System Call Programming

Week 7

Processor Modes

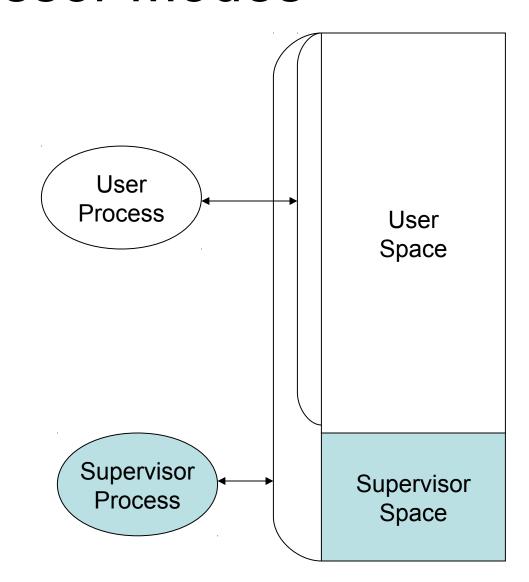
- Mode bit used to distinguish between execution on behalf of OS & execution on behalf of user.
- Supervisor mode: processor executes every instruction in its hardware repertoire.
- User mode: can only use subset of instructions

Processor Modes

- Instructions that can be executed in supervisor mode are supervisor privileges, or protection instruction
 - I/O instructions are protected. If an application needs to do I/O, it needs to get the OS to do it on its behalf.
 - Instructions that can change the protection state of the system are privileges (e.g., process' authorization status, pointers to resources, etc)

Processor Modes

 Mode bit may define areas of memory to be used when the processor is in supervisor mode vs user mode



The Kernel

- Core of OS software executing in supervisor state
- Trusted software:
 - Manages hardware resources (CPU, Memory and I/O)
 - Implements protection mechanisms that could not be changed through actions of untrusted software in user space
- System call interface is a safe way to expose privileged functionality and services of the processor

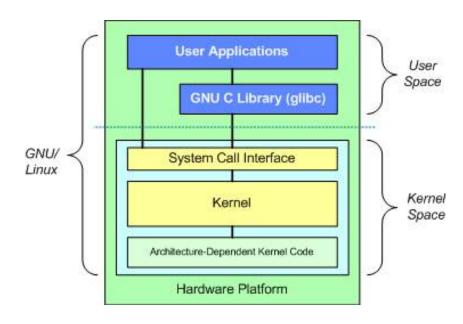
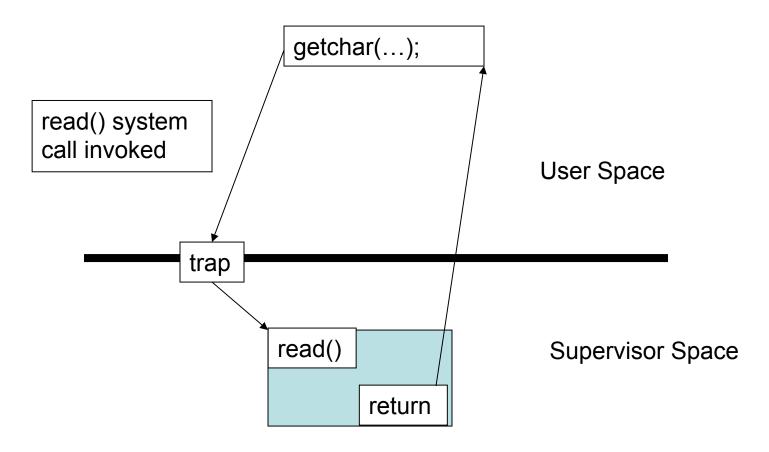


Image by: Tim Jones (IBM)

System Calls



Trap: System call causes a switch from user mode to kernel mode

System calls

- A system call involves the following
 - The system call causes a 'trap' that interrupts the execution of the user process (user mode)
 - The kernel takes control of the processor (kernel mode/privilege switch)
 - The kernel executes the system call on behalf of the user process
 - The user process gets back control of the processor (user mode/privilege switch)
- System calls have to be used with care. Expensive due to privilege switching.

