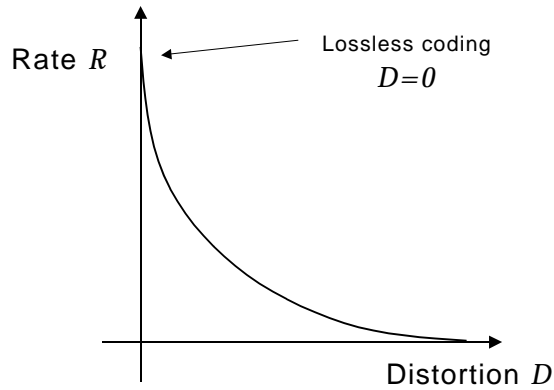


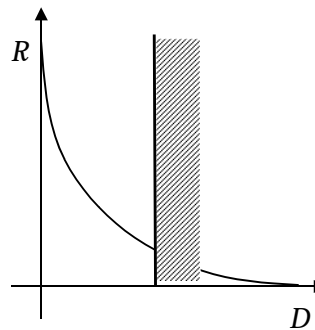
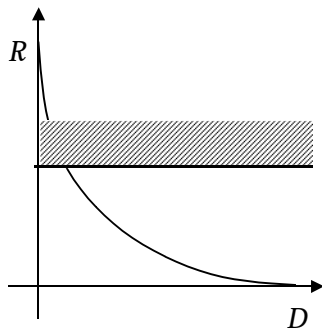
# Lossy compression

- Lower the bit-rate  $R$  by allowing some acceptable distortion  $D$  of the signal.



## Types of lossy compression problems

- Given maximum rate  $R$ , minimize distortion  $D$
- Given distortion  $D$ , minimize rate  $R$



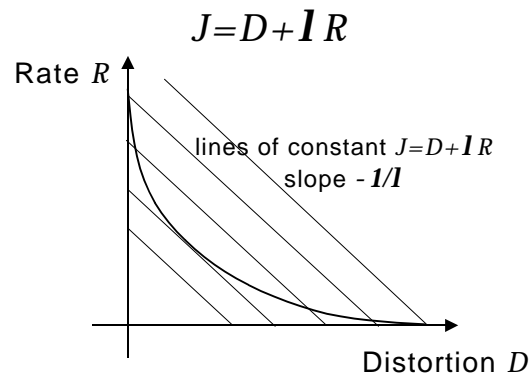
Equivalent constrained optimization problems, often unwieldy due to constraint.



## Lagrangian formulation of the R-D problem

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- Instead of cost function  $D$ , with constrained  $R$ , or  $R$ , with constrained  $D$ , use unconstrained Lagrangian cost function



- Requires convex R-D function



## Topics in lossy compression

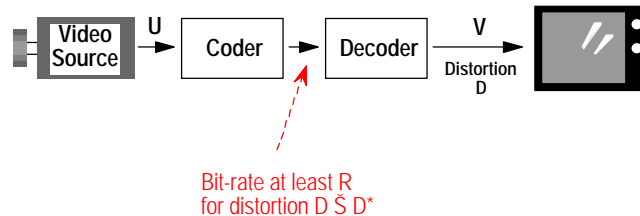
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- Information theoretical bounds for lossy compression: rate distortion theory
  - R-D function for memoryless Gaussian sources with mean-squared error distortion criterion
  - R-D function for Gaussian sources with memory
  - R-D for images
- Practical lossy compression techniques: quantization
  - Scalar quantization and vector quantization
  - Quantizer design for fixed length codes
  - Quantizer design for variable length codes



## Rate distortion theory

- Rate distortion theory calculates the minimum transmission bit-rate  $R$  for a required picture quality.



- Results of rate distortion theory are obtained without consideration of a specific coding method.



## Mutual information

- "Mutual information" is the information that symbols  $u$  and symbols  $v$  convey about each other.
- Average mutual information:

$$\begin{aligned} I(U;V) &= H(V) - H(V|U) = H(U) - H(U|V) \\ &= \sum_u \sum_v P(u,v) \log \frac{P(u,v)}{P(u)P(v)} \end{aligned}$$

- Properties of mutual information:

$$\begin{aligned} 0 &\leq I(U;V) = I(V;U) \\ I(U;V) &\leq H(U) \\ I(V;U) &\leq H(V) \end{aligned}$$

- Channel coding:** channel capacity  $C$  is maximum of mutual information between transmitter and receiver.



