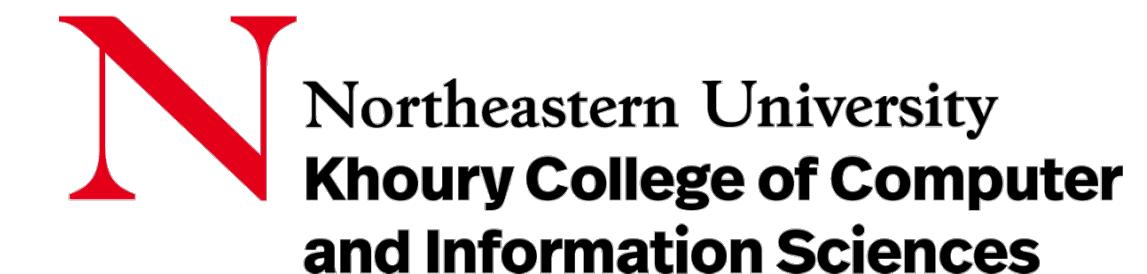


Framing Frames: Bypassing Wi-Fi Encryption by Manipulating Transmit Queues

Domien Schepers, *Aanjhan Ranganathan*, Mathy Vanhoef

WAC6 (colocated with CRYPTO 2023)



KU LEUVEN

Signal Intelligence Lab @ Northeastern

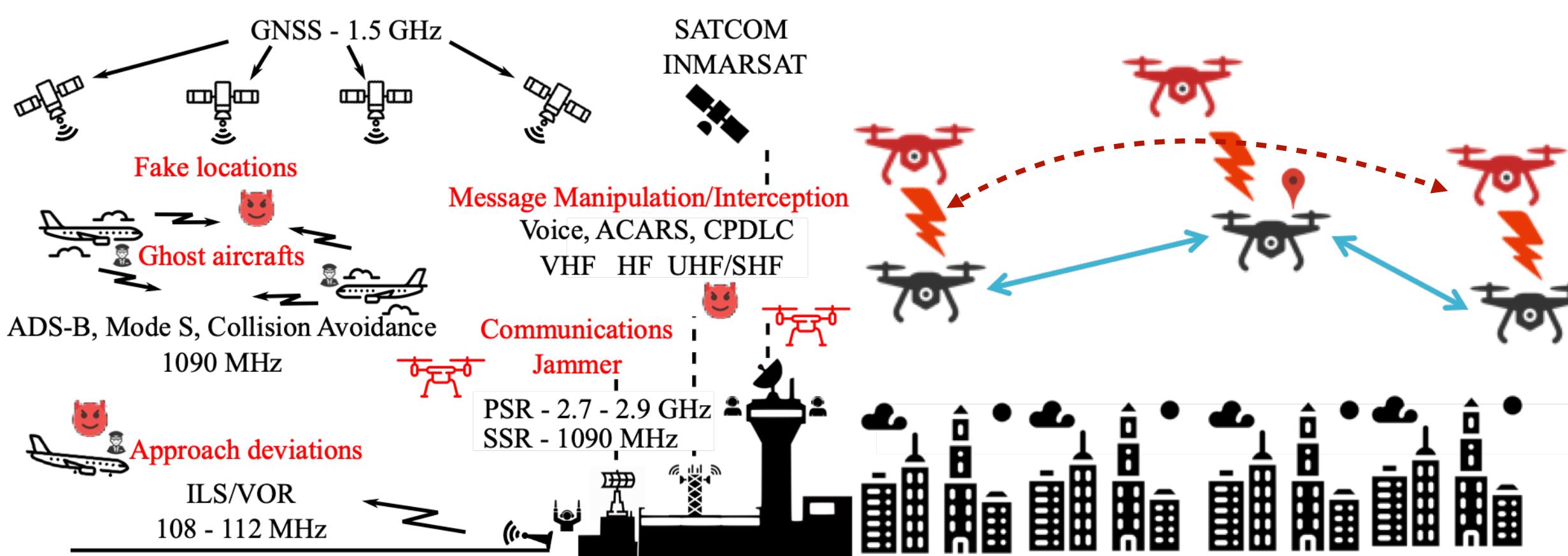
armasuisse
Schweizerische Eidgenossenschaft
Confédération suisse
Confederazione Svizzera
Confederaziun svizra



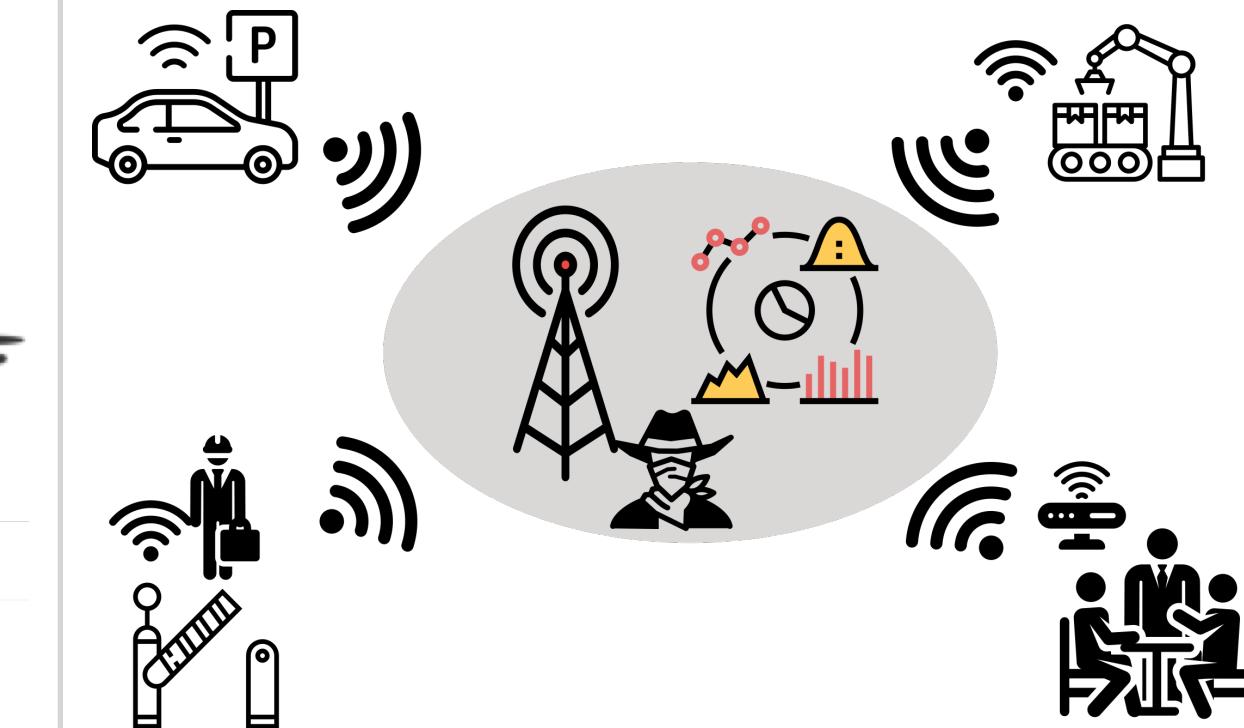
Security and privacy of wireless networks with a focus on
autonomous cyber-physical systems and smart ecosystems.



**Secure and Private
Wide-area Positioning**



Aviation and Aerospace Security

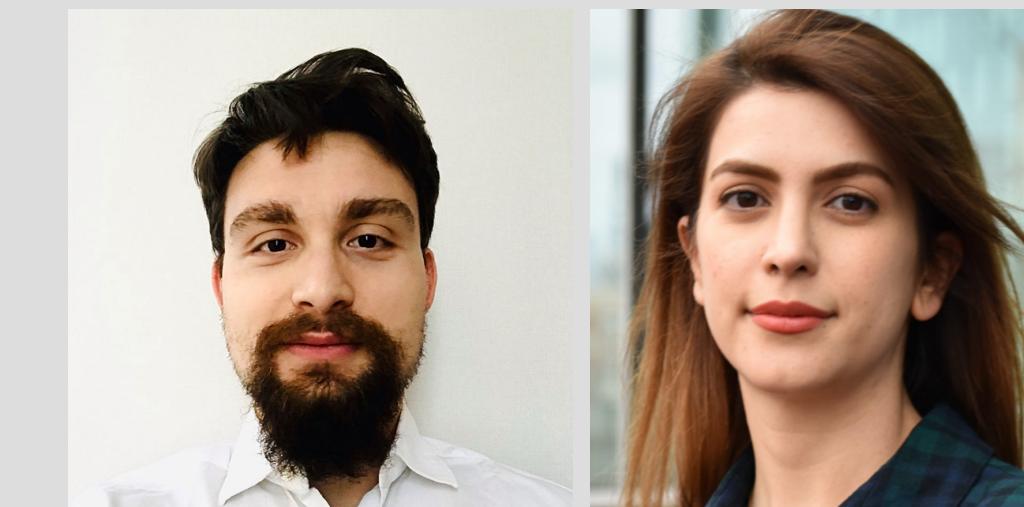


**Security and Privacy
of xIoT**

Faculty



Aanjhan Ranganathan
Assistant Professor
www.aanjhan.com



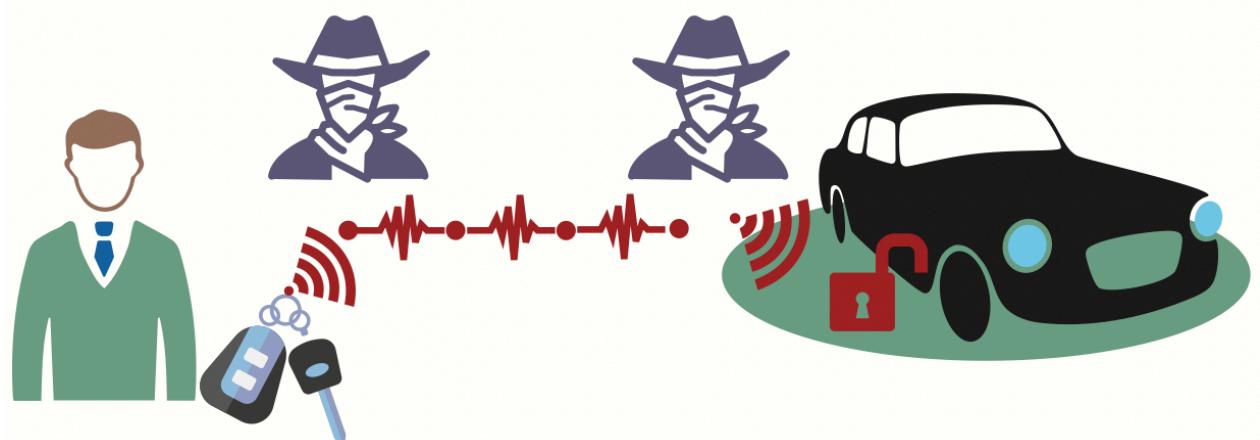
PhD Students



Secure Proximity and Location Verification

Towards Secure and Private Wide-area Positioning

Attacks on Location



The Telegraph

Home Video News World Sport Business Money Comment Culture Travel Life
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HOME > TECHNOLOGY > TECHNOLOGY NEWS
Researchers commandeer £50m superyacht with GPS-spoofing



© 29 June 2012 | Technology



Researchers use spoofing to 'hack' into a flying drone

Selected Research

An Experimental Study of GPS Spoofing and Takeover Attacks on UAVs, Harshad Sathaye, Martin Strohmeier, Vincent Lenders, Aanjan Ranganathan (USENIX Security 2022)

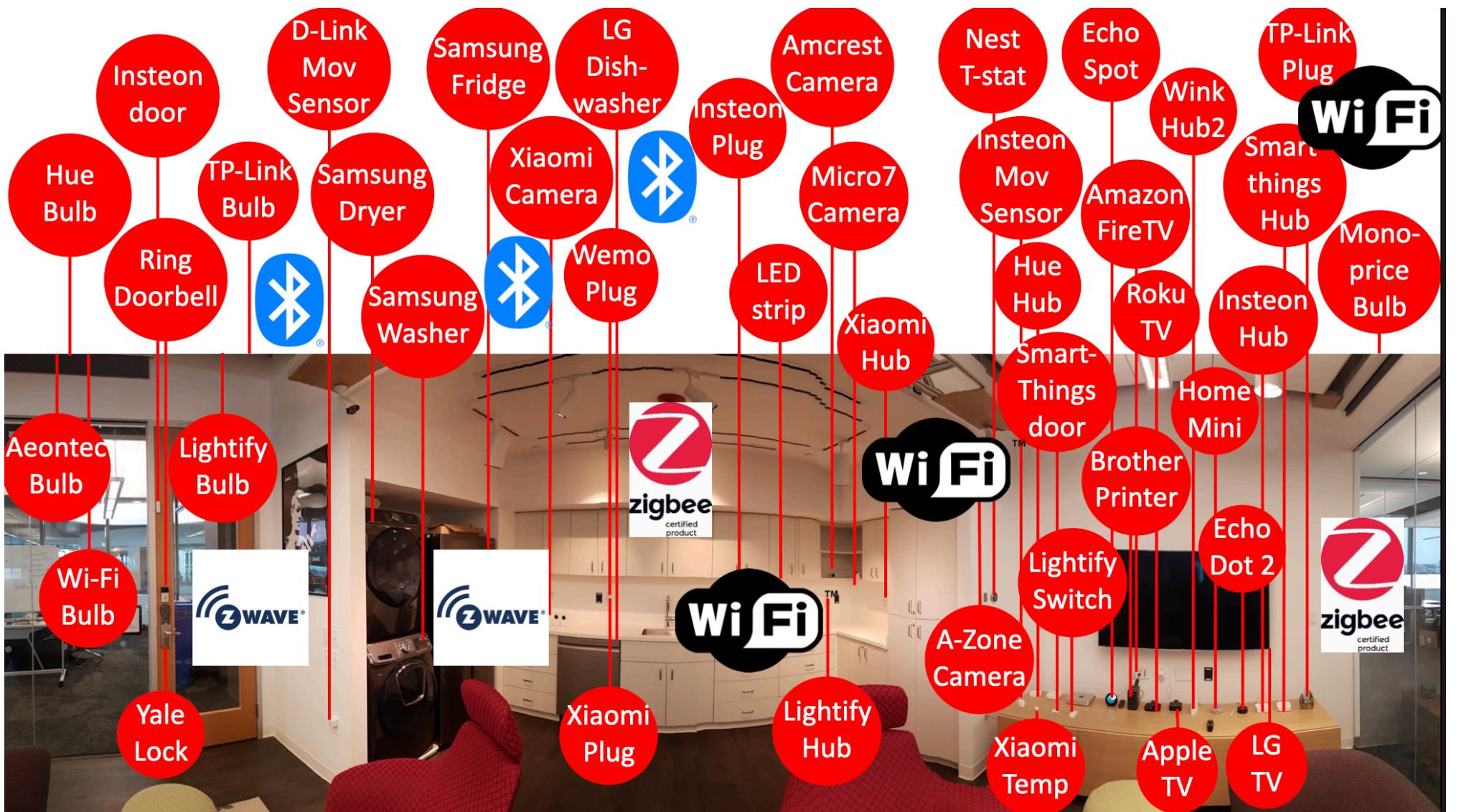
VRange: Enabling Secure Ranging in 5G-NR Wireless Networks, Mridula Singh, Marc Roeschlin, Aanjan Ranganathan, Srdjan Capkun (NDSS 2022)

