



Today's agenda

- ↳ Char and String
- ↳ ASCII
- ↳ Problems

↳ Sunday → optional class (Problem Solving) → 8PM

↳ Alpoor Kumar → 7* on Codechef

masters on Codeforces



AlgoPrep



String: Sequence of Characters
↳ Upper case, Lowercase, Special etc.

String st = "AlgoBsp";

Characters: a) Alphabet

- ↳ a-z (lowercase)
- ↳ A-Z (uppercase)

b) Special characters

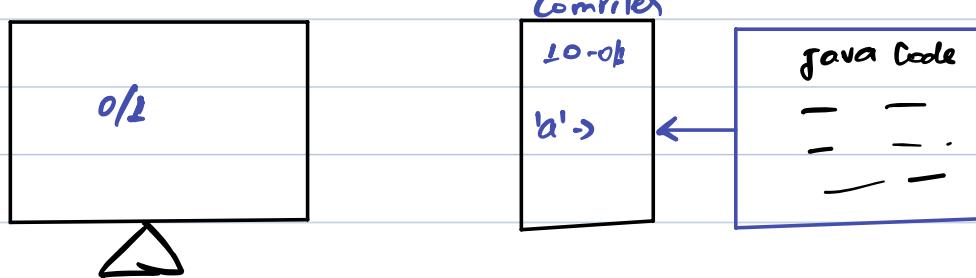
- ↳ @ # ? etc.

c) Numbers

- ↳ '0' - '9'

Syntax:

char ch = 'A';
type, name, character



1. numbers → binary number
 2. Characters → number → binary number
- Predefined*



ASCII → 256 Characters

'A' : 65

'B' : 66

'C' : 67

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* char rules

1. char to int: implicit. ex: int a = 'A'; ↗

2. int to char: Complicated. → (Typecast always)
↳ implicit.

Char ch = (char)66;



AlgoPrep



Quiz 1 :

Char ch1 = 'B';
S.O.P(ch1);

Quiz 2 :

Char ch2 = 66; → typecast automatic → 'B'
S.O.P(ch2);

Quiz 3 :

Char ch3 = 'xyz'; → error
S.O.P(ch3);

Quiz 4 :

int n = 'A'; → 65
n = n + 2; → 67
System.out.println(n); → 67



Quiz 5:

→ mathematical
Operation in
Character allowed

(ASCII)

Char ch4 = 'A';

ch4 = ⁶⁵(₆₅)ch4 + 3;

s.o.p (ch4);

Implicit conversion.

Quiz 6:

Char ch5 = 'A';

⁶⁵if (ch5 >= 90) {

s.o.p ("Greater");

}

else {

s.o.p ("Smaller");

→ Smaller

}

Break till 9:45 PM

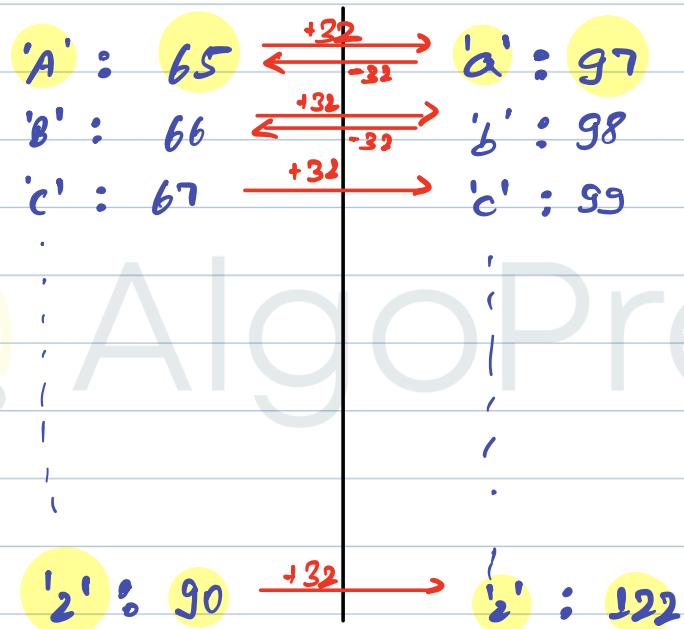


Q) Toggle Character

Given a `char[]` which contains only small and capital letters, toggle them.

