

MATH5425QUESTION 1 [10 marks]

(a) Let $U = V - A$ where $A \subseteq V$ is an independent set. Suppose that edge e of G is disjoint from U , $e = xy$. -1

[1] Then $\{x, y\} \subseteq A$ which contradicts the assumption that A is an indep set.

(b) Let A be a maximum indep set, $|A| = \alpha(G)$. Then

$$n = |A| + |V - A| = \alpha(G) + |V - A| \quad -1$$

$$\geq \alpha(G) + \tau(G) \quad -1$$

Since $V - G$ is a cover but might not be minimal.

(c) Assume G is triangle-free.

[2] (i) If v has degree $\Delta(G)$ then $N(v)$ is an indep set or else G has a triangle involving v & two adjacent neighbours of v . Hence A -1

$$|N(v)| = \Delta(G) \leq \alpha(G). \quad -1$$

(ii) Every edge is incident with at least one vertex of U , where U is a minimal cover. So $\alpha(G) \leq \tau(G)$ by (i)

[3] $|E| \leq \sum_{u \in U} d_G(u) \leq \Delta(G) \cdot |U| = \Delta(G) \tau(G).$ -1

(ii) Hence

$$|E| \leq \Delta(G) \tau(G) \leq \Delta(G) \tau(G) \text{ by (i).}$$

QUESTION 1

[1] (c) (iii) Hence, using fact $ab \leq \left(\frac{a+b}{2}\right)^2$,

[2] $|E| \leq \Delta(G) \tau(G)$

~~$\leq \left(\frac{\Delta(G) + \tau(G)}{2}\right)^2$ by "fact"~~

~~$\leq \Delta(G) \tau(G)$ by (ii)~~

$\leq \left(\frac{\Delta(G) + \tau(G)}{2}\right)^2$ by "fact" \downarrow

$= \left(\frac{n}{2}\right)^2 = \frac{n^2}{4}$ by (b). \downarrow

□.

QUESTION 2

(7)

[10 marks overall]

(a) (i) If G was disconnected then at least [2] one face of G has a disconnected boundary. But every face of G is

bounded by a cycle, which is connected. Hence G is connected.

(ii) If $e \in E(G)$ then e lies on boundary of [2] at least one face, by Lemma 4.2.2. Every face is bounded by a cycle, so e lies on a cycle. Hence, by Lemma 4.2.2 again, e ~~belongs~~ lies on the boundary of exactly 2 faces.

(iii) As G is connected, Euler's formula [3] applies. Write $m = |E(G)|$, $l = |F(G)|$ and note that by assumption and (ii),

$$2m = \frac{3l \cdot 4}{7} + \frac{4l \cdot 5}{7}$$

(double-counting the number of incident face-edge pairs).

$$\text{So } 2m = \frac{12+20}{7} l = \frac{32}{7} l$$

and $l = m \frac{2 \times 7}{32} = \frac{7}{16} m$. Hence $n - m + \frac{7}{16} m$ by Euler's formula, so $\frac{9}{16} m = n - 2$, $m = \frac{16}{9}(n-2)$.

4

Enough to find H_0

$S \subseteq V$ with

$q(G-S) \geq |S|$, so Tutte's condition fails.

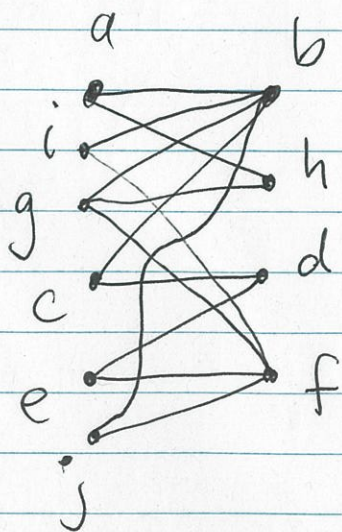
h a i j c d e

Then $q(G-S) = 4 > |S| = 2$ as all components of $G-S$ are odd.

五

1 - state or use Tutte's condition correctly.

OR:



1 - choose
1 - apply
1 - Hall's
(or state)

Let $S = \{i, j, g, a\}$

Then $N(s) = \{b, f, h\}$

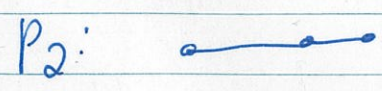
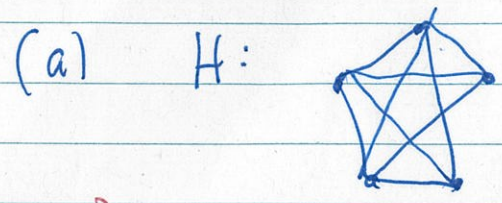
too small

Hall's Th^m fails.

Also
 $S = \{a, c, e, i, g\}$
 $N(S) = \{b, h, d, f\}$
 ie LHS \neq RHS

(BUT) that only says no matching of LITS!!
1 - Also explain why this implies no perf. matching

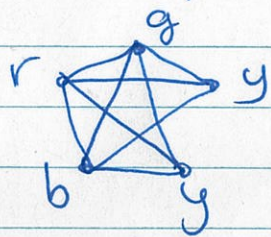
QUESTION 3. [10 marks] (6+4?)



[2]

(i) $R(H, P_2) \geq 7$. Bollobás?

$$(\chi(H)-1)(|P_2|-1) + u(H) = 3 \times 2 + 1 = 7$$

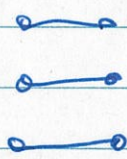


$$\chi(H)=4, u(H)=1.$$

or:

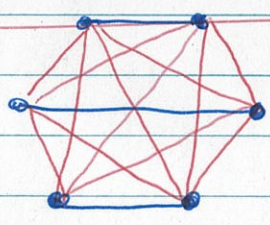
in K_6 ,

blue matching



no blue P_2 ,

-1



rest red: -1

any 5 vertices contain 2 blue edges

But H has 9 edges.

