

Performance evaluation and control improvements for solving optimization problems on Noisy Intermediate-Scale Quantum (NISQ) platforms

PhD Thesis defense

Valentin GILBERT

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Reviewers: Frank PHILLIPSON, Caroline PRODHON

Examiners: Daniel ESTEVE, Jeannette LORENZ



Introduction and Context

Collection criteria

- Score of interest (beauty, color, shape, brand ...)



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$$S_{\text{interest}} = 0.03$$



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$$S_{\text{interest}} = 0.8$$

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- Similarity score: color and shape similarity

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$$S_{\text{similarity}} = 0.95$$

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$$S_{\text{similarity}} = 0.2$$

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$$S_{\text{similarity}} = 0.2$$

Collection representation

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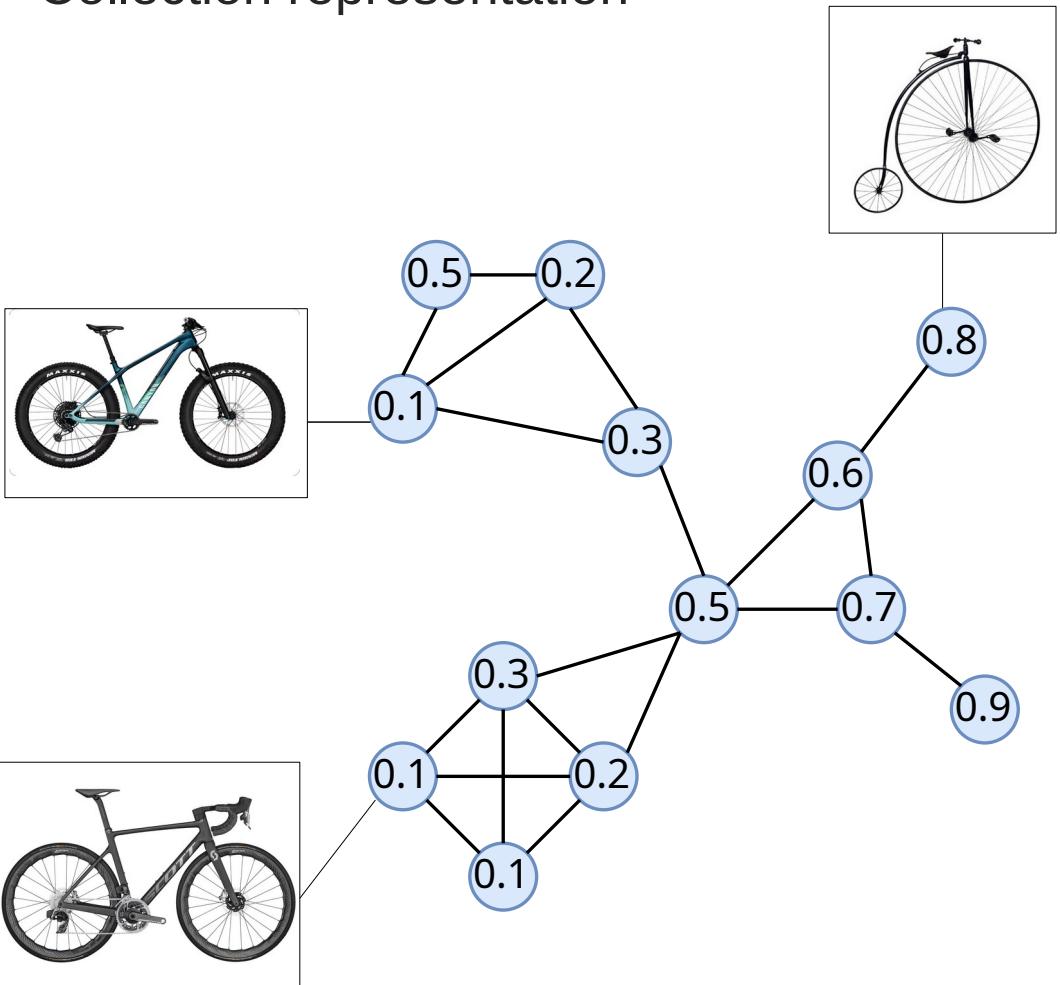
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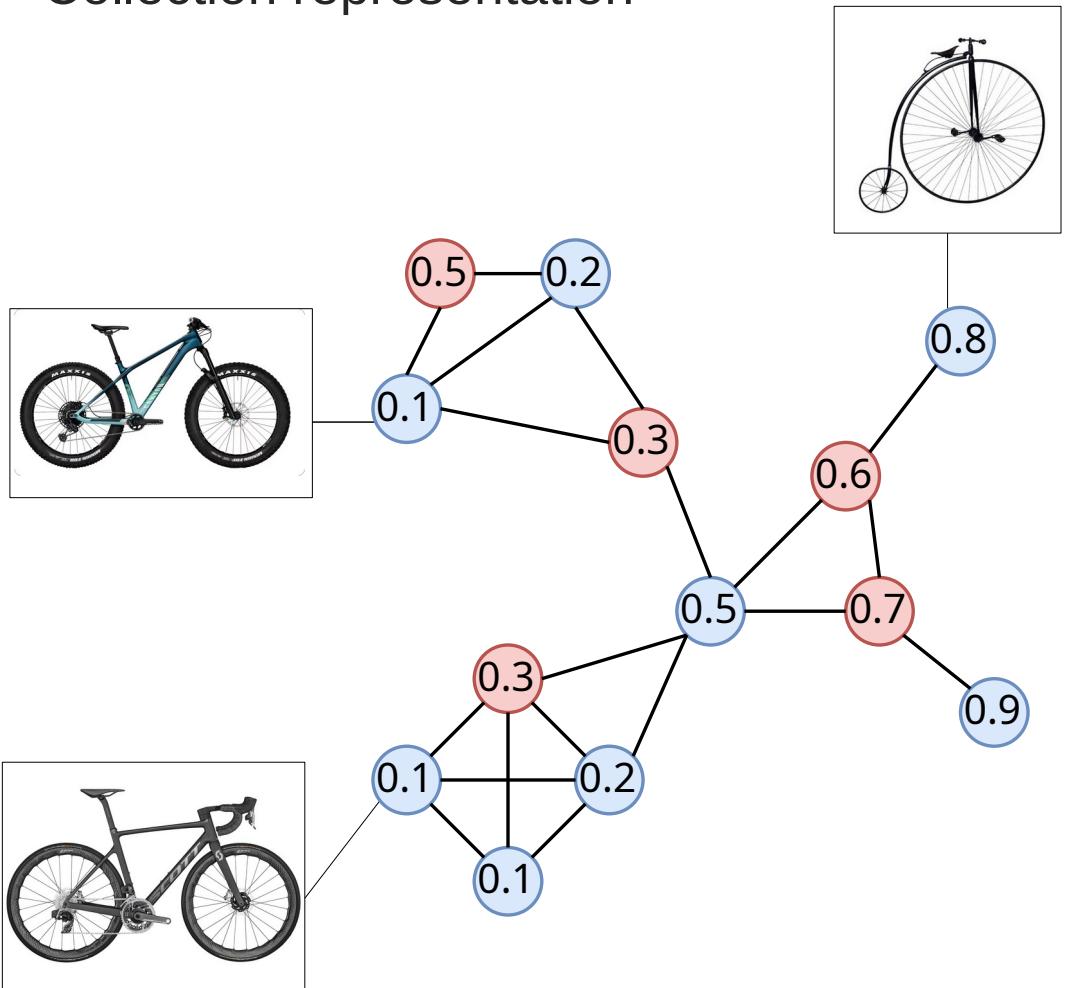
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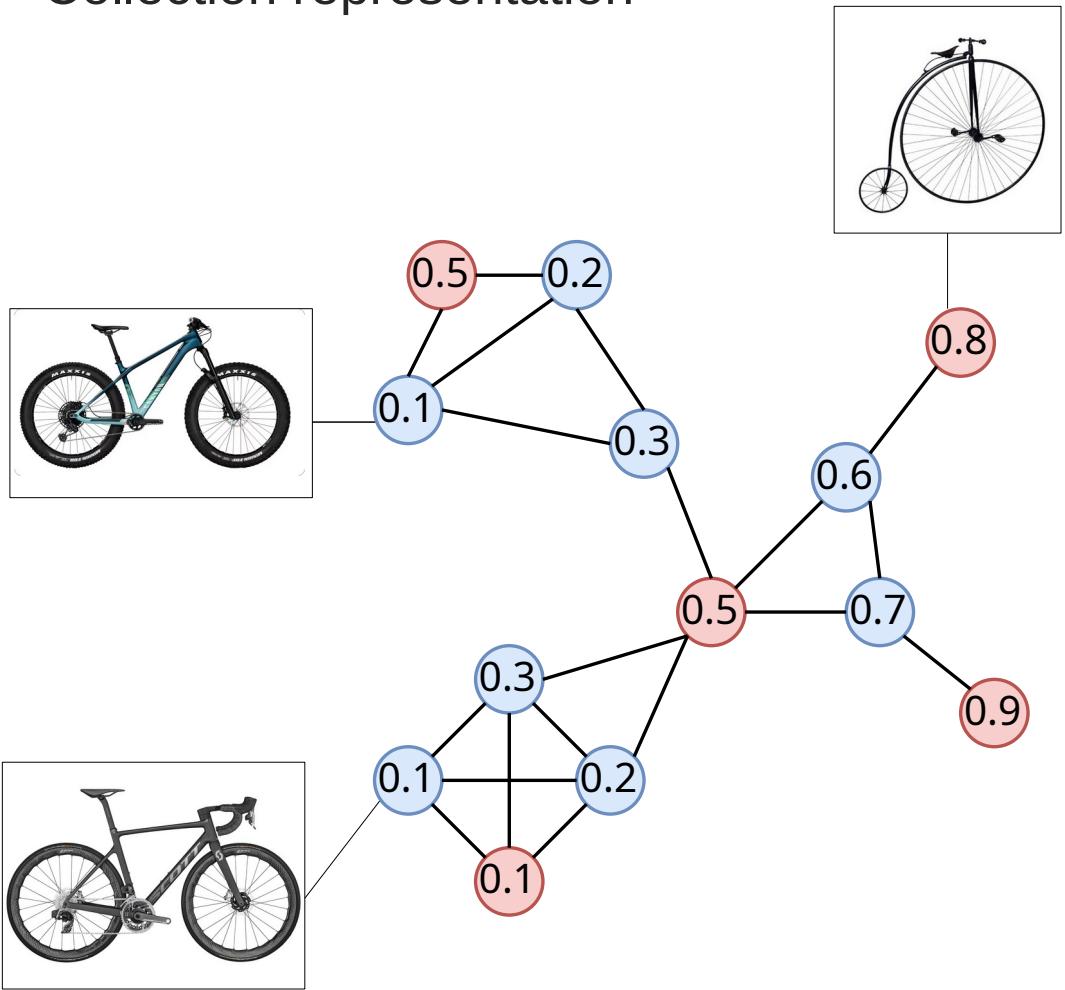
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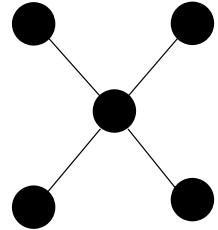


Introduction and Context

Maximum Independent Set (MIS) problem

$$G = (V_s, E_s)$$

$$x_v \in \{0, 1\}$$

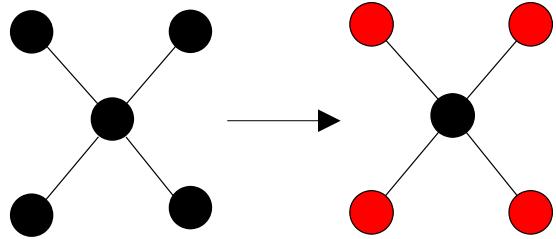


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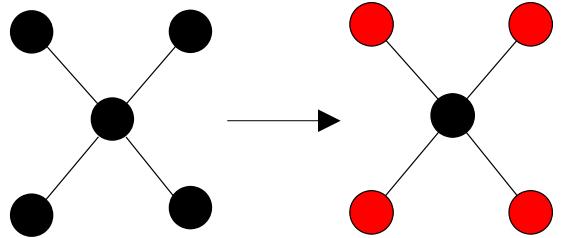


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■ Problem formulation

$$\max \sum_{v \in V_s} x_v$$

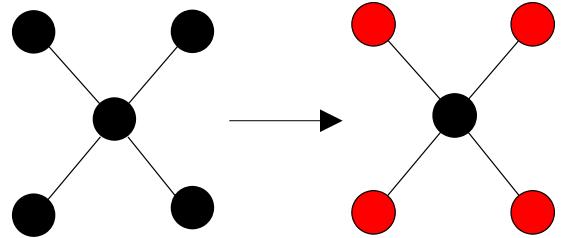
subject to $x_u \neq x_v$ if $(u, v) \in E_s$

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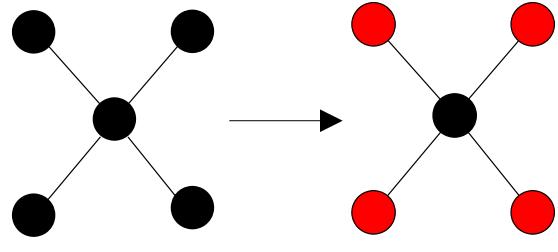
$$\max \sum_{v \in V_s} x_v - 2 \sum_{(u, v) \in E_s} x_u x_v$$

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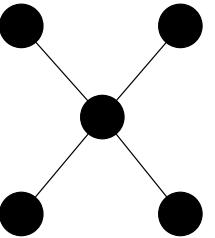
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Max-cut problem

$$G = (V_s, E_s)$$

$$s_v \in \{-1, +1\}$$



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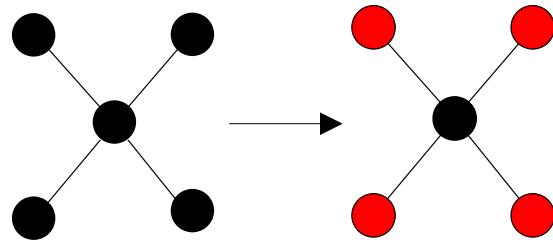
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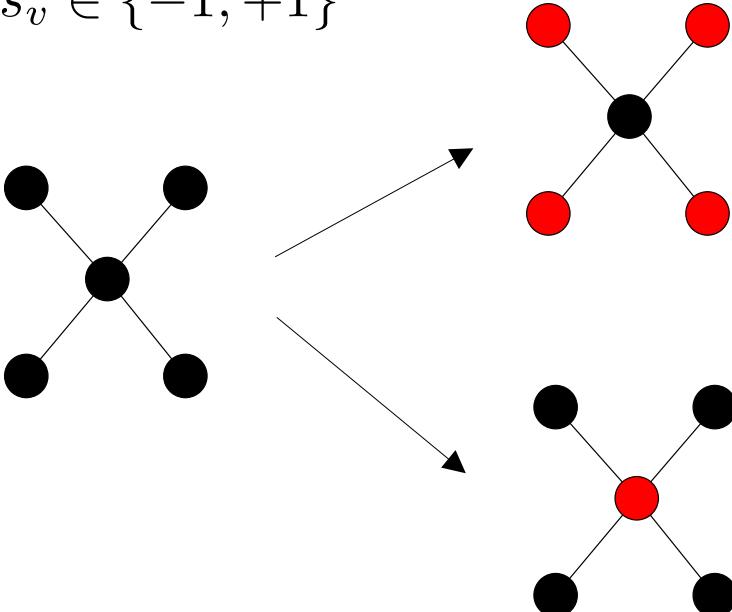
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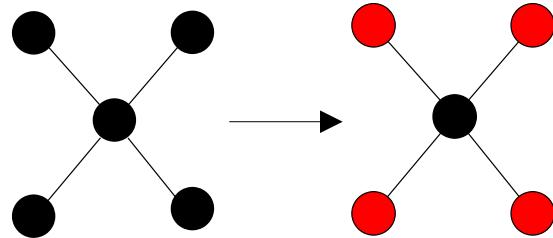


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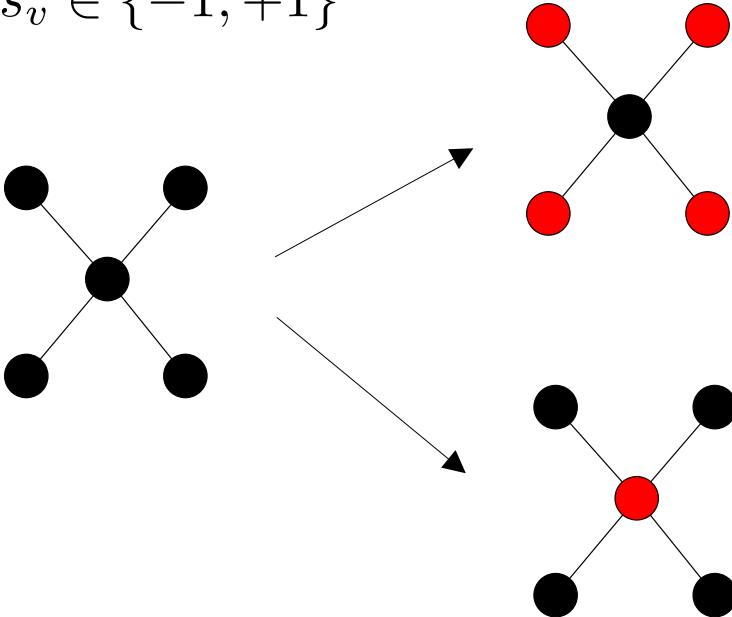
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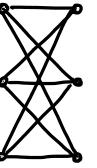
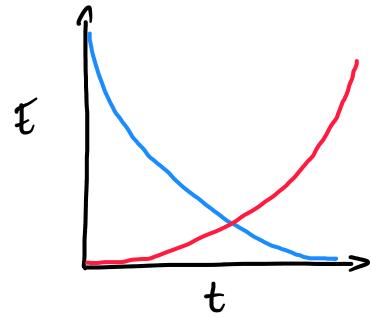
$$\max - \sum_{(u, v) \in E_s} s_u s_v$$



Introduction - Quantum approaches for optimization problems

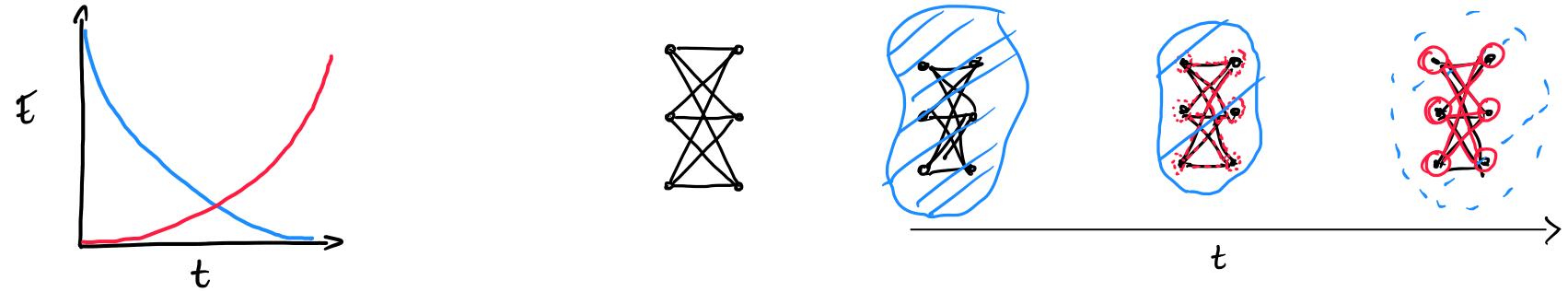
Introduction - Quantum approaches for optimization problems

- Quantum Annealers (analog-based) (NISQ) [KN98]



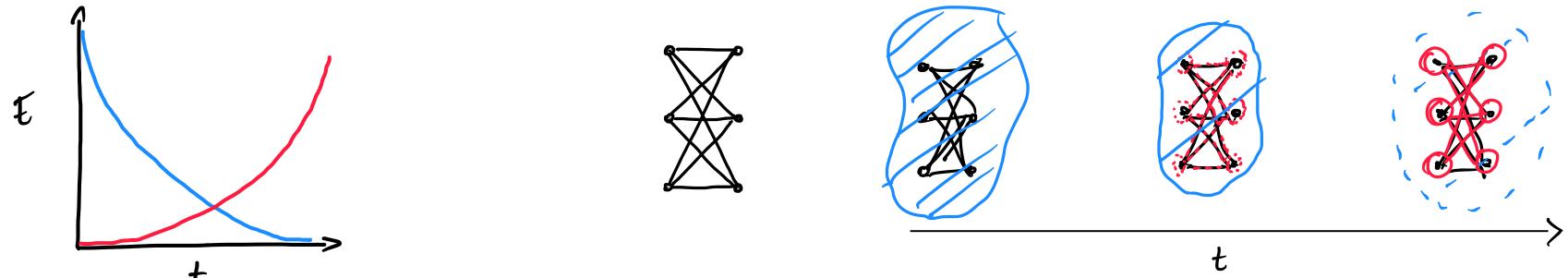
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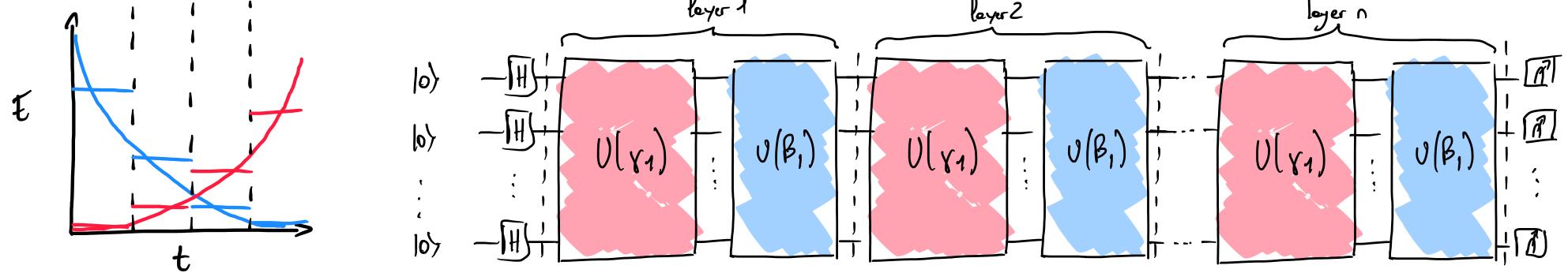


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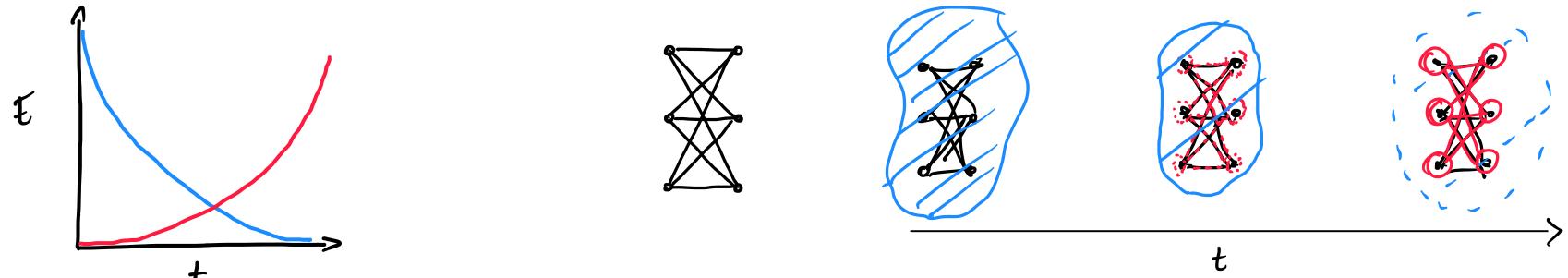


- The Quantum Approximate Optimization Algorithm (gate-based) (NISQ) [FGS14]

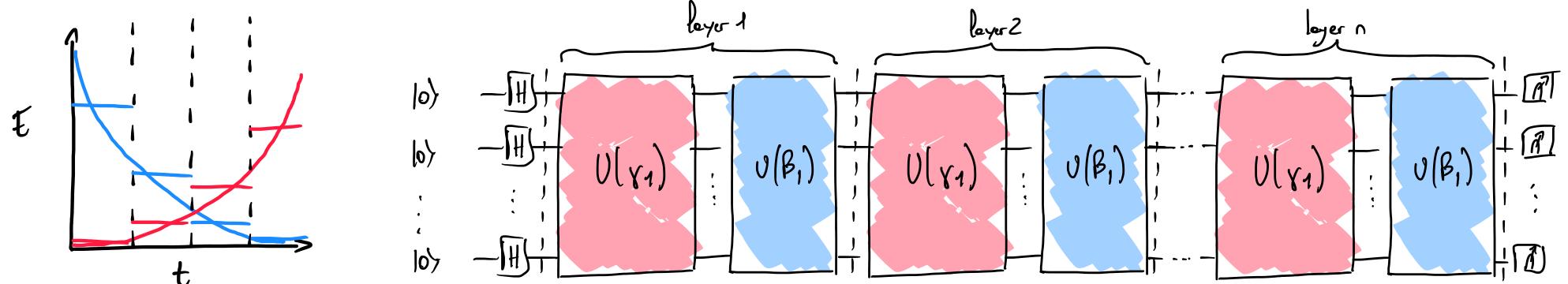


Introduction - Quantum approaches for optimization problems

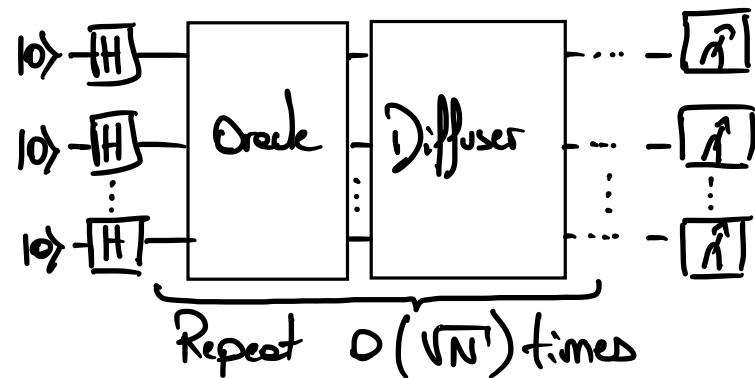
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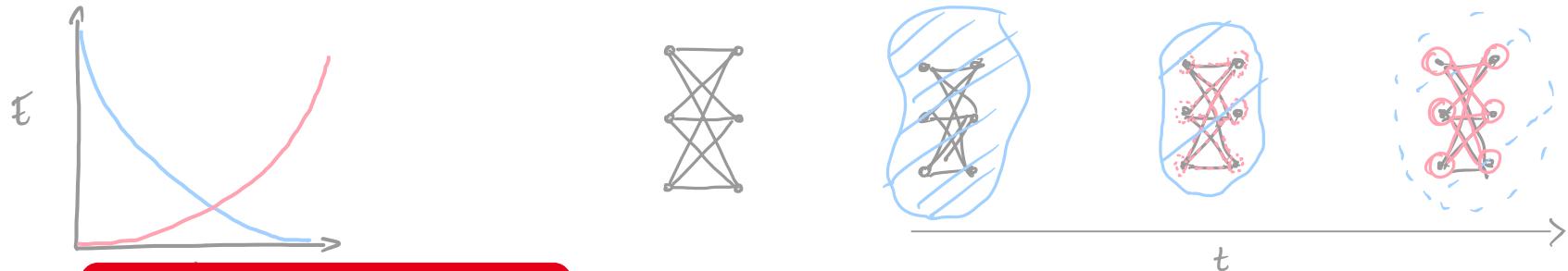


- Exact resolution using quantum circuits (gate-based) (not NISQ) [Gro96]



Introduction - Quantum approaches for optimization problems

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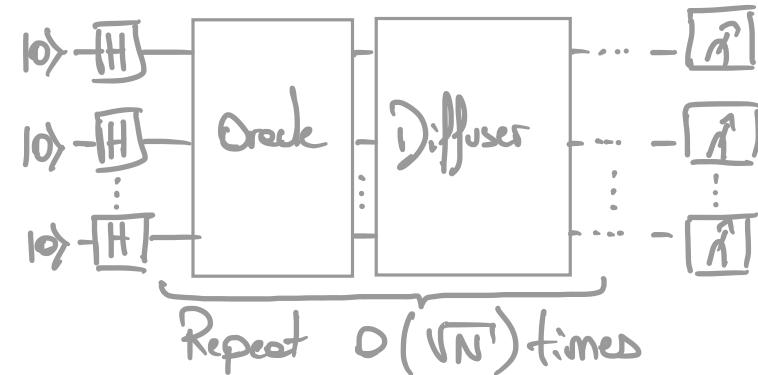


- The Quantum Annealer

Research question

- What makes a Quantum Annealer perform well?
- How to improve their performance?