System Calls

```
#include<unistd.h>
ssize t read(int fildes, void *buf, size t nbyte)

    fildes: file descriptor

    buf: buffer to write to

    nbyte: number of bytes to read

ssize t write(int fildes, const void *buf, size_t nbyte);

    fildes: file descriptor

    buf: buffer to write from

    nbyte: number of bytes to write

int open (const char *pathname, int flags, mode t mode);
int close (int fd);
File descriptors
 - 0 stdin
 1 stdout
 - 2 stderr
```

Why are these system calls and not just regular library functions?

Example System Calls

- pid_t getpid(void)
 - Returns the process ID of the calling process
- int **dup**(int fd)
 - Duplicates a file descriptor fd. Returns a second file descriptor that points to the same file table entry as fd does.
- int **fstat**(int filedes, struct stat *buf)
 - Returns information about the file with the descriptor filedes into buf

```
struct stat {
  ino t     st ino;     /* inode number */
  mode t st mode;
                 /* protection */
  nlink t st nlink; /* number of hard links */
  gid t st gid; /* group ID of owner */
  off t st size;
                 /* total size, in bytes */
  blksize t st blksize; /* blocksize for file system I/O */
  blkcnt t st blocks; /* number of 512B blocks allocated */
         st atime;
                 /* time of last access */
  time t
  time t st mtime;
                 /* time of last modification */
        st ctime;
  time t
                 /* time of last status change */
};
```

Laboratory

• Write tr2b and tr2u programs in 'C' that transliterates bytes. They take two arguments 'from' and 'to'. The programs will transliterate every byte in 'from' to corresponding byte in 'to'

```
- ./tr2b 'abcd' 'wxyz' < bigfile.txt
```

Replace 'a' with 'w', 'b' with 'x', etc

```
- ./tr2b 'mno' 'pqr' < bigfile.txt</pre>
```

- tr2b uses getchar and putchar to read from STDIN and write to STDOUT.
- tr2u uses **read** and **write** to read and write each byte, instead of using getchar and putchar. The nbyte argument should be 1 so it reads/writes a single byte at a time.
- Test it on a big file with 5000000 bytes

Laboratory

- read and write are system calls.
- getchar and putchar are library functions that live in user space but call read and write. And do buffering!
 - getchar() calls read() only if the local buffer is empty
 - getchar() reads multiple bytes into its buffer in one call to read()

Buffering Issues

- What is buffering?
- Why do we buffer?
- Can we make our buffer really big?

strace and time

• **strace**: Intercepts and prints out system calls to stderr or to an output file.

```
strace -o strace_output ./tr2u 'a' 'b' < bigfile.txt</pre>
```

• time: Timing

```
time ./tr2b 'x' 'y' < bigfile.txt
real     0m0.150s (Total time elapsed)
user     0m0.132s (time spent by CPU executing in user space)
sys     0m0.016s (time spent by CPU executing in kernel space)</pre>
```

Buffered versus Unbuffered

- How many read and write calls are there in each?
- Why does one take longer than the other?

Writing to a File vs. Terminal

Do you see any differences in the read/writes?

```
- ./tr2u 'a' 'b' < bigfile.txt (TERMINAL)
- ./tr2u 'a' 'b' < bigfile.txt > out.txt (FILE)
```

- In the buffered version, how many bytes are being read/written when outputting to a file and when outputting to the terminal?
- Why is there a difference?
- What are the performance considerations?

Homework

- Rewrite sfrob.c using system calls
- If stdin is a regular file, your program initially allocate enough memory to hold all data in the file all at once.
 - Otherwise, it should behave like before
- In addition to the functions you used last time, you'll need read, write, and fstat
 - Read the man pages for these functions
- The program should be named sfrobu.c
- Measure differences in performance between sfrobu and sfrob using the time command.
- Write a shell script sfrobs that uses standard 'tr' and 'sort' command to sort encrypted files using a pipeline (no temporary files should be created!). Use the 'time' command to compare its performance to sfrob and sfrobu.