Team notebook

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```
vector <int> res;
res.resize((int)Q.size());

int L = 1, R = 0;
for(query q: Q){
    while (L > q.1) add(--L);
    while (R < q.r) add(++R);

    while (L < q.l) del(L++);
    while (R > q.r) del(R--);

    res[q.pos] = calc(1, R-L+1);
}
return res;
```

1.2 Mo's Algorithms on Trees

```
Given a tree with N nodes and Q queries. Each node has
    an integer weight.
Each query provides two numbers u and v, ask for how
    many different integers weight of nodes
there are on path from u to v.
Modify DFS:
For each node u, maintain the start and the end DFS
    time. Let's call them ST(u) and EN(u).
=> For each query, a node is considered if its
    occurrence count is one.
Query solving:
Let's query be (u, v). Assume that ST(u) \le ST(v).
    Denotes P as LCA(u, v).
Case 1: P = 11
Our query would be in range [ST(u), ST(v)].
Case 2: P != u
Our query would be in range [EN(u), ST(v)] + [ST(p),
     ST(p)]
void update(int &L, int &R, int qL, int qR){
   while (L > qL) add(--L);
   while (R < qR) add(++R);
```

```
while (L < qL) del(L++);
    while (R > qR) del(R--);
}
vector <int> MoQueries(int n, vector <query> Q){
    block_size = sqrt((int)nodes.size());
    sort(Q.begin(), Q.end(), [](const query &A, const
         query &B){
       return (ST[A.1]/block size !=
            ST[B.1]/block_size)? (ST[A.1]/block_size <</pre>
            ST[B.1]/block size) : (ST[A.r] < ST[B.r]):
   });
    vector <int> res:
    res.resize((int)Q.size());
    LCA lca:
    lca.initialize(n):
    int L = 1, R = 0:
    for(query q: Q){
       int u = q.1, v = q.r;
       if(ST[u] > ST[v]) swap(u, v); // assume that
            S[u] \le S[v]
       int parent = lca.get(u, v);
       if(parent == u){
           int qL = ST[u], qR = ST[v];
           update(L, R, qL, qR);
       }else{
           int qL = EN[u], qR = ST[v];
           update(L, R, qL, qR);
           if(cnt_val[a[parent]] == 0)
               res[q.pos] += 1;
       res[q.pos] += cur_ans;
    }
    return res:
```

1.3 Parallel Binary Search

```
int lo[N], mid[N], hi[N];
vector<int> vec[N];

void clear() //Reset
{
        memset(bit, 0, sizeof(bit));
}

void apply(int idx) //Apply ith update/query
```

```
{
       if(ql[idx] <= qr[idx])</pre>
               update(ql[idx], qa[idx]),
                    update(qr[idx]+1, -qa[idx]);
        else
               update(1, qa[idx]);
               update(qr[idx]+1, -qa[idx]);
               update(ql[idx], qa[idx]);
       }
}
bool check(int idx) //Check if the condition is
     satisfied
{
       int req=reqd[idx];
       for(auto &it:owns[idx])
               req-=pref(it);
               if(req<0)
                       break:
       if(req<=0)</pre>
               return 1;
       return 0;
}
void work()
        for(int i=1;i<=q;i++)</pre>
               vec[i].clear();
        for(int i=1;i<=n;i++)</pre>
               if(mid[i]>0)
                       vec[mid[i]].push_back(i);
        clear():
        for(int i=1:i<=a:i++)</pre>
               for(auto &it:vec[i]) //Add appropriate
                    check conditions
                       if(check(it))
                              hi[it]=i:
                       else
                              lo[it]=i+1;
               }
       }
}
void parallel_binary()
       for(int i=1;i<=n;i++)</pre>
               lo[i]=1, hi[i]=q+1;
```

2 Combinatorics

2.1 Factorial Approximate

Approximate Factorial:

$$n! = \sqrt{2.\pi \cdot n} \cdot \left(\frac{n}{e}\right)^n \tag{1}$$

2.2 Factorial

n	123	4	5 6	7	8	9	10	
n!	1 2 6	24 1	.20 720	0 5040	40320	362880	3628800	
n	11	12	13	14	15	16	17	
$\overline{n!}$							13 3.6e14	
n	20	25	30	40	50 10	00 - 15	0 171	
$\overline{n!}$	2e18	2e25	3e32	8e47 3	6e64 9e	157 6e2	$62 > DBL_M$	AX

2.3 Fast Fourier Transform

```
/**
 * Fast Fourier Transform.
 * Useful to compute convolutions.
 * computes:
 * C(f star g)[n] = sum_m(f[m] * g[n - m])
 * for all n.
 * test: icpc live archive, 6886 - Golf Bot
 * */
```

```
using namespace std;
#include <bits/stdc++.h>
#define D(x) cout << #x " = " << (x) << endl
#define endl '\n'
const int MN = 262144 << 1;</pre>
int d[MN + 10], d2[MN + 10];
const double PI = acos(-1.0);
struct cpx {
 double real, image;
 cpx(double _real, double _image) {
   real = real:
   image = _image;
 cpx(){}
cpx operator + (const cpx &c1, const cpx &c2) {
 return cpx(c1.real + c2.real, c1.image + c2.image);
cpx operator - (const cpx &c1, const cpx &c2) {
 return cpx(c1.real - c2.real, c1.image - c2.image);
cpx operator * (const cpx &c1, const cpx &c2) {
 return cpx(c1.real*c2.real - c1.image*c2.image,
      c1.real*c2.image + c1.image*c2.real);
int rev(int id, int len) {
 int ret = 0:
 for (int i = 0; (1 << i) < len; i++) {
   ret <<= 1;
   if (id & (1 << i)) ret |= 1;</pre>
 return ret;
cpx A[1 << 20]:
void FFT(cpx *a, int len, int DFT) {
 for (int i = 0; i < len; i++)</pre>
   A[rev(i, len)] = a[i];
 for (int s = 1; (1 << s) <= len; s++) {
   int m = (1 << s);
   cpx wm = cpx(cos(DFT * 2 * PI / m), sin(DFT * 2 *
   for(int k = 0; k < len; k += m) {
```

```
cpx w = cpx(1, 0);
     for(int j = 0; j < (m >> 1); j++) {
       cpx t = w * A[k + j + (m >> 1)];
       cpx u = A[k + j];
       A[k + j] = u + t;
       A[k + j + (m >> 1)] = u - t;
 if (DFT == -1) for (int i = 0: i < len: i++)
      A[i].real /= len, A[i].image /= len;
 for (int i = 0; i < len; i++) a[i] = A[i];</pre>
 return:
cpx in[1 << 20]:
void solve(int n) {
 memset(d, 0, sizeof d);
 for (int i = 0; i < n; ++i) {</pre>
   cin >> t:
   d[t] = true;
  int m;
 cin >> m;
 vector<int> q(m);
 for (int i = 0; i < m; ++i)</pre>
   cin >> q[i];
  for (int i = 0; i < MN; ++i) {</pre>
   if (d[i])
     in[i] = cpx(1, 0);
     in[i] = cpx(0, 0);
 FFT(in, MN, 1):
 for (int i = 0; i < MN; ++i) {</pre>
   in[i] = in[i] * in[i];
 FFT(in, MN, -1);
 int ans = 0:
 for (int i = 0; i < q.size(); ++i) {</pre>
   if (in[q[i]].real > 0.5 || d[q[i]]) {
  cout << ans << endl;
```

```
int main() {
  ios_base::sync_with_stdio(false);cin.tie(NULL);
  int n;
  while (cin >> n)
    solve(n);
  return 0;
}
```

2.4 General purpose numbers

Bernoulli numbers

EGF of Bernoulli numbers is $B(t)=\frac{t}{e^t-1}$ (FFT-able). $B[0,\ldots]=[1,-\frac{1}{2},\frac{1}{6},0,-\frac{1}{30},0,\frac{1}{42},\ldots]$ Sums of powers:

$$\sum_{i=1}^{n} n^{m} = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_{k} \cdot (n+1)^{m+1-k}$$

Euler-Maclaurin formula for infinite sums:

$$\sum_{i=m}^{\infty} f(i) = \int_{m}^{\infty} f(x)dx - \sum_{k=1}^{\infty} \frac{B_{k}}{k!} f^{(k-1)}(m)$$

$$\approx \int_{m}^{\infty} f(x)dx + \frac{f(m)}{2} - \frac{f'(m)}{12} + \frac{f'''(m)}{720} + O(f^{(5)}(m))$$

Stirling numbers of the first kind

Number of permutations on n items with k cycles.

$$c(n,k) = c(n-1,k-1) + (n-1)c(n-1,k), \ c(0,0) = 1$$
$$\sum_{k=0}^{n} c(n,k)x^{k} = x(x+1)\dots(x+n-1)$$

c(8, k) = 8, 0, 5040, 13068, 13132, 6769, 1960, 322, 28, 1

Stirling numbers of the second kind

Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-j} \binom{k}{j} j^{n}$$

Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(j) \geq j$, k j:s s.t. $\pi(j) > j$.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n, n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} \binom{n+1}{j} (k+1-j)^{n}$$

Bell numbers

Total number of partitions of n distinct elements. B(n) = 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, For <math>p prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

Labeled unrooted trees

Catalan numbers

