

Farewell

Functional Programming 2018/19

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[Faculty of **Science**
Information and Computing **Sciences**]

Final menu

- ▶ Presentations about libraries and structures
- ▶ Q&A session
 - ▶ Ask me more in the break
- ▶ Closing remarks



Participation on research

We would like to gather your DomJudge assignments to perform research on programming education

- ▶ You do not need to do anything else than allowing us to look at the assignments
- ▶ Assignments are anonymized before anybody looks at it

Participation is **completely optional**

- ▶ During the exam I will give you a *formulier toestemming*, which you need to sign if you agree to help us



Presentations

- ▶ Parallelism with `monad-par`
- ▶ Semirings



Q&A session



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Things with symbols: \$, <\$>, <*>, <|>, >>=

$(\$)$:: $(a \rightarrow b) \rightarrow a \rightarrow b$

Rule 1: writing \$ is like writing a parentheses until the end of the expression

Rule 2: nested uses of \$ create nested parentheses

```
f lst = map (+1) (filter even (map read lst))  
-- Too many parentheses!  
f lst = map (+1) $ filter even $ map read lst  
-- In this case, better use composition  
f      = map (+1) . filter even . map read
```



Things with symbols: <\$>, <*>

-- *From Functor and Applicative*

```
fmap    :: (a -> b) -> f a -> f b  
(<$>)   :: (a -> b) -> f a -> f b  
(<*>)   :: f (a -> b) -> f a -> f b
```

If we have a bunch of arguments inside a context/functor/monad

```
x :: m a, y :: m b, z :: m c ...
```

and we want to apply a *pure* function

```
f :: a -> b -> ... -> r
```

we need to *lift* the function using a combination of them

```
f <$> x <*> y <*> z <*> ... :: m r
```



Things with symbols: <\$>, <*>

```
mapM :: (a -> m b) -> [a] -> m [b]
mapM _ []      = return []
mapM f (x:xs) = (:) <$> f x <*> mapM f xs
               ---      ---      -----
               --      m b      m [b]
               --
               -- (:) :: b -> [b] -> [b]
```



Things with symbols: <|>

<|> or `mp1us` model the idea of *trying different possibilities*

- ▶ In the case of `[]`, concatenate all solutions
 $(\langle | \rangle) = (++)$
- ▶ In the case of `Maybe`, get the first one which doesn't fail
 - ▶ That is, obtain the first `Just`

`Just x` $\langle | \rangle$ `_` = `Just x`

`Nothing` $\langle | \rangle$ `y` = `y`



Things with symbols: >>=

(>>=) is the *bind* operation of a monad

```
class Monad m where
  return :: a -> m a
  (>>=)  :: m a -> (a -> m b) -> m b
```

- ▶ If you give me an *a* inside a monad
- ▶ And tell me how to continue if I “unwrap” the *a* for you
- ▶ Then I can apply the continuation to the value for you

```
do x <- thing    ===    thing >>= \x ->
  continue                               continue
```



Monoids

Monoid is the generalization of the properties exhibited by $+$, \times , list concatenation, \vee , \wedge , ...

- ▶ A binary operation `mappend` or $(\langle \rangle)$
- ▶ Which is associative

$$(x \langle \rangle y) \langle \rangle z === x \langle \rangle (y \langle \rangle z)$$

- ▶ And has a neutral element `mempty`

$$x \langle \rangle \text{mempty} = x$$

$$\text{mempty} \langle \rangle x = x$$

- ▶ 0 for $+$, 1 for \times , [] for $(++)$, False for \vee , True for \wedge



Monoids and monads

Are monoids and monads related in any way?

Yes, they are deeply connected.



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Yes, they are deeply connected.

But we don't have time to explain how now.

You need to look at *category theory*, a branch of mathematics (and computer science).

```
return :: a -> m a           -- is like `mempty`  
join   :: m (m a) -> m a    -- is like `mappend`
```



Lecture 10, exercise 1

Finish the proof of $\text{reverse} \circ \text{reverse} = \text{id}$

We need the following lemmas:

-- *Distributivity of (++) over reverse*

$\text{reverse} \ (xs \ ++ \ ys) = \text{reverse} \ ys \ ++ \ \text{reverse} \ xs$

-- *Reverse on singleton lists*

$\text{reverse} \ [x] = [x]$

The second one is simple equational reasoning.



