Lambda Calculus and category theory

Valeran MAYTIE

Contents

1 Introduction

Boole :

- If you consider propositions (no quantifiers) of classical logic : $A ::= P|A \wedge B|\neg A|A \wedge B|\top|\bot$
- Ordered by logical implication $A \leq B \Leftrightarrow A \Rightarrow B$, A implies B or $A \vdash B$

Observation $A \wedge B \leq A, A \wedge B \leq B$. moreover if $C \leq A$ and $C \leq B$ then $C \leq A \wedge B$ (for all proprieties) Which means that $A \wedge B$ define a infimum of A and B greatest lower bound glb

Definition – $A \Rightarrow B = (\neg A) \vee B = \neg(A \wedge \neg B)$.

Observation :

- $A \wedge (A \Rightarrow B) \leq B$
- $A \vee \neg A \leq$ true
- $A \wedge \neg A \geq$ false

Frege : ideography (first proof system)

The idea that a mathematical proof is a mathematical object. In particular there may be different proofs of a proposition A formula.

$$\begin{tikzcd} \text{B} \\ \text{A} \arrow[u, bend left, "\pi_1"] \arrow[u, bend right, "\pi_2"'] \arrow[u, phantom, "\neq"] \end{tikzcd} \qquad\qquad \begin{tikzcd} \text{B} \\ \text{A} \arrow[u, dash, "\leq"'] \end{tikzcd}$$

Lambek Lambek

Lambek understood connection between.

$$\begin{tikzcd} \text{Intuitionists minimal propositional logic} \arrow[rr, leftrightarrow, "\text{Curry, Howard, Lambek}"] \arrow[dr, leftrightarrow, "\text{Curry, Howard}"'] & & \text{Cartesian closed category} \arrow[dl, leftrightarrow, "\text{Lambek}"] \\ & \text{Simply typed lambda calculus} & \end{tikzcd}$$

Definition – A monoid $(M, \bullet, e)$ is a set M equipped with a binary operation $\bullet : M \times M \to M$ with a neutral element $e \in M_e : M^0 \to M$ satisfying two equations :

- (associativity) $\forall x, y, z \in M, x \bullet (y \bullet z) = (x \bullet y) \bullet z$
- (neutrality) $\forall x, \in M, x \bullet e = x = e \bullet x$

Example – $(\mathbb{N}, +, 0), (\mathbb{Z}, +, 0), (\mathbb{N}, \times, 1)$ and any group.
Free monoide on a set (=alphabet) A A^* contains finite sequences of element A $w = [a_1 \dots a_n]$

- multiplication is concatenation
- neutral element is empty word

2 Categories, functors, natural transformations

Definition – A category $\mathcal{C}$ is a graph

- whose node are called objects
- whose edges are called morphism/maps/arrow.

The objects of $\mathcal{C}$ form a class of objects. Every pair of objects A, B
Every pair of object A, B comes with a set $Hom(A, B)$ of morphisms $A \xrightarrow{f} B$, $f \in Home(A, B)$
The graph is equipped with :

- an morphisms $id_A \in Home(A, A)$ for all object A of $\mathcal{C}$
- a composition defined as a function $\circ_{A,B,C} : Hom(B, C) \times Hom(A, B) \to Hom(A, C)$ for every objects A, B, C of category $\mathcal{C}$

  It satisfying the following equation :

  - associativity :

$$\begin{tikzcd} & B \arrow[r, "g"] & C \arrow[dr, "h"] & \\ A \arrow[ur, "f"] \arrow[urr, "g \circ f"] \arrow[rrr, "h \circ g \circ f"'] & & & D \end{tikzcd} \qquad \begin{tikzcd} & B \arrow[r, "g"] \arrow[drr, "h \circ g"] & C \arrow[dr, "h"] & \\ A \arrow[ur, "f"] \arrow[rrr, "h \circ g \circ f"'] & & & D \end{tikzcd}$$

$$h \circ (g \circ f) = (h \circ g) \circ f$$

  - neutrality :

$$\begin{tikzcd} A \arrow[loop above, "Id_A"] \arrow[r, "f"'] & B \arrow[loop above, "Id_B"] \end{tikzcd}$$

$$Id_B \circ f = f = f \circ Id_A$$

Definition – A small category is a category whose class of object is a set. What we defined as a category is called "locally small category".

Example – Ordered Set : Claim every ordered set A defines a category

- object : elements of A
- morphisms : $a \to b \Leftrightarrow a \leq b$

$$Hom(a, b) = \begin{cases} \text{singleton} & a \leq b \\ \emptyset \end{cases}$$

The composition is defined by transitivity :

$$\begin{array}{ccccc} a & \xrightarrow[a \leq b]{} & b & \xrightarrow[b \leq c]{} & c \\ a & & \leq & & c \\ a & & \longrightarrow & & b \end{array}$$

Definition – An order category $\mathcal{C}$ is a category when $Hom(A, B)$ is a singleton for all object A, B of $\mathcal{C}$.

Observation – an order category is the some thing as a pre-order (= trans, refl).

Example – <u>monoid</u>

- A category with one object $*$, $M = Hom(*, *)$ defined a monoid
  - $\circ : Hom(*, *) \times Hom(*, *) \to Hom(*, *)$
  - $id_* \in M = Hom(*, *)$ defined the neutral element
- Conversely every monoid $M = (M, \bullet, e)$ defined a category $\mathcal{B}M$ or ΣM with on object $*$ and $Hom(*, *) = M$ composition defined $y \circ x = y \bullet x$ id_* defined the neutral element.

$$2 = 1 \circ 1$$

$$\mathcal{B}(\mathbb{N}, +, 0) \qquad \begin{tikzcd} * \arrow[loop above, "0"] \arrow[loop below, "1"] \end{tikzcd}$$

$$3 = 2 \circ 1$$