

Lists and recursion

Functional Programming



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Goals

- ▶ More list functions
- ▶ List comprehensions
- ▶ **Recursion**

Chapters 5 and 6 from Hutton's book



From previous lectures

Primitives for building lists

- ▶ `[]` :: `[a]` is the empty list
- ▶ `(:)` :: `a -> [a] -> [a]` (the “cons” operator)
 - ▶ Build a list by putting an element at the front
- ▶ When we write `[1, 2, 3]` the compiler translates it to
`1 : 2 : 3 : []`

Pattern matching over lists

`length []` = 0

`length (_:xs)` = 1 + `length xs`



From previous lectures

Useful list functions

`null :: [a] -> Bool`

`head :: [a] -> a`

`tail :: [a] -> [a]`

`reverse :: [a] -> [a]`

`(++) :: [a] -> [a] -> [a]`

`sum :: Num a => [a] -> a`

`replicate :: Int -> a -> [a]`



Foldable in the interpreter

If you ask for the type of `sum` in `ghci`, you get

```
sum :: (Foldable t, Num a) => t a -> a
```

- ▶ This is a **more generic** version of `sum`
- ▶ “Adding up all elements” works for other containers
 - ▶ Think of sets or (binary) trees



How to obtain the types I show

```
> :t sum
sum :: (Num a, Foldable t) => t a -> a
> :t +d sum
sum :: [Integer] -> Integer
```



List comprehensions



List comprehensions

```
[ expr | x <- list ]
```

Succinct notation for building **new** lists from **old** ones

```
addone :: Num a => [a] -> [a]
```

```
addone xs = [x + 1 | x <- xs]
```

- ▶ “For each x in xs , return $x + 1$ ”
- ▶ Very similar to mathematical notation

$$\{x + 1 \mid x \in xs\}$$



Guards

```
[ expr | x <- list, condition ]  
  
-- Check if a number is divisible by 2  
even :: Integer -> Bool
```

```
sumeven :: [Integer] -> Integer  
sumeven xs = sum [x | x <- xs, even x]
```

- ▶ “Take all x in xs such that x is even”
- ▶ The result of a comprehension is another list
 - ▶ We can further consume it with other functions
 - ▶ In this case, we use `sum`



Pattern matching

```
[ expr | pattern <- list ]
```

```
heads :: [[a]] -> [a]
```

```
heads xs = [y | (y:_) <- xs]
```

- ▶ Only includes those elements which match the pattern
 - ▶ In this case, non-empty lists
 - > heads [[1,2],[],[3,4,5]]
[1,3]
- ▶ We can introduce new names, as we do with usual pattern matching
 - ▶ In this case, we refer to the head in the result



Multiple clauses

We can have multiple generators and guards

- ▶ Generators provide every possible combination

```
> [(x,y) | x <- [1,2], y <- [3,4]]  
[(1,3),(1,4),(2,3),(2,4)]
```

- ▶ Generators and conditions may refer to each other

```
> [(x,y) | x <- [1,2,3], y <- [1,2,3], x <= y]  
[(1,1),(1,2),(1,3),(2,2),(2,3),(3,3)]
```

```
> [(x,y) | x <- [1,2,3], y <- [x .. 3]]  
[(1,1),(1,2),(1,3),(2,2),(2,3),(3,3)]
```



Prime numbers up to a bound

- **Problem:** Compute all primes $\leq n$



Prime numbers up to a bound

► **Problem:** Compute all primes $\leq n$

1. A number x is a prime iff ($x \geq 2$ and) it has exactly two factors
2. f is a factor of x if the remainder of $\frac{x}{f}$ is zero



Prime numbers up to a bound

► **Problem:** Compute all primes $\leq n$

1. A number x is a prime iff ($x \geq 2$ and) it has exactly two factors
2. f is a factor of x if the remainder of $\frac{x}{f}$ is zero

