

# *Data Structures*

## Circular Queue

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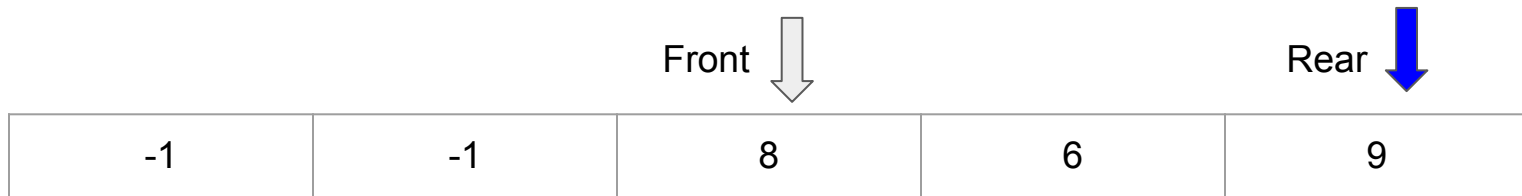
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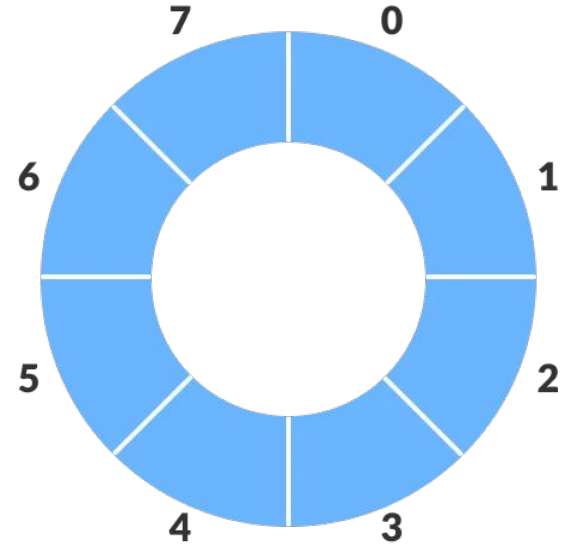
# Array-based: Front-Rear approach

- We'll use two **indices**; front and rear; which represent the start and end of the array respectively
  - When we **enqueue** an element, we add it to the rear  $\Rightarrow O(1)$
  - When dequeue an element, we **shift the front** index to the right  $\Rightarrow O(1)$
- Enqueue 3: **ERROR Queue is full!**
- However, there are empty slots at the beginning!
  - This is a critical **drawback** of this approach



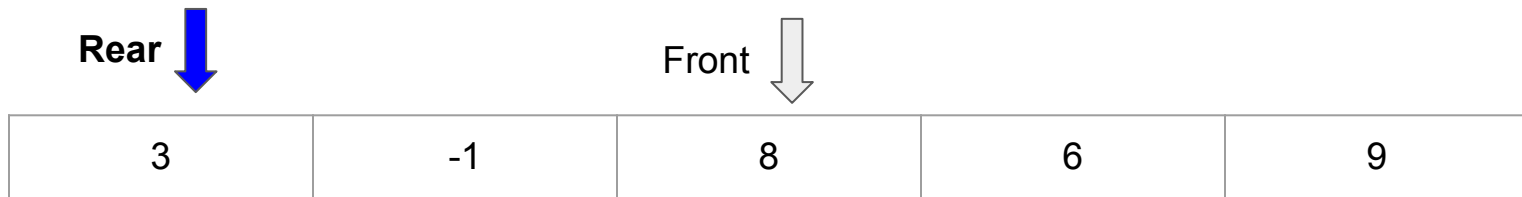
# Circular Queue

- There is a simple way to solve the previous space issue
- Simply, think of the array as a circle
  - On the right, an array of 8 elements is arranged as a circle
- This helps to visualize the circular queue. In the diagram, after the final element in the array, there is another element, which is at position 0
- Now, the queue is full IFF all elements are in use



# Array-based Circular queue

- We will have 2 indices: front and rear representing start to end in array
  - When we enqueue element we add it in rear
  - When we dequeue element we shift front to the right  $\Rightarrow O(1)$
- Enqueue 3
  - Now, move from the last index to index 0 - and add the new element
  - Observe how 'rear' is now **BEFORE** front



# Initial values for rear & front

- There are several approaches for that
  - For all approaches, you must be **consistent** throughout the entire implementation
  - **Careful** conditions for IsEmpty and IsFull
- Possible initializations
  - rear = front = -1 [initially equal]
  - rear = front = 0 [initially equal]
  - rear = -1 and front = 0 [initially !equal]
  - rear = size - 1 and front = 0 [initially !equal]
- **int added\_elements = 0;**
  - To both simplify coding and avoid tricky conditions, maintain a counter for the number of elements added

