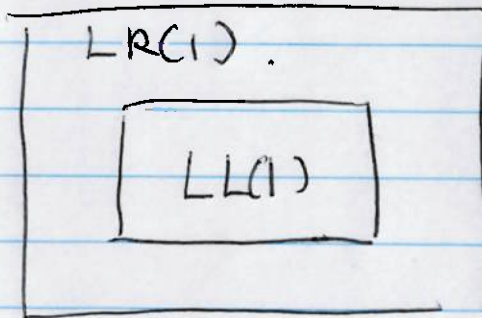


04/07/2025

3.4. Bottom up parsing : build the parse tree from leaves to the root

↓ right most derivation in reverse
 $LR(1)$
 ↑ look ahead. one token
 ↑ left to right scan



$LR(1)$ can recognize a larger set of grammars / languages

$$L = \{a^i b^j, \quad i \geq j \geq 1\}$$

$$S \rightarrow \overbrace{a s b} \\ | \underline{a b} \\ | T S$$

$$T \rightarrow \overbrace{a T} \\ | \underline{\epsilon}$$

$L \notin LL(k)$
 for any k

$L \in LR(1)$

95

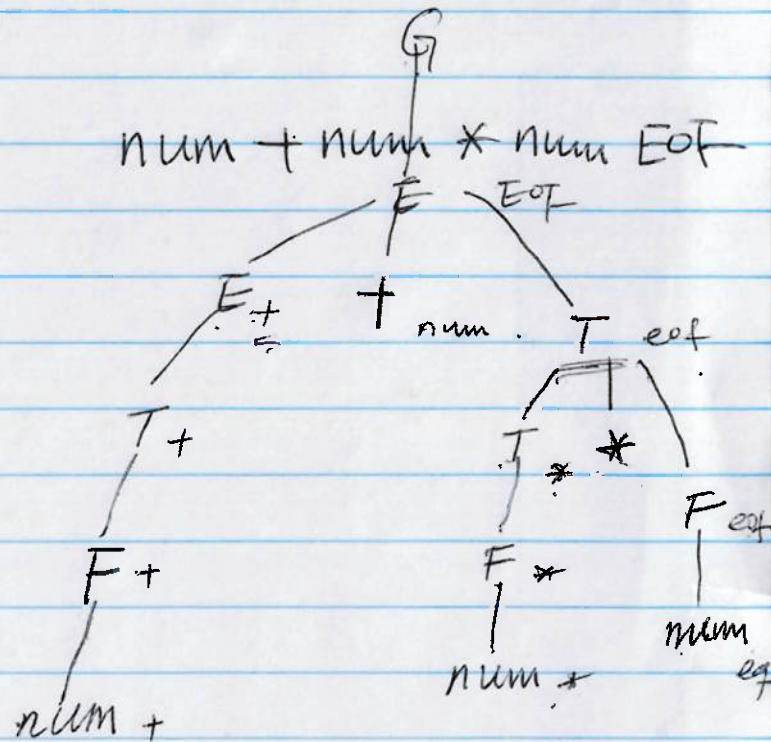
②. build ^a reversed right most derivation tree

③. at each step. the parser either

read/scan next token / shift

~~reduction~~
or replace rhs of a prod. with its lhs.

ex: $G \rightarrow E$
 $E \rightarrow \underline{E} + T$
 $E \rightarrow T$
 $T \rightarrow \underline{T} * F$
 $F \rightarrow F / F$
 F
 $F \rightarrow \text{num}$
 name
 (E)



shift:

reduction.

