

Data Structures

Self-balancing

binary search tree

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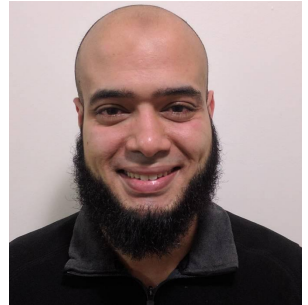
Teaching, Training and Coaching since more than a decade!

Artificial Intelligence & Computer Vision Researcher

PhD from Simon Fraser University - Canada

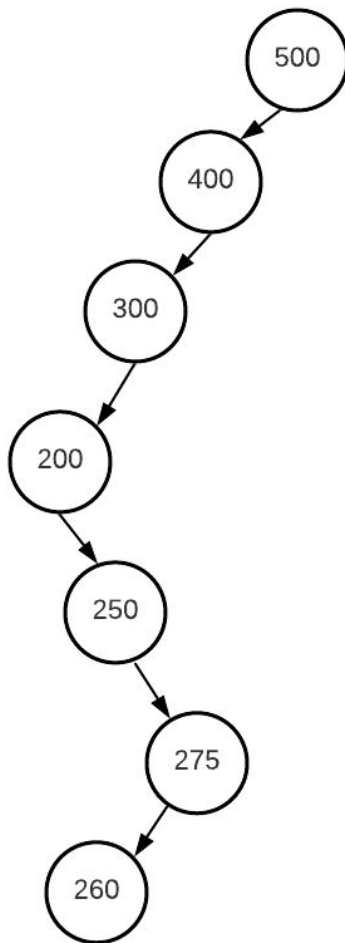
Bachelor / Msc from Cairo University - Egypt

Ex-(Software Engineer / ICPC World Finalist)



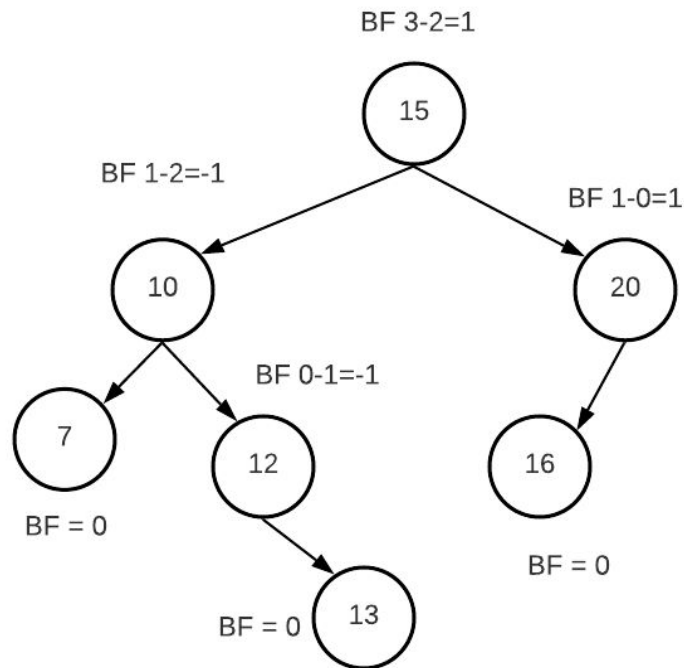
Recall: Degenerate BST

- Each node has 1 child
- From a performance perspective, it is similar to a linked list, with several operations being done in $O(n)$
- However, in any tree, the bigger the gap between $\log(n)$ and the tree height, the less efficient the tree will be!



