

SEMESTER 1 2023-2024

CS370 Computation and Complexity

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Time allowed: 2 hours

Answer at least *three* questions

Your mark will be based on your best three answers

All questions carry equal marks

Instructions

	Yes	No	N/A
Formulae and Tables book allowed (i.e. available on request)	X		
Formulae and Tables book required (i.e. distributed prior to exam commencing)			
Statistics Tables and Formulae allowed (i.e. available on request)			Х
Statistics Tables and Formulae required (i.e. distributed prior to exam commencing)			Х
Dictionary allowed (supplied by the student)		Χ	
Non-programmable calculator allowed			Χ
Students required to write in and return the exam question paper		Χ	

[25 marks]

1 (a) Draw a diagram of a Turing machine and label the key elements [10 marks] it has.

The mathematical description of the Turing machine contains the following terms:

Alphabet States Initial state Accepting state

Explain what each of these terms mean

Give one difference between a Turing machine and a modern computer device?

- (b) For any three problems of your choice, describe those problems [6 marks] as sets.
- (c) Give an explanation of the difference between the **Decidability** [3 marks] and **Undecidability** of any Decision problem.
- (d) Let S_5 be the set of binary strings of length at least 5, i.e. [6 marks] $S_5 = \{x \in \{0,1\}^*: |x| \ge 5\}$
 - (I) Build a recogniser for S_5 that is not a decider for S_5 .
 - (II) Is S_5 decidable?
 - (III) If you had the option to have access to either a decider or a recogniser when solving membership of a set, which would you choose and why?

[25 marks]

2 (a) Briefly describe the Halting problem with the aid of a suitable [9 marks] graphic. (i.e. draw a block diagram that shows the inputs and outputs from a Turing machine and uses arrows to illustrate the flow of control of the program code).

If it is stated that the Halting problem is undecidable. Can it be proven by contradiction that this is the case?

(b) Consider the set $L_5 = \{(M) : M \text{ is a TM such that } |L(M)| \ge 5\}.$ [4 marks]

Describe this set in English. Then, for the following TMs whose inputs are binary strings, state whether they are elements of L_5 or not:

- i. $M_1(x) =$ "if $|x| \le 4$ accept, else reject."
- ii. $M_2(x)$ = "if x = 00000 accept, else reject.
- iii. $M_3(x)$ = "loop."
- (c) In computing, what does the reduction of a problem mean? [6 marks] Support your answer using two different examples.
- (d) Consider the following sets:

[6 marks]

- (a) $A_{TM} = \{ \langle M, w \rangle : M \text{ is a TM and } w \text{ a string so that } M(w) = acc \}$
- (b) $Rej_{010} = \{\langle M, w \rangle : M \text{ is a TM that rejects input 010.} \}$

Perform the following Turing reductions

- (a) $A_{TM} \leq_T Halt_{TM}$.
- (b) $A_{TM} \leq_T \text{Rej}_{010}$

[25 marks] [5 marks]

- **3** (a) Define what an oracle machine is and how oracle machine can be used to solve decision problems.
 - (b) State Rice's theorem and explain why it is important. [5 marks]
 - (c) State Why Rice's theory cannot be applied to the following sets: [9 marks]
 - (a) $A_{TM} = \{\langle M, w \rangle : M \text{ is a TM and } w \text{ a string so that } M(w) = acc\}$ (b) $L = \{\langle M, w \rangle : M \text{ is a TM such that } M \text{ decides } Halt_{TM}\}$ (c) $L = \{\langle M, w \rangle : M \text{ is a TM such that } M' \text{s first two tape cells are } M' \text{ such that } M' \text{ so that$
 - (c) $L = \{\langle M, w \rangle : M \text{ is a TM such that } M \text{'s first two tape cells are read only once} \}$
 - (d) An interesting question about Turing machines is whether they can [6 marks] reproduce themselves. A Turing machine cannot be defined in terms of itself, but can it still somehow print its own source code? With reference to the recursion theorem, explain why this is true.

[25 marks]

4 (a) Write definitions for the following terms:

[10 marks]

- (I) PTime
- (II) EXPtime
- (III) BPPtime
- (IV) PSPACE
- (V) EXPSPACE
- (b) Whether P=NP or P!=NP is something that appears frequently in [8 marks] complexity theory. Explain what this means for the problems that we are attempting to solve using computers?
- (c) Give the details behind two complex problems where the time [7 marks] required to solve them could become infeasible. Use suitable diagrams to illustrate your answer.