$$A \rightarrow B_1 B_2 \dots B_k$$

handle

Q. How to find a handle?

ex: ① $Goal \rightarrow IB+$

② $IB+ \rightarrow IB+$ pair

③ I pair

④ $pair \rightarrow (IB+)$

⑤ $I ()$

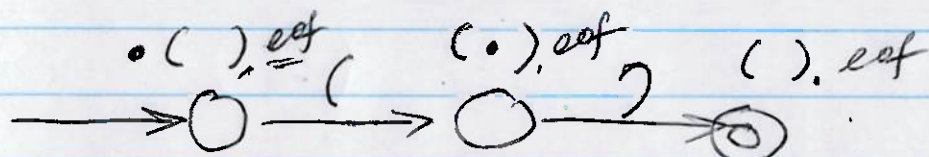
input $X) EOF$

$\times (() EOF$

$\checkmark () () EOF$

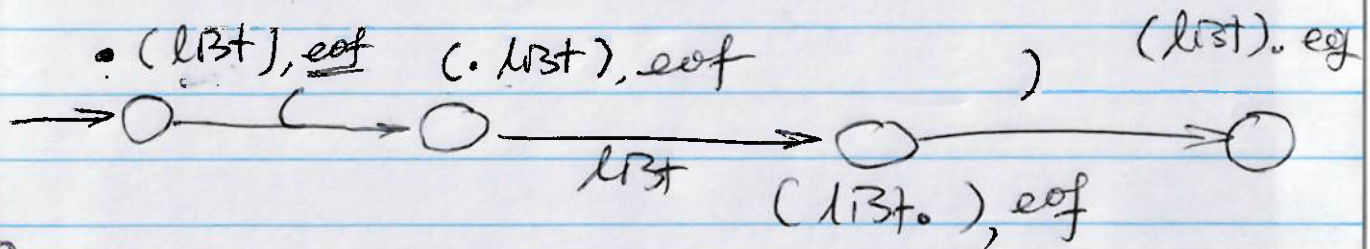
$\checkmark (()) EOF$

⑤. $pair \rightarrow ()$

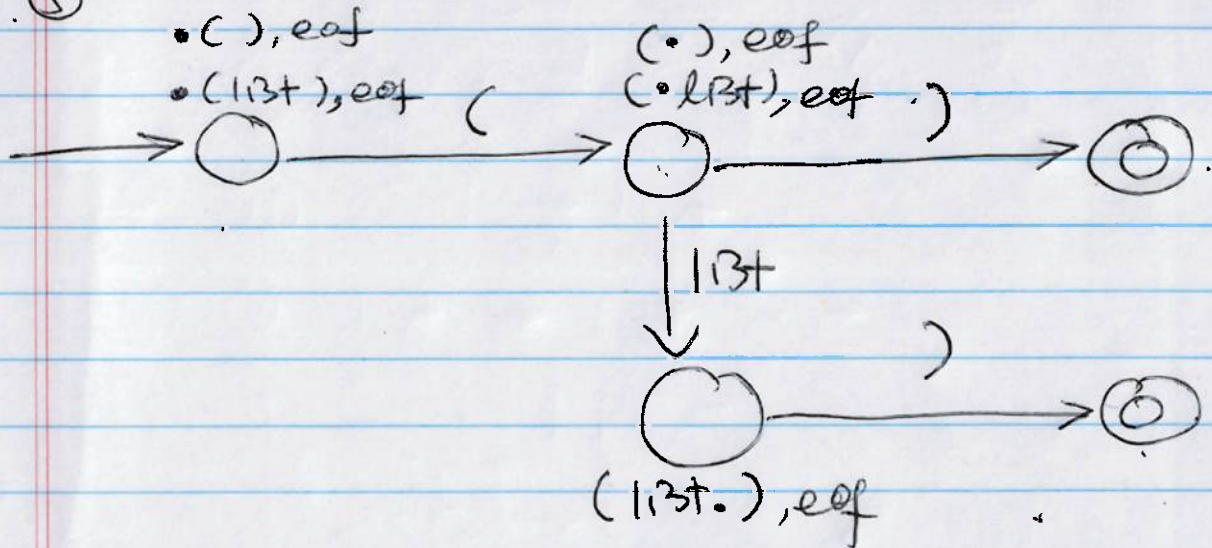


(97)

pair $\rightarrow (list)$



(4) (5)



