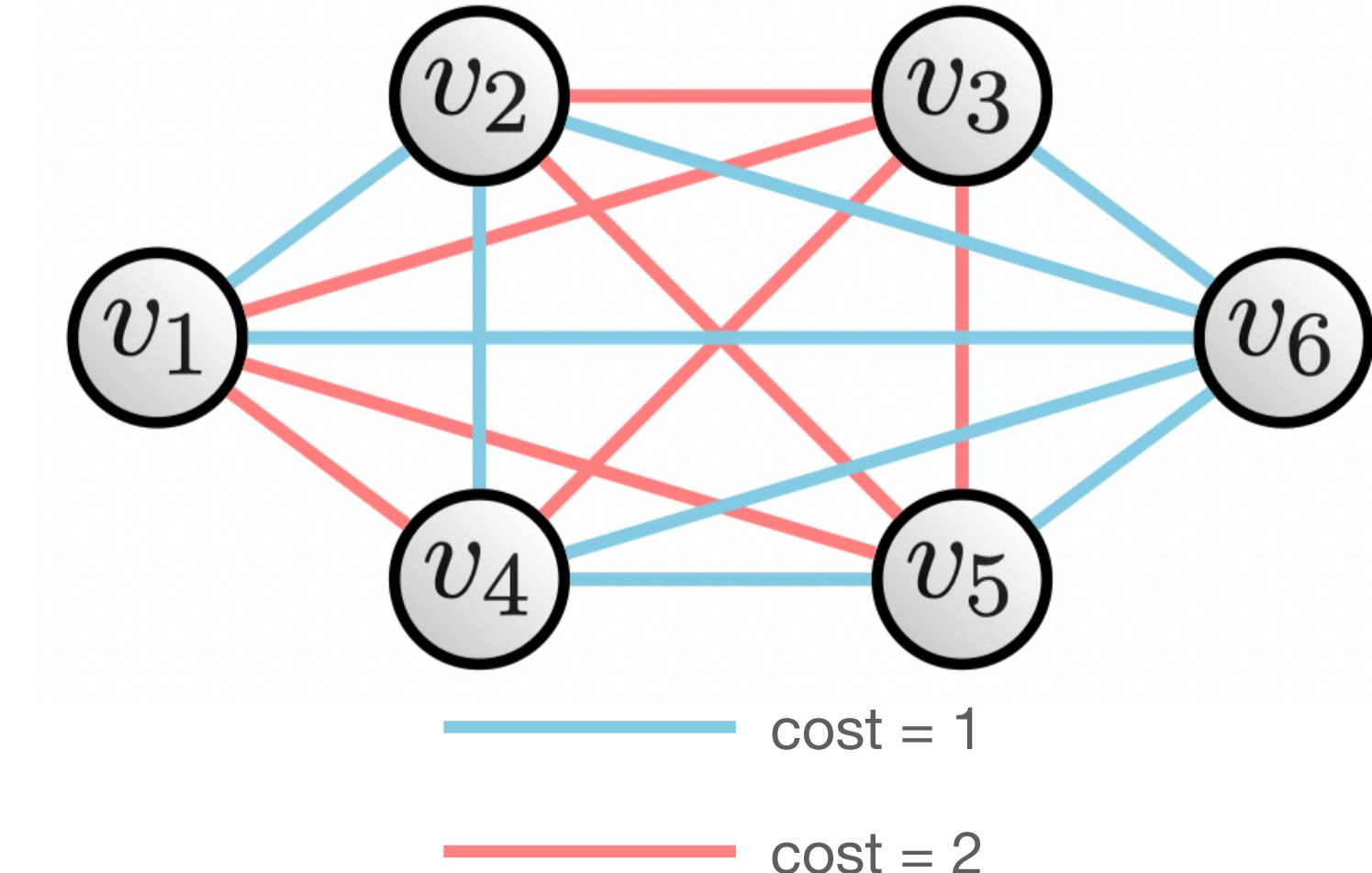
- initial solution is a complete tour
- $k$ -opt neighbourhood: solutions which differ by at most  $k$  edges



- neighbourhood size  $\mathcal{O}(n^k)$

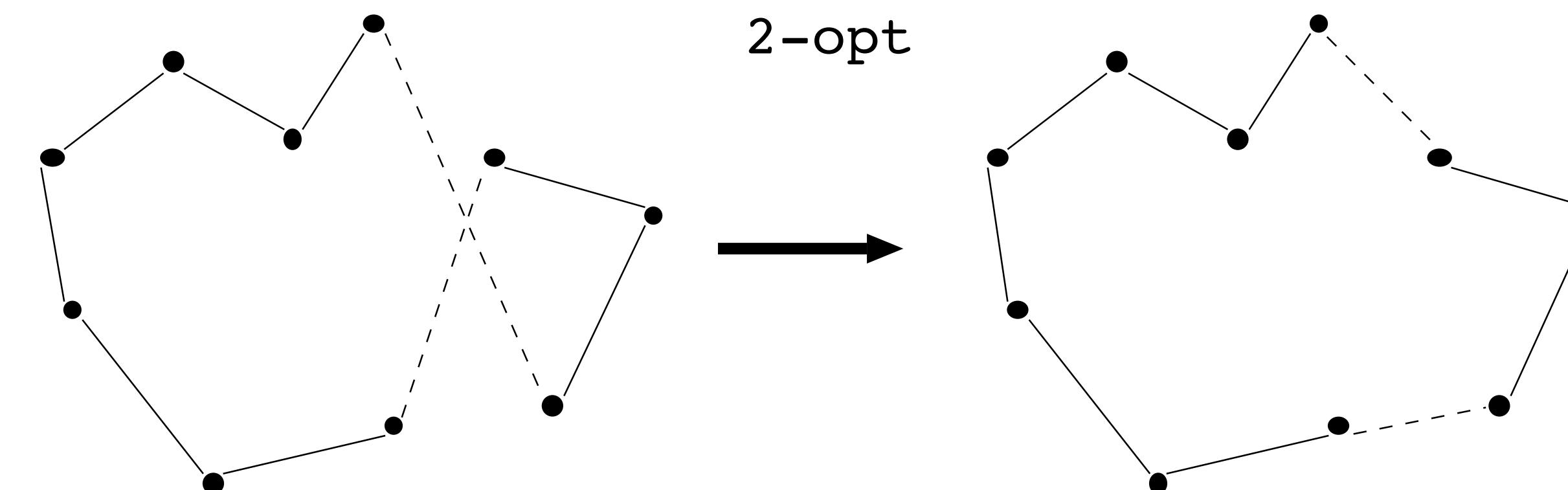
## Nearest Neighbour heuristic for the TSP:

- always choose at the current city the closest unvisited city
  - choose an arbitrary initial city  $\pi(1)$
  - at the  $i$ th step choose city  $\pi(i + 1)$  to be the city  $j$  that minimises  $\{d(\pi(i), j)\}; j \neq \pi(k), 1 \leq k \leq i$



## Iterative Improvement for the TSP

- initial solution is a complete tour
- $k$ -opt neighbourhood: solutions which differ by at most  $k$  edges