# Distortion

- Symbol  $u$  sent,  $v$  received
- Per-symbol distortion:

$$\begin{aligned} d(u, v) &\geq 0 \\ d(u, v) &= 0 \text{ for } u = v \end{aligned}$$

- Average distortion:

$$D(U, V) = E\{d(U, V)\} = \sum_u \sum_v P(u, v) d(u, v)$$

- Distortion criterion:

$$D \leq D^*$$

maximum average distortion



# Rate distortion function

- Definition:

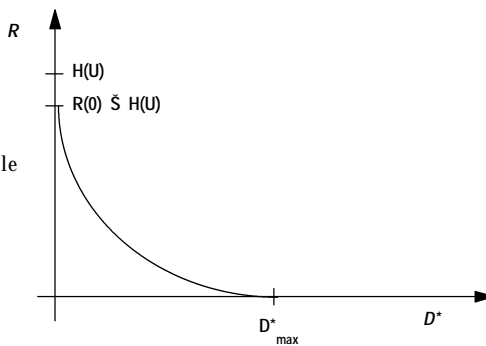
$$R(D^*) = \min_{D \leq D^*} \{I(U; V)\}$$

- For a given maximum average distortion  $D^*$ , the rate distortion function  $R(D^*)$  is the lower bound for the transmission bit-rate.

- Typical  $R(D^*)$  function:

$H(U)$ : entropy of the source

$R(0)$ : rate with just not perceptible loss of quality



## Shannon Lower Bound

- It can be shown that  $H(U - V | V) = H(U | V)$

- Thus

$$\begin{aligned} R(D^*) &= \min_{D \leq D^*} \{H(U) - H(U | V)\} \\ &= H(U) - \max_{D \leq D^*} \{H(U | V)\} \\ &= H(U) - \max_{D \leq D^*} \{H(U - V | V)\} \end{aligned}$$

- Ideally, the source coder would produce distortions  $u-v$  that are statistically independent from the reconstructed signal  $v$  (not always possible!).
- Shannon lower bound:

$$R(D^*) \geq H(U) - \max_{D \leq D^*} H(U - V)$$



## $R(D^*)$ function for a memoryless Gaussian source and MSE distortion

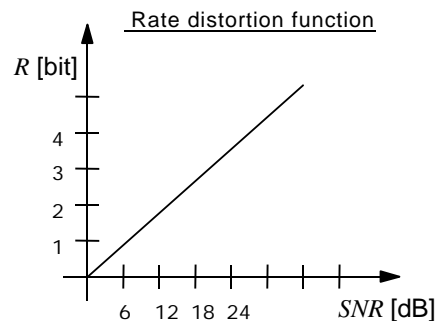
- Gaussian source, variance  $\mathbf{S}^2$

- Mean squared error

$$D = E\{(U - V)^2\}$$

- Rule of thumb: 6 dB  $\cong$  1 bit
- The  $R(D^*)$  for non-Gaussian sources with the same variance  $\mathbf{S}^2$  is always below this Gaussian  $R(D^*)$ -curve.

$$R(D^*) = \frac{1}{2} \log \left( \frac{\mathbf{S}^2}{D^*} \right)$$



$$SNR = 10 \log_{10} \left( \frac{\mathbf{S}^2}{D^*} \right) [\text{dB}]$$



## $R(D^*)$ function for Gaussian source with memory I

- Jointly Gaussian source with power spectrum  $\Phi_{uu}(\omega)$
- Mean squared error distortion  $D = E\{(U - V)^2\}$
- Parametric formulation of the  $R(D^*)$  function

$$D(\mathbf{q}) = \frac{1}{2p} \int_{\omega} \min\{\mathbf{q}, \Phi_{uu}(\omega)\} d\omega$$

$$R(\mathbf{q}) = \frac{1}{2p} \int_{\omega} \max\left\{0, \frac{1}{2} \log \frac{\Phi_{uu}(\omega)}{\mathbf{q}}\right\} d\omega$$