build DFA

encode DFA into 2 tables

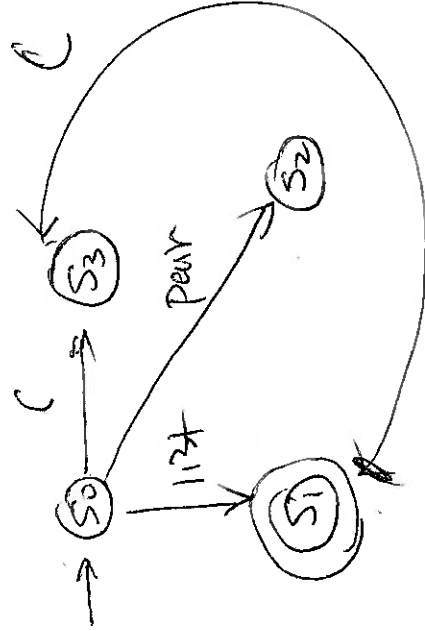
Action table

State	Terminal	EOF
s_0		
s_1		
\vdots		

Action (state, $a \in T$) = $\begin{cases} \text{accept} \\ s_{\text{next state}} \\ \text{rule \#} \\ \text{error/blank} \end{cases}$

LR(1) Tables for Parenthesis Grammar

1	Goal	→	List
2	List	→	List Pair
3			Pair
4	Pair	→	{ List }
5			{ (}



State	ACTION Table				Goto Table	
	EOF	{	}		List	Pair
S ₀		s3			1	2
S ₁	Acc	s3				4
S ₂	r3	r3				
S ₃		s7	s8		5	6
S ₄	r2	r2				
S ₅		s7	s10			9
S ₆		r3	r3			
S ₇		s7	s12		11	6
S ₈	r5	r5				
S ₉		r2	r2			
S ₁₀	r4	r4				
S ₁₁		s7	s13			9
S ₁₂		r5	r5			
S ₁₃		r4	r4			

98

Goto table

Goto(state, A ∈ NT)
= next state

state	NT
s ₀	
s ₁	
s ₂	
⋮	

LR(1) skeleton parser . table driven

// make a stack to keep the prefix
of the upper frontier of the parse tree

push (\$, \$)

push (start symbol, s₀)

word ← next word ()

while (true)

(symbol, state) ← top () of the stack

switch (Action (state, word))

case "accept" word = eof

then break

case s_i

push (word, s_i)

word ← next word ()

case $r \neq$. $(A \rightarrow \underline{B_1 B_2 \dots B_k})$

pop k times

$(symbol, \underline{state}) \leftarrow top()$

push(A , Goto($state$, A))

otherwise

report error

example: () eof

state	lookahead	stack	handle	Action
	($\$(Goal, 0)$		
0	(S3
3)	$\$(Goal, 0) (c, 3)$		S8
8	eof	$\$(Goal, 0) (c, 3), (c), 8)$		$\underline{r5}$ $\underline{A-c}$
2	eof	$\$(Goal, 0) (\underline{pair}, 2)$		r3
1	<u>eof</u>	$\$(Goal, 0) (\underline{list}, 1)$		Acco

())

(()

() ()

((,)