```
# on n vertices: n^{n-2}
# on k existing trees of size n_i: n_1 n_2 \cdots n_k n^{k-2}
# with degrees d_i: (n-2)!/((d_1-1)!\cdots(d_n-1)!)
```

$$C_n = \frac{1}{n+1} {2n \choose n} = {2n \choose n} - {2n \choose n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \ C_{n+1} = \frac{2(2n+1)}{n+2}C_n, \ C_{n+1} = \sum_{i=1}^{n} C_i C_{n-i}$$

 $C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$

[noitemsep]sub-diagonal monotone paths in an $n \times n$ grid. strings with n pairs of parenthesis, correctly nested. binary trees with with n+1 leaves (0 or 2 children). ordered trees with n+1 vertices. ways a convex polygon with n+2 sides can be cut into triangles by connecting vertices with straight lines. permutations of [n] with no 3-term increasing subseq.

2.5 Lucas Theorem

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p},$$

where:

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,$$

and:

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that $\binom{n}{n} = 0$ if $m \le n$.

2.6 Multinomial

```
/**
 * Description: Computes $\displaystyle \binom{k_1 +
      \dots + k_n}{k_1, k_2, \dots, k_n} = \frac{(\sum
      k_i)!}{k_1!k_2!...k_n!}$.
 * Status: Tested on kattis:lexicography
 */
#pragma once
long long multinomial(vector<int>& v) {
      long long c = 1, m = v.empty() ? 1 : v[0];
      for (long long i = 1; i < v.size(); i++) {
      for (long long j = 0; j < v[i]; j++) {
            c = c * ++m / (j + 1);
      }
    }
    return c;
}</pre>
```

2.7 Others

Cycles Let $g_S(n)$ be the number of *n*-permutations whose cycle lengths all belong to the set S. Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp\left(\sum_{n \in S} \frac{x^n}{n}\right)$$

Derangements Permutations of a set such that none of the elements appear in their original position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left\lfloor \frac{n!}{e} \right\rfloor$$

Burnside's lemma Given a group G of symmetries and a set X, the number of elements of X up to symmetry equals

$$\frac{1}{|G|} \sum_{g \in G} |X^g|,$$

where X^g are the elements fixed by g (g.x = x).

If f(n) counts "configurations" (of some sort) of length n, we can ignore rotational symmetry using $G = Z_n$ to get

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k|n} f(k)\phi(n/k).$$

2.8 Permutation To Int

2.9 Sigma Function

The Sigma Function is defined as:

$$\sigma_x(n) = \sum_{d|n} d^x$$

when x = 0 is called the divisor function, that counts the number of positive divisors of n.

Now, we are interested in find

$$\sum_{d|n} \sigma_0(d)$$

If n is written as prime factorization:

$$n = \prod_{i=1}^{k} P_i^{e_k}$$

We can demonstrate that:

$$\sum_{d|n} \sigma_0(d) = \prod_{i=1}^k g(e_k + 1)$$

where q(x) is the sum of the first x positive numbers:

$$g(x) = (x * (x + 1))/2$$

3 Data Structures

3.1 Binary Index Tree

```
struct BIT {
    int n:
    int t[2 * N];
    void add(int where, long long what) {
       for (where++; where <= n; where += where &</pre>
            -where) {
           t[where] += what;
    void add(int from, int to, long long what) {
        add(from, what);
        add(to + 1, -what);
    long long query(int where) {
        long long sum = t[0]:
        for (where++: where > 0: where -= where &
             -where) {
            sum += t[where]:
        return sum;
};
```

3.2 Disjoint Set Uninon (DSU)

```
class DSU{
public:
    vector <int> parent;
    void initialize(int n){
        parent.resize(n+1, -1);
}

int findSet(int u){
    while(parent[u] > 0)
        u = parent[u];
    return u;
}

void Union(int u, int v){
    int x = parent[u] + parent[v];
```

```
if(parent[u] > parent[v]){
        parent[v] = x;
        parent[u] = v;
}else{
        parent[u] = x;
        parent[v] = u;
    }
}
```

3.3 Fake Update

```
vector <int> fake bit[MAXN]:
void fake_update(int x, int y, int limit_x){
   for(int i = x; i < limit_x; i += i\&(-i))
       fake bit[i].pb(v):
}
void fake_get(int x, int y){
   for(int i = x; i >= 1; i -= i&(-i))
       fake_bit[i].pb(y);
vector <int> bit[MAXN];
void update(int x, int v, int limit_x, int val){
    for(int i = x; i < limit_x; i += i&(-i)){
       for(int j = lower_bound(fake_bit[i].begin(),
            fake_bit[i].end(), y) -
            fake_bit[i].begin(); j <</pre>
            fake_bit[i].size(); j += j&(-j))
           bit[i][j] = max(bit[i][j], val);
}
int get(int x, int v){
    int ans = 0:
   for(int i = x: i \ge 1: i = i&(-i)){
       for(int j = lower_bound(fake_bit[i].begin(),
            fake bit[i].end(), v) -
            fake_bit[i].begin(); j \ge 1; j = j\&(-j))
           ans = max(ans, bit[i][j]);
    return ans;
int main(){
    int n; cin >> n;
    vector <int> Sx, Sy;
```

```
for(int i = 1; i <= n; i++){</pre>
   cin >> a[i].fi >> a[i].se;
   Sx.pb(a[i].fi);
   Sy.pb(a[i].se);
unique_arr(Sx);
unique_arr(Sy);
// unique all value
for(int i = 1: i \le n: i++)
   a[i].fi = lower_bound(Sx.begin(), Sx.end(),
        a[i].fi) - Sx.begin():
   a[i].se = lower_bound(Sy.begin(), Sy.end(),
        a[i].se) - Sy.begin();
// do fake BIT update and get operator
for(int i = 1: i <= n: i++){
   fake get(a[i].fi-1, a[i].se-1):
   fake update(a[i].fi, a[i].se, (int)Sx.size());
for(int i = 0; i < Sx.size(); i++){</pre>
   fake_bit[i].pb(INT_MIN); // avoid zero
   sort(fake_bit[i].begin(), fake_bit[i].end());
   fake_bit[i].resize(unique(fake_bit[i].begin(),
        fake_bit[i].end()) - fake_bit[i].begin());
   bit[i].resize((int)fake bit[i].size(). 0):
// real update, get operator
int res = 0:
for(int i = 1; i <= n; i++){</pre>
   int maxCurLen = get(a[i].fi-1, a[i].se-1) + 1;
   res = max(res, maxCurLen);
   update(a[i].fi, a[i].se, (int)Sx.size(),
        maxCurLen):
```

3.4 Fenwick Tree

```
template <typename T>
class FenwickTree{
  vector <T> fenw;
  int n;
public:
  void initialize(int _n){
    this->n = _n;
    fenw.resize(n+1);
}
```

```
void update(int id, T val) {
   while (id <= n) {
      fenw[id] += val;
      id += id&(-id);
      }
}

T get(int id){
   T ans{};
   while(id >= 1){
      ans += fenw[id];
      id -= id&(-id);
   }
   return ans;
}
```

3.5 Hash Table

```
/*
 * Micro hash table, can be used as a set.
 * Very efficient vs std::set
 *
 */

const int MN = 1001;
struct ht {
  int _s[(MN + 10) >> 5];
  int len;
  void set(int id) {
    len++;
    _s[id >> 5] |= (1LL << (id & 31));
  }
 bool is_set(int id) {
    return _s[id >> 5] & (1LL << (id & 31));
  }
};</pre>
```

3.6 Range Minimum Query

3.7 STL Treap

```
struct Node {
       Node *1 = 0, *r = 0;
       int val, y, c = 1;
       Node(int val) : val(val), y(rand()) {}
       void recalc();
};
int cnt(Node* n) { return n ? n->c : 0; }
void Node::recalc() { c = cnt(1) + cnt(r) + 1: }
template<class F> void each(Node* n, F f) {
       if (n) { each(n->1, f); f(n->val); each(n->r,
           f): }
}
pair<Node*. Node*> split(Node* n. int k) {
       if (!n) return {}:
       if (cnt(n->1) >= k) { // "n->val >= k" for
           lower bound(k)
              auto pa = split(n->1, k);
              n->1 = pa.second;
              n->recalc();
              return {pa.first, n};
       } else {
              auto pa = split(n->r, k - cnt(n->1) -
                  1): // and just "k"
              n->r = pa.first;
              n->recalc():
              return {n, pa.second};
```

```
Node* merge(Node* 1, Node* r) {
       if (!1) return r;
       if (!r) return 1;
       if (1->y > r->y) {
              1->r = merge(1->r, r);
              1->recalc();
              return 1:
       } else {
              r->1 = merge(1, r->1):
              r->recalc():
              return r:
       }
Node* ins(Node* t. Node* n. int pos) {
       auto pa = split(t, pos);
       return merge(merge(pa.first, n), pa.second):
}
// Example application: move the range [1, r) to index k
void move(Node*& t, int 1, int r, int k) {
       Node *a, *b, *c;
       tie(a,b) = split(t, 1); tie(b,c) = split(b, r -
       if (k \le 1) t = merge(ins(a, b, k), c);
       else t = merge(a, ins(c, b, k - r));
```

3.8 Segment Tree

