

Reviewing Divisibility, Prime, Sigma, and Pi

■ Divisibility:

Let $n, d \in \mathbb{Z}$ be two arbitrary integers.

n is divisible by d if $\exists k \in \mathbb{Z} (n = dk)$.

We use the notation $d | n$, read as " d divides n " to denote that n is divisible by d .

$6 | 30$

Remark: $\forall n \in \mathbb{Z} \ 0 \nmid n$ (note: $0 | 0$ is defined as false)

Remark: $\forall n \in \mathbb{Z}^{\neq 0} \ n | 0$ (since you can choose $k=0$, i.e., $0 = n \times 0$)

Remark: $d \nmid n$ denotes that d does not divide n

$7 \nmid 30$

Prime Numbers:

Defn: $p \in \mathbb{Z}^{>1}$ is a prime number if p is only divisible by $p, (-p), 1$, and (-1) .

Fundamental Theorem of Arithmetic (FTA):

Every $n \in \mathbb{Z}^{>1}$ is prime itself or is the product of unique combination of prime numbers.

(Everybody knows this Theorem & how to use it.
Today, we will prove part of it by Strong induction)

Theorem: There are infinite number of primes
(we will prove it by Contradiction)

■ Sigma (Σ) and Pi (Π)

$$\sum_{i=n}^m a_i = a_n + a_{n+1} + \dots + a_m$$

$$\sum_{i=n}^{m+1} a_i = \sum_{i=n}^m a_i + a_{m+1}$$

$$\sum_{i=n+1}^m a_i = \sum_{i=n}^m a_i - a_n$$

$a_{n+1} + \dots + a_m$ $a_n + a_{n+1} + \dots + a_m$

EX: $\sum_{i=1}^n \frac{1}{i} = \frac{1}{1} + \frac{1}{2} + \dots = \sum_{i=0}^{n-1} \frac{1}{i+1}$

$$\sum_{i=m}^n a_i = \sum_{i=m+k}^{n+k} a_{i-k}$$

$$\sum_{i=m}^n a_i + \sum_{i=m}^n b_i = \sum_{i=m}^n (a_i + b_i)$$

Notation

$$\prod_{i=n}^m a_i = a_n \times a_{n+1} \times \dots \times a_m$$

$$\prod_{i=n}^{m+1} a_i = \prod_{i=n}^m a_i \times a_{m+1}$$

$$\prod_{i=n+1}^m a_i = \prod_{i=n}^m a_i \times \frac{1}{a_n}$$

EX: $\prod_{i=1}^n \sqrt{i^3} = \sqrt{1^3} \cdot \sqrt{2^3} \cdot \dots = \prod_{i=1}^{n-1} \sqrt{i+1^3}$

$$\prod_{i=m}^n a_i = \prod_{i=m+k}^{n+k} a_{i-k}$$

$$\prod_{i=m}^n a_i \times \prod_{i=m}^n b_i = \prod_{i=m}^n (a_i \cdot b_i)$$

Mathematical Induction

■ Principle of Mathematical Induction:

Let $P(n)$ be a statement defined for integer n , and let k be a fixed integer.

Suppose the following are true.

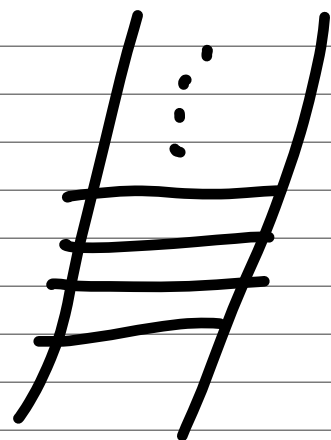
1. $P(k)$ is true

2. For all $m \in \mathbb{Z}$ with $m \geq k$, if $P(m)$ is true, then $P(m+1)$ is true.

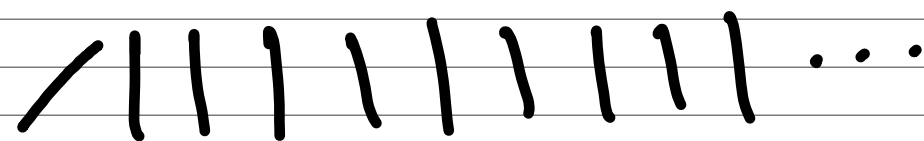
Then, for all $n \geq k$ $P(n)$ is true.

