

Lecture 12. Property Testing with QuickCheck

Functional Programming



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Why testing?

- ▶ Gain confidence in the correctness of your program
- ▶ Show that common cases work correctly
- ▶ Show that **corner** cases work correctly



Why testing?

- ▶ Gain confidence in the correctness of your program
- ▶ Show that common cases work correctly
- ▶ Show that **corner** cases work correctly

Testing cannot prove the absence of bugs



When is a program correct?



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When is a program correct?

- ▶ When it satisfies the specification



When is a program correct?

- ▶ When it satisfies the specification
- ▶ What is a specification?
- ▶ How to establish a relation between the specification and the implementation?
- ▶ What about bugs in the specification?

More in **Software Testing and Verification**, period 4



Property Testing using QuickCheck

QuickCheck, an **automated** testing library/tool for Haskell

- ▶ Describe properties as Haskell programs using an embedded domain-specific language (EDSL)
- ▶ Automatic datatype-driven random test case generation
- ▶ Extensible, e.g. test case generators can be adapted
 - ▶ A default generator for list generates any list, but you may want only sorted lists



Case study: insertion sort



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A buggy insertion sort

```
sort :: [Int] -> [Int]
sort []      = []
sort (x:xs)  = insert x xs
```

```
insert :: Int -> [Int] -> [Int]
insert x []                = [x]
insert x (y:ys) | x <= y   = x : ys
                | otherwise = y : insert x ys
```

Let's try to debug it using QuickCheck



How to write a specification?

A good specification is

- ▶ as precise as necessary
- ▶ but no more precise than necessary

A good specification for a particular problem, such as sorting, should:

1. distinguish sorting from all other operations on lists,
2. without forcing us to use a particular sorting algorithm



A first approximation

Certainly, sorting a list should not change its length

```
sortPreservesLength :: [Int] -> Bool
sortPreservesLength xs =
    length (sort xs) == length xs
```

We can test by invoking the function:

```
> quickCheck sortPreservesLength
Failed! Falsifiable, after 4 tests:
[0,3]
```

QuickCheck gives back a **counterexample**



Correcting the bug

```
sort :: [Int] -> [Int]
sort []      = []
sort (x:xs)  = insert x xs
```

```
insert :: Int -> [Int] -> [Int]
insert x []                        = [x]
insert x (y:ys) | x <= y          = x : ys
                  | otherwise     = y : insert x ys
```

Which branch does not preserve the list length?



A new attempt

```
> quickCheck sortPreservesLength  
OK, passed 100 tests.
```

Looks better. But have we tested enough?



Properties are first-class objects

```
preserves :: Eq b => (a -> a) -> (a -> b)
           -> a -> Bool
(algo `preserves` prop) x = prop (algo x) == prop x

sortPreservesLength = sort `preserves` length
```



Properties are first-class objects

```
preserves :: Eq b => (a -> a) -> (a -> b)
           -> a -> Bool
(algo `preserves` prop) x = prop (algo x) == prop x
```

```
sortPreservesLength = sort `preserves` length
```

```
idPreservesLength   = id `preserves` length
```

id also preserves the lists length:

```
> quickCheck idPreservesLength
OK, passed 100 tests.
```

So we need to refine our specification



When is a list sorted?

We can define a predicate that checks if a list is sorted:

```
isSorted :: [Int] -> Bool
isSorted []      = True
isSorted [x]     = True
isSorted (x:y:xs) = x < y && isSorted (y:xs)
```

And use this to check that sorting a list produces a list that isSorted



Testing again

```
sortEnsuresSorted :: [Int] -> Bool  
sortEnsuresSorted xs = isSorted (sort xs)
```

```
> quickCheck sortEnsuresSorted  
Falsifiable, after 5 tests:  
[5,0,-2]  
> sort [5,0,-2]  
[0,-2,5]
```

We're still not quite there...