```
#include <bits/stdc++.h>
using namespace std;

const int N = 1e5 + 10;

int node[4*N];

void modify(int seg, int 1, int r, int p, int val){
    if(1 == r){
        node[seg] += val;
        return;
    }
    int mid = (1 + r)/2;
    if(p <= mid){
        modify(2*seg + 1, 1, mid, p, val);
    }else{
        modify(2*seg + 2, mid + 1, r, p, val);
    }
    node[seg] = node[2*seg + 1] + node[2*seg + 2];</pre>
```

```
}
int sum(int seg, int 1, int r, int a, int b){
   if(1 > b || r < a) return 0;
   if(1 >= a && r <= b) return node[seg];
   int mid = (1 + r)/2;
   return sum(2*seg + 1, 1, mid, a, b) + sum(2*seg + 2, mid + 1, r, a, b);
}</pre>
```

3.9 Sparse Table

```
template <typename T, typename func = function<T(const</pre>
    T. const T)>>
struct SparseTable {
   func calc:
    int n:
    vector<vector<T>> ans:
    SparseTable() {}
    SparseTable(const vector<T>& a. const func& f) :
        n(a.size()), calc(f) {
       int last = trunc(log2(n)) + 1:
       ans.resize(n):
       for (int i = 0; i < n; i++){
           ans[i].resize(last);
       for (int i = 0; i < n; i++){</pre>
           ans[i][0] = a[i];
       for (int j = 1; j < last; j++){</pre>
           for (int i = 0; i \le n - (1 \le j); i++){
              ans[i][j] = calc(ans[i][j-1], ans[i+
                   (1 << (i - 1))][i - 1]);
   T querv(int 1, int r){
       assert(0 <= 1 && 1 <= r && r < n):
       int k = trunc(log2(r - 1 + 1)):
       return calc(ans[1][k], ans[r - (1 << k) +
            1][k]);
   }
};
```

3.10 Trie

```
const int MN = 26; // size of alphabet
const int MS = 100010; // Number of states.
struct trie{
 struct node{
   int c:
   int a[MN];
 node tree[MS]:
 int nodes:
 void clear(){
   tree[nodes].c = 0:
   memset(tree[nodes].a, -1, sizeof tree[nodes].a):
   nodes++:
  void init(){
   nodes = 0:
   clear():
  int add(const string &s, bool query = 0){
   int cur node = 0:
   for(int i = 0; i < s.size(); ++i){</pre>
     int id = gid(s[i]);
     if(tree[cur_node].a[id] == -1){
       if(query) return 0;
       tree[cur_node].a[id] = nodes;
       clear();
     cur_node = tree[cur_node].a[id];
   if(!query) tree[cur_node].c++;
   return tree[cur node].c:
};
```

4 Dynamic Programming Optimization

4.1 Convex Hull Trick

```
#define long long long
#define pll pair <long, long>
#define all(c) c.begin(), c.end()
```

```
#define fastio ios_base::sync_with_stdio(false);
     cin.tie(0)
struct line{
    long a, b;
    line() {};
    line(long a, long b) : a(a), b(b) {};
    bool operator < (const line &A) const {</pre>
              return pll(a,b) < pll(A.a,A.b);</pre>
}:
bool bad(line A. line B. line C){
    return (C.b - B.b) * (A.a - B.a) \le (B.b - A.b) *
        (B.a - C.a):
void addLine(vector<line> &memo. line cur){
    int k = memo.size():
    while (k \ge 2 \&\& bad(memo[k - 2], memo[k - 1],
        cur)){
       memo.pop_back();
    memo.push_back(cur):
long Fn(line A, long x){
    return A.a * x + A.b:
long query(vector<line> &memo, long x){
    int lo = 0, hi = memo.size() - 1;
    while (lo != hi){
       int mi = (lo + hi) / 2:
       if (Fn(memo[mi], x) > Fn(memo[mi + 1], x))
           lo = mi + 1:
       else hi = mi:
    return Fn(memo[lo], x);
const int N = 1e6 + 1:
long dp[N];
int main()
    fastio;
    int n, c; cin >> n >> c;
    vector<line> memo;
    for (int i = 1; i <= n; i++){</pre>
       long val; cin >> val;
```

4.2 Divide and Conquer

```
/**
* recurrence:
     dp[k][i] = min dp[k-1][i] + c[i][i - 1], for all
     j > i;
 * "comp" computes dp[k][i] for all i in O(n log n) (k
     is fixed)
 * Problems:
 * https://icpc.kattis.com/problems/branch
 * http://codeforces.com/contest/321/problem/E
void comp(int 1, int r, int le, int re) {
 if (1 > r) return;
 int mid = (1 + r) >> 1;
 int best = max(mid + 1, le);
 dp[cur][mid] = dp[cur ^ 1][best] + cost(mid, best -
 for (int i = best: i <= re: i++) {</pre>
   if (dp[cur][mid] > dp[cur ^ 1][i] + cost(mid, i -
        1)) {
     best = i:
     dp[cur][mid] = dp[cur^1][i] + cost(mid, i - 1):
 comp(l, mid - 1, le, best):
 comp(mid + 1, r, best, re);
```

5 Geometry

5.1 Closest Pair Problem

```
struct point {
 double x, y;
 int id;
 point() {}
 point (double a, double b) : x(a), y(b) {}
double dist(const point &o, const point &p) {
 double a = p.x - o.x, b = p.y - o.y;
 return sqrt(a * a + b * b):
double cp(vector<point> &p. vector<point> &x.
    vector<point> &y) {
 if (p.size() < 4) {</pre>
   double best = 1e100:
   for (int i = 0: i < p.size(): ++i)</pre>
    for (int j = i + 1; j < p.size(); ++j)
      best = min(best, dist(p[i], p[i])):
   return best:
 int ls = (p.size() + 1) >> 1;
 double l = (p[ls - 1].x + p[ls].x) * 0.5;
 vector<point> xl(ls), xr(p.size() - ls);
 unordered_set<int> left;
 for (int i = 0; i < ls; ++i) {</pre>
   xl[i] = x[i]:
   left.insert(x[i].id);
 for (int i = ls; i < p.size(); ++i) {</pre>
   xr[i - ls] = x[i];
 vector<point> v1. vr:
 vector<point> pl. pr:
 yl.reserve(ls); yr.reserve(p.size() - ls);
 pl.reserve(ls): pr.reserve(p.size() - ls):
 for (int i = 0; i < p.size(); ++i) {</pre>
   if (left.count(v[i].id))
     yl.push_back(y[i]);
   else
     yr.push_back(y[i]);
   if (left.count(p[i].id))
    pl.push_back(p[i]);
   else
     pr.push_back(p[i]);
 double dl = cp(pl, xl, vl);
 double dr = cp(pr, xr, yr);
 double d = min(dl, dr);
```

```
vector<point> vp; vp.reserve(p.size());
 for (int i = 0; i < p.size(); ++i) {</pre>
   if (fabs(y[i].x - 1) < d)
     yp.push_back(y[i]);
 for (int i = 0; i < vp.size(); ++i) {</pre>
   for (int j = i + 1; j < vp.size() && j < i + 7;</pre>
     d = min(d, dist(yp[i], yp[j]));
 return d;
double closest_pair(vector<point> &p) {
 vector<point> x(p.begin(), p.end());
 sort(x.begin(), x.end(), [](const point &a, const
      point &b) {
   return a.x < b.x:</pre>
 }):
 vector<point> y(p.begin(), p.end());
 sort(y.begin(), y.end(), [](const point &a, const
      point &b) {
   return a.v < b.v;
 return cp(p, x, y);
```

5.2 Convex Diameter

```
struct point{
   int x, y;
};

struct vec{
   int x, y;
};

vec operator - (const point &A, const point &B){
   return vec{A.x - B.x, A.y - B.y};
}

int cross(vec A, vec B){
   return A.x*B.y - A.y*B.x;
}

int cross(point A, point B, point C){
   int val = A.x*(B.y - C.y) + B.x*(C.y - A.y) +
        C.x*(A.y - B.y);
   if(val == 0)
        return 0; // coline
```

```
if(val < 0)
       return 1; // clockwise
   return -1; //counter clockwise
vector <point> findConvexHull(vector <point> points){
   vector <point> convex;
    sort(points.begin(), points.end(), [](const point
        &A, const point &B){
       return (A.x == B.x)? (A.v < B.v): (A.x < B.x):
   vector <point> Up, Down;
   point A = points[0], B = points.back();
   Up.push_back(A);
   Down.push_back(A);
   for(int i = 0: i < points.size(): i++){</pre>
       if(i == points.size()-1 || cross(A, points[i].
            B) > 0){
           while(Up.size() > 2 &&
               cross(Up[Up.size()-2], Up[Up.size()-1],
               points[i]) <= 0)
              Up.pop_back();
           Up.push_back(points[i]);
       if(i == points.size()-1 || cross(A, points[i],
            B) < 0){}
           while(Down.size() > 2 &&
                cross(Down[Down.size()-2],
               Down[Down.size()-1], points[i]) >= 0)
              Down.pop_back();
           Down.push_back(points[i]);
       }
   for(int i = 0: i < Up.size(): i++)</pre>
        convex.push back(Up[i]):
   for(int i = Down.size()-2: i > 0: i--)
        convex.push back(Down[i]);
   return convex:
int dist(point A. point B){
   return (A.x - B.x)*(A.x - B.x) + (A.v - B.v)*(A.v -
double findConvexDiameter(vector <point> convexHull){
   int n = convexHull.size():
   int is = 0, is = 0;
   for(int i = 1; i < n; i++){
       if(convexHull[i].y > convexHull[is].y)
```