↳ lowercase → uppercase
uppercase → lowercase

Ex: AlgoPrep : aLGHOpREP



uppercase to lowercase → +32

lowercase to uppercase → -32



II Pseudo Code

```
P S void main( ) {
    Scanner scn = new Scanner (System.in);
    int n = scn.nextInt();
    char[] ch = new char[n];
    String st = scn.nextLine();

    for (int i=0; i<n; i++) {
        ch[i] = st.charAt(i);
        Toggle (ch);
    }
}
```



AlgoPrep

```
P S void Toggle [char() ch) {
    for (int i=0; i<ch.length; i++) {
        if (ch[i]>=65 && ch[i]<=90) { // uppercase
            ch[i] = (char)(ch[i] + 32);
        } else { // lowercase
            ch[i] = (char)(ch[i] - 32);
        }
    }
}
```



Q) Reverse the given string

Given a string str, reverse and print it.

Ex: algoPrep : prepoga

// Pseudo Code

P S String reverseString (String str){
int n = str.length();

Char [] ch = str.toCharArray();

int SP = 0;

int EP = n-1;

while (SP < EP) {

 char temp = ch[SP];

 ch[SP] = ch[EP];

 ch[EP] = temp;

 SP++; EP--;

String ans = str.valueOf(ch);

return ans;

}



String st: "Hello";

st = st + "e";

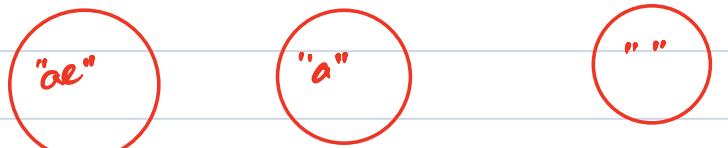
S-O-P (st); → Helloe

Char [] ch = { 'a', 'e', 'o', 'm' };

String ans: "";

i	ans
0	"a"
1	"ac"
2	"aeo"
3	"aeom"

```
for (int i=0; i<ch.length; i++) {  
    ans = ans + ch[i];  
}
```



Toggle Character

Java Code:

```
import java.io.*;
import java.util.*;

public class Solution {

    public static void main(String[] args) {
        Scanner scn=new Scanner(System.in);
        String st=scn.nextLine();

        char[] ch = st.toCharArray();
        toggle(ch);
        for(int i=0;i<ch.length;i++){
            System.out.print(ch[i]);
        }
    }

    private static void toggle(char[]ch) {

        for(int i=0;i<ch.length;i++){
            if(ch[i] >= 65 && ch[i]<= 90){
                ch[i] = (char)(ch[i] + 32);
            }else{
                ch[i] = (char)(ch[i] - 32);
            }
        }
    }
}
```

C++ Code:

```
#include <iostream>
#include <cctype>
using namespace std;

void toggle(char *ch, int length) {
    for (int i = 0; i < length; i++) {
        if (isupper(ch[i])) {
```

```

        ch[i] = tolower(ch[i]);
    } else {
        ch[i] = toupper(ch[i]);
    }
}

int main() {
    string st;
    getline(cin, st);

    char ch[st.length() + 1];
    strcpy(ch, st.c_str());

    toggle(ch, st.length());

    for (int i = 0; i < st.length(); i++) {
        cout << ch[i];
    }

    return 0;
}

```

Python Code:

```

def toggle(ch):
    toggled = []
    for char in ch:
        if char.isupper():
            toggled.append(char.lower())
        else:
            toggled.append(char.upper())
    return ''.join(toggled)

def main():
    st = input()
    toggled_string = toggle(st)
    print(toggled_string)

if __name__ == "__main__":
    main()

```

Reverse String

Java Code:

```
import java.io.*;
import java.util.*;

public class Solution {

    public static void main(String[] args) {
        Scanner scn=new Scanner(System.in);
        String st=scn.nextLine();

        char[] ch = st.toCharArray();
        reverse(ch);
        for(int i=0;i<ch.length;i++){
            System.out.print(ch[i]);
        }
    }

    private static void reverse(char[]ch) {

        int sp =0;
        int ep = ch.length - 1;

        while(sp < ep){
            char temp = ch[sp];
            ch[sp] = ch[ep];
            ch[ep] = temp;
            sp++;
            ep--;
        }
    }
}
```

C++ Code:

```
#include <iostream>
```

```

#include <cstring>
using namespace std;

void reverse(char *ch, int length) {
    int sp = 0;
    int ep = length - 1;

    while (sp < ep) {
        char temp = ch[sp];
        ch[sp] = ch[ep];
        ch[ep] = temp;
        sp++;
        ep--;
    }
}

int main() {
    string st;
    getline(cin, st);

    char ch[st.length() + 1];
    strcpy(ch, st.c_str());

    reverse(ch, st.length());

    for (int i = 0; i < st.length(); i++) {
        cout << ch[i];
    }

    return 0;
}

