

Graph Algorithms

TD1 : Introduction

Valeran MAYTIE

1 To begin

Exercise 1.1 – Show that a graph always has an even number of odd degree vertices

Correction 1.1

Let G a graph. Thanks to *Handshaking Lemma* we have:

$$\sum_{v \in V(G)} \deg(v) = 2|G(E)|$$

So the number $\sum_{v \in V(G)} \deg(v)$ is an even number. To keep this property we must have an even number of odd degree vertices otherwise the sum become odd.

Exercise 1.2 – Show that a graph with at least 2 vertices contains 2 vertices of equal degree

Correction 1.2

Let G a graph with at least 2 vertices.

- If G has no isolated vertex, the degree of a vertex is between 1 and $|V(G)| - 1$ (v a vertex $1 \leq \deg(v) < |V(G)|$). So we have $n := |V(G)| - 1$ different values for the degree of a vertex in G . The graph G contains $n + 1$ vertices so by the Pigeonhole principle we have two vertices with the same degree.
- If G has one isolated vertex, it's the same idea that the previous one but the degree of a vertex is between 0 and $|V(G)| - 2$.
- If G has at least two isolated vertices, we have 2 vertices with the same degree.

Exercise 1.3 – Let G be a graph of minimum degree $\delta(G) \leq 2$. Show that G contains a cycle.

Correction 1.3

Let G a graph with $\delta(G) \leq 2$.

Show that G has a cycle. Let $P = v_1 \dots v_l$ be the maximum path in G . Since it cannot be extended into a larger path, we must have $N(v_l) \subseteq V(P)$. So $\exists i \leq l - 1$ $v_i \in N(v_l)$ which yields that v_i, \dots, v_l is a cycle.

Exercise 1.4 – Let G be a graph of minimum degree d , and of girth $2t + 1$. Given any vertex $v \in V(G)$, show that there are at least $d(d - 1)^{i-1}$ vertices at distance exactly i from v in G , for every $1 \leq i \leq t$. Deduce a lower bound on the number of vertices of G .

Correction 1.4

2 Dense subgraphs

Exercice 2.1 – Show that every graph of average degree d contains a subgraph of minimum degree at least $\frac{d}{2}$.

Correction 2.1

Exercice 2.2 – Can you find a similar relation between the maximum degree and the minimum degree ? And between the maximum degree and the average degree ?

Exercice 2.3 – Show that every graph of average degree d contains a bipartite subgraph of average degree at least $\frac{d}{2}$.

3 Cuts and trees

Exercice 3.1 – If G is connected, and $e = uv$ is a bridge in G , how many connected components does $G \setminus e$ contain ? Show that u and v are cut-vertices.

Exercice 3.2 – Show that a graph G is a tree if and only if there exists a unique path from u to v in G , for every pair of vertices $u, v \in G$.

Exercice 3.3 – Let T a BFS tree of a graph G . Show that every edge of G is contained either within a layer of T , or between two consecutive layers of T .

Exercice 3.4 – Let T be a DFS tree of a graph G . Show that, for every edge $e \in E(G)$, there is a branch of T that contains both extremities of e .