SemperFi: Anti-spoofing GPS receiver for UAVs, Harshad Sathaye, Gerald LaMountain, Pau Closas, Aanjan Ranganathan (NDSS 2022)

Wireless Attacks on Aircraft Instrument Landing Systems, Harshad Sathaye, Domien Schepers, Aanjan Ranganathan, Guevara Noubir (USENIX Security 2019)

Security and Privacy in xIoT

Validating and Building Trustworthy Smart Ecosystems



Mon(IoT)Or Lab at Northeastern University

Selected Research

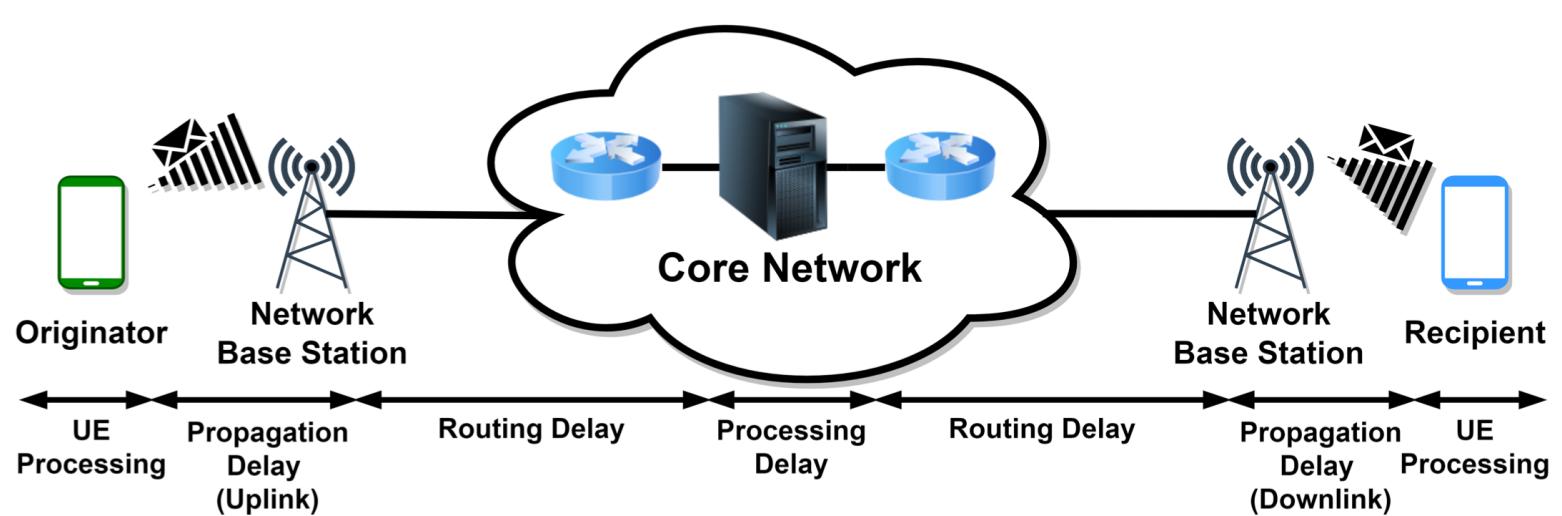
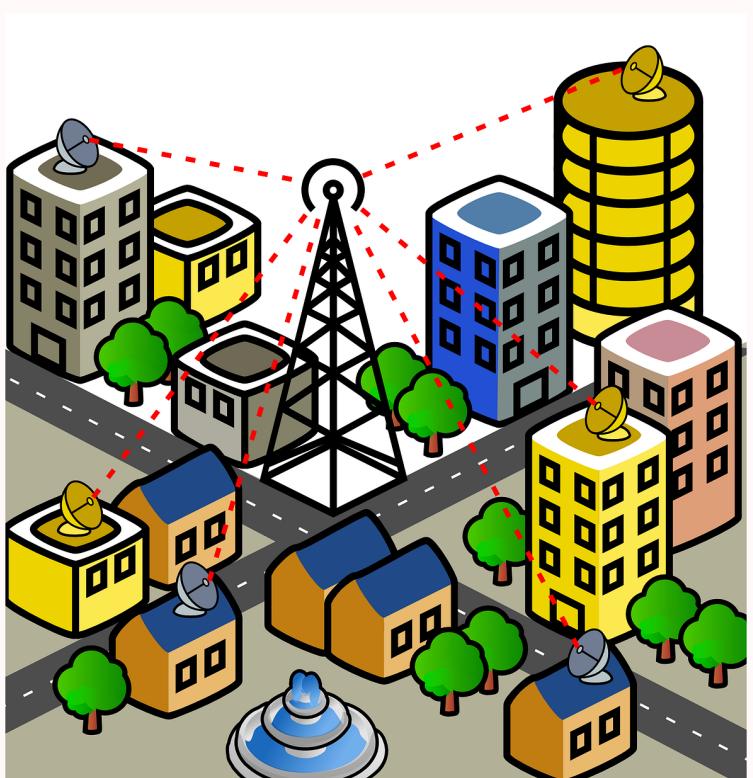
Track You: A Deep Dive into Safety Alerts for Apple AirTags,
Narmeen Shafqat, Nicole Gerzon, Maggie Von Nortwick, Victor
Sun, Alan Mislove, Aanjan Ranganathan (PETS 2023)

**ZLeaks: Passive Inference Attacks on Zigbee based Smart
Homes,** Narmeen Shafqat, Daniel Dubois, Dave Choffnes, Aaron
Schulman, Dinesh Bharadia, Aanjan Ranganathan (ACNS 2022,
Best Student Paper Award)

**Privacy-Preserving Positioning in Wi-Fi Fine Timing
Measurements,** Domien Schepers, Aanjan Ranganathan
(PETS 2022)

**I Send, Therefore I Leak: Information Leakage in Low-Power
Wide Area Networks,** Patrick Leu, Ivan Puddu, Aanjan
Ranganathan, Srdjan Capkun (WiSec 2018)

Wi-Fi and Cellular Security



Selected Research

Framing Frames: Bypassing Wi-Fi Encryption by Manipulating Transmit Queues

Domien Schepers, Aanjhan Ranganathan, Mathy Vanhoef
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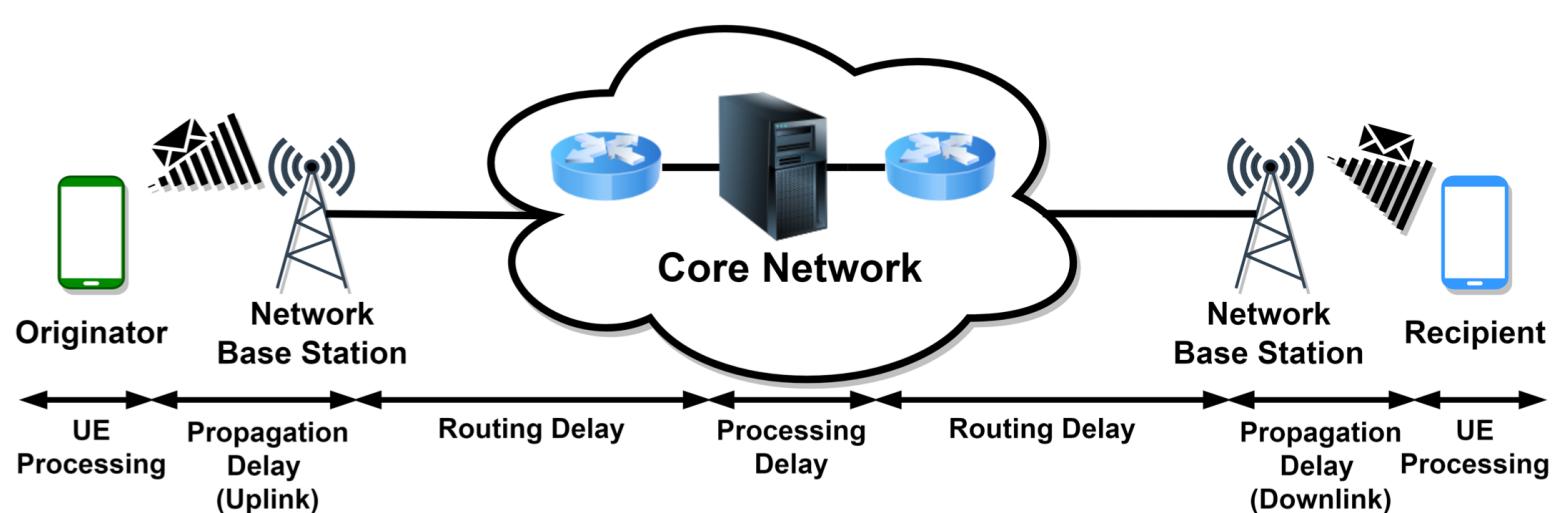
Freaky Leaky SMS: Extracting User Locations by Analyzing SMS Timings

Evangelos Bitsikas, Theo Schnitzler, Christina Poepper, Aanjhan Ranganathan (USENIX Security 2023)

On the Robustness of Wi-Fi Deauthentication

Countermeasures, Domien Schepers, Aanjhan Ranganathan, Mathy Vanhoef (Wisec 2022)

Wi-Fi and Cellular Security



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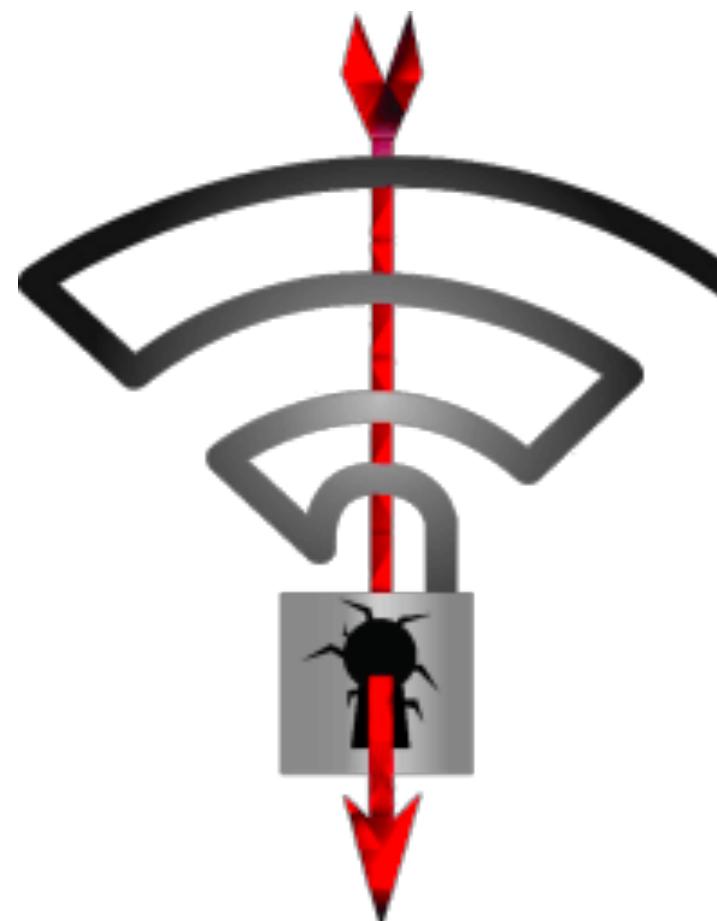
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History of Wi-Fi

