

# *Data Structures*

## Asymptotic Complexity (2)

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How can skilled programmers find the order  
**VERY quickly?**

# Tips

- This code involves ~8 steps
- **Ignore** all constants
- Constants don't affect the overall order when  $n$  is large
- All constants are simply FIXED numbers
- So, the code is just  **$O(1)$**

```
def constant_order1():    #  $O(1)$   
    start = 6  
    end = 100  
    mid = (end - start) // 2  
    if mid % 2 == 0:  
        del mid
```

# Tips

- Too many steps?
- Yes; but a FIXED number!
  - Useless with very large N
- Ignore them
- Tip: Ignore anything that doesn't involve our variables (e.g. n)

```
def constant_order2():  
    start = 7  
    end = 0  
    for i in range(1000):  
        end += end * 2 + start
```

# Tips

- Search for loops that are based on  $n$
- A single loop is  $O(n)$
- Two nested loop is  $O(n^2)$
- Triple nested loops is  $O(n^3)$
- *And so on*
- This code shows a single loop
- inside it is a FIXED number of operations  
⇒ IGNORE!

```
def linear1(n):      #  $O(n)$ 
    sum = 0
    for i in range(n):
        # All below are  $O(1)$ 
        x = 2 + 3 * 4
        sum += i
        sum += 2
        sum += x
```

# Tips

- 2 parallel loops.
- Each is a single loop
- Each loop depends on  $n$ 
  - One is  $10n$  and one is  $5n$
  - Ignore these constants
- Practically:  $10n + 5n = 17n$
- Ignore constants  $\Rightarrow O(n)$
- Tip: what is the **deepest**?
  - A single loop  $\Rightarrow O(n)$

```
def linear2(n):      #  $O(n)$ 
    for i in range(10 * n):
        constant_order1()

    for i in range(5 * n):
        constant_order1()
```

# Tips

- This is  $5n \times 3n$  loop steps
  - Multiplied with some factor from all these FIXED steps
  - Overall  $O(n^2)$
- Tip: nested loops  $\Rightarrow O(n^2)$

```
def quadratic1(n): #  $O(n^2)$ 
    cnt = 0
    for i in range(5 * n):
        for j in range(3 * n):
            cnt += 1
            constant_order1()
```

# Tips

- We have 2 parallel blocks
  - Nested loops:  $O(n^2)$
  - Linear loop:  $O(n)$
- Tip: focus on the biggest
  - As it dominates the operations
  - $n^2 + n \Rightarrow n^2$

```
def quadratic2(n): #  $O(n^2)$ 
    cnt = 0
    for i in range(5 * n):
        for j in range(3 * n):
            cnt += 1
            constant_order1()

    for i in range(10 * n):
        constant_order1()
```



# Tips

- 2 parallel blocks
  - 3 nested loops
  - 1 loop
- But in the 3 nested loops
  - One loop is just a fixed # of operations
  - Again **ignore constant operations**
  - This 3rd loop is useless (ignore)
- Total:  $15000 n^2 + 10n \Rightarrow n^2$

```
def quadratic3(n): #  $O(n^2)$ 
    cnt = 0
    for i in range(5 * n):
        for j in range(3 * n):
            for k in range(1000):
                cnt += 1
                constant_order1()

    for i in range(10 * n):
        constant_order1()
```

# Tips

- 2 parallel blocks
  - A single loop
  - A single loop
- So  $O(n)$ ? No, there is a trick
- The 2nd loop is not linear in  $n$ 
  - It iterates  $3n^2$  steps
- The order of the second loop is  $O(n^2)$
- Tip: observe if the number of loop operations is fixed,  $n$ ,  $n^2$  etc...
  - This value decides the order!

```
def quadratic4(n): #  $O(n^2)$ 
    for i in range(10 * n):
        constant_order1()

    for i in range(3 * n * n):
        constant_order1()
```