```
if(convexHull[is].v > convexHull[i].v)
       js = i;
}
int maxd = dist(convexHull[is], convexHull[js]);
int i, maxi, j, maxj;
i = maxi = is;
i = maxi = is;
   int ni = (i+1)%n, nj = (j+1)%n;
   if(cross(convexHull[ni] - convexHull[i].
        convexHull[ni] - convexHull[i]) <= 0){</pre>
   }else{
       i = ni:
   int d = dist(convexHull[i], convexHull[i]):
   if(d > maxd){
       maxd = d:
       maxi = i:
       maxj = j;
}while(i != is || j != js);
return sqrt(maxd);
```

5.3 Pick Theorem

```
struct point{
   11 x, y;
}:
//Pick: S = I + B/2 - 1
ld polygonArea(vector <point> &points){
   int n = (int)points.size();
   ld area = 0.0:
   int j = n-1:
   for(int i = 0; i < n; i++){
       area += (points[j].x + points[i].x) *
            (points[j].y - points[i].y);
       j = i;
    return abs(area/2.0);
}
11 boundary(vector <point> points){
   int n = (int)points.size();
   11 num_bound = 0;
   for(int i = 0; i < n; i++){</pre>
```

5.4 Square

```
typedef long double ld;
const ld eps = 1e-12:
int cmp(ld x, ld v = 0, ld tol = eps) {
   return ( x \le y + tol) ? (x + tol < y) ? -1 : 0 : 1;
struct point{
 ld x, v;
 point(ld a, ld b) : x(a), y(b) {}
 point() {}
}:
struct square{
 ld x1, x2, y1, y2,
    a, b, c;
 point edges[4];
  square(ld _a, ld _b, ld _c) {
   a = _a, b = _b, c = _c;
   x1 = a - c * 0.5:
   x2 = a + c * 0.5:
   y1 = b - c * 0.5;
   y2 = b + c * 0.5;
   edges[0] = point(x1, y1);
   edges[1] = point(x2, y1);
   edges[2] = point(x2, y2);
   edges[3] = point(x1, v2):
 }
};
ld min_dist(point &a, point &b) {
 1d x = a.x - b.x,
    y = a.y - b.y;
 return sqrt(x * x + y * y);
bool point_in_box(square s1, point p) {
 if (cmp(s1.x1, p.x) != 1 && cmp(s1.x2, p.x) != -1 &&
     cmp(s1.y1, p.y) != 1 \&\& cmp(s1.y2, p.y) != -1)
   return true;
 return false;
```

```
bool inside(square &s1, square &s2) {
 for (int i = 0; i < 4; ++i)</pre>
   if (point_in_box(s2, s1.edges[i]))
     return true;
 return false;
bool inside vert(square &s1, square &s2) {
 if ((cmp(s1.y1, s2.y1) != -1 && cmp(s1.y1, s2.y2) !=
     (cmp(s1.y2, s2.y1) != -1 \&\& cmp(s1.y2, s2.y2) !=
   return true;
 return false:
bool inside_hori(square &s1, square &s2) {
 if ((cmp(s1.x1, s2.x1) != -1 \&\& cmp(s1.x1, s2.x2) !=
     (cmp(s1.x2, s2.x1) != -1 \&\& cmp(s1.x2, s2.x2) !=
         1))
   return true;
 return false;
ld min_dist(square &s1, square &s2) {
 if (inside(s1, s2) || inside(s2, s1))
   return 0;
 ld ans = 1e100;
 for (int i = 0; i < 4; ++i)
   for (int j = 0; j < 4; ++j)
     ans = min(ans, min dist(s1.edges[i].
          s2.edges[i])):
 if (inside_hori(s1, s2) || inside_hori(s2, s1)) {
   if (cmp(s1.y1, s2.y2) != -1)
     ans = min(ans, s1.v1 - s2.v2):
   if (cmp(s2.v1, s1.v2) != -1)
     ans = min(ans. s2.v1 - s1.v2):
 if (inside_vert(s1, s2) || inside_vert(s2, s1)) {
   if (cmp(s1.x1, s2.x2) != -1)
     ans = min(ans, s1.x1 - s2.x2);
    else
    if (cmp(s2.x1, s1.x2) != -1)
     ans = min(ans, s2.x1 - s1.x2);
```

```
}
return ans;
}
```

5.5 Triangle

Let a, b, c be length of the three sides of a triangle.

$$p = (a + b + c) * 0.5$$

The inradius is defined by:

$$iR = \sqrt{\frac{(p-a)(p-b)(p-c)}{p}}$$

The radius of its circumcircle is given by the formula:

$$cR = \frac{abc}{\sqrt{(a+b+c)(a+b-c)(a+c-b)(b+c-a)}}$$

6 Graphs

6.1 Bridges

```
struct Graph {
   vector<vector<Edge>> g:
   vector<int> vi, low, d, pi, is_b; // vi = visited
   int bridges_computed;
   int ticks, edges;
   Graph(int n, int m) {
       g.assign(n, vector<Edge>());
      id b.assign(m. 0):
      vi.resize(n):
      low.resize(n):
      d.resize(n):
      pi.resize(n):
       edges = 0;
      bridges_computed = 0;
   void addEge(int u, int v) {
      g[u].push_back(Edge(v, edges));
       g[v].push_back(Edge(u, edges));
       edges++;
   }
```

```
void dfs(int u) {
   vi[u] = true:
   d[u] = low[u] = ticks++;
   for (int i = 0; i < g[u].size(); i++) {</pre>
       int v = g[u][i].to;
       if (v == pi[u]) continue;
       if (!vi[v]) {
           pi[v] = u;
           dfs(v):
           if(d[u] < low[v]) is_b[g[u][i].id] =</pre>
           low[u] = min(low[u], low[v]);
       } else {
           low[u] = min(low[u], low[v]);
   }
// multiple edges from a to b are not allowerd.
// (they could be detected as a bridge).
// if we need to handle this, just count how many
    edges there are from a to b.
void compBridges() {
   fill(pi.begin(), pi.end(), -1);
   fill(vi.begin(), vi.end(), false);
   fill(d.begin(), d.end(), 0);
   fill(low.begin(), low.end(), 0);
   ticks = 0:
   for (int i = 0; i < g.size(); i++)</pre>
       if (!vi[i]) dfs(i);
   bridges_computed = 1;
map<int, vector<Edge>> bridgesTree() {
   if (!bridges computed) compBridges():
   int n = g.size():
   Dsu dsu(n):
   for (int i = 0: i < n: i++)
       for (auto e : g[i])
          if (!is_b[e.id]) dsu.Join(i, e.to);
   map<int. vector<Edge>> tree;
   for (int i = 0: i < n: i++)
       for (auto e : g[i])
           if (is b[e.id])
              tree[dsu.Find(i)].emplace back(dsu.Find(e.to).
   return tree;
```

};

6.2 Dijkstra

```
struct edge {
    int to:
    long long w;
    edge() {}
    edge(int a, long long b) : to(a), w(b) {}
    bool operator<(const edge &e) const {</pre>
       return w > e.w:
}:
typedef <vector<vector<edge>> graph;
const long long inf = 1000000LL * 10000000LL;
pair<vector<int>, vector<long long>> dijkstra(graph& g,
     int start) {
    int n = g.size();
    vector<long long> d(n, inf);
    vector<int> p(n, -1);
    d[start] = 0:
    priority_queue<edge> q;
    q.push(edge(start, 0));
    while (!q.empty()) {
       int node = q.top().to;
       long long dist = q.top().w;
       q.pop();
       if (dist > d[node]) continue;
       for (int i = 0: i < g[node].size(): i++) {</pre>
           int to = g[node][i].to:
           long long w_extra = g[node][i].w;
           if (dist + w extra < d[to]) {</pre>
               p[to] = node:
               d[to] = dist + w extra:
               a.push(edge(to, d[to])):
       }
    return {p, d};
```

6.3 Directed MST

```
struct Edge { int a, b; ll w; };
struct Node { /// lazy skew heap node
    Edge key;
    Node *1, *r;
    ll delta;
    void prop() {
        key.w += delta;
        if (l) l->delta += delta;
```