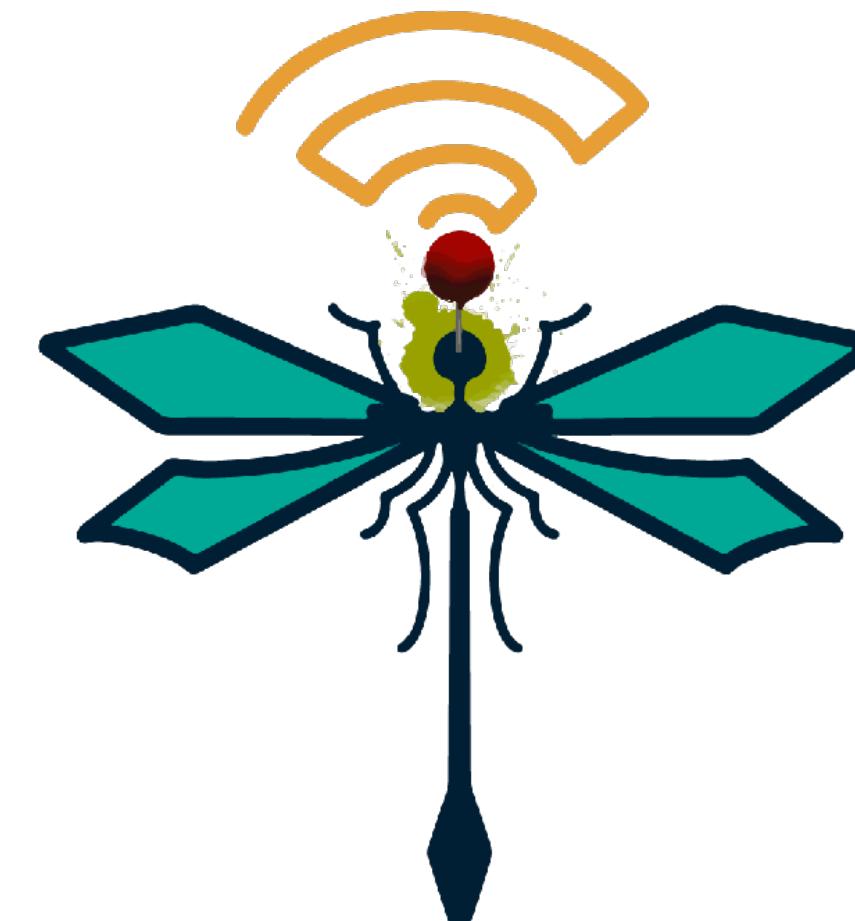
- WEP (1999): quickly broken [FMS01]
- WPA1/2 (~2003)
 - » Offline password brute-force
 - » **KRACK & Kraken** [VP17,VP18]
- WPA3 (2018):
 - » **Dragonblood** side-channels [VR20]



<https://www.eset.com/int/kr00k>



<https://www.krackattacks.com>



<https://wpa3.mathyvanhoef.com>

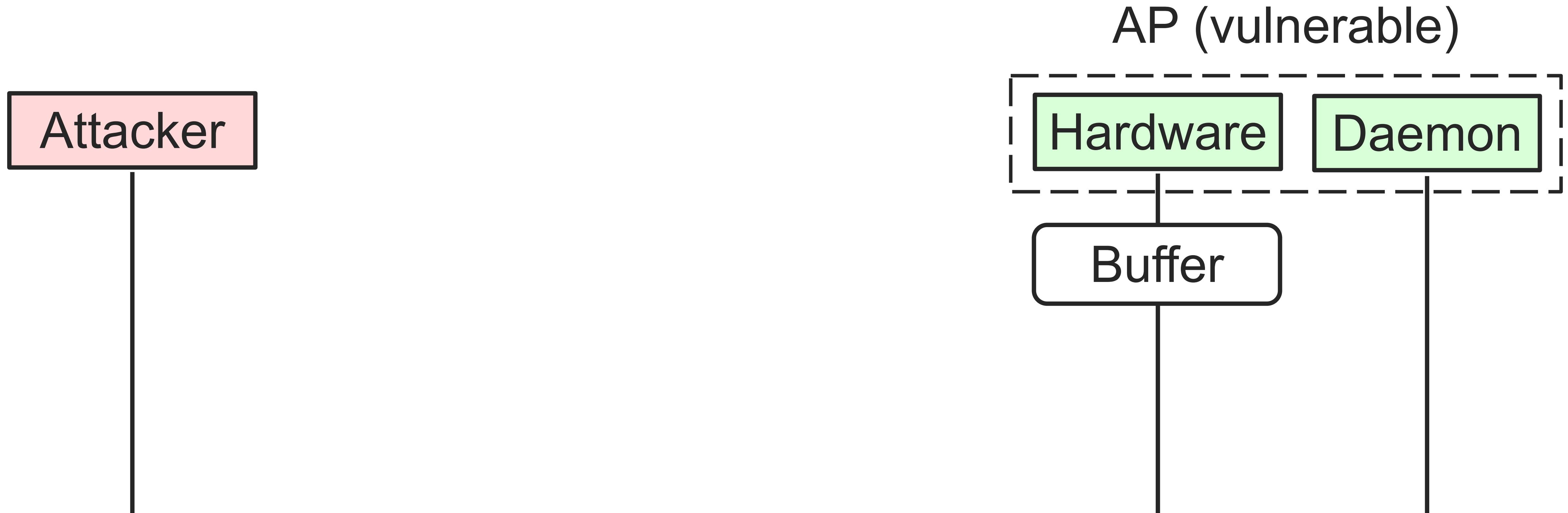


<https://www.fragattacks.com>

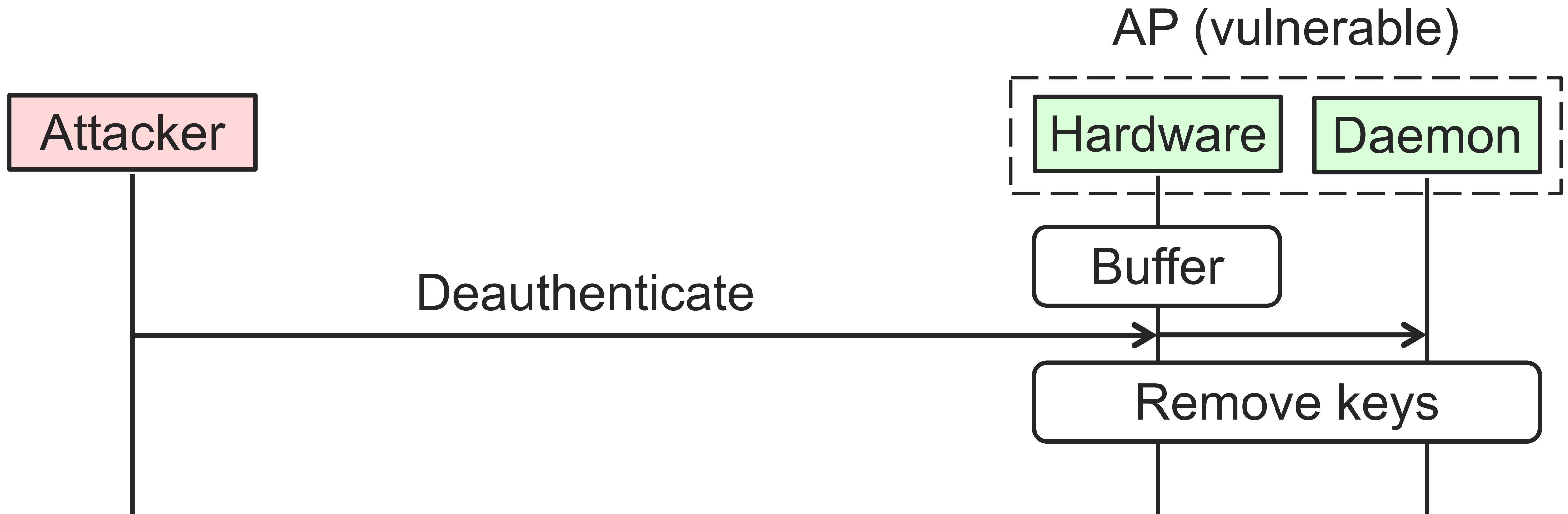
Background: Kr00k implementation flaw



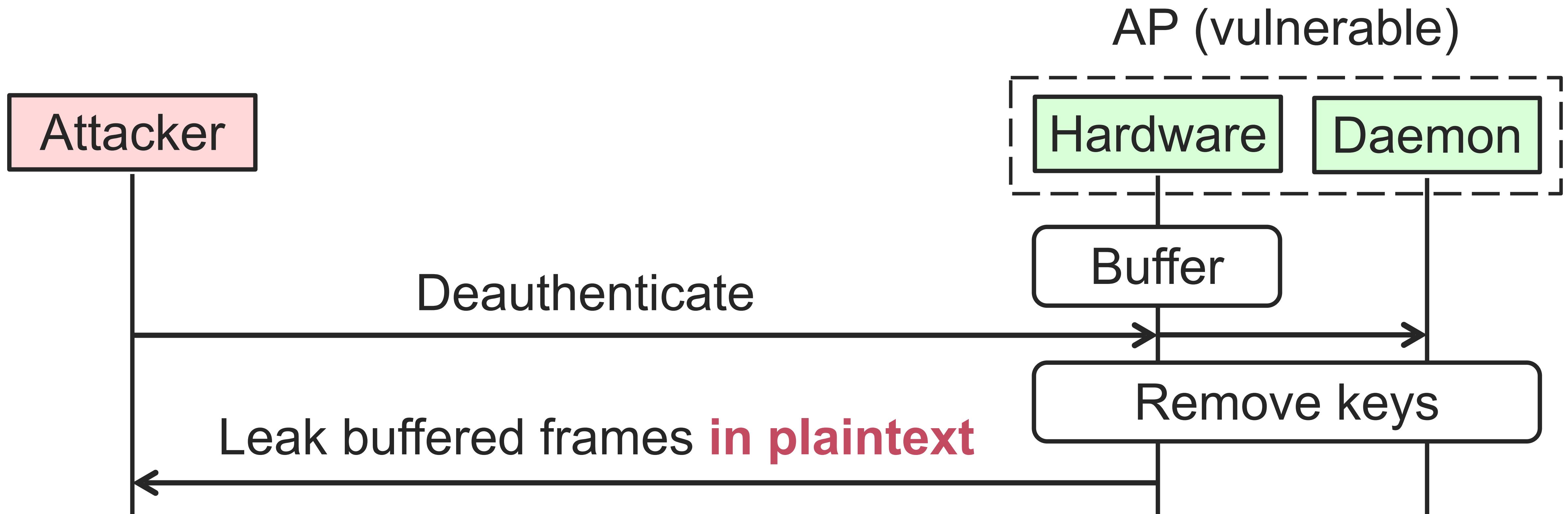
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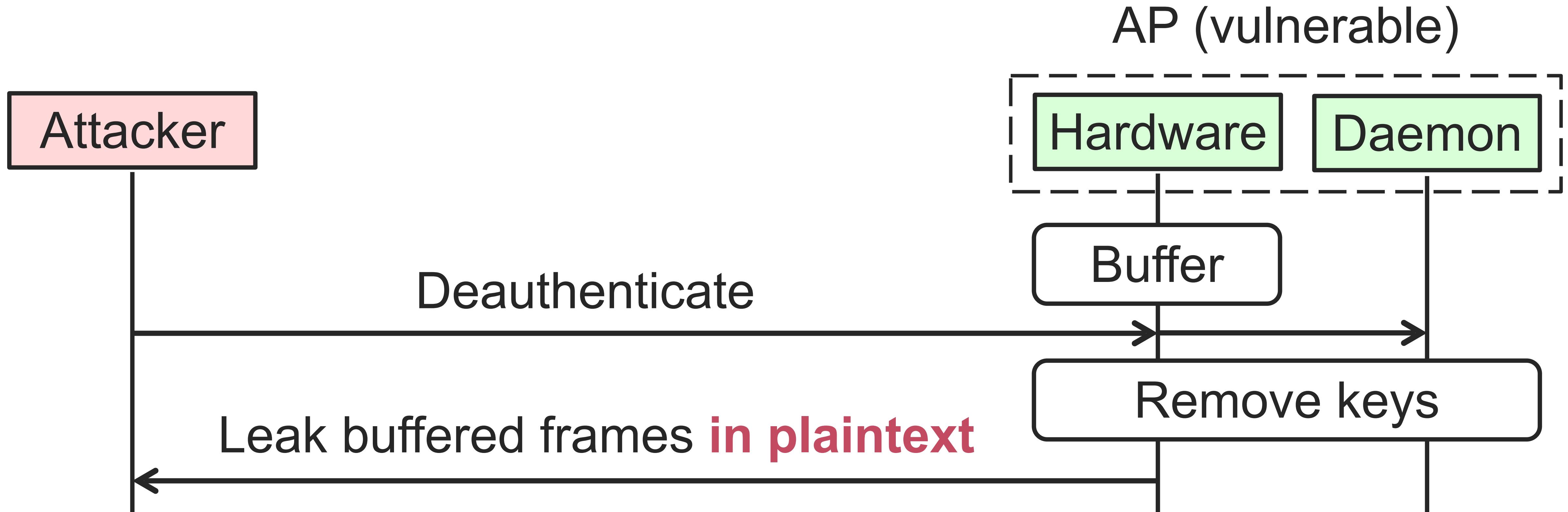
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Research question: how are security contexts managed?

The Security Context

Formally known as the ‘*security association*’ in the IEEE 802.11 standard:

- Protocol suites, negotiated encryption keys, packet counters, ...
- All information needed to securely communicate.

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What is the relation between security context and frames in the transmit queues?

