CSE 321- Introduction into Algorithm Design 141044086

Homework 1:

$$\begin{array}{l}
\overline{1}_{2}(n) = 3n^{4} + 3n^{3} + 1 \\
\overline{1}_{2}(n) = 3^{n} \\
\overline{1}_{3}(n) = (n-2)! \\
\overline{1}_{4}(n) = \ln^{2} n \\
\overline{1}_{5}(n) = 2^{2n}
\end{array}$$

To (n) = 3n

To(n) < To(n) < To(n) < To(n) < To(n) < To(n) < To(n)

$$\lim_{n \to \infty} \frac{T_{4}(n)}{T_{6}(n)} = \lim_{n \to \infty} \frac{2/n \ln n}{1/3} = \lim_{n \to \infty} \frac{1/n^{2/3}}{1/n^{2/3}} = \lim_{n \to \infty} \frac{1}{1/3} = \lim_{n \to \infty}$$

Ty(n) & O (To(n)) => Ty(n) & O (To(n))

$$\lim_{n\to\infty} \frac{T_{4(n)}}{T_{4(n)}} = \lim_{n\to\infty} \frac{\ln^2(n)}{3n^4 + 3n^3 + 1} = \lim_{n\to\infty} \frac{2\ln n}{12n^3 + 9n^2} = \lim_{n\to\infty} \frac{2\ln n}{12n^4 + 9n^3} = \lim_{n\to\infty} \frac{2\ln n}{12n^4 + 9n^4} = \lim_{n\to\infty} \frac{2\ln n}{12n^4 + 9n^4} = \lim_{n\to\infty} \frac{2\ln n}{12n^4 + 9n^4} = \lim_{n\to\infty} \frac{2\ln n}{12n^4 +$$

$$= \lim_{n \to \infty} \frac{2/n}{98n^3 + 24n^2} = \lim_{n \to \infty} \frac{2}{98n^4 + 24n^3} = 0$$

Tylnie O(Tyln)

$$\lim_{n \to \infty} \frac{T_1(n)}{T_2(n)} = \frac{8n^4 + 8n^5 + 1}{3^n} = \lim_{n \to \infty} \frac{12n^3 + 9n^2}{5^n \ln(3)} = \lim_{n \to \infty} \frac{36n^2 + 18n}{6n^2(3)} = \lim_{n \to$$

$$= \lim_{n \to \infty} \frac{42 \, n + 18}{\ell_{n}^{2}(3) \, 3^{n}} = \lim_{n \to \infty} \frac{42}{\ell_{n}^{3}(5) \, 3^{n}} = 0$$

Talnie O (Talni) V

$$\lim_{n \to 8} \frac{T_2(n)}{T_{(n)}} = \lim_{n \to \infty} \frac{3^n}{2^{2n}} = \left(\frac{3}{4}\right)^n = 0$$

Talm & O(Tsim) V

$$\lim_{n \to \infty} \frac{\overline{(s(n))}}{\overline{T_3(n)}} = \lim_{n \to \infty} \frac{y^n}{(n-2)!} = \lim_{n \to \infty} \frac{y^n}{(2\pi/(n-2))} \left(\frac{1}{n-2}\right)^2$$

$$n! = \sqrt{2\pi} n \left(\frac{n}{e}\right)^n$$

$$(n-2)! = \sqrt{2\pi} (n-2) \left(\frac{n-2}{e}\right)^{n-2} = \sqrt{2\pi} (n-2) \left(\frac{n-2}{n-2}\right)^n \left(\frac{e}{n-2}\right)^2$$

$$\Rightarrow \lim_{n \to \infty} \left(\frac{4e}{n-2}\right)^n \cdot \frac{(n-2)^2}{\sqrt{2\pi}(n-2) \cdot e} = 0$$

$$\overline{T_6(n)} \in O(T_3(n)) \checkmark$$

(2) This algorithm finds median value of the given

array. plum - min value of the array watermelon - max value of the given array

orange - median value of the given array orange Time has no particular role, procedure of finding min and mack value's could have been Jone without this unnecessary complication.