04/09/2025

100

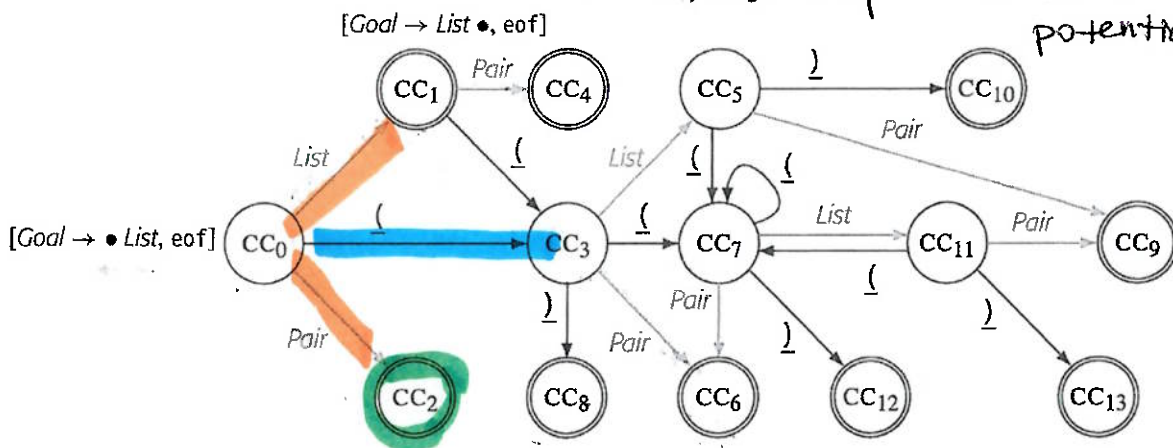
4/9/2025

		State	eof	()	List	Pair
1	Goal \rightarrow List	0		s 3		1	2
2	List \rightarrow List Pair	1	acc	s 3			4
3	Pair	2	r 3	r 3			
4	Pair \rightarrow (List)	3		s 7	s 8	5	6
5	()	4	r 2	r 2			
		5		s 7	s 10		9
		6		r 3	r 3		
		7		s 7	s 12	11	6
		8	r 5	r 5			
		9		r 2	r 2		
		10	r 4	r 4			
		11		s 7	s 13		9
		12		r 5	r 5		
		13		r 4	r 4		

2) Action and Goto Tables for Parentheses Grammar

Item. pair \rightarrow () , eof
 • () , eof
 (•) , eof

CCi : canonical collection of items a state of the parser that contains one or more items that represent a handles or potential handles