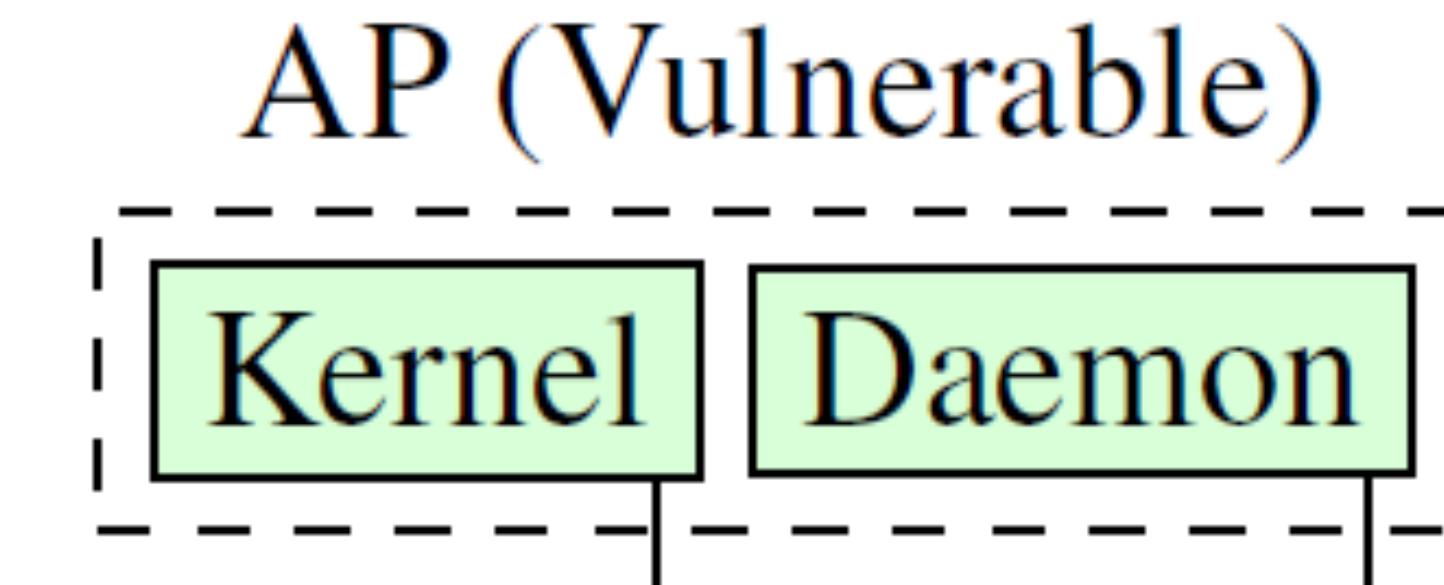
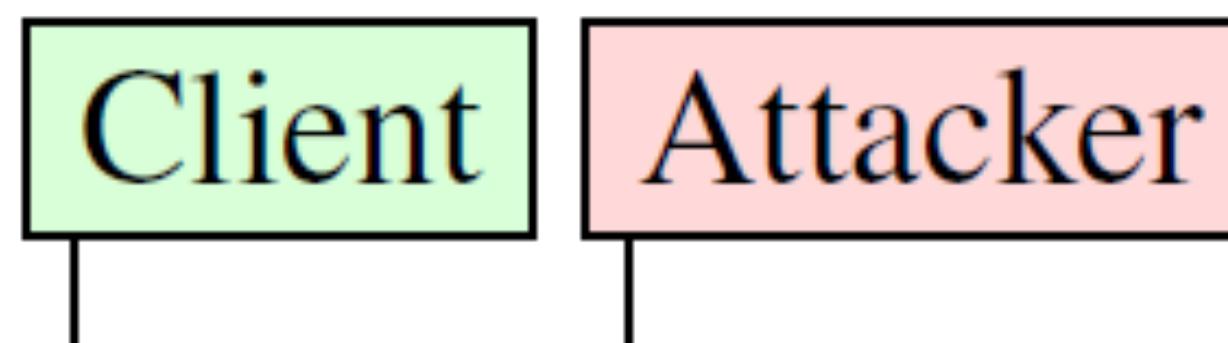
- What happens to a queue if the security context changes?
E.g., reconnection.

1. Can an Adversary Manipulate the Queue
and Security Context?

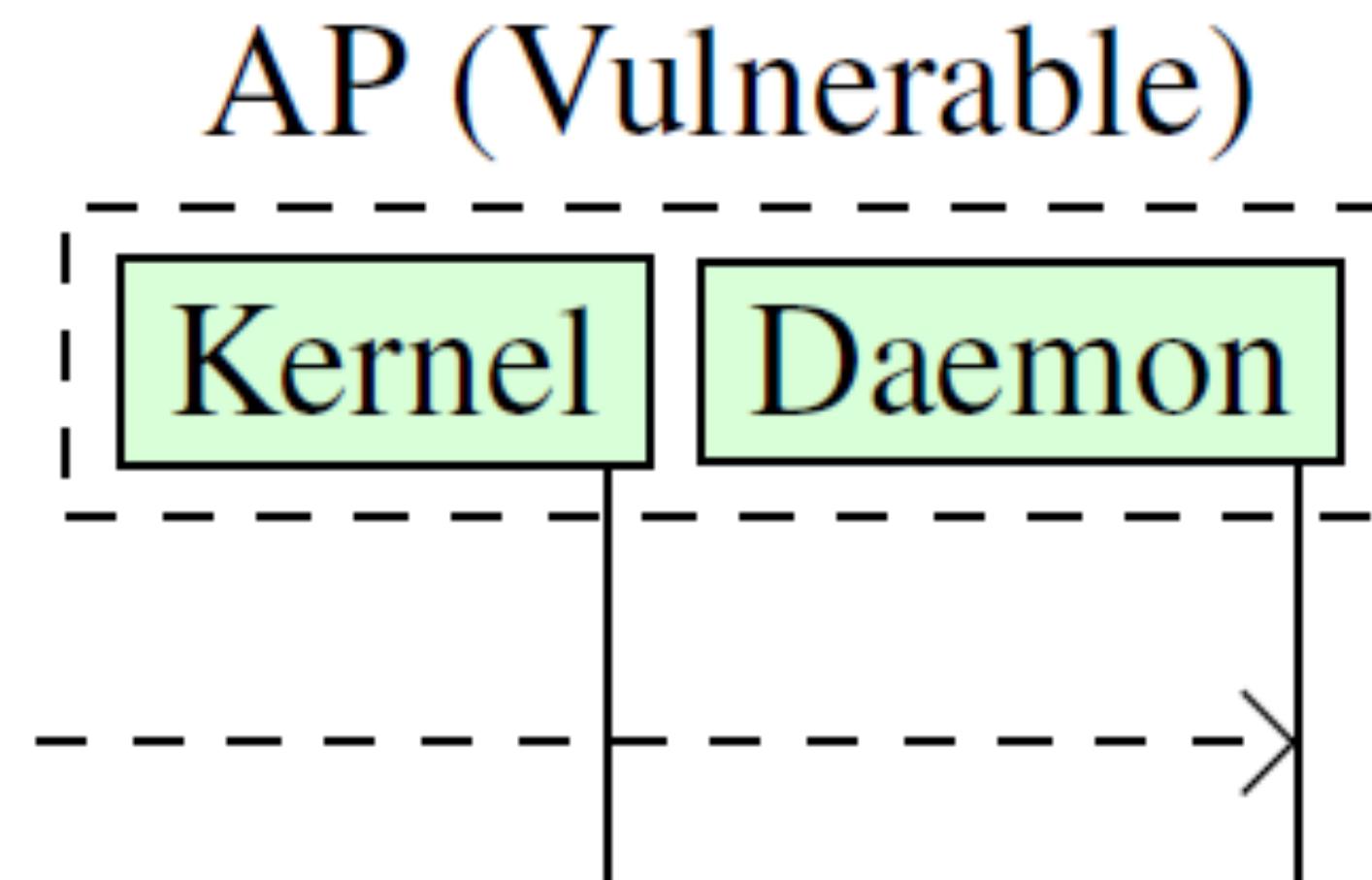
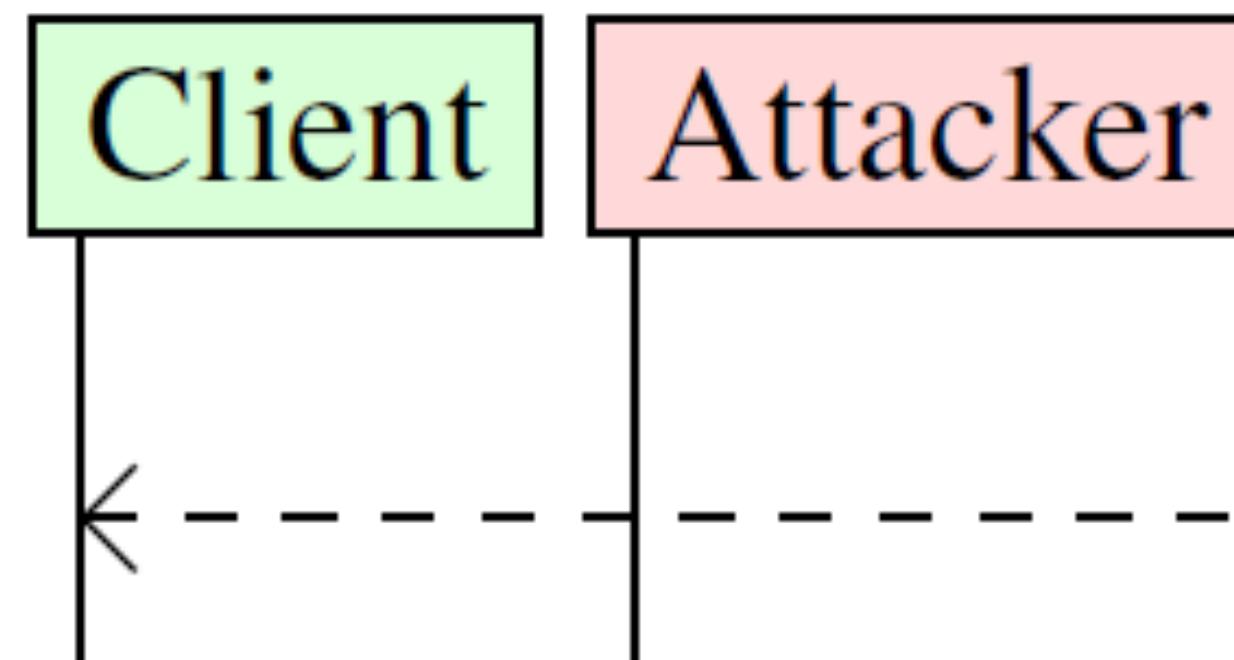
1. Can an Adversary Manipulate the Queue and Security Context?
2. What are the implications?

