

# \* Recursion - Quick Sort

5, 4, 3, 2, 1

pivot :

Choose any element  
 → after first pass all elements  
 $< p$  will be on LHS of  $p$   
 &

→  $> p$  will be at RHS of  $p$

1, 3, 2, (4), 5

pivot correct position

1, 3, 2, (4), 5  
 1, (3), 2, (4), 5  
 (1), 2, (3), (4), 5

recursion

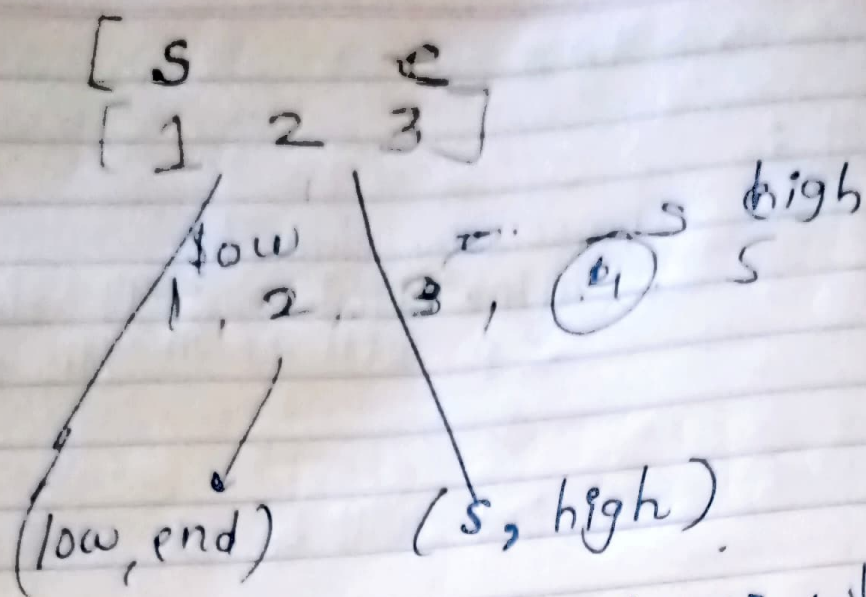
1, 2, 3, 4, 5

\* How to put pivot at correct position?

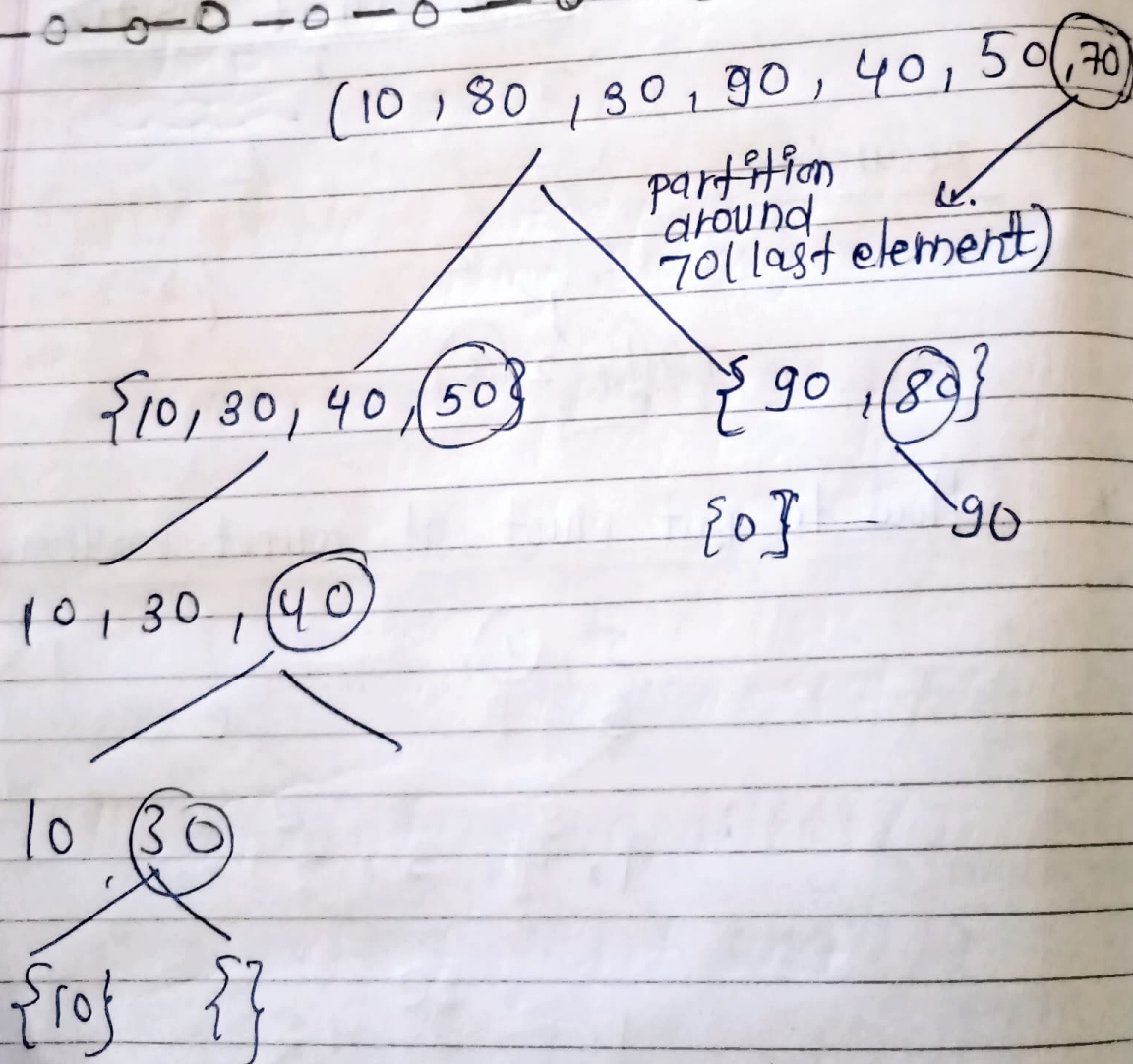
5, (4), 3, 2, 1       $P=4$   
 $s$        $p$        $e$

