Lecture 14. Monadic utilities and traversables

Functional Programming 2018/19

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Goals

- ▶ Look at some utilities for monadic code
 - How to write functions working on monads
- ▶ In particular, learn about *traversable* functors

Chapter 14.3 from Hutton's book

The functor - applicative - monad hierarchy

```
class Functor f where
  fmap :: (a -> b) -> f a -> f b

class Functor f => Applicative f where
  pure :: a -> f a
  (<*>) :: f (a -> b) -> f a -> f b

class Applicative f => Monad f where
  -- return is the same as Applicative's pure
  (>>=) :: f a -> (a -> f b) -> f b
```

Monadic utilities

The "final M" family

Many standard functions have monadic counterparts

```
map :: (a -> b) -> [a] -> [b]
mapM :: (a -> m b) -> [a] -> m [b]
filter :: (a \rightarrow Bool) \rightarrow [a] \rightarrow [b]
filterM :: (a \rightarrow m Bool) \rightarrow [a] \rightarrow m [b]
fold: (b \rightarrow a \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow [b]
foldM :: (b -> a -> m b) -> b -> [a] -> m [b]
zipWith :: (a \rightarrow b \rightarrow c) \rightarrow [a] \rightarrow [b] \rightarrow [c]
zipWithM :: (a \rightarrow b \rightarrow m c) \rightarrow [a] \rightarrow [b] \rightarrow m [c]
```

1. Define the type

```
filterM :: (a -> m Bool) -> [a] -> m [b]
```

2. Enumerate the cases

```
filterM p [] = _
filterM p (x:xs) = _
```

- 3. Define the simple (base) cases
 - Remember that we work in a monadic context

```
filterM _ [] = return []
```

- 4. Define the other (recursive) cases
 - ▶ We cannot use a guard, because p does not return Bool

```
filterM p (x:xs) | p x = _ -- Does not work
```

► For the same reason, we cannot use an if directly

```
filterM p (x:xs) = if p x then _ else _ -- Nope
```

- p returns its value wrapped in a monad
- ▶ We unwrap it by means of <- in a do-block

```
filterM p (x:xs) = do q \leftarrow p x
```

. .

- 4. Define the other (recursive) cases
 - Let us try to write the rest as with filter

```
filterM p (x:xs) = do
  q <- p x
  if q
    then x : filterM p xs
    -- :: a :: m [a]
  else filterM p xs</pre>
```

erse filtern p xs

Wrong: x is pure but the result of filterM is monadic

- 4. Define the other (recursive) cases
 - ► Solution 1: unwrap the result of filterM with <-

```
filterM p (x:xs) = do
  q \leftarrow p x
  r <- filterM p xs
  if q then return (x:r) else return r
 Solution 2: lift the pure part with applicatives
```

```
filterM p (x:xs) = do
  q \leftarrow p x
  if q then (x:) <$> filterM p xs
       else
                       filterM p xs
```



Having to unwrap $p \times distracts$ a bit

- ▶ We just need a "lifted" if-then-else
- ▶ Why not define our own ifM which does just that?

Cooking zipWithM

Using a do-block:

```
zipWithM _ [] _ = return []
zipWithM _ _ [] = return []
zipWithM f (x:xs) (y:ys) = do
 z \leftarrow f x y
  r <- zipWithM f xs ys
  return (z:r)
```

Cooking zipWithM

```
Using a do-block:
zipWithM _ [] _ = return []
zipWithM [] = return []
zipWithM f (x:xs) (y:ys) = do
 z \leftarrow f x y
  r <- zipWithM f xs ys
  return (z:r)
Using applicative style:
zipWithM _ [] _ = return []
zipWithM [] = return []
zipWithM f (x:xs) (y:ys)
  = (:) <$> f x y <*> zipWithM f xs ys
```

Traversables

Functors generalize maps

We started with a map function for lists

which we generalized to arbitrary functors

fmap :: Functor f => (a -> b) -> f a -> f b

Can we do something similar for mapM?



Cooking mapTreeM

- 1. Define the type
 - Remember: we want a monadic function as argument

```
mapTreeM :: Monad m
=> (a -> m b) -> Tree a -> m (Tree b)
```

2. Enumerate the cases

```
mapTreeM f Leaf = _
mapTreeM f (Node l x r) = _
```

- 3. Define the simple (base) cases
 - Remember: we are now in a monadic context

```
mapTreeM _ Leaf = return Leaf
```

Cooking mapTreeM

- 4. Define the other (recursive) cases
 - Solution 1: using do-notation
 mapTreeM f (Node 1 x r) = do
 1' <- mapTreeM f 1
 x' <- f x
 r' <- mapTreeM f r
 return (Node 1' x' r')</pre>
 - Solution 2: using applicative style

Cooking mapTreeM

5. Generalize and simplify

- ► The second implementation only needs Applicative, the first uses do and thus needs Monad
- Remember: Applicative is more general than Monad

Traversables

The generalization of Functor to handle functions of the form a -> f b is called a **traversable** (functor)

- ▶ f defines the context in which the function run
- ▶ t defines the data structure which contains the elements to map over

Cooking printTree

printTree t print the elements of the tree to the screen, in infix order, one per line

1. Define the type

```
printTree :: Show e => Tree e -> IO ()
```

- 2. IO is an applicative, Tree is a traversable
 - We can just "map" the print function!

```
printTree = traverse print
```

Cooking printTree

- 3. We run into a problem, the result of traverse print is IO (Tree ()), not IO ()

 - ► Solution 2: use void :: Functor f => f a -> f () to discard the value
 - printTree = void . traverse print
- 4. This implementation works for any traversable



Summary

- ► Haskell has powerful ways to abstract
 - Code and design patterns become functions
- ► Higher-order functions
 - Maps, folds, filters...
- ► Higher-kinded abstractions
 - Functors, monads and applicatives for "contexts"
 - ► Functors and traversables for "containers"

sequence for arbitrary traversables is not part of 2018/2019 contents

Note that sequence for IO is part of the contents

Another look at zipWithM

What happens if we just zip the function?

Alas, what we require is m [r]

sequence to the rescue

During the lecture on IO, we introduced a function

```
sequence :: [IO a] -> IO [a]
```

This function actually works on any monad

```
sequence :: Monad m \Rightarrow [m a] \rightarrow m [a]
```

We can write zipWithM with its help

```
zipWithM f xs ys = sequence (zipWith f xs ys)
```

Generalizing sequence

A traversable admits a generic version of sequence

Cooking generic sequence

Let us try to find the implementation by looking at the types

```
traverse :: (a \rightarrow fb) \rightarrow ta \rightarrow f(tb)
sequence :: t(fr) \rightarrow f(tr)
```

Cooking generic sequence

Let us try to find the implementation by looking at the types

```
traverse :: ( a \rightarrow f b) \rightarrow t a \rightarrow f (t b) sequence :: t (f r) \rightarrow f (t r)
```

Solution: let us make a = f rand b = r

traverse ::
$$(f r \rightarrow f r) \rightarrow t (f r) \rightarrow f (t r)$$

How do we get a function of type $f r \rightarrow f r$?

Cooking generic sequence

Let us try to find the implementation by looking at the types

```
traverse :: ( a -> f b) -> t a -> f (t b) sequence :: t (f r) -> f (t r)

Solution: let us make a = f r and b = r

traverse :: (f r -> f r) -> t (f r) -> f (t r)

How do we get a function of type f r -> f r?
```