

## Homework #1 (Due date: 2/21/2025)

1. Below is an algorithm for Bounded Buffer problem. Fill in the table of the values of semaphores. (2.5 points)

Shared data

semaphore full, empty, mutex;

Initially:

full = 0, empty = n, mutex = 1

Producer process:

do {

...

produce an item in nextp

...

wait(empty);

wait(mutex);

...

add nextp to buffer

...

signal(mutex);

signal(full);

} while (1);

Consumer process:

do {

wait(full)

wait(mutex);

...

remove an item from buffer to nextc

...

signal(mutex);

signal(empty);

...

consume the item in nextc

...

} while (1);

events	mutex	full	empty
produce 1 item	1 → 0 → 1	0 → 1	n → n-1
produce 1 item	1 → 0 → 1	1 → 2	n-1 → n-2
produce 1 item	1 → 0 → 1	2 → 3	n-2 → n-3
consume 1 item	1 → 0 → 1	3 → 2	n-3 → n-2
consume 1 item	1 → 0 → 1	2 → 1	n-2 → n-1
produce 1 item	1 → 0 → 1	1 → 2	n-1 → n-2
consume 1 item	1 → 0 → 1	2 → 1	n-2 → n-1
consume 1 item	1 → 0 → 1	1 → 0	n-1 → n

For  $n=3$ :

events	mutex	full	empty
produce 1 item	$1 \rightarrow 0 \rightarrow 1$	$0 \rightarrow 1$	$3 \rightarrow 2$
produce 1 item	$1 \rightarrow 0 \rightarrow 1$	$1 \rightarrow 2$	$2 \rightarrow 1$
produce 1 item	$1 \rightarrow 0 \rightarrow 1$	$2 \rightarrow 3$	$1 \rightarrow 0$
consume 1 item	$1 \rightarrow 0 \rightarrow 1$	$3 \rightarrow 2$	$0 \rightarrow 1$
consume 1 item	$1 \rightarrow 0 \rightarrow 1$	$2 \rightarrow 1$	$1 \rightarrow 2$
produce 1 item	$1 \rightarrow 0 \rightarrow 1$	$1 \rightarrow 2$	$2 \rightarrow 1$
consume 1 item	$1 \rightarrow 0 \rightarrow 1$	$2 \rightarrow 1$	$1 \rightarrow 2$
consume 1 item	$1 \rightarrow 0 \rightarrow 1$	$1 \rightarrow 0$	$2 \rightarrow 3$

1 → 0	1 → 0	1 → 0
2 → 1	2 → 1	2 → 1
3 → 2	3 → 2	3 → 2
4 → 3	4 → 3	4 → 3
5 → 4	5 → 4	5 → 4
6 → 5	6 → 5	6 → 5
7 → 6	7 → 6	7 → 6
8 → 7	8 → 7	8 → 7
9 → 8	9 → 8	9 → 8
10 → 9	10 → 9	10 → 9

2. Below is an algorithm for Readers-Writers problem. Fill in the table of the values of semaphores or variables. (2.5 points)

Shared data

semaphore mutex, db;

Initially

mutex = 1, db = 1, readcount = 0

Writer process:

wait(db);

...

writing is performed

...

signal(db);

Reader process:

wait(mutex);

readcount++;

if (readcount == 1)

wait(db);

signal(mutex);

...

reading is performed

...

wait(mutex);

readcount--;

if (readcount == 0)

signal(db);

signal(mutex);

event	mutex = 1	db = 1	readcount = 0
a writer comes	1	1 → 0	0
a reader comes	1 → 0	0 → -1	0 → 1
a writer has done	0 → 1	-1 → 0	1
a reader comes	1 → 0 → 1	0	1 → 2
a reader comes	1 → 0 → 1	0	2 → 3
a writer comes	1	0 → -1	3
a reader has done	1 → 0 → 1	-1	3 → 2
a reader has done	1 → 0 → 1	-1	2 → 1
a reader has done	1 → 0 → 1	-1 → 0	1 → 0
a writer comes	1	0 → -1	0



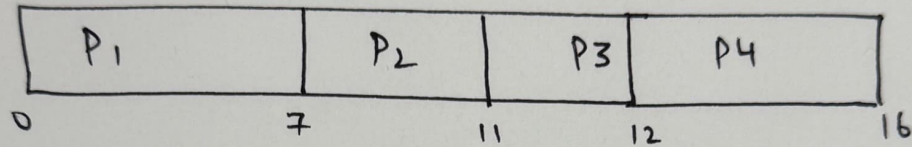
3. Consider the following set of processes, with length of the Burst Time given in milliseconds:

(5 points)

Process	Arrival Time	Burst Time	Priority
P <sub>1</sub>	0.0	7	2
P <sub>2</sub>	2.0	4	1
P <sub>3</sub>	4.0	1	3
P <sub>4</sub>	5.0	4	2

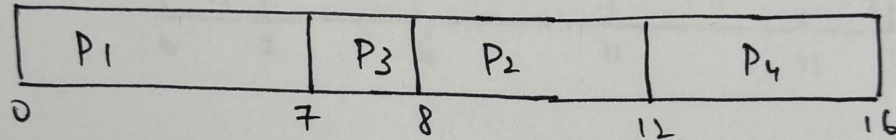
1) Draw Gantt charts **with ready queues** if applicable, that using FCFS, SJF, Priority (a smaller priority number implies a higher priority), and RR (quantum = 1), with non-preemptive for all the algorithms applicable.