# Tips

- As this code has 3 nested loops
  - Each depends on  $n$
  - It is  $O(n^3)$

```
def cubic1(n): #  $O(n^3)$ 
    cnt = 0
    for i in range(n):
        for j in range(n):
            for k in range(n):
                cnt += 1
```

# Tips

- 2 parallel blocks
  - 3 nested loops  $\Rightarrow n^3$
  - 2 nested loops  $\Rightarrow n^2$ 
    - Don't be distracted by the size of the constant (1000)
    - IGNORE constants
- $n^3 + n^2 \Rightarrow O(n^3)$ 
  - Always focus on the **biggest**

```
def cubic2(n): #  $O(n^3)$ 
    cnt = 0
    for i in range(n):
        for j in range(n):
            for k in range(n):
                cnt += 1

    for i in range(1000 * n):
        for j in range(1000 * n):
            cnt += 1
```

# Tips

- Why not  $O(n^3)$ ?
  - Our loops are  $n^2$ ,  $n$ , and  $n^3$  respectively
  - Total  $O(n^6)$
- Again, **double check** if the number of operations in the loop is based on a *fixed constant number,  $n$ ,  $n^2$  etc...*

```
def f(n):    #  $O(n^6)$ 
    cnt = 0
    for i in range(n * n):
        for j in range(n):
            for k in range(n * n * n):
                cnt += 1
```

# Tips

- F1 is  $O(n^3)$ 
  - $n*n$  then  $n$
- F2 has a single loop:  $O(n)$ 
  - But its body is NOT constant!
  - The body contains a call that is  $O(n^3)$
- Overall  $O(n^4)$
- Tip
  - Imagine we copy-pasted f2 in f1
  - You should see clearly it is  $O(n^4)$  steps in total
- Tip
  - Double check each line of code (like a function call): is fixed or variable?

```
def f1(n): #  $O(n^3)$ 
    cnt = 0
    for i in range(n * n):
        for j in range(n):
            cnt += 1

def f2(n): #  $O(n^4)$ 
    for i in range(n):
        f1(i)
```

# Tips

- Sometimes our function depends on several variables
- Total is  $6nm$  steps
- Drop the constant value  $\Rightarrow O(nm)$

```
def f3(n, m):    #  $O(nm)$ 
    cnt = 0
    for i in range(2 * n):
        for j in range(3 * m):
            cnt += 1
```

# Tips

- 2 parallel blocks
  - Block 1:  $O(nm)$
  - Block 2:  $O(n^2)$
- Which is bigger? We don't know
- Total:  $O(nm + n^3)$

```
def f4(n, m):    #  $O(nm + n^3)$ 
    cnt = 0
    for i in range(2 * n):
        for j in range(3 * m):
            cnt += 1

    for i in range(n * n * n):
        cnt += 1
```



# Polynomial Order

- Today we discussed polynomial order functions (format  $n^k$ )
  - $n^0 = 1$  (const),  $n^2$ ,  $n^3$  and so on
- Intuition: code is doing some hundred million steps  $\Rightarrow \sim 1$  second (not really)
- From the table,  
the **bigger** your  $O()$   
the **slower** your code
- There are other worse families
  - E.g.  $O(n^n)$  or  $O(n!)$

	n=100	n=1000	n=1000000
$O(n)$	100	1000	<b>1000000</b>
$O(n^2)$	10000	<b>1000000</b>	Too much
<b><math>O(n^3)</math></b>	<b>1000000</b>	1000000000	Too much
$O(n^4)$	100000000	Too much	Too much

# Overall

- Try to keep these tips in mind
- Don't be overly systematic. Some codes are trickier!
  - E.g. 3 nested while-loops might actually just do  $10n$  steps. It's not necessarily  $O(n^3)$ 
    - We will see several cases during my other courses
- Whenever you write a code from now on, always compute its time order
  - This is how you will develop your skill!

*“Acquire knowledge and impart it to the people.”*

*“Seek knowledge from the Cradle to the Grave.”*