Debugging sort

What's wrong now?

```
sort :: [Int] -> [Int]
sort []      = []
sort (x:xs)  = insert x xs
```

```
insert :: Int -> [Int] -> [Int]
```



Debugging sort

What's wrong now?

```
sort :: [Int] -> [Int]
sort []      = []
sort (x:xs)  = insert x xs
```

```
insert :: Int -> [Int] -> [Int]
```

We are not recursively sorting the tail in sort!



Another bug

```
> quickCheck sortEnsuresSorted  
Falsifiable, after 7 tests:  
[4,2,2]  
> sort [4,2,2]  
[2,2,4]
```

This is correct. What is wrong?



Another bug

```
> quickCheck sortEnsuresSorted
Falsifiable, after 7 tests:
[4,2,2]
> sort [4,2,2]
[2,2,4]
```

This is correct. What is wrong?

```
> isSorted [2,2,4]
False
```



Fixing the specification

The `isSorted` specification reads:

```
sorted :: [Int] -> Bool
sorted []           = True
sorted (x:[])       = True
sorted (x:y:ys)     = x < y && sorted (y : ys)
```

Why does it return `False`? How can we fix it?



Are we done yet?

Is sorting specified completely by saying that

- ▶ sorting preserves the length of the input list,
- ▶ the resulting list is sorted?



Are we done yet?

Is sorting specified completely by saying that

- ▶ sorting preserves the length of the input list,
- ▶ the resulting list is sorted?

Not really...

```
evilNoSort :: [Int] -> [Int]  
evilNoSort xs = replicate (length xs) 1
```

This function fulfills both specifications, but does not sort



Specifying sorting

```
permutes :: ([Int] -> [Int]) -> [Int] -> Bool  
permutes f xs = f xs `elem` permutations xs
```

```
sortPermutes :: [Int] -> Bool  
sortPermutes xs = sort `permutes` xs
```

This completely specifies sorting and our algorithm passes the corresponding tests



QuickCheck in general



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The type of quickCheck

The type of is an **overloaded** type:

```
quickCheck :: Testable prop => prop -> IO ()
```

- ▶ The argument of is a property of type `prop`
- ▶ The only restriction on the type is that it is in the `Testable` **type class**.
- ▶ When executed, prints the results of the test to the screen – hence the `IO ()` result type.



Which properties are Testable?

So far, all our properties have been of type:

```
sortPreservesLength :: [Int] -> Bool  
sortEnsuresSorted  :: [Int] -> Bool  
sortPermutes       :: [Int] -> Bool
```

When used on such properties, QuickCheck generates random integer lists:

- ▶ If the result is `True` for 100 cases, this success is reported in a message
- ▶ If the result is `False` for a test case, the input triggering the failure is printed



Other example properties

```
appendLength :: [Int] -> [Int] -> Bool
appendLength xs ys =
    length xs + length ys == length (xs ++ ys)
```

```
plusIsCommutative :: Int -> Int -> Bool
plusIsCommutative m n = m + n == n + m
```

```
takeDrop :: Int -> [Int] -> Bool
takeDrop n xs = take n xs ++ drop n xs == xs
```

```
dropTwice :: Int -> Int -> [Int] -> Bool
dropTwice m n xs =
    drop m (drop n xs) == drop (m + n) xs
```



Other forms of properties – contd.

```
> quickCheck takeDrop  
OK, passed 100 tests.
```

```
> quickCheck dropTwice  
Falsifiable after 7 tests.  
1  
-1  
[0]
```

```
> drop (-1) [0]  
[0]
```

```
> drop 1 (drop (-1) [0])  
[]
```



Nullary properties

A property without arguments is also possible:

```
lengthEmpty :: Bool
lengthEmpty = length [] == 0
wrong :: Bool
wrong = False
```

```
> quickCheck lengthEmpty
```

OK, passed 100 tests.

```
> quickCheck wrong
```

Falsifiable, after 0 tests.