Good style: divide the problem in parts and refine it

```
primes    :: Int -> [Int]
primes n = [ x | x <- [2 .. n], isPrime x ]
  where isPrime x = _
```



Prime numbers up to a bound

► **Problem:** Compute all primes $\leq n$

1. A number x is a prime iff ($x \geq 2$ and) it has exactly two factors
2. f is a factor of x if the remainder of $\frac{x}{f}$ is zero

Good style: divide the problem in parts and refine it

```
primes    :: Int -> [Int]
primes n = [ x | x <- [2 .. n], isPrime x ]
  where isPrime x = length (factors x) == 2
        factors x = _
```



Prime numbers up to a bound

► **Problem:** Compute all primes $\leq n$

1. A number x is a prime iff ($x \geq 2$ and) it has exactly two factors
2. f is a factor of x if the remainder of $\frac{x}{f}$ is zero

Good style: divide the problem in parts and refine it

```
primes    :: Int -> [Int]
primes n = [ x | x <- [2 .. n], isPrime x ]
  where isPrime x = length (factors x) == 2
        factors x = [f | f <- [1 .. x]
                        , x `mod` f == 0
                      ]
```



(Functional) QuickSort

- ▶ Divide and conquer approach
 1. Pick a pivot
 2. Partition the elements smaller and larger than the pivot
 3. Sort those partitions
 4. Put together the list



(Functional) QuickSort

- ▶ Divide and conquer approach

1. **Pick a pivot**

- ▶ The first element in the list works

2. Partition the elements smaller and larger than the pivot

3. Sort those partitions

4. Put together the list

```
quicksort [] = []
```

```
quicksort (pivot:rest) = undefined
```



(Functional) QuickSort

► Divide and conquer approach

1. Pick a pivot
2. **Partition the elements**
3. Sort those partitions
4. Put together the list

```
quicksort [] = []  
quicksort (pivot:rest) = undefined  
  where smaller = [x | x <- rest, x <= pivot]  
        larger  = [x | x <- rest, x >  pivot]
```



(Functional) QuickSort

- ▶ Divide and conquer approach
 1. Pick a pivot
 2. Partition the elements smaller and larger than the pivot
 3. **Sort those partitions**
 4. **Put together the list**

```
quicksort [] = []  
quicksort (pivot:rest) =  
  quicksort smaller ++ [pivot] ++ quicksort larger  
  where smaller = [x | x <- rest, x <= pivot]  
        larger = [x | x <- rest, x > pivot]
```



Question

Define `replicate` using comprehensions



Question

Define replicate using comprehensions

```
replicate :: Int -> a -> [a]
replicate n x = [x | _ <- [1 .. n]]
```



Recursion



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Recursion on natural numbers

Recursion = defining something in terms of itself

$$\text{fac } 0 = 1$$

$$\text{fac } n = n * \text{fac } (n - 1)$$

$$0 * m = 0$$

$$n * m = m + (n - 1) * m$$

- ▶ A case for 0 or 1
- ▶ A recursive case where the value of n is computed from the same function applied to $n - 1$



Does our product work?

$$0 * m = 0 \quad \text{-- (1)}$$

$$n * m = m + (n - 1) * m \quad \text{-- (2)}$$

$$2 * 4$$

$$= \text{-- apply (2)}$$

$$4 + (2 - 1) * 4$$

$$= \text{-- perform subtraction}$$

$$4 + 1 * 4$$

$$= \text{-- apply (2) and perform subtraction}$$

$$4 + (4 + 0 * 4)$$

$$= \text{-- apply (1)}$$

$$4 + (4 + 0)$$

$$= \text{-- perform additions}$$

$$8$$



Recursion can go wrong

No base case

```
fac n = n * fac (n-1)  -- (1)
-- No more equations
```

```
fac 1
= -- apply (1), what else?
1 * fac 0
= -- apply (1)
1 * 0 * fac (-1)
= -- apply (1)
1 * 0 * (-1) * fac (-2)
= -- apply (1)
...
```