```

Python Code:

```

def reverse(ch):
    sp = 0
    ep = len(ch) - 1

    while sp < ep:
        ch[sp], ch[ep] = ch[ep], ch[sp]
        sp += 1
        ep -= 1

def main():
    st = input()

```

```

ch = list(st)
reverse(ch)
reversed_string = ".join(ch)
print(reversed_string)

if __name__ == "__main__":
    main()

```

Insert Difference_HW

Solution Vid:

https://youtu.be/6u99_gdqtOU

Java Code:

```

import java.io.*;
import java.util.*;

public class Solution {

    public static void main(String[] args) {
        Scanner scn=new Scanner(System.in);
        String st=scn.nextLine();
        System.out.println(insert(st));
    }

    private static String insert(String st) {
        String ans = "";

        for(int i=0;i<st.length()-1;i++){
            char ch1 = st.charAt(i);
            ans = ans + ch1;
            int temp = st.charAt(i+1)-st.charAt(i);
            ans = ans + temp;
        }

        ans = ans + st.charAt(st.length()-1);
        return ans;
    }
}

```

C++ Code:

```
#include <iostream>
#include <string>
using namespace std;

string insert(string st) {
    string ans = "";

    for (int i = 0; i < st.length() - 1; i++) {
        char ch1 = st[i];
        ans += ch1;
        int temp = st[i + 1] - st[i];
        ans += to_string(temp);
    }

    ans += st[st.length() - 1];
    return ans;
}

int main() {
    string st;
    getline(cin, st);

    string result = insert(st);
    cout << result << endl;

    return 0;
}
```

Python Code:

```
def insert(st):
    ans = ""

    for i in range(len(st) - 1):
        ch1 = st[i]
        ans += ch1
        temp = ord(st[i + 1]) - ord(st[i])
        ans += str(temp)

    ans += st[-1]
    return ans

def main():
```

```

st = input()
result = insert(st)
print(result)

if __name__ == "__main__":
    main()

```

Is Palindrome_HW

Solution Vid:

<https://youtu.be/QBsaSTN48kA>

Java Code:

```

import java.io.*;
import java.util.*;

public class Solution {

    public static void main(String[] args) {
        Scanner scn=new Scanner(System.in);
        String st=scn.nextLine();
        int n = st.length();

        char[] arr = new char[n];
        for(int i=0;i<n;i++){
            if(st.charAt(i)>='A' && st.charAt(i)<='Z'){
                arr[i] = (char)(st.charAt(i) + 32);
            }else{
                arr[i] = st.charAt(i);
            }
        }

        int sp = 0;
        int ep = n-1;

        while(sp<ep){
            if(arr[sp]!= arr[ep]){
                System.out.println("false");
                return;
            }
        }
    }
}

```

```

        sp++;
        ep--;
    }

    System.out.println("true");
}
}

```

C++ Code:

```

#include <iostream>
#include <cctype>
using namespace std;

int main() {
    string st;
    getline(cin, st);
    int n = st.length();

    char arr[n];
    for (int i = 0; i < n; i++) {
        if (isupper(st[i])) {
            arr[i] = tolower(st[i]);
        } else {
            arr[i] = st[i];
        }
    }

    int sp = 0;
    int ep = n - 1;

    while (sp < ep) {
        if (arr[sp] != arr[ep]) {
            cout << "false" << endl;
            return 0;
        }

        sp++;
        ep--;
    }

    cout << "true" << endl;
    return 0;
}

```

Python Code:

```
def main():
    st = input()
    n = len(st)

    arr = [0] * n
    for i in range(n):
        if st[i].isupper():
            arr[i] = st[i].lower()
        else:
            arr[i] = st[i]

    sp = 0
    ep = n - 1

    while sp < ep:
        if arr[sp] != arr[ep]:
            print("false")
            return

        sp += 1
        ep -= 1

    print("true")

if __name__ == "__main__":
    main()
```

Today's agenda

- ↳ Prefix and Suffix String
- ↳ LPS of a given String
- ↳ LPS array of a given String
- ↳ Problems

Q) Given a string of size n.

Prefix strings: All the substrings starting from 0th idn.
Suffix strings: " " " ending at last idn.

Ex: a b d b

Prefix strings

a
a b
a b d
a b d b

Suffix strings

b
d b
d b b
a b d b

Q) Given S_n , LPS of the String.

length of longest prefix which is also a suffix.