- b) i) Worst case occures when min=L[n]: $W(n) = 2n+1+n = 3n+1 \in \Theta(n)$
 - i) Best case occurres when min=LE1]: B(n)= n+1+n=2n+1 & O(n)
 - M A(n) = O(n) Because W(n) = B(n) = O(n)

(3) a)
$$\sum_{j=0}^{n-1} (j^2+j)^2 = f(n)$$

$$\int_{0}^{n-1} (x^2+j)^2 dx = f(n) = \int_{0}^{n} (x^2+j)^2 dx$$

$$\int_{0}^{n-1} (x^2+j)^2 dx = \int_{0}^{n-1} (x^4+2x^2+j) dy = (\frac{x^5}{5} + \frac{2}{3}x^3 + x) \Big|_{0}^{n-1} = \frac{(n-1)^5}{5} + \frac{2(n-1)^3}{3} + (n-1)$$

$$\int_{0}^{n-1} (x^2+j)^2 dx = \int_{0}^{1} f(n^5 + \frac{2}{3}n^3 + n)$$

$$\int_{0}^{n-1} (x^2+j)^2 dx = \int_{0}^{1} f(n^5 + \frac{2}{3}n^3 + n)$$

$$\int_{0}^{n-1} (x^2+j)^2 dx = \int_{0}^{1} f(n-1) = f(n) = \int_{0}^{1} f(n^5 + \frac{2}{3}n^3 + n)$$

$$\int_{0}^{n-1} f(n-1)^5 + \frac{2}{3}(n-1)^3 + (n-1) = f(n) = \int_{0}^{1} f(n) = \int_{$$

$$\frac{\sum_{i=1}^{n} (i+1)2^{i-1}}{\sum_{i=1}^{n} (i+1)2^{i-1}} = f(n)$$

$$\frac{\sum_{i=1}^{n} (i+1)2^{i-1}}{$$

 $f(n) \in \Theta(n2^n)$

$$n = 2^{K} - 1$$

$$\sum_{i=0}^{K} 2^{i} = 2^{K+1} - 1 = f(x)$$

$$\int_{0}^{\infty} 2^{x} dx \leq f(x) \leq \int_{0}^{\infty} 2^{x} dx$$

$$\int_{0}^{1} 2^{x} dx = \left(\frac{2^{x}}{\ln 2}\right) \left| \frac{\log_{2}(n+1)}{\log_{2}(n+1)} - \frac{1}{\ln 2} \right| = \frac{n+1}{\ln 2} - \frac{1}{\ln 2} = \frac{n+1}{\ln 2}$$

$$= \frac{n}{\ln 2} = \frac{1}{\ln 2} \cdot n$$

$$\int_{0}^{\log(n+2)} 2^{x} dx = \frac{n+2}{\ln 2} - \frac{1}{\ln 2} = \frac{1}{\ln 2} (n+1) = \frac{n}{\ln 2} + \frac{1}{\ln 2}$$

$$p(x) \in \Phi(h)$$

(b) a)
$$n^{3} \in O(3^{2n})$$

$$\lim_{n \to \infty} \frac{n^{3}}{3^{2n}} = \lim_{n \to \infty} \frac{n^{3}}{g^{n}} = \lim_{n \to \infty} \frac{3n^{2}}{g^{n} \ln g} = \lim_{n \to \infty} \frac{6n}{g^{n} (\ln g)^{2}} = \lim_{n \to \infty} \frac{6}{g^{n} (\ln g)^{3}} = 0$$

$$n^{3} \in O(3^{2n}) \to n^{3} \in O(3^{2n}) \vee$$

b)
$$n \in o(\log \log n) \times$$
 $\lim_{n \to \infty} \frac{h}{\log \log n} = \lim_{n \to \infty} \frac{1}{n \log n (\log n)^2} = \lim_{n \to \infty} n \log n (\log n)^2 = \infty$
 $n \in V(\log \log n)$

d)
$$\lim_{n \to \infty} \frac{\sqrt{10n^2 + 7n + 3}}{n} = \lim_{n \to \infty} \frac{\sqrt{n^2(10 + \frac{7}{n} + \frac{3}{n^2})}}{n} = \lim_{n \to \infty} \frac{\sqrt{10 + \frac{7}{n} + \frac{3}{n^2}}}{n}$$

$$= \lim_{n \to \infty} \sqrt{10 + \frac{7}{n} + \frac{3}{n^2}} = \sqrt{10}$$

$$\sqrt{10n^2 + 7n + 3} \in \Theta(n)$$