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[TUTORIAL 5]

Ans+1

BFS

- · uses queue DS
- · Stands for Breadth First Keasch
- Can be used to find Single source shortest path in an unweighted graph, I we reach a vertex with max. no. of edges from a source vertex.
- · Siblings are visited before the children.

Applications -

- · Shortest path & min spanning bee for unweighted graph.
- · Peer to peer networks
- · Social Networking sites

DFS

- · we stack Ds
- · Stands for Depth First Seasch.
- · We might boverse through more edger to reach a dutination vertex from a source

· Children and visited before the siblings.

Applications

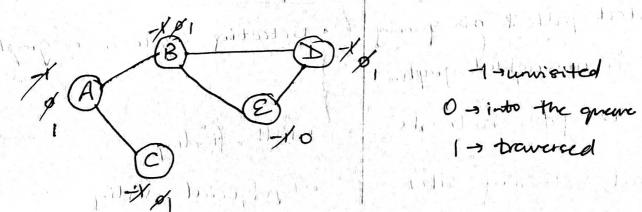
- · Detecting cycle in a graph.
- · Path finding.
- · Topolgocial sorting

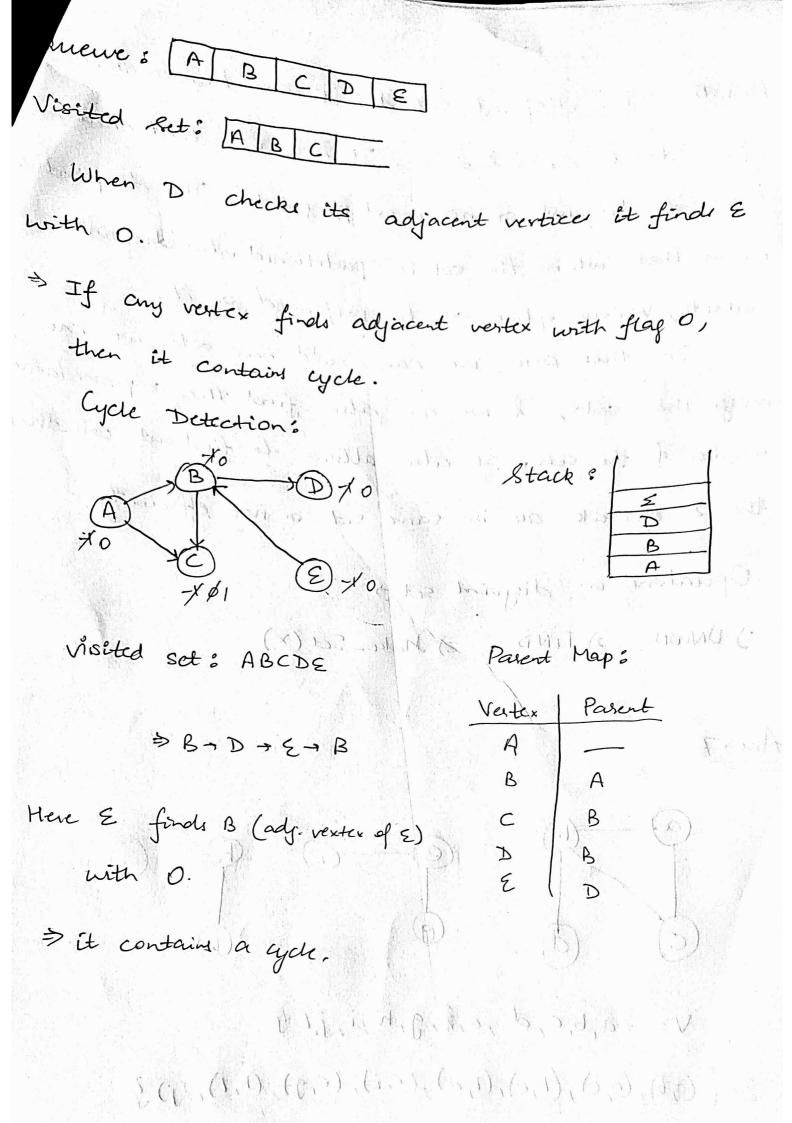
Ans 2 In BFS we we are Queue DS as queue is wed when things don't have to be processed immediately, but have to be processed in PIFO order lake BFS.

In DFS stack is used as DFS were backtrackyr. For DFS, we retrieve it from root to the fasthest mode as much as possible, this is the same idea as LIFO.

Ann = 3 Denne graph is a graph in which the no. of edges is close to the maximal no. of edges. Sparse graph is a graph in which the na of edges is close to the minimal no. of edges. It can be disconnected graph. Adjacency lists are preferred for sparse graph. I adjacency matrix for dense graph.

Ans-4 Cycle detection in Undirected graph (BFS)



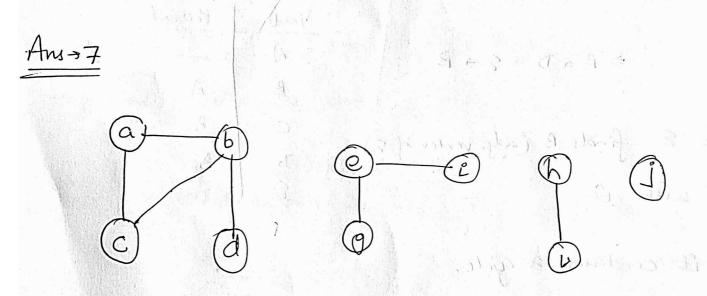


Ans The disjoint set ds. is also known as unid find ds & merge-find set. It is a ds that contains coll of disjoint or non-overlapping sets. The disjoint set means that when the set is partitioned into disjoint subsets, various option can be performed on it.