Note: other than entire string

(longest proper prefix which is also proper suffix)

Ex: $S = b \text{ } d \text{ } a \text{ } b \text{ } d$

Prefix

b
b d
b d a
b d a b
~~b d~~ ~~a~~ ~~b d~~

Suffix

d
b d → 2
a b d
d a b d
~~b d~~ ~~a~~ ~~b d~~

$S = S_0 S_1 S_2 S_3 S_4$

	index:
S_0	S_4 1
$S_0 S_1$	$S_3 S_4$ 2
$S_0 S_1 S_2$	$S_2 S_3 S_4$ 3
$S_0 S_1 S_2 S_3$	$S_1 S_2 S_3 S_4$ 4
	⋮
	n+1

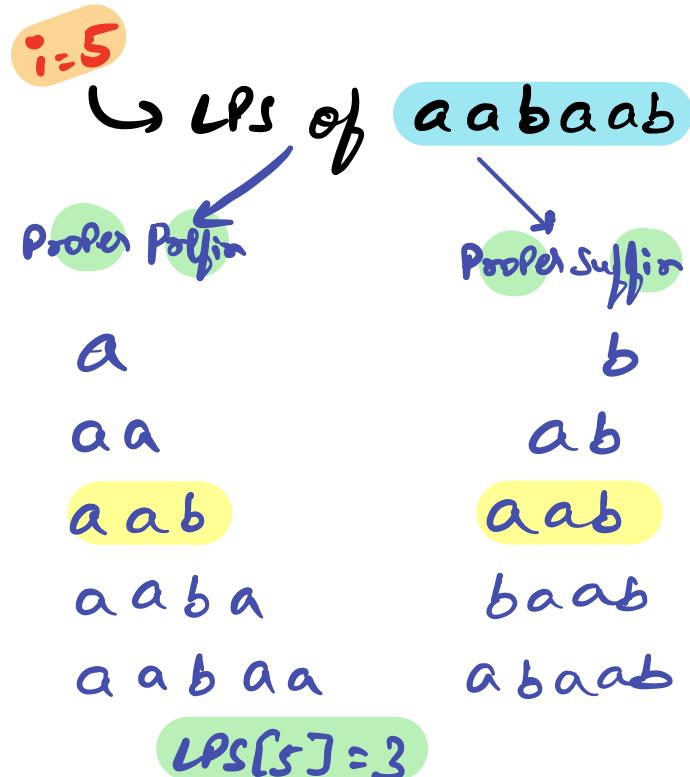
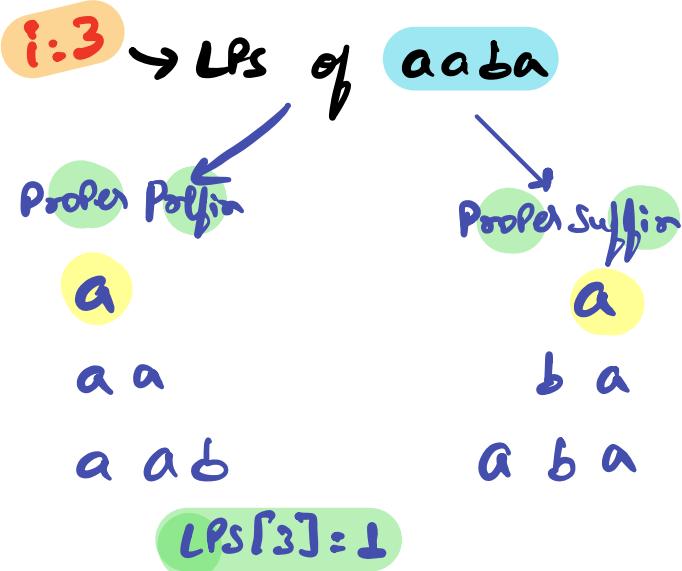
T.C: $\frac{n + O(n)}{2}$

$\frac{(n+1)n}{2} \approx O(n^2)$

Q) Given S_n , return $LPS[n]$

Ex: $s[7]: a \underset{0}{a} \underset{1}{b} \underset{2}{a} \underset{3}{a} \underset{4}{b} \underset{5}{a} \underset{6}{a}$
 $LPS[7]: 0 \perp 0 \perp 2 3 4$

$LPS[i] = LPS$ value of substring l_0, \dots, i



T.C: $n * n^2 \approx O(n^3)$

int $[n]$ LPS (string s)
"already written
in $O(n)$ "

enforced T.C: $O(n)$ \rightarrow we will cover this in

2nd class of string.