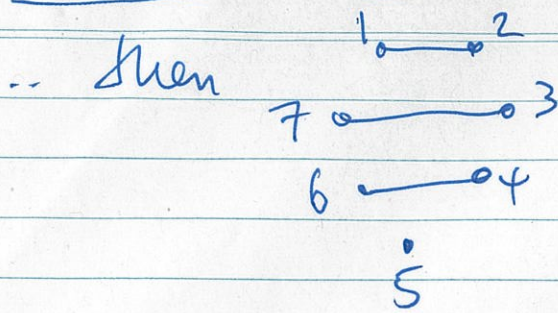
So, no red H .

$$\binom{5}{2} - 2 = 8 \text{ red edges.}$$

(ii) $R(H, P_2) \geq 7$. If 7 vertices, colour edges of K_n blue or red, try to avoid a red H or blue P_2 . Then blue edges form a matching. If 3 blue edges (max possible)

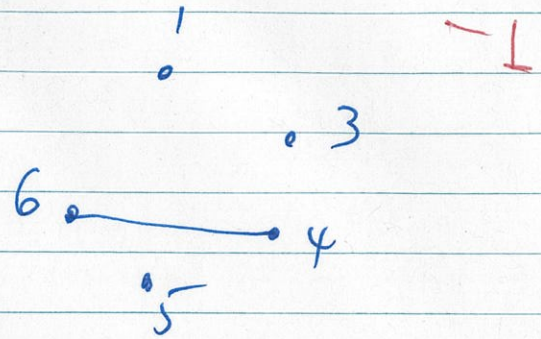
-1

(6)

QUESTION 3 (a), continued

Take vertices 1, 3, 4, 5, 6,

say:



This includes exactly one blue edge & so all the rest are red, giving a red copy of H .

Note, if less than 3 blue edges then we still get a red copy of H (even more easily).

[4]?(b) $H_1 \subseteq H_2$. True or false? $R(H_1, H_1) \subseteq R(H_2, H_2)$

TRUE. Proof let $n = R(H_2, H_2)$.

Then any red-blue colouring of the edges of K_n must contain a monochromatic

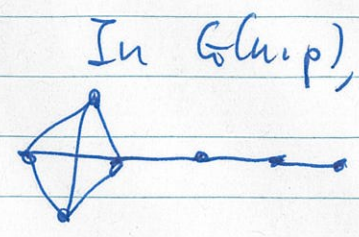
H_2 , by defⁿ of graph Ramsey numbers.

But $H_1 \subseteq H_2$, so we also have a monochromatic copy of H_1 , always.

Therefore $R(H_1, H_1) \leq n = R(H_2, H_2)$. \square

QUESTION 4 [10 marks]

$X = \# K_4$'s, $Y = \# \text{copies of } H$



(a) $\mathbb{E}X = \frac{\binom{n}{4}}{4!} p^6 = \frac{\binom{n}{4}}{24} p^6$ using result [2] └

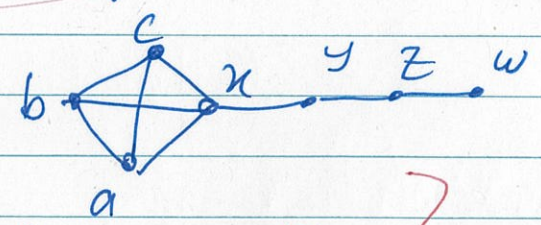
of Prob Sheet 8, as K_4 has 4 vertices, $\binom{4}{2} = 6$ edges

$\& 4! = 24$ automorphisms. └

~~or quote Prob Sheet~~

(b) $\mathbb{E}Y = \frac{1}{6} \binom{n}{7} p^9$ └

Well H has 7 vertices, 9 edges and $3! = 6$ automorphisms,



as a, b, c can be permuted in all ways and no other automorphisms are present. └

Hence by Prob Sheet 8 Q5 again,

$\mathbb{E}Y = \frac{1}{|Aut(H)|} \binom{n}{|H|} p^{|E(H)|} = \frac{1}{6} \binom{n}{7} p^9$ as stated.

(8)

(c) Put $p = n^{-\alpha}$, $\alpha > 0$. Want

(2) $\mathbb{E}X \approx \frac{1}{24} n^{4-6\alpha} \rightarrow 0$ so $4-6\alpha < 0$
 $\frac{2}{3} < \alpha$.

Want

$\mathbb{E}Y = \frac{1}{6} n^{7-9\alpha} \rightarrow \infty$ so $7-9\alpha > 0$,
 $\alpha < 7/9$.

eg $\frac{6}{9} = \frac{2}{3} < \alpha < 7/9$ ($\alpha = 3/4$ also works!)

so $\alpha = \frac{13}{18}$ would do.

Note $p(n) = n^{-13/18} \in (0,1)$ for $n \geq 2$.
 anything that works.

Take $p(1) = 1/2$, say. 1-reason.

Other conditions hold by choice of α .

(d)(i) By Markov's lemma, for $p(n) = n^{-13/18}$,

(2) $0 \leq P(X > 0) \leq \mathbb{E}X \rightarrow 0$ so

$P(X=0) = 1 - P(X > 0) \rightarrow 1$.

Hence $X=0$ a.s. -1

(ii) No this is false. If $Y > 0$ then $X > 0$
 as H has a subgraph which is a copy of K_4 .

So $0 \leq \Pr(Y > 0) \leq \Pr(X > 0) \rightarrow 0$

& hence a.s. $Y=0$ in fact. -1 reason