1,  $s$  4, 3,  $e$  2, 5

1, 2,  $s$  3,  $e$  4, 5  
 $s$        $p$        $e$   
 [ 1 2 3 ]



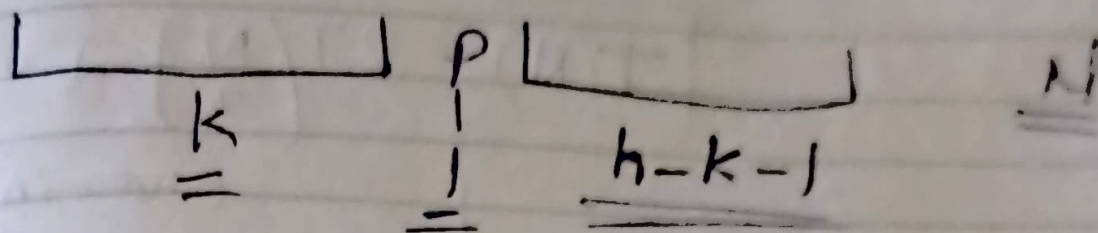
low, high  $\rightarrow$  tells us which part of array you working on — swap.





How to pick pivot:

- \* Random element
- \* Other element
- \* pick the middle element



Recurrence relation of quick sort

$$T(N) = T(k) + T(N-k-1) + O(N)$$

Worst case (common sense) General.

$k=0$       $3, 8, 15, 20, 31, 41, (86)_p$

$h-1$

$$T(N) = T(0) + T(N-1) + O(N)$$

$$T(N) = T(N-1) + O(N)$$

worse case  
\* Complexity is  $O(N^2)$   $\Rightarrow$  complexity lec



\*

Best case :

$$k = \frac{N}{2}$$

$$T(N) = T\left(\frac{N}{2}\right) + T\left(\frac{N}{2}\right) + o(1)$$

$$T(N) = 2T\left(\frac{N}{2}\right) + o(N)$$

~~covered in d~~

Best case =  $O(N \log N)$

Imp

- 1) Not stable: (already covered in initial sorting videos)
- 2) In place: why preferred for arrays instead of ms. ms takes  $O(N)$  extra space.
- 3) ms is better in linked lists due to memory allocation  $\rightarrow$  not continuous

# Hybrid sorting algorithms (Tim Sort, merge sort + insert sort)