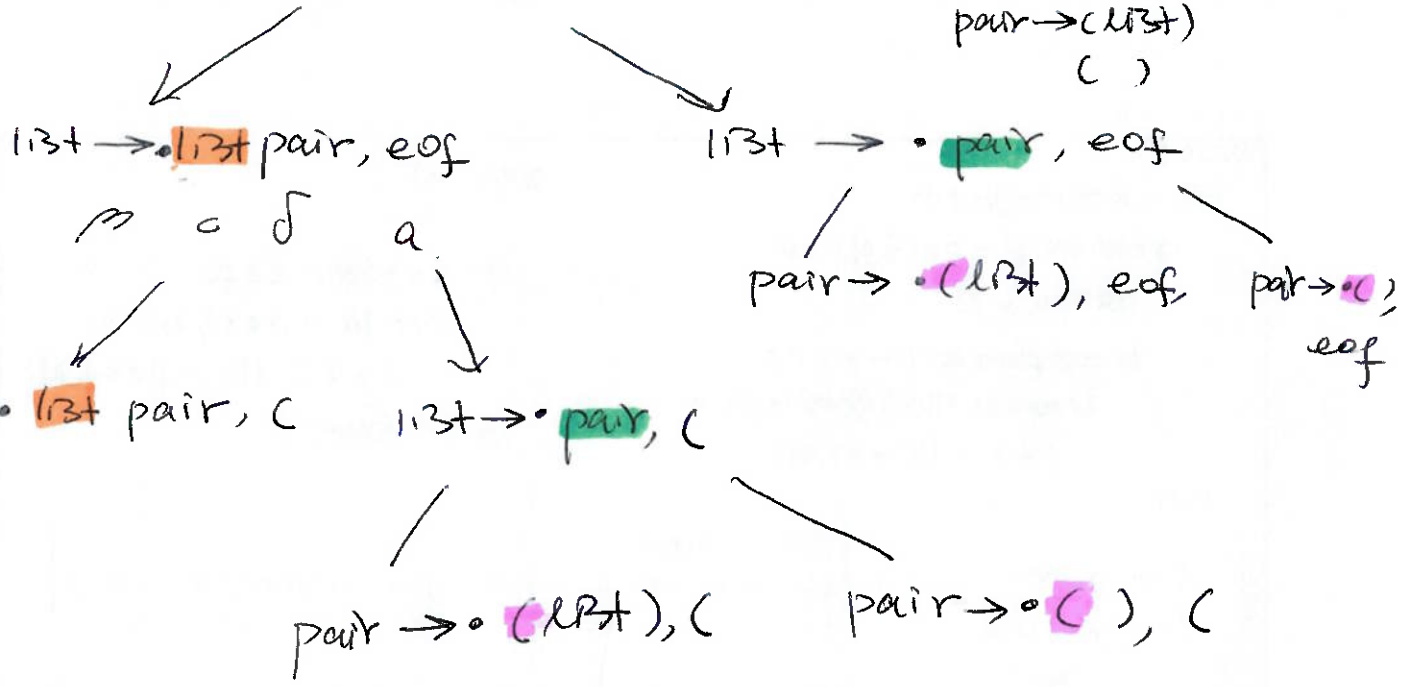
Review subset construction

initial state *: core item.

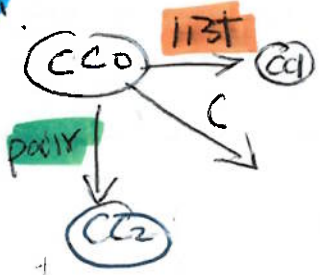
(10)

Goal \rightarrow list
 list \rightarrow list pair
 | pair
 pair \rightarrow (list)
 ()

CC0 * Goal \rightarrow • list, eof



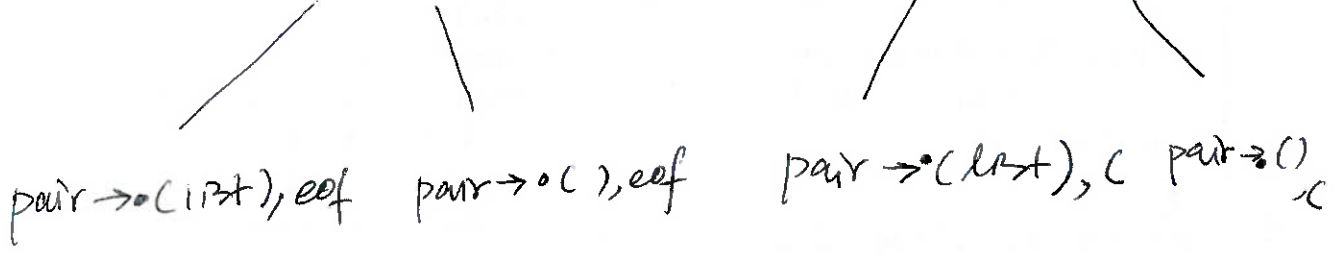
transition



CC1

* Goal \rightarrow list •, eof.

* list \rightarrow list • pair, eof * list \rightarrow list • pair, (



CC2

core * list \rightarrow pair •, (

* list \rightarrow pair; eof.

accepting state

compute a state starting core states
 β is a string
 left context

closure(s)

```

while (s is still changing) do
  for each item  $[A \rightarrow \beta \bullet C \delta, a] \in s$  do
    lookahead  $\leftarrow \delta a$ 
    for each production  $C \rightarrow \gamma \in P$  do
      for each  $b \in \text{FIRST}(\text{lookahead})$  do
         $s \leftarrow s \cup \{[C \rightarrow \bullet \gamma, b]\}$ 
return s
        
```

goto(s, x) computes a transition

```

t  $\leftarrow \emptyset$ 
for each item  $i \in s$  do
  if  $i$  is  $[\alpha \rightarrow \beta \bullet x \delta, a]$  then
     $t \leftarrow t \cup \{[\alpha \rightarrow \beta x \bullet \delta, a]\}$ 
return closure(t)
        
```

