1. Exercise 1

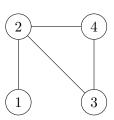
1.1 Consider the graph described by the following adjacency matrix:

$$W = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{pmatrix}.$$

- **a.** Are there any self-loops in this graph? What are the in-degree and the out-degree of node 2? Does the graph have any sinks (i.e., nodes whose out-degree is 0) or sources (i.e., nodes whose in-degree is 0)? List the elements of the out-neighborhood of node 2.
- **b.** Assume the links are directed as follows: (1,2),(2,3),(2,4) and (3,4). Give the adjacency matrix of the obtained directed graph. What are the in-degree and the out-degree of node 2 now? Does the directed graph have any sinks or sources? List the elements of the out-neighborhood of node 2.

Solution

a. Self-loops correspond to non-zero diagonal entries of W, so there are no self-loops. $W_{ij}=1$ if there is a link from i to j. Hence, the in-degree of node 2 is $\sum_i W_{i,2}=3$ and its out-degree is $\sum_j W_{2,j}=3$. This is an undirected graph, since the adjacency matrix is symmetric. There are no sinks or sources. The out-neighborhood of node 2 is $\{1,3,4\}$. The graph is depicted in the figure below:

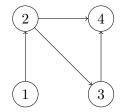


b. The adjacency matrix of the directed version of the graph is

$$W = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{pmatrix}.$$

Now, the in-degree of node 2 is $\sum_i W_{i,2} = 1$ and the out-degree is $\sum_j W_{2,j} = 2$. This graph has one source, node 1, and one sink, node 4. The out-neighborhood of node 2 is $\{3,4\}$. The graph is depicted in the figure below:

1



1.2 Consider the weighted graph in Figure 1.1.

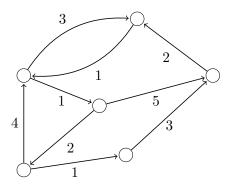


Figure 1.1: The graph for Problem 1.2.

- a. Determine its weight matrix.
- **b.** Compute the average degree.

Solution

a. With the following numbering of the nodes

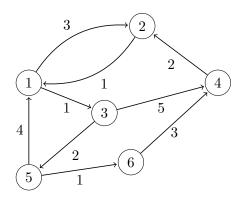


Figure 1.2: The graph for Problem 1.2.

the adjacency matrix can be written as

$$W = \begin{pmatrix} 0 & 3 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 5 & 2 & 0 \\ 0 & 2 & 0 & 0 & 0 & 0 \\ 4 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 3 & 0 & 0 \end{pmatrix}.$$

- **b.** The average degree of the graph, computed either based on the outdegree $\bar{w} = \sum_i w_i/n = \mathbb{1}'W\mathbb{1}/n$ or on the in-degree $\bar{w} = \sum_i w_i^-/n = \mathbb{1}'W'\mathbb{1}/n$, where n = 6 is the number of nodes, is equal to 3.7 (in both cases).
- 1.3 a. Give an example of a graph that is balanced but not undirected.
 - **b.** Give an example of an undirected graph that is not regular.

Solution

a. See Figure 1.3 below for an example.



Figure 1.3: Balanced but not undirected

b. See Figure 1.4 for an example.



Figure 1.4: Undirected but not regular

1.4 Consider the graph described by the following adjacency matrix:

$$W = \begin{pmatrix} 0 & 2 & 1 \\ 2 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix}.$$

- **a.** Determine the Laplacian matrix L and the normalized weight matrix P of this graph and verify that $L\mathbb{1} = 0$, $P\mathbb{1} = \mathbb{1}$.
- **b.** Check if the graph is balanced and/or regular.
- **c.** Determine the number of connected components in the graph.

Solution

a. The out-degree vector, Laplacian and normalized weight matrix are

$$w = \begin{pmatrix} 3 \\ 3 \\ 2 \end{pmatrix} \quad L = \begin{pmatrix} 3 & -2 & -1 \\ -2 & 3 & -1 \\ -1 & -1 & 2 \end{pmatrix} \text{ and } P = \begin{pmatrix} 0 & 2/3 & 1/3 \\ 2/3 & 0 & 1/3 \\ 1/2 & 1/2 & 0 \end{pmatrix}.$$

b. w is an invariant vector of P' if and only if the graph is balanced. Hence, since in this case P'w = w, the considered graph is balanced. Also, since the graph is undirected (because its adjacency matrix is symmetric), it can be immediately concluded that it is also balanced. $\mathbb{1}$ is an invariant vector of P' if and only if the graph is regular. Hence, since in this case $P'\mathbb{1} \neq \mathbb{1}$, the graph is not regular.