Recursion can go wrong

Argument does not get smaller

```
replicate 0 _ = [] -- (1)
```

```
replicate n x = x : replicate n x -- (2)
```

```
replicate 2 'a'
```

```
= -- apply (2)
```

```
'a' : replicate 2 'a'
```

```
= -- apply (2)
```

```
'a' : 'a' : replicate 2 'a'
```

```
= -- apply (2)
```

```
...
```



Recursion on Lists

```
length []          = 0
length (_ : xs) = 1 + length xs
```

```
[] ++ ys = ys
(x:xs) ++ ys = x : (xs ++ ys)
```



Does our concatenation work?

```
[] ++ ys = ys -- (1)
```

```
(x:xs) ++ ys = x : (xs ++ ys) -- (2)
```

```
[1, 2] ++ [3, 4]
```

```
= -- remove syntactic sugar for [1, 2]
```

```
(1 : 2 : []) ++ [3, 4]
```

```
= -- apply (2)
```

```
1 : ((2 : []) ++ [3, 4])
```

```
= -- apply (2)
```

```
1 : (2 : ([] ++ [3, 4]))
```

```
= -- apply (1)
```

```
1 : 2 : [3, 4]
```

```
= -- resugar the resulting list
```

```
[1, 2, 3, 4]
```



Hutton's recipe for recursion

1. Define the type
2. Enumerate the cases
3. Define the simple (base) cases
4. Define the other (recursive) cases
 - ▶ This part involves most of the thinking
 - ▶ The main question: can I obtain the value of the function if I know its result for a smaller part?
 - ▶ The tail of the list, or $n - 1$ for numbers
5. Generalize and simplify
 - ▶ Remove duplicate equations
 - ▶ Pattern match only as necessary
 - ▶ Infer a more general type



Cooking sum

1. Define the type

```
sum :: [Int] -> Int
```

2. Enumerate the cases

```
sum [] = _
```

```
sum (x:xs) = _
```



Cooking sum

3. Define the simple (base) cases

```
sum [] = 0
```

4. Define the other (recursive) cases

- ▶ If I know the result of `sum xs`, can I get `sum (x:xs)`?
- ▶ Just add the head element to that result!

```
sum (x:xs) = x + sum xs
```

5. Generalize and simplify

- ▶ In this case our definition works for any numeric type

```
sum :: Num a => [a] -> a
```



Cooking elem

`elem x xs` tells you whether `x` is an element of `xs`

```
> 1 `elem` [1,2]
```

```
True
```

```
> 3 `elem` [1,2]
```

```
False
```

```
> 2 `elem` []
```

```
False
```

We usually write `elem` infix to make it look like $1 \in [1, 2]$



Cooking elem

1. Define the (approximate) type

```
elem :: Int -> [Int] -> Bool
```

2. Enumerate the cases

```
elem x [] = _
```

```
elem x (y:ys) = _
```

3. Define the simple (base) cases

```
elem x [] = False
```



Cooking elem

4. Define the other (recursive) cases

- ▶ We need to distinguish between x equal to y or not
 - ▶ Remember: we cannot repeat a variable in a pattern
- ▶ If it is, we stop; otherwise, we continue further

```
elem x (y:ys) | x == y      = True
               | otherwise = elem x ys
```

5. Generalize and simplify

- ▶ We only use $(==)$ to inspect values, so `Eq` is enough

```
elem :: Eq a => a -> [a] -> Bool
```



Cooking take

`take n xs` gets the first `n` elements of list `xs`, or the entire list if there are less than those

```
> take 2 [1,2,3]
```

```
[1,2]
```

```
> take 0 [1,2,3]
```

```
[]
```

```
> take 4 [1,2,3]
```

```
[1,2,3]
```



Cooking take

1. Define the type

- ▶ The type of the elements of the list does not matter

```
take :: Int -> [a] -> [a]
```