Finding 1: Leaking Frames

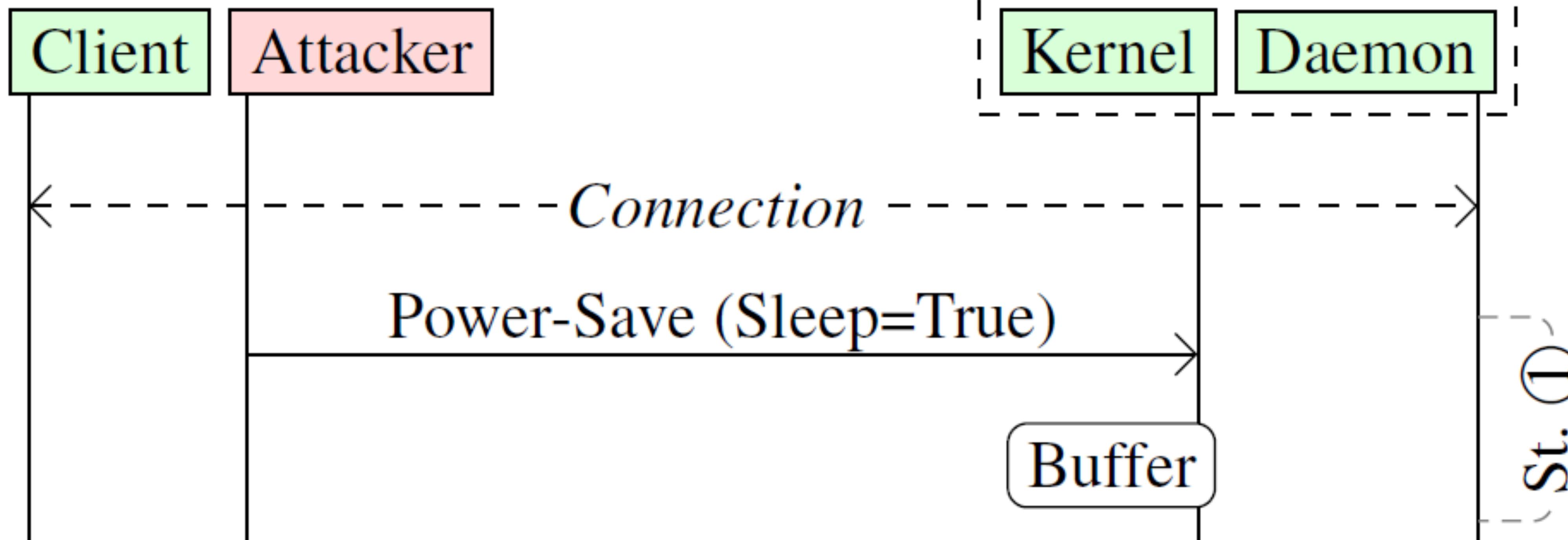
Attack 1: leaking frames



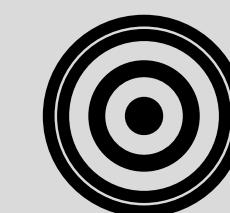
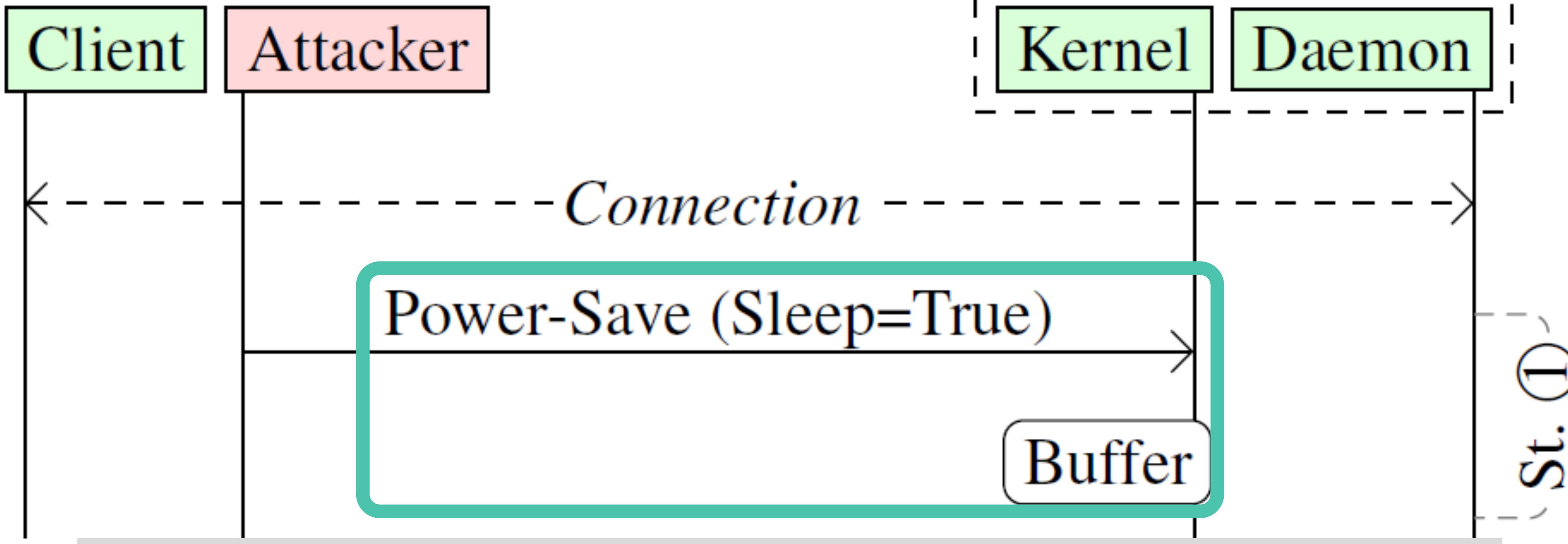
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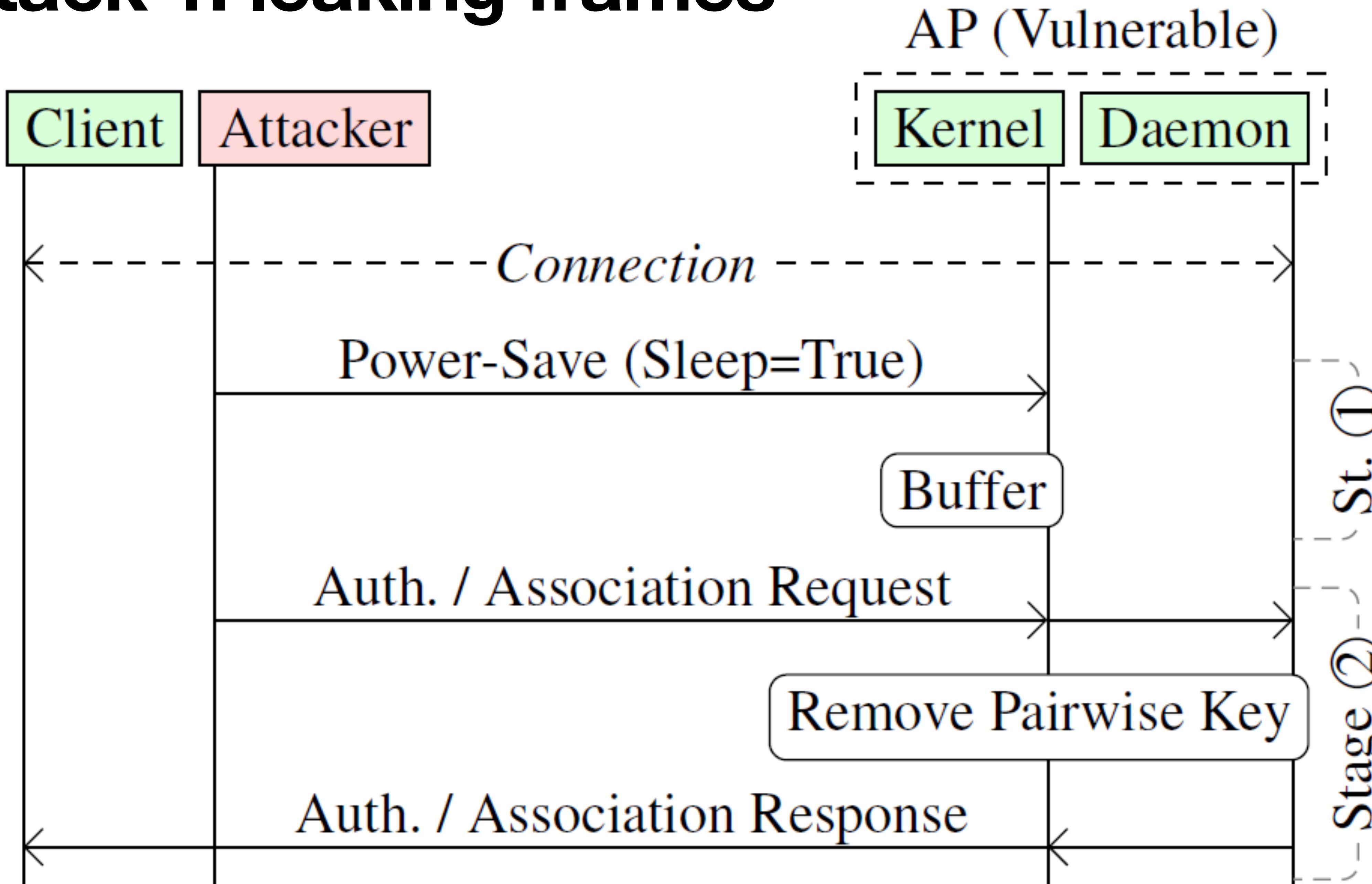


Attack 1: leaking frames

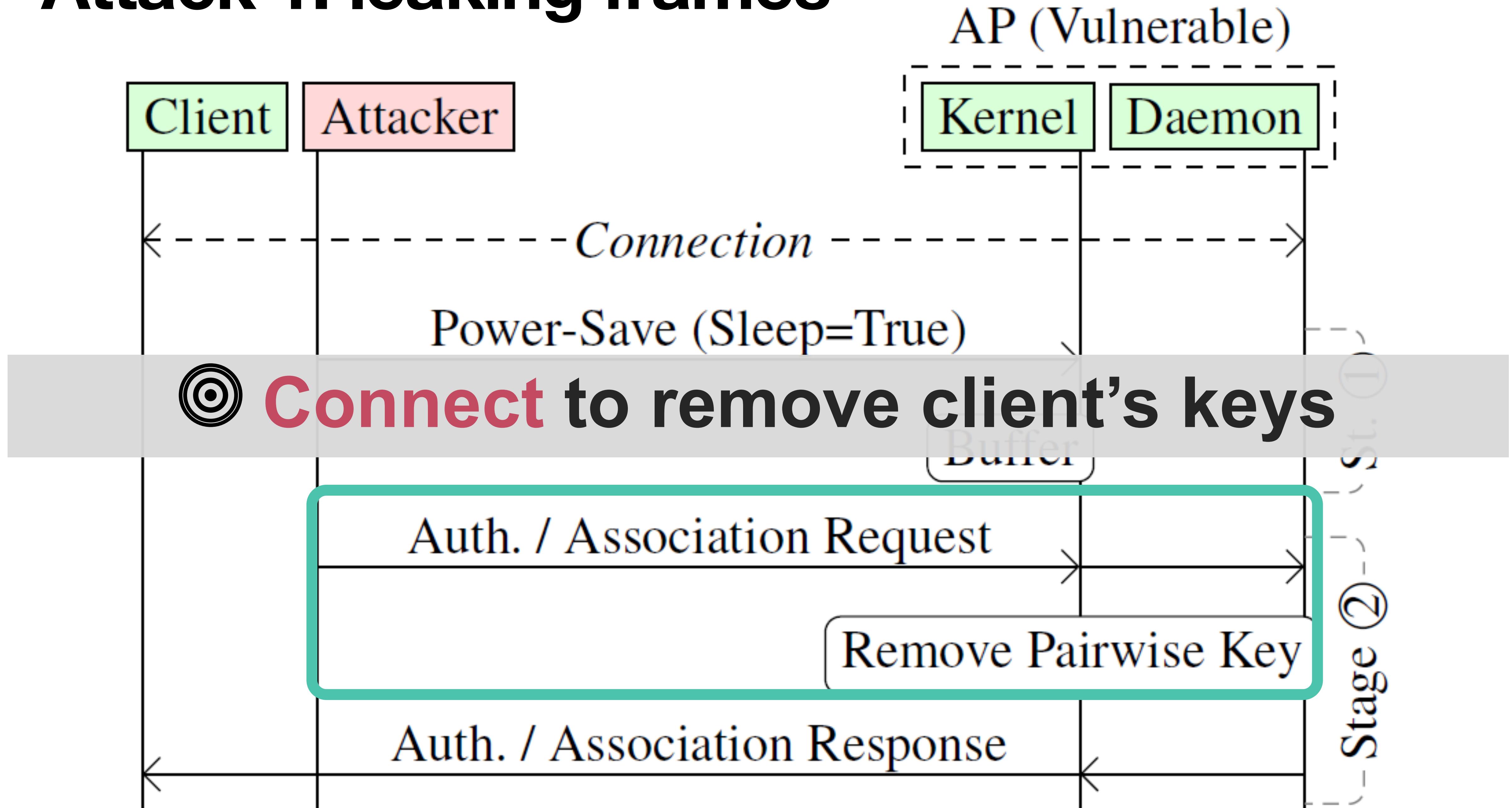


Controlled buffering

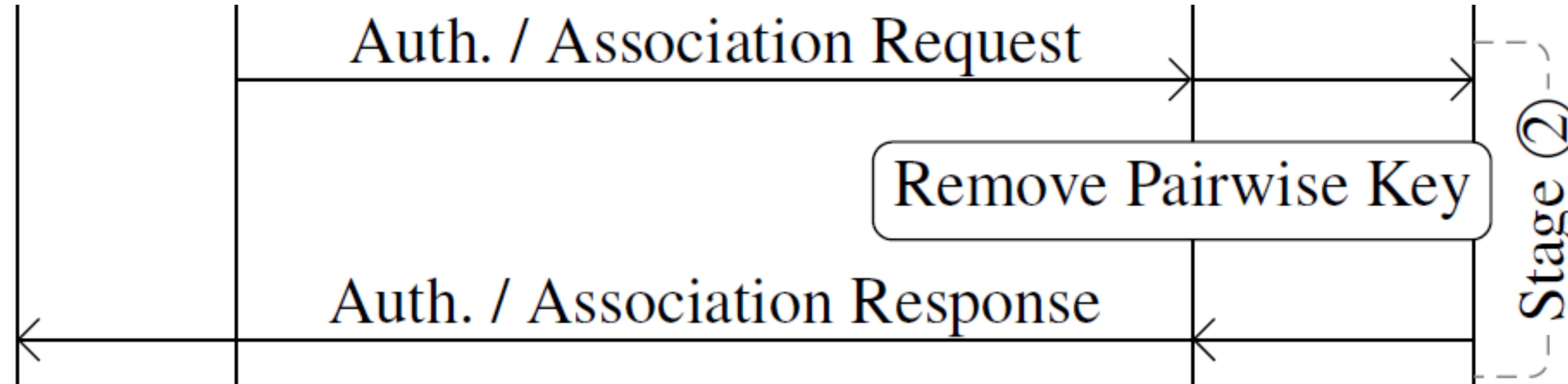
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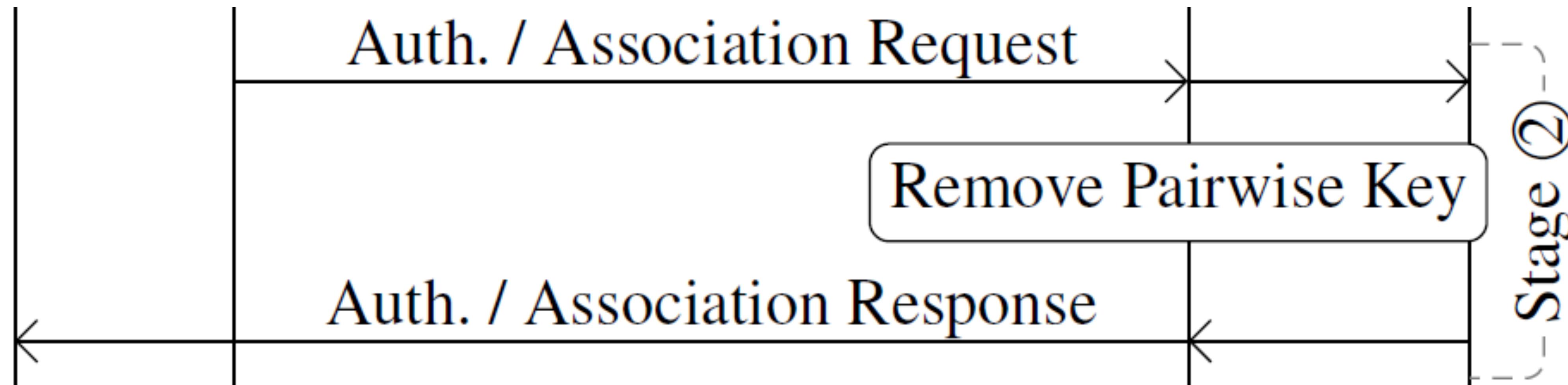
Attack 1: leaking frames