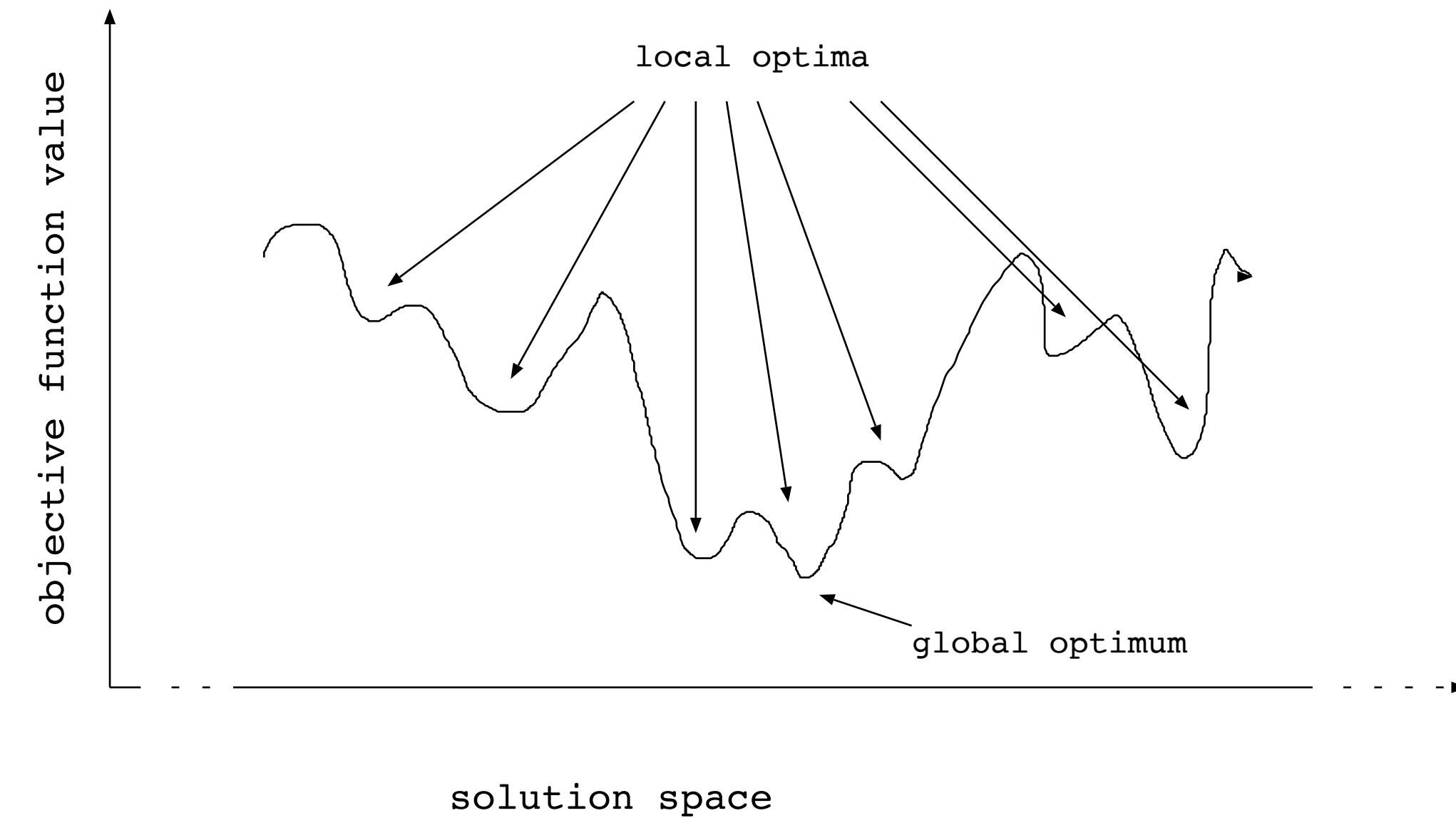


Problems with local search?

- neighbourhood size  $\mathcal{O}(n^k)$

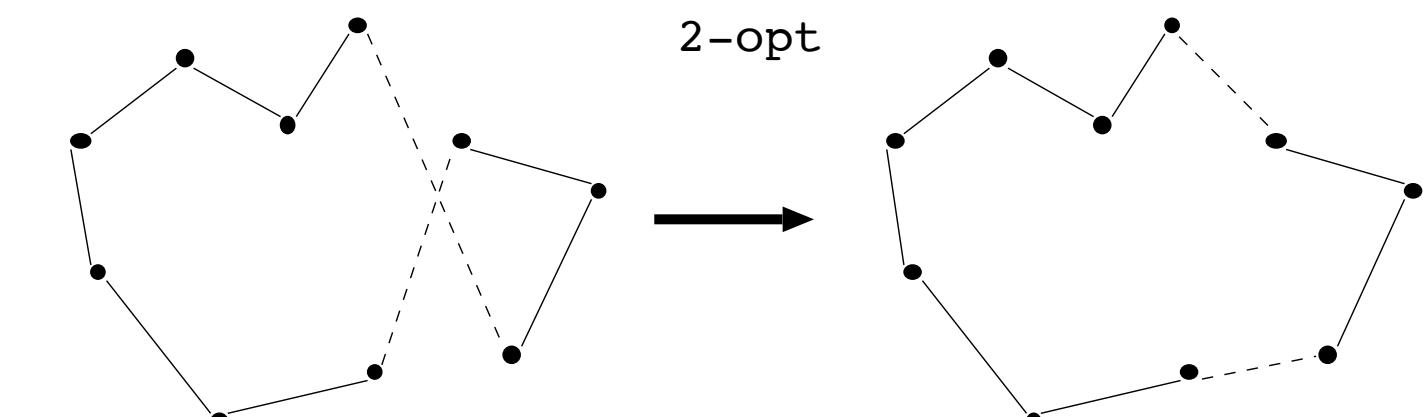
## Stochastic Local Search:

- randomise initialisation step
  - random initial solutions
  - randomised construction heuristics
- randomise search steps such that suboptimal/worsening steps are allowed  
~ improved performance & robustness
- typically, degree of randomisation controlled by noise parameter
- allows to invest arbitrary computation times



## Iterative Improvement for the TSP

- initial solution is a complete tour
- $k$ -opt neighbourhood: solutions which differ by at most  $k$  edges



# 0-1 Knapsack Problem

There is a budget  $b$  available for investment in projects during the coming year and  $n$  projects are under consideration, where  $a_j$  is the outlay for project  $j$  and  $c_j$  is its expected return. The goal is to choose a set of projects so that the budget is not exceeded and the expected return is maximized.

*Definition of the variables.*

$x_j = 1$  if project  $j$  is selected, and  $x_j = 0$  otherwise.

*Definition of the constraints.*

The budget cannot be exceeded:

$$\sum_{j=1}^n a_j x_j \leq b.$$

The variables are 0-1:

$$x_j \in \{0, 1\} \quad \text{for } j = 1, \dots, n.$$

*Definition of the objective function.*

The expected return is maximized:

$$\max \sum_{j=1}^n c_j x_j.$$

Can you find a  
feasible solution  
greedily?

# Knapsack Greedy algorithm

*Definition of the objective function.*

The expected return is maximized:

$$\max \sum_{j=1}^n c_j x_j.$$

The budget cannot be exceeded:

$$\sum_{j=1}^n a_j x_j \leq b.$$

The variables are 0-1:

$$x_j \in \{0, 1\} \quad \text{for } j = 1, \dots, n.$$

Sort items in increasing order of  $\frac{c_j}{a_j}$

While budget  $< b$  and there is an item that fits:

Insert next best item from sorted list into knapsack; update budget

# Knapsack

## Greedy algorithm

$$\max \sum_{j=1}^n c_j x_j. \quad \text{s.t. } \sum_{j=1}^n a_j x_j \leq b. \quad x_j \in \{0, 1\} \quad \text{for } j = 1, \dots, n.$$

Sort items in increasing order of  $\frac{c_j}{a_j}$

While budget  $< b$  and there is an item that fits:

Insert next best item from sorted list into knapsack; update budget

$$b = 5$$

$j$	$c_j$	$a_j$	$c_j/a_j$
1	6	1	6
2	10	2	5
3	12	3	4

# Knapsack

## Greedy algorithm

$$\max \sum_{j=1}^n c_j x_j. \quad \text{s.t. } \sum_{j=1}^n a_j x_j \leq b. \quad x_j \in \{0, 1\} \quad \text{for } j = 1, \dots, n.$$

Sort items in increasing order of  $\frac{c_j}{a_j}$

While budget  $< b$  and there is an item that fits:

Insert next best item from sorted list into knapsack; update budget

$$b = 5$$

$j$	$c_j$	$a_j$	$c_j/a_j$	$c_j/a_j^{0.5}$
1	6	1	6	6
2	10	2	5	7.07
3	12	3	4	6.9

# Branch & Bound for Integer Optimization

LP-based

$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

Land & Doig, 1960

Repeat:

- 1 Select Node
- 2 Solve LP Relaxation
- 3 Prune?
- 4 Add Cuts
- 5 Run Heuristics
- 6 Branch

# Branch & Bound for Integer Optimization

LP-based

$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

Land & Doig, 1960

Repeat:

**1 Select Node**



**2 Solve LP Relaxation**

**3 Prune?**

**4 Add Cuts**

**5 Run Heuristics**

**6 Branch**

# Branch & Bound for Integer Optimization

LP-based

$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

Land & Doig, 1960

Repeat:

- 1 **Select Node**
- 2 **Solve LP Relaxation**
- 3 **Prune?**
- 4 **Add Cuts**
- 5 **Run Heuristics**
- 6 **Branch**



Solve LP Relaxation  
→ Lower Bound on OPT

# Branch & Bound for Integer Optimization

LP-based

$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

Land & Doig, 1960

Repeat:

1 **Select Node**



Solve LP Relaxation  
→ Lower Bound on OPT

2 Solve LP Relaxation

worse than best solution?  
Prune!

3 Prune?

4 Add Cuts

5 Run Heuristics

6 Branch

# Branch & Bound for Integer Optimization

LP-based

$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

Land & Doig, 1960

Repeat:

- 1 **Select Node**
- 2 **Solve LP Relaxation**
- 3 **Prune?**
- 4 **Add Cuts**
- 5 **Run Heuristics**
- 6 **Branch**



Add Cuts:  
Tightening Constraints

# Branch & Bound for Integer Optimization

LP-based

$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

Land & Doig, 1960

Repeat:

1 **Select Node**

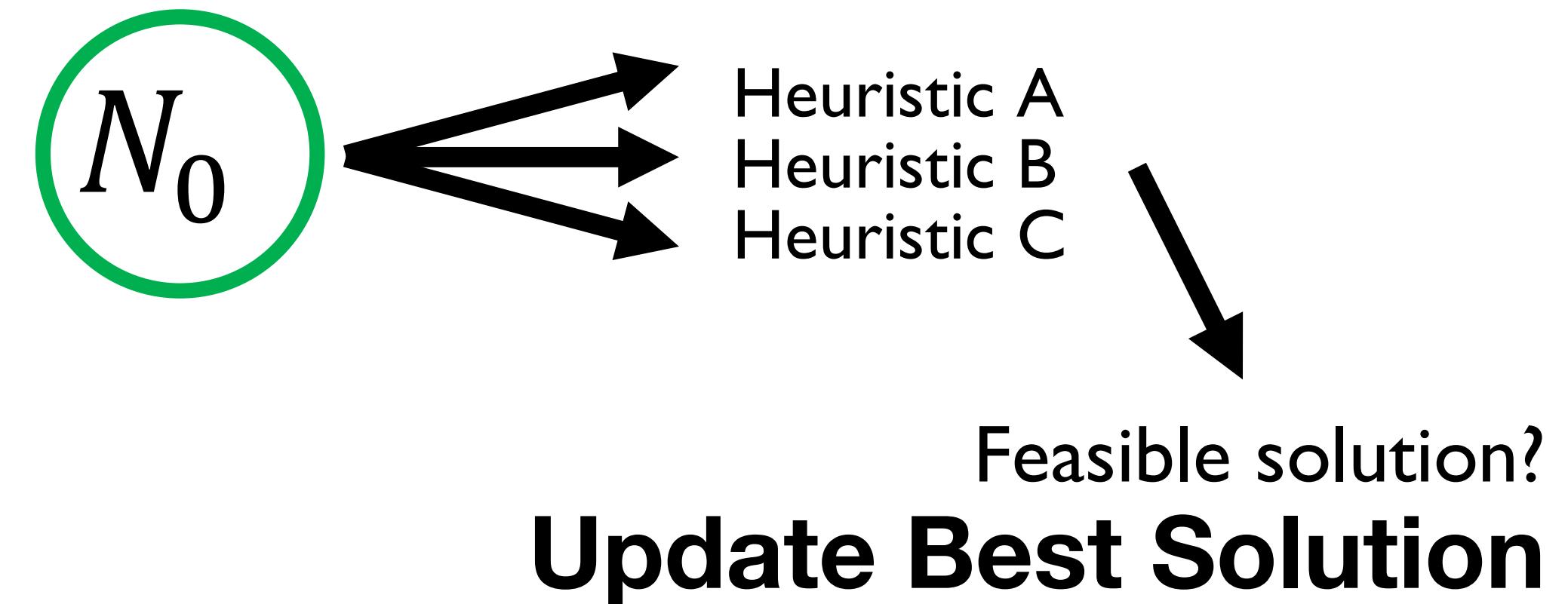
2 **Solve LP Relaxation**

3 **Prune?**

4 **Add Cuts**

5 **Run Heuristics**

6 **Branch**



# Branch & Bound for Integer Optimization

LP-based

$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

Land & Doig, 1960

Repeat:

**1 Select Node**

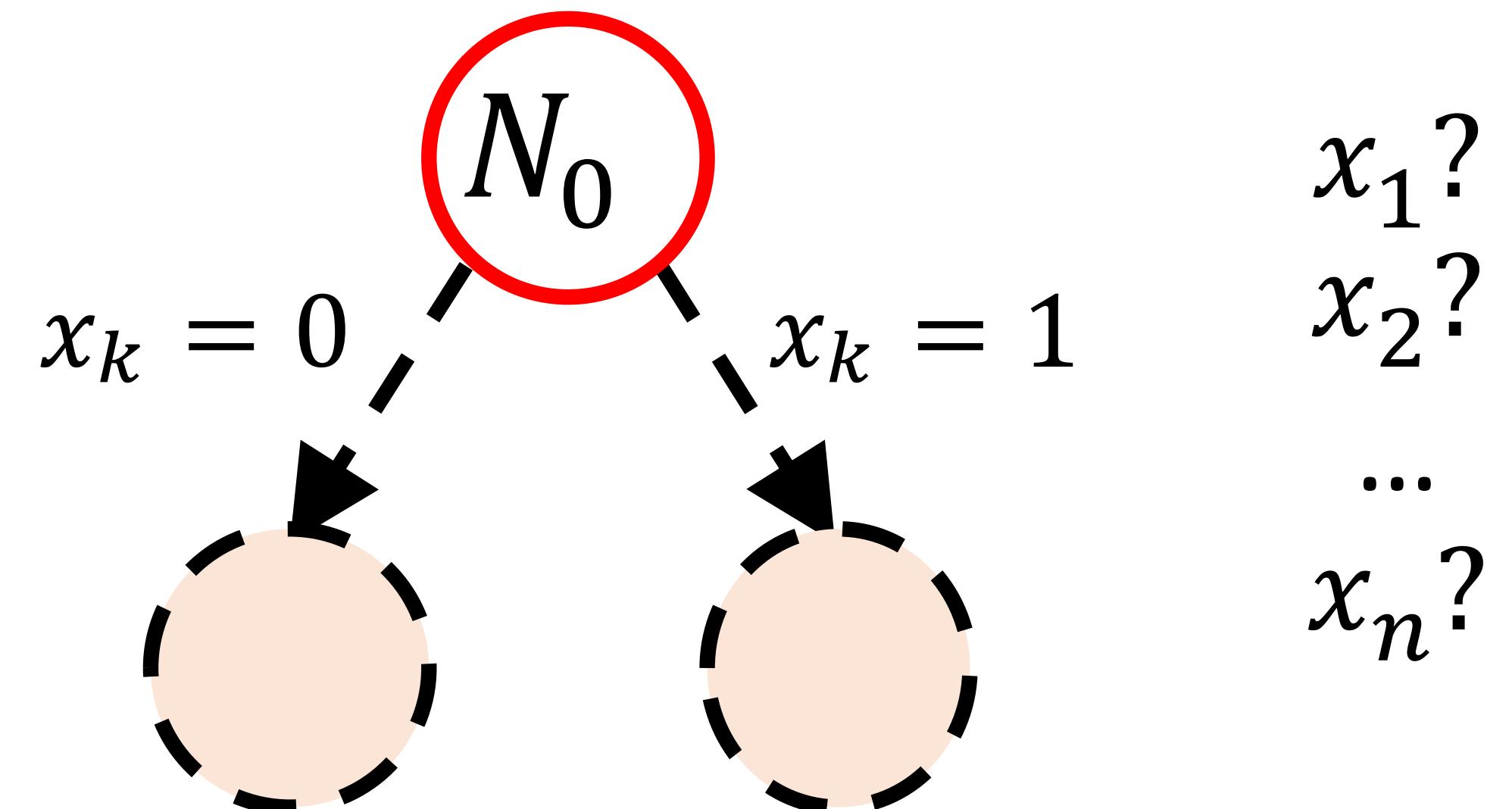
**2 Solve LP Relaxation**

**3 Prune?**

**4 Add Cuts**

**5 Run Heuristics**

**6 Branch**



# Branch & Bound for Integer Optimization

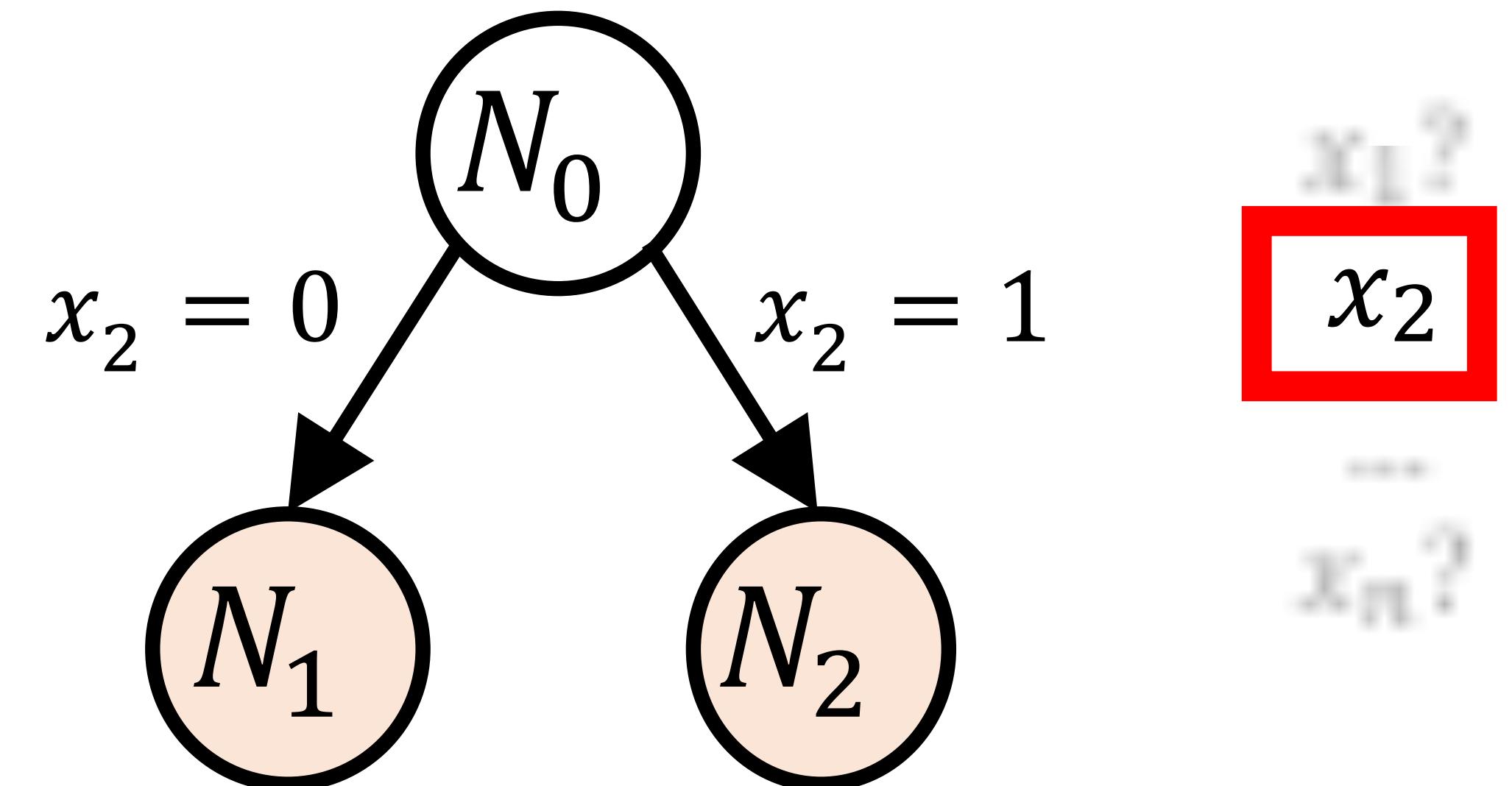
LP-based

$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

Land & Doig, 1960

Repeat:

- 1 Select Node**
- 2 Solve LP Relaxation**
- 3 Prune?**
- 4 Add Cuts**
- 5 Run Heuristics**
- 6 Branch**



# Branch & Bound for Integer Optimization

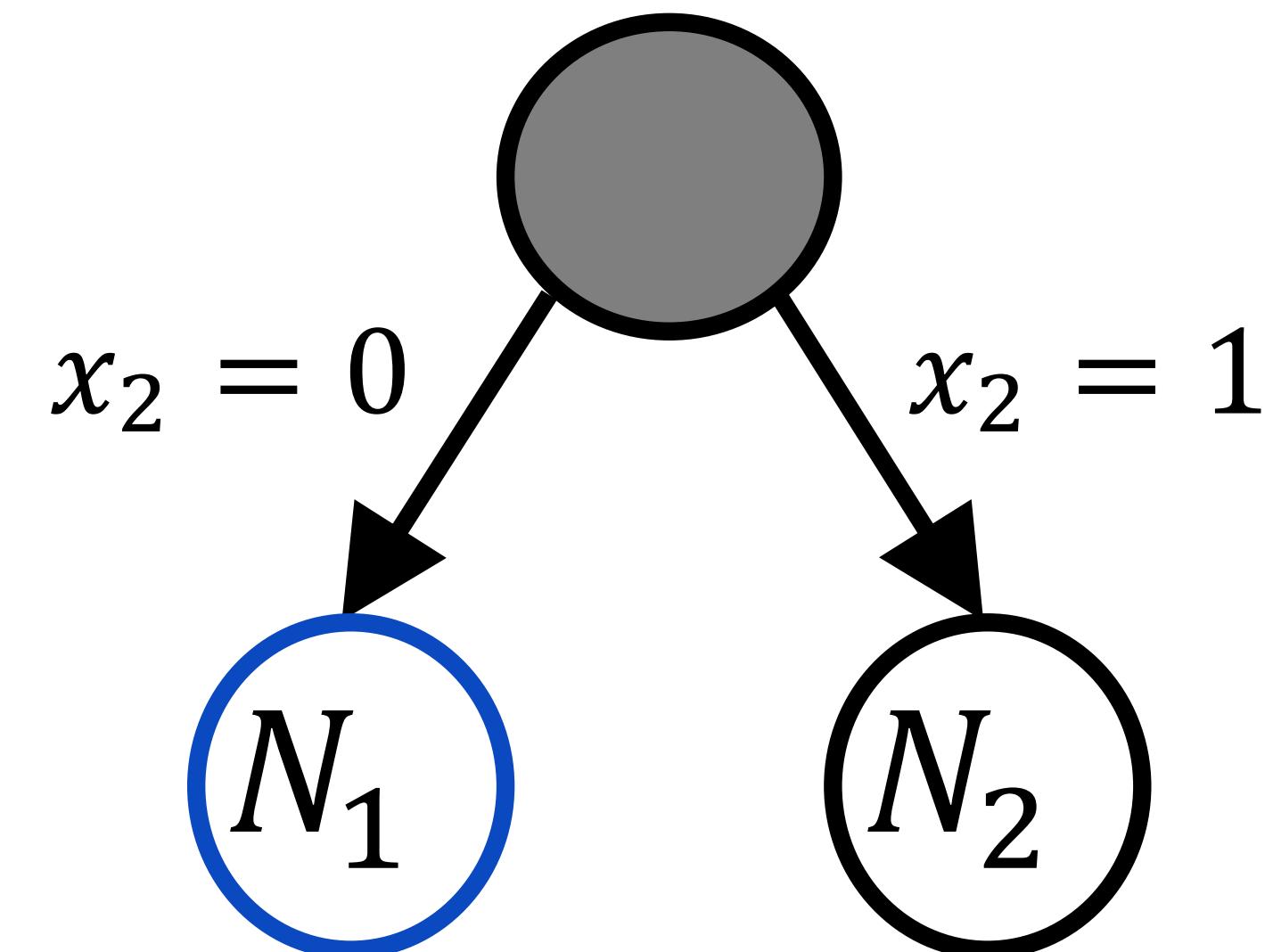
LP-based

$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

Land & Doig, 1960

Repeat:

- 1 Select Node**
- 2 Solve LP Relaxation**
- 3 Prune?**
- 4 Add Cuts**
- 5 Run Heuristics**
- 6 Branch**



# Branch & Bound for Integer Optimization

LP-based

$$\min_x c^T x \text{ s.t. } Ax \leq b, x \in \{0,1\}^n$$

Land & Doig, 1960

Repeat:

1 Select Node

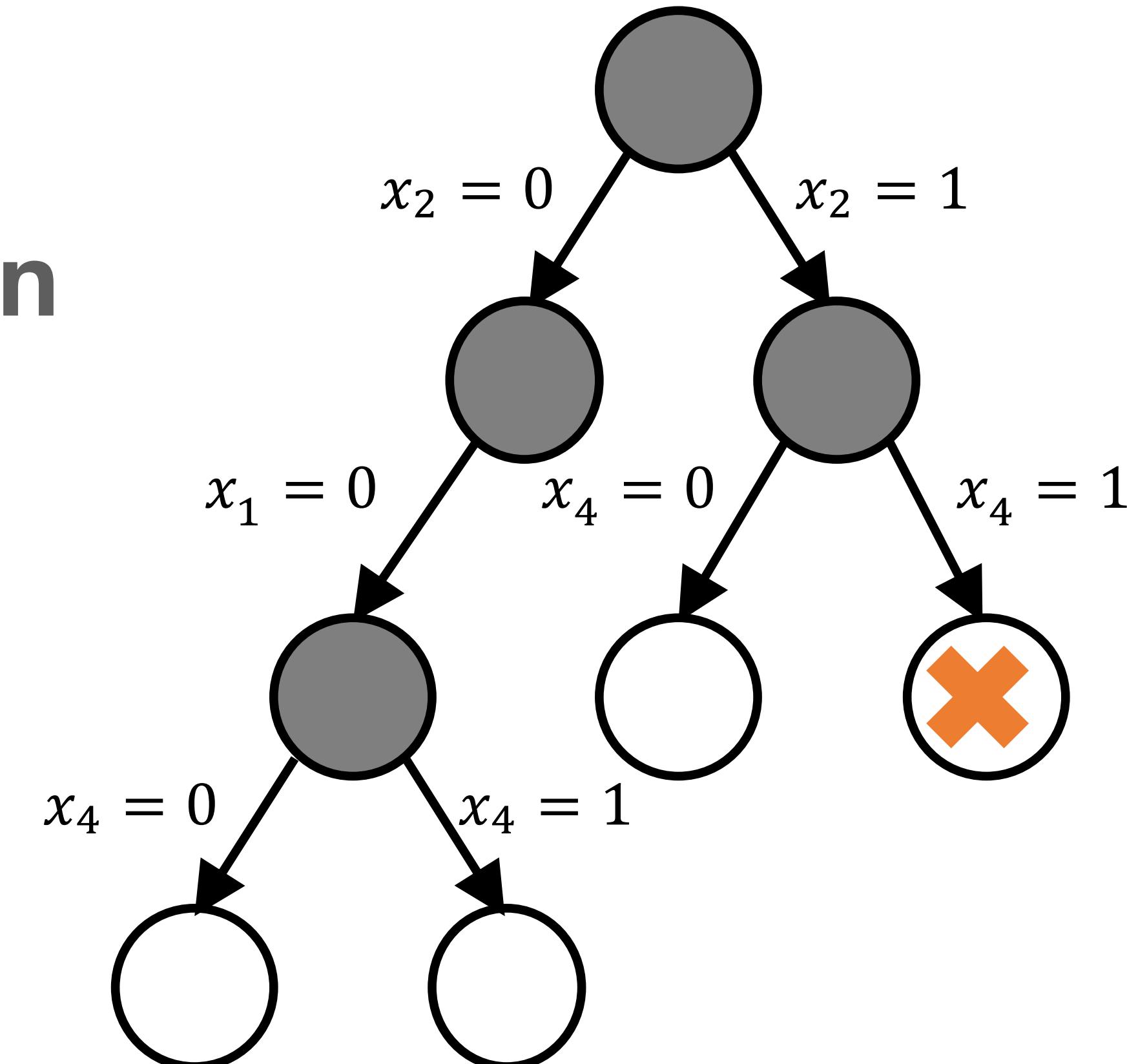
2 Solve LP Relaxation

3 Prune?

4 Add Cuts

5 Run Heuristics

6 Branch



# Towards Tailored Algorithms



IBM Knowledge Center

Managing sets of parameters

Parameter names

Correspondence of parameters  
between APIs

Saving parameter settings to a file  
in the C API

## – Topical list of parameters

Barrier

Benders algorithm

Distributed MIP

## – MIP

MIP general

**MIP strategies**

MIP cuts

MIP tolerances

MIP limits

Here are links to parameters controlling MIP strategies.

