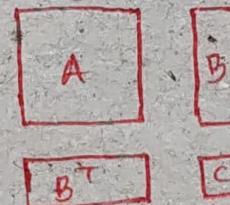


Week 8

Schur's Complement

Consider the Block matrix;

$$X = \begin{bmatrix} A & B \\ B^T & C \end{bmatrix}$$



$$A \in S^{n \times n}$$

$$C \in S^{m \times m}$$

$$X \in S^{n+m}$$

- $A = A^T$, $C = C^T$, A, C are Symmetric

- $X = X^T \in S^n$, X is also Symmetric.

Schur's Complement,

$$S_1 = C - B^T A^{-1} B, \text{ when } A \text{ is invertible}$$

$$S_2 = A - B C^{-1} B^T, \text{ when } C \text{ is invertible.}$$

$B \rightarrow$ Rectangular Matrix

Key result :

$$(a) \underline{x} > 0 \Leftrightarrow S_1 > 0 \quad A > 0 \Leftrightarrow S_2 > 0 \quad C > 0$$

(b) $\det(x) = \det(A) \det(S_1) = \det(C) \det(S_2)$. Positive Definite (PD)

Example

$$D = \begin{bmatrix} tI & \underline{x} \\ \underline{x}^T & t \end{bmatrix}$$

$$\begin{array}{c|c} tI & \underline{x} \\ \hline \underline{x}^T & t \end{array}$$

$\underline{x} \rightarrow$ column vector
 $\underline{x}^T \rightarrow$ Row vector
 $t \rightarrow$ Single element.

Relate with Block matrix,

$$A = tI, B = \underline{x}, C = t, t \geq 0.$$

$$S_1 = t - \underline{x}^T t^{-1} \underline{x} = t - \frac{\underline{x}^T \underline{x}}{t}$$

$$\text{so, } D \geq 0 \Leftrightarrow \underline{x}^T \underline{x} \leq t^2. \quad (\text{or}) \quad \|\underline{x}\| \leq t$$

Second order cone (SOC)
constraint

Semi definite Programs (SDP):

- Linear Matrix Inequalities (LMI) form.

Let us look into the SDPs, in particular we'll look at a form of SDPs, which is called LMI.

The problem is of the form,

$$\min \underline{C}^T \underline{x} \quad \leftarrow \text{Linear objective}$$

$$\circ F(\underline{x}) : G + F_1 \underline{x}_1 + F_2 \underline{x}_2 + \dots + F_n \underline{x}_n \leq 0$$

$$\circ A \underline{x} = b$$

Negative semidefinite

This is Linear Matrix Inequality (LMI).

This is more general than SOCP.

$$\underline{x} = \begin{bmatrix} \underline{x}_1 \\ \vdots \\ \underline{x}_m \end{bmatrix} \quad G, F_i \in \mathbb{S}^m$$

$$\text{Example: } [G]_{jj} = g_j \quad [F_i]_{jj} = f_{ij}$$

G, F_i are diagonal.

$$F(x) = \begin{bmatrix} g_1 + \sum_{i=1}^n f_{i1} x_i & 0 & \cdots & 0 \\ 0 & g_2 + \sum_{i=1}^n f_{i2} x_i & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & g_m + \sum_{i=1}^n f_{im} x_i \end{bmatrix}$$

$$F(x) \leq 0 \iff g_j + \sum_{i=1}^m f_{ij} x_i \leq 0, \forall j=1, \dots, m.$$

This is a very special case in which SDP reduces to a Linear Program (LP).

Is SDP convex?

$$F(x) \leq 0 \iff \lambda_{\max}(F(x)) \leq 0$$

$$(u) F(x) \text{ is PSD.} \quad (v) \lambda_{\max}(F(x)) \leq 0$$

Maximum Eigen value.

Claim that we are going to make:

$\lambda_{\max}(F(x))$ convex function in x , for $G, F_i \in S^m$

Note: $\lambda_{\max}(A) = \max_y y^T A y$
 $\|y\| \leq 1$

$$\max_{\|y\| \leq 1} y^T F(x) y = \max_{\|y\| \leq 1} y^T G y + \sum_{i=1}^m (y^T F_i y) x_i$$

Affine function of x .

$$\lambda_{\max}(F(\underline{x})) = \max_{\|\underline{y}\|=1} \underline{y}^T F(\underline{x}) \underline{y}$$

(pointwise maximum of affine)

$\Rightarrow \lambda_{\max}(F(\underline{x}))$ is convex in \underline{x} .

\Rightarrow SDP are convex.

Note: $\lambda_{\min}(F(\underline{x}))$ is concave in \underline{x} $\Rightarrow F(\underline{x}) \geq 0$ also valid.

$$F(\underline{x}) \geq 0 \Leftrightarrow -F(\underline{x}) \leq 0 \quad (\text{or}) \quad \lambda_{\min}(F(\underline{x})) \geq 0$$

$$F(\underline{x}) = G + \sum F_i \underline{x}_i \leq 0$$

$$-F(\underline{x}) = -G - \sum F_i \underline{x}_i = H(\underline{x}) \geq 0$$

Both are valid convex form.

