Thread-Safe Cache Study Guide

Overview

The Thread-Safe LRU Cache demonstrates advanced concurrency patterns for high-performance caching systems with reader-writer synchronization and lazy LRU management.

Key Concepts Covered

1. Reader-Writer Synchronization

```
std::shared_mutex mutex_;

// Multiple concurrent readers
std::shared_lock<std::shared_mutex> read_lock(mutex_);

// Exclusive writer access
std::unique_lock<std::shared_mutex> write_lock(mutex_);
```

2. LRU (Least Recently Used) Algorithm

```
// Doubly-linked list + hash map = O(1) operations
std::unordered_map<Key, Iterator> cache_map_;
std::list<std::pair<Key, Value>> lru_list_;
```

- 3. Lazy vs Eager Updates
 - Eager: Update LRU order on every read (high contention)
 - Lazy: Update LRU order only on writes (better performance)

Data Structure Design

Hybrid Architecture

```
struct CacheNode {
    Key key;
    Value value;
    // Implicit position in linked list
};

// Fast lookup: O(1)
std::unordered_map<Key, list_iterator> map_;

// LRU ordering: O(1) insert/remove
std::list<CacheNode> lru_list_;
```

Operations Complexity

Operation	Time	Space	Contention
get() put() eviction	O(1)	O(1)	Shared lock
	O(1)	O(1)	Exclusive lock
	O(1)	O(1)	During put()

Real-World Applications

System Caches

CPU caches: L1/L2/L3 cache hierarchies
 Database buffer pools: Page caching

Web caches: HTTP response cachingCDNs: Content delivery networks

Application Caches

• Redis/Memcached: Distributed caching

• In-memory caches: Application-level caching

File system caches: OS page cacheBrowser caches: Resource caching

Interview Questions & Answers

Q: "Why use shared_mutex instead of regular mutex?"

A: Performance optimization for read-heavy workloads: - Concurrent reads: Multiple threads can read simultaneously - Exclusive writes: Only one writer, no readers during writes - Typical cache ratio: 90% reads, 10% writes - Speedup: 5-10x improvement in read-heavy scenarios

Q: "How do you handle cache eviction?"

A: LRU eviction strategy:

```
if (cache_map_.size() >= capacity_) {
    // Remove least recently used (back of list)
    auto lru_key = lru_list_.back().first;
    cache_map_.erase(lru_key);
    lru_list_.pop_back();
}
```

Q: "What about thread safety during eviction?"

A: Write lock protects entire operation: - Check capacity while holding exclusive lock - Eviction and insertion are atomic - No partial states visible to readers

Q: "Why not update LRU on every read?"

A: Lazy LRU trade-off: - Pros: Better performance, less contention - Cons: Slightly less accurate LRU ordering - Real-world: Most caches use approximations anyway

Design Patterns Demonstrated

1. Adaptive Locking Strategy

```
// Readers use shared locks
std::shared_lock<std::shared_mutex> lock(mutex_);
// Writers use exclusive locks
std::unique_lock<std::shared_mutex> lock(mutex_);
```

2. RAII Lock Management

- Automatic lock release
- Exception safety
- Clear lock scope boundaries

3. Template Specialization

```
template<typename Key, typename Value>
class ThreadSafeCache {
    // Generic for any key-value types
    // Requires Key to be hashable
};
```

Performance Optimization Techniques

- 1. Lock Granularity
 - Coarse-grained: One lock for entire cache
 - Fine-grained: Bucket-level locking (more complex)
 - Trade-off: Simplicity vs scalability

2. Memory Layout

```
// Cache-friendly access patterns
std::list<std::pair<Key, Value>> lru_list_; // Sequential access
std::unordered_map<Key, Iterator> map_; // Random access
```

3. Move Semantics

```
void put(Key key, Value value) {
    // Move to avoid copies
```

```
lru_list_.emplace_front(std::move(key), std::move(value));
}
```

Cache Replacement Policies

LRU Alternatives

- LFU (Least Frequently Used): Count-based eviction
- FIFO: Simple queue-based eviction
- Random: Minimal overhead
- ARC (Adaptive Replacement Cache): Balances recency and frequency

When to Use Each

- LRU: General-purpose, good temporal locality
- LFU: Frequency matters more than recency
- FIFO: Simple, predictable behavior
- Random: Very low overhead, surprisingly effective

Test Scenarios Covered

- 1. BasicOperations: put/get functionality
- 2. EvictionBehavior: LRU ordering verification
- 3. ConcurrentAccess: Multiple readers/writers
- 4. CapacityLimits: Proper eviction handling
- 5. MemoryLeaks: Resource cleanup verification
- 6. StressTest: High contention scenarios

Common Interview Challenges

1. "Implement LRU cache in 15 minutes"

Key points to hit: - Hash map for O(1) lookup - Doubly-linked list for O(1) eviction - Move recently used to front - Evict from back when full

2. "How would you scale this to multiple machines?"

Distributed cache considerations: - **Consistent hashing**: Distribute keys across nodes - **Replication**: Handle node failures - **Cache coherence**: Keep replicas synchronized - **Hot spotting**: Load balancing strategies

3. "What if the cache is too big for memory?"

Hybrid storage strategies: - **Tiered caching**: Memory + disk - **Compression**: Reduce memory footprint - **Eviction to disk**: Spillover storage - **Memory mapping**: OS-managed paging

Production Considerations

Monitoring & Metrics

• Hit rate: Percentage of cache hits

Eviction rate: How often items are removed
Memory usage: Cache size monitoring

• Contention: Lock wait times

Configuration Tuning

Cache size: Balance memory vs hit rate
Eviction policy: Match access patterns
Concurrency level: Readers vs writers ratio
Warmup strategy: Pre-populate important data

Advanced Extensions

Lock-Free Cache

Hazard pointers: Safe memory reclamation
Atomic operations: CAS-based updates
Memory ordering: Sequential consistency

• Complexity: Much harder to implement correctly

Partitioned Cache

• Segment locks: Reduce contention

Hash-based partitioning: Distribute load
NUMA awareness: Thread-local caches

This cache implementation demonstrates sophisticated concurrent data structure design - a key skill for high-performance systems and distributed architectures!