

Q2

merge_sort function:

The function divides the array into halves recursively until each subarray contains only one element. This operation takes $O(\log n)$ time because the array is divided into halves in each recursive call until it reaches a size of 1.

Then, the merge operation is called on each pair of subarrays. Since there are $\log n$ levels of recursion,

and at each level, each element in the array is merged exactly once, the total time complexity for merging all levels is $O(n)$.

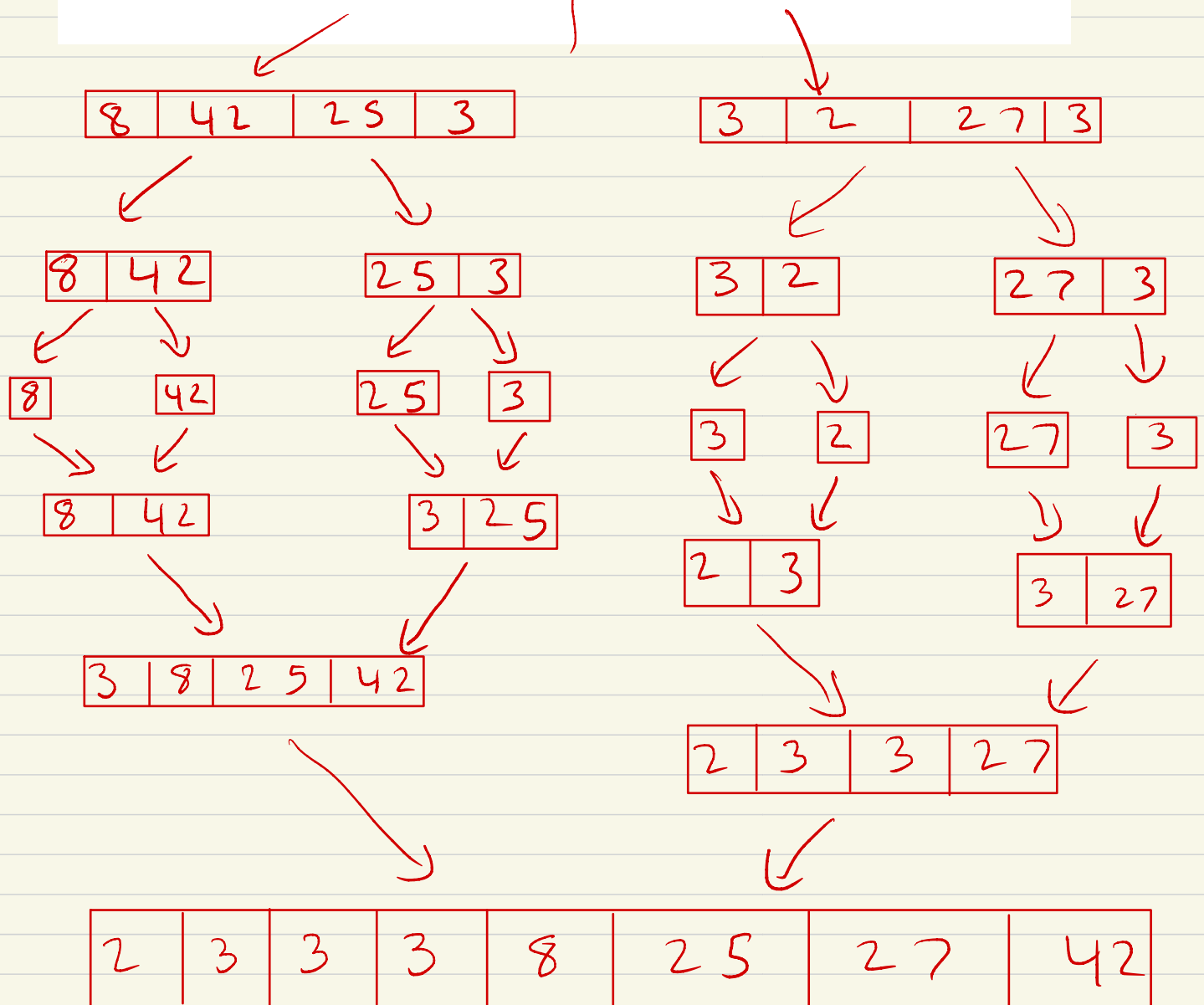
merge function:

The merge function iterates through each element of the two subarrays exactly once, comparing and merging

them into a single sorted array. This process takes $O(n)$ time, where n is the total number of elements in the two subarrays.

3)

1	2	3	4	5	6	7	8
8	42	25	3	3	2	27	3



The initial array is divided into smaller sub-arrays:

[8 42 25 3] [3 2 27 3]

These are further divided into [8 42] then [8] [42], [25 3] then [25] [3] and

$[3 \ 2]$ then $[3]$ $[2]$ and
 $[27 \ 3]$ then $[27]$ $[3]$.

These sub arrays are
then sorted individually
and merged. After merge
they become $[3 \ 8 \ 25 \ 42]$
and $[2 \ 3 \ 3 \ 27]$. The two
sorted sub arrays are
then merged this gives
us final sorted array
 $[2 \ 3 \ 3 \ 3 \ 8 \ 25 \ 27 \ 42]$.

Q4)

The divide phase would require at most $\log_2(8) = 3$ steps to create sub-arrays of size 1 which is consistent with the diagram in Q3. In the worst case scenario, each level of the recursion during the merge phase would require processing all 8 elements at once. Therefore, the total number of steps would be 3 (divide phase) * 8 (merge phase) = 24, which is consistent with the expected time complexity of $O(n \log n)$. So, yes, the number of steps aligns with the expected time complexity analysis of Merge Sort. Each element is visited at most $\log(n)$ times during the divide phase and n times during the merge phase, leading to an overall time complexity of $O(n \log n)$.