- ${f c.}$ The graph is strongly connected, hence the number of connected components is 1.
- 1.5 Consider two undirected graphs described by the following two adjacency matrices W_1 and W_2 :

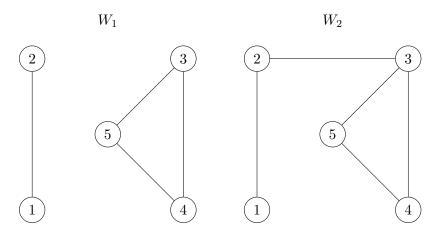
$$W_1 = \begin{pmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 & 0 \end{pmatrix}, \quad W_2 = \begin{pmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 & 0 \end{pmatrix}.$$

Draw the graphs. For each graph, answer the following questions:

- **a.** Is the graph strongly connected or not?
- **b.** What is the multiplicity of 0 as an eigenvalue of the Laplacian matrix, and of 1 as an eigenvalue of the normalized adjacency matrix?

Solution

a. The graphs \mathcal{G}_1 and \mathcal{G}_2 with adjacency matrices W_1 and W_2 , respectively, are depicted below:



The graph \mathcal{G}_1 is not strongly connected and has two connected components, $\mathcal{C}_1 = \{1,2\}$ and $\mathcal{C}_2 = \{3,4,5\}$. The graph \mathcal{G}_2 is strongly connected, hence it only has one connected component. The following adjacency matrix, A, gives the same graph as W_1 , just with different ordering of the nodes,

$$A = \begin{pmatrix} 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \end{pmatrix}$$

It is not as obvious to determine if the graph is strongly connected or not from A as it was from W_1 ; indeed, matrix W_1 is in block-diagonal form (from which we can clearly deduce that the graph is not connected), while matrix A is not.

- **b.** If we form the condensation graphs for \mathcal{G}_1 and \mathcal{G}_2 , we see that \mathcal{G}_1 is formed by two isolated nodes, while \mathcal{G}_2 has a single node. The algebraic and geometric multiplicities of the eigenvalue 1 for P and the eigenvalue 0 for L is the same as the number of sinks in the condensation graph, hence 2 for \mathcal{G}_1 and 1 for \mathcal{G}_2 .
- **1.6** Consider a k-regular simple graph (where all nodes have degree k).
 - **a.** Show that the vector $\mathbb{1}$ is an eigenvector of the adjacency matrix, associated with eigenvalue k.
 - **b.** Recall that the eigenvectors of a symmetric matrix are orthogonal. Exploit this property to show that for this graph, no eigenvector x of the adjacency matrix that is associated to an eigenvalue $\lambda \neq k$ can have all nonnegative entries.

Solution

- **a.** In each row, the adjacency matrix W of an unweighted k-regular graph has k entries equal to one, while the other entries are zero. Hence, for every node i, one has that $\sum_{j} W_{ij} = k$, so that $W\mathbb{1} = \mathbb{1}k$, i.e., $\mathbb{1}$ is an eigenvector of W associated to eigenvalue k.
- **b.** Since W is symmetric, the eigenvector v associated with $\lambda \neq k$ is orthogonal to $\mathbb{1}$ i.e. $v'\mathbb{1} = 0$. Thus $\sum_i v_i = 0$ which makes it clear that not all v_i can be non-negative, apart from the trivial solution $v_i = 0 \ \forall i$,
- 1.7 Every tree is a bipartite graph (i.e., a graph such that the node set can be partitioned into two nonempty subsets so that there are no links between nodes in the same subset). Propose a method to split the nodes in two subsets, thus proving the statement.

Solution

Consider two sets A and B. For any tree graph, pick one node: assign the node itself to set A and its offspring to set B. Further, assign the offspring of the offspring to set A and their offspring to set B. Continue in this manner for the entire graph. In this way, the nodes in set A are only connected to the nodes in set B, and vice versa.

- **1.8** A graph with $n \ge 5$ nodes has the out-degree distribution $(p_0, p_1, p_2) = (0, \frac{4}{5}, \frac{1}{5})$ and the in-degree distribution $(p_0^-, p_1^-, p_2^-, p_3^-) = (\frac{2}{5}, \frac{1}{5}, \frac{1}{5}, \frac{1}{5})$.
 - **a.** Does the graph have any sink nodes?
 - **b.** Does the graph have any source nodes?
 - **c.** Draw a graph with n=5 nodes and with the given degree distribution.

Solution

- **a.** No, because the number of nodes with out-degree 0 is $np_0 = 0$.
- **b.** Yes, because the number of nodes with in-degree 0 is $np_0^- = 2n/5 > 0$.
- ${f c.}$ A graph with 5 nodes and the given degree distribution is shown in Figure 1.5.

