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- Tasks: $F[0][0] = C[0][0] \equiv 1, \text{read}(F[0][k-1]) \equiv 1, \text{read}(C[0][k]) \equiv 1, \text{add}(F[0][k-1], C[0][k]) \equiv 3, \text{write}(F[0][k]) \equiv 1, \text{read}(F[i-1][0]) \equiv 1, \text{read}(C[i][0]) \equiv 1, \text{add}(F[i-1][0], C[i][0]) \equiv 3, \text{write}(F[i][0]) \equiv 1, \text{read}(F[i-1][j]) \equiv 1, \text{read}(F[i][j-1]) \equiv 1, \max(F[i-1][j], F[i][j-1]) \equiv 3, \text{read}(C[i][j]) \equiv 1, \text{add}(\max(F[i-1][j], F[i][j-1]), C[i][j]) \equiv 3, \text{write}(F[i][j]) \equiv 1, \text{read}(F[n-1][m-1]) \equiv 1, \Theta(N) =$

$$6m + 6n + 10mn + 2$$

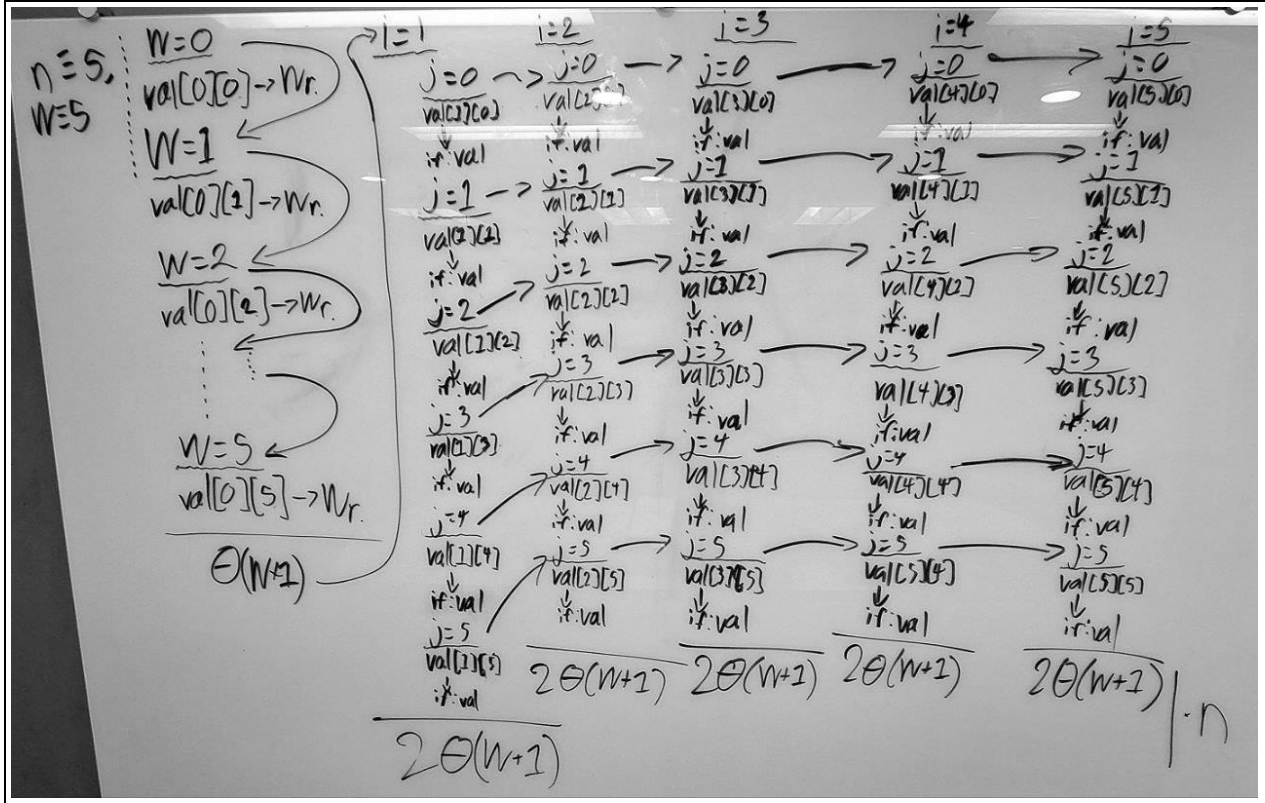
- **Dependencies:**

$$\begin{aligned} F[0][0] = C[0][0] &\rightarrow \text{read}(F[0][k-1]) + \text{read}(C[0][k]) + \text{write}(F[0][k]) \rightarrow \\ \text{read}(F[i-1][j]) + \text{read}(F[i][j-1]) + \text{read}(C[i][j]) &\rightarrow \max(F[i-1][j], F[i][j-1]) + \\ \text{read}(C[i][j]) &\rightarrow +\text{add}(\max(F[i-1][j], F[i][j-1]), C[i][j]) \rightarrow \text{write}(F[i][j]) \rightarrow \\ &\text{read}(F[n-1][m-1]) \end{aligned}$$