SDP Examples

① Express the two LMI's into single LMI.

$$F_1(\underline{x}) \leq 0 \quad \Rightarrow \quad \begin{bmatrix} F_1(\underline{x}) & 0 \\ 0 & F_2(\underline{x}) \end{bmatrix} \leq 0$$

$$\Rightarrow \begin{bmatrix} G_1 & 0 \\ 0 & G_2 \end{bmatrix} + \sum x_i \begin{bmatrix} [F_1]_i & 0 \\ 0 & [F_2]_i \end{bmatrix} \leq 0.$$

② $\min \lambda_{\max}(F(\underline{x}))$

Applying epigraph trick,

$$\Rightarrow \min t$$

$$\text{s.t. } \circ \lambda_{\max}(F(\underline{x})) \leq t$$

(or)

$$\circ \lambda_{\max}(F(\underline{x}) - t \mathbb{I}) \leq 0$$

(or)

$$\min t$$

$$F(\underline{x}) \leq t \mathbb{I}$$

$$\Rightarrow \underbrace{G + \sum F_i \underline{x}_i - t \mathbb{I}}_{\text{LMI constraint}} \leq 0, \underline{x}_1, \dots, \underline{x}_n, t$$

$$\textcircled{3} \quad \text{LP} \subseteq \text{QP} \subseteq \text{QCQP} \subseteq \text{SOCP} \subseteq \text{SDP}$$

Most Specialized
(Easiest to solve)

Most General
(Hardest to solve)

$$\begin{aligned} & \min c^T x \\ \text{s.t. } & \|A_i x + b_i\| \leq c_i^T x + d_i \quad (\text{Second order cone (soc) constraint}) \\ & Ax = b. \end{aligned}$$

$$\text{Recall } \|x\| \leq t \Rightarrow \begin{bmatrix} tI & x \\ x^T & t \end{bmatrix} \leq 0.$$

$$\Rightarrow \begin{bmatrix} (c_i^T x + d_i)I & A_i x + b_i \\ (A_i x + b_i)^T & c_i^T x + d_i \end{bmatrix} \leq 0$$

$\underbrace{}_{\text{LMI}}$

$$\Rightarrow \begin{bmatrix} d_i I & -b_i \\ b_i^T & -d_i \end{bmatrix} + \sum_{j=1}^n x_j \begin{bmatrix} C_{ij} I & a_{ij} \\ a_{ij}^T & C_{ij} \end{bmatrix} \leq 0$$

$$\textcircled{4} \quad \min \frac{(c^T x)^2}{d^T x}$$

$$Ax = b.$$

Applying epigraph trick,

$$\min t.$$

$$(c^T x)^2 \leq (d^T x)t$$

Applying Schur's complement,

$$\begin{bmatrix} t & c^T x \\ c^T x & d^T x \end{bmatrix} \geq 0.$$

This is a LMI (SDP).

SDP Dual form

Let us see the Dual problem corresponding to the LMI, which yield different ways of expressing the problems and also different ways of taking the dual.

Consider the problem,

$$\begin{aligned} & \min c^T x \\ \text{s.t. } & \lambda_{\max}(F(x)) \leq 0 \quad (\text{or}) \quad \min c^T x \\ & F(x) \leq 0 \quad (\text{not in standard convex form}) \end{aligned}$$

Matrix Dual Variable : $Y \in S^n$

$$\begin{aligned} L(x, Y) &= c^T x + \langle F(x), Y \rangle \\ &= c^T x + \sum_{i,j} [F(x)]_{ij} Y_{ij} \\ &= c^T x + \text{Tr}(Gy) + \sum_{i=1}^n x_i \text{Tr}(F_i Y) \end{aligned}$$

Observe that $L(x, Y)$ is affine in x and Y .

Therefore,

$$\min_x L(x, Y) = \sum_{i=1}^n \min_{x_i} x_i (c_i + \text{Tr}(F_i Y)) + \text{Tr}(Gy)$$

$$\min_{x_i} x_i (c_i + \text{Tr}(F_i Y)) = \begin{cases} 0, & \text{only if } c_i + \text{Tr}(F_i Y) \geq 0 \\ -\infty, & \text{otherwise} \end{cases}$$

The dual of SDP is given by

$$\begin{aligned} & \max -\text{Tr}(Gy) \\ \text{s.t. } & \circ \quad c_i + \text{Tr}(F_i Y) = 0 \\ & \circ \quad Y \succeq 0. \end{aligned}$$

PSD

Dual of SDP is also an SDP.

$$y = \underbrace{\text{vec}(Y)}_{\text{Vectorized version of } Y} \in \mathbb{R}^{n^2}$$

$$\text{Tr}(Gy) = g^T y, \quad g = \text{vec}(G)$$

$$\text{Tr}(F_i y) = f_i^T y, \quad f_i = \text{vec}(F_i)$$

then, $y = \sum_{i,j} \underline{\lambda}_{ij} Y_{ij} \geq 0$

LMI form

$$\underline{\lambda}_{ij} = \begin{bmatrix} 0 & \dots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \dots & 0 \end{bmatrix}$$

$(i,j)^{\text{th}}$ entry

$$\max g^T y$$

$$f_i^T y + c_i = 0$$

$$H(y) = \sum \underline{\lambda}_{ij} Y_{ij} \geq 0. \quad \text{also LMI.}$$