1	Goal \rightarrow List	$CC_0 = \{ [Goal \rightarrow \bullet List, eof] [List \rightarrow \bullet List Pair, eof] [List \rightarrow \bullet List Pair, _] [List \rightarrow \bullet Pair, eof] [List \rightarrow \bullet Pair, _] [Pair \rightarrow \bullet (List), eof] [Pair \rightarrow \bullet (List), _] [Pair \rightarrow \bullet (), eof] [Pair \rightarrow \bullet (), _]]$
2	List \rightarrow List Pair	
3	Pair	
4	Pair \rightarrow (List)	$CC_1 = \{ [Goal \rightarrow List \bullet, eof] [List \rightarrow List \bullet Pair, eof] [List \rightarrow List \bullet Pair, _] [Pair \rightarrow \bullet (List), eof] [Pair \rightarrow \bullet (List), _] [Pair \rightarrow \bullet (), eof] [Pair \rightarrow \bullet (), _]]$
5	()	$CC_2 = \{ [List \rightarrow Pair \bullet, eof] [List \rightarrow Pair \bullet, _]]$

if $[A \rightarrow \beta \bullet a] \in CC_i$, then CC_i

// build collection DFA

```

CC0  $\leftarrow \emptyset$ 
for each production of the form  $Goal \rightarrow \alpha$  do
   $CC_0 \leftarrow CC_0 \cup \{ [Goal \rightarrow \bullet \alpha, eof] \}$ 
CC0  $\leftarrow$  closure(CC0)
CC  $\leftarrow \{ CC_0 \}$ 
while (new sets are still being added to CC) do
  for each unmarked set  $CC_i \in CC$  do
    mark  $CC_i$  as processed
    for each  $x$  following  $a$  in an item in  $CC_i$  do
      temp  $\leftarrow$  goto( $CC_i, x$ )
      if temp  $\notin CC$  then
        CC  $\leftarrow CC \cup \{ temp \}$ 
    record transition from  $CC_i$  to temp on  $x$ 
        
```

// table filling Action Goto

```

for each  $CC_i \in CC$  do
  for each item  $l \in CC_i$  do
    ① if  $l$  is  $[A \rightarrow \beta \bullet x \delta, a]$  and goto( $CC_i, x$ ) =  $CC_j$  then
      Action[i, x]  $\leftarrow$  "shift j"
    ② else if  $l$  is  $[A \rightarrow \beta \bullet, a]$  then
      Action[i, a]  $\leftarrow$  "reduce A  $\rightarrow$   $\beta$ "
    ③ else if  $l$  is  $[Goal \rightarrow \beta \bullet, eof]$  then
      Action[i, eof]  $\leftarrow$  "accept"
  for each  $n \in NT$  do
    if goto( $CC_i, n$ ) =  $CC_j$  then
      Goto[i, n]  $\leftarrow$  j
        
```

```

graph LR
  CC0((CC0)) -- List --> CC1((CC1))
  CC0 -- Pair --> CC2((CC2))
  CC1 -- Pair --> CC2
  CC1 -- List --> CC0
  CC2 -- List --> CC1
  CC2 -- Pair --> CC0
        
```