```
if (r) r->delta += delta;
              delta = 0:
       Edge top() { prop(); return key; }
};
Node *merge(Node *a, Node *b) {
       if (!a || !b) return a ?: b;
       a->prop(), b->prop();
       if (a->kev.w > b->kev.w) swap(a, b):
       swap(a->1, (a->r = merge(b, a->r)));
       return a:
}
void pop(Node*& a) { a \rightarrow prop(): a = merge(a \rightarrow 1, a \rightarrow r): }
pair<11, vi> dmst(int n, int r, vector<Edge>& g) {
       RollbackUF uf(n):
       vector<Node*> heap(n):
       for (Edge e : g) heap[e.b] = merge(heap[e.b],
            new Node(e):
       11 res = 0:
       vi seen(n, -1), path(n), par(n):
       seen[r] = r;
       vector<Edge> Q(n), in(n, \{-1,-1\}), comp;
       deque<tuple<int, int, vector<Edge>>> cycs;
       rep(s,0,n) {
              int u = s, qi = 0, w;
              while (seen[u] < 0) {
                      if (!heap[u]) return {-1,{}};
                      Edge e = heap[u]->top();
                      heap[u]->delta -= e.w,
                           pop(heap[u]);
                      Q[qi] = e, path[qi++] = u,
                           seen[u] = s;
                      res += e.w, u = uf.find(e.a);
                      if (seen[u] == s) { /// found
                           cvcle, contract
                             Node* cvc = 0:
                             int end = gi, time =
                                  uf.time():
                              do cyc = merge(cyc, heap[w
                                   = path[--qi]]);
                              while (uf.join(u, w)):
                             u = uf.find(u), heap[u] =
                                   cvc. seen[u] = -1:
                             cycs.push_front({u, time,
                                   {&O[ai], &O[end]}}):
                      }
              rep(i,0,qi) in[uf.find(Q[i].b)] = Q[i];
       for (auto& [u,t,comp] : cycs) { // restore sol
            (optional)
```

6.4 Edge Coloring

```
vi edgeColoring(int N, vector<pii> eds) {
       vi cc(N + 1), ret(sz(eds)), fan(N), free(N),
       for (pii e : eds) ++cc[e.first], ++cc[e.second]:
       int u. v. ncols = *max element(all(cc)) + 1:
       vector<vi> adj(N, vi(ncols, -1));
       for (pii e : eds) {
              tie(u, v) = e;
              fan[0] = v;
              loc.assign(ncols, 0);
              int at = u, end = u, d, c = free[u], ind
                  = 0. i = 0:
              while (d = free[v], !loc[d] && (v =
                   adi[u][d]) != -1)
                     loc[d] = ++ind, cc[ind] = d,
                          fan[ind] = v:
              cc[loc[d]] = c:
              for (int cd = d; at != -1; cd ^= c ^ d,
                   at = adj[at][cd])
                     swap(adj[at][cd], adj[end =
                          at][cd ^ c ^ d]);
              while (adi[fan[i]][d] != -1) {
                     int left = fan[i], right =
                          fan[++i], e = cc[i]:
                     adi[u][e] = left:
                     adi[left][e] = u:
                     adi[right][e] = -1:
                     free[right] = e:
              adj[u][d] = fan[i];
              adi[fan[i]][d] = u;
              for (int y : {fan[0], u, end})
                     for (int& z = free[y] = 0;
                          adi[v][z] != -1; z++);
       rep(i,0,sz(eds))
              for (tie(u, v) = eds[i]; adj[u][ret[i]]
                   != v;) ++ret[i];
       return ret;
```

}

6.5 Eulerian Path

```
struct DirectedEulerPath
       int n:
       vector<vector<int> > g;
       vector<int> path;
       void init(int n){
              n = n:
              g = vector < vector < int > > (n + 1,
                   vector<int> ());
              path.clear():
       void add_edge(int u, int v){
              g[u].push_back(v);
       void dfs(int u)
              while(g[u].size())
                      int v = g[u].back();
                      g[u].pop_back();
                      dfs(v):
              path.push back(u):
       }
       bool getPath(){
              int ctEdges = 0:
              vector<int> outDeg, inDeg;
              outDeg = inDeg = vector<int> (n + 1, 0):
              for(int i = 1: i <= n: i++)
                      ctEdges += g[i].size():
                      outDeg[i] += g[i].size();
                      for(auto &u:g[i])
                             inDeg[u]++;
              int ctMiddle = 0, src = 1;
              for(int i = 1; i <= n; i++)</pre>
                      if(abs(inDeg[i] - outDeg[i]) > 1)
                             return 0;
                      if(inDeg[i] == outDeg[i])
                             ctMiddle++;
                      if(outDeg[i] > inDeg[i])
```

6.6 Floyd - Warshall

6.7 Ford - Bellman

6.8 Gomory Hu

6.9 Karp Min Mean Cycle

```
/**
 * Finds the min mean cycle, if you need the max mean
    cycle
 * just add all the edges with negative cost and print
 * ans * -1
 *
 * test: uva, 11090 - Going in Cycle!!
 * */
const int MN = 1000;
struct edge{
  int v;
```

```
long long w;
 edge()\{\} edge(int v, int w) : v(v), w(w) \{\}
long long d[MN][MN];
// This is a copy of g because increments the size
// pass as reference if this does not matter.
int karp(vector<vector<edge> > g) {
 int n = g.size();
  g.resize(n + 1): // this is important
  for (int i = 0: i < n: ++i)</pre>
   if (!g[i].empty())
     g[n].push_back(edge(i,0));
  for(int i = 0:i < n:++i)
   fill(d[i].d[i]+(n+1).INT MAX):
  d[n - 1][0] = 0:
  for (int k = 1; k \le n; ++k) for (int u = 0; u \le n;
      ++u) {
    if (d[u][k - 1] == INT_MAX) continue;
   for (int i = g[u].size() - 1; i >= 0; --i)
     d[g[u][i].v][k] = min(d[g[u][i].v][k], d[u][k -
          1] + g[u][i].w);
  bool flag = true;
  for (int i = 0; i < n && flag; ++i)</pre>
   if (d[i][n] != INT_MAX)
     flag = false:
  if (flag) {
   return true: // return true if there is no a cycle.
  double ans = 1e15:
  for (int u = 0: u + 1 < n: ++u) {
   if (d[u][n] == INT MAX) continue:
   double W = -1e15:
   for (int k = 0; k < n; ++k)
     if (d[u][k] != INT MAX)
       W = \max(W, (double)(d[u][n] - d[u][k]) / (n -
            k)):
    ans = min(ans, W);
```

```
// printf("%.21f\n", ans);
cout << fixed << setprecision(2) << ans << endl;
return false;
}</pre>
```

6.10 Konig's Theorem

In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover

6.11 LCA

```
#include "../Data Structures/RMQ.h"
struct LCA {
       int T = 0:
       vi time, path, ret;
       RMQ<int> rmq;
       LCA(vector<vi>& C) : time(sz(C)),
           rmq((dfs(C,0,-1), ret)) {}
       void dfs(vector<vi>& C, int v, int par) {
              time[v] = T++;
              for (int y : C[v]) if (y != par) {
                     path.push back(v).
                          ret.push_back(time[v]);
                     dfs(C, y, v);
              }
       int lca(int a, int b) {
              if (a == b) return a;
              tie(a, b) = minmax(time[a], time[b]):
              return path[rmq.query(a, b)];
       //dist(a,b){return depth[a] + depth[b] -
            2*depth[lca(a,b)];}
};
```

6.12 Math

Number of Spanning Trees

Create an $N \times N$ matrix mat, and for each edge $a \rightarrow b \in G$, do mat[a][b]--, mat[b][b]++ (and mat[b][a]--,

mat[a][a]++ if G is undirected). Remove the ith row and column and take the determinant; this yields the number of directed spanning trees rooted at i (if G is undirected, remove any row/column).

Erdős-Gallai theorem

A simple graph with node degrees $d_1 \geq \cdots \geq d_n$ exists iff $d_1 + \cdots + d_n$ is even and for every $k = 1 \dots n$,

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k).$$

6.13 Push Relabel

```
struct PushRelabel {
       struct Edge {
              int dest. back:
              11 f, c;
       vector<vector<Edge>> g;
       vector<ll> ec;
       vector<Edge*> cur;
       vector<vi> hs; vi H;
      PushRelabel(int n) : g(n), ec(n), cur(n),
           hs(2*n), H(n) {}
       void addEdge(int s, int t, ll cap, ll rcap=0) {
              if (s == t) return:
              g[s].push_back(\{t, sz(g[t]), 0, cap\});
              g[t].push back({s. sz(g[s])-1. 0. rcap}):
      }
       void addFlow(Edge& e, ll f) {
              Edge &back = g[e.dest][e.back];
              if (!ec[e.dest] && f)
                   hs[H[e.dest]].push back(e.dest):
              e.f += f: e.c -= f: ec[e.dest] += f:
              back.f -= f: back.c += f: ec[back.dest]
                   -= f:
      11 calc(int s. int t) {
              int v = sz(g); H[s] = v; ec[t] = 1;
              vi co(2*v); co[0] = v-1;
              rep(i,0,v) cur[i] = g[i].data();
              for (Edge& e : g[s]) addFlow(e, e.c);
              for (int hi = 0::) {
                     while (hs[hi].empty()) if (!hi--)
                          return -ec[s]:
                     int u = hs[hi].back();
                          hs[hi].pop_back();
```