- Exact resolution using quantum circuits (gate-based) (not NISQ) [Gro96]

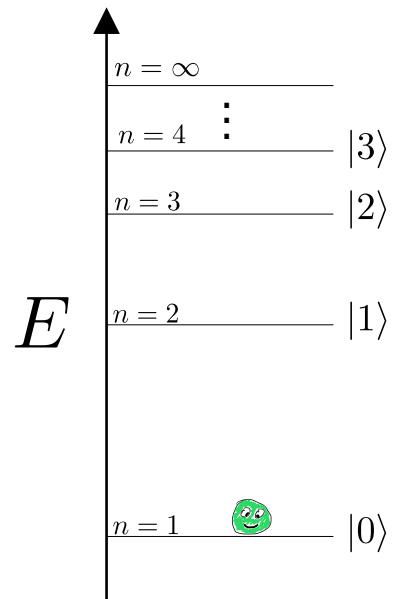




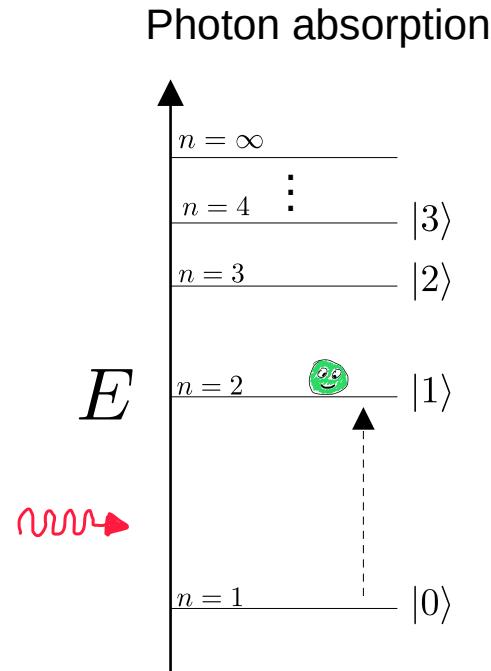
1 ■ Introduction to Quantum Annealing



1- Ground state



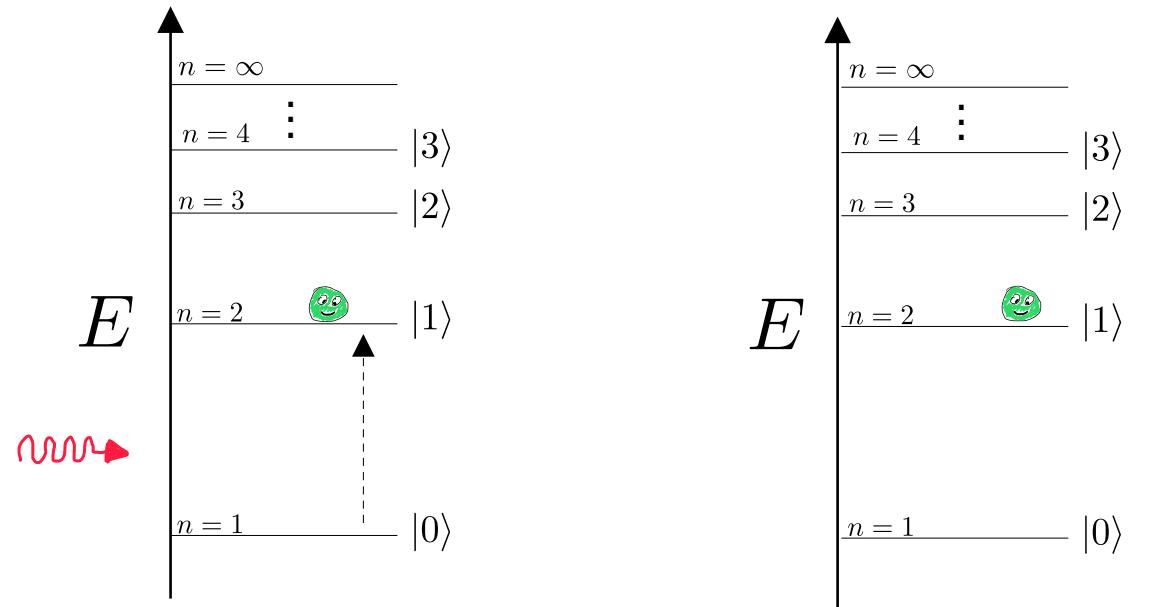
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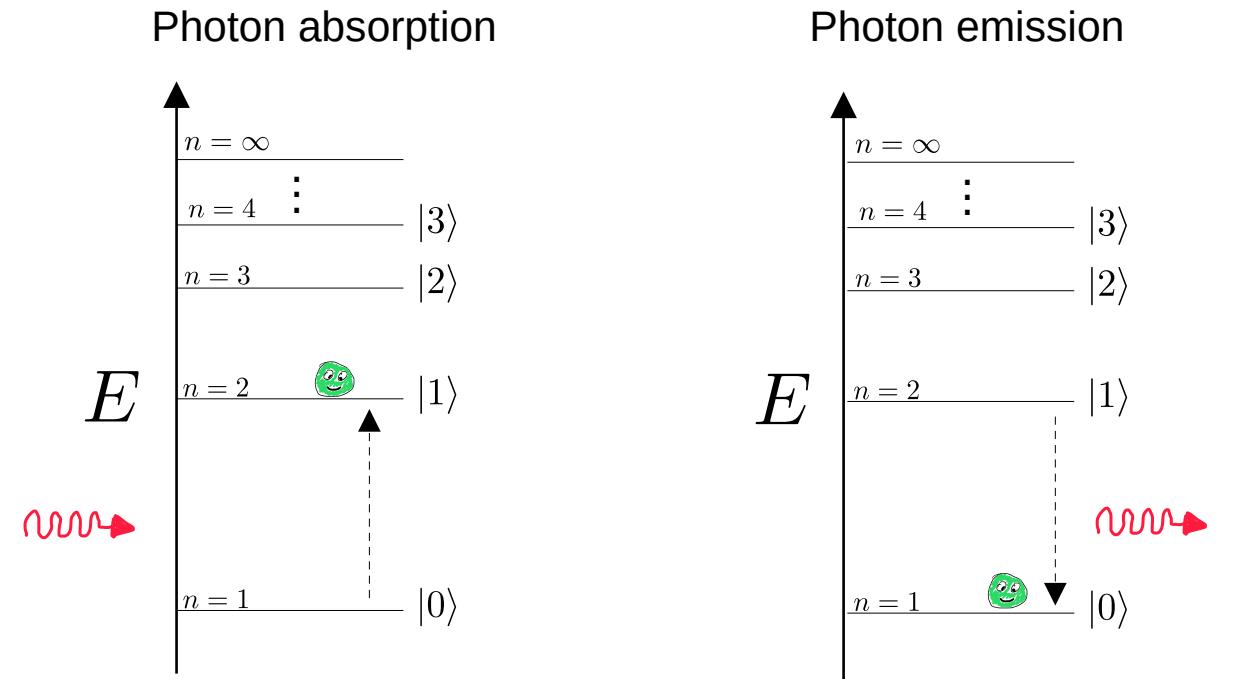


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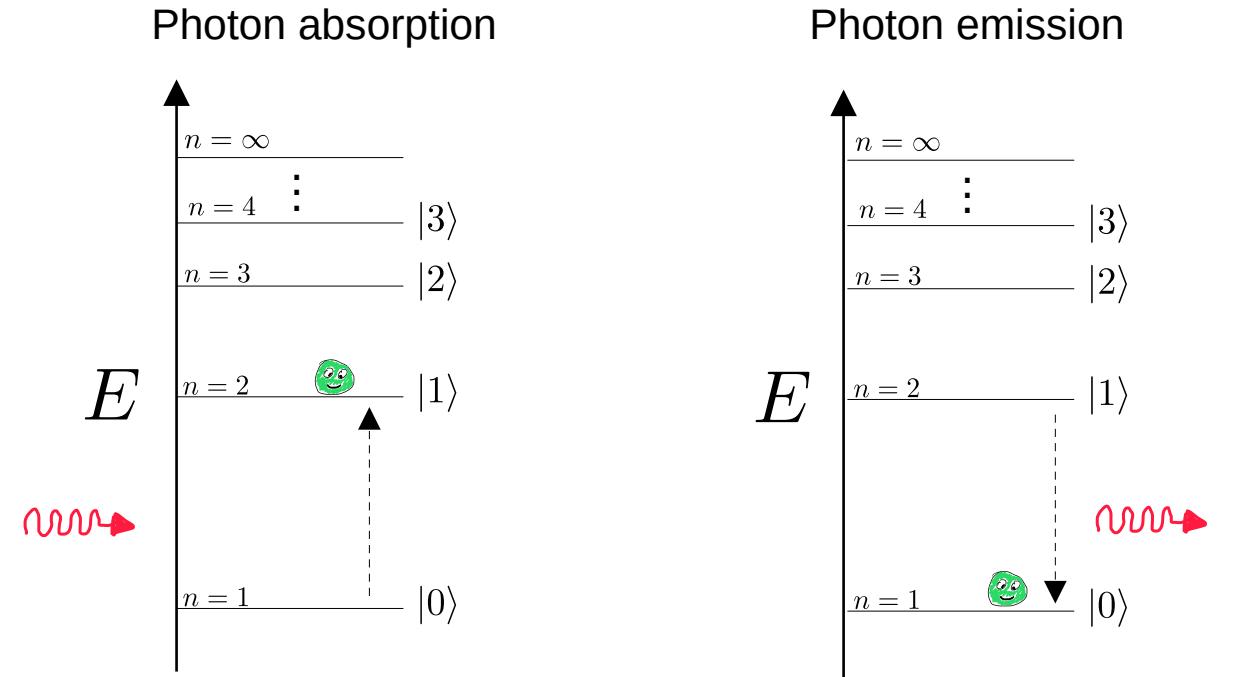
Photon absorption



1- Ground state



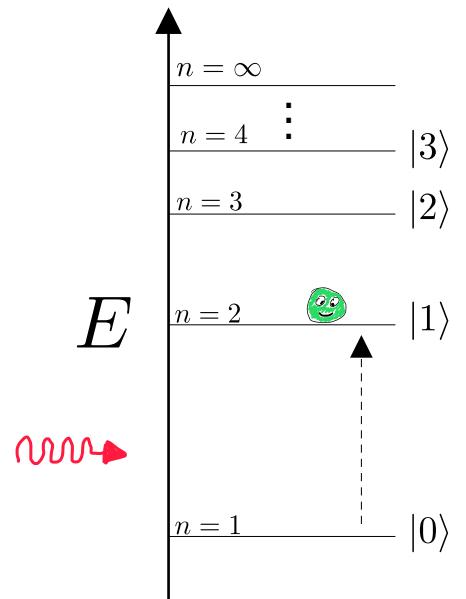
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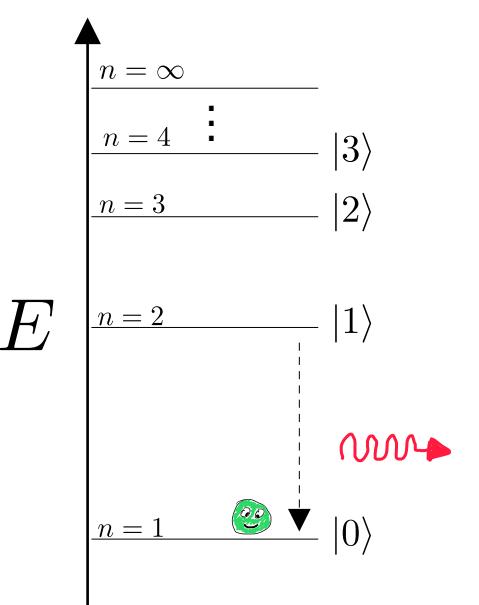
1- Ground state

- Hamiltonian
Energy of distinguishable states

Photon absorption

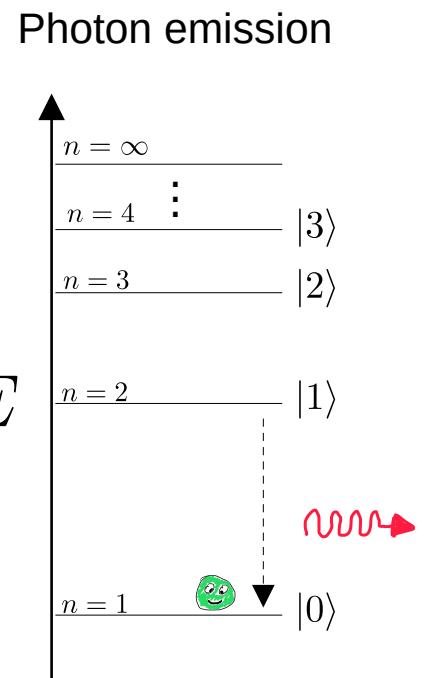
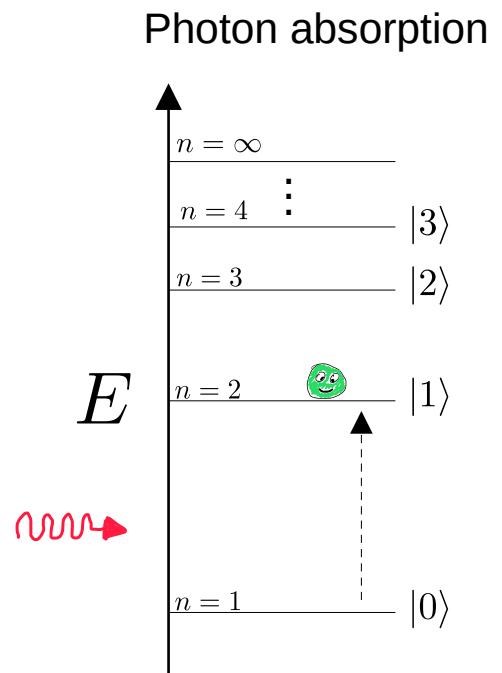


Photon emission





1- Ground state

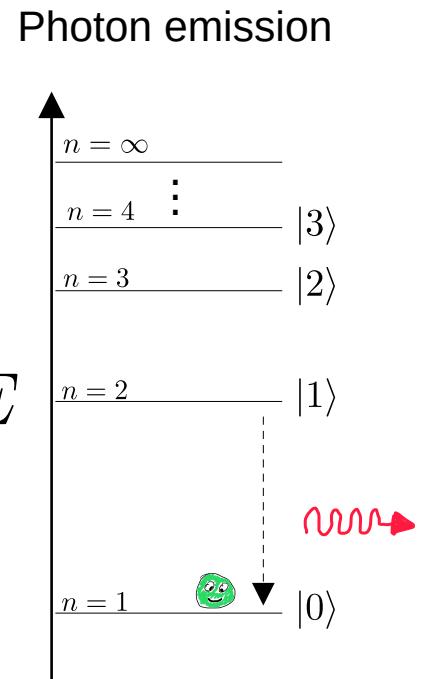
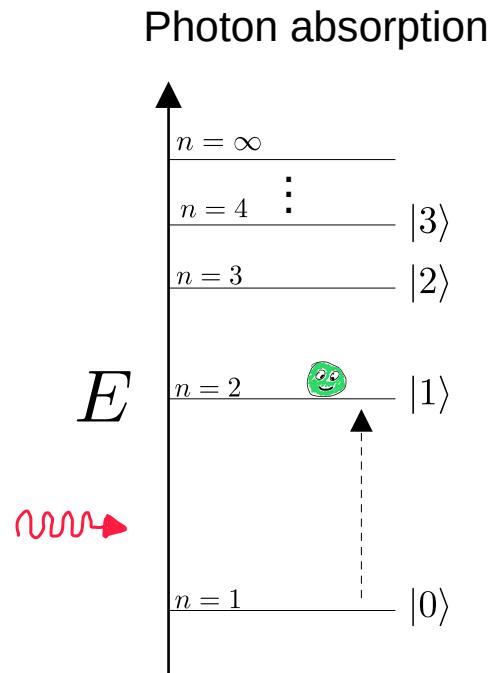


- Hamiltonian
Energy of distinguishable states

$$H |0\rangle = E_0 |0\rangle \quad |0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$



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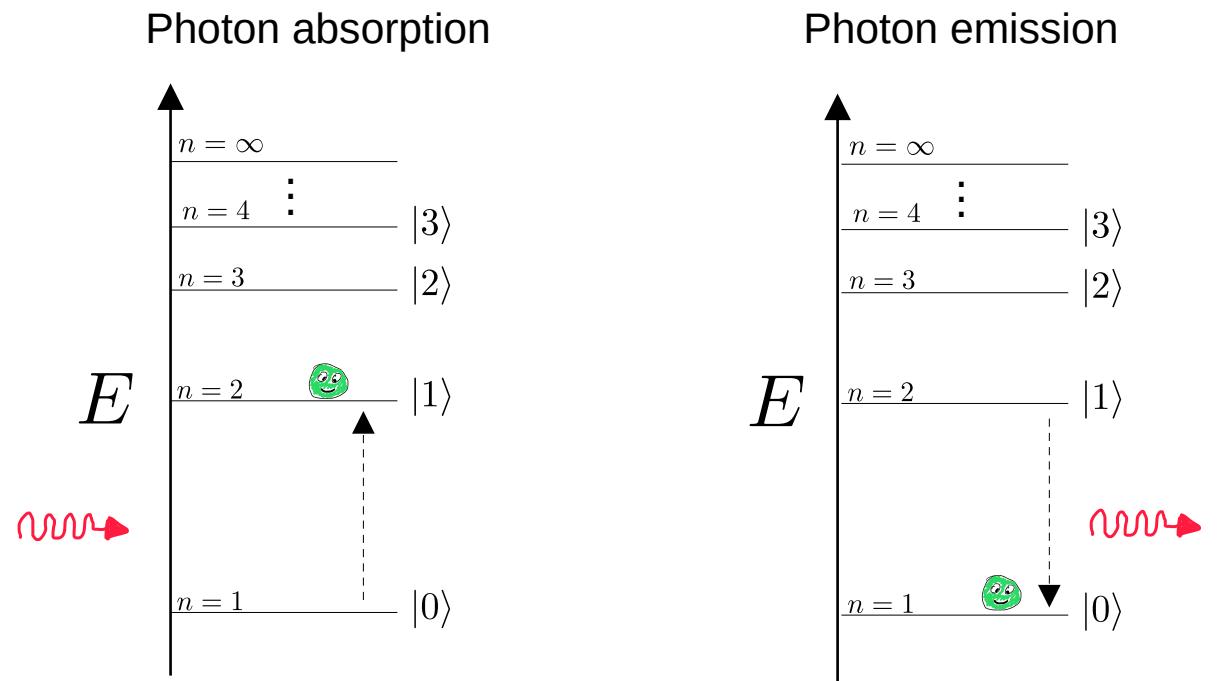


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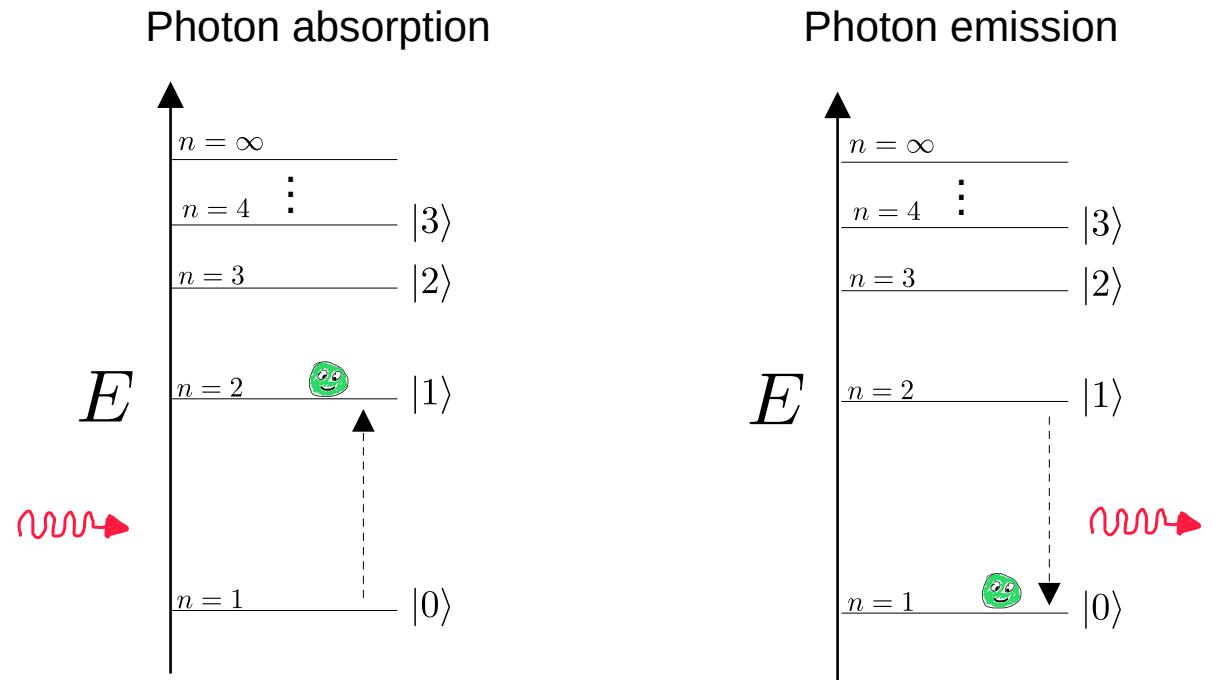
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- n -qubit Hamiltonian

1- Ground state



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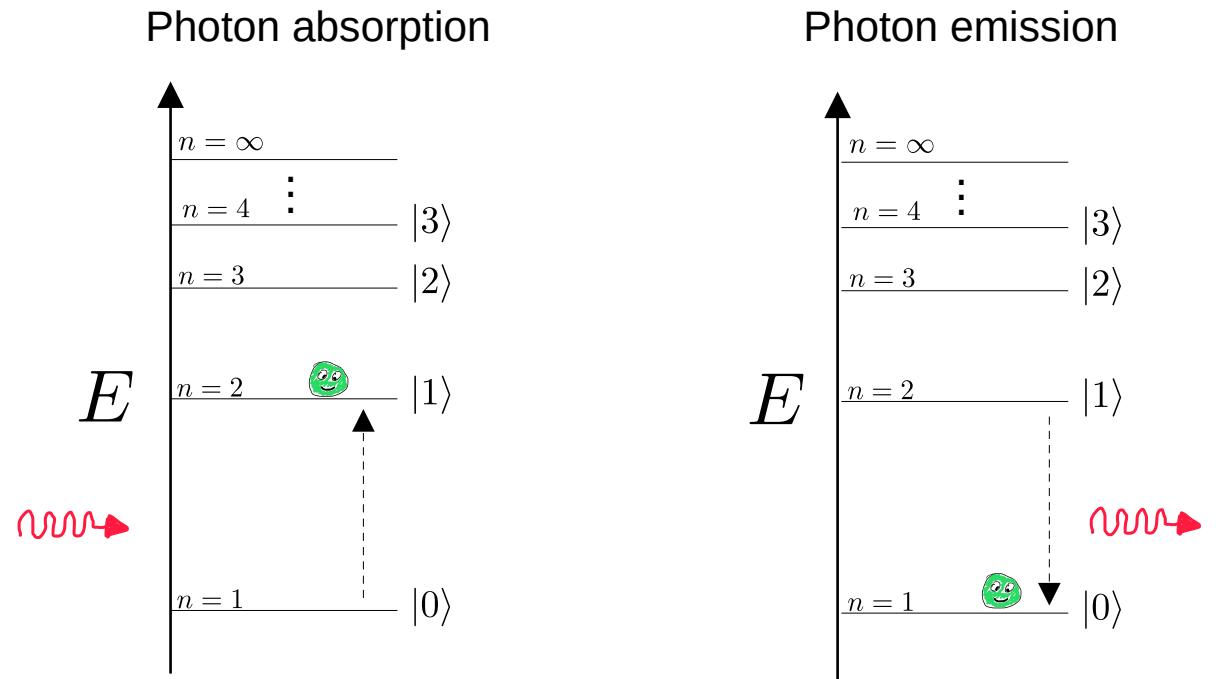
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$$H |\psi_p\rangle = E_p |\psi_p\rangle$$

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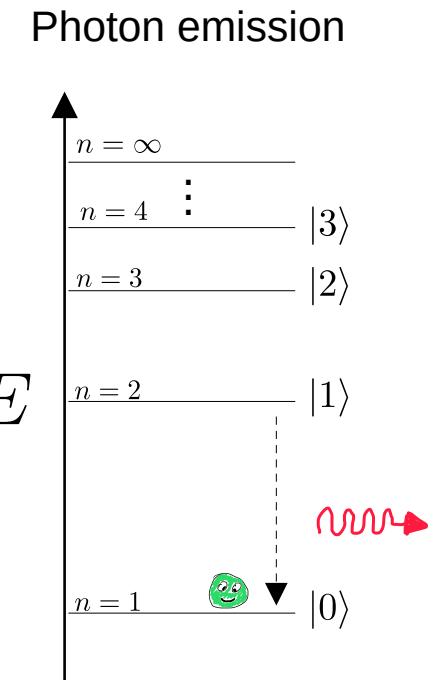
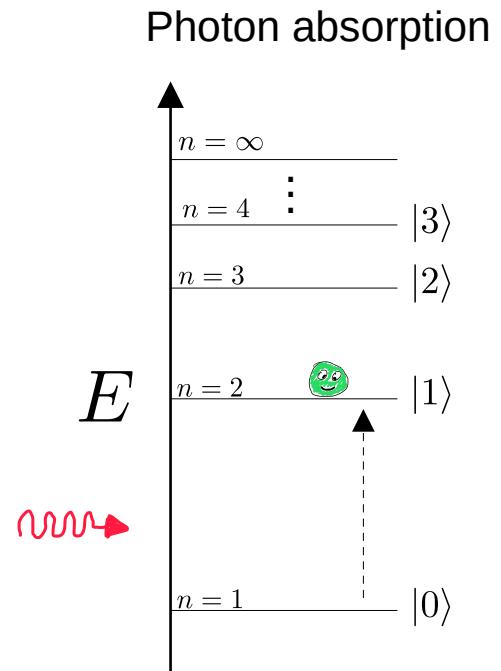
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- n -qubit Hamiltonian

$$H |\psi_p\rangle = \boxed{E_p} \boxed{|\psi_p\rangle}$$

eigenvalue eigenvector

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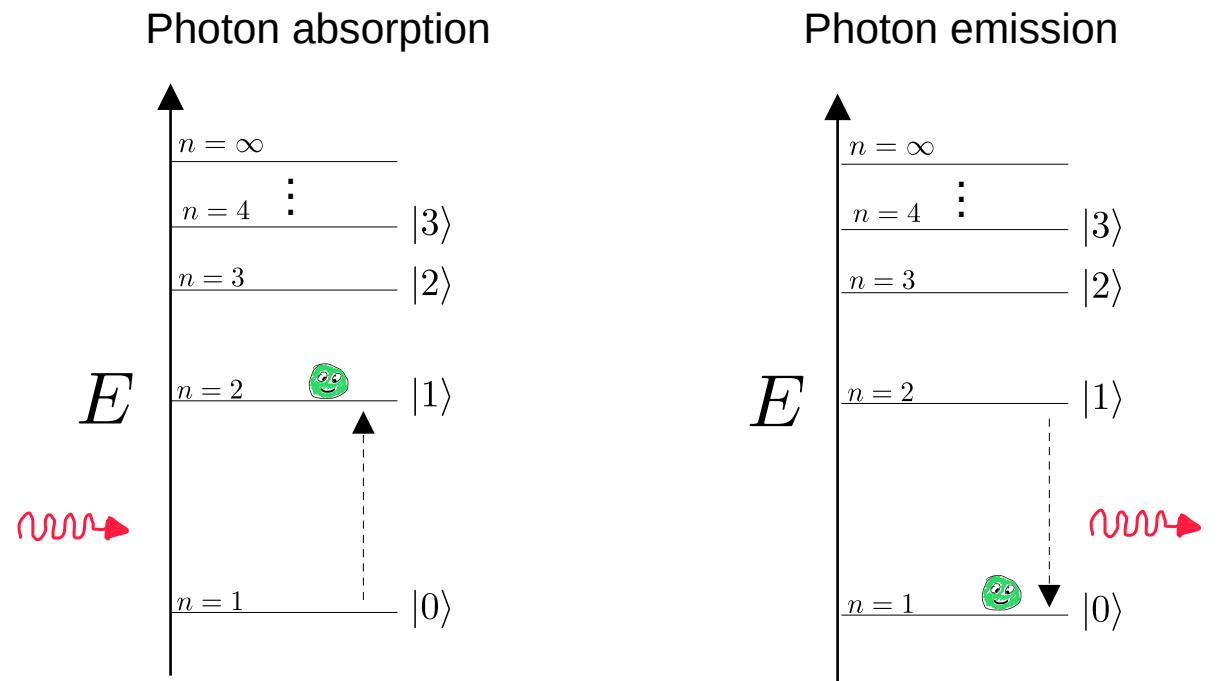
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- Time-dependent Hamiltonian

1- Ground state



- Hamiltonian
 - Energy of distinguishable states

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- ### ■ n -qubit Hamiltonian

$$H |\psi_p\rangle = E_p |\psi_p\rangle$$

The diagram illustrates the decomposition of a state vector. A red bracket groups the term $E_p |\psi_p\rangle$. Two red lines point from this bracket to the words "eigenvalue" and "eigenvector" located below the equation.

- ## ■ Time-dependent Hamiltonian

$$H(t) |\psi_p(t)\rangle = E_p(t) |\psi_p(t)\rangle$$



1- Adiabatic theorem

Quantum Adiabatic Theorem

“A quantum system described by a time-dependent Hamiltonian $H(t)$ initially prepared in an eigenstate $|\psi_p(0)\rangle$ of $H(0)$ (e.g., the ground state), will approximately remain in the instantaneous eigenstate $|\psi_p(t)\rangle$, given that t varies sufficiently slowly”



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- Linear interpolation of Hamiltonians:

$$H(t) = -A(t)H_{init} + B(t)H_{final}$$

$$t \in [0, T]$$



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- Assumption

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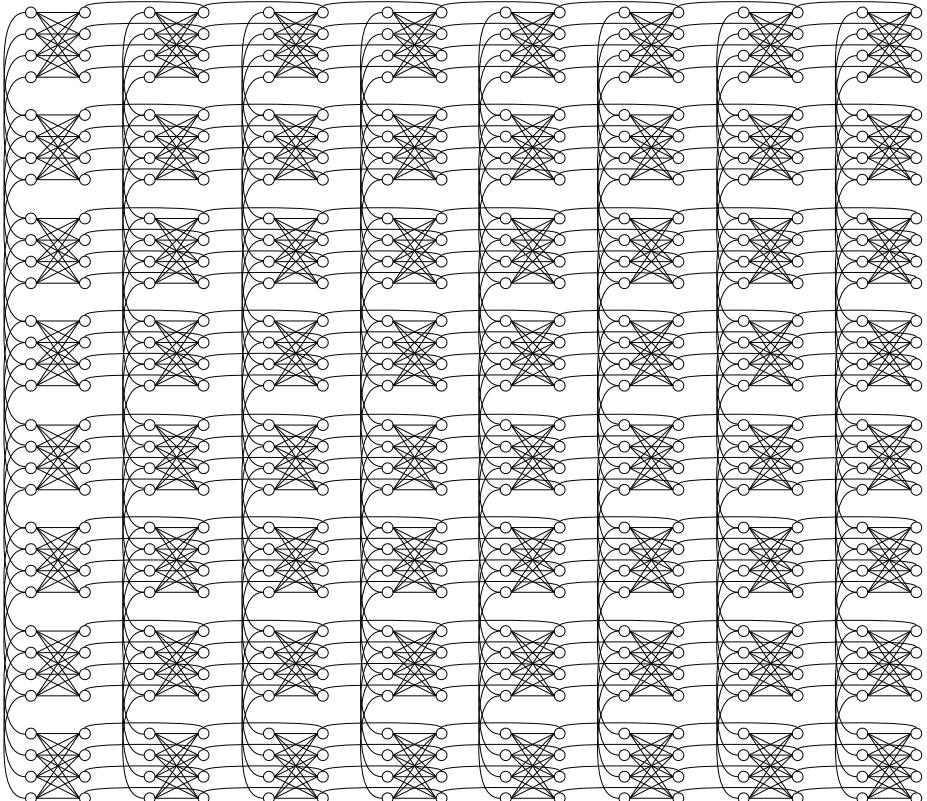
- Ideal closed quantum system

- Theoretical advantage:

- Universal (as quantum circuit)
 - Performance depends on the « sufficiently slowly »

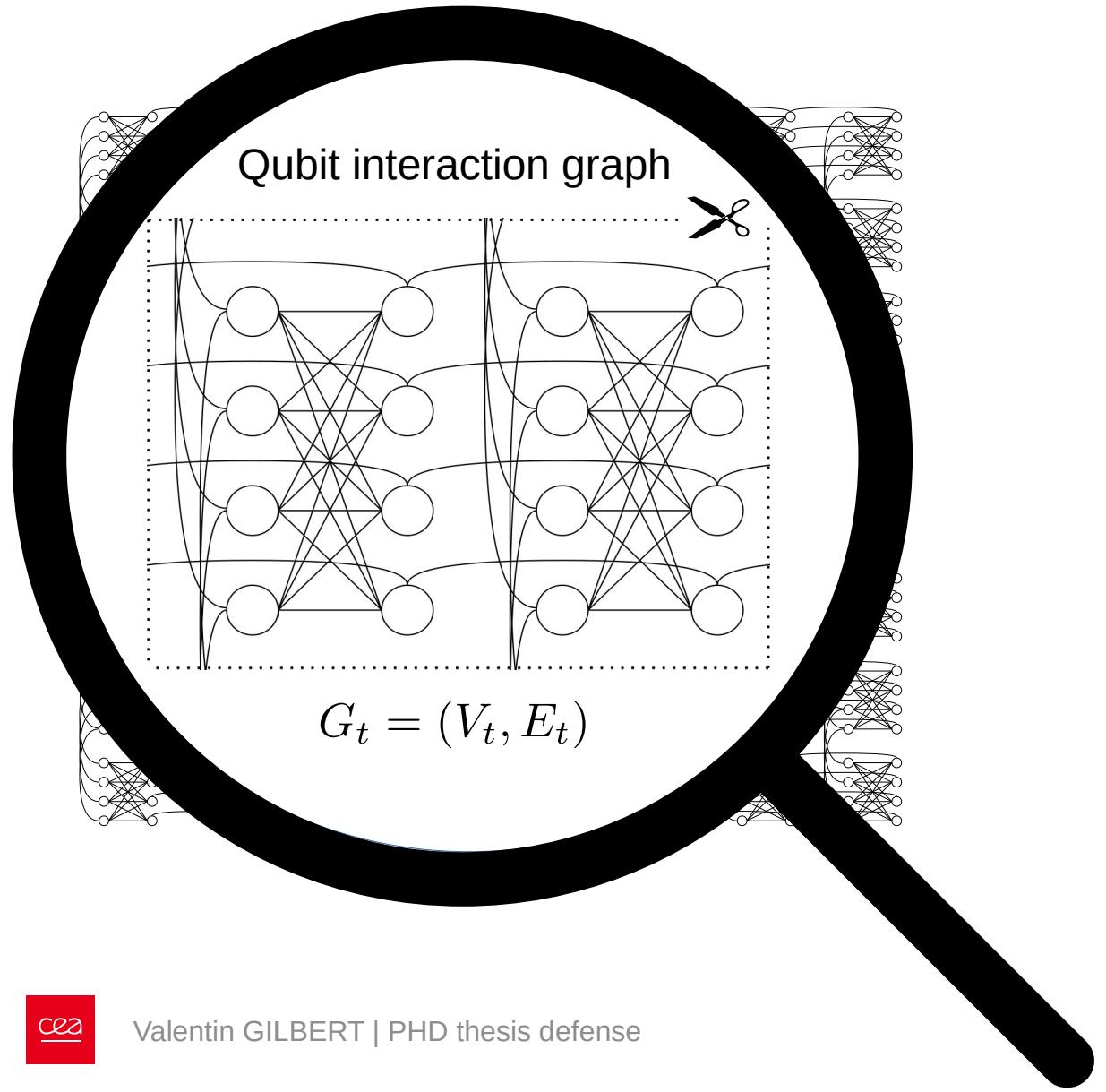


1- Quantum Annealers (D-Wave systems)



