

NFA = DFA

Recall: A nondeterministic finite automaton (NFA) is a machine

$$M = (Q, \Sigma, \Delta, Q_0, F)$$

ternary relation set of initial states
 $\Delta \subseteq Q \times \Sigma \times Q$

Today: The power of nondeterminism

We said that every language in Reg is accepted by a DFA.

NFAs add the power of ϵ -transitions and choice.

Can they recognize languages outside Reg ?

Thm: The class of languages recognized by NFAs is **Reg**

We show that every DFA has an equivalent NFA.
What is it?

We show that for any NFA, there is an equivalent DFA.
What does a correct guess boil down to, operationally?
Try all possibilities, choose the right one!

How can one simulate an NFA using a DFA?

An NFA can be in one of multiple different states
on reading the same word!

Subset construction:

Each state of the DFA tracks some subset of states of the NFA
For now, consider NFAs without ϵ -transitions.

Consider an NFA $M = (Q, \Sigma, \Delta, Q_0, F)$

Want a DFA $M' = (Q', \Sigma, \delta', q'_0, F')$ s.t. $L(M) = L(M')$.

Each state of M' needs to track (potentially) multiple states of M

$$\text{So } Q' = 2^Q. \quad q'_0 = Q_0 \subseteq Q.$$

We now have to define δ' .

Consider a situation where $q_1, q_2, q_3 \in Q$, $a \in \Sigma$, and
 $\Delta(q_1, a, q_1)$ and $\Delta(q_1, a, q_2)$.

What should $\delta'(\{q_1\}, a)$ be? $\{q_1, q_2\} \in Q'$

We first extend Δ to a function $\hat{\Delta}: 2^Q \times \Sigma^* \rightarrow 2^Q$.

Any set of states of the NFA:

- stays the same on no input
- On a nonempty string, what?

$$\hat{\Delta}(S, \varepsilon) = S$$

$$\hat{\Delta}(S, xa) = \{q \mid \text{there is a } p \in \hat{\Delta}(S, x) \text{ s.t. } \Delta(p, a, q)\}$$

When does M accept s ?

There is a run from some $q_0 \in Q_0$ to some $f \in F$ on s .

$$\hat{\Delta}(Q_0, s) \cap F \neq \emptyset$$

Lemma: For $S \subseteq Q$ and $x, y \in \Sigma^*$,

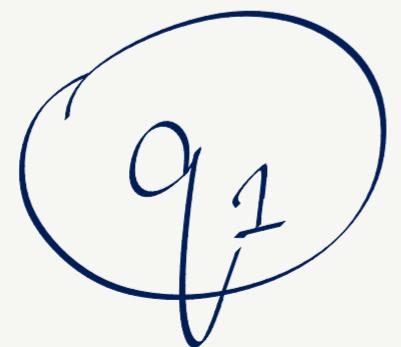
$$\hat{\Delta}(S, xy) = \hat{\Delta}(\hat{\Delta}(S, x), y)$$

Lemma: Suppose there exists an indexed set of subsets $S_i \subseteq Q$. Then,

for any $x \in \Sigma^*$,

$$\hat{\Delta}\left(\bigcup_i S_i, x\right) = \bigcup_i \hat{\Delta}(S_i, x)$$

Exercise: Prove these lemmas



So back to our DFA. $M' = (Q', \Sigma, \delta', q_0, F')$

$$Q' = 2^Q \quad q_0 = Q_0 \subseteq Q$$

$$\delta'(q, a) = \hat{\Delta}(q, a) \quad F' = \{q \in Q' \mid q \cap F \neq \emptyset\}$$

diff δ' to $\hat{\delta}: Q' \times \Sigma^* \rightarrow Q'$ ($\hat{\delta}(q, xa) = \delta'(\hat{\delta}(q, x), a)$)

Thm: For any $S \subseteq Q$ and $\omega \in \Sigma^*$,

$$\hat{\delta}(S, \omega) = \hat{\Delta}(S, \omega).$$

Thm: $L(M) = L(M')$ — consider which strings are accepted
and what those conditions are

Exercise: Prove these statements

How do we handle ϵ -transitions?

Thm: If L is accepted by an NFA with ϵ -transitions,
it is accepted by an NFA without any ϵ -transitions

Suppose $L = L(M)$ where $M = (Q, \Sigma_\epsilon, \Delta, Q_0, F)$.

Define ϵ -closure: $Q \rightarrow 2^Q$ as follows:

ϵ -closure(q) = $\{q'\} \mid \text{there is a path from } q \text{ to } q' \text{ in } M$
 $\text{consisting only of } \epsilon\text{-transitions}\}$

Now, $\hat{\Delta}(q, \epsilon) = \epsilon\text{-closure}(q)$

$\hat{\Delta}(q, xa) = \{p \mid \text{there is some } r \in \hat{\Delta}(q, x) \text{ s.t. } \Delta(r, a, p)\}$

Define a new NFA

$$M' = (Q', \Sigma, \Delta', Q'_0, F'), \text{ where}$$

$$Q' = Q \text{ and } Q'_0 = Q_0.$$

For any $a \in \Sigma$, $\Delta'(q, a, q')$ holds iff $q' \in \epsilon\text{-closure}(\hat{\Delta}(q, a))$

$$F' = F \cup \{q \in Q \mid \epsilon\text{-closure}(q) \cap F \neq \emptyset\}$$

Thm: $L(M) = L(M')$

The proof proceeds by induction on the length of an input word.
What is the base case? Finish this proof.