```
while (ec[u] > 0) // discharge u
                             if (cur[u] == g[u].data()
                                  + sz(g[u])) {
                                    H[u] = 1e9;
                                    for (Edge& e :
                                         g[u]) if (e.c
                                         && H[u] >
                                         H[e.dest]+1)
                                           H[u] =
                                                H[e.dest]+1.
                                                cur[u]
                                                = &e:
                                    if (++co[H[u]].
                                         !--co[hi] &&
                                         hi < v)
                                           rep(i.0.v)
                                                if (hi
                                                < H[i]
                                                && H[i]
                                                < v)
                                                   --co[H[i]]
                                                       H[i]
                                    hi = H[u]:
                             } else if (cur[u]->c &&
                                  H[u] ==
                                  H[cur[u]->dest]+1)
                                    addFlow(*cur[u],
                                         min(ec[u],
                                         cur[u]->c));
                             else ++cur[u]:
              }
       bool leftOfMinCut(int a) { return H[a] >=
            sz(g); }
};
6.14 SCC Kosaraju
```

```
// SCC = Strongly Connected Components
struct SCC {
   vector<vector<int>> g, gr;
   vector<bool> used:
   vector<int> order, component;
   int total_components;
   SCC(vector<vector<int>>& adj) {
```

```
g = adi;
       int n = g.size();
       gr.resize(n);
       for (int i = 0; i < n; i++)</pre>
          for (auto to : g[i])
              gr[to].push_back(i);
       used.assign(n, false);
       for (int i = 0: i < n: i++)
       if (!used[i])
           GenTime(i):
       used.assign(n. false):
       component.assign(n, -1);
       total components = 0:
       for (int i = n - 1; i \ge 0; i--) {
           int v = order[i]:
           if (!used[v]) {
              vector<int> cur component:
              Dfs(cur component, v):
              for (auto node : cur component)
                  component[node] = total_components;
   void GenTime(int node) {
       used[node] = true;
       for (auto to : g[node])
           if (!used[to])
              GenTime(to);
       order.push_back(node);
   void Dfs(vector<int>& cur. int node) {
       used[node] = true:
       cur.push_back(node);
       if (!used[to])
           Dfs(cur. to):
   vector<vector<int>> CondensedGraph() {
       vector<vector<int>> ans(total components):
       for (int i = 0: i < int(g.size()): i++) {</pre>
           for (int to : g[i]) {
              int u = component[i]. v = component[to]:
              if (u != v)
              ans[u].push_back(v);
       }
       return ans;
};
```

6.15 Topological Sort

```
vi topoSort(const vector<vi>& gr) {
    vi indeg(sz(gr)), ret;
    for (auto& li : gr) for (int x : li) indeg[x]++;
    queue<int> q; // use priority_queue for lexic.
        largest ans.
    rep(i,0,sz(gr)) if (indeg[i] == 0) q.push(i);
    while (!q.empty()) {
        int i = q.front(); // top() for priority
            queue
        ret.push_back(i);
        q.pop();
        for (int x : gr[i])
            if (--indeg[x] == 0) q.push(x);
    }
    return ret;
}
```

7 Misc

7.1 Dates

```
// Time - Leap years
// A[i] has the accumulated number of days from months
    previous to i
const int A[13] = \{ 0, 0, 31, 59, 90, 120, 151, 181, \dots \}
    212, 243, 273, 304, 334 };
// same as A. but for a leap year
const int B[13] = \{ 0, 0, 31, 60, 91, 121, 152, 182, \dots \}
     213, 244, 274, 305, 335 }:
// returns number of leap years up to, and including, v
int leap_years(int y) { return y / 4 - y / 100 + y /
    400: }
bool is_leap(int y) { return y % 400 == 0 || (y % 4 ==
    0 && v % 100 != 0); }
// number of days in blocks of years
const int p400 = 400*365 + leap_years(400);
const int p100 = 100*365 + leap_years(100);
const int p4 = 4*365 + 1;
const int p1 = 365;
int date_to_days(int d, int m, int y)
 return (y - 1) * 365 + leap_years(y - 1) +
      (is_{pap}(y) ? B[m] : A[m]) + d;
```

```
void days_to_date(int days, int &d, int &m, int &y)
 bool top100; // are we in the top 100 years of a 400
      block?
 bool top4; // are we in the top 4 years of a 100
      block?
 bool top1; // are we in the top year of a 4 block?
 top100 = top4 = top1 = false:
 y += ((days-1) / p400) * 400;
 d = (davs-1) \% p400 + 1:
 if (d > p100*3) top100 = true, d = 3*p100, y += 300;
 else v += ((d-1) / p100) * 100, d = (d-1) % p100 + 1:
 if (d > p4*24) top4 = true, d = 24*p4, y += 24*4;
 else v += ((d-1) / p4) * 4, d = (d-1) % p4 + 1:
 if (d > p1*3) top1 = true, d = p1*3, y += 3;
 else y += (d-1) / p1, d = (d-1) \% p1 + 1;
 const int *ac = top1 && (!top4 || top100) ? B : A;
 for (m = 1; m < 12; ++m) if (d \le ac[m + 1]) break;
 d = ac[m];
```

8 Number Theory

8.1 Chinese Remainder Theorem

8.2 Convolution

```
typedef long long int LL:
typedef pair<LL, LL> PLL;
inline bool is_pow2(LL x) {
 return (x & (x-1)) == 0:
inline int ceil log2(LL x) {
 int ans = 0:
 while (x != 0) {
  x >>= 1:
   ans++:
 return ans;
/* Returns the convolution of the two given vectors in
    time proportional to n*log(n).
* The number of roots of unity to use nroots_unity
     must be set so that the product of the first
* nroots_unity primes of the vector nth_roots_unity is
     greater than the maximum value of the
* convolution. Never use sizes of vectors bigger than
     2^24, if you need to change the values of
* the nth roots of unity to appropriate primes for
     those sizes.
vector<LL> convolve(const vector<LL> &a. const
    vector<LL> &b. int nroots unitv = 2) {
 int N = 1 << ceil_log2(a.size() + b.size());</pre>
 vector<LL> ans(N.O), fA(N), fB(N), fC(N);
 LL modulo = 1:
 for (int times = 0: times < nroots unity: times++) {</pre>
   fill(fA.begin(), fA.end(), 0);
   fill(fB.begin(), fB.end(), 0);
   for (int i = 0: i < a.size(): i++) fA[i] = a[i]:</pre>
   for (int i = 0; i < b.size(); i++) fB[i] = b[i];</pre>
   LL prime = nth_roots_unity[times].first;
   LL inv_modulo = mod_inv(modulo % prime, prime);
   LL normalize = mod_inv(N, prime);
   ntfft(fA, 1, nth_roots_unity[times]);
   ntfft(fB, 1, nth_roots_unity[times]);
   for (int i = 0; i < N; i++) fC[i] = (fA[i] * fB[i])</pre>
   ntfft(fC, -1, nth_roots_unity[times]);
   for (int i = 0; i < N; i++) {</pre>
```

```
LL curr = (fC[i] * normalize) % prime;
LL k = (curr - (ans[i] % prime) + prime) % prime;
k = (k * inv_modulo) % prime;
ans[i] += modulo * k;
}
modulo *= prime;
}
return ans;
```

8.3 Diophantine Equations

```
long long gcd(long long a, long long b, long long &x,
    long long &y) {
 if (a == 0) {
   x = 0:
   y = 1;
   return b;
 long long x1, v1;
 long long d = gcd(b \% a, a, x1, y1);
 x = v1 - (b / a) * x1;
 y = x1;
 return d;
bool find_any_solution(long long a, long long b, long
    long c, long long &x0,
   long long &y0, long long &g) {
 g = gcd(abs(a), abs(b), x0, y0);
 if (c % g) {
   return false:
 x0 *= c / g;
 y0 *= c / g;
 if (a < 0) x0 = -x0:
 if (b < 0) y0 = -y0;
 return true:
void shift_solution(long long &x, long long &y, long
    long a, long long b,
   long long cnt) {
 x += cnt * b;
 y -= cnt * a;
long long find_all_solutions(long long a, long long b,
    long long c,
   long long minx, long long maxx, long long miny,
```

```
long long maxy) {
long long x, y, g;
if (!find_any_solution(a, b, c, x, y, g)) return 0;
b /= g;
long long sign_a = a > 0 ? +1 : -1;
long long sign_b = b > 0 ? +1 : -1;
shift_solution(x, y, a, b, (minx - x) / b);
if (x < minx) shift_solution(x, y, a, b, sign_b);</pre>
if (x > maxx) return 0:
long long lx1 = x;
shift_solution(x, y, a, b, (maxx - x) / b);
if (x > maxx) shift solution(x, v, a, b, -sign b):
long long rx1 = x:
shift solution(x, v, a, b, -(minv - v) / a):
if (y < miny) shift_solution(x, y, a, b, -sign_a);</pre>
if (y > maxy) return 0;
long long 1x2 = x;
shift_solution(x, y, a, b, -(maxy - y) / a);
if (y > maxy) shift_solution(x, y, a, b, sign_a);
long long rx2 = x;
if (1x2 > rx2) swap(1x2, rx2);
long long lx = max(lx1, lx2);
long long rx = min(rx1, rx2);
if (lx > rx) return 0;
return (rx - lx) / abs(b) + 1;
```

8.4 Discrete Logarithm

```
// Computes x which a ^ x = b mod n.
long long d_log(long long a, long long b, long long n) {
  long long m = ceil(sqrt(n));
  long long aj = 1;
  map<long long, long long> M;
  for (int i = 0; i < m; ++i) {
    if (!M.count(aj))
        M[aj] = i;
        aj = (aj * a) % n;
  }
long long coef = mod_pow(a, n - 2, n);</pre>
```

