# Algorithms with Predictions: learning challenges

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# a.k.a. Learning-Augmented Algorithms

- Lykouris and Vassilvitskii (ICML'18, JACM)
- · Kraska, Beutel, Chi, Dean, Polyzotis (SIGMOD'18)

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- What are the "typical" inputs?
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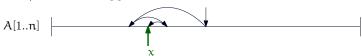


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- What are the "typical" inputs?
  - stochastic properties, specific structure/patterns
- Chapter 30 on ALPS [Mitzenmacher, Vassilvitskii] 1/13

# Classical Binary Search:

- Sorted array A
- Given x, find i s.t. A[i] = x



- Sorted array A
- Given x and untrusted prediction p(x), find i s.t. A[i] = x



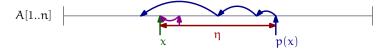
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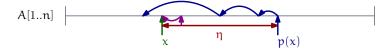


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### Binary Search with Predictions [Kraska et al. '18]:

- Sorted array A
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#### **Properties**

- consistency: O(1) steps with perfect predictions
- smoothness:  $O(\log \eta)$  steps with error  $\eta$
- robustness: never worse than  $O(\log n)$  steps  $(\eta \le n)$

# Pick up an algorithmic problem

- · online: caching, scheduling, routing, ...
- · offline: matching, clustering, flows, ...

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- primal solution, dual solution, centers of clusters, ...

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### Design algorithms

- for learning the predictions with small error (efficient)
- for using the predictions (consistent, smooth, robust)

# Outline

- (1) Learning the predictions
  - · PAC learning and beyond
- (2) Achieving smoothness
  - · how to detect small errors?
- (3) Adapting algorithmic strategy online
  - learning from the best algorithm

# (1) Learning the predictions

# PAC learning paradigm

- $I_1, \ldots, I_k \sim \mathcal{D}$
- find prediction  $\hat{p}$  s.t.

$$\mathbb{E}_{I \sim \mathcal{D}}[\eta(\boldsymbol{\hat{p}}, I)] \leqslant \min_{\boldsymbol{p}} \mathbb{E}_{I \sim \mathcal{D}}[\eta(\boldsymbol{p}, I)] + \varepsilon.$$

· related to Data-Driven Algorithm Design [Balcan et al.]

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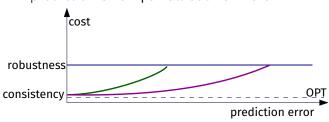
#### We want:

- · k polynomial
- polytime algorithm for minimizing empirical error
  - often easy: scheduling [Lattanzi et al.'20], matching [Dinitz et al.'21]
  - but not always: shortest path [Lattanzi et al.'23]

# (2) Smoothness

# Typical situation

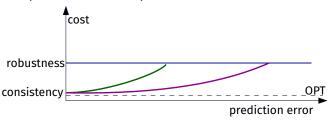
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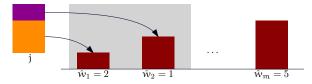
# How to deal with imprecise predictions?

- $\cdot$   $\hat{p}$  contains useful information but a few errors
- we need to find the errors (and correct them)

# (2) Smoothness: Example

# Online Load Balancing (Restricted Assignment) [Moseley et al. '20]

- · predictions: machine weights  $\hat{w}_1, \ldots, \hat{w}_m$ ,
  - $\cdot$  jobs assigned to machines proportionally to the weights

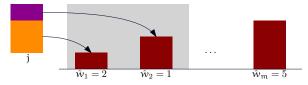


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$$\eta = \max_i \frac{\hat{w}_i}{w_i}$$

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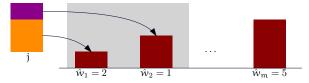
#### How to fix incorrect weights?

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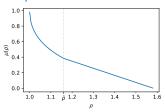
# Smoothness bound [Moseley et al. '20]

There is an algorithm with competitive ratio  $O(\log \eta)$  for the fractional restricted assignment.

# (2) Smoothness: Estimating $\eta$ online

# Online Dynamic Power Management

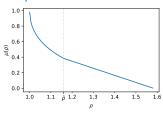
•  $\mathsf{ALG}_{\rho}\leqslant \rho\,\mathsf{OPT}\,+\mu(\rho)\eta\,$  [Antoniadis, Coester, Eliáš, Polak, Simon '21]



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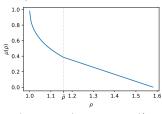


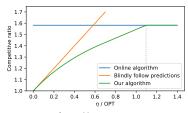
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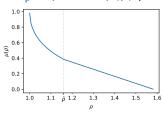


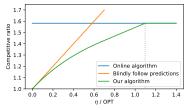
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# Learning question

How much can we trust the prediction?

# Algorithmic questions

How important is the current decision? What is the alternative to following the prediction?

# (3) Adapting strategy online

# Two scenarios in online setting

- algorithm A uses predictor trained on distribution  ${\cal D}$
- $I \sim \mathcal{D} \Rightarrow A$  works great
- I  $\sim \mathcal{D} \Rightarrow$  some worst-case A' still works ok
- I arriving online: we want cost  $\approx \min\{A(I), A'(I)\}$

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### Multiple scenarios in online setting

- ·  $A_i$  uses predictor trained on  $\mathcal{D}_i$ ,  $i=1,\ldots,\ell$
- $\cdot$  I is arriving online, we do not know what  $\mathcal{D}_i$  it comes from
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#### Challenge:

- ullet past actions o current state of the algorithm
- if we take an action of a bad algorithm, we cannot take it back

9/13

#### Related works

### Online choice of online algorithms

- · k-server [Fiat, Rabani, Ravid 1990]
- · caching [Fiat, Karp, Luby, McGeoch, Sleator, Young 1991]
- · general MTS [Azar, Broder, Manasse 1993]

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#### Algorithms with multiple predictions

- Online Covering [Anand et al. 2022], [Kevi, Nguyen 2023]
- Offline Matching, Scheduling [Dinitz et al. 2022]
- Ski Rental [Gollapudi, Panigrahi 2019], [Wang et al. 2020]
- Online Facility Location [Almanza et al. 2021]
- offline algorithms [Srinivas, Blum 2024]

#### **MTS**

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#### **Full information**

- ·  $min{A_1, ..., A_\ell} + Regret$  [Blum, Burch '00]
- $\cdot$  O( $\ell^2$ ) dyn.comb( $A_1,\ldots,A_\ell$ ) [Antoniadis, Coester, Eliáš, Polak, Simon '23]
  - · Layered Graph Traversal [Bubeck, Coester, Rabani '22]

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#### **Bandit setting**

- $min{A_1, ..., A_\ell} + Regret$  [Cosa, Eliáš '25]
- · Dynamic combination: open

# Further perspectives

# Improve the current techniques to:

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### End-to-end learning of online algorithms

- Algorithms with Explicit Predictors [Eliáš, Kaplan, Mansour, Moran '23]
- algorithm has direct access to the dataset of past inputs
- · algorithm learns while processing the input

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# Offline and Approximation algorithms

- currently understudied
- many techniques do not translate to offline setting

#### **Further information**

#### **Posters**

- Xizhi Tan: Learning-Augmented Mechanism Design
- · Kaito Fujii: The Secretary Problem with Predictions

# https://algorithms-with-predictions.github.io/

- by Alexander Lindermayr
- database of the papers in the area

