# TP QPE and Shor

#### Valeran MAYTIE

### 1 Small Practice

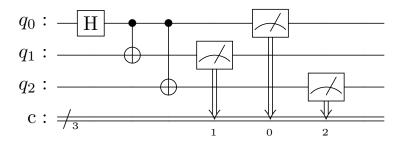


Figure 1: Circuit that calculate  $\frac{1}{\sqrt{2}}(|000\rangle + |111\rangle)$ 

To compute  $\frac{1}{\sqrt{2}}(|000\rangle + |111\rangle)$ , we created the circuit shown in Figure 1. Firstly we apply an Hadamard Gate on the first qubit to have a 50/50 chance of having it at 1 or 0. If it's equal to one then with a *CNot* gate controlled by the first qubit we inverse  $q_1$  and  $q_2$ , so they are all equal to 1. Else nothing change and it remains at 0.

When you run it, you'll find the right measurements the result is roughly 50/50, 000 or 111. The code for this exercise is shown in Listing 1

```
1 q = QuantumRegister(3)
2 c = ClassicalRegister(3)
3 qc = QuantumCircuit(q,c)
4
5 qc.h(q[0])
6
7 qc.cnot(q[0],q[1])
8 qc.cnot(q[0],q[2])
9
10 qc.measure(q, c)
```

Listing 1: Code that generates the circuit in Figure 1

## 2 QPE

We construct the operator  ${\bf U}$  with this matrix :

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{2i\pi\frac{6}{8}} \end{pmatrix}$$

## 2.1 Math Questions

- 1. What is doing this operator ? ('2j' is in Python the complex number  $2 \cdot i$ ) The operator **U** compute this :
  - $\mathbf{U}|00\rangle = |00\rangle$
  - $\mathbf{U}|01\rangle = |01\rangle$
  - $\mathbf{U}|10\rangle = |10\rangle$
  - $\mathbf{U}|11\rangle = e^{2i\pi\frac{6}{8}}|11\rangle$
- 2. On how many qubits does it act? This operator act on 2 qubits.
- 3. What are its eigenvalues/eigenvectors?

eigenvectors	eigenvalues
$ 00\rangle$	1
$ 01\rangle$	1
$ 10\rangle$	1
$ 11\rangle$	$e^{2i\pi\frac{6}{8}}$

4. For each eigenvector, what should QPE return with 3 bits of precisions, as seen in the course?

eigenvectors	QPE return
$ 00\rangle$	000
$ 01\rangle$	000
$ 10\rangle$	000
$ 11\rangle$	110

2

### 2.2 Implementing QPE

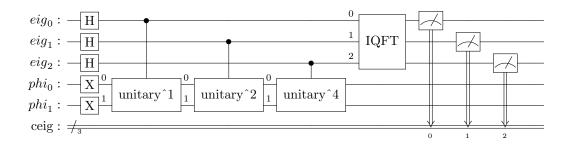


Figure 2: QPE circuit

```
1 eig = QuantumRegister(size_eig, name="eig")
2 phi = QuantumRegister(size_phi, name="phi")
3 ceig = ClassicalRegister(size_eig, name="ceig")
4 qc = QuantumCircuit(eig,phi,ceig)
5
6 qc.x(phi[0])
7 qc.x(phi[1])
8 for i in range(0, size_eig):
9 qc.h(eig[i])
10 qc.append(U.power(2**i).control(), [eig[i], phi[0], phi[1]])
11
12 qc.append(QFT(size_eig).inverse(), eig)
13 qc.measure(eig, ceig)
```

Listing 2: Code that generate the QPE circuit

#### 2.3 Exact result

1. Is it the expected result?

```
Yes, we calculated 110 \equiv \frac{1}{2} + \frac{1}{4} = \frac{6}{8} so \mathbf{U}|11\rangle = e^{2i\pi\theta}|11\rangle.
```

With 3 bits precision  $\theta$  is equal to  $\frac{6}{8}$  seen in Exercise 2.1.1.

- 2. Change the  $\frac{6}{8}$  of the phase of U: use  $\frac{1}{8}$ , then  $\frac{2}{8}$ ... Is QPE returning the correct answer? Yes, we have 001, 010, ... and 111, when we tested up to  $\frac{7}{8}$ , because there is enough precision. But then we get the right result modulo 8.
- 3. Change the precision : use 4 qubits for "eig", and change the fraction in the phase of **U** to  $\frac{10}{16}$  : is QPE indeed returning 10 in binary ?