Q) Search for a given Pattern P in Text T .
↳ KMP

ex: $T_n = b c a c d a \rightarrow \text{true}$
 $P_k = a c d a$

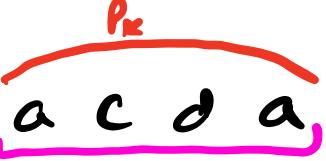
1/idea 1

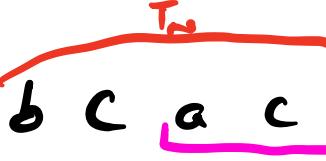
$T_n = b c a c d a$
 $P_k = a c d a$

T.C: $O(n+k) = k = n/2 \rightarrow O(n^2)$

1/idea 2

$T_n = b c a c d a$
 $P_k = a c d a$

$S:$ 
 $LPS[i]:$ 

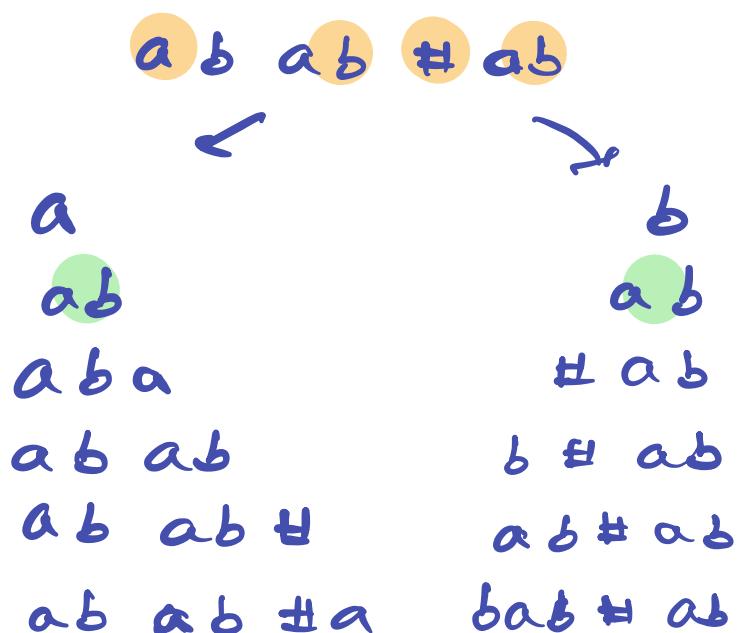
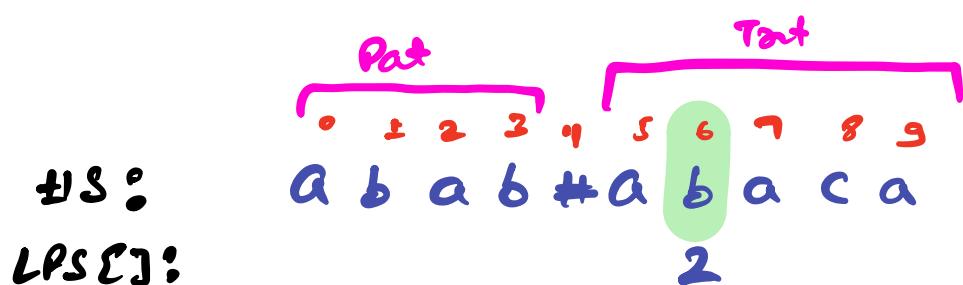
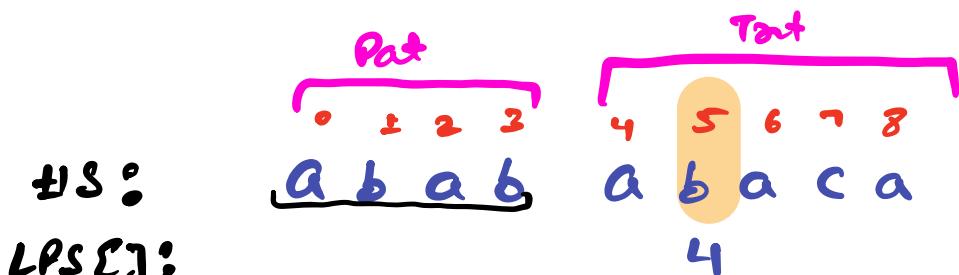
T_n 

If $LPS[i] == P.length \rightarrow \text{return true.}$
T.C: $O(n+k)$

edge case

$T_m: a b a c a \rightarrow \text{false}$

$P_m: a b a b$



Steps:

↳ 1. $S \in Pat + "\#"$ + $text \rightarrow O(n+k)$

2. Calculate $LPS[i]$ of $S[i]$ $\rightarrow O(n+k)$

3. if $LPS[i] == Pat.length$ { return true; }
⋮

if none of the idm equals Pat.length
↳ return false.