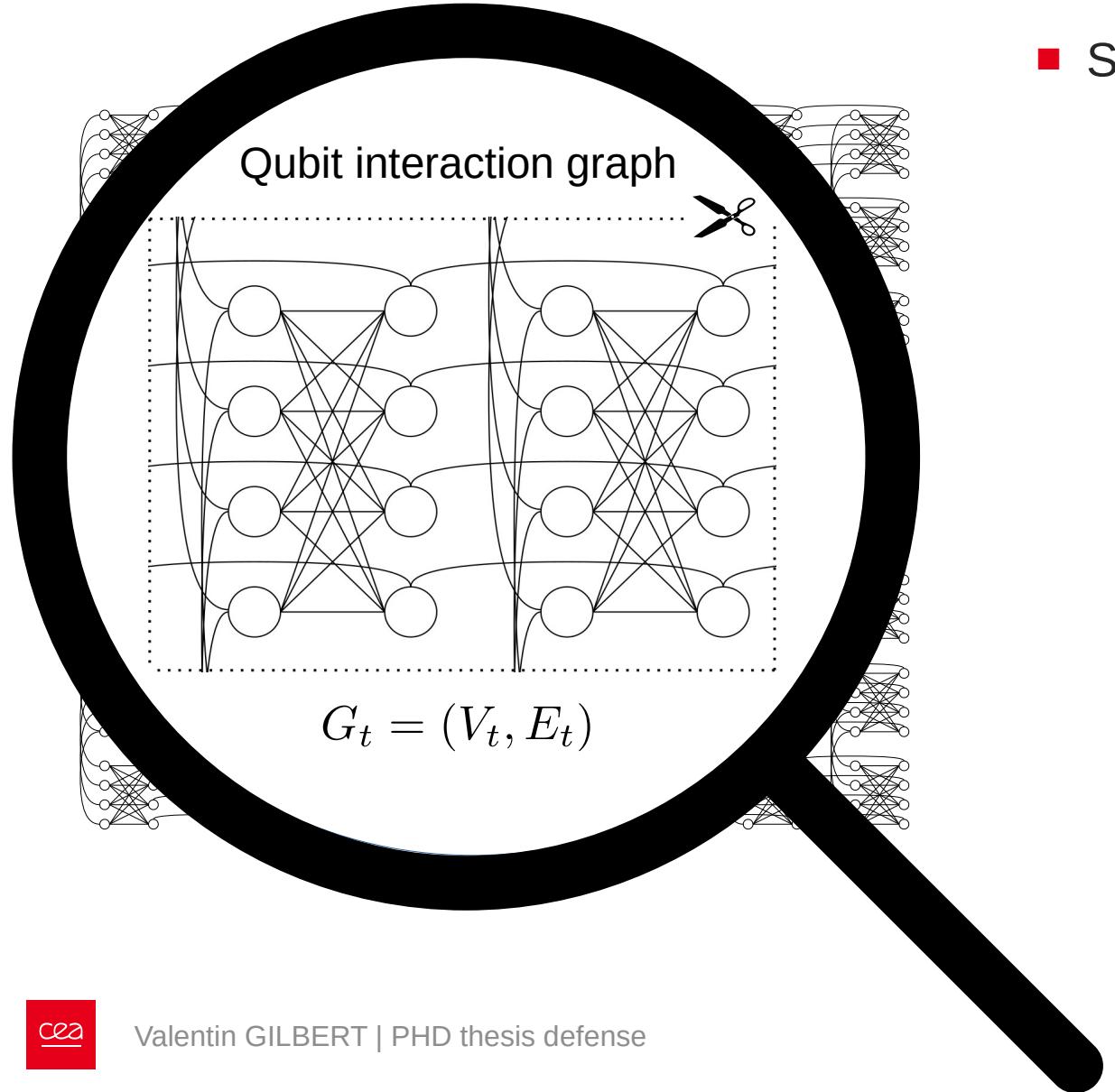


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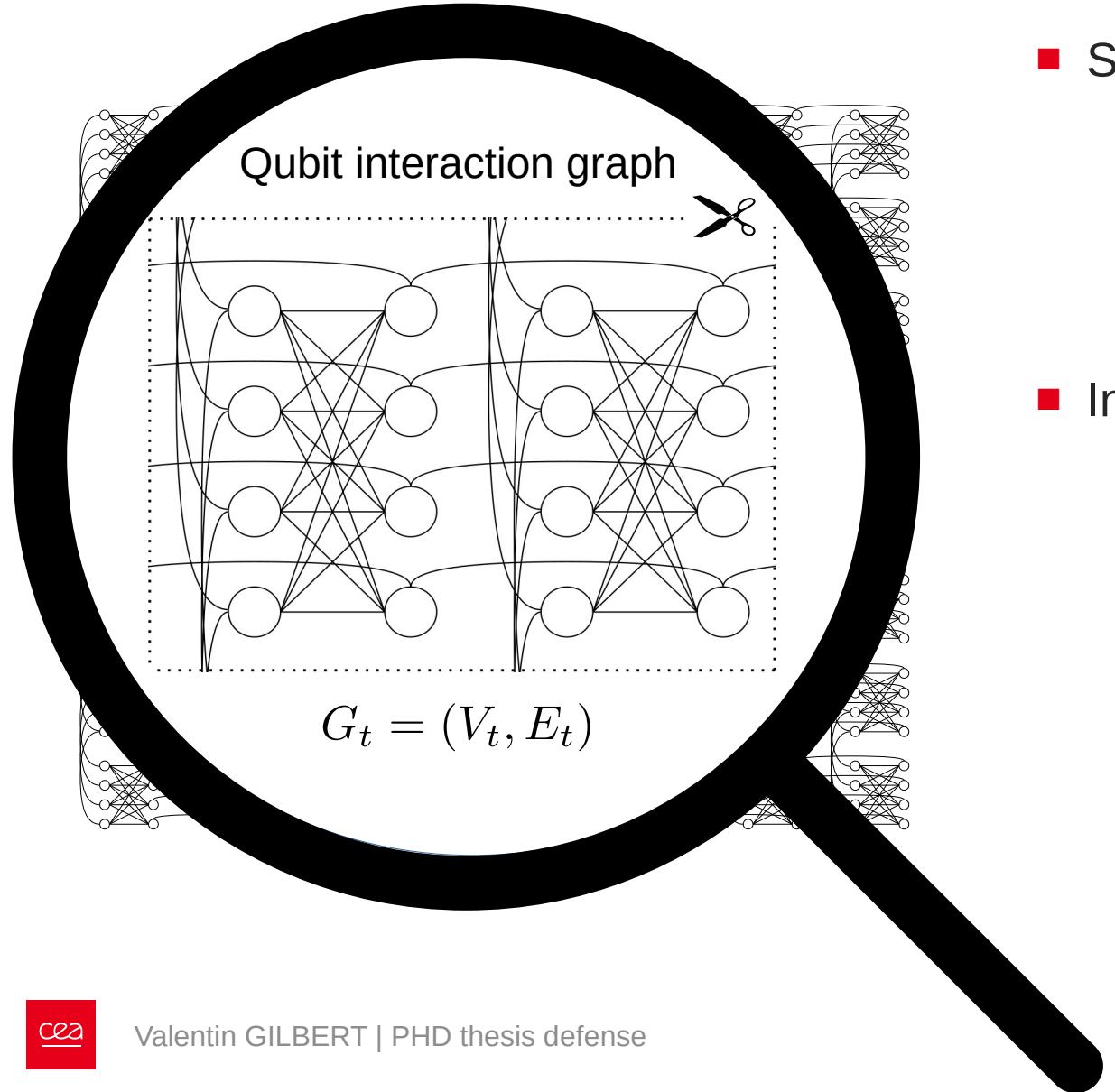


- System Hamiltonian (Transverse Field Ising model)

$$H(s) = -A(s)H_{init} + B(s)H_{final}$$

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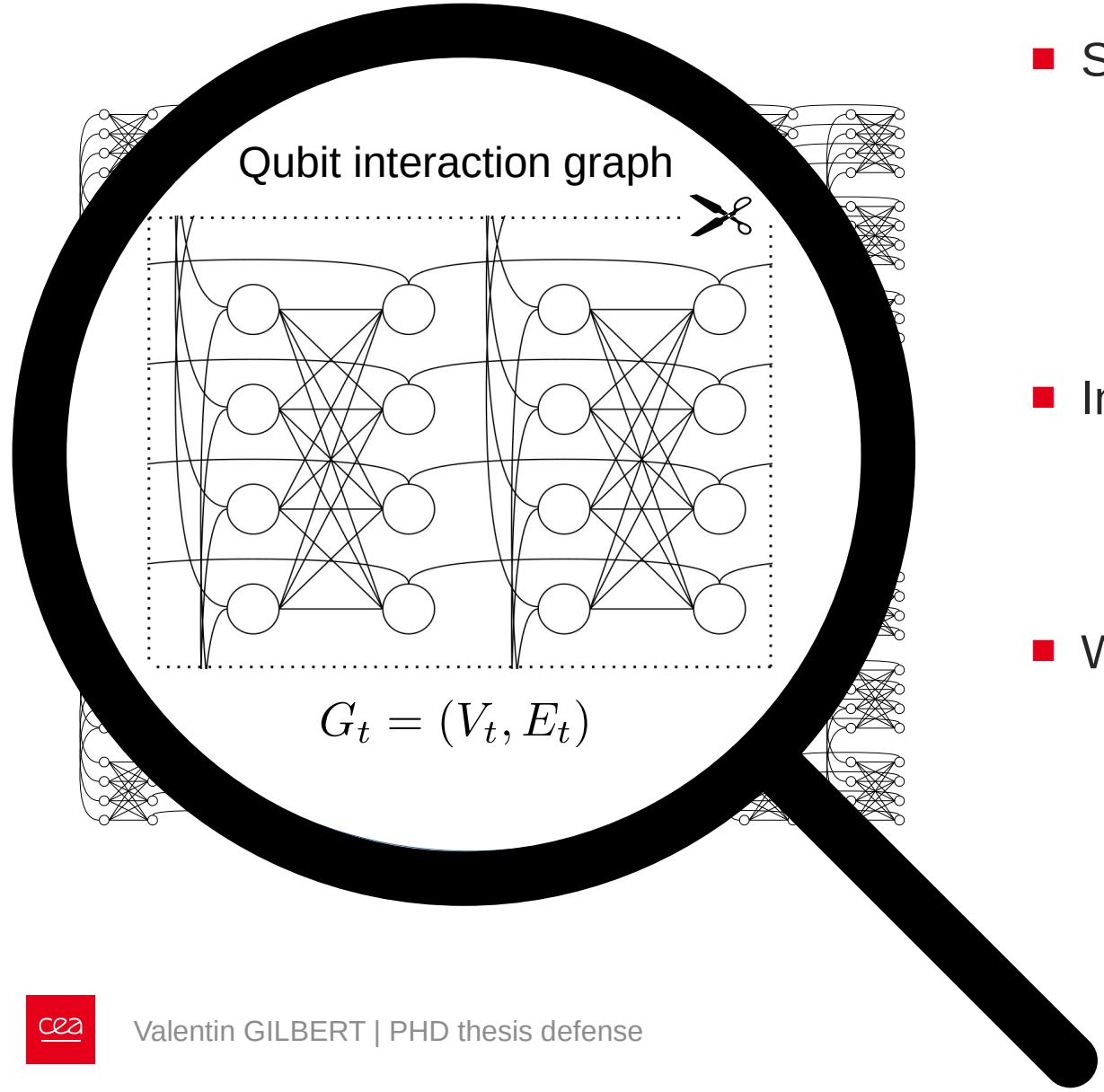
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- What eigenstate minimizes the expression $-A(s) H_{init}$?

$$|\psi\rangle = \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} |x\rangle$$



1- Quantum Annealers

■ Problem Hamiltonian

$$H_{final} = \sum_{v \in V_t} h_v \sigma_v^z + \sum_{(u,v) \in E_t} J_{(u,v)} \sigma_u^z \sigma_v^z$$

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- Ising cost function correspondence

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$$s_v \in \{-1, +1\}$$

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$$H_{final} = \sum_{v \in V_t} h_v \sigma_v^z + \sum_{(u,v) \in E_t} J_{(u,v)} \sigma_u^z \sigma_v^z$$

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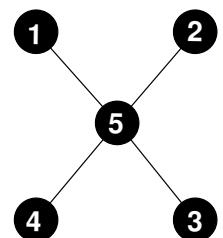
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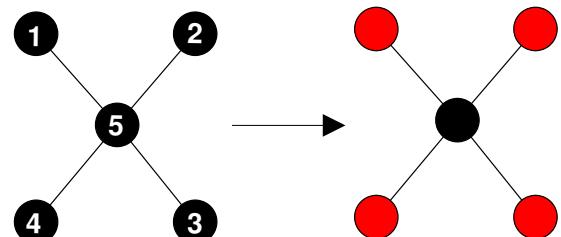
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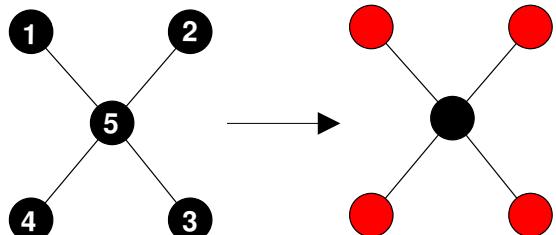
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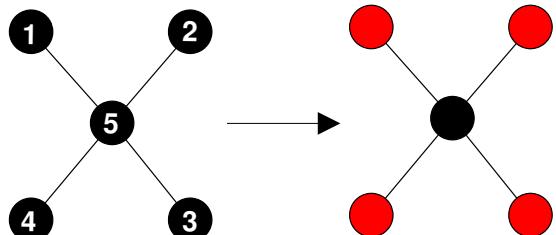
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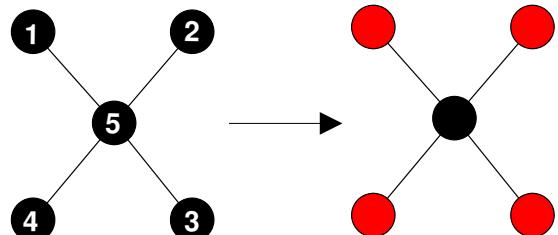
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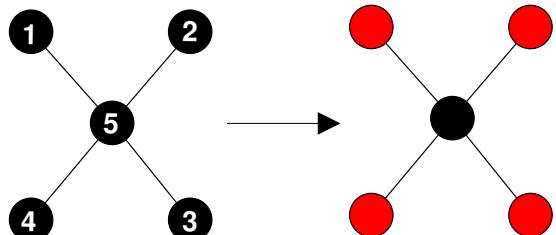
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$$H_{final} = -\sigma_5^z + \frac{1}{2} (\sigma_1^z \sigma_5^z + \sigma_2^z \sigma_5^z + \sigma_3^z \sigma_5^z + \sigma_4^z \sigma_5^z)$$



1- Quantum Annealers

- “Sufficiently slow ?” Analysis of the spectral gap [AL18]:

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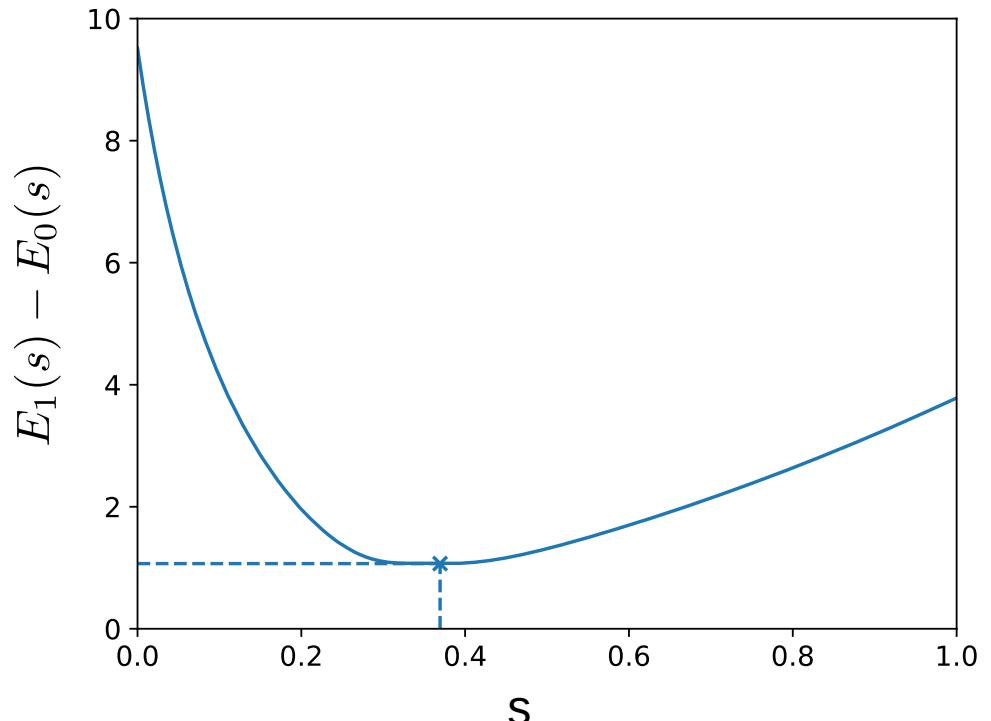
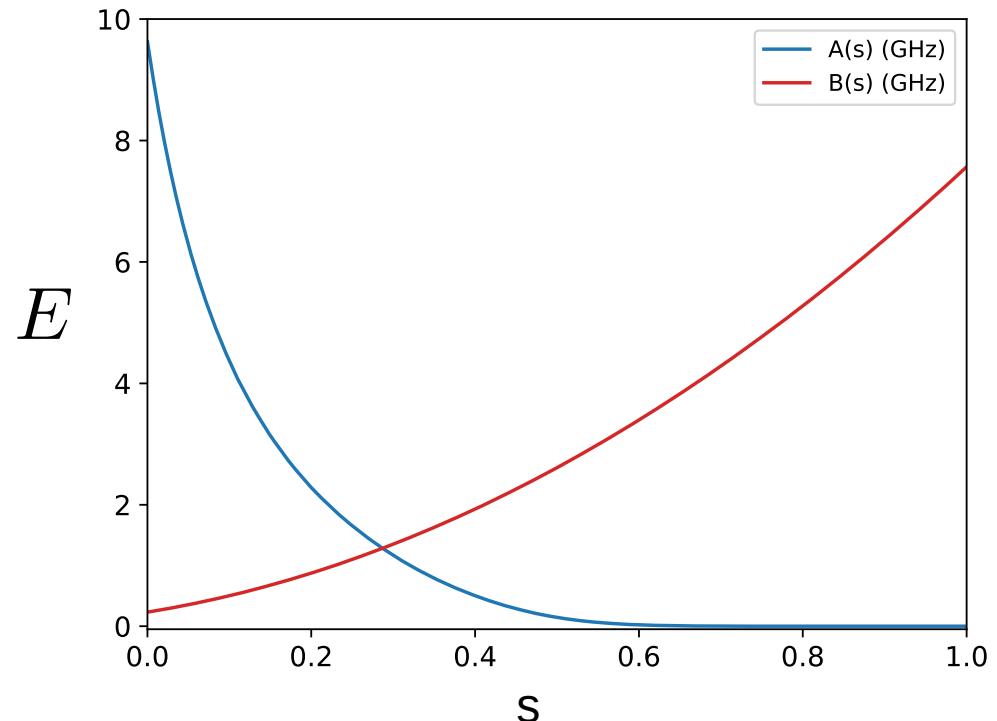


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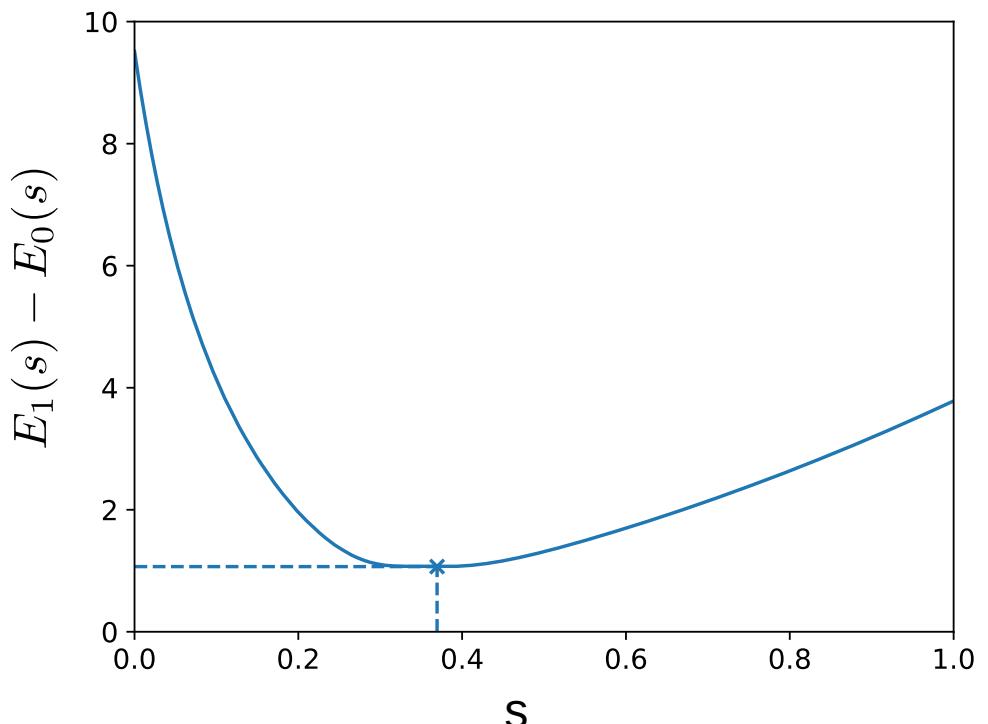
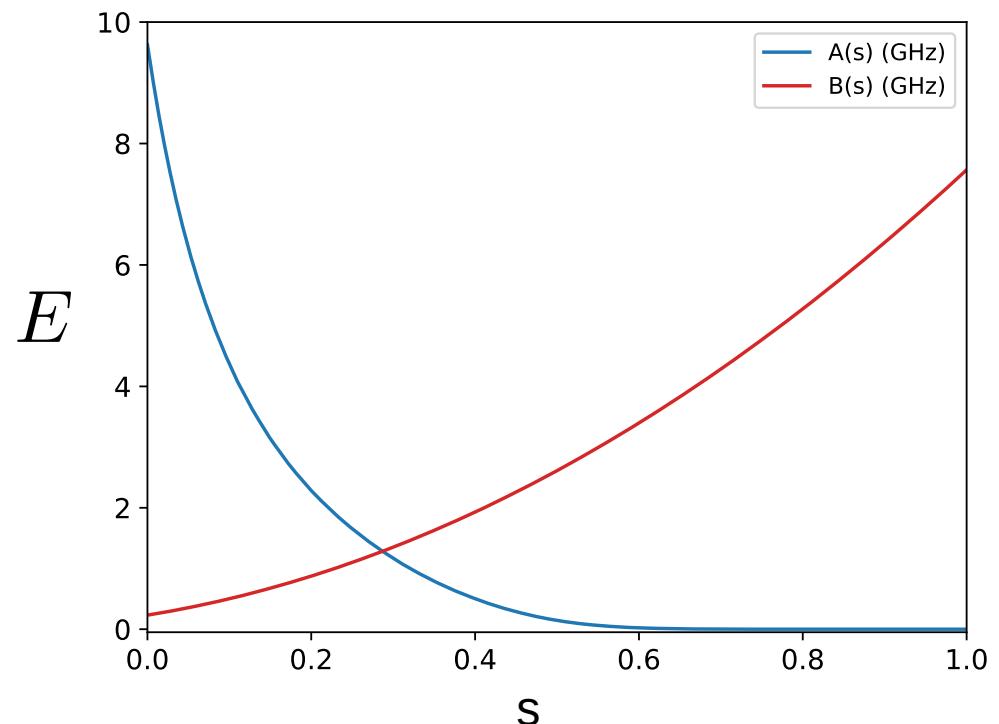
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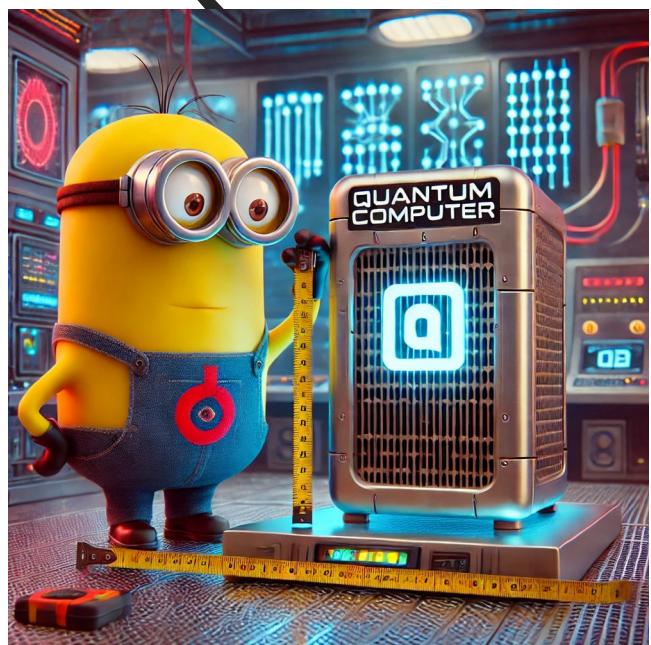
$$g_{\min} = \min_{0 \leq s \leq 1} E_1(s) - E_0(s) \quad T \propto O(1/g_{\min}^2)$$



Contribution #1

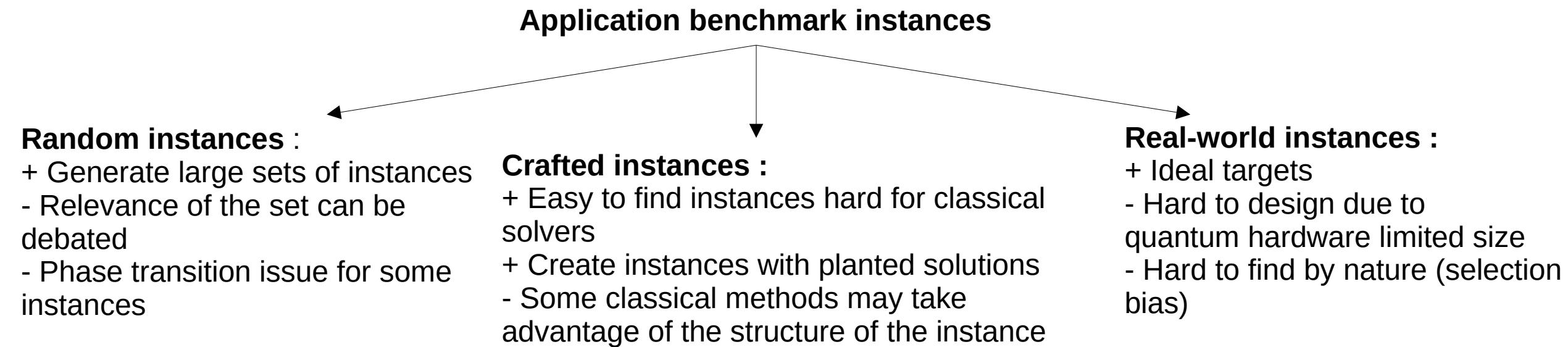
Performance evaluation of Quantum Annealers

2



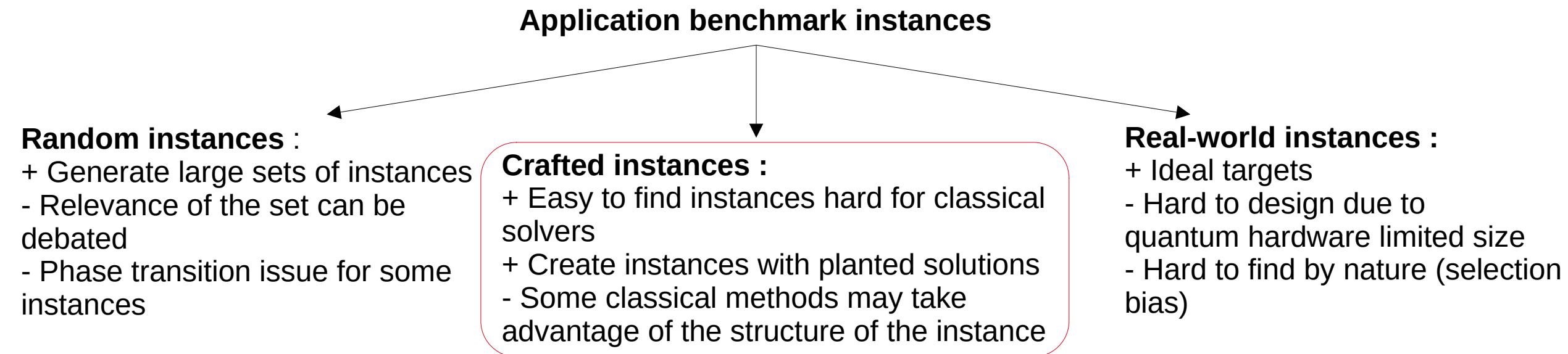


2- Performance evaluation of Quantum Annealers

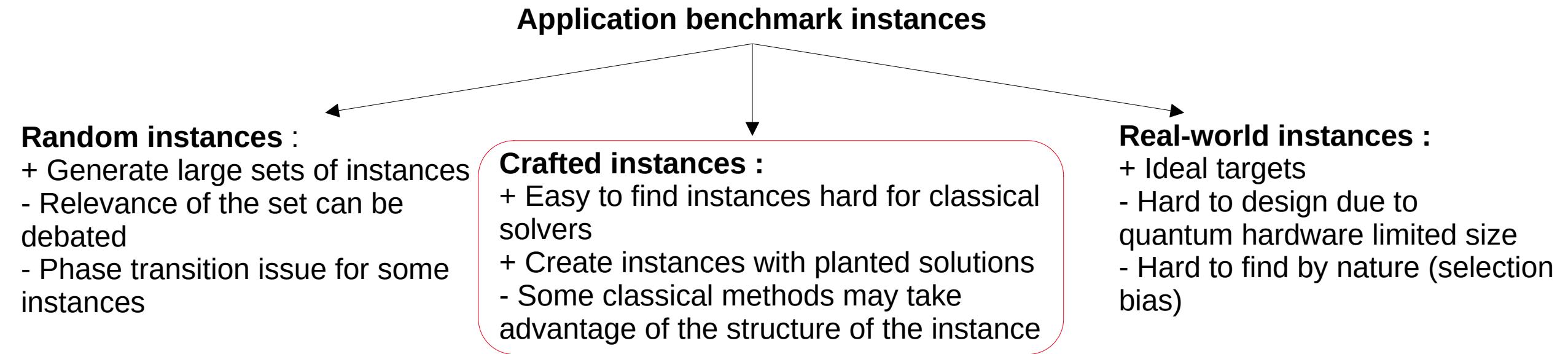




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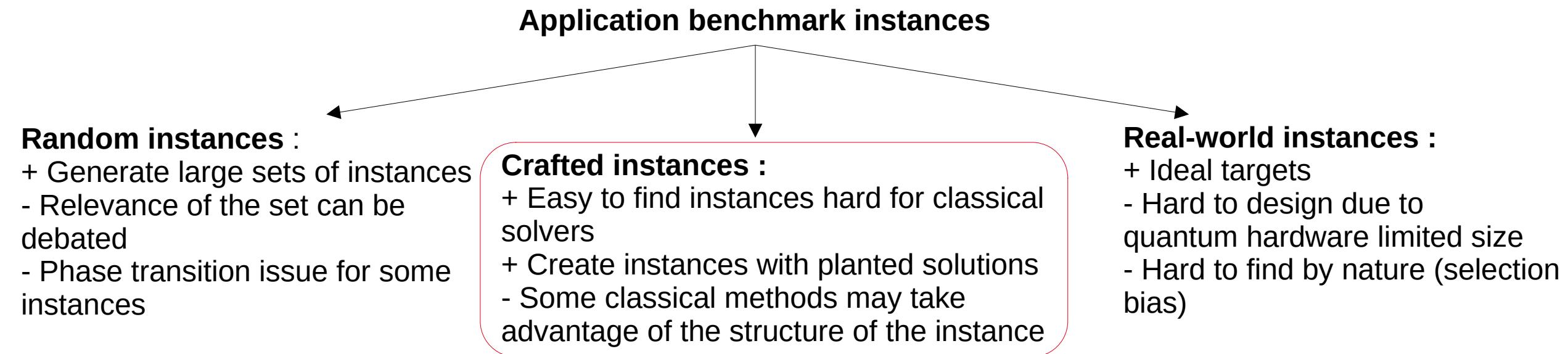
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Source	Instance topology	#Variables	QPU	Embedding
2021 [PCL ⁺ 21]	Chimera graph	2032	2000Q	QPU chip sub-graph
2022 [TAM ⁺ 22]	Pegasus graph	5387	Adv4.1	QPU chip-subgraph
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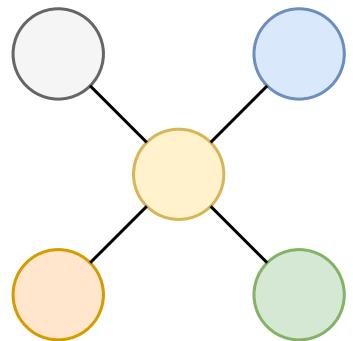
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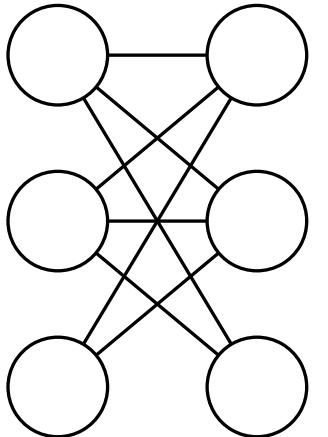
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2- Generation of instances

- Minor-embedding (Graph Minor Theory [RS95])