[algorithm for initial MIP relaxation](#)

[Benders strategy](#)

[MIP subproblem algorithm](#)

[MIP variable selection strategy](#)

[MIP strategy best bound interval](#)

[MIP branching direction](#)

[backtracking tolerance](#)

[MIP dive strategy](#)

[MIP heuristic effort](#)

**CPLEX Documentation**

# Towards Tailored Algorithms

MIP variable selection strategy	Value	Symbol	Meaning
	-1	CPX_VARSEL_MININFEAS	Branch on variable with minimum infeasibility
	0	CPX_VARSEL_DEFAULT	Automatic: let CPLEX choose variable to branch on; <b>default</b>
	1	CPX_VARSEL_MAXINFEAS	Branch on variable with maximum infeasibility
	2	CPX_VARSEL_PSEUDO	Branch based on pseudo costs
	3	CPX_VARSEL_STRONG	Strong branching
	4	CPX_VARSEL_PSEUDOREDUCED	Branch based on pseudo reduced costs

MIP heuristic frequency	Value	Meaning
	-1	None
	0	Automatic: let CPLEX choose; <b>default</b>
	Any positive integer	Apply the periodic heuristic at this frequency

- an algorithm  $A$  with parameters  $p_1, \dots, p_k$  that affect its behaviour,
- a space  $C$  of parameter settings (configurations), where  $c \in C$  specifies values for  $p_1, \dots, p_k$ ,
- a set of problem instances  $I$ ,
- a performance metric  $m$  that measures the performance of  $A$  on instance set  $I$  for a given configuration  $c$ ,

find a configuration  $c^* \in C$  such that running algorithm  $A$  on instance set  $I$  maximizes metric  $m$

## Greedy 0-1 knapsack

- an algorithm  $A$  with parameters  $p_1, \dots, p_k$  that affect its behaviour,

$$p_1 \in (0,1]$$

- a space  $C$  of parameter settings (configurations), where  $c \in C$  specifies values for  $p_1, \dots, p_k$ ,
- a set of problem instances  $I$ ,
- a performance metric  $m$  that measures the performance of  $A$  on instance set  $I$  for a given configuration  $c$ ,

$$\sum_{i \in I} \sum_{j=1}^n c_j x_j$$

find a configuration  $c^* \in C$  such that running algorithm  $A$  on instance set  $I$  maximizes metric  $m$

# Issues to consider

- Generalization
- Time-outs!
- **Optimization!**

# Racing Procedures

- Assumptions: small, finite configuration space  $C$
- Basic idea:
  - sample instance,
  - test remaining configs.,
  - eliminate really bad ones relative to current best config.,
  - repeat.



123rf.com

**procedure** *F-Race*

**input** target algorithm *A*, set of configurations *C*, set of problem instances *I*,  
performance metric *m*;

**parameters** integer  $ni_{min}$ ;

**output** set of configurations  $C^*$ ;

$C^* := C$ ;  $ni := 0$ ;

**repeat**

randomly choose instance *i* from set *I*;

run all configurations of *A* in  $C^*$  on *i*;

$ni := ni + 1$ ;

**if**  $ni \geq ni_{min}$  **then**

perform rank-based Friedman test on results for configurations in  $C^*$  on all instances  
in *I* evaluated so far;

**if** test indicates significant performance differences **then**

$c^* :=$  best configuration in  $C^*$  (according to *m* over instances evaluated so far);

**for all**  $c \in C^* \setminus \{c^*\}$  **do**

perform pairwise Friedman post hoc test on *c* and  $c^*$ ;

**if** test indicates significant performance differences **then**

eliminate *c* from  $C^*$ ;

**end if**;

**end for**;

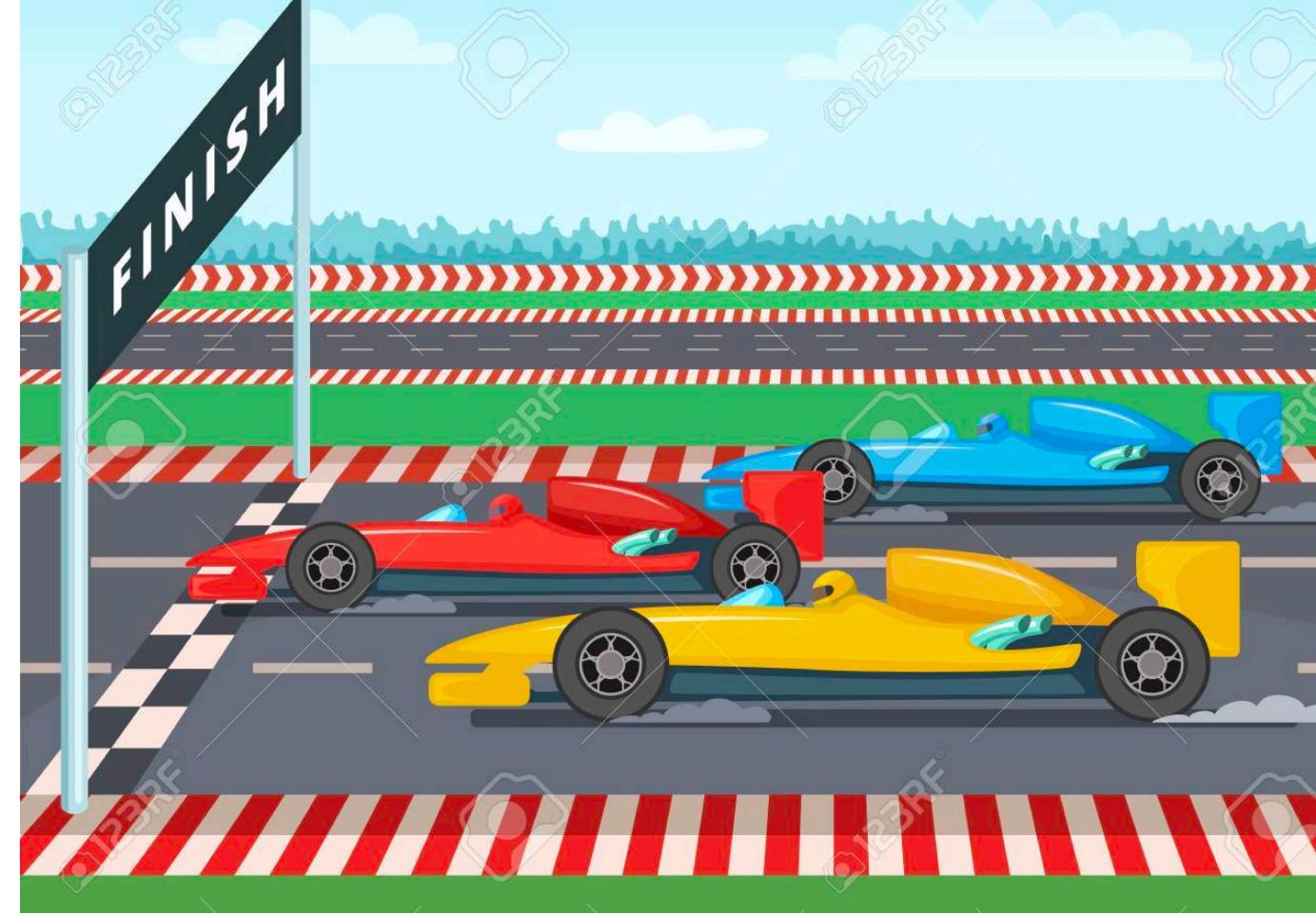
**end if**;

**end if**;

**until** termination condition met;

**return**  $C^*$ ;

**end** *F-Race*



123rf.com

Birattari, Mauro, et al. "A Racing Algorithm  
for Configuring Metaheuristics." GECCO.  
Vol. 2. No. 2002.