T.C: $O(n+k)$

Q) No. of Substrings of given Pattern P in Text T .

Ex: $T_n : \overset{0 \ 1 \ 2 \ 3}{a \ a \ a \ a} \rightarrow \text{ans} = 3$
 $P_k : a \ a$

Ideas

$S : \overset{P_k}{\underset{0 \ 1 \ 2}{a \ a}} \ # \ \overset{T_n}{\underset{3 \ 4 \ 5 \ 6}{a \ a \ a \ a}}$
 $LPS[] : 0 \ 1 \ 0 \ 1 \ 2 \ 2 \ 2$

$$\hookrightarrow \text{ans} = 0 + 1 + 1 + 1 = 3$$

→ Count of P .length in your LPS array of S .

T.C: $O(N+K)$

Q) Given a string s_n , min characters to be added at the start of string to make entire string Palindrome.

Ex: dc a d a c d \rightarrow ans = 2

Ex: dca a b b a a c d \rightarrow ans = 3

Ex: abad a a d a b a \rightarrow ans = 4

I/ideal

$s = \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ a & b & c & b & a & d & e & f \end{matrix}$
 $\rightarrow O(n)$
isPalindrome

$Sto[0-7]$ NO

$Sto[0-6]$ NO

$Sto[0-5]$ NO

$Sto[0-4]$ Yes \rightarrow ans: $n-1-i$

T.C: $N * O(N) \approx O(N^2)$

Idea 2

Palind. String

reversed

$$a b c b a \rightarrow a b c b a$$

$$\underline{a b c b a} d \rightarrow d \underline{a b c b a}$$

$$\underline{a b c b a} d e \rightarrow e d \underline{a b c b a}$$

$$\rightarrow s = \begin{matrix} & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ a & b & c & b & a & d & e & f \end{matrix}$$

$$s_1: \begin{matrix} & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ a & b & c & b & a & d & e & f \end{matrix} \# \begin{matrix} & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ f & e & d & a & b & c & b & a \end{matrix}$$

$$LPS \text{ of } s_1 = s$$

$\rightarrow LPS$ of whole string = length of longest possible palindrome starting from 0th idr.

$$\text{Ans} = \text{length} - LPS \text{ of whole string}$$
$$8 - 5 = 3$$

T.C: $O(2n) \approx O(n)$

→ String S

$S = "Hello"$

$S = S + "e" \rightarrow O(n)$

→ If add N strings one by one.
↳ $O(n^2)$

e.g.: String S: "

```
for (int i=0; i<n; i++) {  
    S = S + "I";  
}
```

→ $O(n^2)$

→ StringBuilder

↳ `StringBuilder sb = new StringBuilder();`

`sb = "Hello";`] → $O(1)$
`sb.append("e");`]

↳ `String S = "Hello";`

$O(n)$ [↳ `StringBuilder sb = new StringBuilder(s);`
[`for (int i=0; i<n; i++) {`
[`sb.append("I");`]]] → $O(n)$
String S = sb.toString()

Pattern Searching

Java Code:

```
class Solution {  
    int search(String text, String pat) {  
        String s1 = pat+"#" +text;  
        int[] lps = LPS(s1);  
        for(int i = 0; i<s1.length(); i++){  
            if(lps[i] == pat.length()) return 1;  
        }  
        return 0;  
    }  
    public int[] LPS(String s){  
        int n = s.length();  
        int[] lps = new int[n];  
        for(int i = 1; i<n; i++){  
            int len = lps[i-1];  
            if(s.charAt(i) == s.charAt(len)){  
                lps[i] = len +1;  
            } else {  
  
                while(s.charAt(i) != s.charAt(len)){  
                    if(len == 0){  
                        len = -1;  
                        break;  
                    }  
                    len = lps[len - 1];  
                }  
                lps[i] = len+1;  
            }  
        }  
        return lps;  
    }  
};
```

C++ Code:

```
#include <string>  
#include <vector>
```

```

class Solution {
public:
    int search(std::string text, std::string pat) {
        std::string s1 = pat + "#" + text;
        std::vector<int> lps = LPS(s1);
        for (int i = 0; i < s1.length(); i++) {
            if (lps[i] == pat.length()) return 1;
        }
        return 0;
    }

    std::vector<int> LPS(std::string s) {
        int n = s.length();
        std::vector<int> lps(n);
        for (int i = 1; i < n; i++) {
            int len = lps[i - 1];
            if (s[i] == s[len]) {
                lps[i] = len + 1;
            } else {
                while (s[i] != s[len]) {
                    if (len == 0) {
                        len = -1;
                        break;
                    }
                    len = lps[len - 1];
                }
                lps[i] = len + 1;
            }
        }
        return lps;
    }
};