$$\begin{aligned}
 CC_0 &= \left\{ \begin{aligned} &[Goal \rightarrow \bullet List, eof] \\ &[List \rightarrow \bullet List Pair, eof] [List \rightarrow \bullet List Pair, _] [List \rightarrow \bullet Pair, eof] [List \rightarrow \bullet Pair, _] \\ &[Pair \rightarrow \bullet _ List _, eof] [Pair \rightarrow \bullet _ List _, _] [Pair \rightarrow \bullet _ _, eof] [Pair \rightarrow \bullet _ _, _] \end{aligned} \right\} \\
 CC_1 &= \left\{ \begin{aligned} &[Goal \rightarrow List \bullet, eof] [List \rightarrow List \bullet Pair, eof] [List \rightarrow List \bullet Pair, _] \\ &[Pair \rightarrow \bullet _ List _, eof] [Pair \rightarrow \bullet _ List _, _] [Pair \rightarrow \bullet _ _, eof] [Pair \rightarrow \bullet _ _, _] \end{aligned} \right\} \\
 CC_2 &= \{ [List \rightarrow Pair \bullet, eof] [List \rightarrow Pair \bullet, _] \} \\
 CC_3 &= \left\{ \begin{aligned} &[Pair \rightarrow _ \bullet List _, eof] [Pair \rightarrow _ \bullet List _, _] [Pair \rightarrow _ \bullet _, eof] [Pair \rightarrow _ \bullet _, _] \\ &[List \rightarrow \bullet List Pair, _] [List \rightarrow \bullet List Pair, _] [List \rightarrow \bullet Pair, _] [List \rightarrow \bullet Pair, _] \\ &[Pair \rightarrow \bullet _ List _, _] [Pair \rightarrow \bullet _ List _, _] [Pair \rightarrow \bullet _ _, _] [Pair \rightarrow \bullet _ _, _] \end{aligned} \right\} \\
 CC_4 &= \{ [List \rightarrow List Pair \bullet, eof] [List \rightarrow List Pair \bullet, _] \} \\
 CC_5 &= \left\{ \begin{aligned} &[List \rightarrow List \bullet Pair, _] [List \rightarrow List \bullet Pair, _] [Pair \rightarrow _ List \bullet _, eof] [Pair \rightarrow _ List \bullet _, _] \\ &[Pair \rightarrow \bullet _ List _, _] [Pair \rightarrow \bullet _ List _, _] [Pair \rightarrow \bullet _ _, _] [Pair \rightarrow \bullet _ _, _] \end{aligned} \right\} \\
 CC_6 &= \{ [List \rightarrow Pair \bullet, _] [List \rightarrow Pair \bullet, _] \}
 \end{aligned}$$

$$\begin{aligned}
 CC_4 &= \{ [List \rightarrow List Pair \bullet, eof] [List \rightarrow List Pair \bullet, _] \} \\
 CC_7 &= \left\{ \begin{aligned} &[Pair \rightarrow _ \bullet List _, _] [Pair \rightarrow _ \bullet List _, _] [Pair \rightarrow _ \bullet _, _] [Pair \rightarrow _ \bullet _, _] \\ &[List \rightarrow \bullet List Pair, _] [List \rightarrow \bullet List Pair, _] [List \rightarrow \bullet Pair, _] [List \rightarrow \bullet Pair, _] \\ &[Pair \rightarrow \bullet _ List _, _] [Pair \rightarrow \bullet _ List _, _] [Pair \rightarrow \bullet _ _, _] [Pair \rightarrow \bullet _ _, _] \end{aligned} \right\} \\
 CC_8 &= \{ [Pair \rightarrow _ _, \bullet, eof] [Pair \rightarrow _ _, \bullet, _] \} \\
 CC_9 &= \{ [List \rightarrow List Pair \bullet, _] [List \rightarrow List Pair \bullet, _] \} \\
 CC_{10} &= \{ [Pair \rightarrow _ List _, \bullet, eof] [Pair \rightarrow _ List _, \bullet, _] \} \\
 CC_{11} &= \left\{ \begin{aligned} &[List \rightarrow List \bullet Pair, _] [List \rightarrow List \bullet Pair, _] [Pair \rightarrow _ List \bullet _, _] [Pair \rightarrow _ List \bullet _, _] \\ &[Pair \rightarrow \bullet _ List _, _] [Pair \rightarrow \bullet _ List _, _] [Pair \rightarrow \bullet _ _, _] [Pair \rightarrow \bullet _ _, _] \end{aligned} \right\} \\
 CC_{12} &= \{ [Pair \rightarrow _ _, \bullet, _] [Pair \rightarrow _ _, \bullet, _] \} \\
 CC_{13} &= \{ [Pair \rightarrow _ List _, \bullet, _] [Pair \rightarrow _ List _, \bullet, _] \}
 \end{aligned}$$

Is the grammar LR(1)?

reduce reduce conflict
reduce-shift conflict

when filling
the table