# Circular Queue: Data Structure

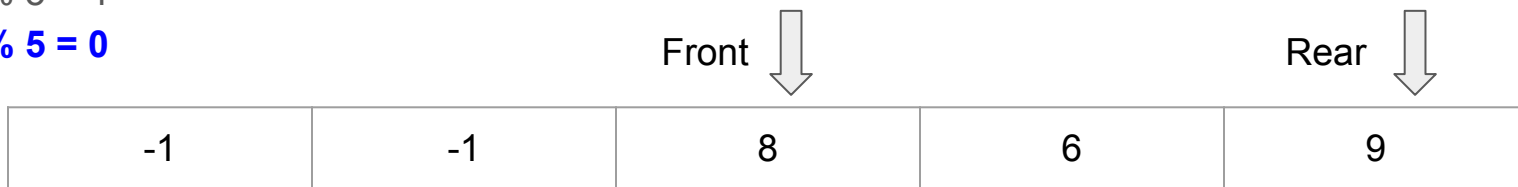
- We initially use **front = rear = 0**
  - To add a new element, **add** at **rear** and **move rear**
  - To dequeue an element, get the **front**, and then move **front**

```
1 class Queue {  
2     int size { };  
3     int front { 0 };  
4     int rear { 0 };  
5     int added_elements { };  
6     int *array { };  
7 }
```

# Circular Queue: Move index

- To move the index forward, consider:
  - If this is the **last element** in the array, make the next position = 0
  - We can do this with an if condition (efficient)
  - Or using the modulus operator (%)
- Assume size = 5. Let's try positions from 0 to 5
  - $0 \% 5 = 0$
  - $1 \% 5 = 1$
  - $2 \% 5 = 2$
  - $3 \% 5 = 3$
  - $4 \% 5 = 4$
  - **$5 \% 5 = 0$**

```
int next(int pos) {  
    //return (pos + 1) % size;  
  
    ++pos;  
    if (pos == size)  
        pos = 0;  
    return pos;  
}
```



# Let's Simulate: Queue of size 5

- Initially an empty queue. Initially,  $\text{rear} = \text{front} = 0$
- Observe: The EMPTY queue has  $\text{rear} == \text{front}$





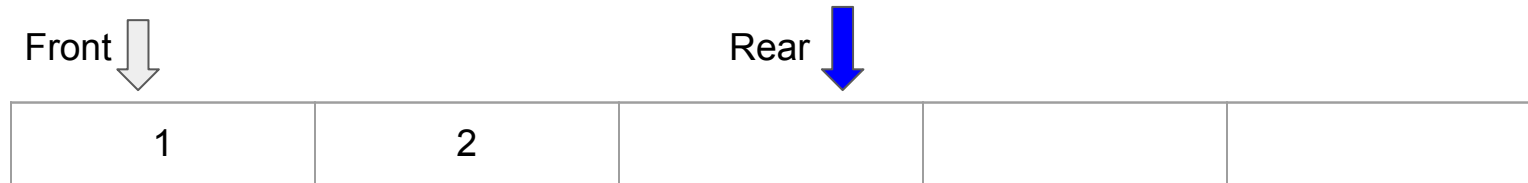
# Let's Simulate: Queue of size 5

- Enqueue (1)  $\Rightarrow$  Add to the rear position and move it



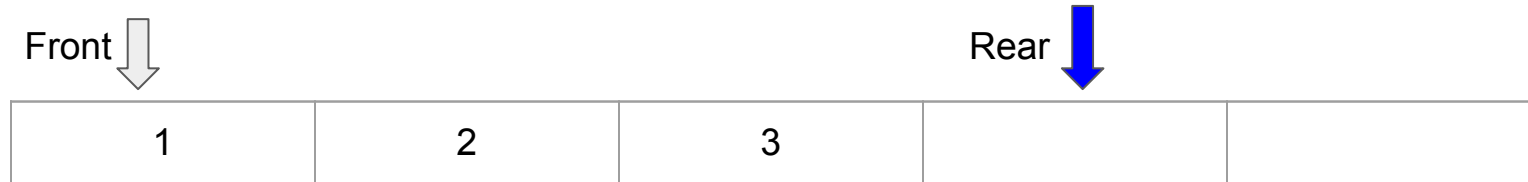
# Let's Simulate: Queue of size 5

- Enqueue (1), Enqueue (2)



# Let's Simulate: Queue of size 5

- Enqueue (1), Enqueue (2), Enqueue (3)