```

Python Code:

```

class Solution:
    def search(self, text: str, pat: str) -> int:
        s1 = pat + "#" + text
        lps = self.LPS(s1)
        for i in range(len(s1)):
            if lps[i] == len(pat):
                return 1
        return 0

    def LPS(self, s: str) -> list:
        n = len(s)

```

```

lps = [0] * n
for i in range(1, n):
    len_ = lps[i - 1]
    if s[i] == s[len_]:
        lps[i] = len_ + 1
    else:
        while s[i] != s[len_]:
            if len_ == 0:
                len_ = -1
                break
            len_ = lps[len_ - 1]
        lps[i] = len_ + 1
return lps

```

Minimum characters to be added at front to make string palindrome

Java Code:

```

class Solution {
    public static int minChar(String str) {
        StringBuilder sb = new StringBuilder(str);
        sb.append("#");
        for(int i = str.length()-1; i>=0; i--){
            sb.append(str.charAt(i));
        }
        String s1 = sb.toString();
        int lps = LPS(s1);

        return str.length()-lps;
    }
    public static int LPS(String s){
        int n = s.length();
        int[] lps = new int[n];
        for(int i = 1; i<n; i++){
            int len = lps[i-1];
            if(s.charAt(i) == s.charAt(len)){
                lps[i] = len +1;
            } else {

                while(s.charAt(i) != s.charAt(len)){
                    if(len == 0){

```

```

        len = -1;
        break;
    }
    len = lps[len - 1];
}
lps[i] = len+1;
}

}

return lps[n-1];
}
}

```

C++ Code:

```

#include <string>
#include <vector>

class Solution {
public:
    static int minChar(std::string str) {
        std::string s1 = str + "#";
        for (int i = str.length() - 1; i >= 0; i--) {
            s1 += str[i];
        }
        int lps = LPS(s1);
        return str.length() - lps;
    }

    static int LPS(std::string s) {
        int n = s.length();
        std::vector<int> lps(n);
        for (int i = 1; i < n; i++) {
            int len = lps[i - 1];
            if (s[i] == s[len]) {
                lps[i] = len + 1;
            } else {
                while (s[i] != s[len]) {
                    if (len == 0) {
                        len = -1;
                        break;
                    }
                    len = lps[len - 1];
                }
                lps[i] = len + 1;
            }
        }
    }
}

```

```

    }
    return lps[n - 1];
}
};

```

Python Code:

```

class Solution:
    @staticmethod
    def minChar(str_: str) -> int:
        s1 = str_ + "#"
        for i in range(len(str_) - 1, -1, -1):
            s1 += str_[i]
        lps = Solution.LPS(s1)
        return len(str_) - lps

    @staticmethod
    def LPS(s: str) -> int:
        n = len(s)
        lps = [0] * n
        for i in range(1, n):
            len_ = lps[i - 1]
            if s[i] == s[len_]:
                lps[i] = len_ + 1
            else:
                while s[i] != s[len_]:
                    if len_ == 0:
                        len_ = -1
                        break
                    len_ = lps[len_ - 1]
                lps[i] = len_ + 1
        return lps[-1]

```



Today's agenda

↳ LPS[] construction

↳ multiply 2 strings



AlgoPrep



↳ calculate LPS

ex: $S = \underline{\underline{b\ c\ a\ d\ c\ b\ c\ a\ d}}$

OBS: The LPS value at the $(i+1)$ th index can never be more than 1 compared to i th index.

↳ The minimum LPS value can be 0.

$S = \underline{\underline{b\ c\ a\ d\ c\ b\ c\ a\ d\ c}}$

II General

$S_n = [S_0 | S_1 | S_2 | \dots | S_{n-1} | S_n | \dots | S_{i-5} | S_{i-4} | S_{i-3} | S_{i-2} | S_{i-1} | S_i]$

if $(S_i = S_n)$ {
 $LPS[i] = LPS[i-1] + 1$
 ↗ n



6

 $s =$

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
c	a	c	y	c	a	c	a	b	c	a	c	y	c	a	c	y
0	0	1	0	1	2	3	2	0	1	2	3	4	5	6	7	8

 $i = 16 \quad \text{len} = 7$

ans

$$S[i] = S[\text{len}]$$

$$8 \quad S[16] = S[7] \rightarrow \text{not same} \rightarrow LPS[i:j] \neq 8$$

$$4 \quad S[16] = S[3]$$

4 matches $\rightarrow LPS[i:j] = 4$

8 length $\times \times$

7 length

6 length

5 length

4 length



prefix

$i = 23$	$len = 11$
S: a b c a b d a b c a b e a b c a b d a b c a b c	0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23

