

## Number theory in cryptography

## – Exercise set 7 –

The exercises T7.1, T7.3 have to be handed in on Tuesday, 16th April 2024, 8:30 at latest.

## THEORETICAL QUESTIONS

**T 7.1** Let  $\psi(x, y)$  be the number of  $y$ -smooth integers in the interval  $[1, x]$  as introduced in the lecture notes on integer factorization. Let  $f$  be a real-valued function defined for  $y \geq 2$  and satisfying  $f(y) \geq 1$  for all  $y$  and  $f(y) = y^{1+o(1)}$  as  $y \rightarrow \infty$ .

We let  $y = L_x(1/2, v)$  for some parameter  $v > 0$ . Prove that as  $x \rightarrow \infty$ , we have

$$\frac{xf(y)}{\psi(x, y)} \sim L_x(1/2, g(v) + o(1))$$

for some function  $g : \mathbb{R}_{>0} \rightarrow \mathbb{R}_{>0}$  satisfying  $g(v) \geq \sqrt{2}$  for all  $v > 0$ .

**T 7.2** Various algorithms for both integer factorization and discrete log (to be seen later in class) produce auxiliary numbers up to some bound  $x$  via random samples and then test these numbers for smoothness. Suppose that our goal is to find  $y^{1+o(1)}$  numbers that are  $y$ -smooth for some  $y < x$  (that depends on  $x$ ).

How would you choose  $y$  (using the  $L$ -notation for the complexity parameter  $x$ ) so that you minimize the number of samples needed to achieve the goal? Show that there is a choice of  $y$  for which  $L_x(1/2, \sqrt{2})$  samples are sufficient for achieving the goal.

**T 7.3** Define  $L_{\alpha, C}(X) := \exp(C \log(X)^\alpha \cdot \log(\log(X))^{1-\alpha})$  for  $C > 0, X > 1$  and  $0 < \alpha < 1$ . The goal is to show that  $L_{\alpha, C}(X)$  is subexponential in  $\log(X)$ , namely prove the following claims:

- a) For all  $\varepsilon > 0$ ,  $L_{\alpha, C}(X)$  is smaller than a constant times  $X^\varepsilon$  for  $X$  large enough.
- b) For all  $N > 0$ ,  $L_{\alpha, C}(X)$  is bigger than a constant times  $(\log X)^N$  for  $X$  large enough.

**T 7.4** Let  $\alpha, \beta, r, s \in \mathbb{R}_{>0}$  be given with  $s < r \leq 1$ . Show that the probability that a random positive integer less than or equal to  $L_x(r, \alpha)$  is  $L_x(s, \beta)$ -smooth is

$$L_x(r - s, -\alpha \cdot (r - s)/\beta)$$

as  $x \rightarrow \infty$ .

## PROGRAMMING EXERCISES

**P 7.1 Bonus question: Easter Egg Hunt!** You intercept the Easter bunnies' communications. You know the bunnies are using RSA with  $e = 11$  and

$N = 17850620114655432894259040410860055152615689657922061545941753524468448$   
 $52662729862255910375678524308427988269101598453525050708892379497516176$   
 $84565873359332289011212047453797368854717523411001544538018037006070502$   
 $00128340099072001$

and you intercepted the ciphertext

$c = 2958069616077860652412945389416718089451490811802433593048131259639033693$   
78512074906959493801275205949299424164997057607126394065736514321370364367  
84439921229177447569160651700234332868383974579267263248532933176771637594  
24201496

You also know the message was encoded using a shift of ASCII: every block of 2 plaintext digits (which will be a number  $k$  between 10 and 99) encodes the character with ASCII code  $k + 22$  if  $k \leq 73$  and  $k + 23$  otherwise. (e.g., the 12-character string **Hello World!** corresponds to the 24-digit integer 507885858810658891857711). Can you recover the plaintext?