QuickCheck subsumes unit tests



Properties

Recall the type of `quickCheck`:

```
quickCheck :: Testable prop => prop -> IO ()
```

We can now say more about when types are `Testable`:

- ▶ testable properties usually are functions (with any number of arguments) resulting in a `Bool`

What argument types are admissible?

- ▶ `QuickCheck` has to know how to produce random test cases of such types



Properties – continued

A Testable thing is something which can be turned into a Property:

```
class Testable prop where
  property :: prop -> Property
```

A Bool is testable:

```
instance Testable Bool where ...
```

If a type is testable, we can add a function argument, as long as we know how to generate and print test cases:

```
instance (Arbitrary a, Show a, Testable b) =>
  Testable (a -> b) where
```



Information about test data

We can show the actual data that is tested:

```
> quickCheck (\ xs -> collect xs (sPL xs))  
OK, passed 100 tests:  
6% []  
1% [9,4,-6,7]  
1% [9,-1,0,-22,25,32,32,0,9,...  
...
```

Why is it important to have access to the test data?



Implications

The function insert preserves an ordered list:

```
implies :: Bool -> Bool -> Bool
```

```
implies x y = not x || y
```

```
insertPreservesOrdered :: Int -> [Int] -> Bool
```

```
insertPreservesOrdered x xs =  
    sorted xs `implies` sorted (insert x xs)
```



Implications – contd.

```
> quickCheck insertPreservesOrdered  
OK, passed 100 tests.
```

But:

```
> let iPO = insertPreservesOrdered  
> quickCheck (\x xs -> collect (sorted xs)  
                                     (iPO x xs))
```

OK, passed 100 tests.

88% False

12% True

For **88** test cases, insert has not actually been relevant!



Implications – contd.

The solution is to use the QuickCheck implication operator:

```
(==>) :: Testable prop => Bool -> prop -> Property
```

```
iP0 :: Int -> [Int] -> Property
```

```
iP0 x xs = sorted xs ==> sorted (insert x xs)
```

Now, lists that are not sorted are discarded and do not contribute towards the goal of 100 test cases



Implications – contd.

We can now easily run into a new problem:

```
iP0 :: Int -> [Int] -> Property
iP0 x xs = length xs > 2 && sorted xs ==>
           sorted (insert x xs)
```

We try to ensure that lists are not too short, but:

```
> quickCheck (\x xs -> collect (sorted xs)
                               (iP0 x xs))
```

Arguments exhausted after 20 tests (100% True).

The chance that a random list is sorted is extremely small



Custom generators



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Generators

- ▶ Generators belong to an abstract data type `Gen`
 - ▶ The only effect available to us is access to random numbers
 - ▶ Think of as a restricted version of IO
- ▶ We can define our own generators using another domain-specific language
 - ▶ The default generators for datatypes are specified by defining instances of class `Arbitrary`

```
class Arbitrary a where  
  arbitrary :: Gen a  
  ...
```



Generator combinators

```
choose      :: Random a => (a,a) -> Gen a
oneof       :: [Gen a] -> Gen a
frequency   :: [(Int, Gen a)] -> Gen a
elements    :: [a] -> Gen a
sized       :: (Int -> Gen a) -> Gen a
```



Simple generators

```
instance Arbitrary Bool where
    arbitrary = choose (False, True)

instance (Arbitrary a, Arbitrary b)
    => Arbitrary (a,b) where
    arbitrary = do x <- arbitrary
                  y <- arbitrary
                  return (x,y)
    -- arbitrary = (,) <$> arbitrary <*> arbitrary

data Dir = North | East | South | West
instance Arbitrary Dir where
    arbitrary = elements [North, East, South, West]
```



Generating random numbers

- ▶ A simple possibility:

```
instance Arbitrary Int where  
    arbitrary = choose (-20,20)
```

- ▶ Better:

```
instance Arbitrary Int where  
    arbitrary = sized (\n -> choose (-n,n))
```