We have 10 written in binary. It works because we have enough precision to get the real eigenvalues.

4. Move to 5 bits of precision is it still working?

It works, we have  $\frac{1}{2} + \frac{1}{8} = \frac{10}{16}$ 

#### 2.4 Approximate result

The QPE approach calculations with 3-bits precision are given in the table below

value	number of times obtained
000	21
001	34
010	178
011	708
100	39
101	20
110	10
111	14
average / 2 <sup>3</sup>	0.365

We can see that the average divided by two power of precision is close to the eigenvalue  $(\frac{1}{3})$ . And the more you increase the accuracy, the closer you get to  $\frac{1}{3}$ .

#### 2.5 Superposition

By changing the phi initialization to  $\frac{1}{\sqrt{2}}(|\phi_1\rangle + |\phi_2\rangle)$  Two calculations are performed in parallel, one to calculate the eigenvalue of eigenvectors  $|11\rangle$  and  $|00\rangle$ .

For the phase it works very well we have:

$$\begin{array}{c|c|c} phi & eig & \\ \hline 00 & 000 & 498 \\ 11 & 011 & 526 \\ \end{array}$$

## 3 Implementing Shor's algorithm

#### 3.1 Oracle synthesis

$$Mult_{a^p \mod N} : x \mapsto \left\{ \begin{array}{ll} (a^p \cdot x) \mod N & \text{si } x < N \\ x & \text{si} N \le x < 2^n \end{array} \right.$$

n	$_{ m time}$	circuit size (gates)
3	28.2  ms	61
4	$120~\mathrm{ms}$	296
5	$527~\mathrm{ms}$	1341
6	$2.46 \mathrm{\ s}$	5633
7	$11.5 \mathrm{\ s}$	23044

Figure 3: Generation of  $gateMult(3, 3, 2^n, n)$ 

- 1. What are the sizes of the generated circuits? It is written on Figure 3.1.
- 2. What is the complexity of the circuit size in term of number of qubits? It is exponential.
- 3. Can you explain why?

The matrix as an exponential size in function of the number of qubits  $(2^n*2^n)$ , so this size has repercussions on the final circuit.

4. What alternate method could you suggest, with what potential drawbacks?

You can encode certain types of number, such as powers of two. However, we lose expressiveness.

#### 3.2 Plugging everything together : Shor

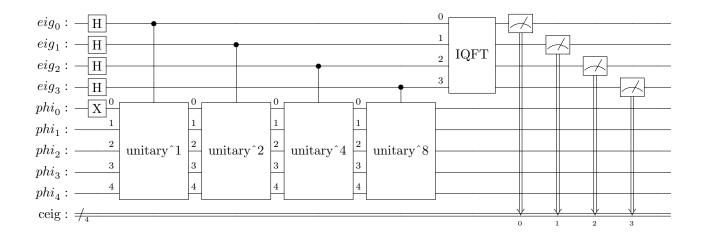


Figure 4: Shor algorithm

- 1. What is the order r of  $a \mod N$  (here  $7 \mod 30$ )? The order r of  $7 \mod 30$  is 4.
- 2. On the drawing, where are we supposed to see the values  $\frac{s}{r}$ ? The horizontal axis is graded with integers... To what real numbers between 0 and 1 these correspond to?
- 3. Can you infer from the graph the value of r? Where do you see it on the graph? Yes r = 4, it is presented by the period of the keys.
- 4. Change a and N respectively to 20 and 29. Can you read the value r ? Is it correct ? No I can't read the value r
- 5. The drawing is not very precise... How to make it better? Try it!

  We need to increase the precision. But the calculation is very expensive.
- 6. Is it still working if you change the value of 'a' and/or 'N' to other values? Beware not to use too large values for 'N'... To get some inspiration, below is the list of possibilities up to 31.

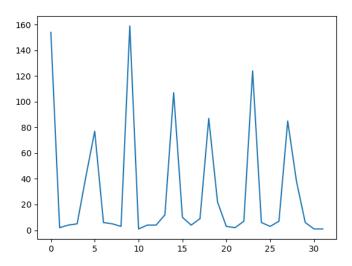


Figure 5: Result of Shor execution