2. Enumerate the cases

- ▶ We can match on both the number and list

```
take 0 [] = _
```

```
take 0 (x:xs) = _
```

```
take n [] = _
```

```
take n (x:xs) = _
```



Cooking take

3. Define the simple (base) cases

- ▶ If there are no elements to take, we obtain an empty list

```
take 0 [] = []
```

```
take 0 (x:xs) = []
```

```
take n [] = []
```

4. Define the other (recursive) cases

- ▶ If we have taken 1 element from $x:xs$, there are only $n-1$ left to take from xs

```
take n (x:xs) = x : take (n-1) xs
```



Cooking take

4. We have the following until now

```
take 0 []          = []  
take 0 (x:xs)     = []  
take n []          = []  
take n (x:xs)     = x : take (n-1) xs
```

5. Generalize and simplify

- ▶ When the number is 0, the list does not matter
- ▶ If the list is empty, the number does not matter

```
take 0 _          = []  
take _ []         = []  
take n (x:xs)     = x : take (n-1) xs
```



Question

Define list difference

$(\backslash\backslash) :: \text{Eq } a \Rightarrow [a] \rightarrow [a] \rightarrow [a]$

- Return all elements in the first list **except** if they appear in the second

```
> [1,2] \ \ [1]  
[2]
```

```
> [1,2] \ \ [2,3,4]  
[1]
```

```
> [] \ \ [1,2,3]  
[]
```



Question

Define list difference

$(\backslash\backslash) :: \text{Eq } a \Rightarrow [a] \rightarrow [a] \rightarrow [a]$

- Return all elements in the first list **except** if they appear in the second

```
> [1,2] \ \ [1]  
[2]
```

```
> [1,2] \ \ [2,3,4]  
[1]
```

```
> [] \ \ [1,2,3]  
[]
```

Hint: use `elem` to detect if an element appears in the second



Cooking init

`init xs` gives you all the elements except for the last

```
> init [1,2,3]
```

```
[1,2]
```

```
> init []
```

```
*** Exception: Prelude.init: empty list
```

1. Define the type

```
init :: [a] -> [a]
```

2. Enumerate the cases

► The empty list should yield an error

```
init []      = error "empty list in init"
```

```
init (x:xs) = _
```



Cooking init

- ▶ Here is the trick, we need to distinguish whether we have just one element in the list – and we are finished – or we need to get more elements
 - ▶ We do this by further pattern matching

2. Enumerate the cases

```
init (x:[]) = _  
init (x:xs) = _
```

3. Define the simple (base) cases

```
init (x:[]) = []
```

4. Define the other (recursive) cases

```
init (x:xs) = x : init xs
```



5. Generalize and simplify

- ▶ We can use `[x]` to match a one-element list
- ▶ We do not care about that single element → use `_`

```
init :: [a] -> [a]
init []      = error "empty list in init"
init [_]     = []
init (x:xs)  = x : init xs
```



Cooking sorted

`sorted xs` returns `True` if and only if the elements in the list are in ascending order

```
> sorted [1,2,3]
```

`True`

```
> sorted [2,1,3]
```

`False`

```
> sorted []
```

`True`

1. Define the type

```
sorted :: [Int] -> Bool
```

2. Enumerate the cases

```
sorted []      = _  
sorted (x:xs) = _
```



Cooking sorted

3. Define the simple (base) cases

```
sorted [] = True
```

4. Define the other (recursive) cases

- ▶ We need to compare the first and second elements
- ▶ We need further pattern matching
- ▶ If they are in the right relation, we check further

```
sorted (x:[]) = True
```

```
sorted (x:y:ys) | x <= y = sorted (y:ys)  
                | otherwise = False
```



5. Generalize and simplify

- ▶ As before, we can use `[x]` instead of `x: []`
- ▶ We are reusing the whole `y:ys` in the right-hand side
 - ▶ We can give it a name using `@`
 - ▶ We avoid matching and rebuilding the list

```
sorted [] = True
sorted [_] = True
sorted (x : xs@(y : _))
  | x <= y = sorted xs
  | otherwise = False
```



