Yet Another SAT Encoding for Car Sequencing

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September 12, 2013

The encoding is wrt. one object (car or option) for which we omit the subscript. We present an encoding for the constraint

$$\bigwedge_{i=0}^{n-q}(\sum_{l=1}^{q} x_{i+l} \leq u)$$

The encoding avoids a separate counter encoding for each cardinality constraint and at the same time is an alternative to [1].

- x_i is true (1) if the object is at position i.
- $s_{i,j}$ is true if in window $i-q+1, \ldots i$ there are at least j objects.

The idea behind the clauses is to express the relationship between x_{i-q}, x_i and the counter $s_{i,j}$.

$$\neg s_{i,j} \vee s_{i+1,j-1} \tag{1}$$
$$\neg s_{i+1,j+1} \vee s_{i,j} \tag{2}$$
$$\neg s_{i,j+1} \vee s_{i,j} \tag{3}$$
$$\neg x_i \vee \neg s_{i-1,j} \vee s_{i,j} \tag{4}$$
$$x_i \vee \neg s_{i,j} \vee s_{i-1,j} \tag{5}$$
$$\neg x_{i-q} \vee \neg s_{i,j} \vee s_{i-1,j} \tag{6}$$
$$x_{i-q} \vee \neg s_{i-1,j} \vee s_{i,j} \tag{7}$$
$$\neg x_i \vee x_{i-q} \vee \neg s_{i-1,j-1} \vee s_{i,j} \tag{8}$$
$$x_i \vee \neg x_{i-q} \vee \neg s_{i,j} \vee s_{i-1,j+1} \tag{9}$$

References

[1] Valentin Mayer-Eichberger and Toby Walsh. SAT Encodings for the Car Sequencing Problem. *Pragmatics of SAT*, Workshop at SAT, Helsinki, 2013.