**procedure** *F-Race*

**input** target algorithm *A*, set of configurations *C*, set of problem instances *I*,  
performance metric *m*;

**parameters** integer  $ni_{min}$ ;

**output** set of configurations  $C^*$ ;

$C^* := C$ ;  $ni := 0$ ;

**repeat**

randomly choose instance *i* from set *I*;

run all configurations of *A* in  $C^*$  on *i*;

$ni := ni + 1$ ;

**if**  $ni \geq ni_{min}$  **then**

perform rank-based Friedman test on results for configurations in  $C^*$  on all instances  
in *I* evaluated so far;

**if** test indicates significant performance differences **then**

$c^* :=$  best configuration in  $C^*$  (according to *m* over instances evaluated so far);

**for all**  $c \in C^* \setminus \{c^*\}$  **do**

perform pairwise Friedman post hoc test on *c* and  $c^*$ ;

**if** test indicates significant performance differences **then**

eliminate *c* from  $C^*$ ;

**end if**;

**end for**;

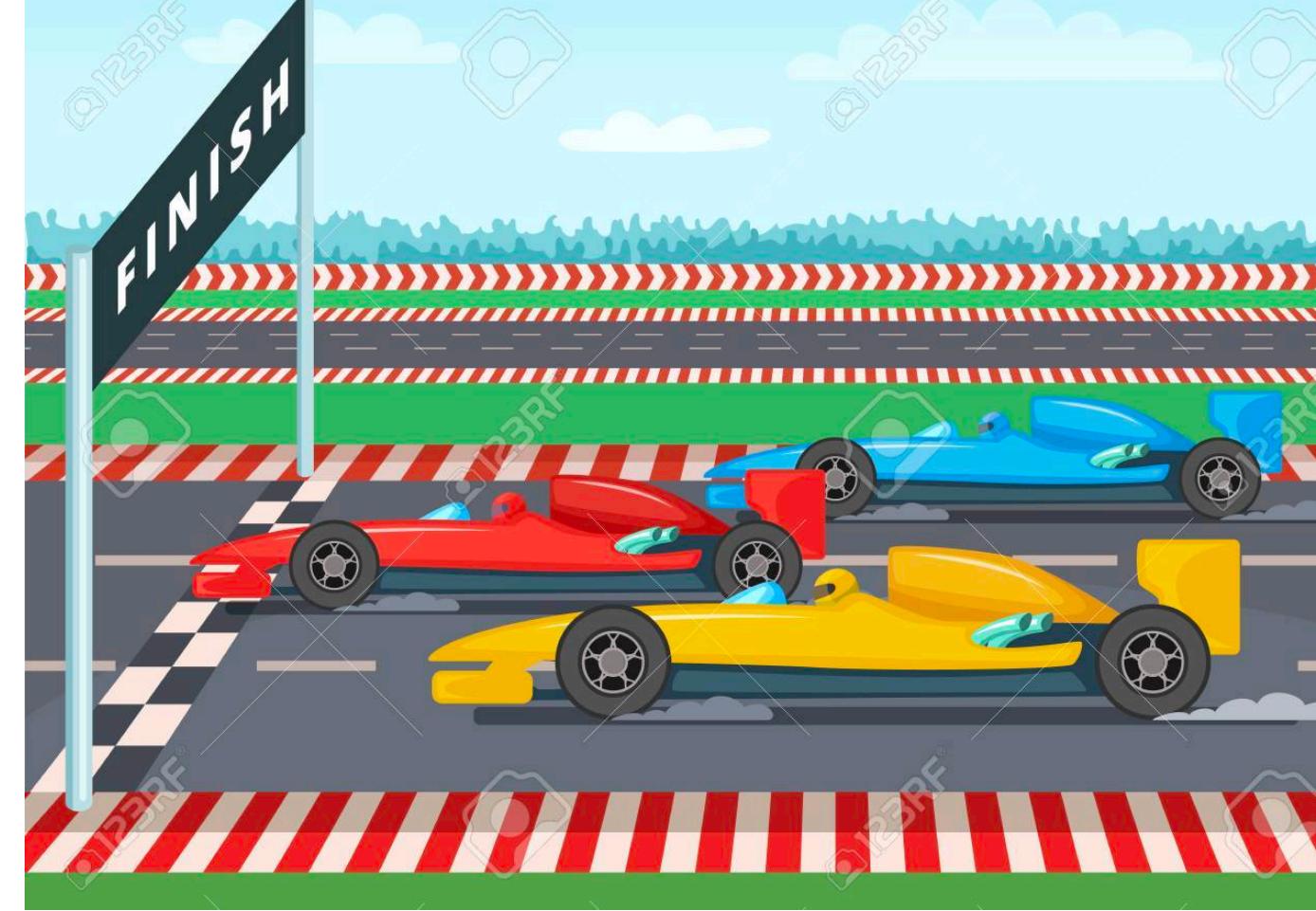
**end if**;

**end if**;

**until** termination condition met;

**return**  $C^*$ ;

**end** *F-Race*



123rf.com

Birattari, Mauro, et al. "A Racing Algorithm  
for Configuring Metaheuristics." GECCO.  
Vol. 2. No. 2002.

**procedure** *F-Race*

**input** target algorithm *A*, set of configurations *C*, set of problem instances *I*,  
 performance metric *m*;  
**parameters** integer  $ni_{min}$ ;  
**output** set of configurations  $C^*$ ;  
 $C^* := C$ ;  $ni := 0$ ;  
**repeat**  
 randomly choose instance *i* from set *I*;  
 run all configurations of *A* in  $C^*$  on *i*;  
 $ni := ni + 1$ ;  
**if**  $ni \geq ni_{min}$  **then**  
 perform rank-based Friedman test on results for configurations in  $C^*$  on all instances  
 in *I* evaluated so far;  
**if** test indicates significant performance differences **then**  
 $c^* :=$  best configuration in  $C^*$  (according to *m* over instances evaluated so far);  
**for all**  $c \in C^* \setminus \{c^*\}$  **do**  
 perform pairwise Friedman post hoc test on *c* and  $c^*$ ;  
**if** test indicates significant performance differences **then**  
 eliminate *c* from  $C^*$ ;  
**end if**;  
**end for**;  
**end if**;  
**end if**;  
**until** termination condition met;  
**return**  $C^*$ ;  
**end F-Race**

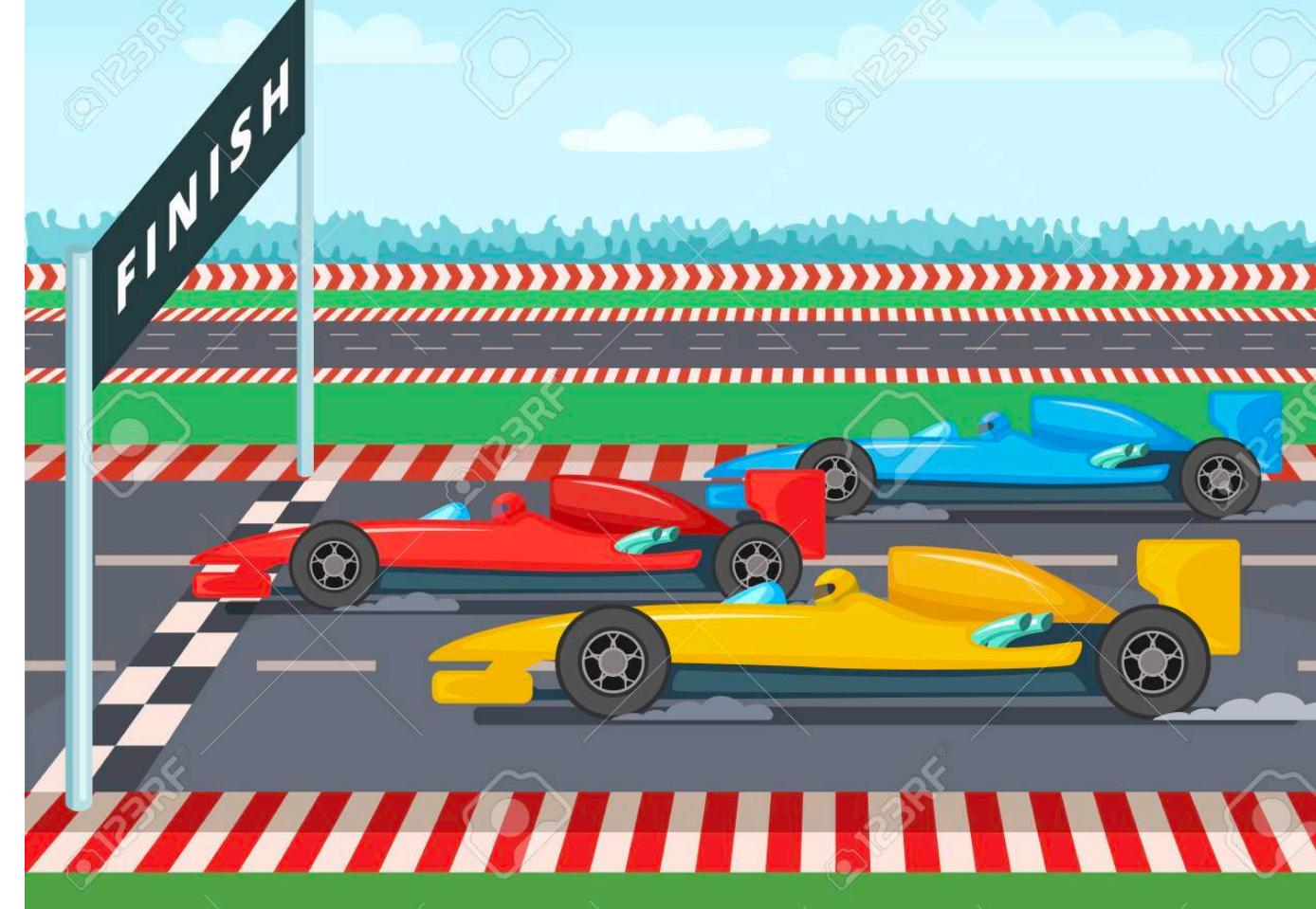
Are these results statistically “similar”?