Cooking zip

`zip xs ys` turns two lists into a list of tuples

```
> zip [1,2] [3,4]
```

```
[(1,3),(2,4)]
```

```
> zip [1,2] [3,4,5]
```

```
[(1,3),(2,4)]
```

If one of the lists runs out of elements, we stop



Cooking zip

`zip xs ys` turns two lists into a list of tuples

```
> zip [1,2] [3,4]
```

```
[(1,3),(2,4)]
```

```
> zip [1,2] [3,4,5]
```

```
[(1,3),(2,4)]
```

If one of the lists runs out of elements, we stop

Try yourself!



Cooking zip

1. Define the type

```
zip :: [a] -> [b] -> [(a,b)]
```

2. Enumerate the cases

```
zip [] [] = _
```

```
zip [] (y:ys) = _
```

```
zip (x:xs) [] = _
```

```
zip (x:xs) (y:ys) = _
```

3. Define the simple (base) cases

```
zip [] [] = []
```

```
zip [] (y:ys) = []
```

```
zip (x:xs) [] = []
```



Cooking zip

4. Define the other (recursive) cases

```
zip (x:xs) (y:ys) = (x,y) : zip xs ys
```

5. Generalize and simplify

► If one of the lists is empty, we don't care about the other

```
zip :: [a] -> [b] -> [(a,b)]
```

```
zip [] _ = []
```

```
zip _ [] = []
```

```
zip (x:xs) (y:ys) = (x,y) : zip xs ys
```



Cooking merge

Given two **sorted** lists `xs` and `ys`, `merge xs ys` produces a new sorted list from those elements

► This is the basis of a sorting algorithm called MergeSort

```
> merge [1,4] [2,3,5]
[1,2,3,4,5]
> merge [] [2,3,5]
[2,3,5]
```



Cooking merge

1. Define the type

```
merge :: [Int] -> [Int] -> [Int]
```

2. Enumerate the cases

```
merge [] [] = _
```

```
merge (x:xs) [] = _
```

```
merge [] (y:ys) = _
```

- In the last case we have to decide which number is larger

```
merge (x:xs) (y:ys)
```

```
  | x <= y = _
```

```
  | otherwise = _
```



Cooking merge

3. Define the simple (base) cases

```
merge [] [] = []  
merge (x:xs) [] = x:xs  
merge [] (y:ys) = y:ys
```

4. Define the other (recursive) cases

- Choose the smallest one and merge the rest

```
merge (x:xs) (y:ys)  
  | x <= y      = x : merge xs (y:ys)  
  | otherwise   = y : merge (x:xs) ys
```



5. Generalize and simplify

- ▶ This function works for any type which can be ordered
- ▶ In the case of an empty list, we just return the other list
- ▶ We can give names to complete lists to avoid duplication

```
merge :: Ord a => [a] -> [a] -> [a]
merge [] ys      = ys
merge xs []      = xs
merge xss@(x:xs) yss@(y:ys)
  | x <= y        = x : merge xs yss
  | otherwise     = y : merge xss ys
```



Cooking reverse

`reverse xs` gives the same elements in reverse order

```
> reverse [1,2,3]
[3,2,1]
```

1. Define the type

```
reverse :: [a] -> [a]
```

2. Enumerate the cases

```
reverse []      = _
reverse (x:xs) = _
```



Cooking reverse

3. Define the simple (base) cases

```
reverse [] = []
```

4. Define the other (recursive) cases

- ▶ Suppose you get $[1, 2, 3]$, which you split as 1 and $[2, 3]$
- ▶ The reverse of $[2, 3]$ is $[3, 2]$, where do you put the 1?
- ▶ At the end of the reversed list!

```
reverse (x:xs) = reverse xs ++ [x]
```