● Metaphor:

Ladder



dominos



$P(1) \rightarrow P(2) \rightarrow P(3) \rightarrow \dots$

Structure of Mathematical Induction Proof:

Goal: prove that $P(n)$ is true $\forall n \geq n_0$.

We proceed by induction.

$P(n)$ is called the induction hypothesis

- Base case ($n = n_0$)

Show that $P(n_0)$ is true.

- Inductive step:

Assume $k \geq n_0$ is arbitrary. Assume $P(k)$ is true.

Show that $P(k+1)$ is true.

This completes the inductive step \square

EX: Prove that $\sum_{i=1}^n i^3 = \left(\frac{n(n+1)}{2}\right)^2$ for $n \geq 1$.

Proof: We proceed by induction. Let $P(n)$ denote the statement that $\sum_{i=1}^n i^3 = \left(\frac{n(n+1)}{2}\right)^2$.

- Base Case ($n=1$):

Observe that $\sum_{i=1}^1 i^3 = 1^3 = 1 = \left(\frac{1(1+1)}{2}\right)^2$.

- Inductive step: Assume $K \geq 1$ is arbitrary.

Assume, $P(K)$ is true, i.e., $\sum_{i=1}^K i^3 = \left(\frac{K(K+1)}{2}\right)^2$.

Observe that $\sum_{i=1}^{K+1} i^3 = \sum_{i=1}^K i^3 + (K+1)^3 = \frac{K^2(K+1)^2}{4} + (K+1)^3$

$$= \frac{K^2(K+1)^2 + 4(K+1)^3}{4} = \frac{(K+1)^2(K^2 + 4K + 4)}{4} = \frac{(K+1)^2(K+2)^2}{4}$$

Hence, $P(K+1)$ holds. This completes the inductive step.

EX: Prove that $\sum_{i=1}^n \frac{1}{i^2} < 2 - \frac{1}{n}$ for $n > 1$.

Strong Induction

Principle of Strong Induction.

Let $P(n)$ be a statement defined for integer n and let k be a fixed integer. Suppose the following is true.

1. $P(k)$

2. for all $m \geq k$, if $P(l)$ is true for all $k \leq l \leq m$, then $P(k+1)$ is true.

Then we can conclude that $P(n)$ is true for all $n \geq k$.

● The process of Strong Induction:

$$P(1) \rightarrow P(2)$$

$$P(1) \wedge P(2) \rightarrow P(3)$$

\vdots

Or equivalently:

$$P(1) \rightarrow P(1) \wedge P(2)$$

$$P(1) \wedge P(2) \rightarrow P(1) \wedge P(2) \wedge P(3)$$

\vdots

$$\left\{ \begin{array}{l} Q(1) \rightarrow Q(2) \\ Q(2) \rightarrow Q(3) \\ \vdots \end{array} \right. \left(Q(n) = P(1) \wedge \dots \wedge P(n) \right) \quad Q(n) = P(1) \wedge \dots \wedge P(n)$$

Strong induction is in fact a regular induction in which the hypothesis of the induction is

Structure of Strong Induction Proof.

We proceed by strong induction.

- Base case ($n = n_0$): show that $p(n_0)$ holds
- Inductive step: Assume $k \geq n_0$ is arbitrary.
Assume $p(l)$ is true for all $n_0 \leq l \leq k$.
Show that $p(k+1)$ holds.
That completes the inductive step \square .

EX. Any integer $n \geq 2$ is a prime or can be written as product of primes.

We proceed by strong induction.

- Base Case ($n = 2$): observe that 2 is a prime

- Inductive step: Assume arbitrary $k \geq 2$.

Assume all integers $2 \leq l \leq k$ are either prime or can be written as product of primes.

We show that $k+1$ is a prime or product of prime through proof by

Cases.

Case 1 ($k+1$ is prim). In this case, $k+1$ is a prime.

Case 2 ($k+1$ is composit): Choose $2 \leq a < k+1$ and $2 \leq b < k+1$ s.t. $k+1 = ab$. By induction step assumption, a and b are either primes or can be written as product of primes.

Choose p_1, \dots, p_i and q_1, \dots, q_j s.t.
 $a = p_1 \times \dots \times p_i$ and $b = q_1 \times \dots \times q_j$.

Observe that $k+1 = ab = p_1 \times \dots \times p_i \times q_1 \times \dots \times q_j$.

So, $k+1$ can be written as product of primes. This completes the inductive step.

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