	Config X	Config Y	Config Z
Instance 1	7	5	8
Instance 2	1	4	2
Instance 3	3	3	3
Instance 4	2	7	8

Sum of ranks	4	7	9
--------------	---	---	---

$$T = \frac{(n-1) \sum_{j=1}^n \left( R_j - \frac{k(n+1)}{2} \right)^2}{\sum_{l=1}^k \sum_{j=1}^n R_{lj}^2 - \frac{kn(n+1)^2}{4}}$$

Birattari, Mauro, et al. "A Racing Algorithm for Configuring Metaheuristics." GECCO. Vol. 2. No. 2002.



123rf.com

**procedure** *I/F-Race*

**input** target algorithm  $A$ , set of configurations  $C$ , set of problem instances  $I$ , performance metric  $m$ ;

**output** set of configurations  $C^*$ ;

initialise probabilistic model  $M$ ;

$C' := \emptyset$ ; // later,  $C'$  is the set of survivors from the previous F-Race

**repeat**

    based on model  $M$ , sample set of configurations  $\hat{C} \subseteq C$ ;

    perform F-Race on configurations in  $\hat{C} \cup C'$  to obtain set of configurations  $C^*$ ;

    update probabilistic model  $M$  based on configurations in  $C^*$ ;

$C' := C^*$ ;

**until** termination condition met;

**return**  $c^* \in C^*$  with best performance (according to  $m$ ) over all instances evaluated;

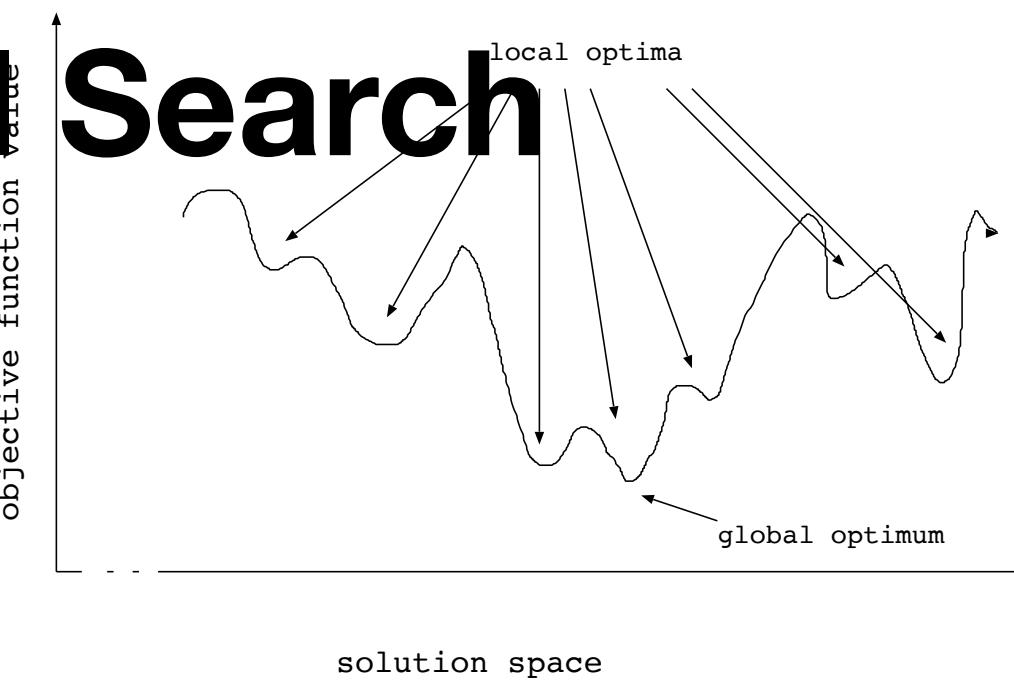
**end** *I/F-Race*

Balaprakash, Prasanna, Mauro Birattari, and Thomas Stützle. "Improvement strategies for the F-Race algorithm: Sampling design and iterative refinement." International workshop on hybrid metaheuristics. Springer, Berlin, Heidelberg, 2007.

# ParamILS

```
procedure ParamILS
  input target algorithm A, set of configurations C, set of problem instances I,
         performance metric m;
  parameters configuration  $c_0 \in C$ , integer r, integer s, probability pr;
  output configuration  $c^*$ ;
   $c^* := c_0$ ;
  for  $i := 1$  to r do
    draw  $c$  from  $C$  uniformly at random;
    assess  $c$  against  $c^*$  based on performance of A on instances from  $I$  according to metric  $m$ ;
    if  $c$  found to perform better than  $c^*$  then
       $c^* := c$ ;
    end if;
  end for;
   $c := c^*$  ;
  perform subsidiary local search on  $c$ ;
  while termination condition not met do
     $c' := c$ ;
    perform s random perturbation steps on  $c'$ 
    perform subsidiary local search on  $c'$ ;
    assess  $c'$  against  $c$  based on performance of A on instances from  $I$  according to metric  $m$ ;
    if  $c'$  found to perform better than  $c$  then // acceptance criterion
      update overall incumbent  $c^*$ ;
       $c := c'$ ;
    end if;
    with probability pr do
      draw  $c$  from  $C$  uniformly at random;
    end with probability;
  end while;
  return  $c^*$ ;
end ParamILS
```

## ILS: Iterated Local Search



### Initial sampling phase

Random perturbation + local search  
Evaluation  
Update incumbent config.

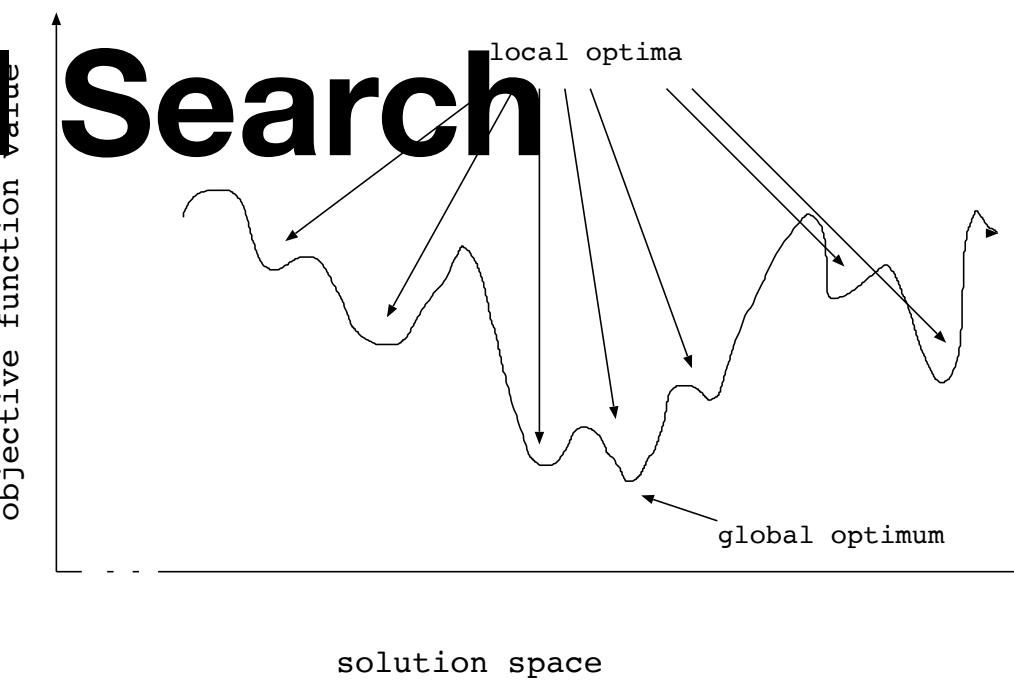
Random restart!

Hutter, Frank, et al. "ParamILS: an automatic algorithm configuration framework." *Journal of Artificial Intelligence Research* 36 (2009): 267-306.

# ParamILS

```
procedure ParamILS
  input target algorithm A, set of configurations C, set of problem instances I,
         performance metric m;
  parameters configuration  $c_0 \in C$ , integer r, integer s, probability pr;
  output configuration  $c^*$ ;
   $c^* := c_0;$ 
  for  $i := 1$  to r do
    draw  $c$  from  $C$  uniformly at random;
    assess  $c$  against  $c^*$  based on performance of A on instances from  $I$  according to metric  $m$ ;
    if  $c$  found to perform better than  $c^*$  then
       $c^* := c;$ 
    end if;
  end for;
   $c := c^*;$ 
  perform subsidiary local search on  $c$ ;
  while termination condition not met do
     $c' := c;$ 
    perform  $s$  random perturbation steps on  $c'$ 
    perform subsidiary local search on  $c'$ ;
    assess  $c'$  against  $c$  based on performance of A on instances from  $I$  according to metric  $m$ ;
    if  $c'$  found to perform better than  $c$  then // acceptance criterion
      update overall incumbent  $c^*$ ;
       $c := c'$ ;
    end if;
    with probability  $pr$  do
      draw  $c$  from  $C$  uniformly at random;
    end with probability;
  end while;
  return  $c^*$ ;
end ParamILS
```

## ILS: Iterated Local Search



### Initial sampling phase

### Random perturbation + local search

### Evaluation

### Update incumbent config.

### Random restart!

Hutter, Frank, et al. "ParamILS: an automatic algorithm configuration framework." *Journal of Artificial Intelligence Research* 36 (2009): 267-306.