- $\text{width} = 3mn$
- $\sum p_i = 6m + 6n + 10mn + 2$
- $T_\infty = 6m + 6n + 10mn + 2, \text{length} = 7$

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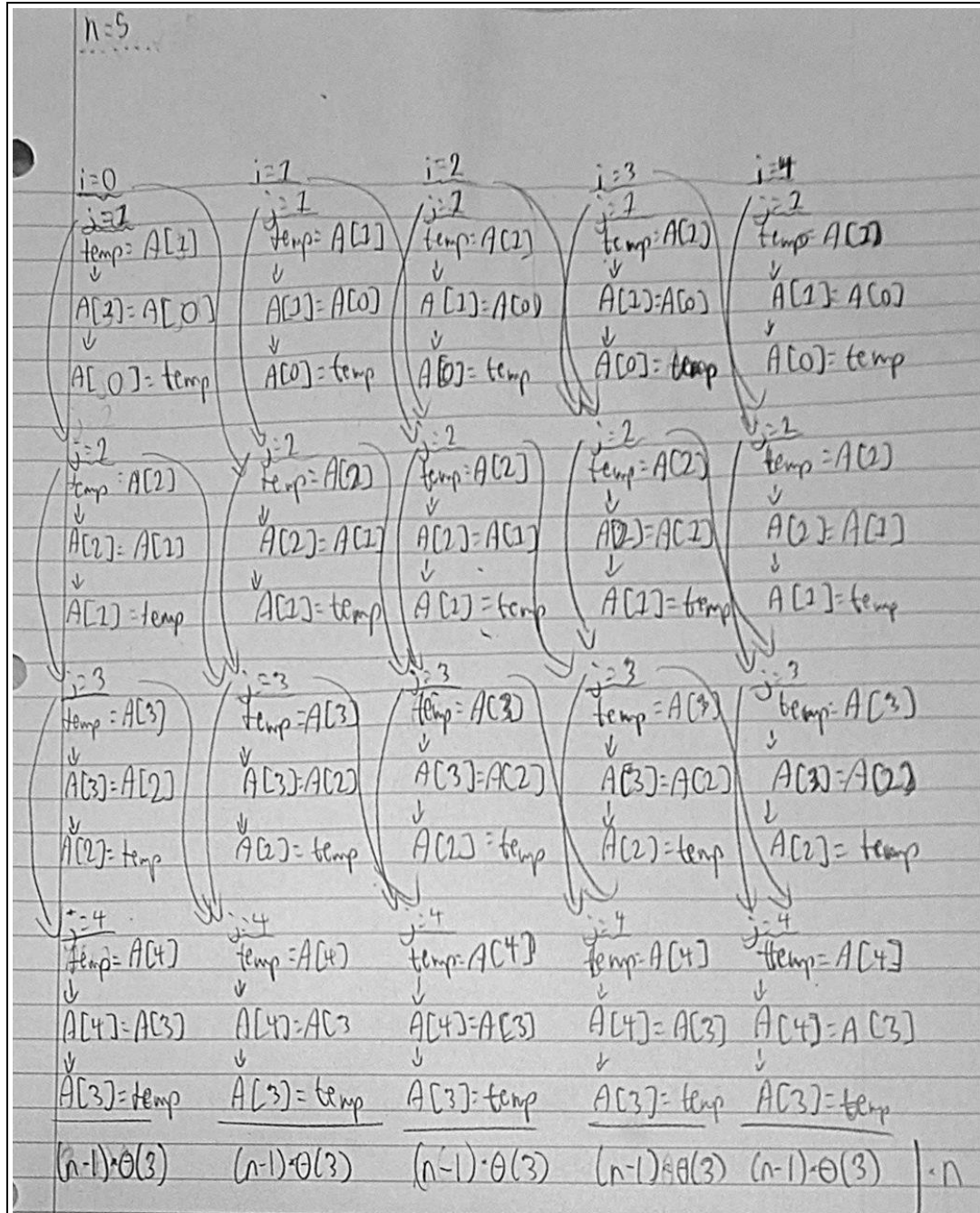
- Complexity: $\Theta(N) = (10n + 1)\Theta(W + 1)$
- Dependencies:



- width = W
- $\sum p_i = (10n + 1)\Theta(W + 1)$
- $T_\infty = (10n + 1)\Theta(W + 1)$, length = $W + n$

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- **Complexity:** $\Theta(N) = \boxed{n(n-1)\Theta(3)}$
- **Dependencies:**



- **width** = n
- $\sum p_i = \boxed{n(n-1)\Theta(3)}$
- $T_\infty = \boxed{(n-1)\Theta(3)}$, **length** = $\boxed{n-1}$