$i = 23$

$len = 11$

len

ans

$s[i] == s[len]$

11

12

$s[23] == s[11] \rightarrow \text{not matching}$

5

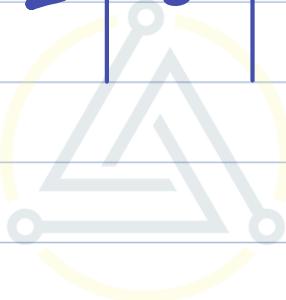
6

$s[23] == s[5] \rightarrow \text{not matching}$

2

3

$s[23] == s[2] \rightarrow \text{matching} \Rightarrow LPS[i] = 3$



AlgoPrep



II Pseudo Code

```
int[] LPS (String s[n]) {  
    int LPS[n];
```

LPS[0] = 0;

```
for (int i=1; i<n; i++) {
```

```
    int len = LPS[i-1];
```

```
    if (s[i] == s[len]) {
```

LPS[i] = len + 1;

}

```
else {
```

```
    while (s[i] != s[len]) {
```

```
        if (len == 0) { len = -1; break; }
```

len = LPS[len-1];

LPS[i] = len + 1;

}

return LPS;

}

}

T.C: O(n)

S.C: O(1)/O(n)



```

for (int i=0; i<n; i++) {
    for (int j=0; j<i; j++) {
        → O(n^2)
    }
}

```

i	j	Count
0	--	0
1	[0,0]	1
2	[0,1]	2
3	[0,2]	3
.	.	.
N-1	[0, n-1]	N-1
$\frac{n + (n-1)}{2}$		$= O(n^2)$

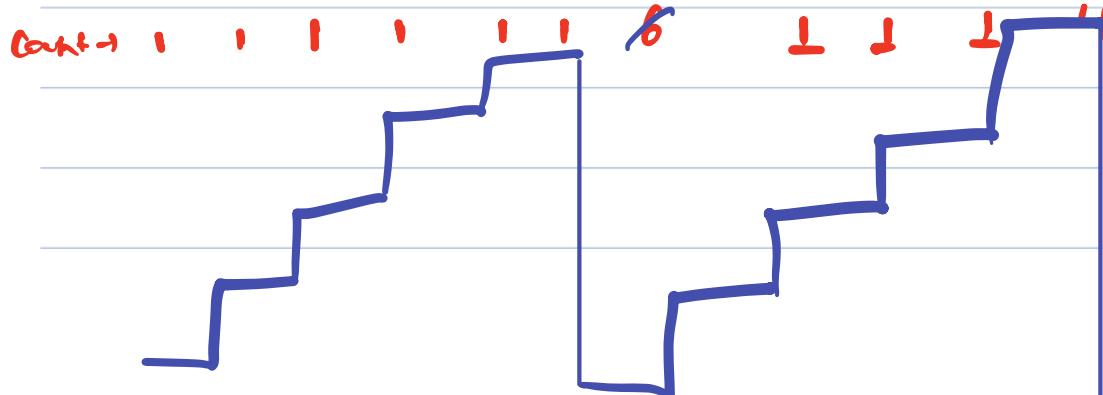
```

for (int i=1; i<n; i++) {
    int len = LPS[i-1];
    if (s[i] == s[len]) {
        LPS[i] = len+1;
    } else {
        while (s[i] != s[len]) {
            if (len==0) { len=-1; break; }
            len = LPS[len-1];
        }
        LPS[i] = len+1;
    }
}

```

i	j	Count
0	--	0
1	--	1
2	--	2/2
3	--	1/2/3
.	.	.
N	--	-

$LPS \rightarrow 0 \perp 2 \ 3 \ 4 \ 5 \ [6] \perp \ 1 \ 2 \ 3 \ [4] \Rightarrow O(n)$





S:	a	b	c	a	b	d	a	b	c	a	b	e
	0	0	0	1	2	0	1	2	3	4	5	0

```
for(int i=1; i<n; i++) {
    int len = LPS[i-1];
    if (s[i] == s[len]) {
        LPS[i] = len+1;
    }
    else {
        while (s[i] != s[len]) {
            if (len==0) {len=-1; break;}
            len = LPS[len-1];
        }
        LPS[i] = len+1;
    }
}
```

i	1	len	0-1
	2		0-1
	3		0
	4		1
	5		2 → 0-1
	6		0
	7		1
	8		2
	9		3
	10		4
	11		5 → 2 → 0-1