# Sequential Model-Based Optimization

**procedure** *SMBO*

**input** *target algorithm A, set of configurations C, set of problem instances I, performance metric m;*

**output** *configuration  $c^*$ ;*

    determine initial set of configurations  $C_0 \subset C$ ;

    for all  $c \in C_0$ , measure performance of *A* on *I* according to metric *m*;

    build initial model *M* based on performance measurements for  $C_0$ ;

    determine incumbent  $c^* \in C_0$  for which best performance was observed or predicted;

**repeat**

        based on model *M*, determine set of configurations  $C' \subseteq C$ ;

        for all  $c \in C'$ , measure performance of *A* on *I* according to metric *m*;

        update model *M* based on performance measurements for  $C'$ ;

        update incumbent  $c^*$ ;

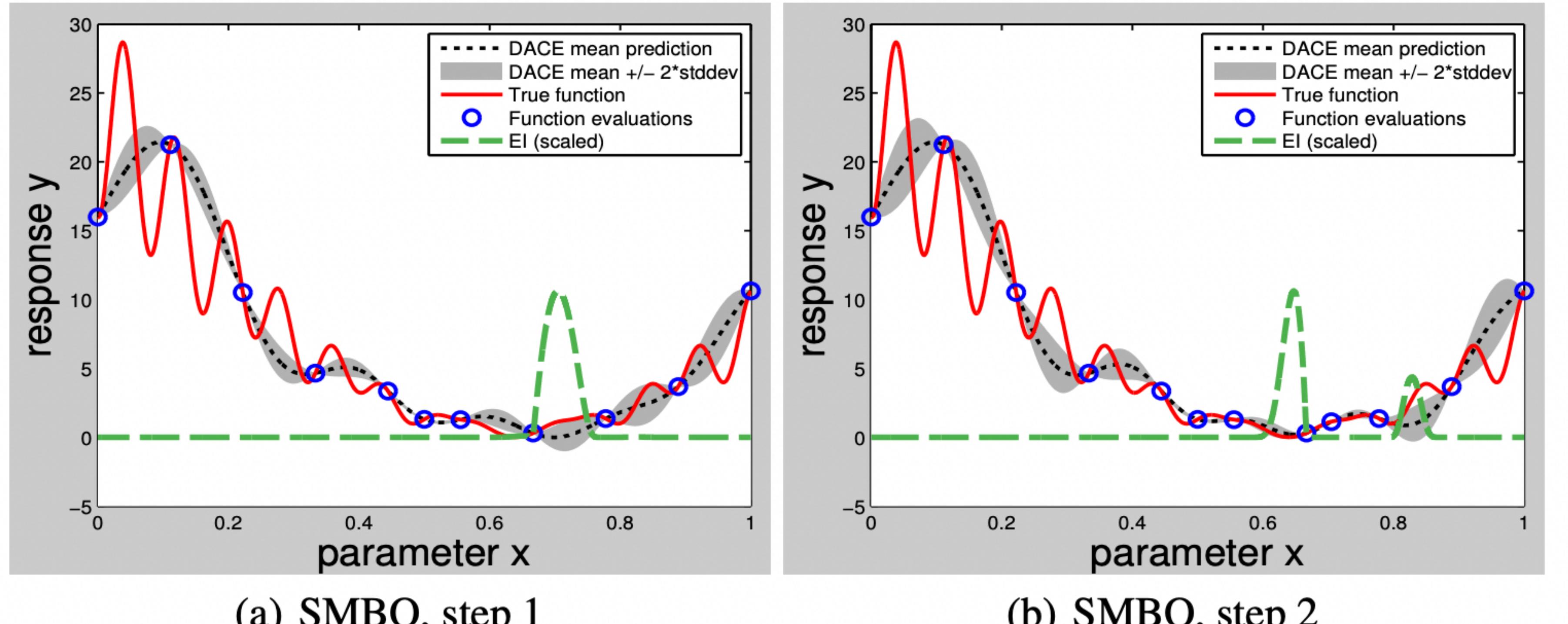
**until** termination condition met;

**return**  $c^*$ ;

**end** *SMBO*

# Sequential Model-Based Optimization

Hutter, Frank, Holger H. Hoos, and Kevin Leyton-Brown. "Sequential model-based optimization for general algorithm configuration." *International conference on learning and intelligent optimization*. Springer, Berlin, Heidelberg, 2011.



**Fig. 1.** Two steps of SMBO for the optimization of a 1D function. The true function is shown as a solid line, and the circles denote our observations. The dotted line denotes the mean prediction of a noise-free Gaussian process model (the “DACE” model), with the grey area denoting its uncertainty. Expected improvement (scaled for visualization) is shown as a dashed line.

# Sequential Model-Based Optimization

**procedure** *SMBO*

**input** target algorithm  $A$ , set of configurations  $C$ , set of problem instances  $I$ ,  
        performance metric  $m$ ;

**output** configuration  $c^*$ ;

    determine initial set of configurations  $C_0 \subset C$ ;

    for all  $c \in C_0$ , measure performance of  $A$  on  $I$  according to metric  $m$ ;

    build initial model  $M$  based on performance measurements for  $C_0$ ;

    determine incumbent  $c^* \in C_0$  for which best performance was observed or predicted;

**repeat**

        based on model  $M$ , determine set of configurations  $C' \subseteq C$ ;

$$EI(x) := E[\max\{f_{min} - \hat{F}(x), 0\}]$$

        for all  $c \in C'$ , measure performance of  $A$  on  $I$  according to metric  $m$ ;

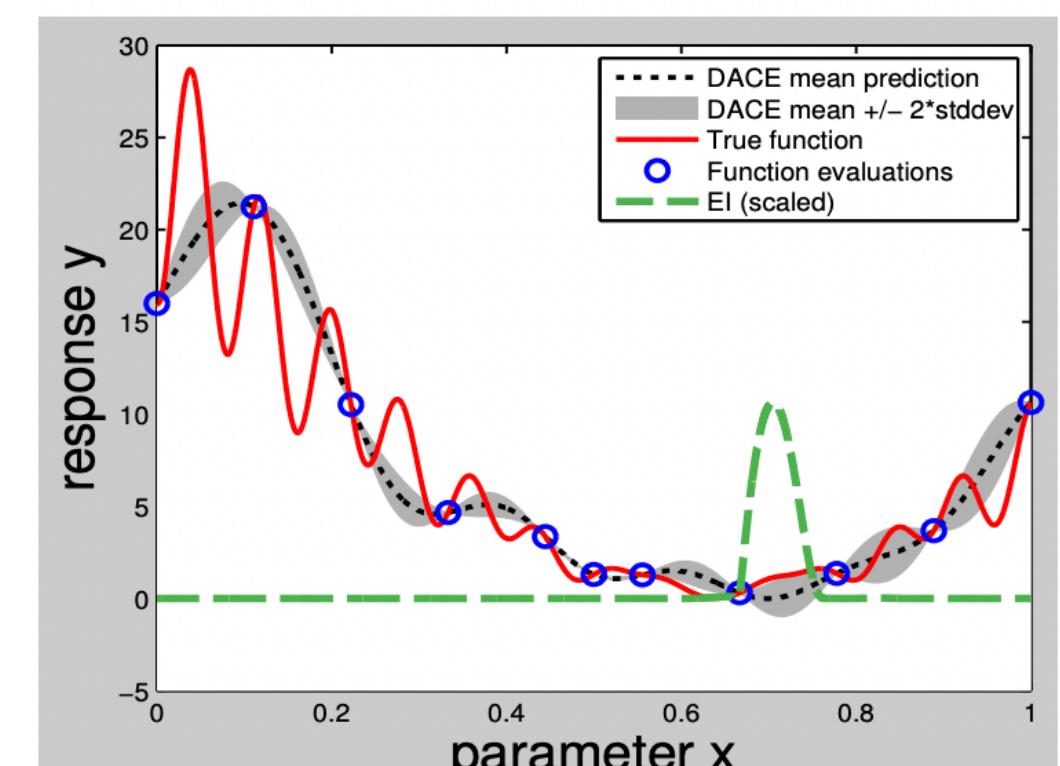
        update model  $M$  based on performance measurements for  $C'$ ;

        update incumbent  $c^*$ ;

**until** termination condition met;

**return**  $c^*$ ;

**end** *SMBO*



# Other flavours of algo. config.

“

- **per-instance algorithm selection** methods choose one of several target algorithms to be applied to a given problem instance based on properties of that instance determined just before attempting to solve it
- **Reactive search procedures**, on-line algorithm control methods and adaptive operator selection techniques switch between different algorithms, heuristic mechanisms and parameter configurations while running on a given problem instance
- **dynamic algorithm portfolio** approaches repeatedly adjust the allocation of CPU shares between algorithms that are running concurrently on a given problem instance

”