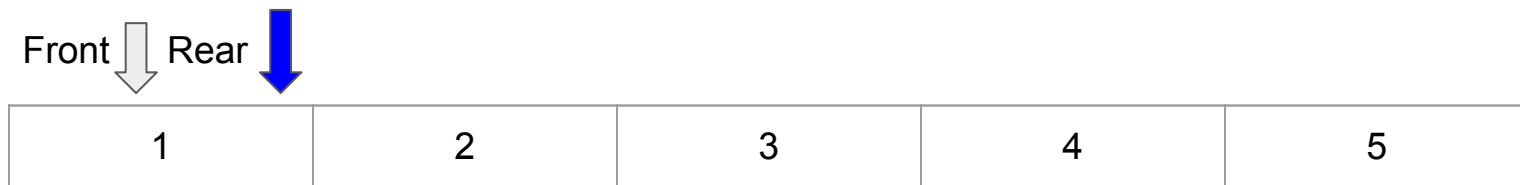
# Let's Simulate: Queue of size 5

- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4)
- Observe: rear is now in the last array position
  - One more enqueue and it moves the index to 0



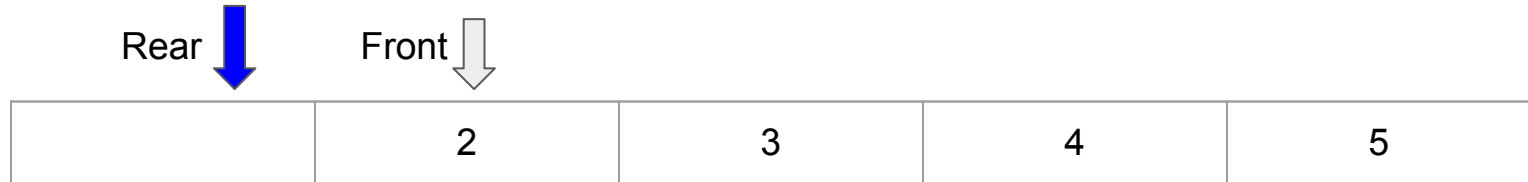
# Let's Simulate: Queue of size 5

- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Observe: the queue is full, but  $\text{rear} == \text{front}$ 
  - How can we know whether the array is empty or full?
  - We **can't!**
  - Use the added\_elements variable
    - 0 = empty
    - 5 = full



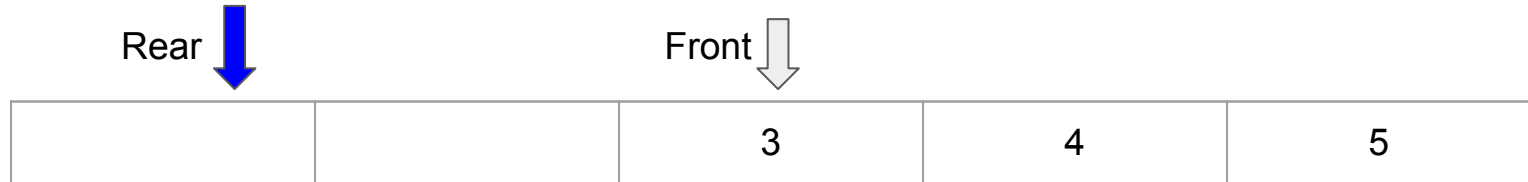
# Let's Simulate: Queue of size 5

- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue  $\Rightarrow$  1
- Observe: Front after Rear



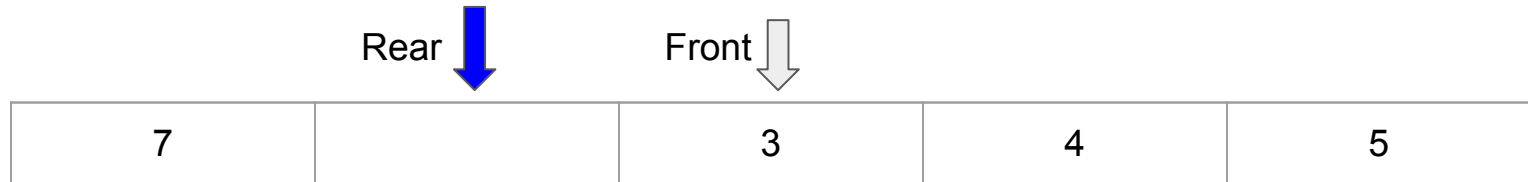
# Let's Simulate: Queue of size 5

- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue  $\Rightarrow$  1
- Dequeue  $\Rightarrow$  2



# Let's Simulate: Queue of size 5

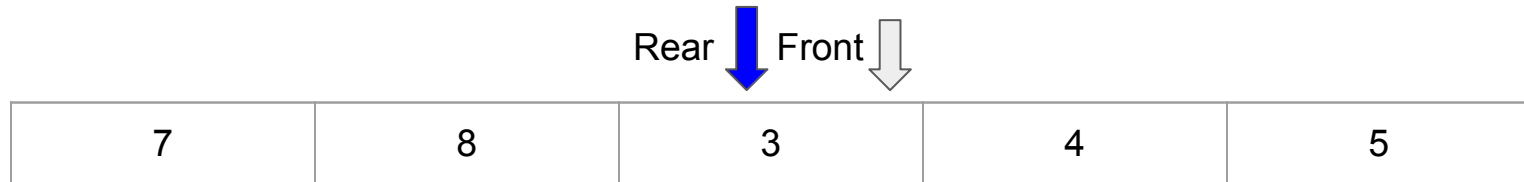
- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue  $\Rightarrow$  1
- Dequeue  $\Rightarrow$  2
- Enqueue (7)