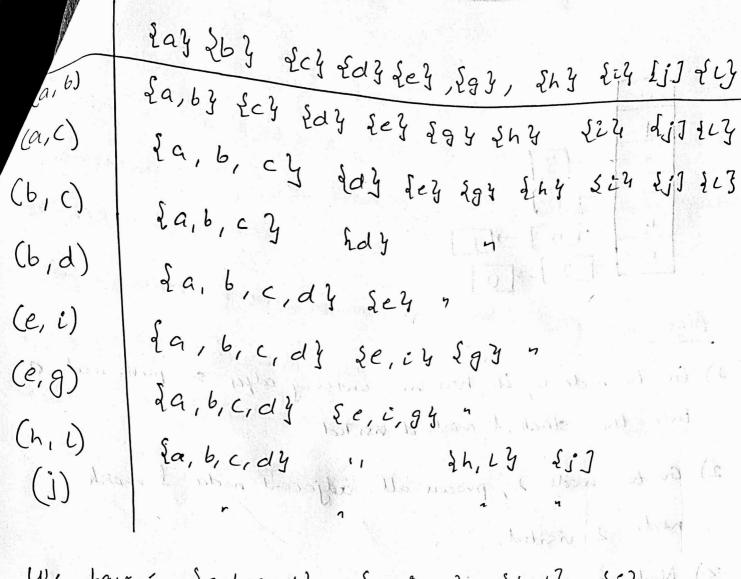
In this case, we can add new sets, we can meege the sets, I we can also find the representative member of the sets. It also allows to find out whether the 2 elements are in same set or not efficiently.

Operations on disjoint set

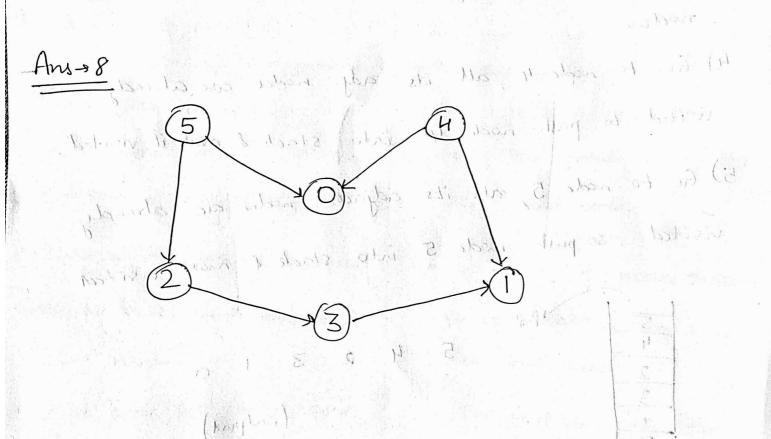
1) UNION 2) FIND 3) Make_Set(x)

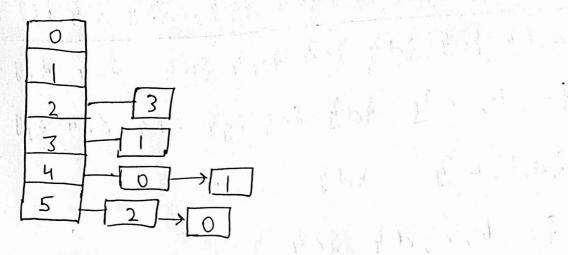


 $V = \{a,b,c,d,e,k,g,h,i,j,l\}$ $\mathcal{E} = \{(9,b),(a,c),(b,c),(b,d),(e,i),(e,g),(h,l),(j)\}$



We have: {a,b,c,d} {e,i,g} {h,l} 2j3





1) Go to node 0, it has no ontgoing edges so puth node o into the stack 4 mark it visited

Branch Branch F

- 2) Go to node 2, procen all adjacent nodes I mark node 2 visited.
- 3) Node 3 is already visited to continue with next
- 4) Go to node 4, all its adj. nodes are already visited to push node by into stack & mail it visited.
- 5) Go to node 5, all its adjacent nodes are already visited so puth mode 5 into stack & mark visited

15	1 Pop			·
4 2		5 (4 2	3 1 0
3				(output)

Heap is generally preferred for priority queue. implementation on heaps provide better performance compared to arrays or linkedlist.

Algos where priority quene is used ,

- 1) Dijkstra's shortet poth algo
- 2) Prim's algorithm -> to store keys of nodes & extract min key node at every step.

Ans -> 10

Min Heap

- · For every pair of the parent + descendent child node, the parent node always has lower vertue than des. child node.
- The value of nodes inc. as we traverse from root to leaf node.
- · Root node has Lowert value.

Max Heap

- · For every pair of pair & des. child node, the parent node has greater value than des. child node.
 - The value of nodes decreams as we traverse from root to leaf node.
 - . The root nock has the greatest value.