

# Linked Lists

## ① Basic struct, class & pointer program

- we create a class (a user defined datatype) named node (now a datatype node is created)
- containing 2 items i.e. - int data  
- node \*next } with public access specifier  
(next is a pointer of type node)
- now we create a constructor to assign values for data & next  
data = d  
next = NULL ptr; } Public

main function

arr = { 2, 5, 7, 9 }

- here we create a node using new keyword for an element in the array. it creates a node in a random place in heap memory
- we store this newly created nodes address in arr.



\* new Node (v[2])

- Basically what happens here is v[2] is passed to the constructor
- Stores v[2] as data in node & next as null

- new if we print n - it prints address of the new node

- n → data - this basically means it points & gets access to the class member data

- Simplified version of (\*n).data

(data = (\*n).data)

## Class Definition

```
Struct Node {  
Public:
```

```
    int Data;  
    Node *next;
```

Struct all members  
are Public by default  
where as Class members  
are Private by Default  
unless specified

\*next is a pointer  
Variable of type node  
i.e \*next holds the  
address of the next  
Node & each node has  
Data & next

```
Public:
```

```
Node (int data)  
{  
    Data = data;  
    next = null ptr;
```

### Constructor

For the Defined class  
you're basically assigning  
value

```
int main()  
{
```

```
    vector<int> arr = {2, 5, 8, 7};
```

```
Node * y1 = new Node (arr[0]);  
Node * y2 = new Node (arr[1]);  
Node * y3 = new Node (arr[2]);  
Node * y4 = new Node (arr[3]);
```

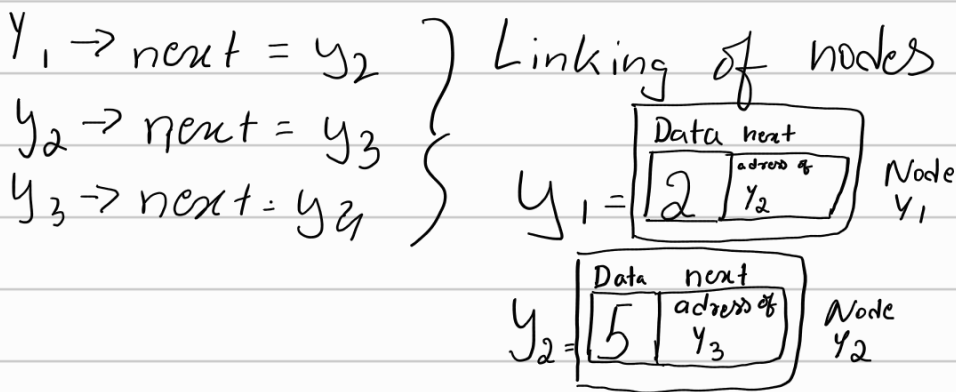
} Assigning values  
- y1 is a pointer  
variable of type node  
which means

$y_1$  has data & next  
 - we create new Node  
 & pass the data which  
 gets assigned in constructor  
 now  $*y_1$  has the  
 address of the new  
 Node

$y_1 =$ 

data	Next
------	------

  
 ↪ 0x2c3d83  
 address



$\text{cout} \ll y_1 \rightarrow \text{Data} \ll " " \ll y_1 \rightarrow \text{next};$  } 2 address of  $y_2$   
 $\text{cout} \ll y_2 \rightarrow \text{Data} \ll " " \ll y_2 \rightarrow \text{next};$  } 5 address of  $y_3$   
 .  
 .  
 .



### ③ Array to linked list

St 1- create a class with data & node \*next

- & here \*next itself is a node type so it will have data & next.

St 2- function to convert array to LL

- returns the head node

- ~~initially~~ assign head node to the first ele of the array  
Node \*head = new node(a[0]);

- keep a current pointer used to link the nodes initially set to head (so current now has head node)
- iterate through the array from  $a[1] \rightarrow N$

- for each element you create a new node & store that in a temp pointer.

- how to store this temp (i.e address of new node in previous nodes next)

- we ~~access~~ access current nodes next & store temp init.

current->next = temp

so after linking we move the current ahead  
repeat for all elements of arr.

S+3- traversal through the created linked list

- we get the head node.
  - we keep a current pointer to initially point at head.
  - now we create a while loop till the current points to null
  - until then we can print the data & its next.
  - then update current to move to next node
    - ↳ create temp node get the current node's next & store that in temp
    - ↳ now  $current = temp$
- ```
node * temp;
temp = current -> next;
current = temp;
```

OR

- ↳ directly update current
- ```
current = current -> next;
```



— / — / —

④ Deletion of head;

- Store the current head node in a temp node
- move the head to the next node
- free up the space by deleting temp.

⑤ Deletion of tail

- Start traversing from the head node
- do the traversal till last but 2nd element  
ie head → [1] → [2] → [3] → NULL  

|  
till here

So when you get  $current \rightarrow next \rightarrow next = null$   
then come out of the loop

- just delete current's next node (i.e. last node)
- $\&$  set  $\text{current} \rightarrow \text{next} = \text{Null}$ .

⑥ Delete  $K^{th}$  node

Edge cases

- if list is empty return null
- if  $K = 1$  then delete head
- basically you have 2 trackers  
one is current & another for prev
- keep a counter
- traversal till current points to null
  - ↳ if count ==  $K$   
then link previous node to current's next node & del current
  - ↳ & break
- if not just keep traversing by  
storing current in prev  
move current ahead  
& count ++.

⑦ Delete node with value  $x$ .

- Same as previous one  
but just compare data with  $x$



\_/\_/\_

### ⑧ Insertion at head

- ~~edge case~~
- just create a new node
- just point the new node's next to head
- return new node.

### ⑨ Insertion at tail

Edge case

- if list is empty then just create new node & return

Else

- create new node
- set last node's next to new node
- & new node's next to null.

### ⑩ Insertion at kth position

Edge case

- if given pos is  $\leq 1$  then return null.
- if given pos = 1 then insert at head.

- just start counting the nodes while tracking current node & previous node

- once the count matches the given position then

- set new node's next to current or prev → next

- then  $\text{prev node} \rightarrow \text{next}$  to new node.

- Edge case

- if at last ~~pos~~ is last place i.e. insertion at tail

(11) insert value ~~at~~ before  $X$  value

- Same as previous one

Edge case

- if  $X$  is first node value or head's value then insert node before head.

- & in while loop just check if  $\text{value} = \text{current data}$ .

(12) Array to DLL

- Create head node

- traverse through the array from 1 to  $n$  creating new node for each element.

- & initially current is initialized to head node as & when we create new nodes we initialize  $\text{current} \rightarrow \text{next}$  to new node & new node's prev as current & then move the current ahead.  
& finally return the head.