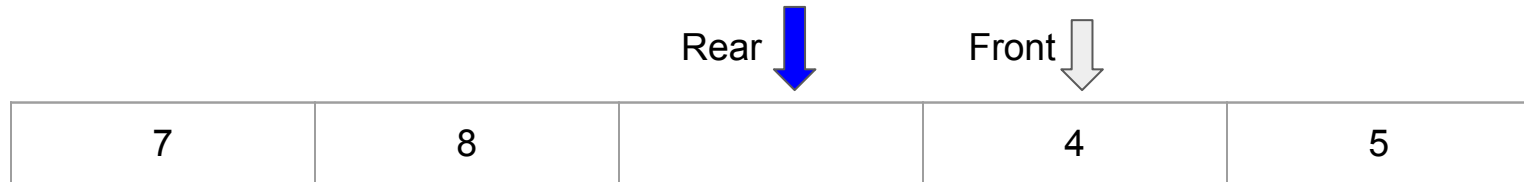
# Let's Simulate: Queue of size 5

- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue  $\Rightarrow$  1
- Dequeue  $\Rightarrow$  2
- Enqueue (7), Enqueue (8)
- Again full but rear = front = index 2



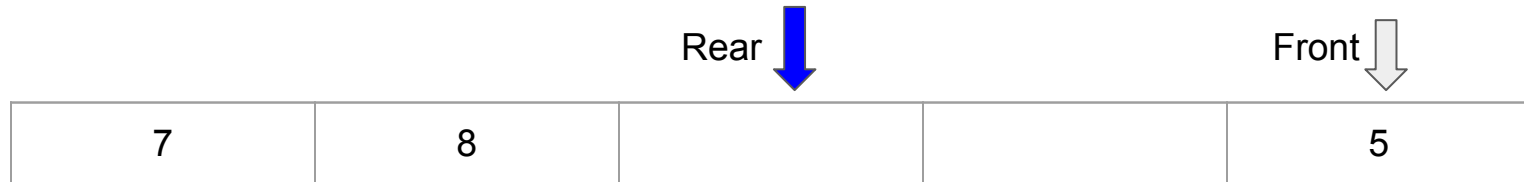
# Let's Simulate: Queue of size 5

- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue  $\Rightarrow$  1
- Dequeue  $\Rightarrow$  2
- Enqueue (7), Enqueue (8)
- Dequeue  $\Rightarrow$  3



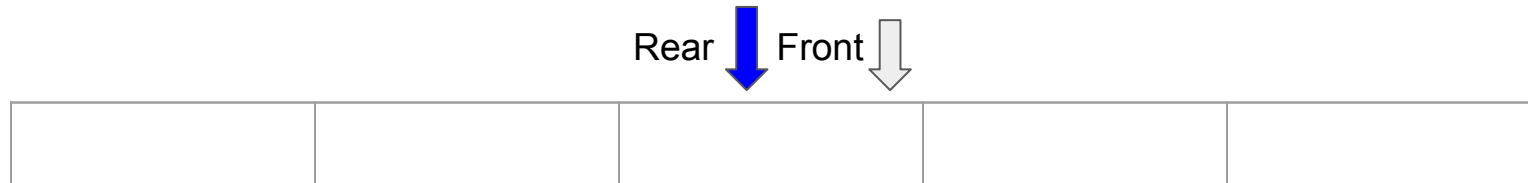
# Let's Simulate: Queue of size 5

- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue  $\Rightarrow$  1
- Dequeue  $\Rightarrow$  2
- Enqueue (7), Enqueue (8)
- Dequeue  $\Rightarrow$  3
- Dequeue  $\Rightarrow$  4



# Let's Simulate: Queue of size 5

- Enqueue (1), Enqueue (2), Enqueue (3), Enqueue (4), Enqueue (5)
- Dequeue  $\Rightarrow$  1
- Dequeue  $\Rightarrow$  2
- Enqueue (7), Enqueue (8)
- Dequeue  $\Rightarrow$  3
- Dequeue  $\Rightarrow$  4
- Dequeue, Dequeue, Dequeue  $\Rightarrow$  5, 7, 8
  - Observe: empty with front = rear = 2



*“Acquire knowledge and impart it to the people.”*

*“Seek knowledge from the Cradle to the Grave.”*