

## **Project 2**

### **Simple Key-Value Storage**



#### **Team**

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**A. Components designed:**

Our group has in general has made the following design components:

1. **Content Provider** – It provides all the functionality such as Chord Creation and Node Join, Insertion, Deletion and Querying the local storage. The Content Provider abstracts the entire functionality from the activity layer above by making available the interfacing methods like insert and query.
2. **Chord Join Function** – This component provides the logic to introduce a new node to the chord and insert it in the correct position. This also takes care of boundary conditions for zero crossings and single and two-node chords.
3. **Lookup Function** – This component provides logic to find the position where the node can be inserted into the chord. Each new message position is verified individually and on receiving the node position and insertion or query is made.

**B. Algorithm Description****a. Chord Join Component**

5554 is the first emulator to be initialized. All joining nodes first send a join request to the 5554. 5554 checks whether the node comes between itself and successor, each node follows the pseudo code mentioned above. Once the position to insert the node is found, a message is sent to the joiner to update its successor and predecessor. And another message is sent to the node's successor to update their predecessor to the joiner's node value.

**b. Lookup for Insert and Query**

The lookup function is called by the Content Provider before every Insert and Query operation. The input to the lookup function is the key and the port number of the initiator node. With every new message, the key is identified and the lookup function is called. The lookup function first checks if the key is stored with the node that has called it.

If yes, the lookup function makes the function calls to actually insert or retrieve data from the data storage. In a case where the lookup determines cannot belong to the current node, it passes it on to the next node by means of a socket connection.

C. Pseudo Algorithm:

Chord Join Component	DHT Insert
<p>Every node runs the algorithm mentioned below until the place where to insert the node is found. Here curr = current, succ = successor, joiner, pred = predecessor all are Port numbers.</p> <p><b>if</b> hash(curr) &lt; hash(succ) <b>then</b>              <b>if</b> hash(joiner) &gt; hash(curr) &amp;&amp; hash(joiner) &lt; hash(succ) <b>then</b>                  // position to insert node is found                  joiner node's succ = curr node's succ                  joiner node's pred = curr node                  curr node's successor's pred = joiner                  current node's successor = joiner              <b>else</b>                  go to next node and check same condition              <b>else</b>                  <b>if</b> hash(joiner) &gt; hash(curr)    hash(joiner) &lt; hash(succ) <b>then</b>                      // position to insert node is found                      joiner node's succ = curr node's succ                      joiner node's pred = curr node                      curr node's succ's pred = joiner                      curr node's succ = joiner</p>	<p>Returns true if node has to insert key and false if key has to be stored in another node</p> <p><b>if</b> there is only one node in the network <b>then</b>              return true</p> <p><b>if</b> myHash &gt; predecessorHash <b>then</b>              <b>if</b> keyHash &gt; predecessorHash &amp;&amp; keyHash &lt;= myHash <b>then</b>                  return true              <b>else</b>                  forward to successor                  return false</p> <p><b>else</b>              <b>if</b> keyHash &gt; predecessorHash    keyHash &lt;= myHash <b>then</b>                  return true              <b>else</b>                  forward to successor                  return false</p>

## DHT Query