```
coef = mod_pow(coef, m, n);
// coef = a ^ (-m)
long long gamma = b;
for (int i = 0; i < m; ++i) {
   if (M.count(gamma)) {
      return i * m + M[gamma];
   } else {
      gamma = (gamma * coef) % n;
   }
}
return -1;
}</pre>
```

8.5 Ext Euclidean

```
void ext_euclid(long long a, long long b, long long &x,
        long long &y, long long &g) {
    x = 0, y = 1, g = b;
    long long m, n, q, r;
    for (long long u = 1, v = 0; a != 0; g = a, a = r) {
        q = g / a, r = g % a;
        m = x - u * q, n = y - v * q;
        x = u, y = v, u = m, v = n;
    }
}
```

8.6 Highest Exponent Factorial

```
int highest_exponent(int p, const int &n){
  int ans = 0;
  int t = p;
  while(t <= n){
    ans += n/t;
    t*=p;
  }
  return ans;
}</pre>
```

8.7 Miller - Rabin

```
// check as in Miller Rabin Primality Test described
 long long u = n - 1;
 int t = 0;
 while (u % 2 == 0) {
   t++:
   u >>= 1;
 long long next = mod_pow(a, u, n);
 if (next == 1) return false:
 long long last;
 for (int i = 0: i < t: ++i) {
   last = next:
   next = mod_mul(last, last, n);
   if (next == 1) {
    return last != n - 1:
 return next != 1:
// Checks if a number is prime with prob 1 - 1 / (2 ^)
// D(miller_rabin(999999999999997LL) == 1);
// D(miller_rabin(999999999971LL) == 1);
// D(miller_rabin(7907) == 1);
bool miller_rabin(long long n, int it = rounds) {
 if (n <= 1) return false;
 if (n == 2) return true;
 if (n % 2 == 0) return false;
 for (int i = 0; i < it; ++i) {</pre>
   long long a = rand() \% (n - 1) + 1;
   if (witness(a, n)) {
     return false:
 return true;
```

8.8 Mod Integer

```
template < class T, T mod>
struct mint_t {
   T val;
   mint_t() : val(0) {}
   mint_t(T v) : val(v % mod) {}

mint_t operator + (const mint_t& o) const {
   return (val + o.val) % mod;
  }
  mint_t operator - (const mint_t& o) const {
```

```
return (val - o.val) % mod;
}
mint_t operator * (const mint_t& o) const {
  return (val * o.val) % mod;
}
};

typedef mint_t<long long, 998244353> mint;
```

8.9 Mod Inv

```
long long mod_inv(long long n, long long m) {
  long long x, y, gcd;
  ext_euclid(n, m, x, y, gcd);
  if (gcd != 1)
   return 0;
  return (x + m) % m;
}
```

8.10 Mod Mul

```
// Computes (a * b) % mod
long long mod_mul(long long a, long long b, long long
    mod) {
    long long x = 0, y = a % mod;
    while (b > 0) {
        if (b & 1)
            x = (x + y) % mod;
        y = (y * 2) % mod;
        b /= 2;
    }
    return x % mod;
}
```

8.11 Mod Pow

```
// Computes ( a ^ exp ) % mod.
long long mod_pow(long long a, long long exp, long long
    mod) {
    long long ans = 1;
    while (exp > 0) {
        if (exp & 1)
            ans = mod_mul(ans, a, mod);
        a = mod_mul(a, a, mod);
        exp >>= 1;
```

```
}
return ans;
}
```

8.12 Number Theoretic Transform

```
typedef long long int LL;
typedef pair<LL, LL> PLL;
/* The following vector of pairs contains pairs (prime.
 * where the prime has an Nth root of unity for N being
     a power of two.
 * The generator is a number g s.t g^(p-1)=1 (mod p)
 * but is different from 1 for all smaller powers */
vector<PLL> nth roots unity {
 {1224736769,330732430},{1711276033,927759239},{167772161,1674
  {469762049,343261969},{754974721,643797295},{1107296257,8838}
PLL ext_euclid(LL a, LL b) {
 if (b == 0)
   return make_pair(1,0);
 pair<LL,LL> rc = ext_euclid(b, a % b);
 return make_pair(rc.second, rc.first - (a / b) *
      rc.second);
//returns -1 if there is no unique modular inverse
LL mod_inv(LL x, LL modulo) {
 PLL p = ext_euclid(x, modulo);
 if ((p.first * x + p.second * modulo) != 1)
   return -1:
 return (p.first+modulo) % modulo;
//Number theory fft. The size of a must be a power of 2
void ntfft(vector<LL> &a, int dir, const PLL
    &root unity) {
 int n = a.size();
 LL prime = root unity.first:
 LL basew = mod_pow(root_unity.second, (prime-1) / n,
 if (dir < 0) basew = mod_inv(basew, prime);</pre>
 for (int m = n; m >= 2; m >>= 1) {
   int mh = m >> 1;
   LL w = 1:
   for (int i = 0; i < mh; i++) {</pre>
     for (int j = i; j < n; j += m) {
      int k = j + mh;
       LL x = (a[j] - a[k] + prime) % prime;
```

```
a[j] = (a[j] + a[k]) % prime;
a[k] = (w * x) % prime;
}
w = (w * basew) % prime;
}
basew = (basew * basew) % prime;
}
int i = 0;
for (int j = 1; j < n - 1; j++) {
  for (int k = n >> 1; k > (i ^= k); k >>= 1);
  if (j < i) swap(a[i], a[j]);
}</pre>
```

8.13 Pollard Rho Factorize

```
long long pollard_rho(long long n) {
 long long x, y, i = 1, k = 2, d;
 x = y = rand() \% n;
 while (1) {
   ++i:
   x = mod_mul(x, x, n);
   x += 2;
   if (x \ge n) x = n;
   if (x == y) return 1;
   d = \_gcd(abs(x - y), n);
   if (d != 1) return d;
   if (i == k) {
    y = x;
     k *= 2:
 return 1;
// Returns a list with the prime divisors of n
vector<long long> factorize(long long n) {
 vector<long long> ans:
 if (n == 1)
   return ans:
 if (miller_rabin(n)) {
   ans.push_back(n);
 } else {
   long long d = 1;
   while (d == 1)
     d = pollard_rho(n);
   vector<long long> dd = factorize(d);
   ans = factorize(n / d):
   for (int i = 0; i < dd.size(); ++i)</pre>
     ans.push_back(dd[i]);
```

```
}
return ans;
}
```

8.14 Primes

```
namespace primes {
 const int MP = 100001:
 bool sieve[MP]:
 long long primes[MP]:
 int num p:
 void fill sieve() {
   num_p = 0;
   sieve[0] = sieve[1] = true:
   for (long long i = 2; i < MP; ++i) {</pre>
     if (!sieve[i]) {
       primes[num p++] = i:
       for (long long j = i * i; j < MP; j += i)
         sieve[i] = true;
   }
 }
 // Finds prime numbers between a and b, using basic
      primes up to sqrt(b)
 // a must be greater than 1.
 vector<long long> seg_sieve(long long a, long long b)
   long long ant = a;
   a = max(a, 3LL):
   vector<bool> pmap(b - a + 1);
   long long sqrt_b = sqrt(b);
   for (int i = 0; i < num_p; ++i) {</pre>
     long long p = primes[i];
     if (p > sqrt_b) break;
     long long j = (a + p - 1) / p;
     for (long long v = (i == 1) ? p + p : i * p : v <=
          f(\alpha =+ v : d)
       pmap[v - a] = true:
   vector<long long> ans;
   if (ant == 2) ans.push_back(2);
   int start = a % 2 ? 0 : 1;
   for (int i = start, I = b - a + 1; i < I; i += 2)</pre>
     if (pmap[i] == false)
       ans.push_back(a + i);
   return ans;
 vector<pair<int, int>> factor(int n) {
```

```
vector<pair<int, int>> ans;
if (n == 0) return ans;
for (int i = 0; primes[i] * primes[i] <= n; ++i) {
    if ((n % primes[i]) == 0) {
        int expo = 0;
        while ((n % primes[i]) == 0) {
            expo++;
            n /= primes[i];
        }
        ans.emplace_back(primes[i], expo);
    }
}

if (n > 1) {
    ans.emplace_back(n, 1);
}
    return ans;
}
```

8.15 Totient Sieve

```
for (int i = 1; i < MN; i++)
  phi[i] = i;

for (int i = 1; i < MN; i++)
  if (!sieve[i]) // is prime
  for (int j = i; j < MN; j += i)
    phi[j] -= phi[j] / i;</pre>
```

8.16 Totient

```
long long totient(long long n) {
   if (n == 1) return 0;
   long long ans = n;
   for (int i = 0; primes[i] * primes[i] <= n; ++i) {
      if ((n % primes[i]) == 0) {
        while ((n % primes[i]) == 0) n /= primes[i];
        ans -= ans / primes[i];
    }
   }
   if (n > 1) {
      ans -= ans / n;
   }
   return ans;
}
```

9 Strings

9.1 Incremental Aho Corasick