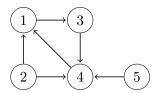
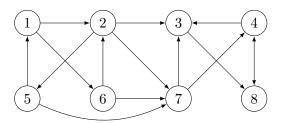


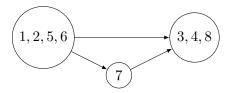
Figure 1.5: Graph in the answer to Problem 1.8.

1.9 Consider the graph \mathcal{G} represented in the figure below. Merge the nodes of the connected components and draw the condensation graph.

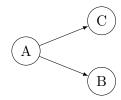


Solution

In the condensation graph, each node corresponds to a connected component of the original graph. A connected component is a set of nodes such that there is a path between each pair of nodes in this set. The condensation graph thus becomes the following:



1.10 Let \mathcal{H} be the directed graph displayed below.



For each of the following statements, say if it is true or false, providing motivation.

a. Every graph \mathcal{G} with condensation graph \mathcal{H} is not strongly connected.

- **b.** Every graph \mathcal{G} with condensation graph \mathcal{H} is acyclic.
- **c.** Every graph \mathcal{G} with condensation graph \mathcal{H} does not possess a globally reachable node.
- **d.** Every graph \mathcal{G} with condensation graph \mathcal{H} is a tree.

Solution

- **a.** True. That a graph is strongly connected means that there is a path between any pair of nodes i and j in the graph. Since there are no paths between nodes in \mathcal{G} corresponding to nodes C and B in the condensation graph and nodes in \mathcal{G} corresponding to node A in the condensation graph, the graph \mathcal{G} is not strongly connected.
- **b.** False. The graph \mathcal{G} is not necessarily acyclic, since any of the nodes in the condensation graph in general corresponds to several nodes in graph \mathcal{G} between which there can be cycles.
- **c.** True. Any node in C or B can be reached from any node in A, but nodes in A and C cannot be reached from nodes in B, and nodes in A and B cannot be reached from nodes in C. So there is no node that is reachable from every other node in the graph.
- **d.** False. Any node in the condensation graph corresponds to a collection of nodes in \mathcal{G} which are fully connected. Thus, they do not in general have a tree structure.
- **1.11** Let \mathcal{G} be a directed ring graph with n nodes.
 - a. Find the eigenvector centrality.
 - **b.** Find the Katz centrality with $\mu = 1$.
 - **c.** Find the Katz centrality with $\mu_1 = 1$ and $\mu_i = 0$, $\forall i > 1$ and $\beta = 0.5$. Reflect on the difference between **b** and **c**.

Solution

a. Recall that the eigenvector centrality vector is the eigenvector of W' corresponding to the largest-in-magnitude eigenvalue λ_W , which is real and non-negative. With $\mathcal G$ being a directed ring the adjacency matrix is

$$W = \begin{pmatrix} 0 & 1 & 0 & \dots & \dots & 0 \\ \vdots & 0 & 1 & 0 & \dots & 0 \\ \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \vdots & & \ddots & \ddots & 0 \\ 0 & 0 & \dots & \dots & 0 & 1 \\ 1 & 0 & \dots & \dots & 0 & 0 \end{pmatrix},$$

which implies that $W'\mathbb{1} = \mathbb{1}$. For the normalized weight matrix, $P = D^{-1}W$, it always holds that $\lambda_P = 1$, i.e., that 1 is the largest-in-magnitude eigenvalue. Since, in this case, W = P (the row sums of

W are already 1), we know that $\lambda_W = \lambda_P = 1$. Thus, 1 is the sought eigenvector.

b. The Katz centrality vector is the solution of

$$x = \left(\frac{1-\beta}{\lambda}\right) W'x + \beta\mu,$$

where $\beta \in (0,1]$, $\mu = 1$, and $\lambda = 1$ is again the largest-in-magitude eigenvalue of W'. For the directed ring graph we thus have

$$x_i = (1 - \beta)x_{i-1} + \beta, \quad i = 1, \dots, n,$$

where $x_0 = x_n$. From this expression we can make an educated guess that the centrality should be equal in each node. Inserting $x_i = k \, \forall i$ into the system of equations gives

$$k = (1 - \beta)k + \beta,$$

which yields k = 1 independently of β . As we found a valid solution the Katz centrality, which is unique, is therefore given by $x_i = 1$ for all i.

c. From the previous problem, we know that the Katz centrality for this ring graph fulfills

$$x_i = (1 - \beta)x_{i-1} + \beta\mu_i, \quad i = 1, \dots, n,$$

which for $\beta = 0.5$ and $\mu_1 = 1, \mu_i = 0 \forall i > 1$ yields

$$x_1 = 0.5x_N + 0.5$$
$$x_i = 0.5x_{i-1} \quad i > 1$$

Let $x_1 = c$, then $x_i = 0.5^{i-1}c \ \forall i > 1$. This gives that

$$x_1 = c = 0.5 \cdot 0.5^{N-1}c + 0.5$$
$$c = \frac{0.5}{1 - 0.5^N}$$

Thus the Katz centrality is given by

$$x_i = \frac{0.5^i}{1 - 0.5^N} \quad \forall i$$

1.12 Consider a k-regular undirected graph (i.e., an undirected graph in which every node has degree k). Find the eigenvector centrality, as well as the Katz centrality vector for $\mu = 1$.