FCFS :-



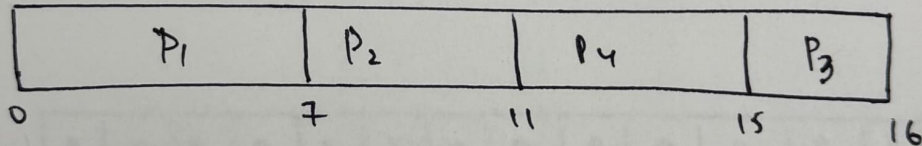
~~7~~ 0  
~~4~~ 0  
~~1~~ 0  
~~4~~ 0

SJF :-



~~7~~ 0  
~~4~~ 0  
~~1~~ 0  
~~4~~ 0

Priority :-



~~7~~ 0  
~~4~~ 0  
~~1~~ 0  
~~4~~ 0

Round Robin :

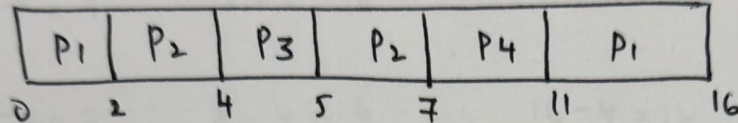
→ RR is an example for pre-emptive algorithm.  
Here ignored !!  
→ As it is always pre-emptive.

2) Repeat 1) with preemptive for the applicable algorithms only.

FCFS :-  $\longrightarrow$  is always non-pre-emptive  
 $\longrightarrow$  Thus ignored Here !!

~~X~~ 0  
~~X~~ 0  
~~X~~ 0  
~~X~~ 0

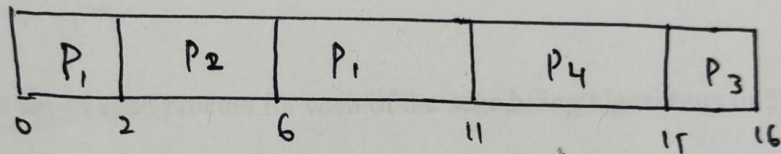
SJF Preemptive :-



uses  
 AT & BT

~~X~~ 0  
~~X~~ 0  
~~X~~ 0  
~~X~~ 0

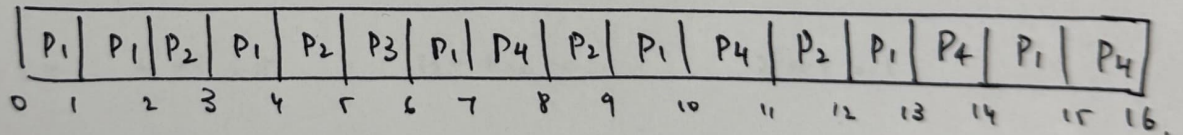
Priority Preemptive :



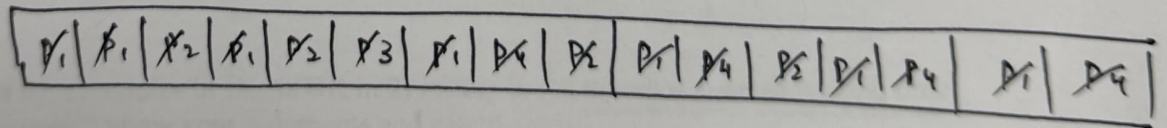
uses  
 - Priority  
 - AT

Round Robin (RR) :

Gantt chart :



Ready Queue :





3) What is the Turnaround time of each process for each of the scheduling algorithms in 1) and RR in 2)?

	<u>FCFS</u>	<u>SJF</u>	<u>Priority</u>	<u>RR</u>
P1	$7 - 0 = 7$	$7 - 0 = 7$	$7 - 0 = 7$	$15 - 0 = 15$
P2	$11 - 2 = 9$	$12 - 2 = 10$	$11 - 2 = 9$	$12 - 2 = 10$
P3	$12 - 4 = 8$	$8 - 4 = 4$	$16 - 4 = 12$	$6 - 4 = 2$
P4	$16 - 5 = 11$	$16 - 5 = 11$	$15 - 5 = 10$	$16 - 5 = 11$

4) What is the Waiting time of each process for each of the scheduling algorithms in 1) and RR in 2).

	<u>FCFS</u>	<u>SJF</u>	<u>Priority</u>	<u>RR</u>
P1	$7 - 7 = 0$	$7 - 7 = 0$	$7 - 7 = 0$	$15 - 7 = 8$
P2	$9 - 4 = 5$	$10 - 4 = 6$	$9 - 4 = 5$	$10 - 4 = 6$
P3	$8 - 1 = 7$	$4 - 1 = 3$	$12 - 1 = 11$	$2 - 1 = 1$
P4	$11 - 4 = 7$	$11 - 4 = 7$	$10 - 4 = 6$	$11 - 4 = 7$

5) Which of the schedules in 1) and RR in 2) results in the minimal average waiting time over all processes? Show your judgments and calculations.

	<u>FCFS</u>	<u>SJF</u>	<u>Priority</u>	<u>RR</u>
Avg wait Time	$\frac{0+5+7+7}{4}$ $= 4.7\text{ms}$	$\frac{0+6+3+7}{4}$ $= 4\text{ms}$	$\frac{0+5+11+6}{4}$ $= 5.5\text{ms}$	$\frac{8+6+1+7}{4}$ $= 5.5\text{ms}$

Conclusion / Judgement : SJF has the lowest average waiting time.