```
class IncrementalAhoCorasic {
 static const int Alphabets = 26:
 static const int AlphabetBase = 'a';
 struct Node {
   Node *fail:
   Node *next[Alphabets];
   int sum:
   Node(): fail(NULL), next{}, sum(0) { }
 struct String {
   string str;
   int sign;
 }:
public:
 //totalLen = sum of (len + 1)
 void init(int totalLen) {
   nodes.resize(totalLen):
   nNodes = 0;
   strings.clear();
   roots.clear();
   sizes.clear();
   que.resize(totalLen);
 void insert(const string &str. int sign) {
   strings.push_back(String{ str, sign });
   roots.push_back(nodes.data() + nNodes);
   sizes.push_back(1);
   nNodes += (int)str.size() + 1:
   auto check = [&]() { return sizes.size() > 1 &&
        sizes.end()[-1] == sizes.end()[-2]; }:
   if(!check())
     makePMA(strings.end() - 1, strings.end(),
         roots.back(), que):
   while(check()) {
     int m = sizes.back();
     roots.pop_back();
     sizes.pop_back();
     sizes.back() += m;
     if(!check())
      makePMA(strings.end() - m * 2, strings.end(),
           roots.back(), que);
 }
 int match(const string &str) const {
```

```
int res = 0;
   for(const Node *t : roots)
     res += matchPMA(t, str);
   return res:
 }
private:
 static void makePMA(vector<String>::const_iterator
      begin, vector<String>::const_iterator end, Node
      *nodes, vector<Node*> &que) {
   int nNodes = 0:
   Node *root = new(&nodes[nNodes ++]) Node():
   for(auto it = begin: it != end: ++ it) {
     Node *t = root:
     for(char c : it->str) {
       Node *&n = t->next[c - AlphabetBase]:
       if(n == nullptr)
        n = new(&nodes[nNodes ++]) Node():
       t = n:
     t->sum += it->sign;
   int qt = 0;
   for(Node *&n : root->next) {
     if(n != nullptr) {
      n->fail = root;
       que[qt ++] = n;
     } else {
       n = root;
   for(int qh = 0; qh != qt; ++ qh) {
     Node *t = que[qh];
     int a = 0:
     for(Node *n : t->next) {
      if(n != nullptr) {
         que[qt ++] = n;
         Node *r = t-fail:
         while(r->next[a] == nullptr)
          r = r->fail:
         n->fail = r->next[a];
         n->sum += r->next[a]->sum:
       ++ a;
 static int matchPMA(const Node *t, const string &str)
   int res = 0;
   for(char c : str) {
     int a = c - AlphabetBase;
```

```
while(t->next[a] == nullptr)
      t = t->fail:
     t = t->next[a]:
     res += t->sum:
   return res;
  vector<Node> nodes:
 int nNodes:
 vector<String> strings:
 vector<Node*> roots:
 vector<int> sizes:
 vector<Node*> que:
int main() {
 int m:
 while("scanf("%d", &m)) {
   IncrementalAhoCorasic iac:
   iac.init(600000);
   rep(i, m) {
     int ty;
     char s[300001];
     scanf("%d%s", &tv, s);
     if(tv == 1) {
      iac.insert(s, +1);
     } else if(ty == 2) {
      iac.insert(s, -1);
     } else if(tv == 3) {
      int ans = iac.match(s);
       printf("%d\n", ans);
      fflush(stdout);
     } else {
       abort():
   }
 }
 return 0:
```

9.2 KMP

```
vi pi(const string& s) {
    vi p(sz(s));
    rep(i,1,sz(s)) {
        int g = p[i-1];
        while (g && s[i] != s[g]) g = p[g-1];
        p[i] = g + (s[i] == s[g]);
}
```

9.3 Minimal String Rotation

```
// Lexicographically minimal string rotation
int lmsr() {
 string s;
 cin >> s;
 int n = s.size();
 vector<int> f(s.size(), -1);
 for (int j = 1; j < 2 * n; ++j) {
   int i = f[j - k - 1];
   while (i != -1 && s[j] != s[k + i + 1]) {
    if (s[j] < s[k + i + 1])
      k = j - i - 1;
    i = f[i];
   if (i == -1 \&\& s[j] != s[k + i + 1]) {
     if (s[i] < s[k + i + 1]) {
      k = j;
    f[j - k] = -1;
   } else {
     f[j - k] = i + 1;
 }
 return k;
```

9.4 Suffix Array

```
const int MAXN = 200005;

const int MAX_DIGIT = 256;
void countingSort(vector<int>& SA, vector<int>& RA, int
    k = 0) {
    int n = SA.size();
}
```

```
vector<int> cnt(max(MAX_DIGIT, n), 0);
   for (int i = 0; i < n; i++)</pre>
       if (i + k < n)
           cnt[RA[i + k]]++;
           cnt[0]++;
    for (int i = 1; i < cnt.size(); i++)</pre>
       cnt[i] += cnt[i - 1];
    vector<int> tempSA(n);
   for (int i = n - 1; i >= 0; i--)
       if (SA[i] + k < n)
           tempSA[--cnt[RA[SA[i] + k]]] = SA[i];
           tempSA[--cnt[0]] = SA[i];
    SA = tempSA:
vector <int> constructSA(string s) {
   int n = s.length();
   vector <int> SA(n);
   vector <int> RA(n):
    vector <int> tempRA(n);
   for (int i = 0; i < n; i++) {</pre>
       RA[i] = s[i];
       SA[i] = i;
   for (int step = 1; step < n; step <<= 1) {</pre>
       countingSort(SA, RA, step);
       countingSort(SA, RA, 0);
       int c = 0;
       tempRA[SA[0]] = c;
       for (int i = 1; i < n; i++) {
           if (RA[SA[i]] == RA[SA[i - 1]] \&\& RA[SA[i] +
                step] == RA[SA[i - 1] + step])
                  tempRA[SA[i]] = tempRA[SA[i - 1]];
               tempRA[SA[i]] = tempRA[SA[i - 1]] + 1;
       RA = tempRA;
       if (RA[SA[n-1]] == n-1) break:
   return SA:
vector<int> computeLCP(const string& s, const
    vector<int>& SA) {
    int n = SA.size();
    vector<int> LCP(n), PLCP(n), c(n, 0);
   for (int i = 0; i < n; i++)</pre>
       c[SA[i]] = i;
   int k = 0;
    for (int j, i = 0; i < n-1; i++) {
       if(c[i] - 1 < 0)
```

9.5 Suffix Automation

```
* Suffix automaton:
 * This implementation was extended to maintain
      (online) the
 * number of different substrings. This is equivalent
 * the number of paths from the initial state to all
     the other
 * The overall complexity is O(n)
 * can be tested here:
     https://www.urionlinejudge.com.br/judge/en/problems/view/
struct state {
 int len, link;
 long long num_paths;
 map<int, int> next;
const int MN = 200011:
state sa[MN << 1]:
int sz. last:
long long tot_paths;
void sa_init() {
 sz = 1;
 last = 0;
 sa[0].len = 0;
 sa[0].link = -1;
 sa[0].next.clear();
 sa[0].num_paths = 1;
 tot_paths = 0;
```

```
void sa_extend(int c) {
 int cur = sz++:
 sa[cur].len = sa[last].len + 1;
 sa[cur].next.clear();
 sa[cur].num_paths = 0;
 for (p = last; p != -1 && !sa[p].next.count(c); p =
      sa[p].link) {
   sa[p].next[c] = cur;
   sa[cur].num_paths += sa[p].num_paths;
   tot_paths += sa[p].num_paths;
 }
 if (p == -1) {
   sa[cur].link = 0;
 } else {
   int q = sa[p].next[c];
   if (sa[p].len + 1 == sa[q].len) {
     sa[cur].link = q;
   } else {
     int clone = sz++:
     sa[clone].len = sa[p].len + 1;
     sa[clone].next = sa[q].next;
     sa[clone].num_paths = 0;
     sa[clone].link = sa[q].link;
     for (; p!= -1 && sa[p].next[c] == q; p =
         sa[p].link) {
```

```
sa[p].next[c] = clone;
sa[q].num_paths -= sa[p].num_paths;
sa[clone].num_paths += sa[p].num_paths;
}
sa[q].link = sa[cur].link = clone;
}
last = cur;
}
```

9.6 Z Algorithm

```
vector<int> compute_z(const string &s){
  int n = s.size();
  vector<int> z(n,0);
  int l,r;
  r = l = 0;
  for(int i = 1; i < n; ++i){
    if(i > r) {
        l = r = i;
        while(r < n and s[r - l] == s[r])r++;
        z[i] = r - l;r--;
  }else{
    int k = i-l;
    if(z[k] < r - i +1) z[i] = z[k];</pre>
```

```
else {
      1 = i:
       while (r < n \text{ and } s[r - 1] == s[r])r++;
       z[i] = r - 1;r--;
   }
 return z;
int main(){
 //string line;cin>>line;
 string line = "alfalfa";
 vector<int> z = compute_z(line);
 for(int i = 0: i < z.size(): ++i ){</pre>
   if(i)cout<<" ":</pre>
   cout<<z[i]:
 cout<<endl;
 // must print "0 0 0 4 0 0 1"
 return 0;
```