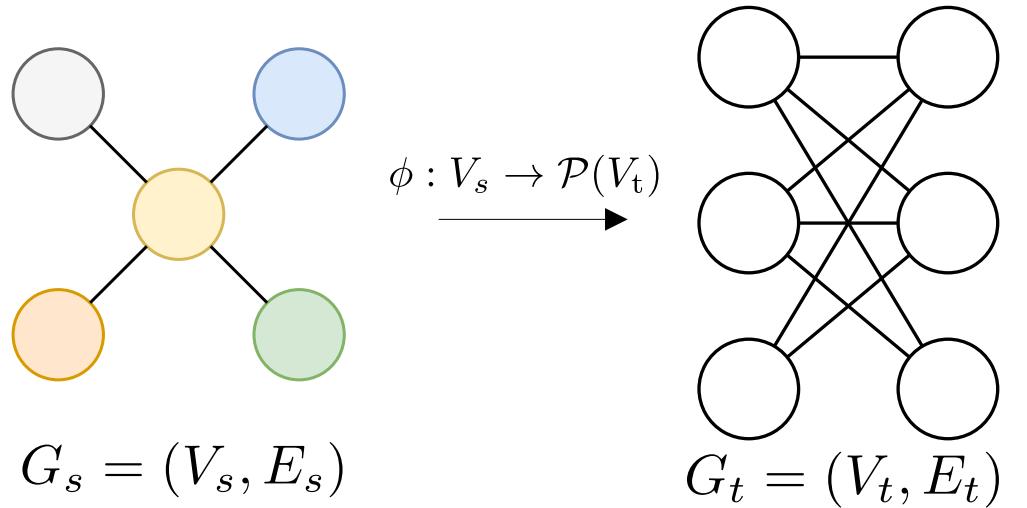
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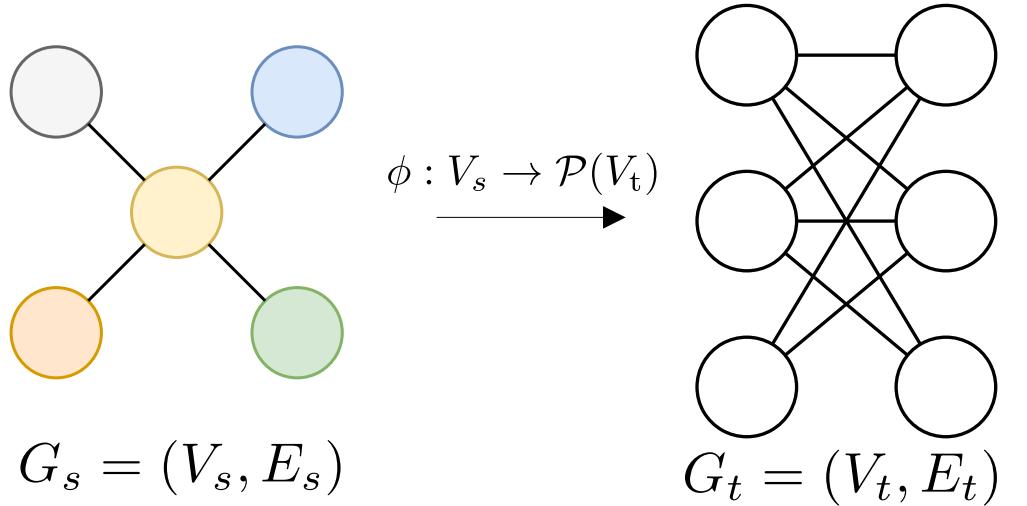
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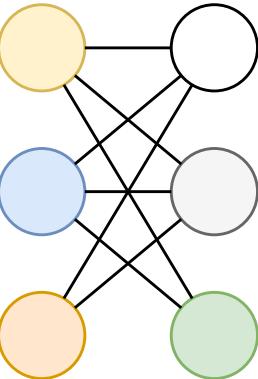


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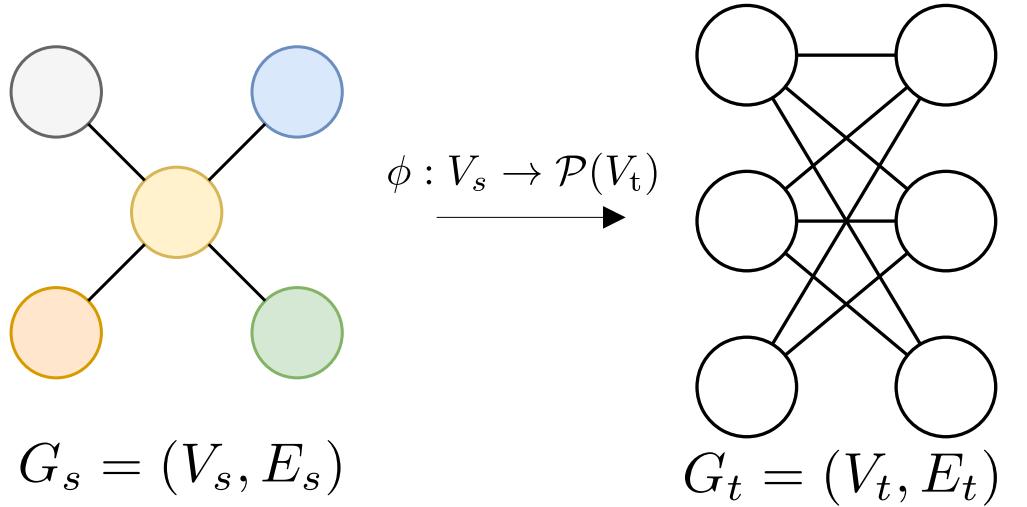
Rule 1:
All the edges in G_s must be represented in G_t



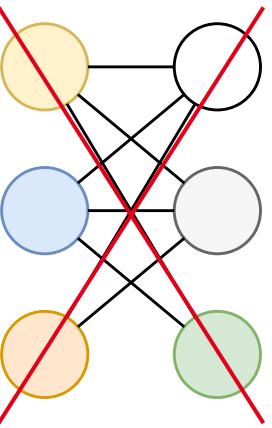


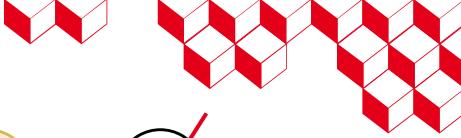
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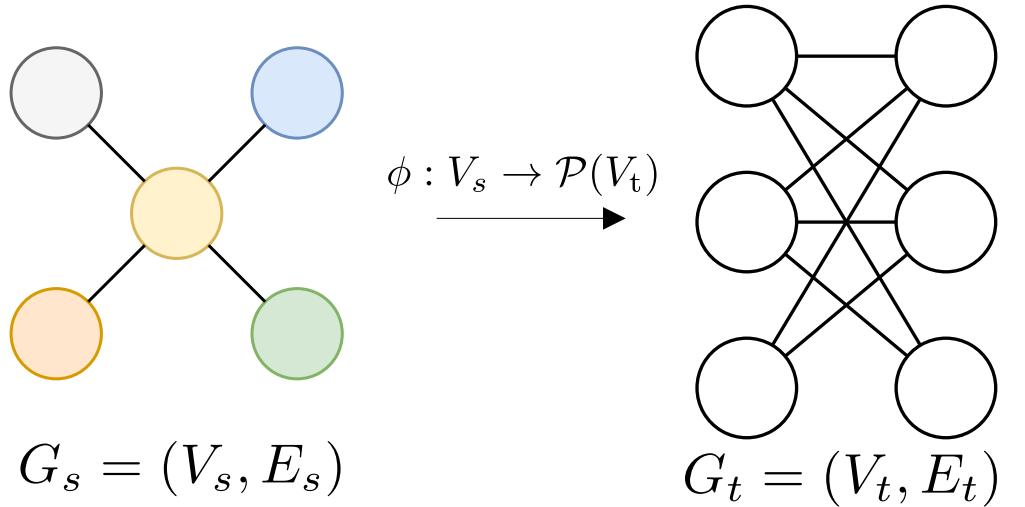
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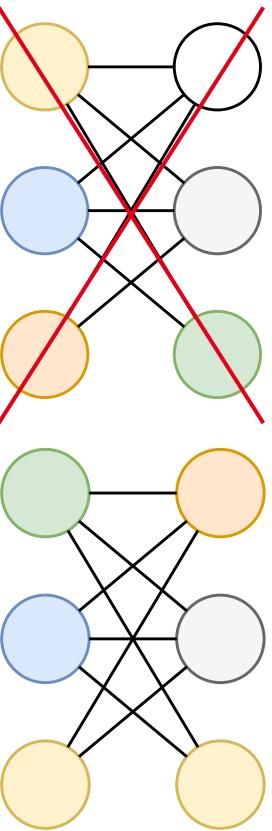


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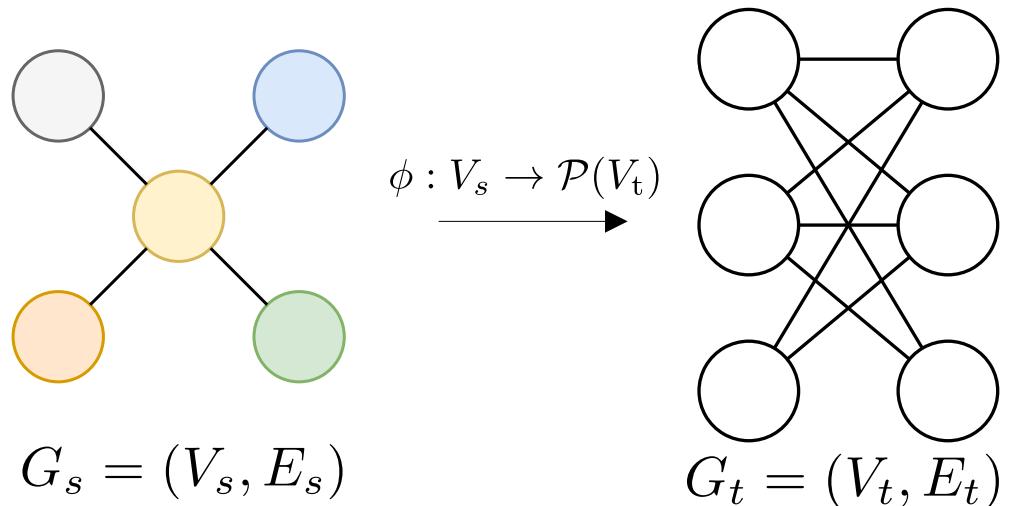
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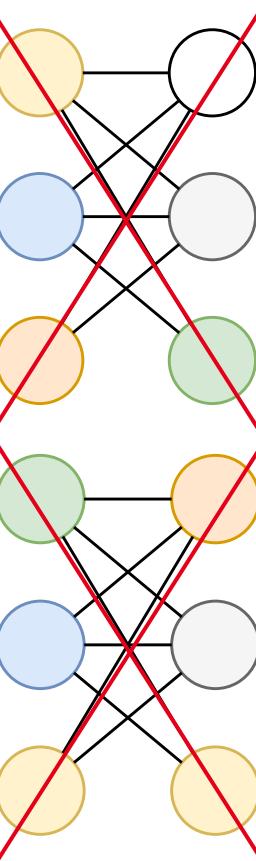


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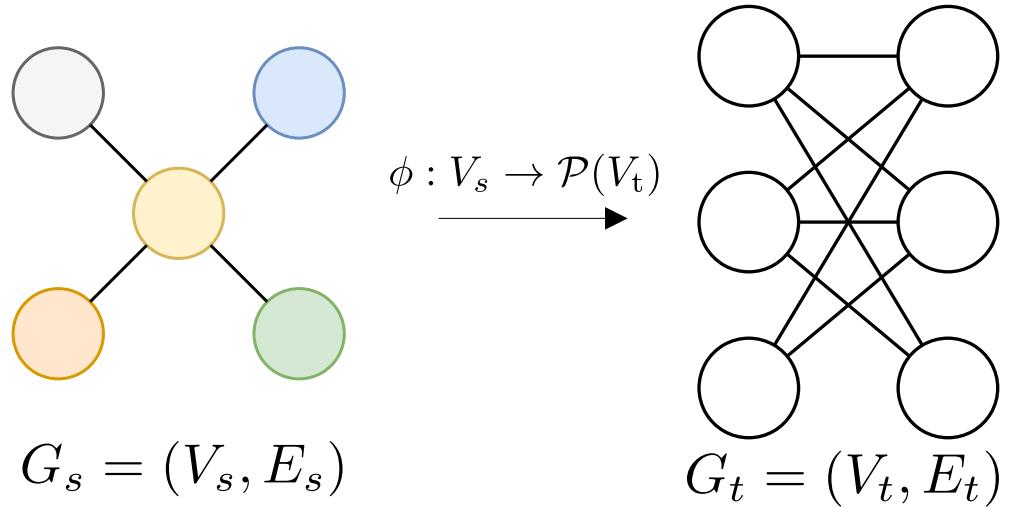
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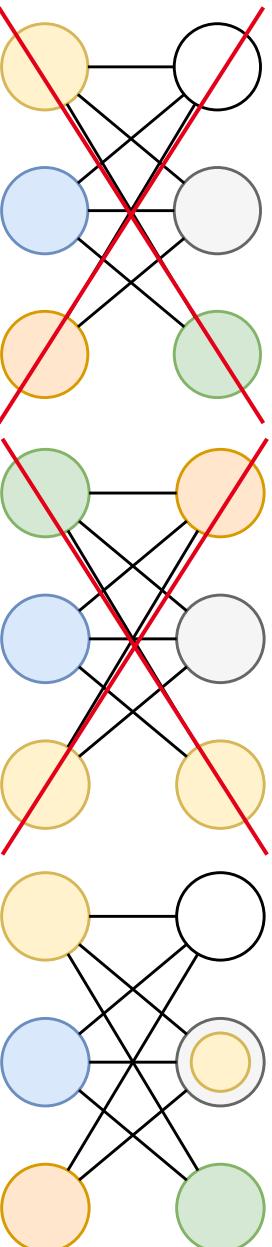
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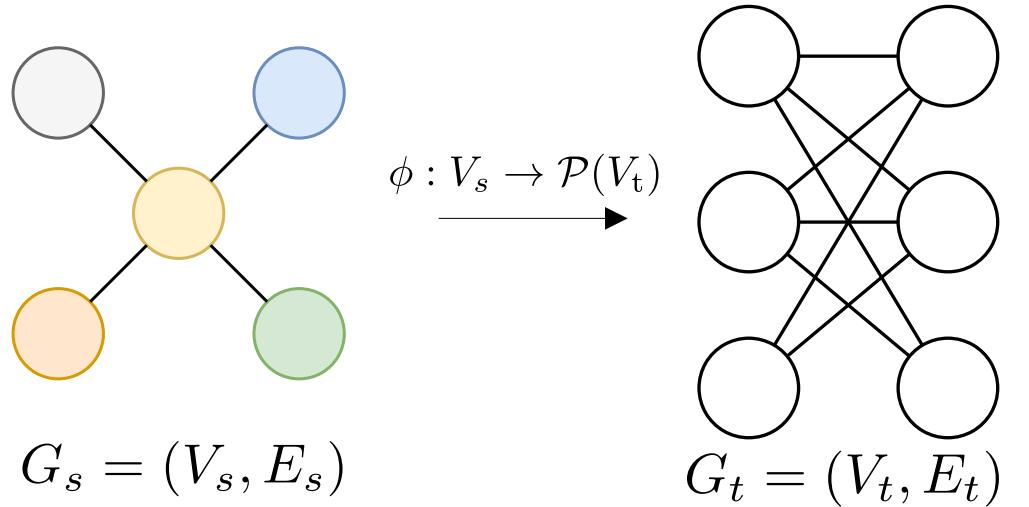
No overlap





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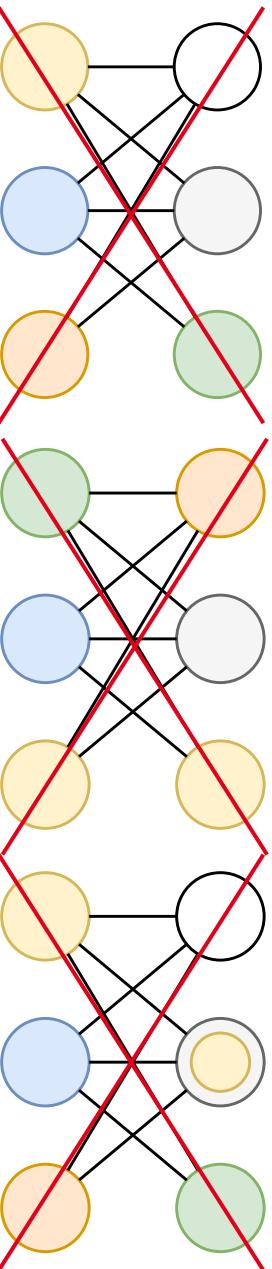
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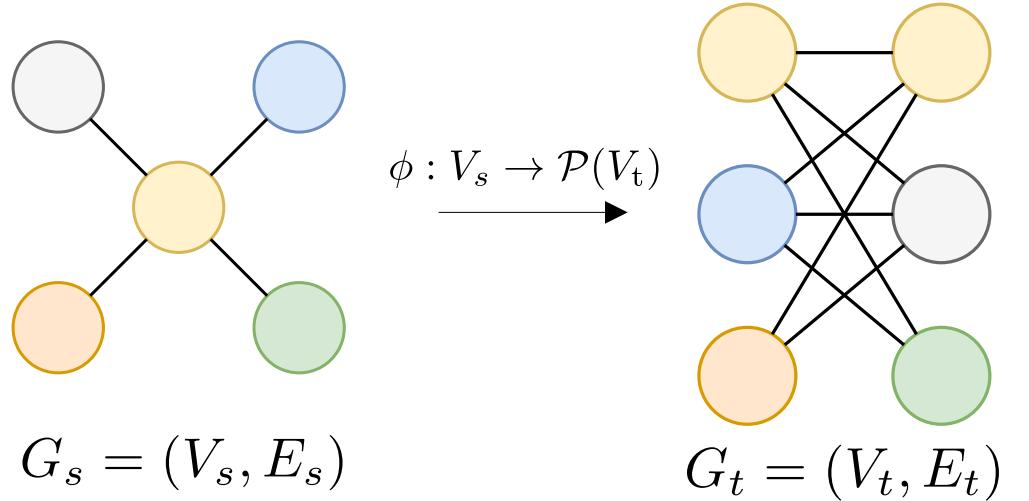
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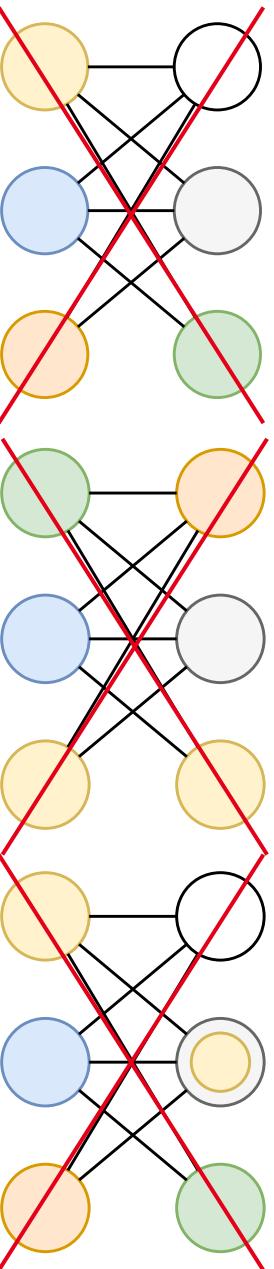
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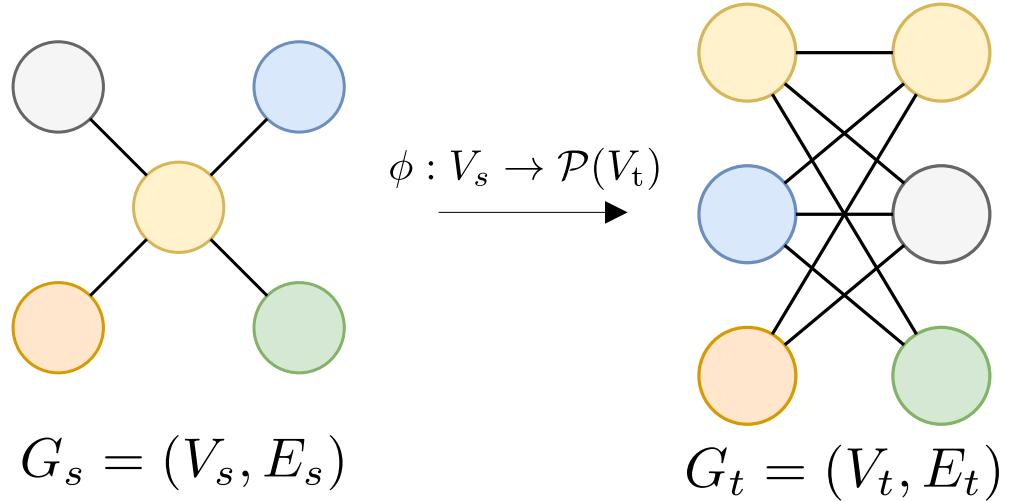
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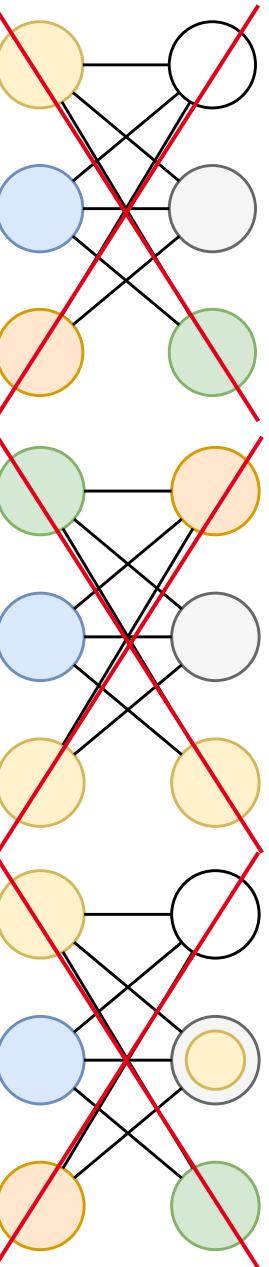