Solution

For a k-regular graph we know that the normalized adjacency matrix $P = \frac{1}{k}W$. We know from Proposition 2.4 in the lecture notes that, for a regular graph, the largest-in-magitude eigenvalue for P and P' is 1 with

corresponding eigenvector $\mathbb{1}$ for P'. Thus, the dominating eigenvalue of W' is k, with corresponding eigenvector $\mathbb{1}$.

The Katz centrality is given as the solution x of

$$x = \frac{1 - \beta}{k} W' x + \beta \mu.$$

Since every row of W' has the same sum and every element of μ is 1, intuitively we assume that all entries of x are the same, i.e. $x_i = \gamma$ for all $i \in \mathcal{V}$. Inserting this assumption into the equation system gives

$$x = \frac{1-\beta}{k}W'x + \beta\mu \quad \Leftrightarrow \quad \gamma\mathbb{1} = \left(\frac{1-\beta}{k}\right)k\gamma\mathbb{1} + \beta\mathbb{1},$$

which gives $\gamma = 1$. The Katz centrality vector is unique, which makes x = 1 the only solution.

1.13 Consider the graph in Figure 1.6.

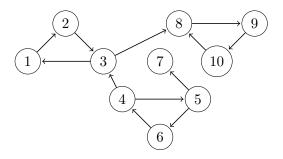
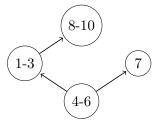


Figure 1.6: Graph in the answer to Problem 1.13.

- **a.** Draw the associated condensation graph. How many sinks does it have?
- **b.** Is the original graph connected? If you were able to delete one node (and all the links connected to it), would your answer change?
- **c.** Determine all possible invariant probability distributions $\pi = P'\pi$ for the given network.

Solution

a. The condensation graph is drawn in the figure below.



Hence, the condensation graph has two sinks.

- **b.** The original graph is not connected: indeed, the condensation graph is not formed by a single node. No single node can be deleted in the original graph in order to turn it into a connected graph.
- c. Recall that for a graph $\mathcal{G} = (\mathcal{V}, \mathcal{E}, W)$ the following holds: (i) Any convex combination of invariant probability distributions of \mathcal{G} is an invariant probability distribution of \mathcal{G} . (ii) For every sink in the condensation graph \mathcal{H} of \mathcal{G} , there exists an invariant probability distribution supported on the connected component of \mathcal{G} corresponding to such a sink of \mathcal{H} . These distributions are referred to as extremals. (iii) Every invariant probability distribution can be obtained as a convex combination of the extremal invariant probability distributions associated to the sinks of the condensation graph.

Since the condensation graph \mathcal{H} of our graph \mathcal{G} has exactly two sinks, there should thus be two extremal probability distributions, $\pi^{(1)}$ and $\pi^{(2)}$. These are supported on the connected components corresponding to the sinks, i.e.,

$$\pi^{(1)} = (0, 0, 0, 0, 0, 0, 0, 0, \pi_8^{(1)}, \pi_9^{(1)}, \pi_{10}^{(1)})' \quad \text{and}$$

$$\pi^{(2)} = (0, 0, 0, 0, 0, 0, \pi_7^{(2)}, 0, 0, 0)'.$$

The only way that the distribution $\pi^{(1)}$ can be invariant is if the non-zero components are chosen to be the same, since the last three equations of the system $\pi^{(1)} = P'\pi^{(1)}$ read $\pi_8^{(1)} = \pi_9^{(1)}$, $\pi_9^{(1)} = \pi_{10}^{(1)}$, and $\pi_{10}^{(1)} = \pi_8^{(1)}$, respectively. Furthermore, probability distributions require that the elements of the distribution vectors sum up to 1. Thus, the extremal distributions are given by the probability vectors

$$\pi^{(1)} = (0, 0, 0, 0, 0, 0, 0, 1/3, 1/3, 1/3)', \quad \pi^{(2)} = (0, 0, 0, 0, 0, 0, 1, 0, 0, 0)'.$$

Finally, point (iii) gives that any invariant probability distribution is given by

$$\pi = \alpha \pi^{(1)} + (1 - \alpha)\pi^{(2)}, \quad \alpha \in [0, 1].$$