Lecture 10, exercise 1

`reverse (xs ++ ys) = reverse ys ++ reverse xs`

By induction on `xs`:



Lecture 10, exercise 1

`reverse (xs ++ ys) = reverse ys ++ reverse xs`

By induction on `xs`:

► Case `xs = []`

```
reverse ([] ++ ys)
= {- defn of (++) -}
reverse ys
```

```
reverse ys ++ reverse []
= {- defn of reverse -}
reverse ys ++ []
= {- WE ARE STUCK AGAIN!!! -}
```

We need to prove a separate lemma: `ts ++ [] = ts`



Lecture 10, exercise 1

► Case $xs = z:zs$

► IH: $\text{reverse } (zs ++ ys) = \text{reverse } ys ++ \text{reverse } zs$

```
reverse ((z:zs) ++ ys)
= {- defn of (++) -}
reverse (z : (zs ++ ys))
= {- defn of reverse -}
reverse (zs ++ ys) ++ [z]
```

```
reverse ys ++ reverse (z:zs)
= {- defn of reverse -}
reverse ys ++ (reverse zs ++ [z])
= {- associativity of (++) -}
(reverse ys ++ reverse zs) ++ [z]
= {- IH -}
reverse (zs ++ ys) ++ [z]
```



Lecture 10, exercise 1

`ts ++ [] = ts`

By induction on `ts`:



Lecture 10, exercise 1

$ts \ ++ \ [] = ts$

By induction on ts :

► Case $ts = []$:

$[] \ ++ \ []$
 $= \{- \text{defn of } (++) \ -\}$
 $[]$

► Case $ts = p:ps$

► IH: $ps \ ++ \ [] = ps$

$(p:ps) \ ++ \ []$
 $= \{- \text{defn of } (++) \ -\}$
 $p : (ps \ ++ \ [])$
 $= \{- \text{IH} \ -\}$
 $p : ps$



Lectures 12/13, exercise 2

Define a function:

```
tuple :: Monad m => m a -> m b -> m (a, b)
```

using explicit ($\gg=$), `do`-notation and applicative operators.



Lectures 12/13, exercise 2

```
tuple :: Monad m => m a -> m b -> m (a, b)
```

Using do-notation

```
tuple x y = do x' <- x  
              y' <- y  
              return (x', y')
```



Lectures 12/13, exercise 2

```
tuple :: Monad m => m a -> m b -> m (a, b)
```

Using do-notation

```
tuple x y = do x' <- x  
              y' <- y  
              return (x', y')
```

Using (>>=)

```
tuple x y = x >>= \x' ->  
              y >>= \y' ->  
              return (x', y')
```



Lectures 12/13, exercise 2

```
tuple :: Monad m => m a -> m b -> m (a, b)
```

Using applicative operators

We need to find a function $a \rightarrow b \rightarrow (a, b)$ and lift it

```
pair x y = (x, y)           -- define it yourself  
(,) :: a -> b -> (a, b)    -- constructor for pairs
```

To lift it, use a combination of $\langle \$ \rangle$ and $\langle * \rangle$

```
tuple x y = (,) <$> x <*> y
```



Functor-applicative-monad hierarchy

Functor which is not applicatives?

Applicative which is not monad?



Functor-applicative-monad hierarchy

Functor which is not applicatives?

Applicative which is not monad?

Thanks, StackOverflow!



Functor which is not applicative

```
data Pair a b = Pair a b  -- like a tuple
```

```
instance Functor (Pair a) where  
    fmap f (Pair x y) = Pair x (f y)
```

```
instance Applicative (Pair a) where  
    pure y = ({- what here?? -}, y)
```

```
-- You can fix it, but not for every "a"  
instance Monoid a => Applicative (Pair a) where  
    pure y = (mempty, y)  
    (<*>) = ...
```



Applicative which is not monad

```
data ZipList a = ZL [a]
```

```
instance Applicative ZipList where
```

```
  pure x = ZL (repeat x)  -- infinite list of "x"s  
  ZL fs <*> ZL xs = ZL (zipWith (\f x -> f x) fs xs)  
  -- This obeys all the laws!
```

```
-- You can use the "Monad" from lists  
-- But then the following does not hold:
```

```
f <*> x  ==  do f' <- f  
              x' <- x  
              return (f' x')
```



Closing remarks



Goals for the course

- ▶ Learn the **functional** paradigm and **style**
 - ▶ You can apply FP techniques everywhere!
 - ▶ Every (serious) language has H-O functions
- ▶ Experience a **strong static type system**
- ▶ **Reason** about programs
 - ▶ Correct software is our ultimate goal



Courses about or using FP at UU

- ▶ *Functioneel Programmeren*
- ▶ Talen en Compilers: year 3, period 2
 - ▶ Haskell applied to compiler writing
- ▶ Software Testing en Verificatie: year 3, period 4
 - ▶ More reasoning about programs



If you want to know more

More Haskell?

- ▶ *Pearls of Functional Algorithm Design*, by Bird
 - ▶ Puzzles with a nice functional solution
- ▶ *the fun of programming*, by Gibbons and de Moor
 - ▶ Even more niceties in a functional style
- ▶ *Haskell from First Principles*, by Allen and Moronuki
 - ▶ Covers additional topics, like transformers
- ▶ *Beginning Haskell*, by, ehmmm... me
 - ▶ Which happens to be an intermediate book



If you want to know more

Learn other functional languages

- ▶ *F#* for the .NET platform
 - ▶ *Beginning F# 4.0* and *Expert F# 4.0*
- ▶ *Kotlin* and *Scala* for the Java platform
 - ▶ *Functional Kotlin*
 - ▶ *Functional Programming in Scala*
- ▶ *Swift* for iOS development
 - ▶ *Functional Swift*



If you want to know more

Or just drop by my office



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Success with your exams!