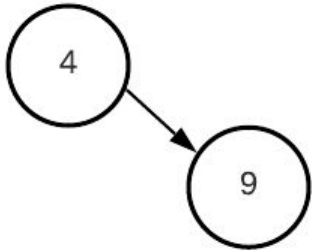
Recall: Balanced BST

- The **difference** in height between the left and right subtrees is never greater than 1.
 - $\text{Height}(\text{left}) - \text{Height}(\text{right})$
 - **For visualization**, assume $\text{height}(\text{leaf}) = 1$
- Let's compute this difference for each node (**Balance factor**)
 - -7 means = right height greater by 7
 - -1, 0, 1 \Rightarrow BBST
 - $|\text{bf}| > 1 \Rightarrow$ imbalanced tree



Consider height and balance_factor()

- In the last line of insert() function, we call:
`update_height()`



```
class Node:
    def __init__(self, val=None, left=None, right=None):
        self.val = val
        self.left = left
        self.right = right
        self.height = 0 # 0 for leaf

    def ch_height(self, node): # child height
        if not node:
            return -1 # -1 for null
        return node.height # 0 for leaf

    def update_height(self): # call in end of insert function
        self.height = 1 + max(self.ch_height(self.left), self.ch_height(self.right))

    def balance_factor(self):
        return self.ch_height(self.left) - self.ch_height(self.right)

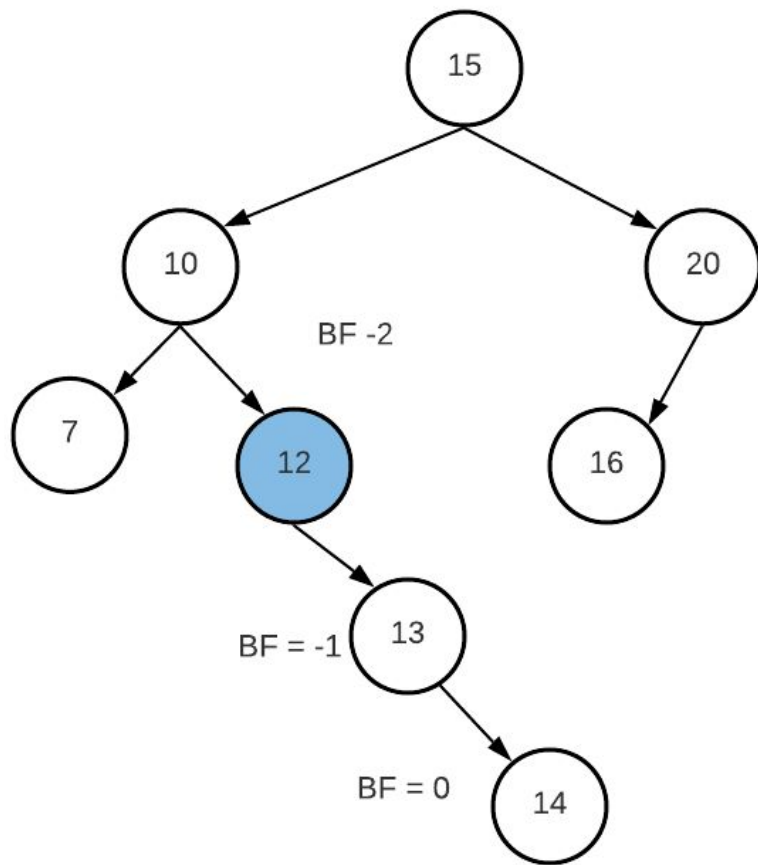
    def is_leaf(self):
        return not self.left and not self.right
```

Maintaining BBST

- Insertion in BST depends on input order - and will generate unbalanced trees.
- To maintain themselves as a BBST, **self-balance** BST trees follow **a change and fix** approach
 - E.g. Insert a new element in BST
 - Do we have an unbalanced BST?
 - Yes \Rightarrow Fix the tree so that the balance factor is once again between -1 to 1 (i.e. $|BF| \leq 1$)
- There are several trees that maintain themselves to remain as a BBST
 - **AVL Trees**: one of the oldest and simplest ways
 - Red-Black Trees, Splay Trees, Treaps
- This section focuses on AVL Trees