NP-Hard problem for arbitrary graphs

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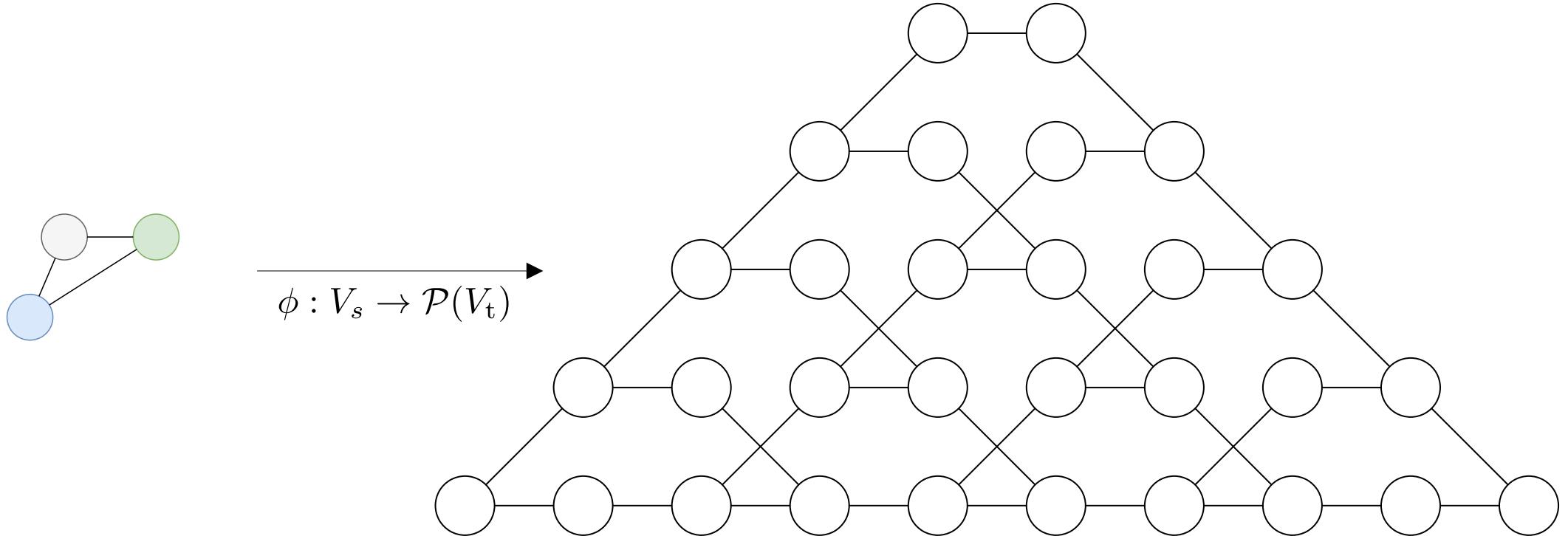
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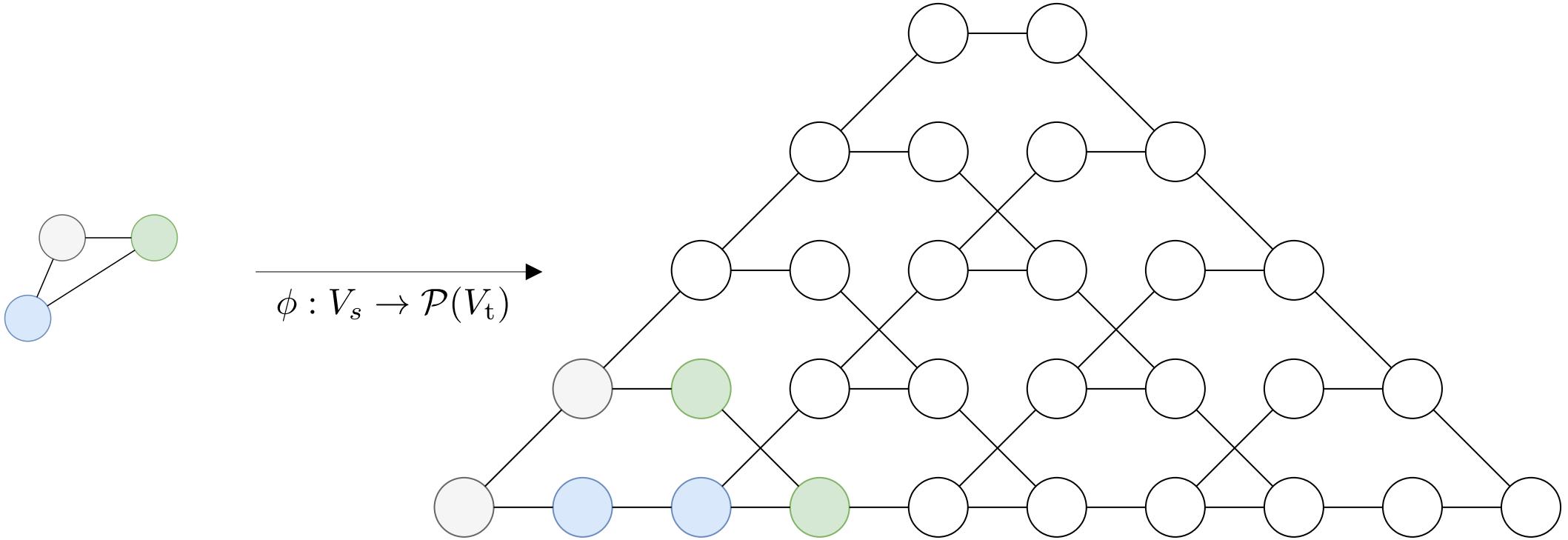
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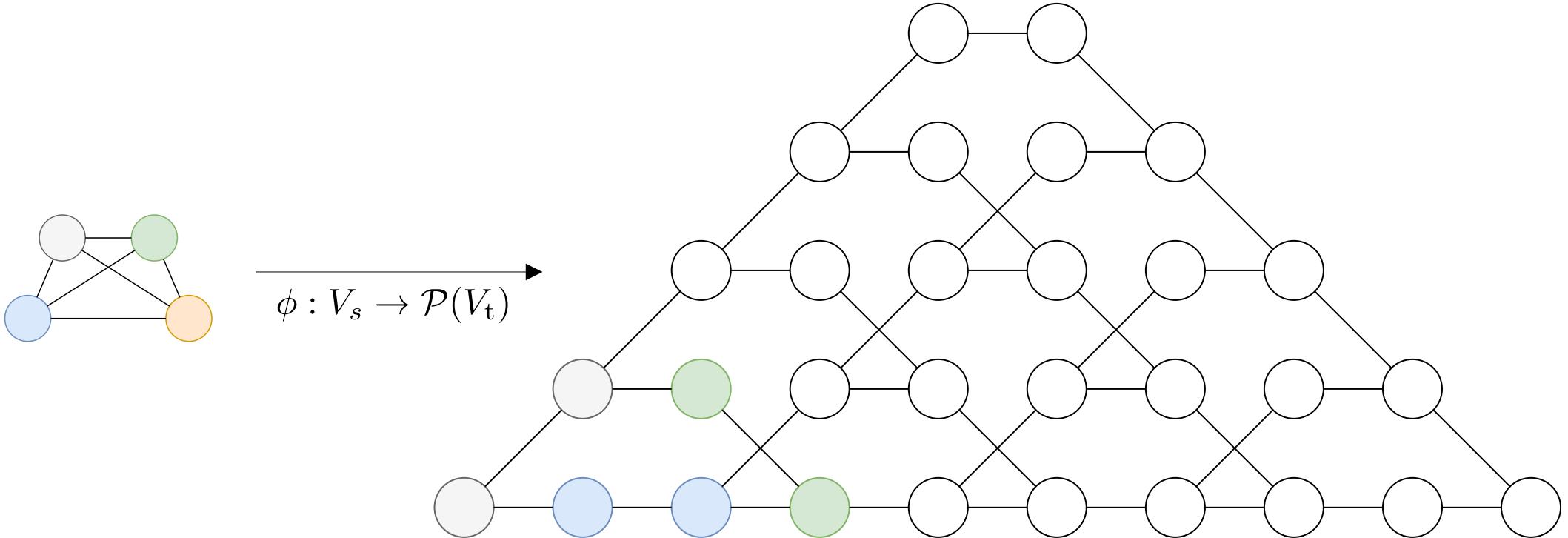
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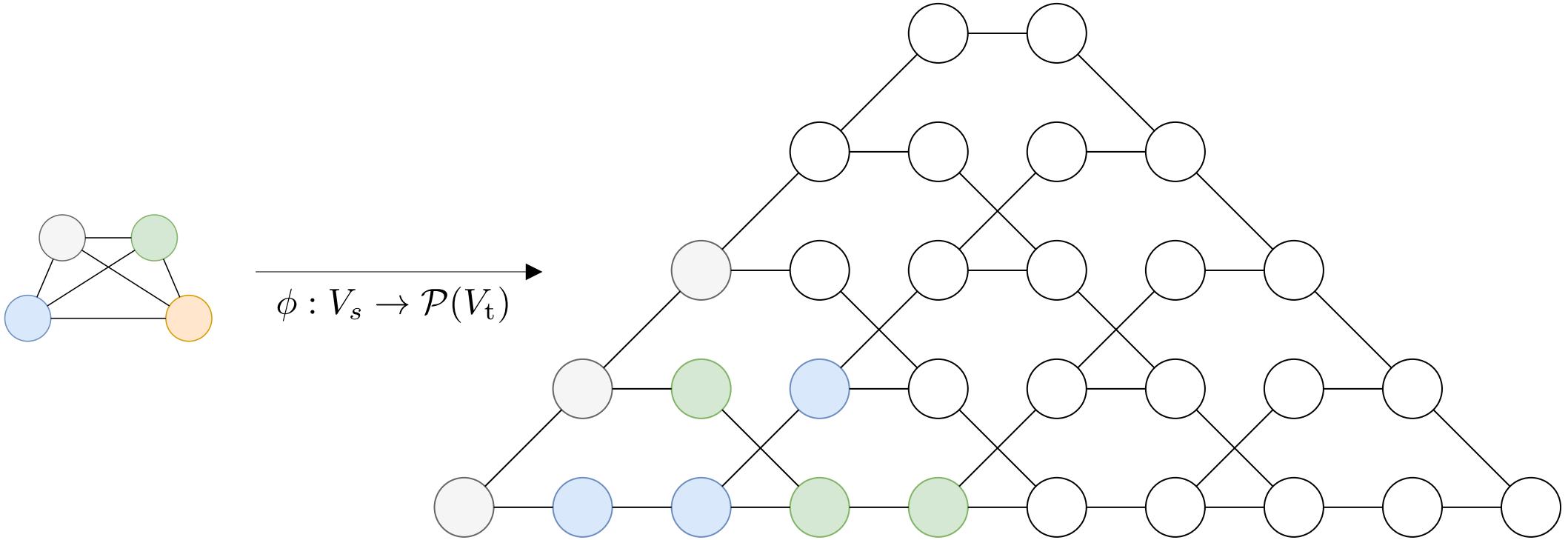
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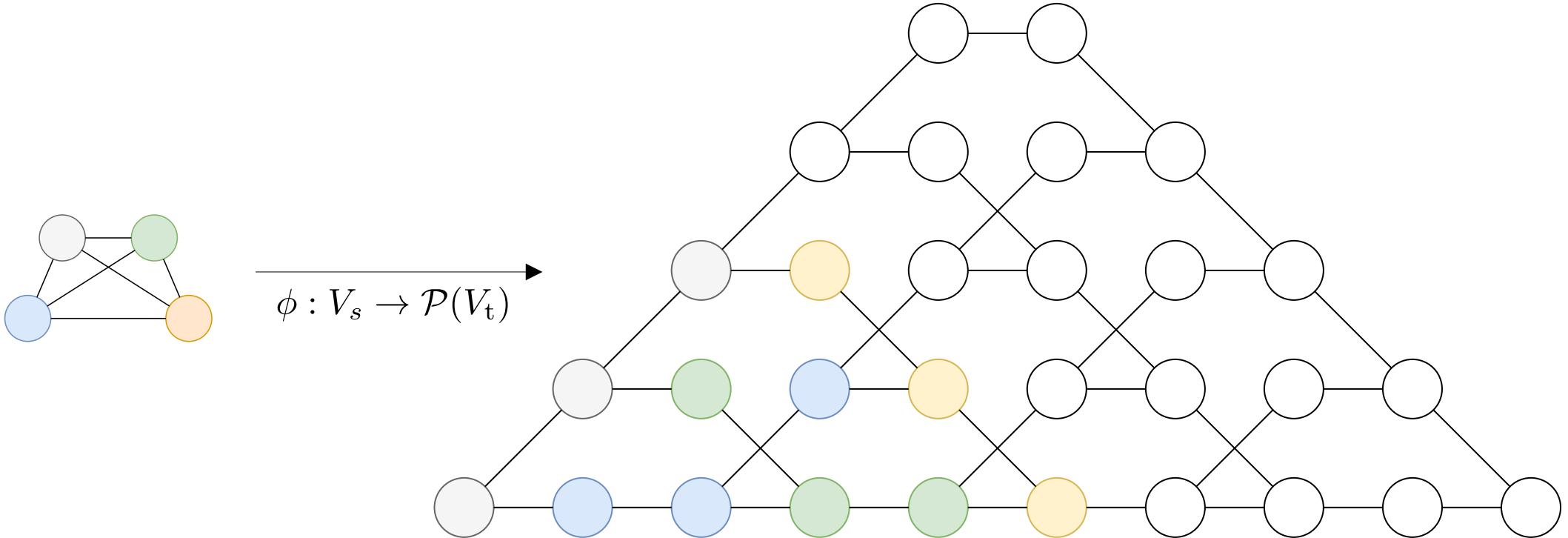
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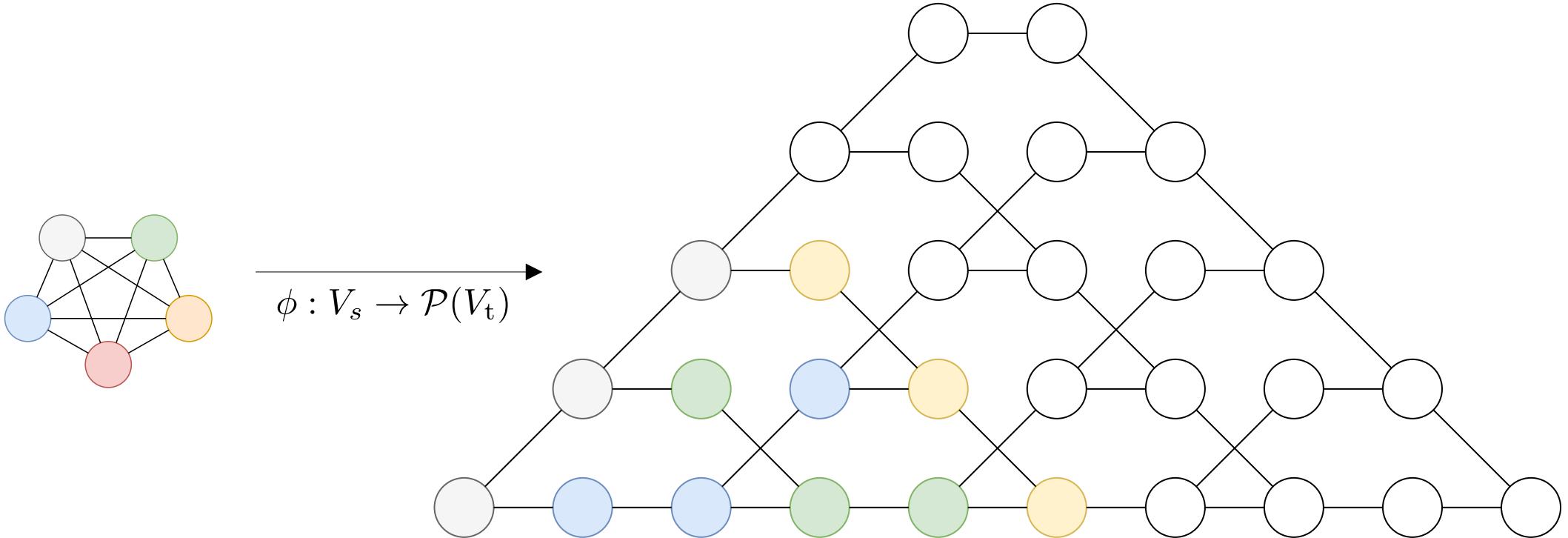
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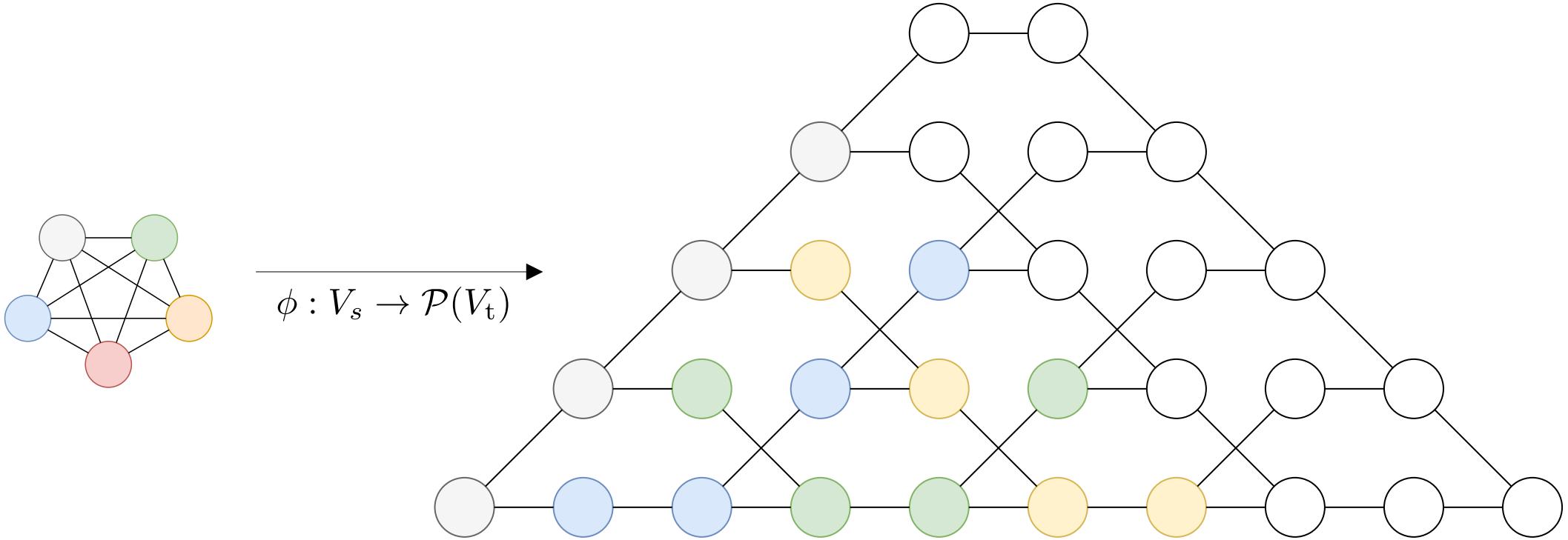
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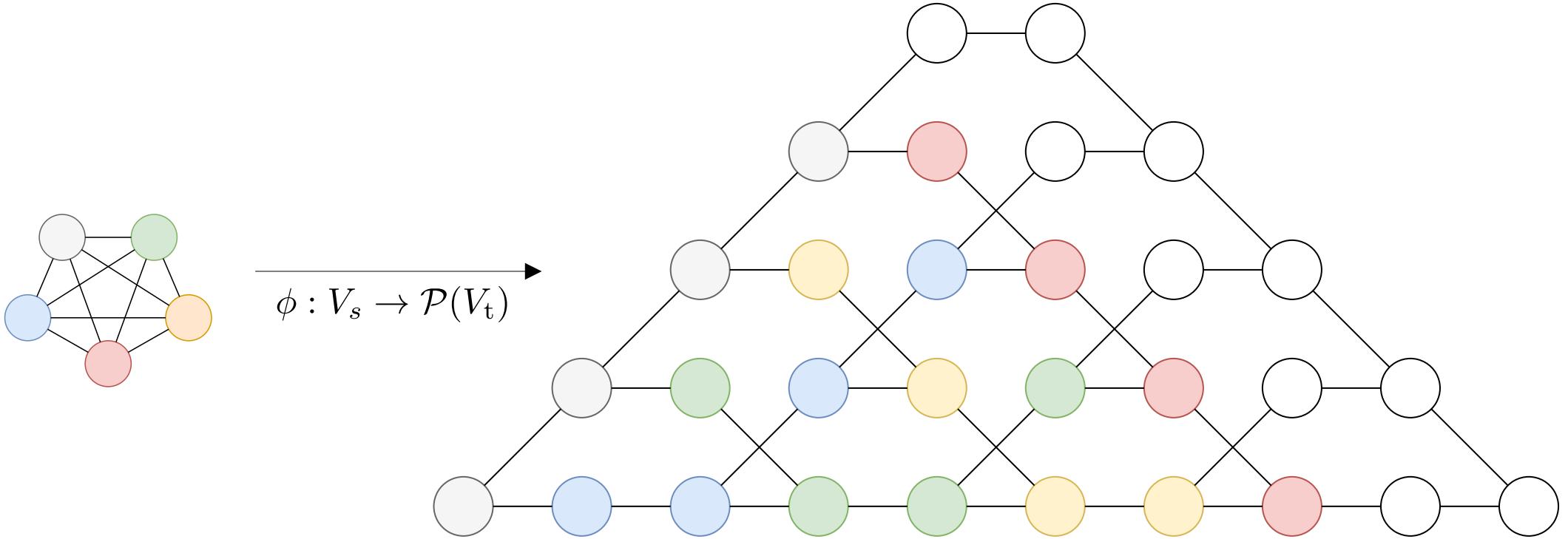
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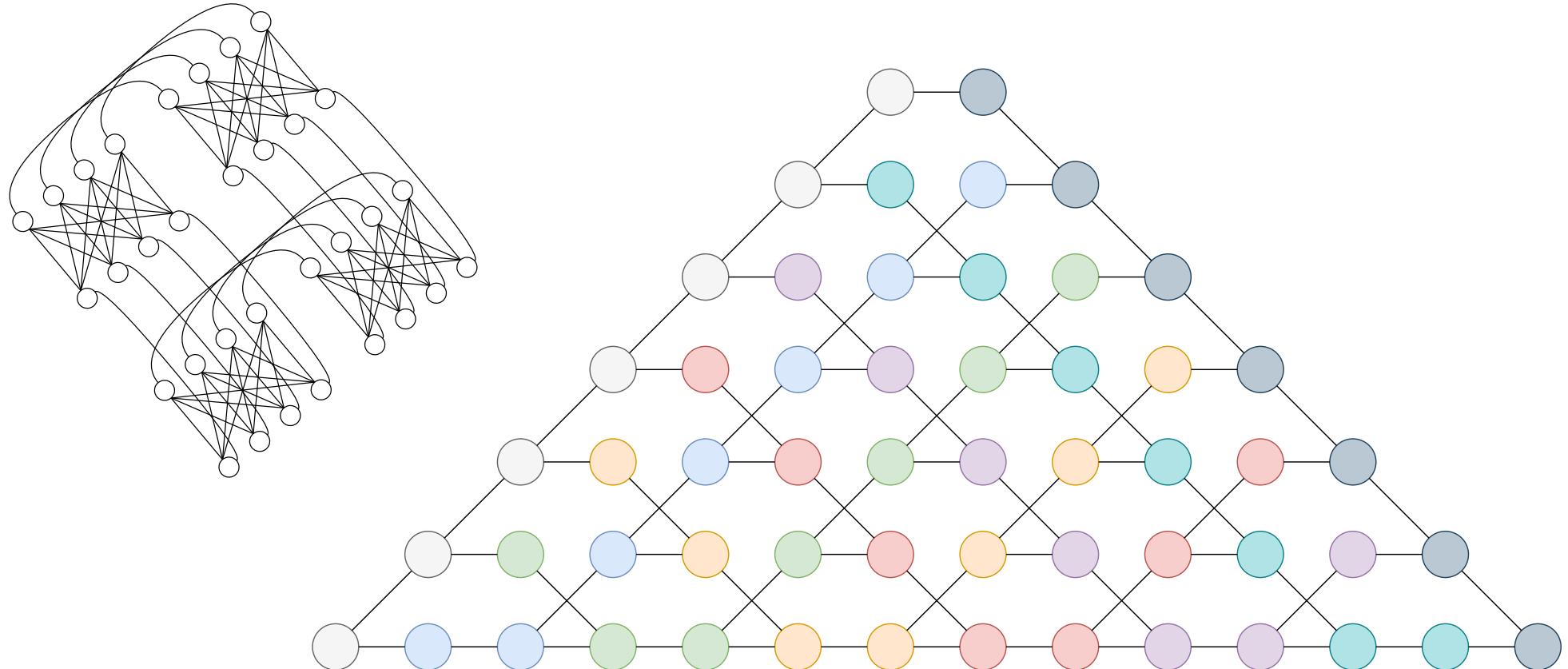
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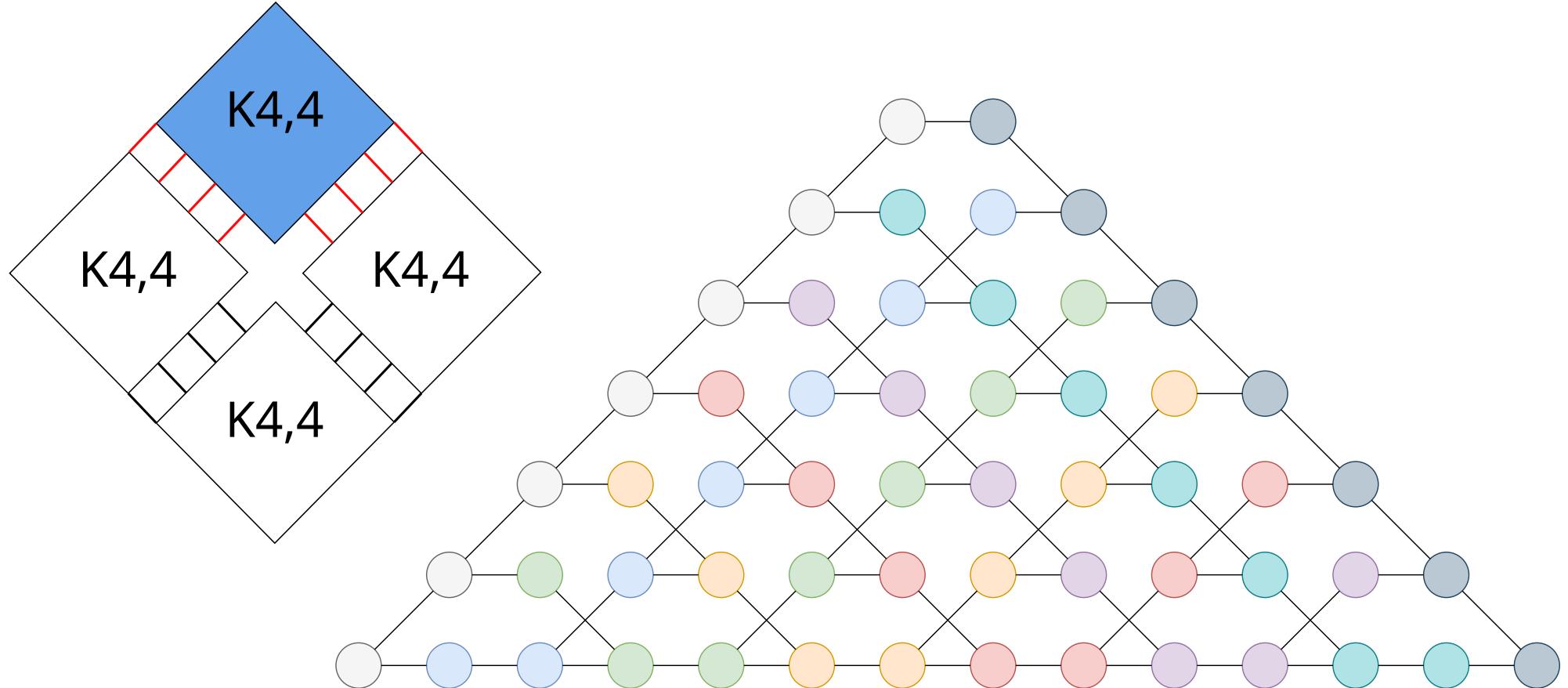
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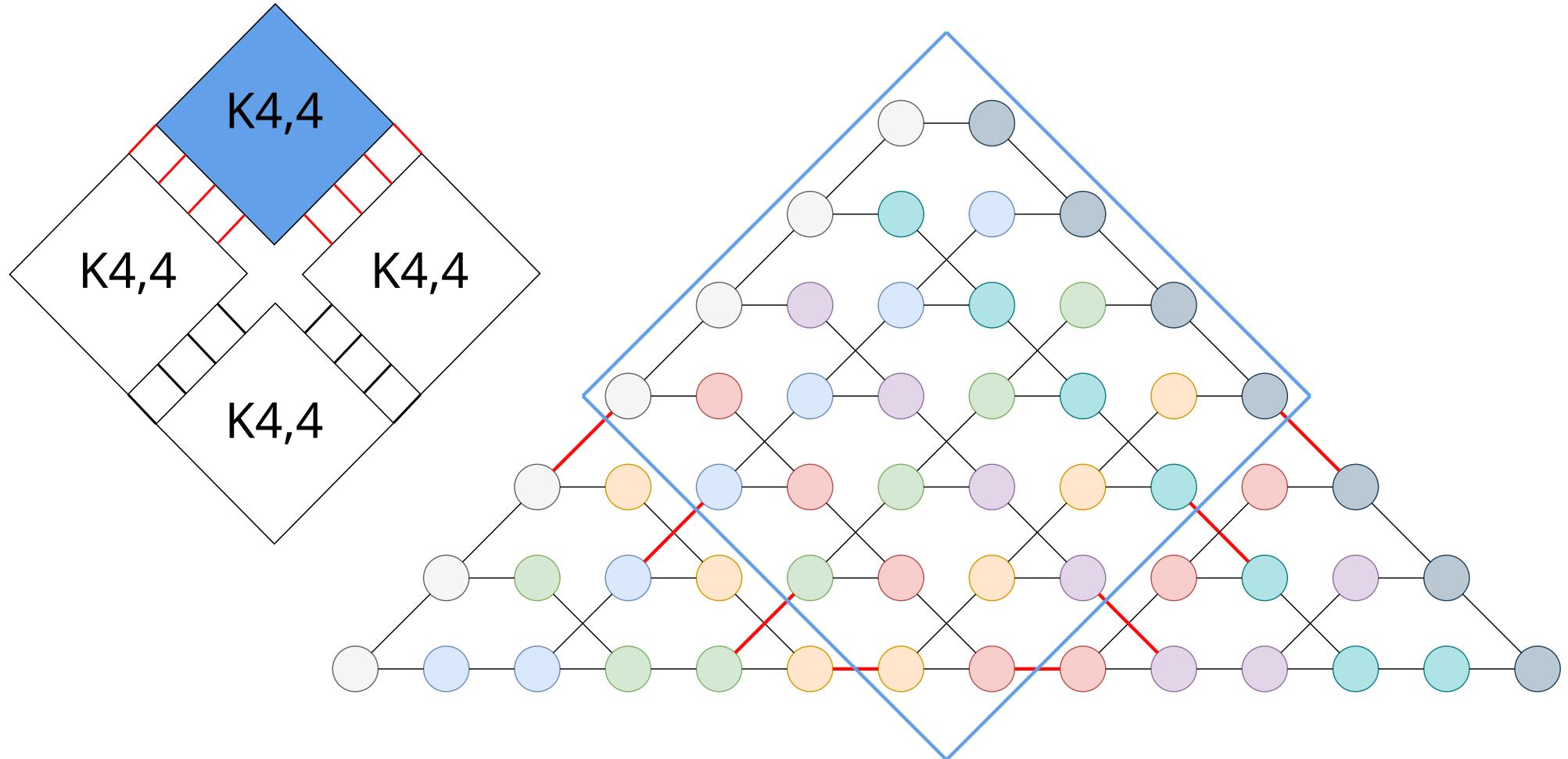
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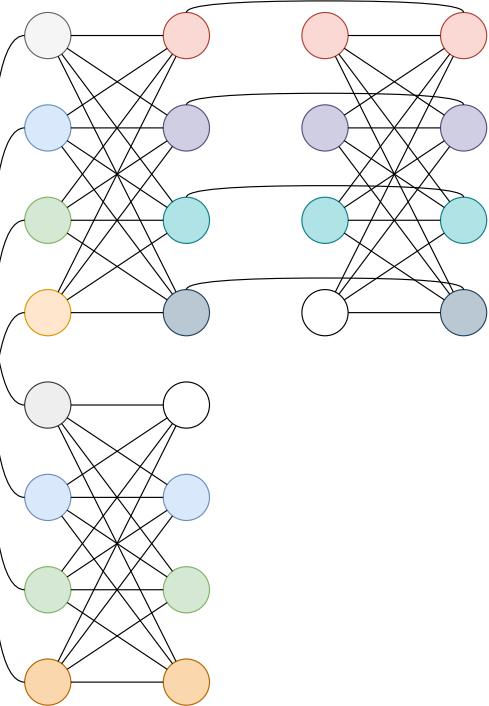
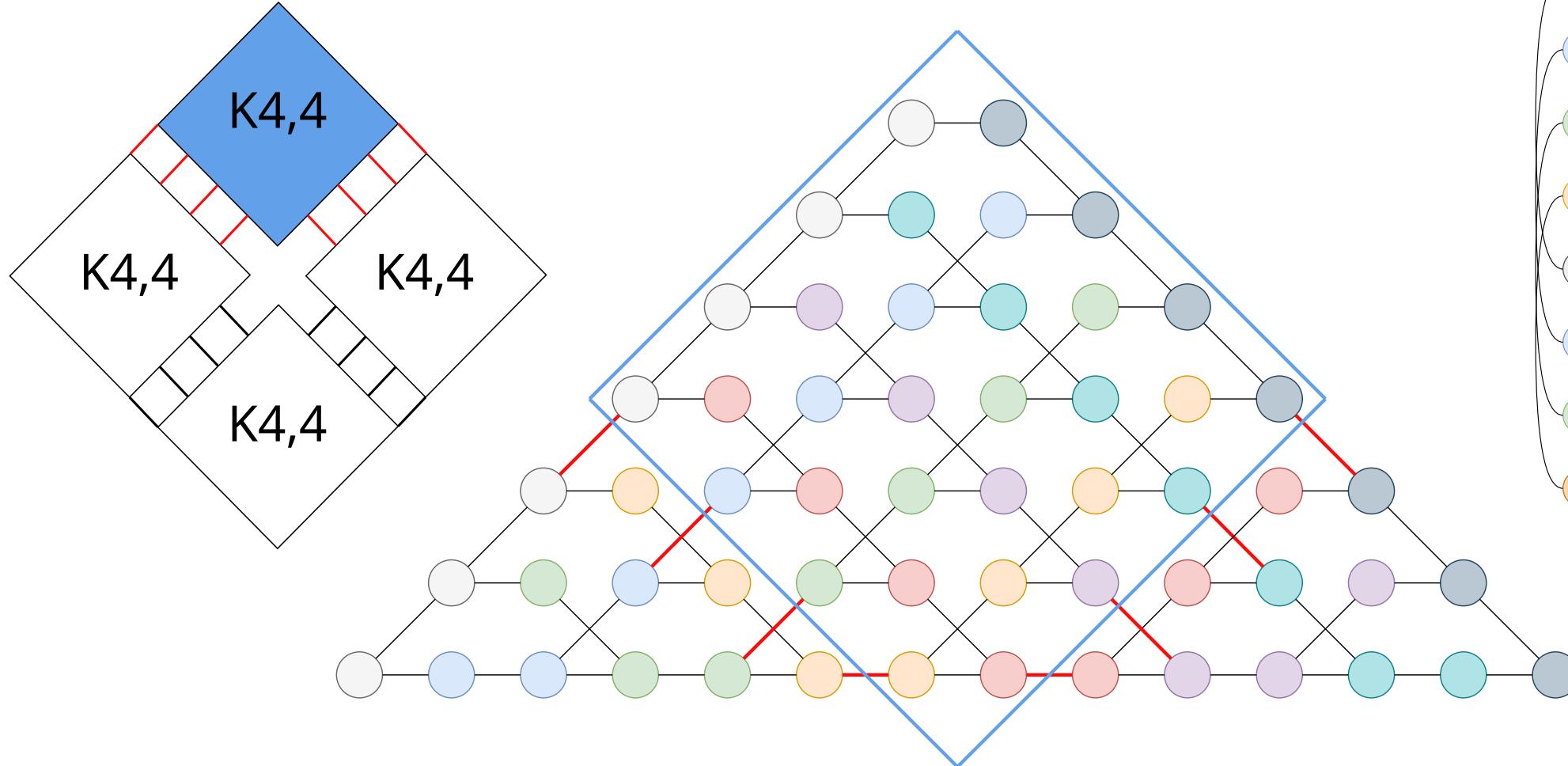
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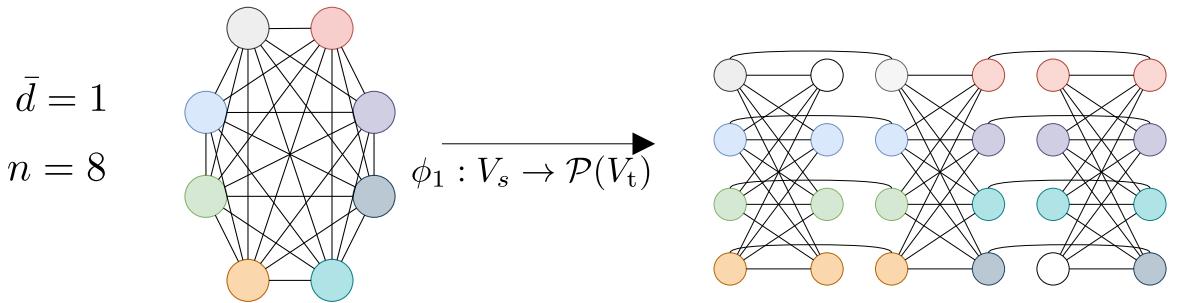
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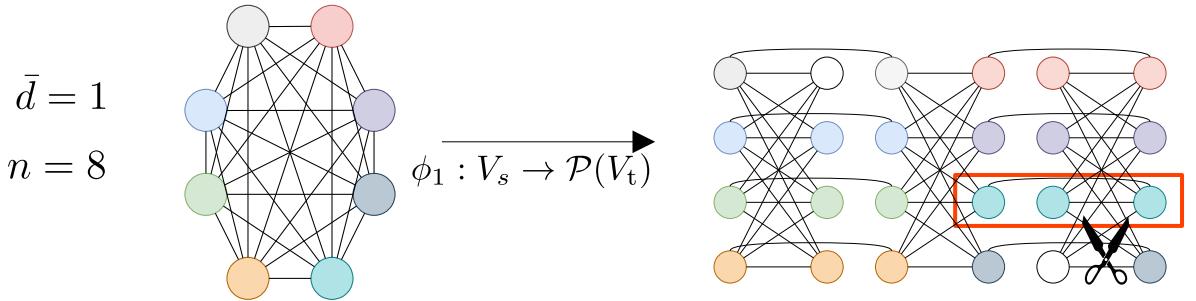
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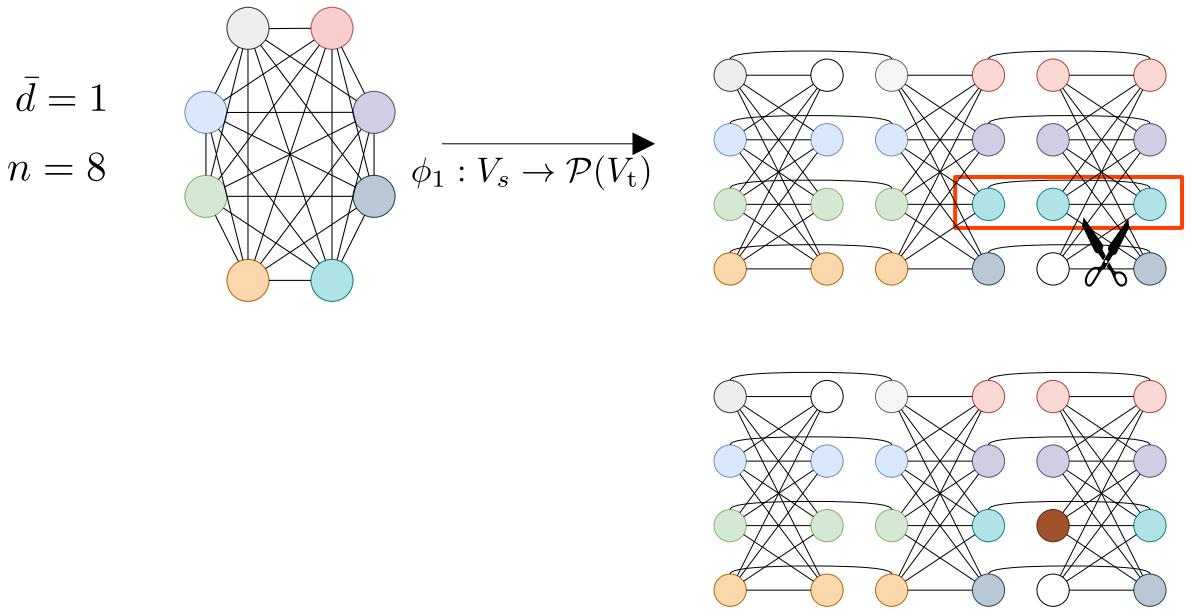
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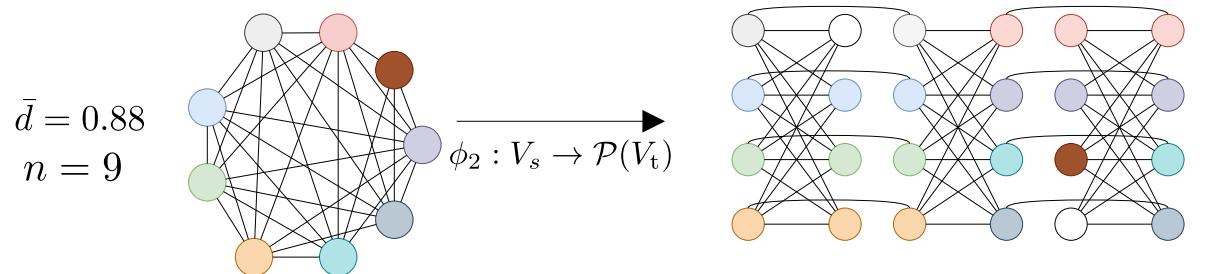
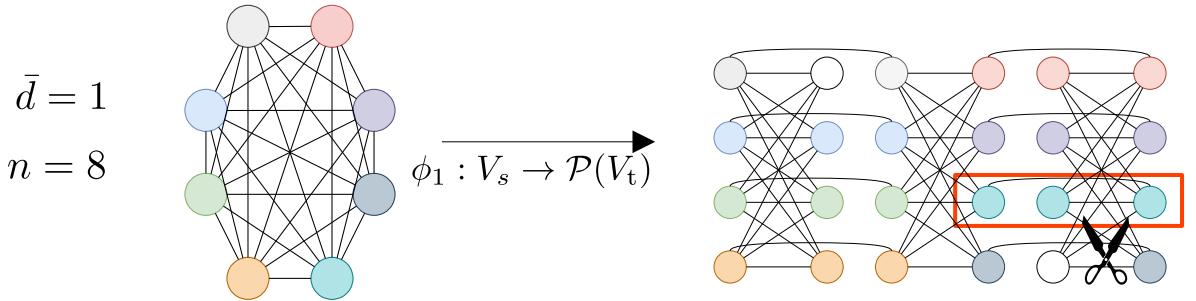
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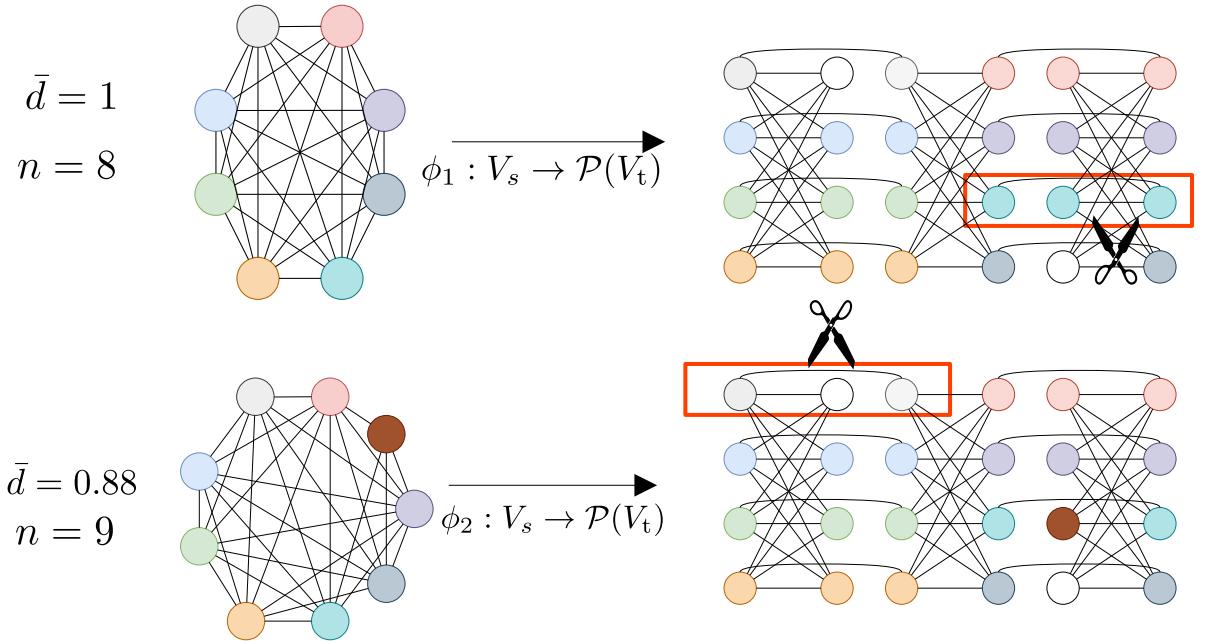
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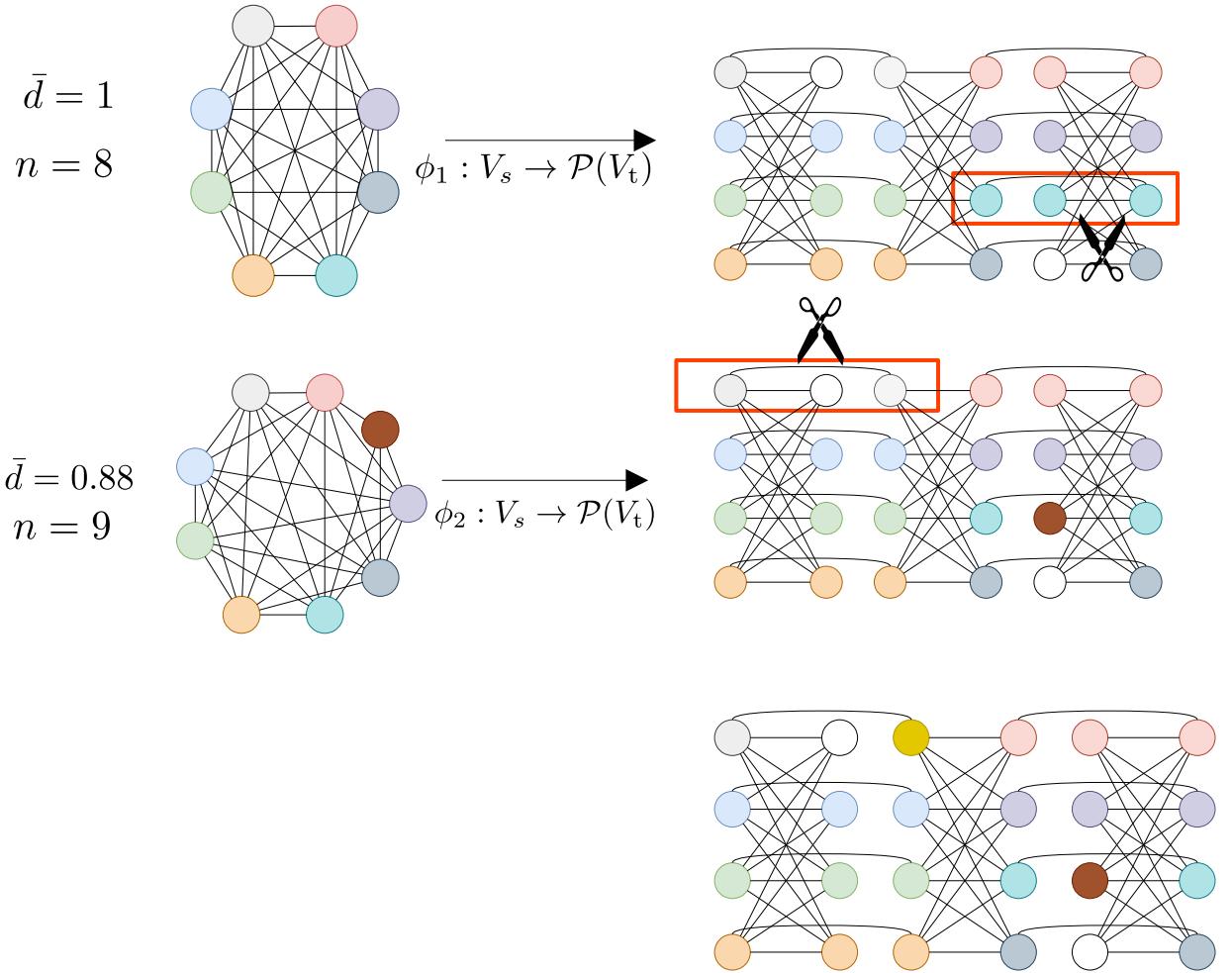
2- Generation of instances