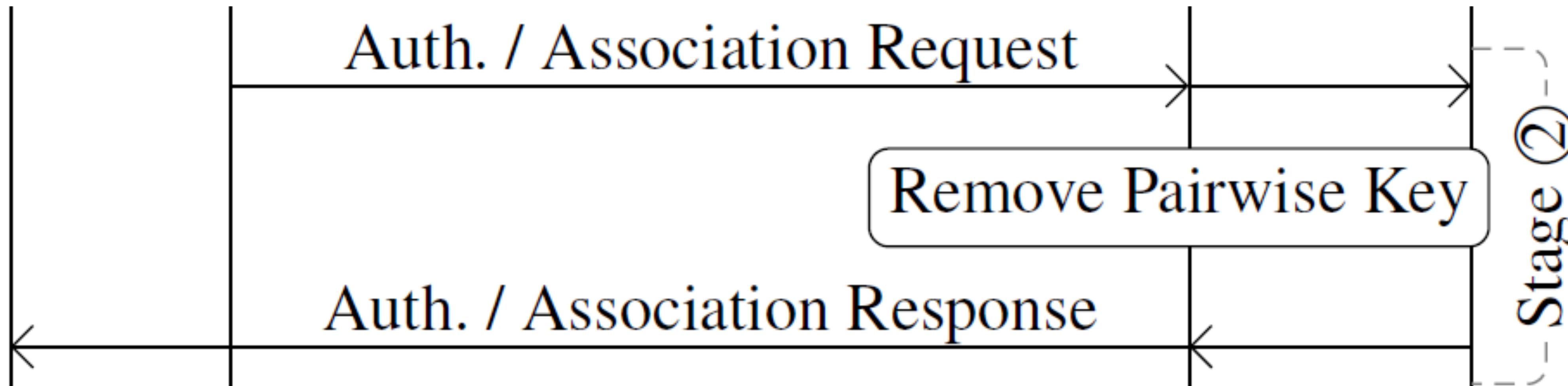
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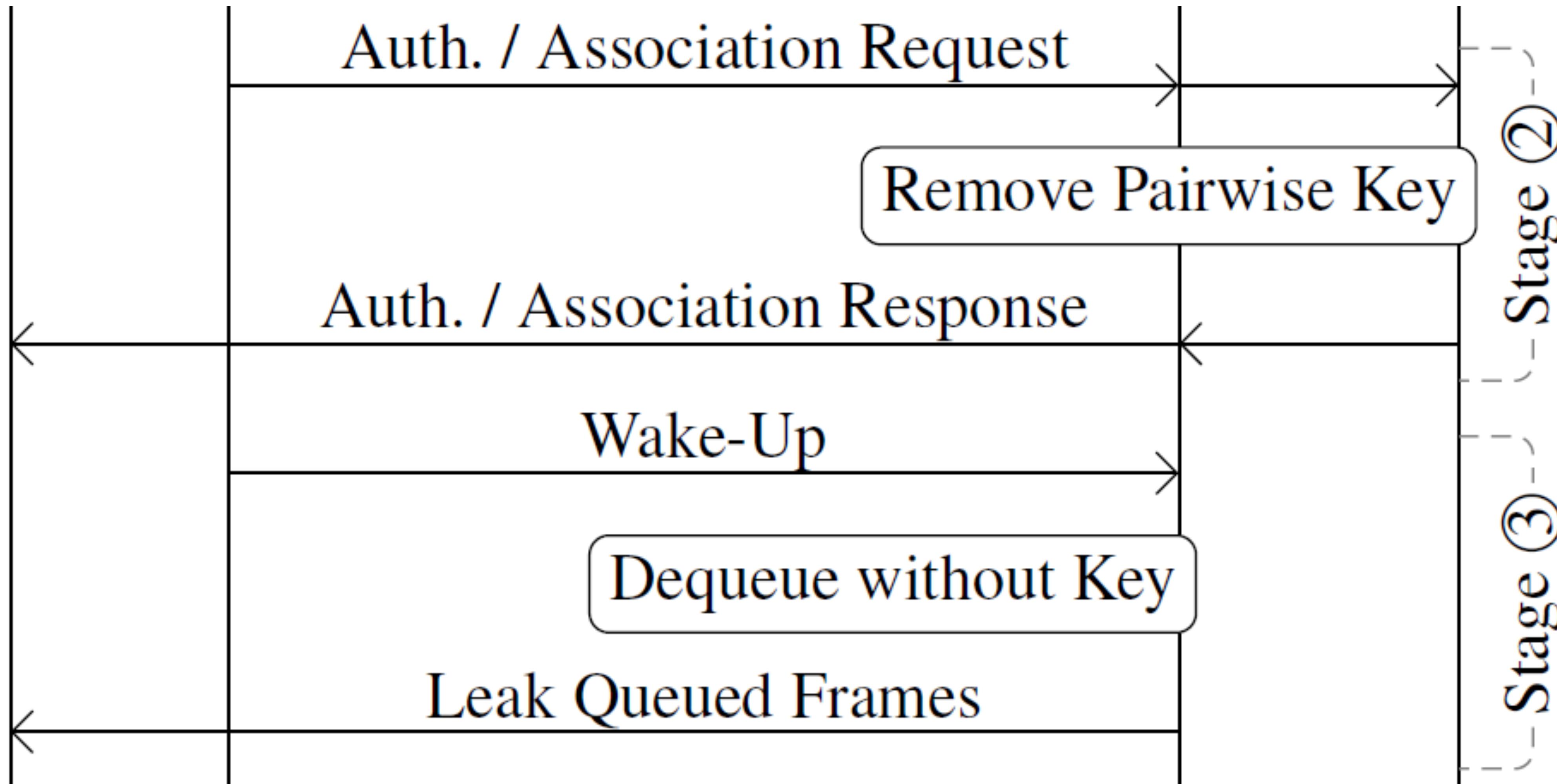
Attack 1: leaking frames



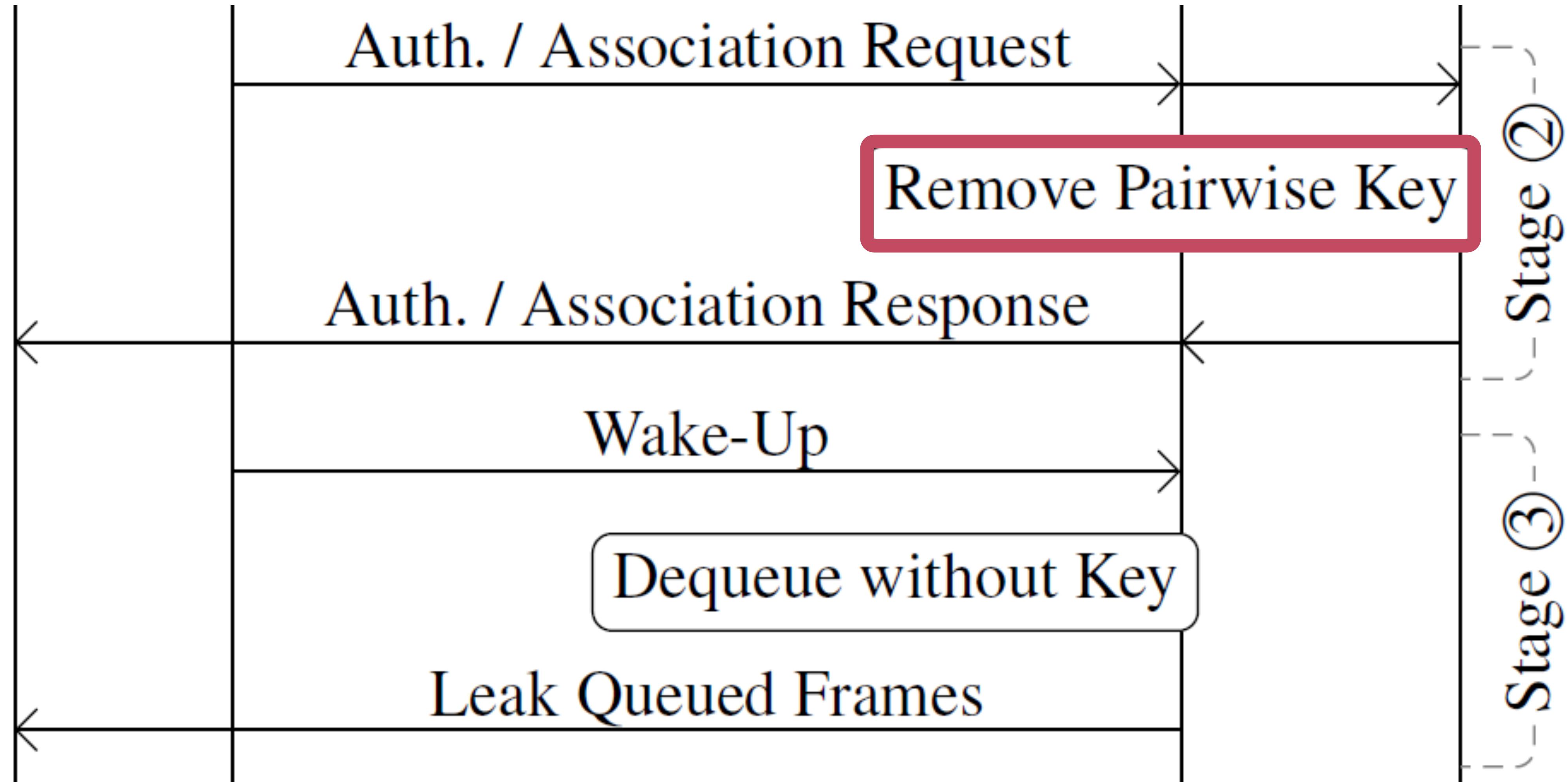
Attack 1: leaking frames



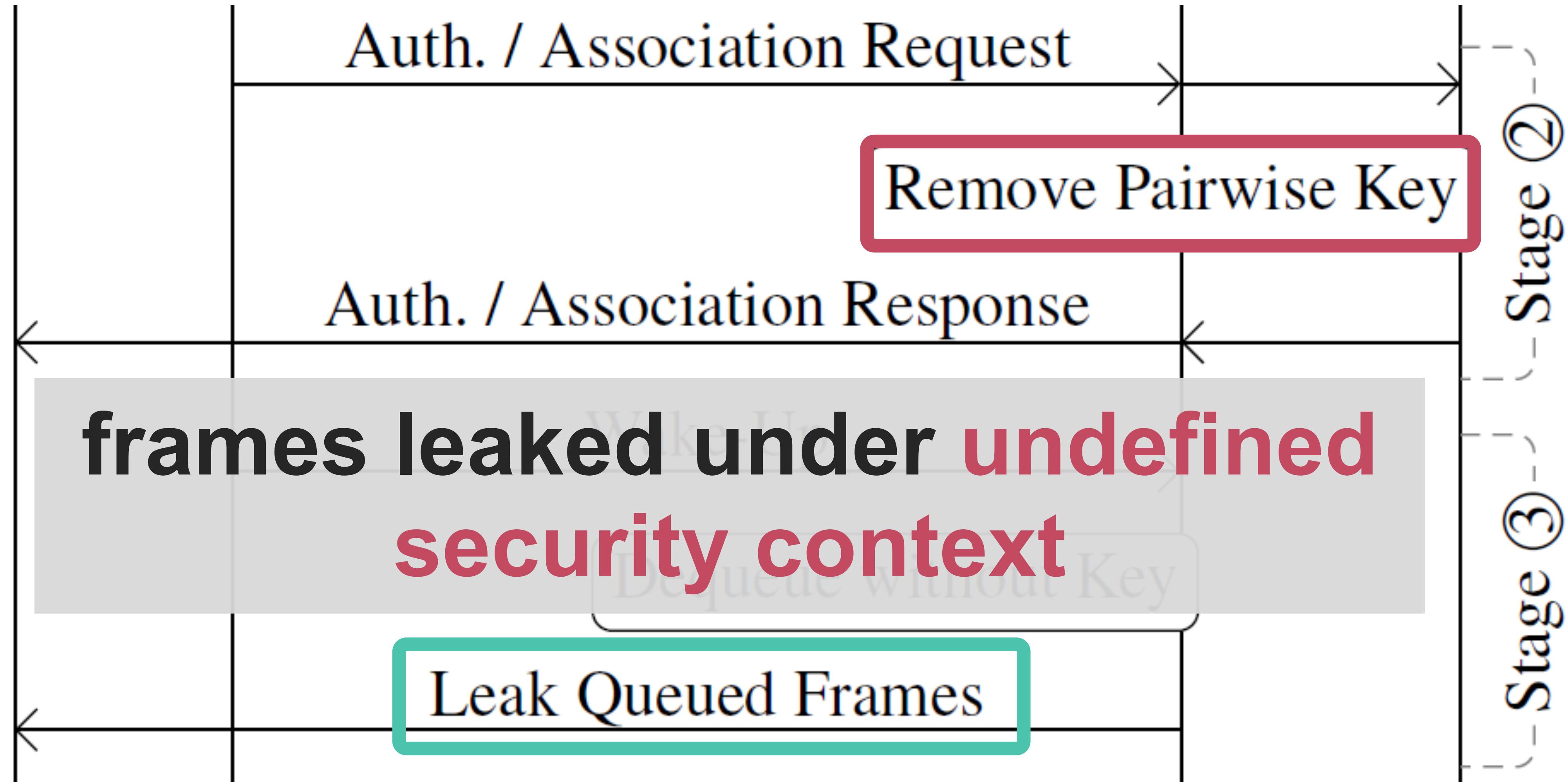
Attack 1: leaking frames



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Attack 1: leaking frames



Undefined security context: FreeBSD example

How the frame is leaked depends on kernel version & driver:

Version	driver (vendor)	Leakage
13.0	run (Ralink)	Plaintext
13.1	run (Ralink)	WEP with all-zero key
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- › Malicious insiders know the group key!
- › Linux, NetBSD, open Atheros firmware also affected

Root cause

Standard isn't explicit on how to manage buffered frames

- Should drop buffered frames when refreshing/deleting keys

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Standard isn't explicit on how to manage buffered frames

- Should drop buffered frames when refreshing/deleting keys

Lesson: include transmit queue in formal Wi-Fi models

- Because buffered frames are not yet encrypted (unlike TLS)
- [CKM20] modelled transmit queue but not key deletion!

Finding 2: Bypassing Client Isolation

Attack 2: Bypassing Wi-Fi Client Isolation

Attack targets networks that use client isolation:

- Defense mechanism against malicious or compromised inside clients.
- Typically networks in large organizations, universities, public hotspots.



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Attacker can connect to the network, but not communicate with others.