Problem with reverse reverse

- ▶ This definition is **very inefficient**
 - ▶ Each time you call `(++)`, you need to traverse the whole list, since the new element goes at the end
 - ▶ If the list has n elements, the amount of steps is

$$n - 1 + n - 2 + n - 3 + \dots + 1 = \frac{n \cdot (n - 1)}{2} = \mathcal{O}(n^2)$$



Solution: use an accumulator

- ▶ There is a standard technique to solve this problem:
using an **accumulator**
 1. Introduce a local definition with an additional parameter
(the accumulator)
invariant: accumulator contains solution for all
elements seen so far.
 2. Initialize the accumulator in the main call
 3. Follow Hutton's recipe, but
 - ▶ Do not pattern match on the accumulator
 - ▶ Return the accumulator in the base case
 - ▶ Update the accumulator in the recursive steps



sum with accumulator

Define `sum` using an accumulator



sum with accumulator

Define `sum` using an accumulator

```
sum [1,2,3,4] = 1 + sum [2,3,4]
               = 1 + 2 + sum [3,4]
               = 1 + 2 + 3 + sum [4]
               = 1 + 2 + 3 + 4 + sum []
```



sum with accumulator

Define `sum` using an accumulator

```
sum [1,2,3,4] = 1 + sum [2,3,4]
              = 1 + 2 + sum [3,4]
              = 1 + 2 + 3 + sum [4]
              = 1 + 2 + 3 + 4 + sum []
```

- **Observation:** Always of the form 'a + sum xs'
- Introduce the function `sum'` that has as invariant:

```
sum' acc xss = acc + sum xs
```



Implementing `sum'`

► invariant: `'sum' acc xs = acc + sum xs`

```
sum' :: Int -> [Int] -> Int
sum' acc [] = _
sum' acc (x:xs) = _
```



Implementing `sum'`

► invariant: `'sum' acc xs = acc + sum xs`

```
sum' :: Int -> [Int] -> Int
sum' acc [] = _
sum' acc (x:xs) = _
```

Invariant tells us that:

```
sum' :: Int -> [Int] -> Int
sum' acc [] = acc
sum' acc (x:xs) = sum' (acc + x) xs
```



Implementing `sum'`

► invariant: `'sum' acc xs = acc + sum xs`

```
sum'           :: Int -> [Int] -> Int
sum' acc []    = _
sum' acc (x:xs) = _
```

Invariant tells us that:

```
sum'           :: Int -> [Int] -> Int
sum' acc []    = acc
sum' acc (x:xs) = sum' (acc + x) xs
```

SO:

```
sum    :: [Int] -> Int
sum xs = sum' 0 xs
```



sum with accumulator

Define `sum` using an accumulator.

We can also apply η -reduction and use a `case` expression.

```
sum :: [Int] -> Int
sum = sum' 0
  where
    sum'      :: Int -> [Int] -> Int
    sum' acc xs = case xs of
      []      -> acc
      (x:xs)  -> sum' (acc+x) xs
```



reverse with an accumulator

1. Introduce a local definition with an additional parameter to hold the interim result

```
reverse xs = _  
  where  
    reverse'      :: [a] -> [a] -> [a]  
    reverse' acc xs = _
```

2. Initialize the accumulator in the main call

► When we start, we haven't accumulated any element yet

```
reverse xs = reverse' [] xs  
  where  
    reverse' acc xs = _
```



reverse with an accumulator

3. Follow Hutton's recipe, but

- ▶ Do not pattern match on the accumulator
- ▶ Return the accumulator in the base case
- ▶ Update the accumulator in the recursive steps

```
reverse xs = reverse' [] xs
  where
    reverse' acc []      = acc
    reverse' acc (x:xs) = reverse' (x:acc) xs
```



reverse with an accumulator

```
reverse xs = reverse' [] xs
  where
    reverse' acc []      = acc
    reverse' acc (x:xs) = reverse' (x:acc) xs
```