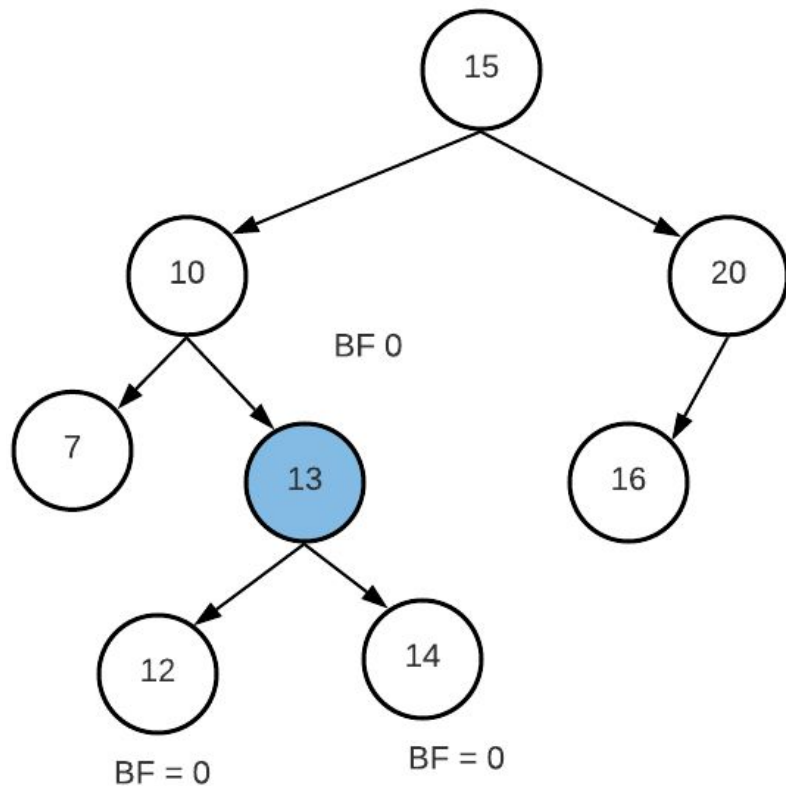
Change (insert 14)

- Let's insert 14
- Now, recompute the Balance Factor
- We reach leaf nodes recursively, then go back to the parents, and so on
- Once you detect $|BF| > 1$, then we know this particular subtree is no longer BBST
- Think for 2 minutes about how to restructure subtree (12) to make it a balanced BST



Fix (Tree Rotation)

- If we pushed node 13 **up** and node 12 **down** left, this subtree is fixed
 - Observe, we did not change other subtrees
 - Observe, the tree remains BBST
- This kind of systematic change is called **tree rotation**
- If, after insertion, we keep fixing corrupted sub-trees from the bottom up, we can maintain a BBST



AVL Demo

AVL Tree

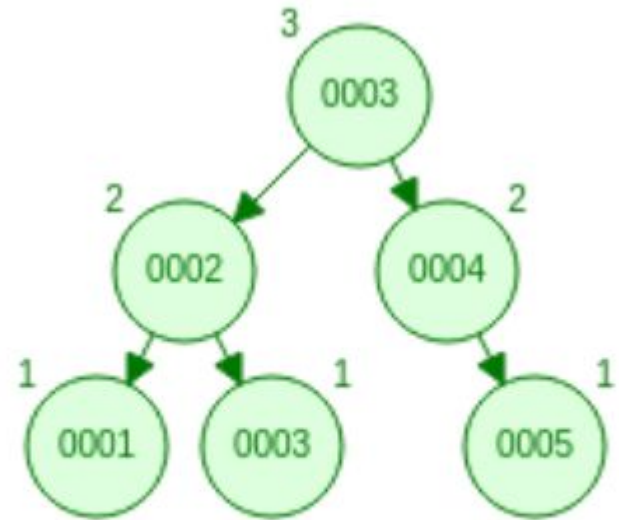
Insert

Delete

Find

Print

- Many websites help you to visualize insertion/deletion in AVL trees with real-time online demos
- Try inserting a batch of unsorted numbers, and watch the tree **restructure** itself to accommodate them



“Acquire knowledge and impart it to the people.”

“Seek knowledge from the Cradle to the Grave.”