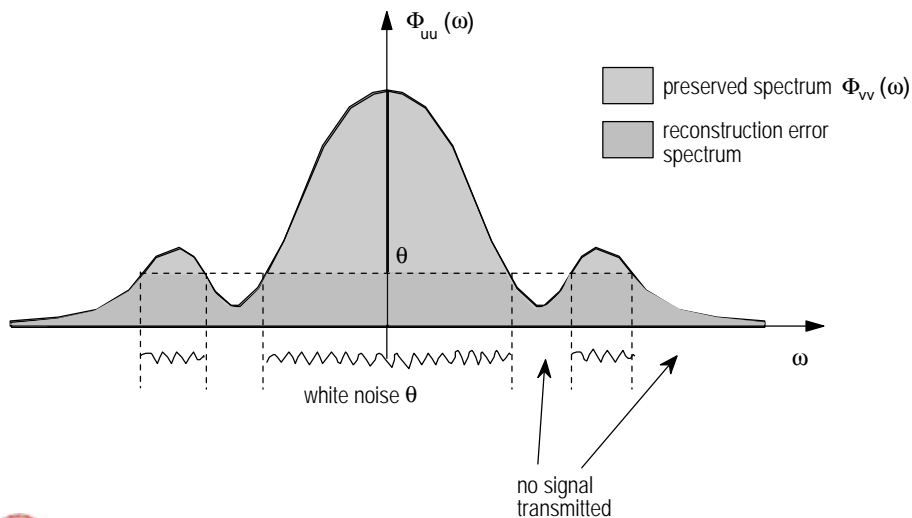
- $R(D^*)$  for non-Gaussian sources with the same power spectral density is always lower.



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Rate Distortion Theory no. 11

## $R(D^*)$ function for Gaussian source with memory II



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Rate Distortion Theory no. 12

## Rate distortion function for images I

- Signal model: Gaussian source with acf

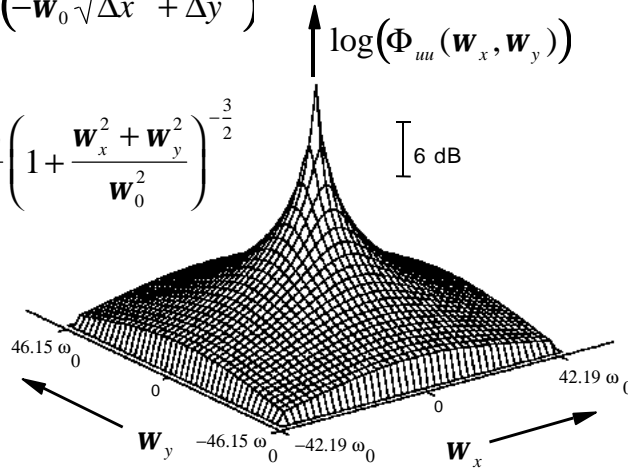
$$R_{uu}(\Delta x, \Delta y) = \exp(-w_0 \sqrt{\Delta x^2 + \Delta y^2})$$

- Power spectral d

$$\Phi_{uu}(\mathbf{w}_x, \mathbf{w}_y) = \frac{2p}{w_0^2} \left( 1 + \frac{w_x^2 + w_y^2}{w_0^2} \right)^{-\frac{3}{2}}$$

$$w_0 = -\ln(0.93)$$

correlation between adjacent pixels

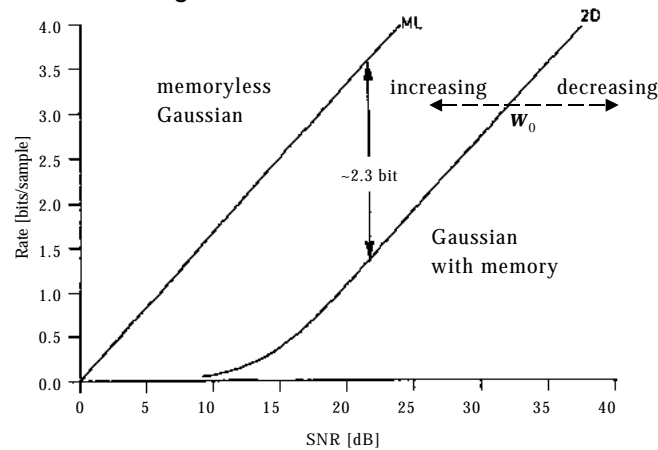


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Rate Distortion Theory no. 13

## Rate distortion function for images II

- Mean squared error criterion:  $D = E\{(U - V)^2\}$
- After numerical integration:



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Rate Distortion Theory no. 14

## Summary: rate distortion theory

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- Lagrangian cost function  $J=D+\lambda R$  can be used instead of solving constrained rate or distortion minimization.
- Rate-distortion theory: minimum transmission bit-rate for given distortion
- Shannon Lower Bound assumes statistical independence between distortion and reconstructed signal
- $R(D^*)$  for memoryless Gaussian source and MSE: 6 dB/bit
- $R(D^*)$  for Gaussian source with memory and MSE: encode spectral components independently, introduce white noise, suppress small spectral components
- Theoretical gain ~2.3 bits/sample by exploiting spatial redundancy in the video signal