- ▶ QuickCheck automatically increases the size gradually



How to generate sorted lists

Idea: Adapt the default generator for lists

The following function turns a list of integers into a sorted list of integers:

```
mkSorted :: [Int] -> [Int]
mkSorted []           = []
mkSorted [x]          = [x]
mkSorted (x:y:ys) = x : mkSorted ((x + abs y : ys))
```

For example:

```
> mkSorted [1,2,-3,4]
[1,3,6,10]
```



Random generator

The generator can be adapted as follows:

```
genSorted :: Gen [Int]
genSorted = do xs <- arbitrary
              return (mkSorted xs)
-- genSorted = mkSorted <$> arbitrary
```



Using a custom generator

There is another function to construct properties provided by QuickCheck, passing an explicit generator:

```
forall :: (Show a, Testable b)
        => Gen a -> (a -> b) -> Property
```

This is how we use it:

```
iP0 :: Int -> Property
iP0 x = forall genSorted
        (\xs -> length xs > 2 && sorted xs ==>
            sorted (insert x xs))
```



Shrinking

The other method in `Arbitrary` is:

```
shrink :: (Arbitrary a) => a -> [a]
```

- ▶ Maps each value to structurally smaller values
 - ▶ `[2,3]` is structurally smaller than `[1,2,3]`
- ▶ When a failing test case is discovered, QuickCheck shrinks repeatedly until no smaller failing test case can be obtained



Loose ends

- ▶ Haskell can deal with infinite values, and so can QuickCheck
 - ▶ Properties must **not** inspect infinitely many values
 - ▶ Solution: only inspect finite parts
- ▶ QuickCheck can also generate functional values
 - ▶ Requires defining an instance of another class `Coarbitrary`
 - ▶ Showing functional values is still problematic
- ▶ QuickCheck has facilities for testing properties that involve IO



Summary

QuickCheck is a great tool:

- ▶ A domain-specific language for writing properties
- ▶ Test data is generated automatically and randomly
- ▶ Another domain-specific language to write custom generators

However, keep in mind that writing good tests still requires practice, and that tests can have bugs, too



Correctness



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Correctness as a goal

Testing can**not** prove the absence of bugs

- ▶ Only point at failing cases

Are there ways to prove your code correct?



Equational reasoning

1. Write a bunch of properties that specify your algorithm
2. Prove that they hold using equational reasoning
3. You are done!



Equational reasoning

1. Write a bunch of properties that specify your algorithm
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Caveats

- ▶ Time-consuming, needs lots of manual work
- ▶ Laziness and exceptions are not taken care of
 - ▶ Proofs only work for finite values



Interactive theorem proving

Help you proving properties about your program

- ▶ Check that every inference step is correct
- ▶ Fill in boring and obvious proofs

Some interactive theorem provers:

- ▶ Coq (blame the French for the name!)
- ▶ Isabelle/HOL



More expressive types

Define the type of your function in such a way that only correct implementations are allowed

```
append :: List n a -> List m a -> List (n + m) a
```

1. Dependent types

- ▶ Allow values to appear in types
- ▶ Examples: Agda, Idris, Coq

2. Refinement types

- ▶ Attach predicates to types
- ▶ Example: LiquidHaskell



More expressive types

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 - ▶ Examples: Agda, Idris, Coq
2. Refinement types
 - ▶ Attach predicates to types
 - ▶ Example: LiquidHaskell

Learn about them in Advanced FP!



Theorems for free

How many implementations are of these signatures?

$f :: a \rightarrow a$

$g :: (a, b) \rightarrow (b, a)$



Theorems for free

How many implementations are of these signatures?

```
f :: a -> a
```

```
g :: (a, b) -> (b, a)
```

Only one!

```
f x      = x      -- identity function
```

```
g (x, y) = (y, x) -- swap pair
```

Types are enough to determine many properties of the implementation

► We call those **free theorems**

