

Stacks & queues

- ① stack using array
- 1- in a class stack (Private)
- define an array size = arr[100]
 - ~~also~~ keep a variable to track index of top element [Top]
 - a variable to assign a max capacity

1.1- Public constructor

- initialize capacity to take in user input size
- initialize top = -1.

1.2 member functions

- ~~bool~~ is Empty ()
if top = -1 then return -1
- bool is full ()
if top = capacity - 1
then return full
- void push (int x)
if stack isn't full
then move top ahead
and insert element.

//_

- void pop()
if stack is not empty
then top --
now return arr[top]

- void peek()
if stack is not empty
return top element

- void display()
- if stack empty return empty
- else iterate from 0 to top
 & print eles.

Q.

Queue using Arrays

- in enqueue & dequeue just remember that you could just do $\text{rear}++$ or $\text{front}++$ to move the pointers
- but if front or rear is at last & you have to enqueue something in rear which is at last but empty space is at start then you do $\text{rear} = (\text{rear} + 1) \% \text{capacity}$ that will give bring back rear to start.
- & after inserting if front is at last you have to bring it back then something again.

Stacks & queues

⑤ Balanced Parenthesis.

- just iterate through the given String
- if current char is an opening bracket just push it to stack.
- if the current char is a closing bracket then look into the stack if the top matches if it doesn't you can return unbalanced directly.
- but if it does match keep going till the end matching the closing ones with the opening ones until the stack is empty & then return true (Balanced).

monotonic stack

① Next greater element
(a simple array)

op - iterate through the array backwards

- for each element check if stack is
 \hookrightarrow stack is ^{not} empty & $\text{st.top} < \text{curr ele}$
 then keep popping till you find a greater ele
 than curr
- now check if stack is empty if yes then
 store -1 in ans vector
- else store st.top in the ans vector
- then just push the current ele to stack
- finally return ans vector

(when you're given 2 arrays)

where arr1 is a subset of arr2 .

- #
- just iterate through the arr2 from back
 - for every element check stack ~~for~~ for greater
 element (if smaller pop till you find greater or becomes
 empty)
 - now if stack is empty NGE for current number
 is -1 so in map store $\{\text{curr} : -1\}$
 - if you find a greater element than current
 then in map store $\{\text{curr} : \text{st.top}\}$
 \hookrightarrow NGE

Now the map would have all NGEs of
 arr2 . ~~as~~ $\{\text{arr2}[\text{ele}] : \text{NGE}\}$

- _/_/_
- Now just iterate through arr & for each element access NGE in map and push it to ans vector
 - now return ans.

(when given a circular array)

- # - basically iterating from $2n-1 \rightarrow 0$. Hypothetically extends the array
- so till $i \geq n$ i.e. in hypothetically part we just store all arr elements into stack
- now in actual iteration $i < n$ we look for NGEs & store in ans usual way.

③ Min Stack

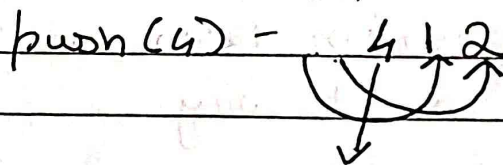
- just declare a stack that stores a pair $\{cur, min\}$
- while pushing into a stack just if it's empty then push $\{cur, cur\}$
 \hookrightarrow initial min
- if not empty then $\{cur, \min(cur, st.top().second)\}$

⑧ Stack using Queue

- basically you declare a queue & that should behave as a stack
- so ~~to~~ just while pushing
 - get the size of the current queue
 - push a value into queue
 - now loop through the size and basically reverse the queue

ex: $q.size = 2$ $[1, 2]$

$i = 0 \rightarrow 1$



new front

i.e the last pushed ele [LIFO]

⑨ Queue using Stack

- basically you declare ² Stack & that should work like a queue
- 1st stack takes in new eles and 2nd stack is reverse of 1st stack
- while pushing directly push in the stack 1
- while popping or getting to p then push all elements from S_1 to S_2 & then pop or top.