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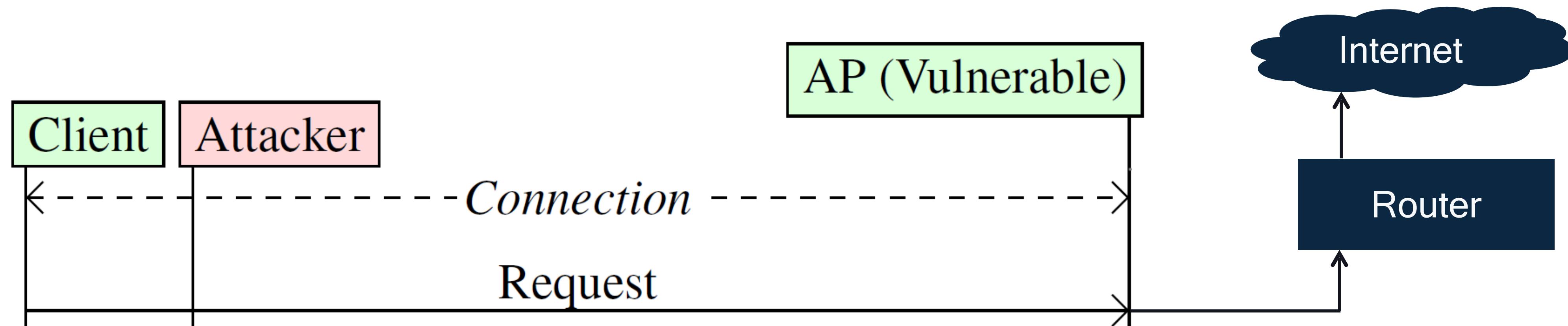
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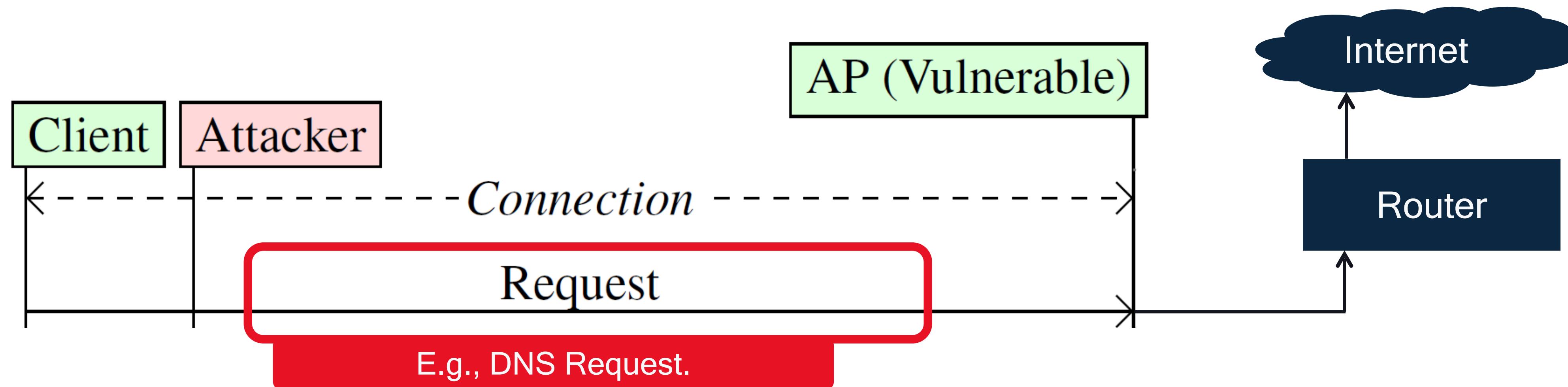
Attacker can connect to the network, but not communicate with others.

... unless we can manipulate the security context!

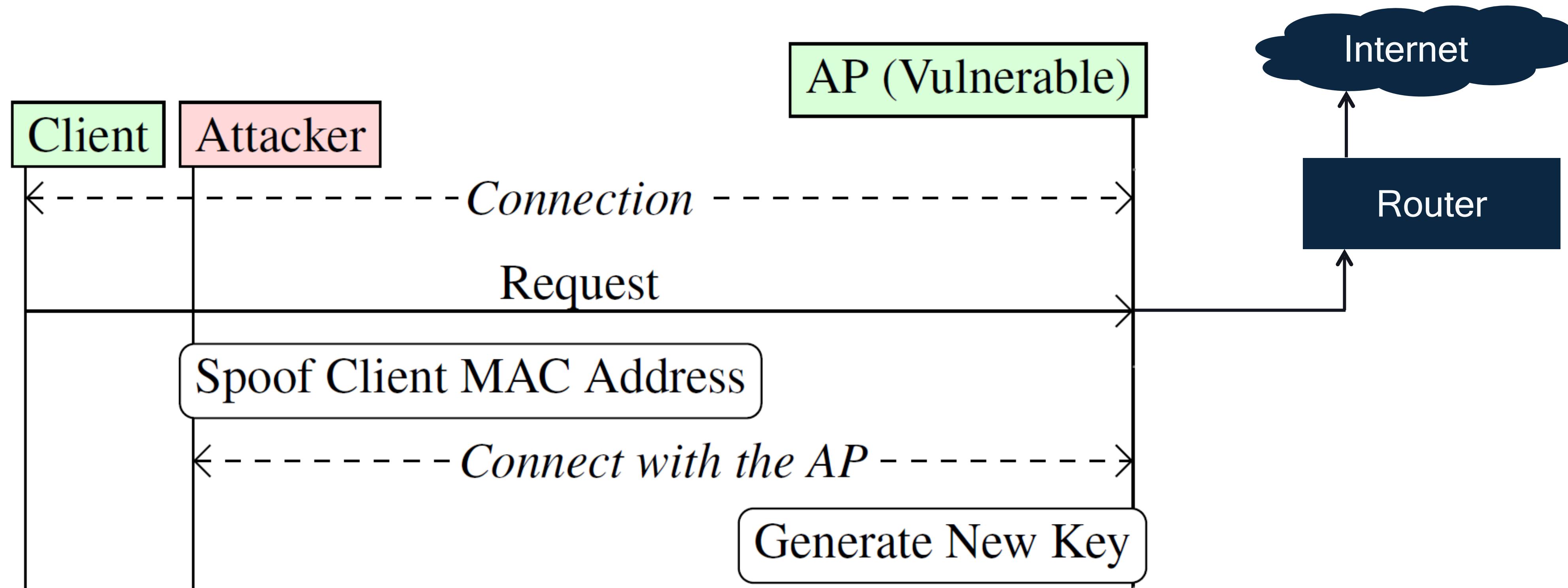
Attack 2: Bypassing Wi-Fi Client Isolation



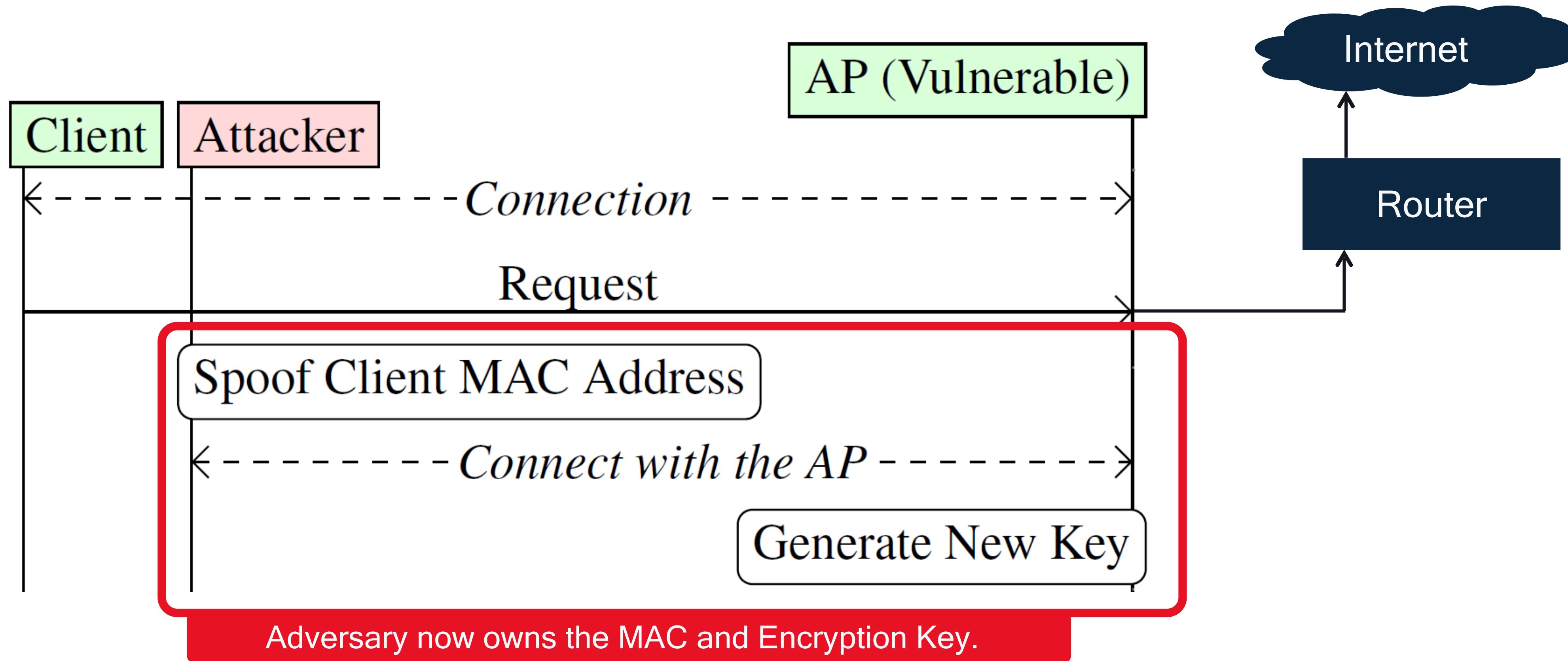
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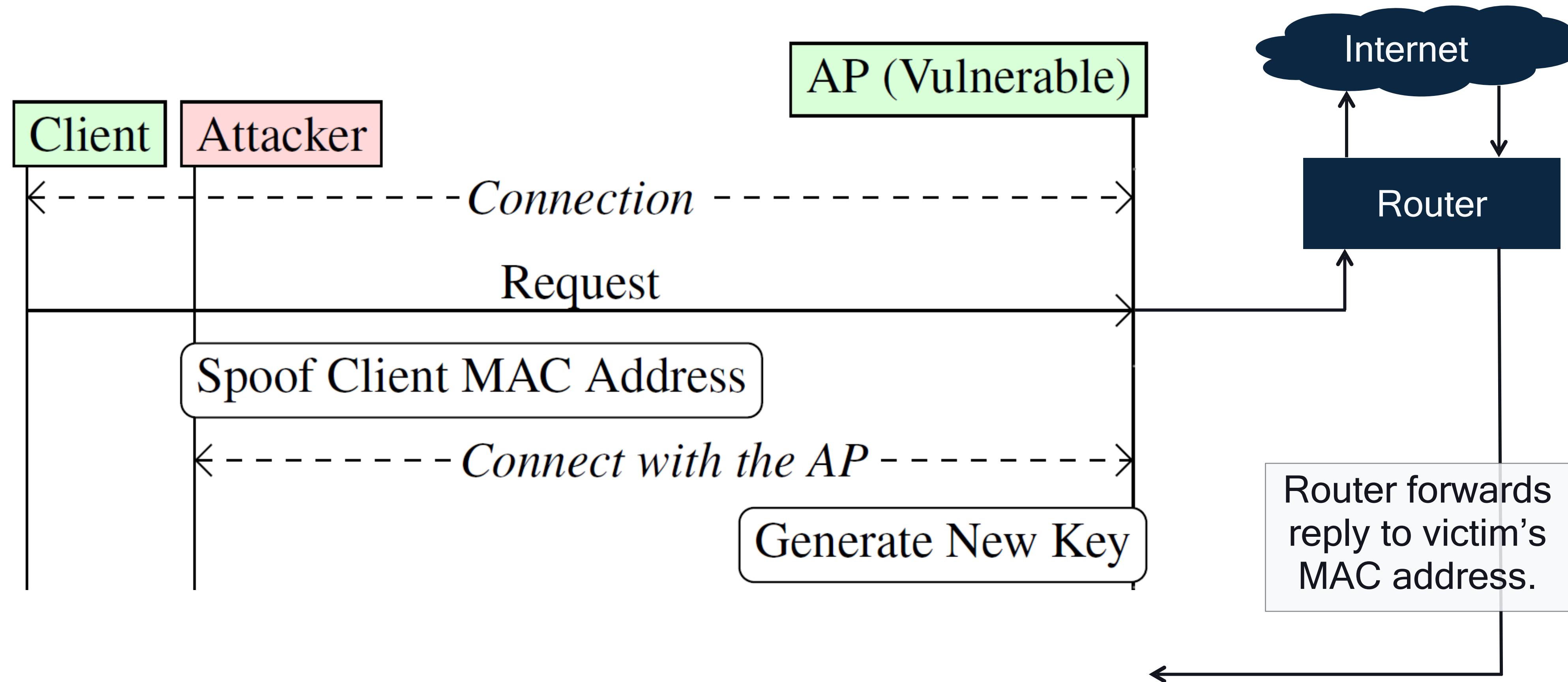
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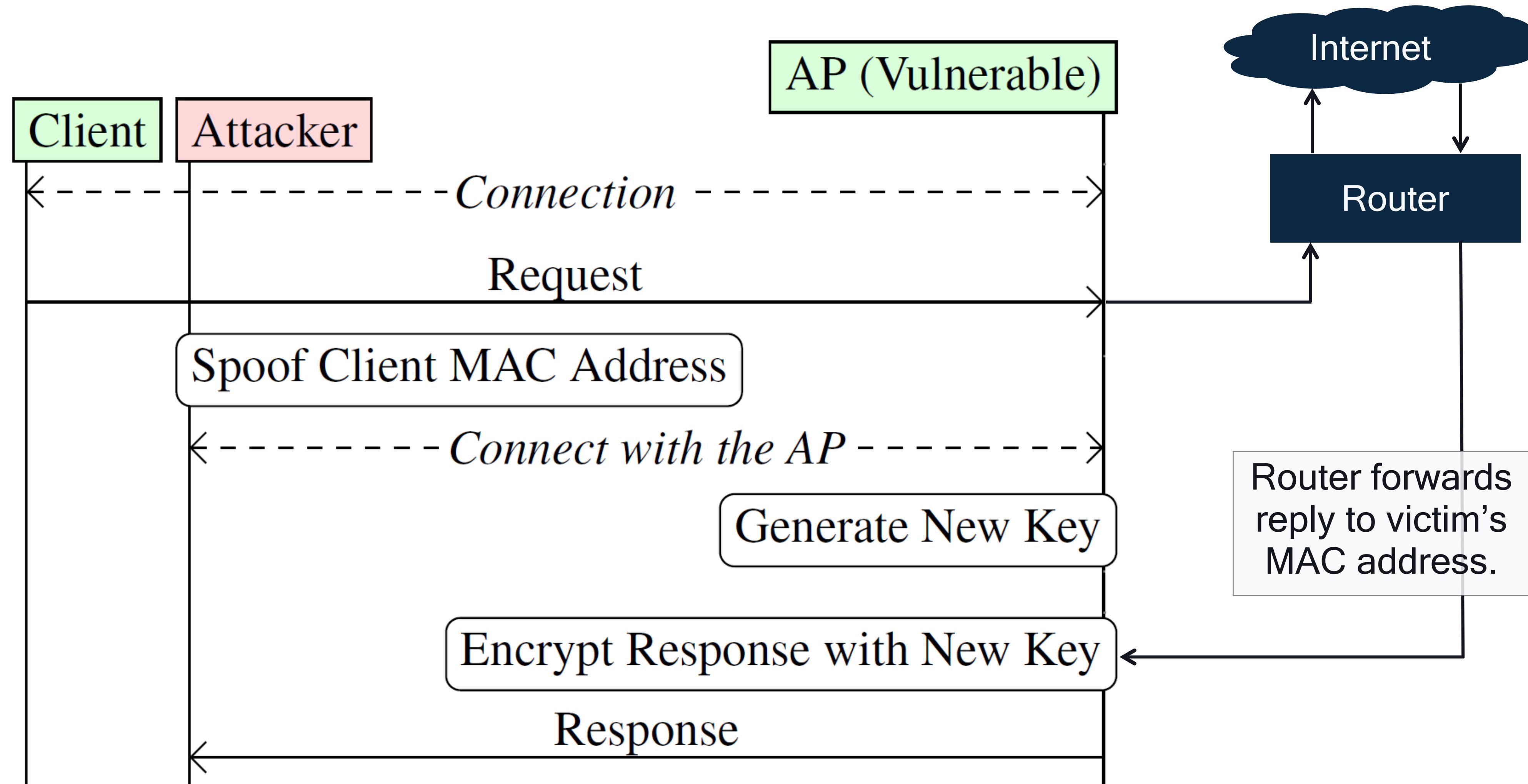
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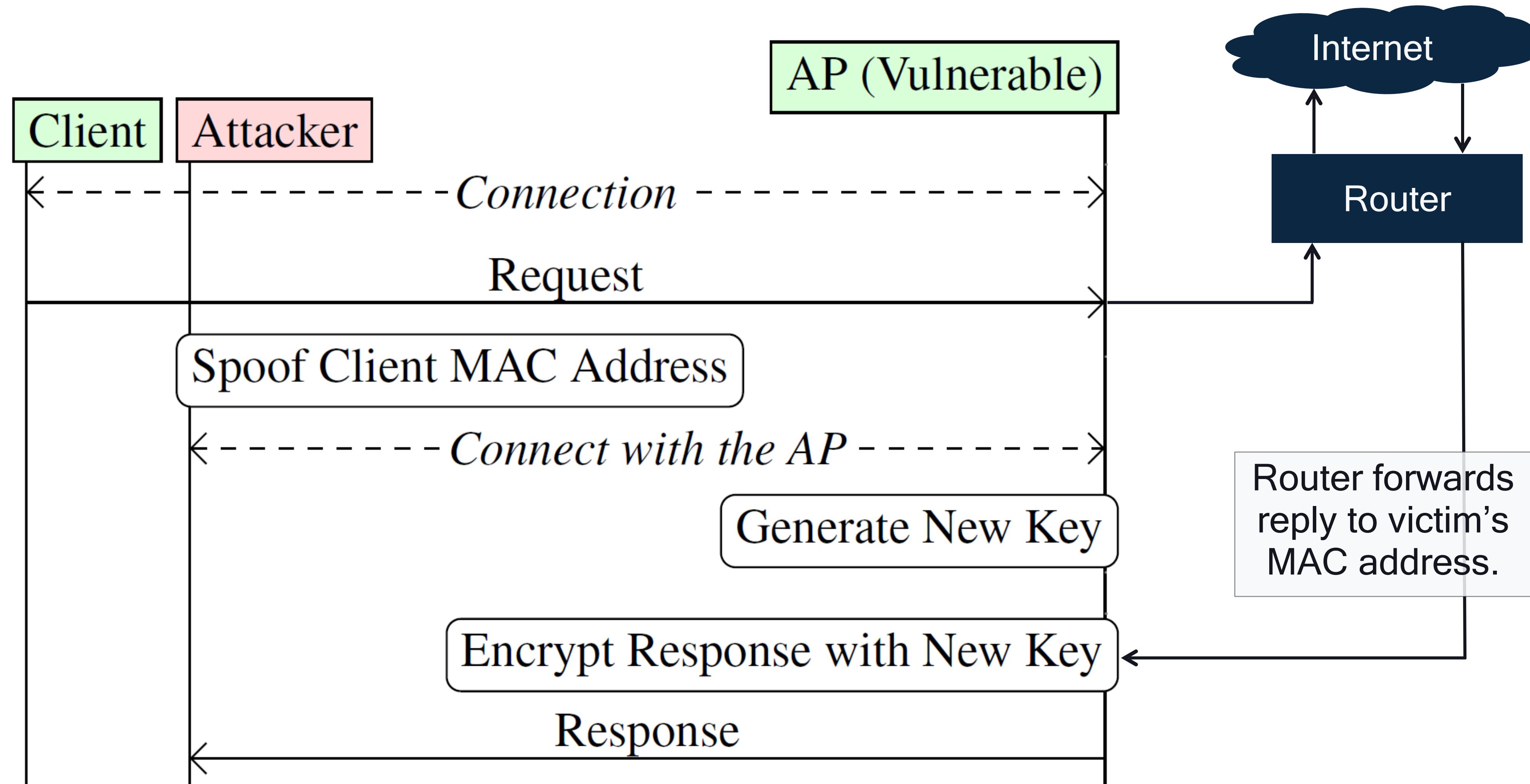
Experiments: home APs