Cooking initial segments

`inits xs` returns the initial segments of `xs`, that is, all the lists which are prefixes of the original one

```
> inits [1,2,3]
[[],[1],[1,2],[1,2,3]]
> inits []
[[]]
```

1. Define the type

```
inits :: [a] -> [[a]]
```

2. Enumerate the cases

```
inits []      = _
inits (x:xs) = _
```



Cooking initial segments

3. Define the simple (base) cases

```
inits [] = [[]]
```

4. Define the other (recursive) cases

- ▶ Suppose you have $[1, 2, 3]$, that is, $1 : [2, 3]$
- ▶ The initial segments of $[2, 3]$ are $[], [2], [2, 3]$, what do you do with the 1?
- ▶ If you put the 1 in front of every list, you get $[1], [1, 2], [1, 2, 3]$
- ▶ We are almost there! We are just missing the extra empty list at the front

```
inits (x:xs) = [] : [x:rs | rs <- inits xs]
```



Cooking final segments

`tails xs` returns the final segments of `xs`, that is, all the lists which are suffixes of the original one

```
> tails [1,2,3]
[[1,2,3],[2,3],[3],[]]
> tails [2,3]
[[2,3],[3],[]]
> tails [3]
[[3],[]]
```

```
tails :: [a] -> [[a]]
tails [] = [[]]
tails ts@(_:xs) = ts : tails xs
```



Final segments using initial segments

Final segments of `xs` seem related to initial segments of `reverse xs`

```
> tails [1,2,3]
[[1,2,3],[2,3],[3],[]]
> inits [3,2,1]
[[],[3],[3,2],[3,2,1]]
```

- ▶ There are two problems with the second result
 1. Each of the inner lists is reversed
 2. The whole outer list is reversed
- ▶ Let's fix this and give an alternative definition of `tails`



Final segments using initial segments

- ▶ To reverse **each** of the inner lists we use a list comprehension

```
> [reverse i | i <- inits [3,2,1]]  
[[], [3], [2,3], [1,2,3]]
```

- ▶ This leads to this final definition

```
tails xs = reverse [reverse i  
                    | i <- inits (reverse xs)]
```



Fizzbuzz

- ▶ Write `fizzbuzz` using direct recursion; test if some element is divisible by `n` (and by `m`) only once.

```
fizzbuzz :: (Int, Int) -> [Int]
          -> ([Int], [Int], [Int])
```

A call of the form `fizzbuzz (m, n) xs` should return a triple with a list in each element:

- ▶ The first list contains elements of `xs` divisible by `m`
- ▶ The second list those divisible by `n` (and not by `m`)
- ▶ The third list should contain the rest



Fizzbuzz

```
fizzbuzz (m,n) xs = fb xs
  where
    fb []      = ([],[],[])
    fb (x:xs) = case ( x `mod` m == 0
                        , x `mod` n == 0
                      ) of
                        (True, _   ) -> (x:ms,ns,  rs)
                        (_   , True) -> (ms,  x:ns,rs)
                        (_   , _   ) -> (ms,  ns,  x:rs)
    where
      (ms,ns,rs) = fb xs
```



Fizzbuzz

```
fizzbuzz (m,n) xs = fb xs
  where
    fb []      = ([],[],[])
    fb (x:xs) = case ( x `mod` m == 0
                      , x `mod` n == 0
                      ) of
      (True, _   ) -> (x:ms,ns,  rs)
      (_,  True) -> (ms,  x:ns,rs)
      (_,  _   ) -> (ms,   ns,  x:rs)
    where
      (ms,ns,rs) = fb xs
```

- Exercise: write fizzbuzz using a comprehensions



Final words

Defining recursive functions is like riding a bicycle: it looks easy when someone else is doing it, may seem impossible when you first try to do it yourself, but becomes simple and natural with practice.

– From "Programming in Haskell"

- ▶ On the other hand, don't get too attached to recursion
- ▶ Many of these examples have better implementations using **higher-order functions**
 - ▶ Which happens to be the topic for next lecture!