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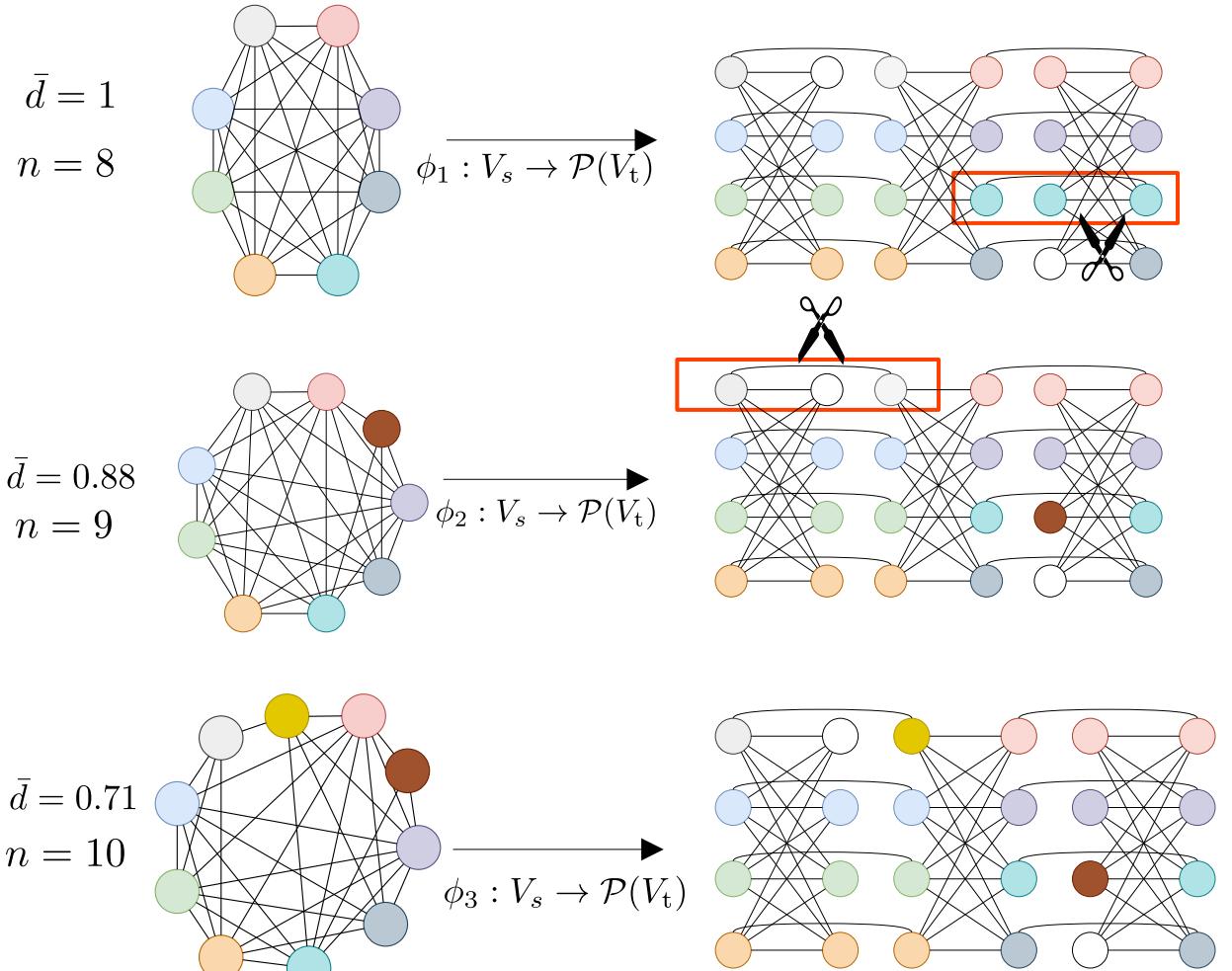
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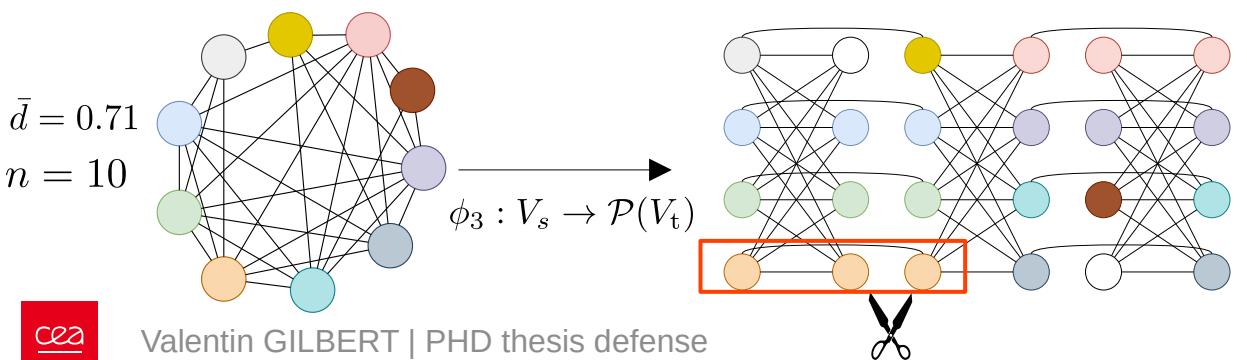
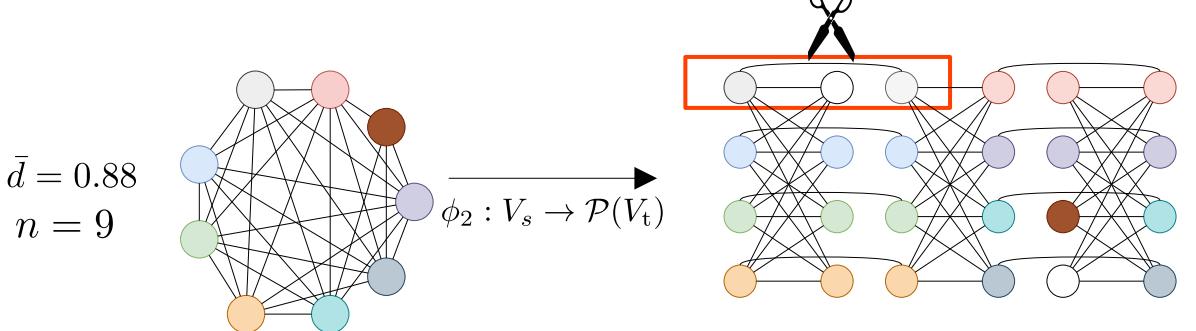
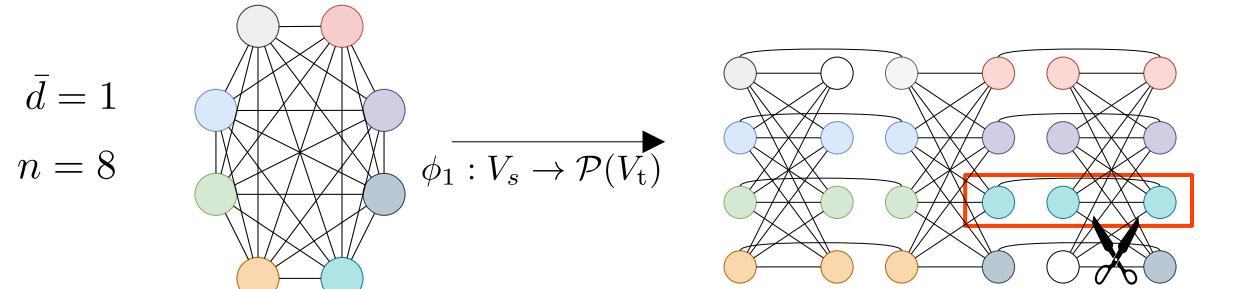
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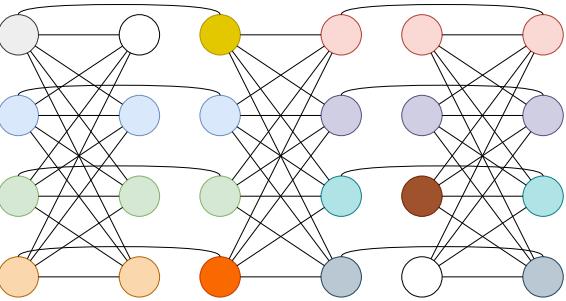
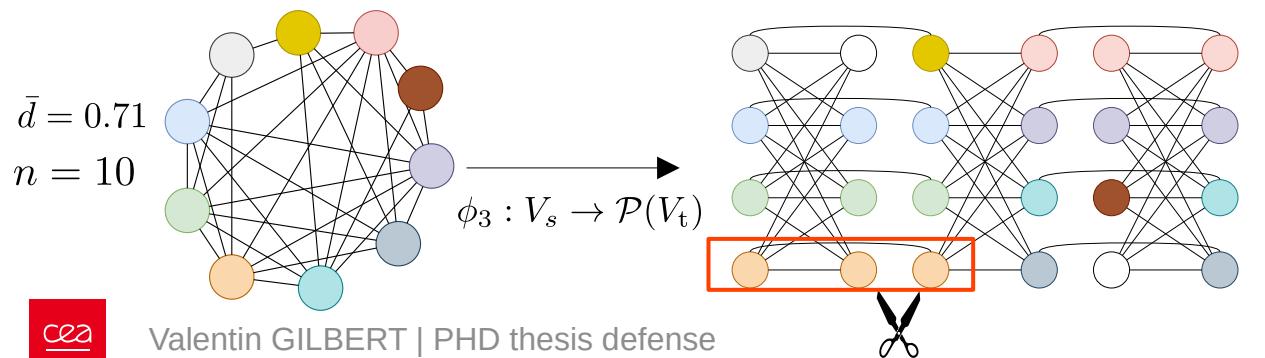
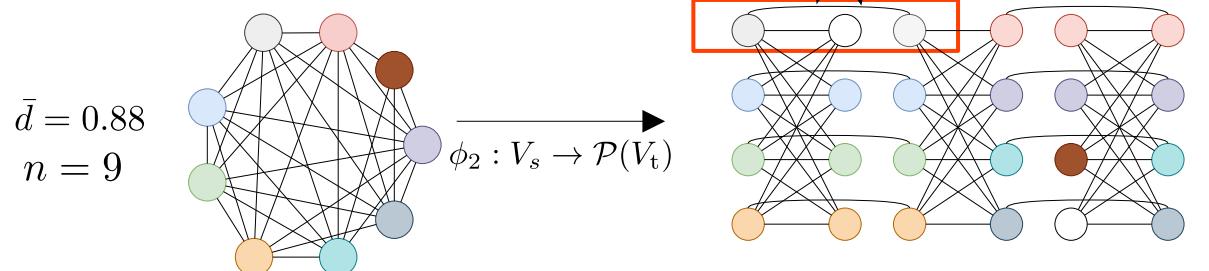
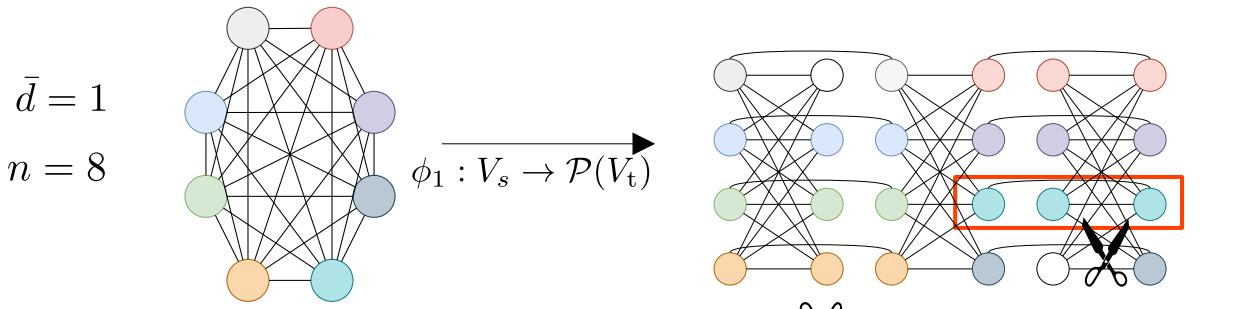
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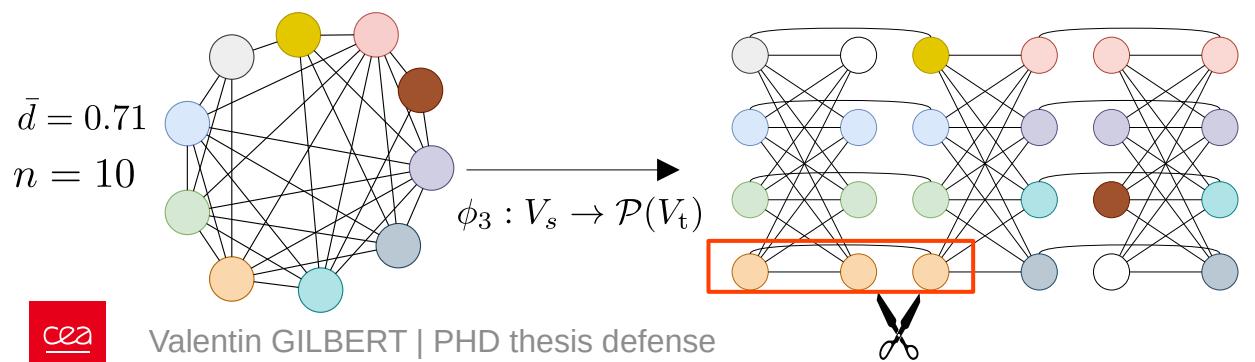
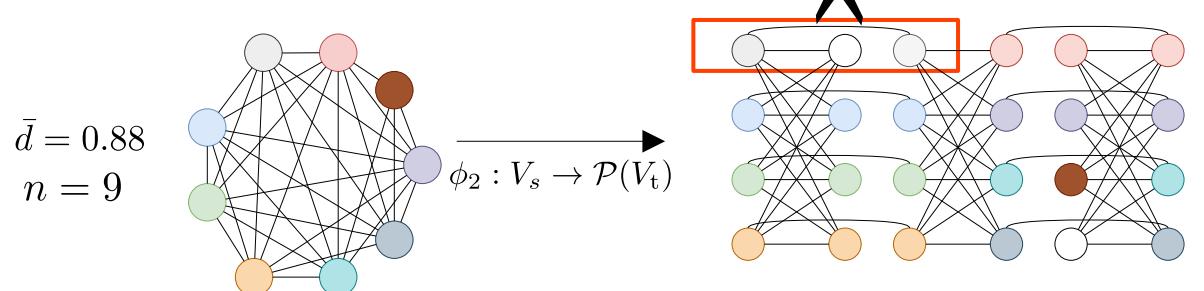
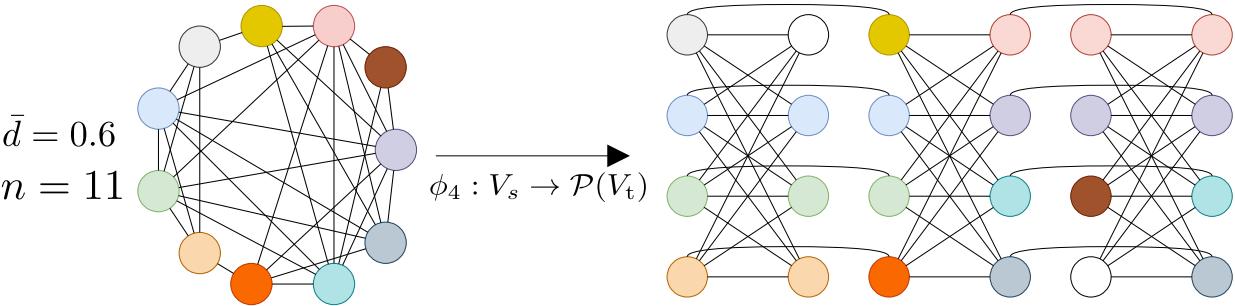
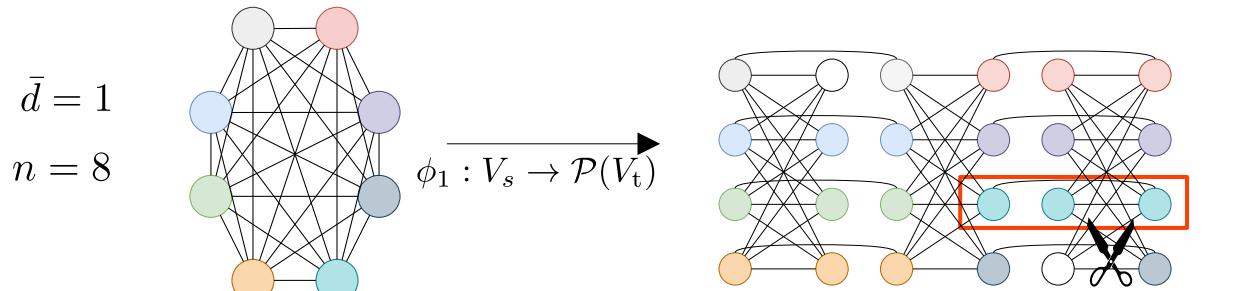
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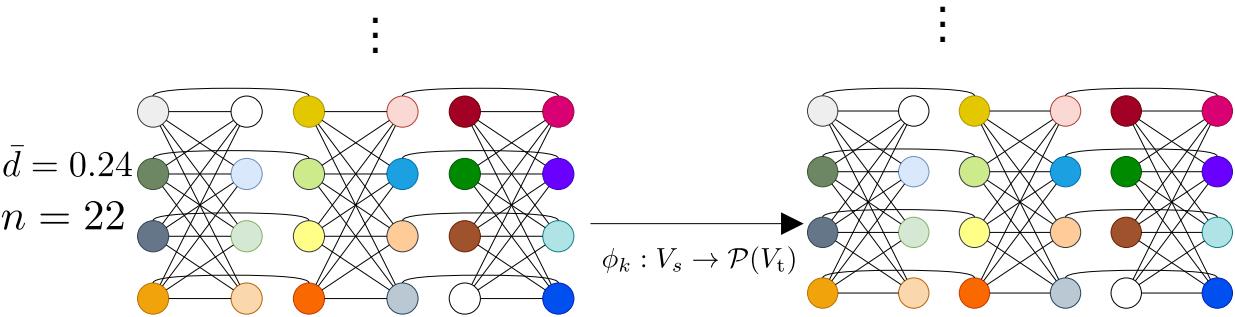
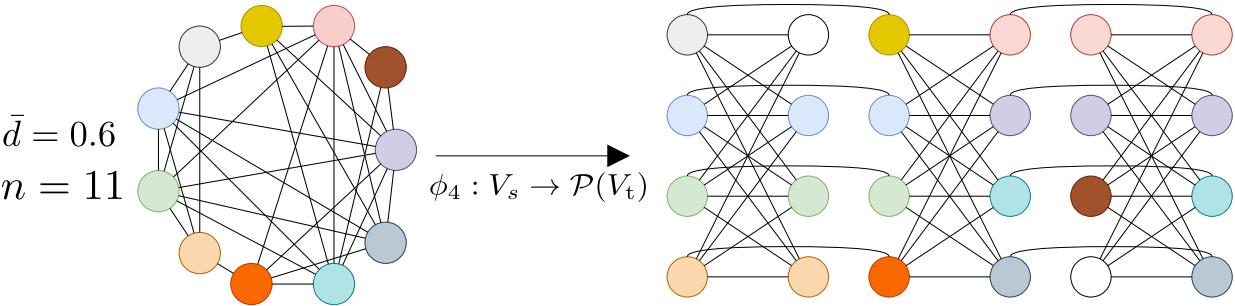
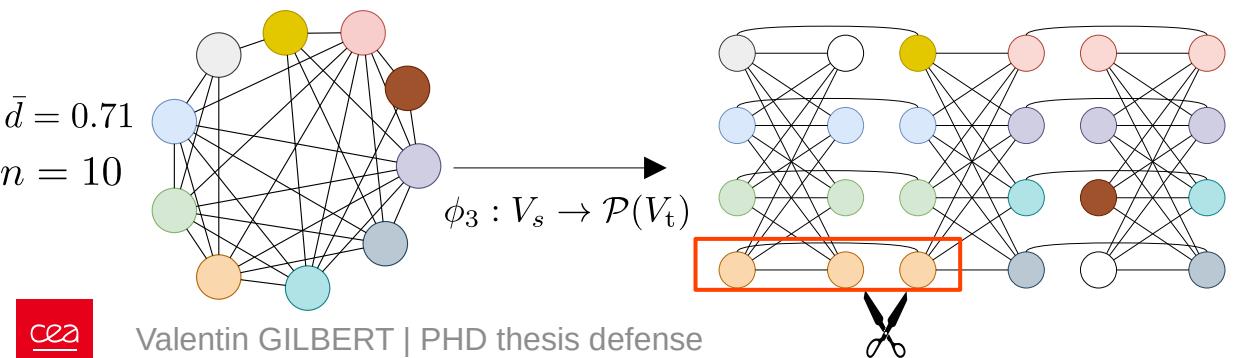
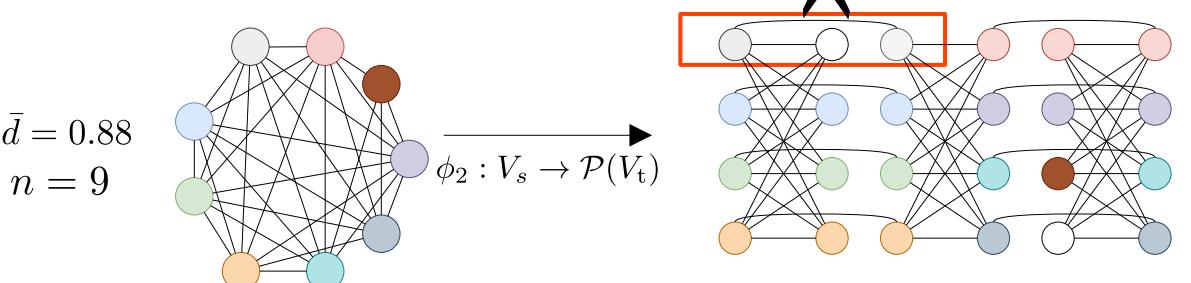
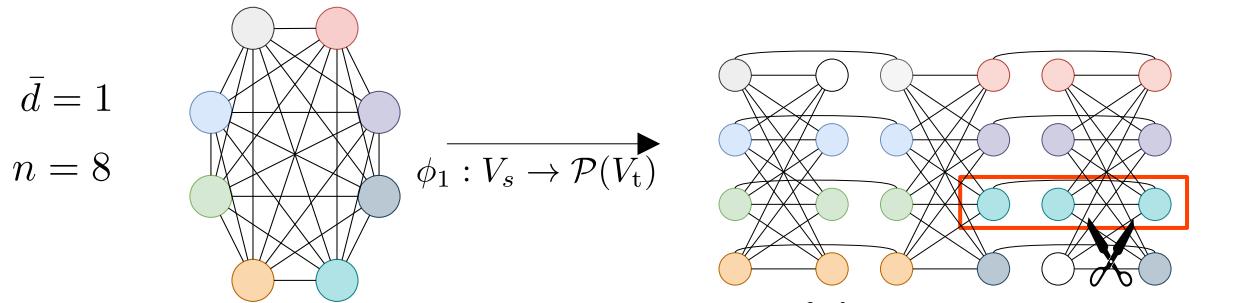
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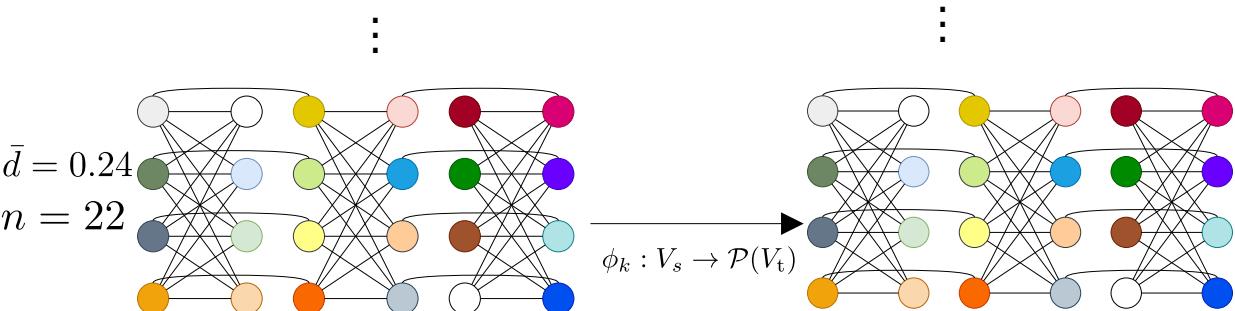
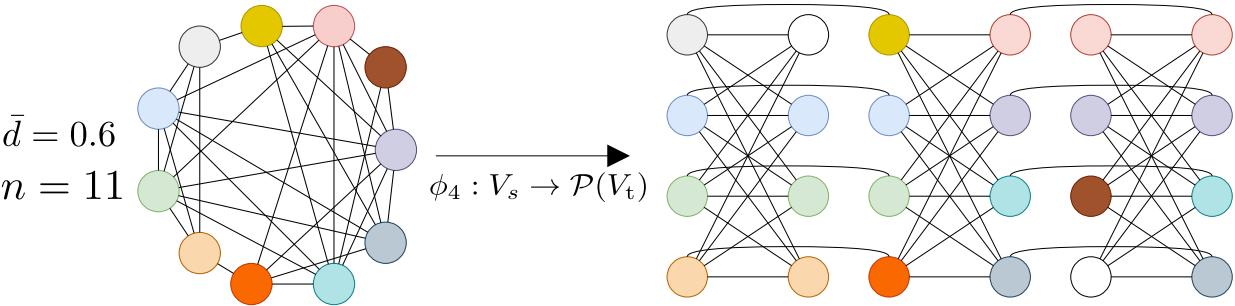
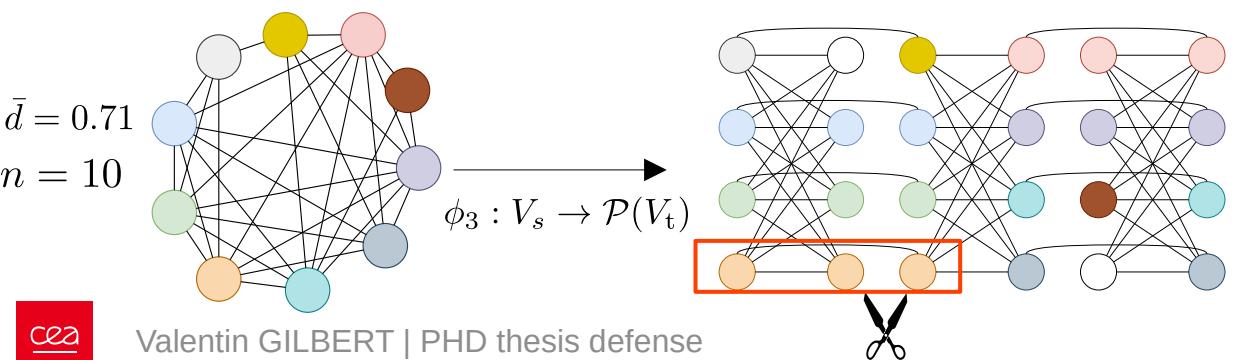
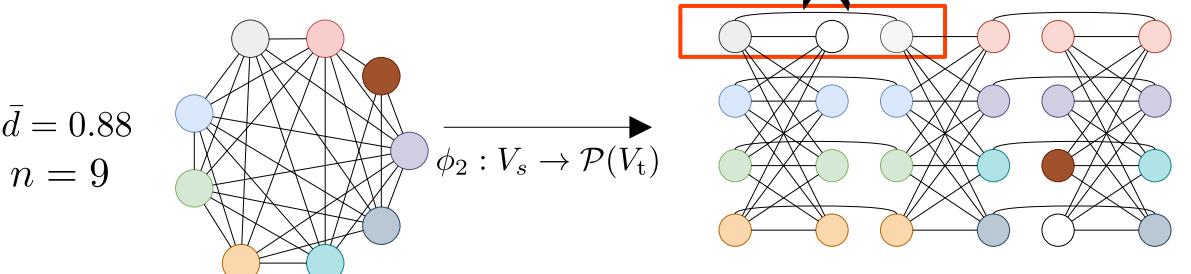
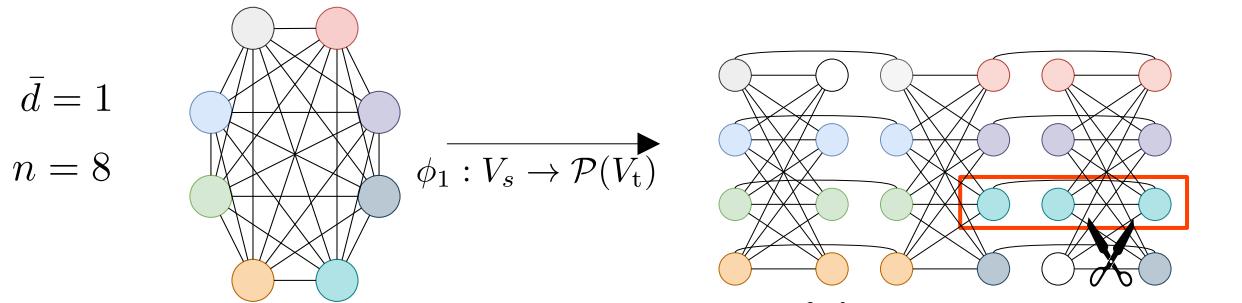
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2- Generation of instances

