# Heart of the SwarmKit: Topology Management

Docker Distributed Systems Summit 10.07.2016

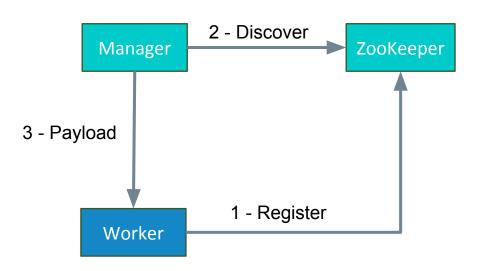


## Push vs Pull Model

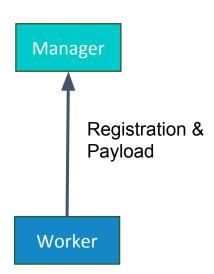


### Push vs Pull

#### **Push**



#### **Pull**





### Push vs Pull

#### Push

- Pros: Provides better control over communication rate
  - Managers decide when to contact Workers

- Cons: Requires a discovery mechanism
  - More failure scenarios
  - Harder to troubleshoot

#### Pull

- Pros: Simpler to operate
  - Workers connect to Managers and don't need to bind
  - Can easily traverse networks
  - Easier to secure
  - Less moving parts
- Cons: Workers must maintain connection to Managers at all times



### Push vs Pull

- SwarmKit adopted the Pull model
- Favored operational simplicity
- Engineered solutions to provide rate control in pull mode

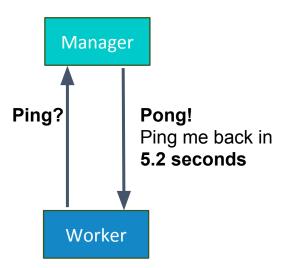


## Rate Control

Controlling communication rate in a Pull model



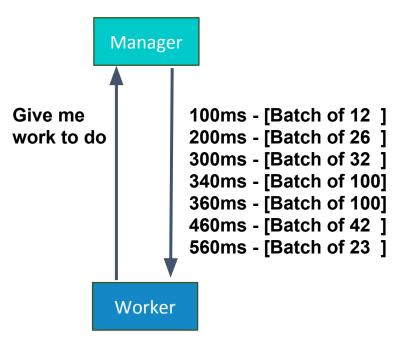
### Rate Control: Heartbeats



- Manager dictates heartbeat rate to Workers
- Rate is Configurable
- Managers agree on same Rate by Consensus (Raft)
- Managers add jitter so pings are spread over time (avoid bursts)



### Rate Control: Workloads



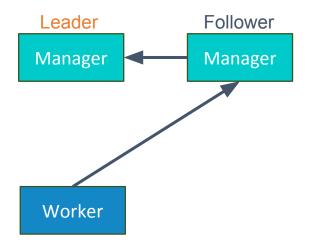
- Worker opens a gRPC stream to receive workloads
- Manager can send data whenever it wants to
- Manager will send data in batches
- Changes are buffered and sent in batches of 100 or every 100 ms, whichever occurs first
- Adds little delay (at most 100ms) but drastically reduces amount of communication



Running multiple managers for high availability

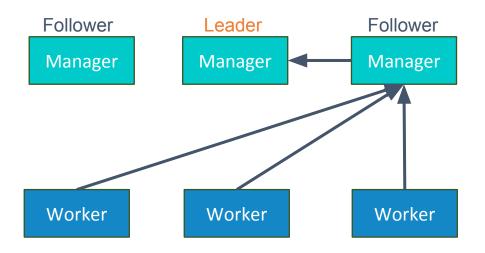


Follower Manager



- Worker can connect to any Manager
- Followers will forward traffic to the Leader

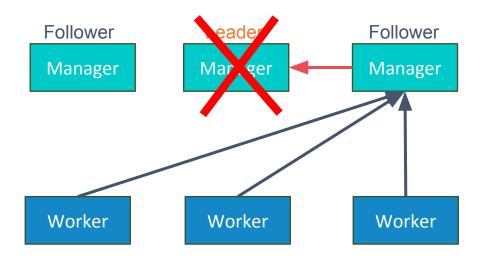




Example: On a cluster with 10,000 workers and 5 managers, each will only have to handle about 2,000 connections. Each follower will forward its 2,000 workers using a single socket to the leader.

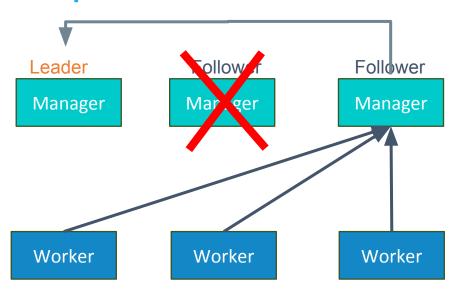
- Followers multiplex all workers to the Leader using a single connection
- Backed by gRPC channels (HTTP/2 streams)
- Reduces Leader networking load by spreading the connections evenly





- Upon Leader failure, a new one is elected
- All managers start redirecting worker traffic to the new one
- Transparent to workers

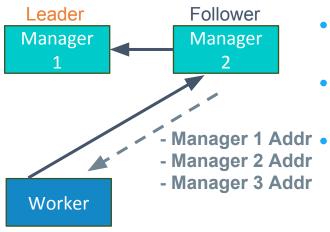




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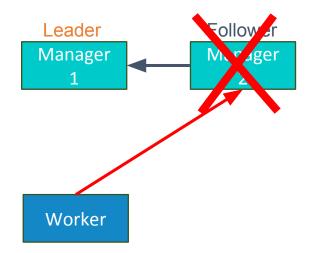
Follower Manager 3



- Manager sends list of all managers' addresses to Workers
- When a new manager joins, all workers are notified
  - Upon manager failure, workers will reconnect to a different manager

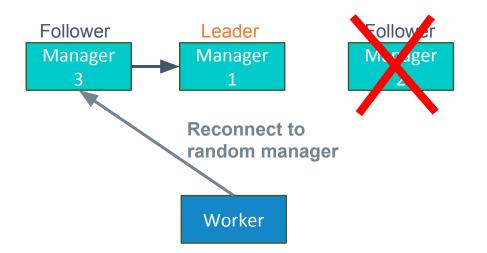


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- gRPC handles connection management
  - Exponential backoff, reconnection jitter, ...
  - Avoids flooding managers on failover
  - Connections evenly spread across Managers
- Manager Weights
  - Allows Manager prioritization / de-prioritization
  - Gracefully remove Manager from rotation



## Presence

Scalable presence in a distributed environment

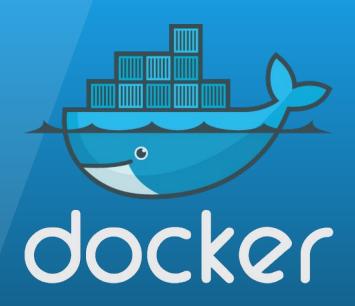


#### Presence

- Leader commits Worker state (Up vs Down) into Raft
  - Propagates to all managers
  - Recoverable in case of leader re-election.

- Heartbeat TTLs kept in Leader memory
  - Too expensive to store "last ping time" in Raft
    - Every ping would result in a quorum write
  - Leader keeps worker<->TTL in a heap (time.AfterFunc)
  - Upon leader failover workers are given a grace period to reconnect
    - Workers considered **Unknown** until they reconnect
    - If they do they move back to Up
    - If they don't they move to Down





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