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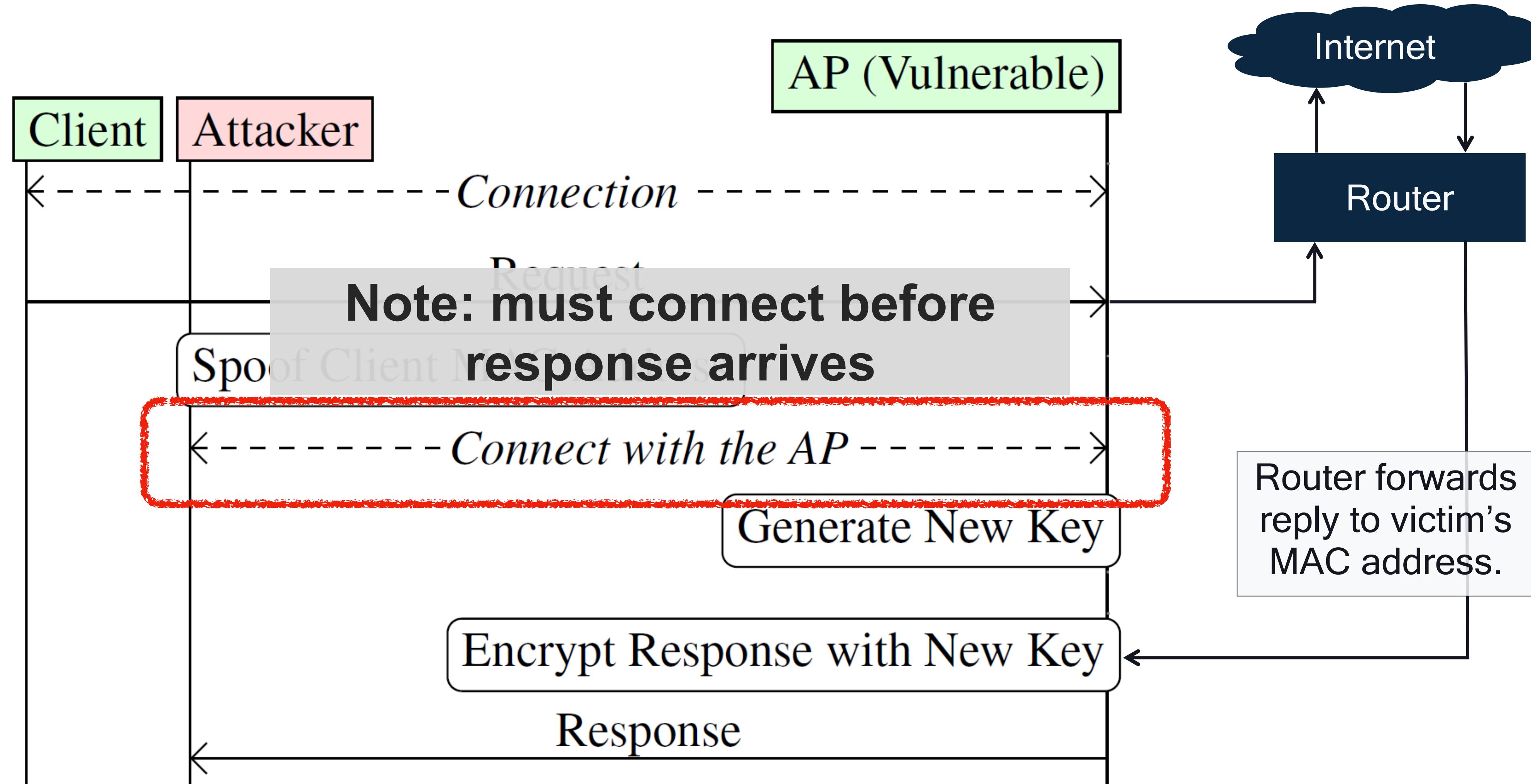
All tested professional & home APs were vulnerable

→ Design flaw in Wi-Fi client isolation!

Attack 2: Bypassing Wi-Fi Client Isolation



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Attack 2: Bypassing Wi-Fi Client Isolation

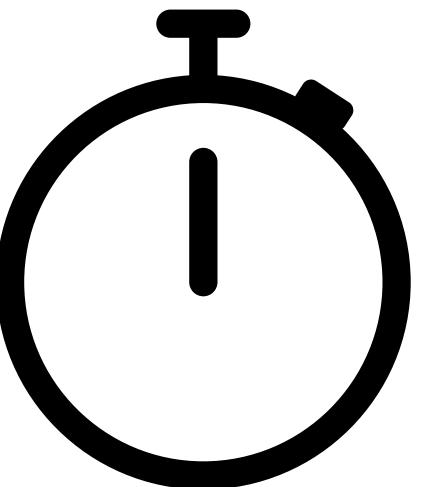
Think of it as a **fast security context override**.



- Requires the attacker to reconnect within certain time restrictions.
- Timing restrictions no concern within transatlantic connections (UDP ~ 70 ms), reasonable within European connections (UDP ~13 ms).
- Protocols such as TCP retransmit when not acknowledged, thus trivial to intercept.

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Adversary can spoof MAC address of a server or gateway in the LAN.

Attack 2: Bypassing Wi-Fi Client Isolation

Why?

Client identities are not bound to each other:

- IEEE 802.1X Identity (username), and
- IP/MAC Addresses.

No concept of ‘protected ownership of a MAC address’ (as is the case in IEEE 802 LANs).

Thus, an adversary can spoof the client’s identity on other layers.

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Design shortcomings/limitations in the standard, network.

Attack 2: Bypassing Wi-Fi Client Isolation

- This is not a simple (or difficult) code fix for anyone.
- Needs to be addressed within multiple network components, beyond an access point.

Solutions? Probably not realistic, practical, or sufficient:

- Reject recently-used MAC addresses (e.g., a ten second delay if client isolation is configured).
- Network configurations to use separate (un)trusted clients (e.g., different SSIDs, usage of VLANs).
- Require connection establishments to use a cached key if recently-used MAC address.

Summary

- Standard is vague and requires explicit elaboration on managing buffered frames
 - Can **leak frames** under different security context
 - Important to **model/define transmit queues**
- Can **bypass client isolation**
 - All devices vulnerable -> **design flaw**
 - Hard to fully prevent
- Some DoS attacks also possible (paper has details)

March 2023

doc.: IEEE 802.11-23/537r0

IEEE P802.11
Wireless LANs

Reassociating STA recognition

Date: 2023-03-27

Author(s):

Name	Affiliation	Address	Phone	email
Jouni Malinen	Qualcomm Technologies, Inc.			jouni@qca.qualcomm.com

Abstract

This document discusses issues related to secure recognition of a reassociating STA by an AP and proposed new mechanism to allow this to be done. This is related to the association comeback in management frame protection and how the use of SA Query can result in undesired latency in being able to negotiate new parameters for an association in the reassociate-to-same-BSS case. Furthermore, the proposed design can provide some help in addressing recently reported security vulnerabilities in MAC address “ownership” and potential insider attacks.

<https://mentor.ieee.org/802.11/doc/23/11-23-0537-00-000m->

GitHub

<https://github.com/vanhoefm/macstealer>

<https://github.com/domienschepers/wifi-framing>

CVE-2022-47522

Thank you!