- Embedding of complete graphs (TRIAD Pattern)



Largest clique size: 174 (Advantage6.4)
=> Use approximately 3000 qubits



2- Performance assessment

- Time To Solution metric (Gold standard) [RWJ⁺14]

Number of runs:

$$R = \left\lceil \frac{\log(1 - p)}{\log(1 - s)} \right\rceil$$

p : probability of getting the ground state in R runs

s : Empirical success probability

$$TTS = t_a \times R$$

t_a : Time to perform a single quantum run

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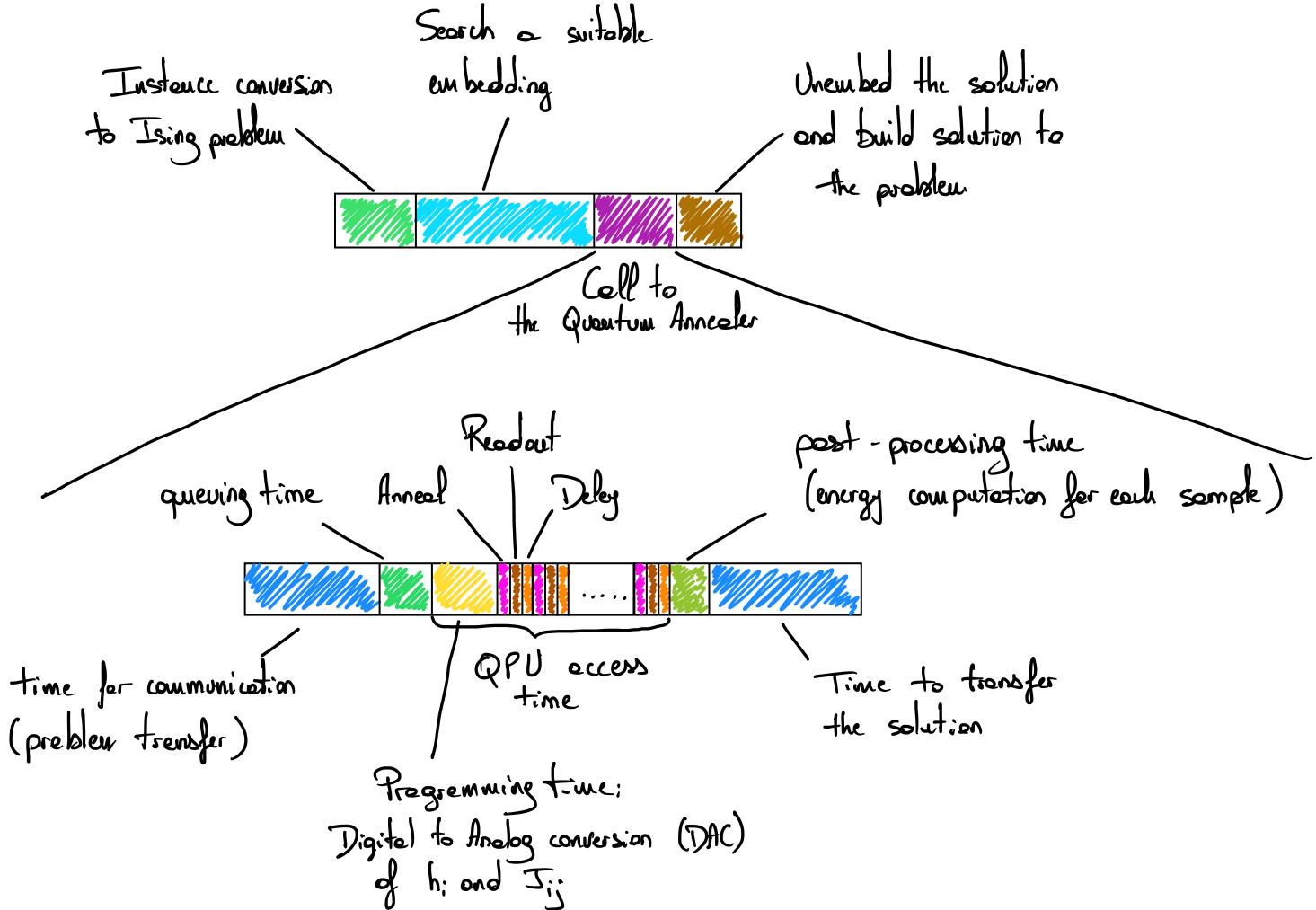
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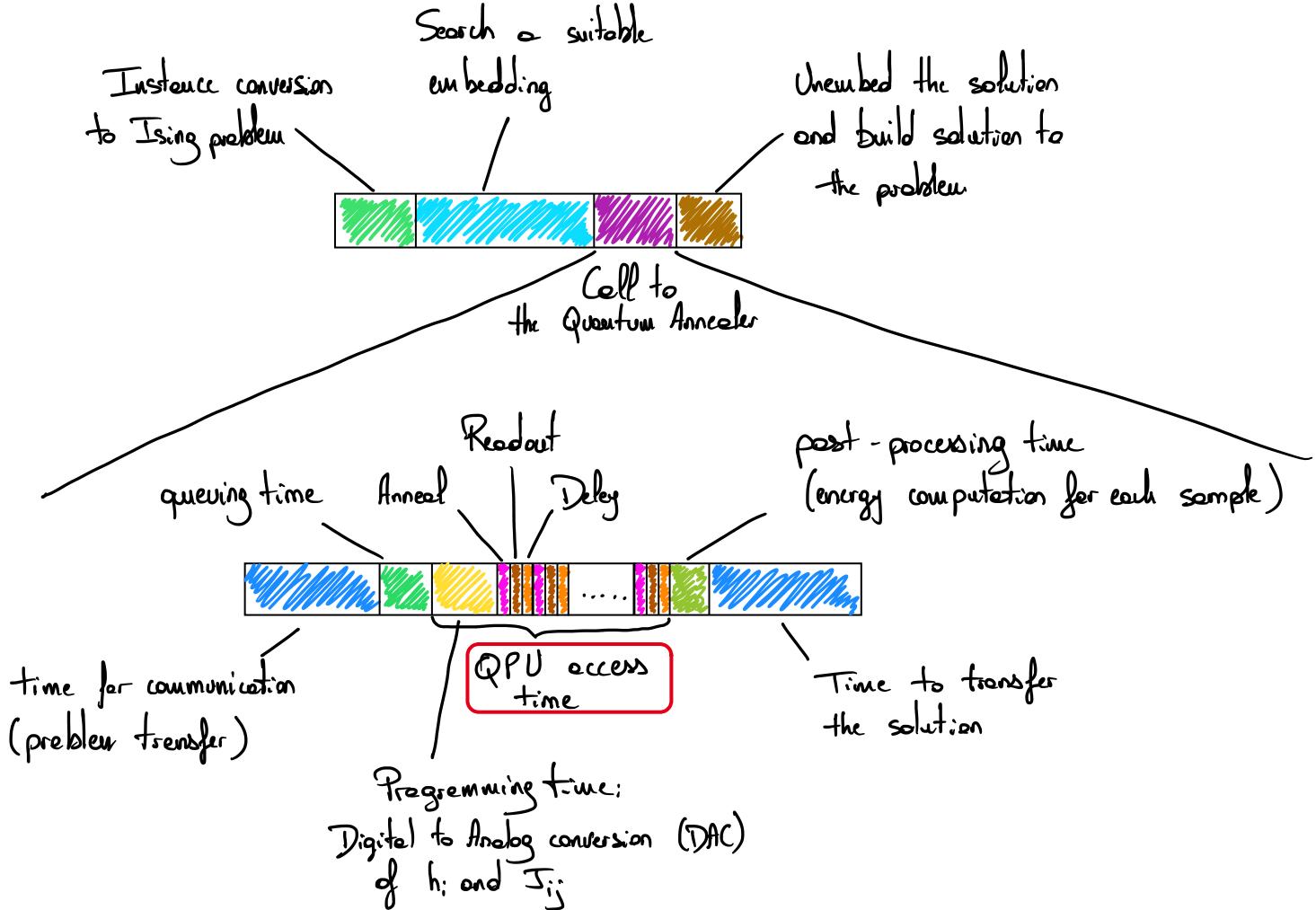
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2- Results – Max-cut problem

Density	0.02	0.03	0.04	0.05	0.06	0.07	0.08	0.09	0.1	0.2	0.3	0.4	0.5	0.6	0.7	0.8	0.9
Avg V _s	1318	1062	912	810	737	680	635	597	565	395	321	277	248	226	209	195	184



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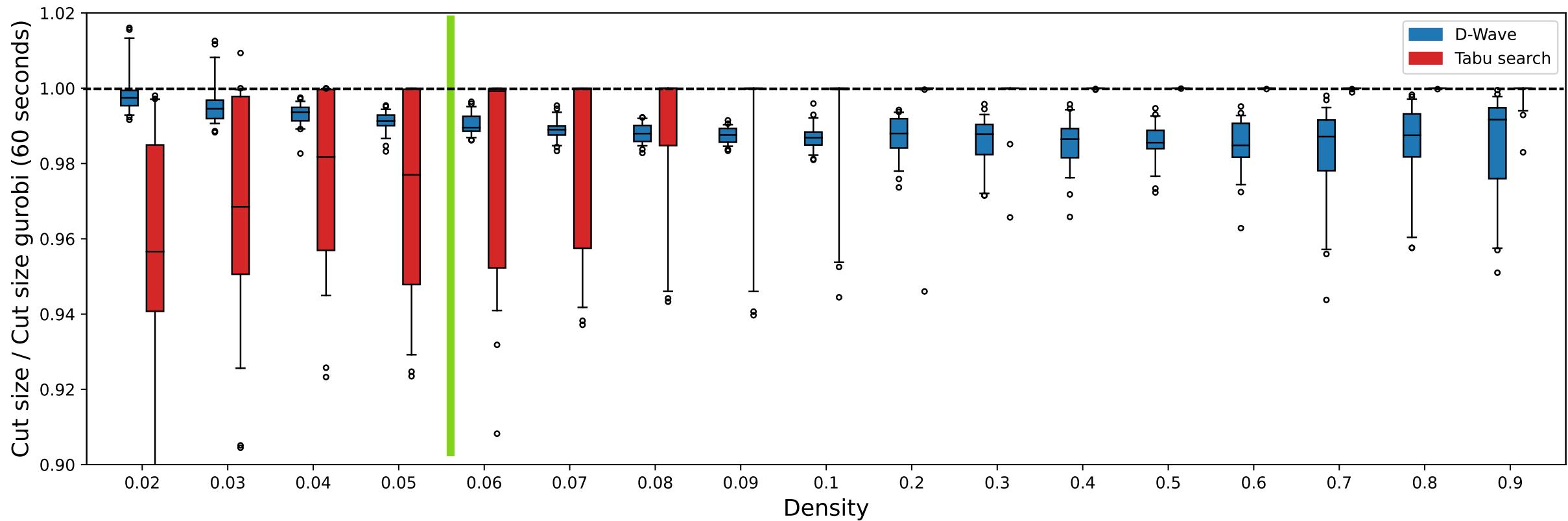
Time windows:

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1s: Tabu Search

█ Performance intersection with Tabu Search
— Reference solution





2- Results – MWIS problem

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Time windows: 1s: D-Wave (5000 shots)
60s: Gurobi (reference solution) 1s: Tabu Search
Random Greedy Search (5000 runs)



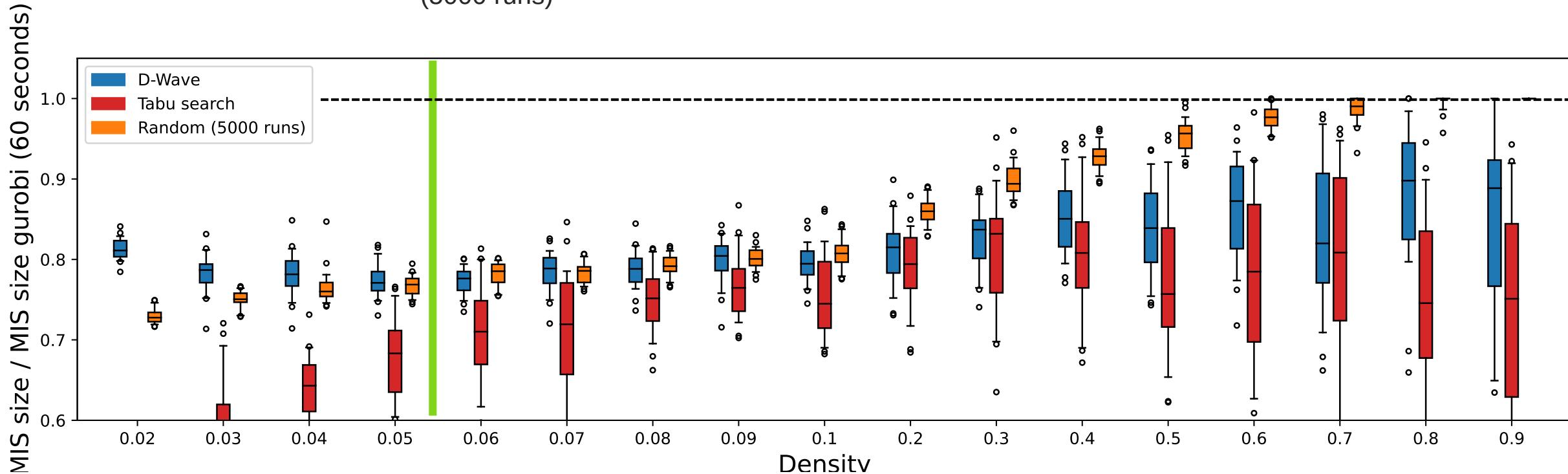
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— Performance intersection with Random Greedy search
- - - Reference solution



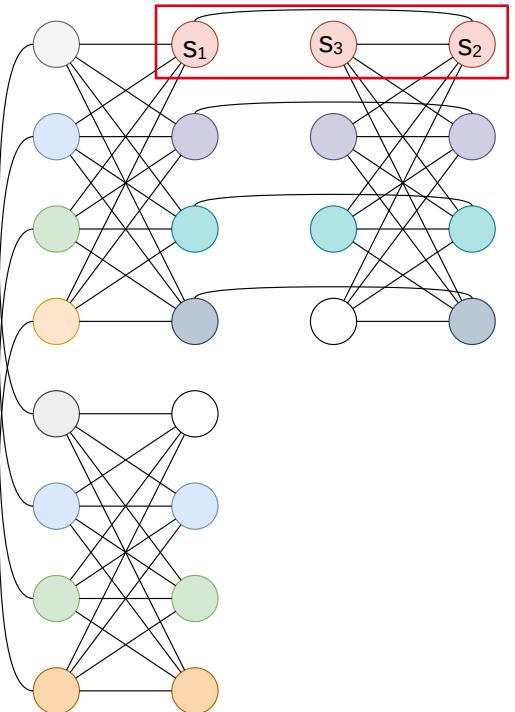


3 ■ Contribution #2 Increasing the performance of Quantum Annealers



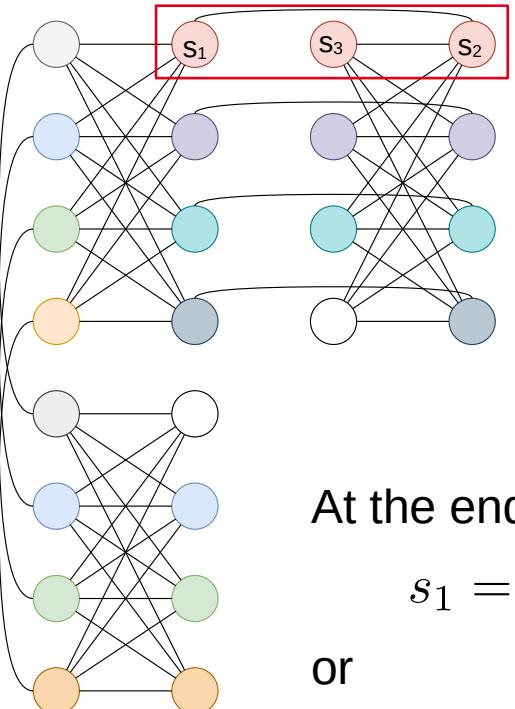
3- The embedding problem II

- Embedding step produces chains of qubits



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At the end of the annealing:

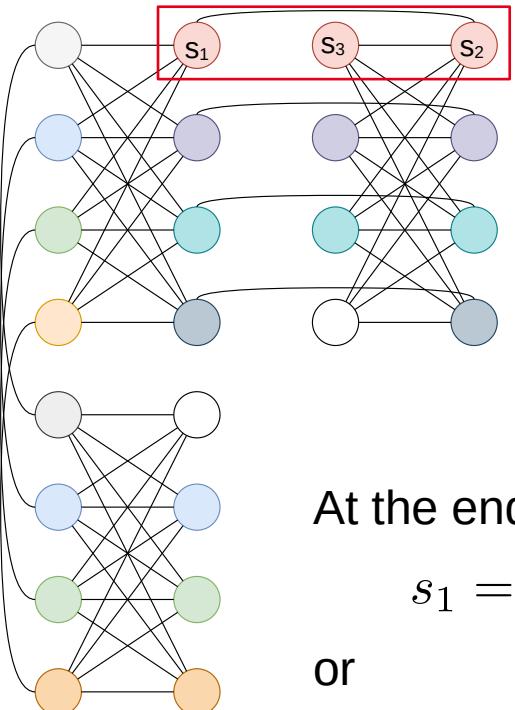
$$s_1 = s_2 = s_3 = +1$$

or

$$s_1 = s_2 = s_3 = -1$$

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- Embedding step produces chains of qubits



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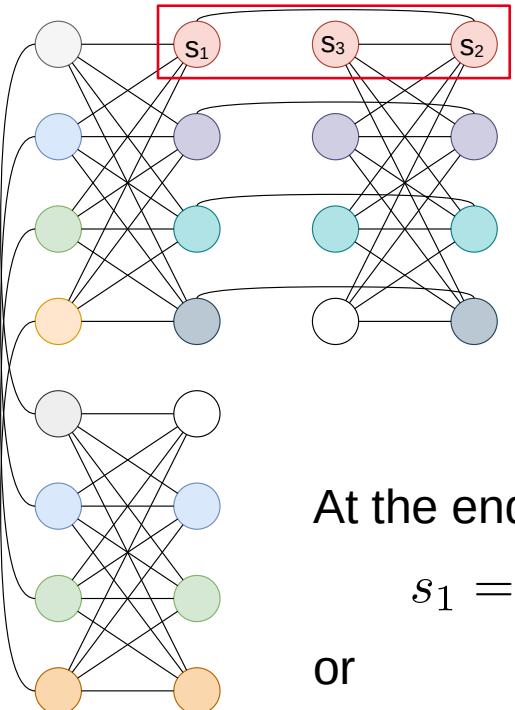
$$s_1 = s_2 = s_3 = -1$$

Let's set:

$$J_{12} = -\infty \quad J_{23} = -\infty$$

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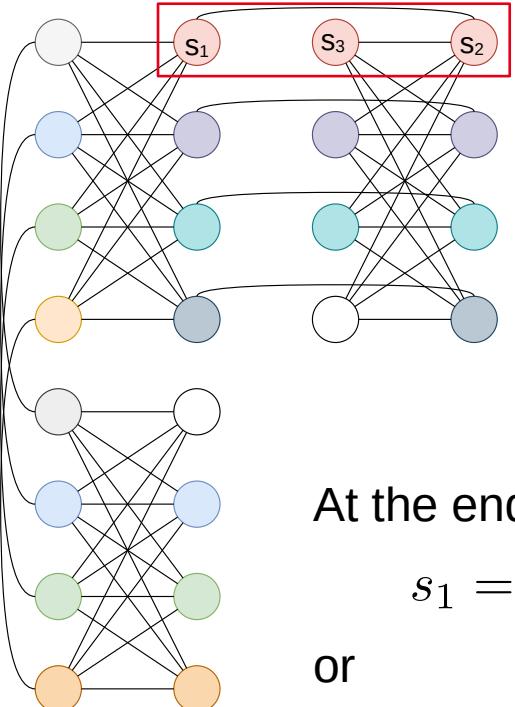
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Limited programming
range

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- Embedding step produces chains of qubits
- How to set the chain between these qubits ?



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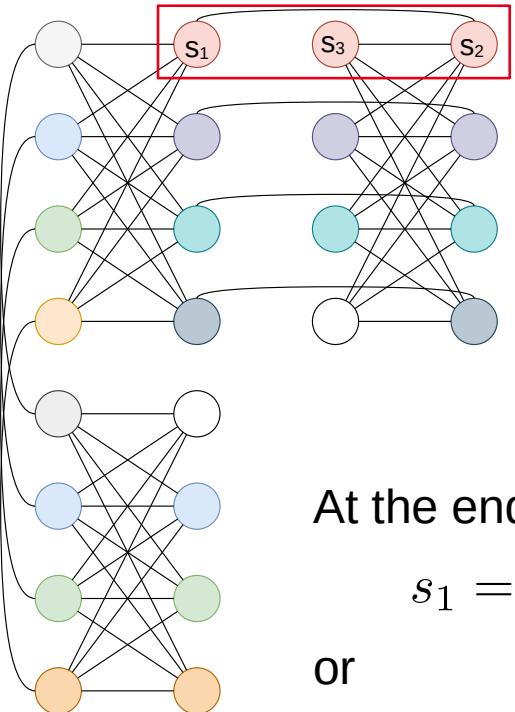
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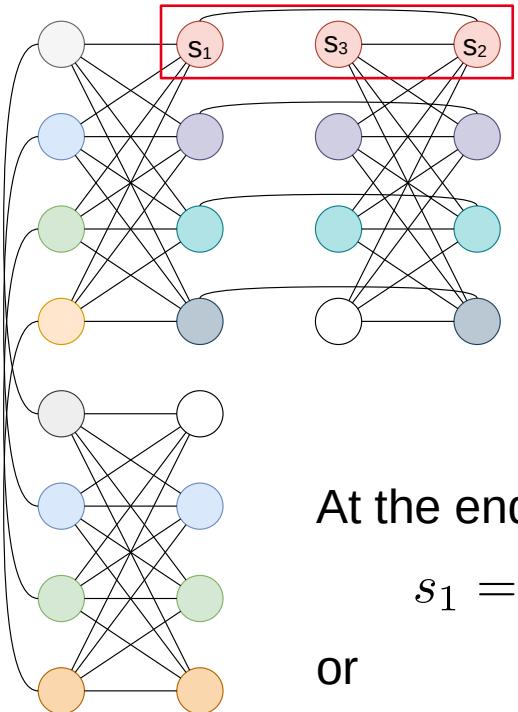
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- How to set the chain between these qubits ?

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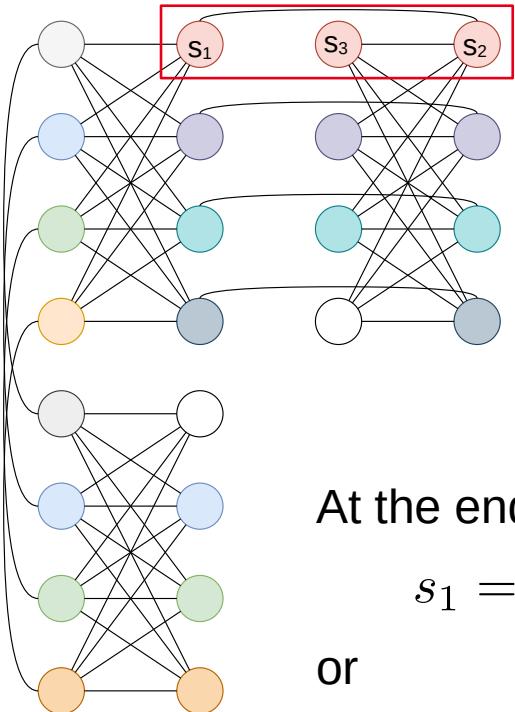
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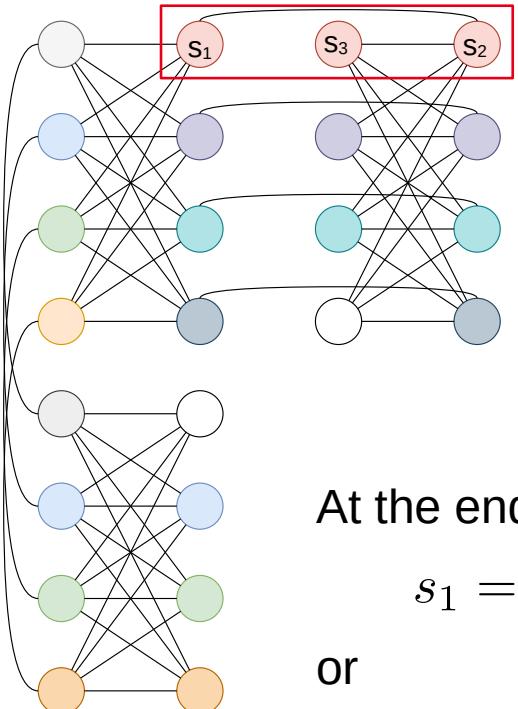
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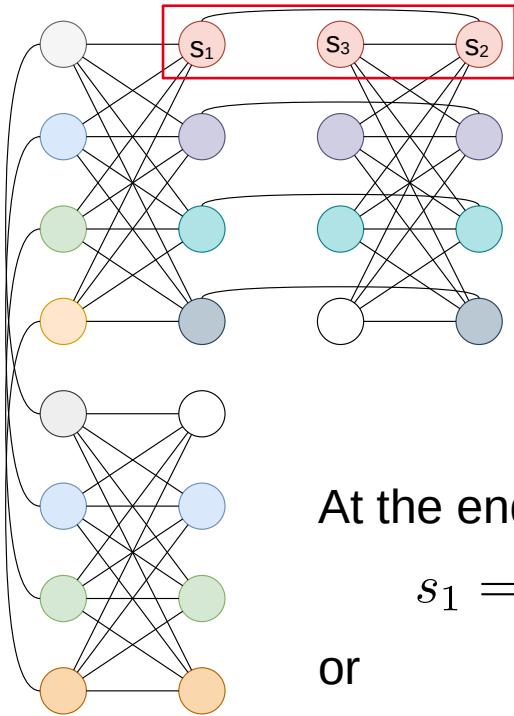
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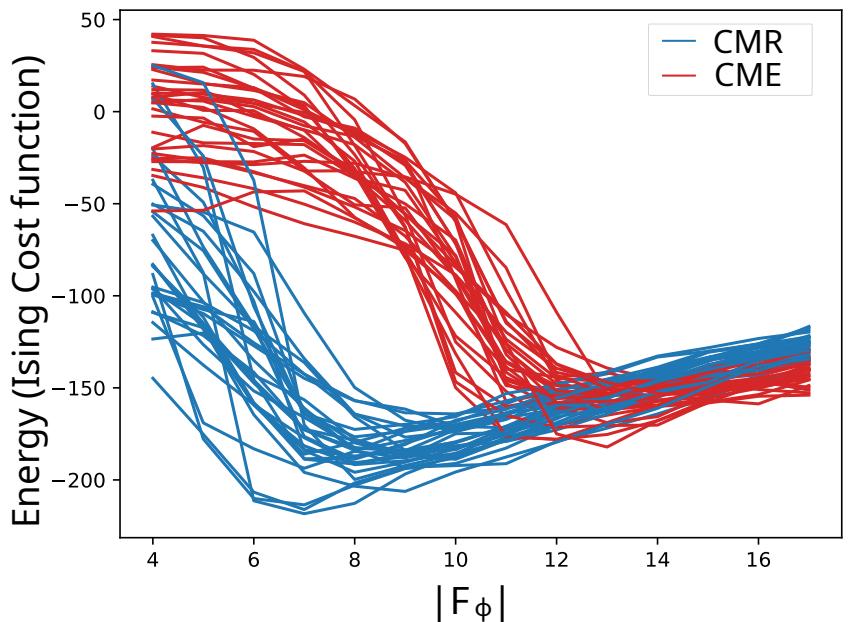


Limited programming range

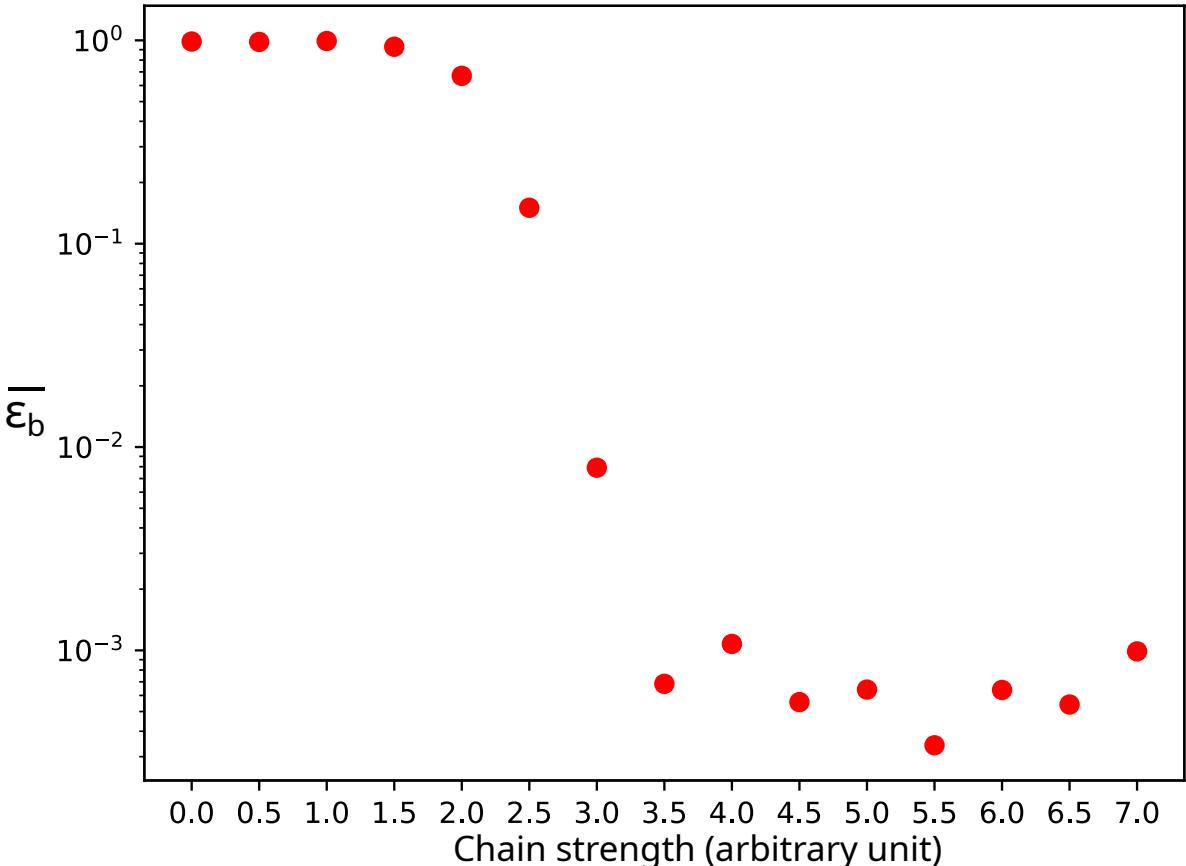
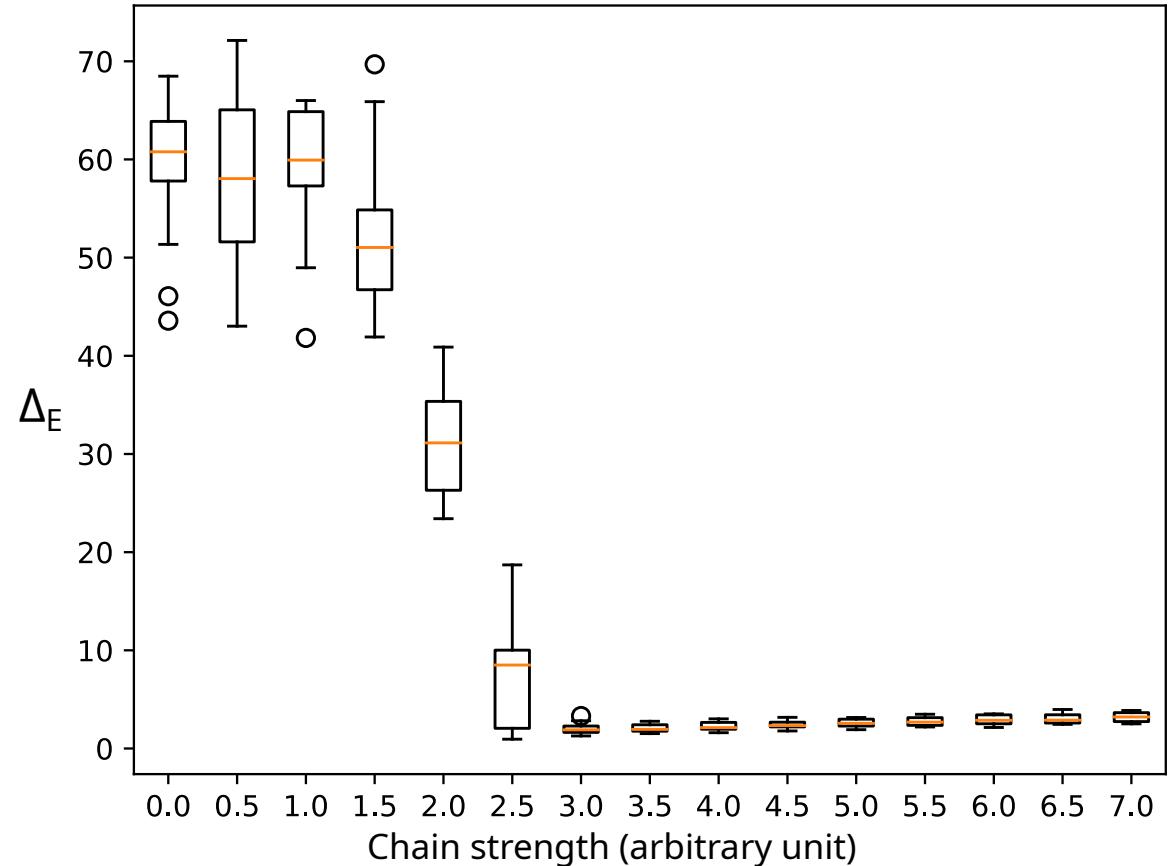
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- Global chain strength: $F_\phi < 0$

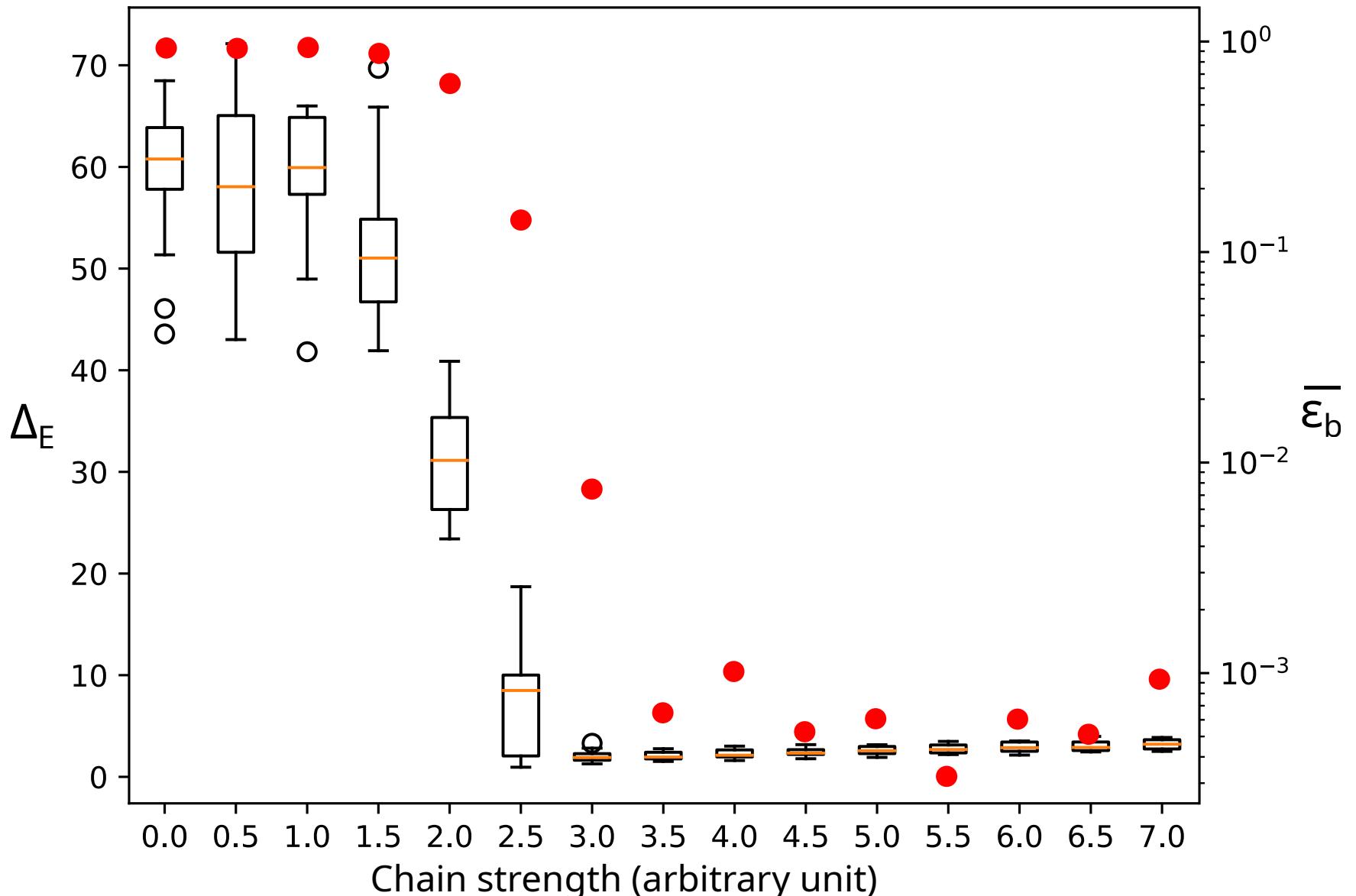


3- Chain scan & Chain breaks



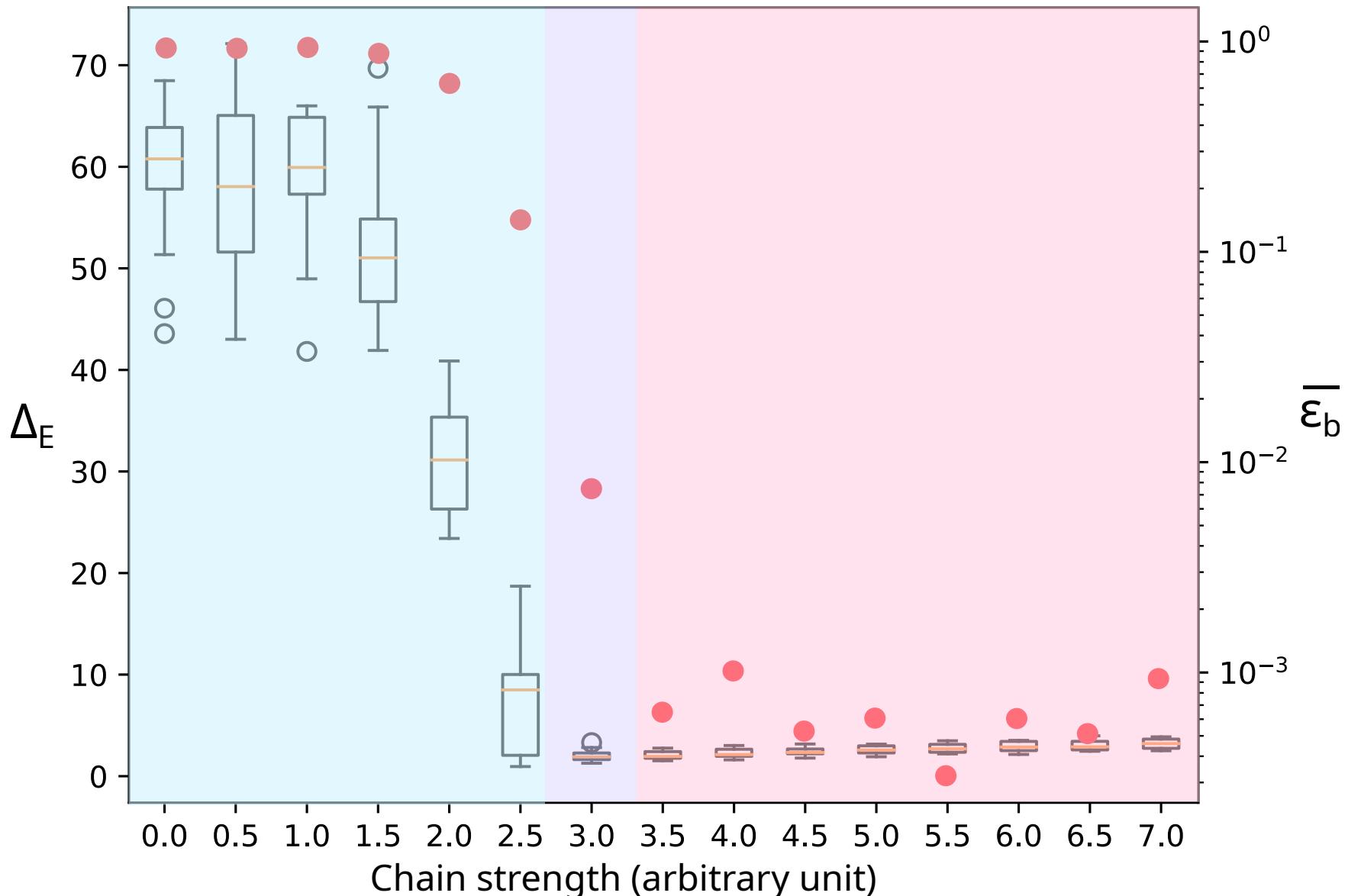


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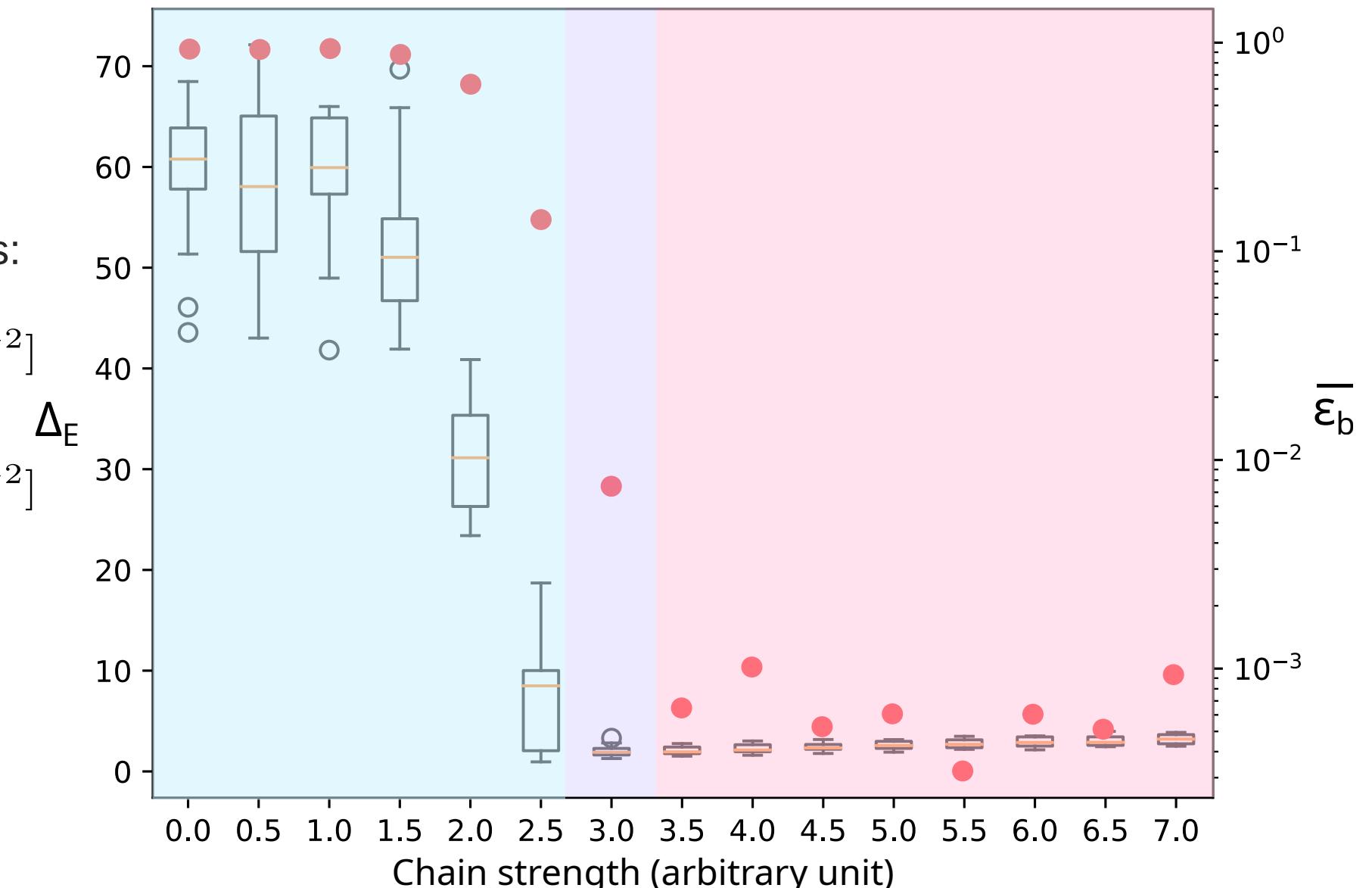
Chain break intervals:

Advantage2:

$$[6 \times 10^{-3}, 2 \times 10^{-2}]$$

Advantage6.4:

$$[2 \times 10^{-2}, 5 \times 10^{-2}]$$



3- Results

- 30 instances of unweighted max-cut for each density
- Shot / pre-processing step: 128
- Final run number of shots:
Advantage2: 3072
Advantage6.4: 4096

Advantage2_prototype2.2			Best cut size	Cut size improvement				Step
Instance size	Density	Embedding		min	max	mean	std	
$n = 40$	0.1	CMR	66.4	+0%	+0%	+0%	0%	5.4
	0.5	CMR	243	+0%	+2%	+0.2%	0%	5.8
	0.9	CMR	362.8	+2.1%	+8.2%	+5%	1.6%	4.6
$n = 80$	0.1	CMR	235.5	+0%	+0%	+0%	0%	4.7
	0.5	CME	804	+9.8%	+17.2%	+12.5%	0.2%	4.2
	0.9	CME	1435	+2%	+4.7%	+3.2%	0.6%	4.2
Advantage6.4			Best Cut size	Cut size improvement				Step
Instance size	Density	Embedding		min	max	mean	std	
$n = 100$	0.1	CMR	355.9	+0%	+0.3%	+0%	0%	4.5
	0.5	CME	1271.4	+5.6%	+14.5%	+8.8%	1.8%	2.7
	0.9	CME	2243	+1.4%	+3.7%	+2.5%	0.5%	3.7
$n = 170$	0.1	CMR	950.8	-2.1%	+0.6%	-0.5%	0.5%	2.1
	0.5	CME	3631.4	+2.8%	+6.2%	+4.5%	0.7%	2.1
	0.9	CME	6519.4	+0.4%	+1.4%	+0.8%	0.2%	3.2

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Conclusion

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 - Generation of Instances of **gradual difficulty**



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Conclusion

- Methodology to benchmark Quantum Annealers
 - Generation of Instances of **gradual difficulty**
 - D-Wave Quantum Annealers are mostly adapted for **very sparse** and **unconstrained** problems
 - **Fairness** evaluation of benchmarking methods:
 - Fairness of instances set (diversity) [GLS23]
 - Fairness of the benchmarking protocol (quantumbenchmarkzoo.org)



- Improvements of Quantum Annealers' parameter setting
 - The parameters list is long and complex (interdependence between parameters)
 - The optimization of these parameters should be included in the TTS metric



Perspectives

■ Benchmarking



Perspectives

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 - Possible extension of our approach to benchmark the QAOA (Optimal planted swapping network)



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“NISQ will not change the world by itself, at least not right away; instead we should regard it as a step toward more powerful quantum technologies we hope to develop in the future.”

J. Preskill [Pre21]



Thank you !



A big thank to ACTIF (Association CEA des thésard.e.s d'Île-de-France)



New president: Lise.jolicoeur@cea.fr



Publications

G. Bettonte, V. Gilbert, D. Vert, S. Louise, and R. Sirdey, “Quantum approaches for wcet-related optimization problems,” in Lecture Notes in Computer Science - ICCS 2022, p. 202-217, Springer International Publishing, 2022

V. Gilbert, J. Rodriguez, S. Louise, and R. Sirdey, “Solving higher order binary optimization problems on nisq devices: experiments and limitations,” in Lecture Notes in Computer Science - ICCS 2023, p.224-232, Springer Nature Switzerland, 2023

V. Gilbert, S. Louise, and R. Sirdey, “Taqos: A benchmark protocol for quantum optimization systems,” in Lecture Notes in Computer Science – ICCS 2023, p.168-176, Springer Nature Switzerland, 2023

V. Gilbert, and S. Louise, “Quantum annealers chain strengths: A simple heuristic to set them all,” in Lecture Notes in Computer Science - ICCS 2024, p.292-306, Springer Nature Switzerland, 2024

V. Gilbert, J. Rodriguez, and S. Louise “Benchmarking quantum annealers with near-optimal minor-embedded instances,” in 2024 IEEE International Conference on Quantum Computing and Engineering (QCE), IEEE, 2024



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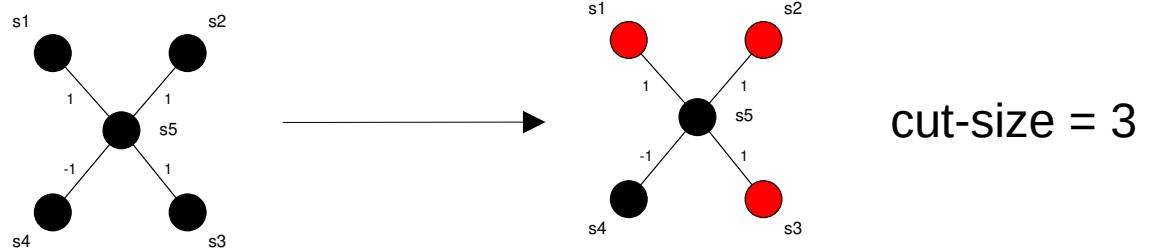
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4 ■ Appendix

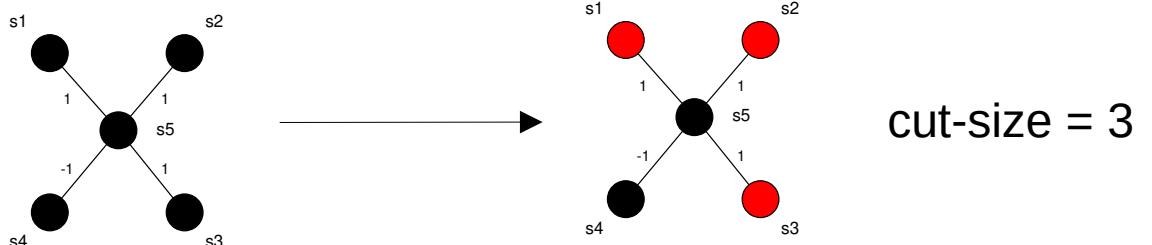
Selection of best Annealing time

- Quality metrics are problem-dependent
 - Max-cut problem: The cut size

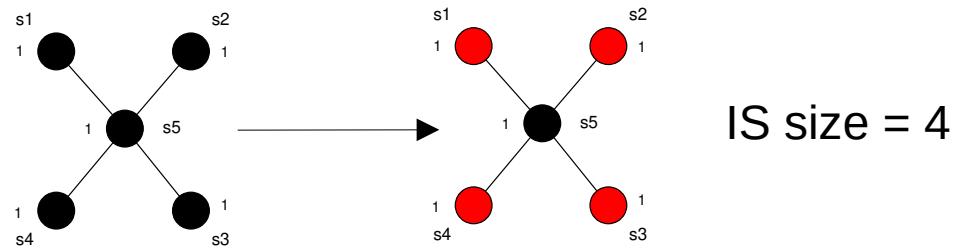


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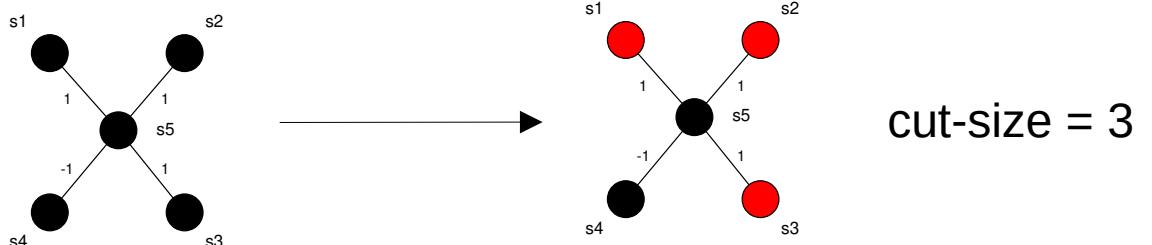


- Maximum Independent set problem: The size of the independent set.

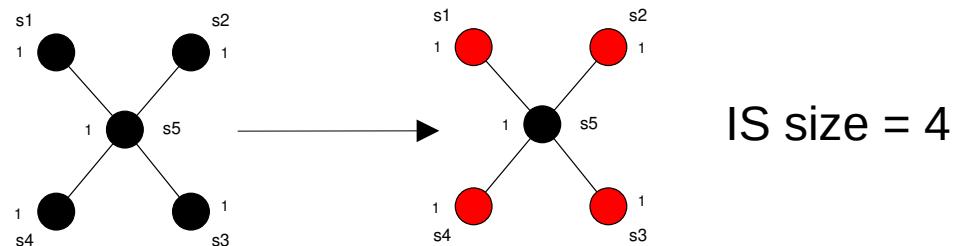


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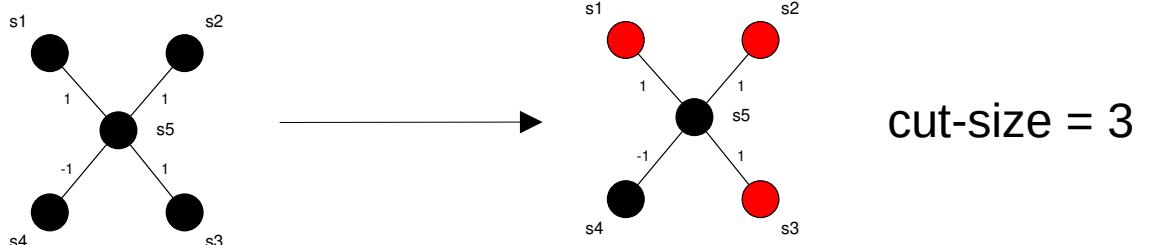


- Time measurement
 - 60s time window for branch & bound algorithm
 - 1s time window for Tabu Search (C implementation)
 - 1s time window for D-Wave Q

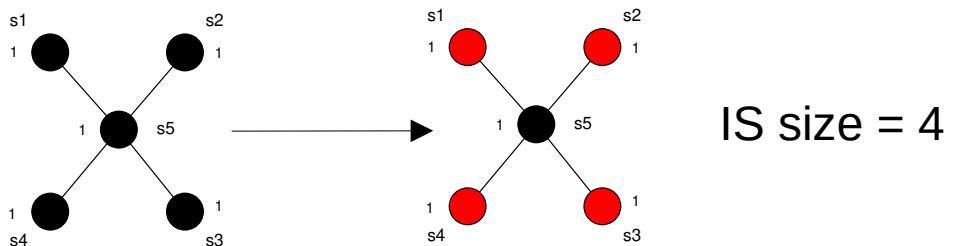


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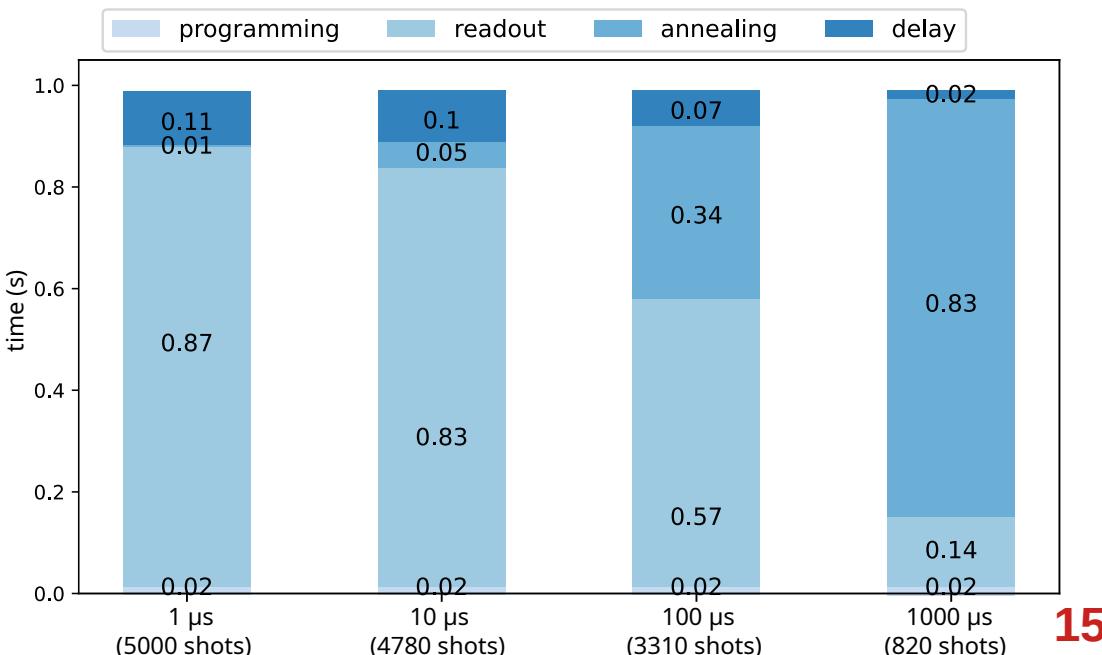
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Optimal mapping

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$$n_{\phi(v)^*} = \begin{cases} 1 & \text{if } \deg(v) \leq c_{\text{phys}} \\ 2 & \text{if } c_{\text{phys}} < \deg(v) \leq (2c_{\text{phys}} - 2) \\ \left\lceil \frac{\deg(v) - (2c_{\text{phys}} - 2)}{c_{\text{phys}} - 2} \right\rceil + 2 & \text{otherwise} \end{cases}$$



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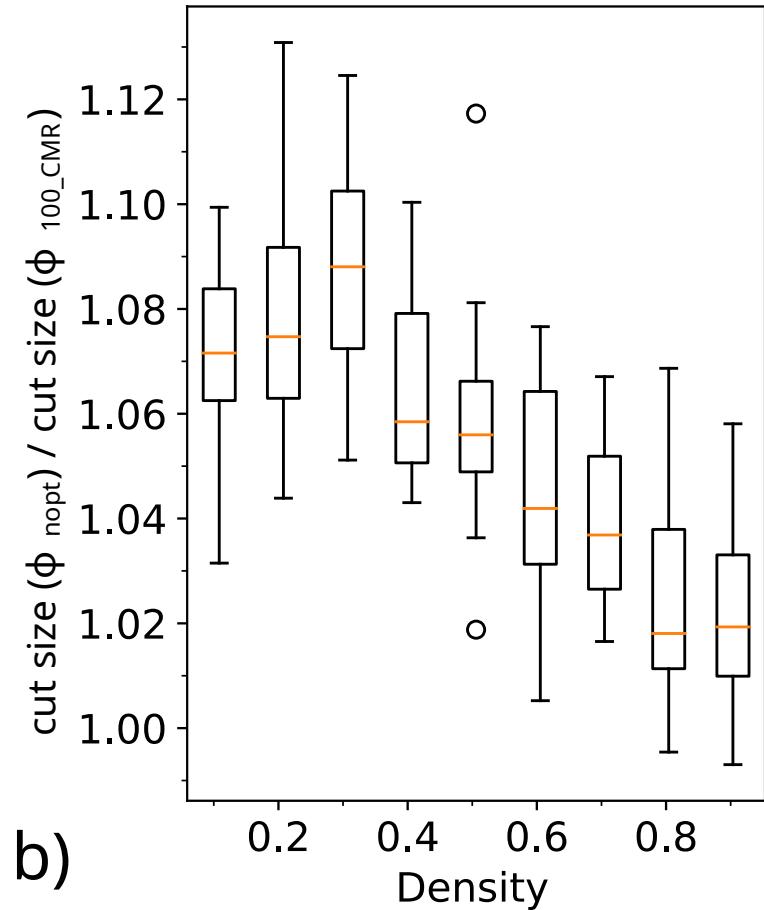
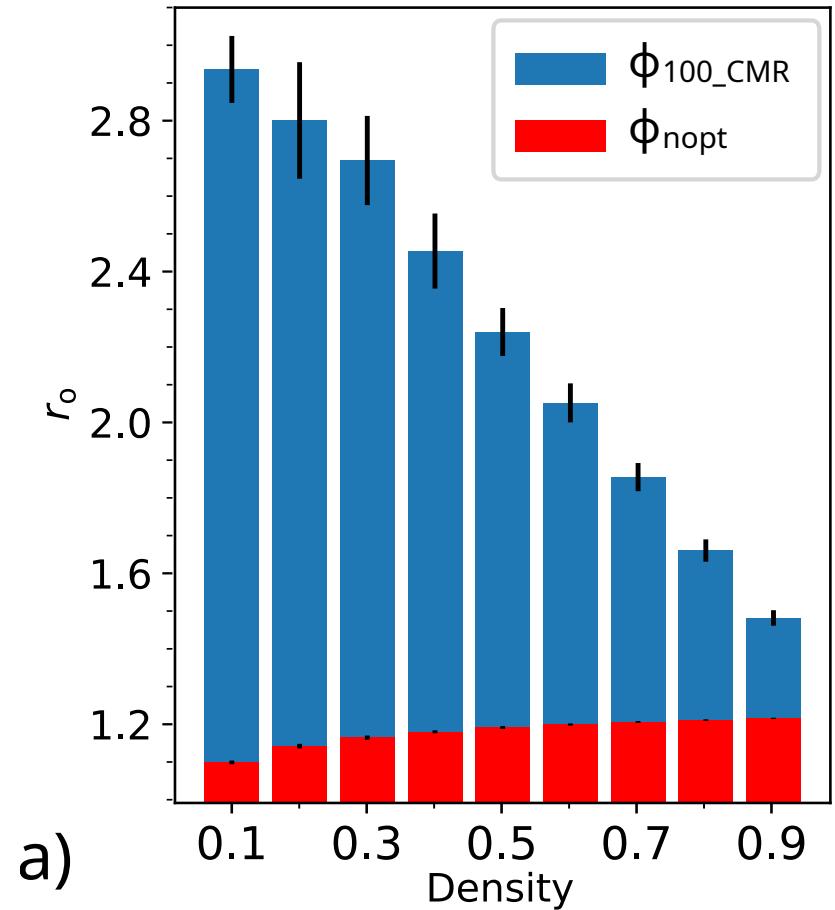
- Compute the overhead ratio considering this bound:

$$r_o = \frac{n_\phi}{\sum_{v \in V_s} n_{\phi(v)^*}}$$

Optimal mapping

- Comparison of the performance of our generation method against state of the art embedding method

at $d = 1$
 $|V_s| = 100$
 $|V_t| = 982$

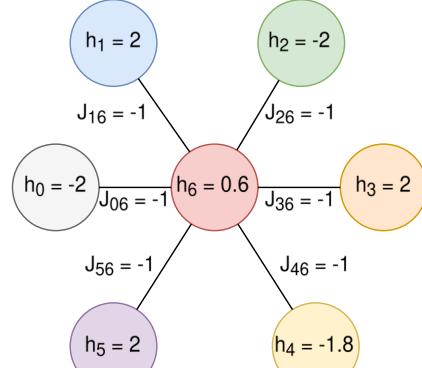


Assumption: Instances with less duplicated qubits are more easily solved by QA => Seems to be true

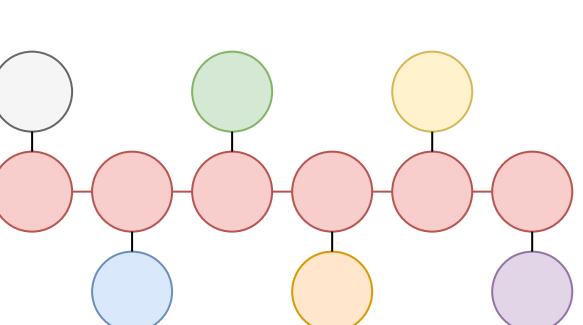


III- Shape of the logical qubit

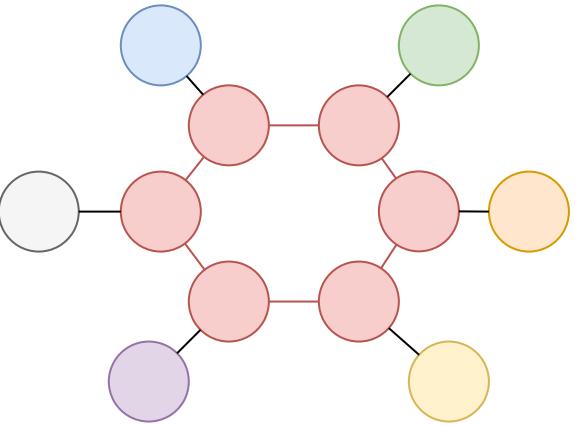
Source graph



Chain encoding



Cycle encoding



Clique encoding

