

(1)

MATHS425 2019.

Exam draft (d) G = (V, E) connected.

[5]

(a) $e \in E$, prove \exists spanning tree containing e.

Proof Since G is connected, it has a ~~max~~ spanning tree T . \neg If $e \in E(T)$ then we are done. \neg

Otherwise $T \cup e$ has a cycle C, since G is maximally acyclic. \neg let f $\in e$ be an edge of C. Then $(T \cup e) - f$ has $n-1$ edges, n vertices and is connected as f cannot be a bridge since it lies on a cycle. Hence $T' = (T \cup e) - f$ is a tree and $e \in E(T')$, as required. \neg

(b) Let T, T' distinct spanning trees in G.

[5] Set

$T_0 = T$ & suppose by induction that we have a sequence

$T = T_0, T_1, \dots, T_i$ with T_{j+1} obtained from T_j by deleting an edge of T_j & inserting an edge of T_{j+1} .

all distinct
spanning trees
of G.

If $T_i = T'$ then we are done.
So suppose that $T_i \neq T'$.

Consider the set difference $E(T_i - T') \cup E(T' - T_i)$.

Since $|E(T_i)| = |E(T')| = n-1$ both

$E(T_i - T')$ and $E(T' - T_i)$ \neg have $k \geq 1$ elements (same size).

Let $f \in E(T' - T_i)$. Then $T_i + f$

has a cycle, as T_i is a tree & so maximally acyclic. This cycle C must contain an edge $e \in E(T_i - T)$, \neg otherwise C is contained in T' .

(2)

contradicting the fact that T' is a tree.

Then $T_{i+1} = (T_i + f) - e$ is a tree,

since \cancel{e} is not a bridge in $T_i + f$ (e belongs to cycle C)
 so T_{i+1} is connected,
 n vertices, $n-1$ edges.

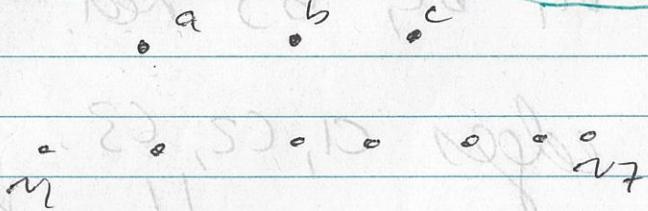
Also $|E(T_{i+1} - T')| = h-1 = |E(T' - T_{i+1})|$

so this process will terminate after $|E(T - T')|$ steps, a finite number. \square

Q2

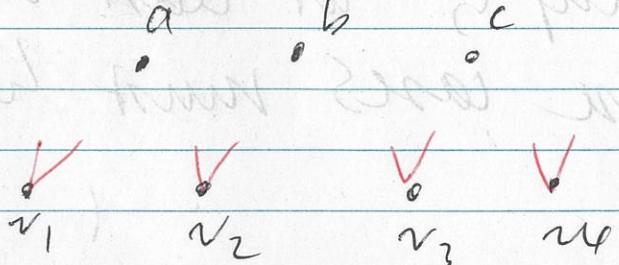
(a)

(5)



Call v_j red if ≥ 2 incident edges red,
 o'wise blue. ($d(v_j) = 3$). \square

At least 4 vertices $\{v_1, v_2\}$ with same colour, wlog red + v_3, v_4 are red.

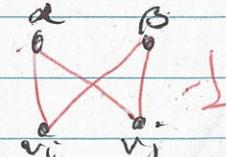


For each of v_1, v_4 , specify a pair in $\{a, b, c\}$ st v_j is incident

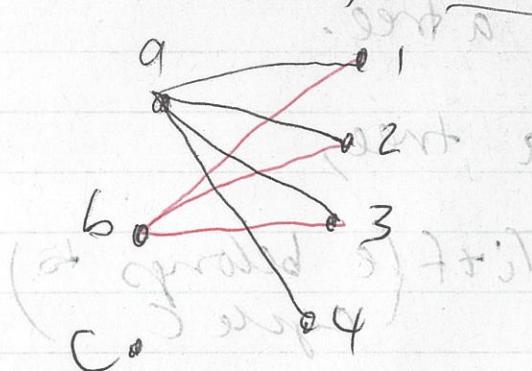
via a red edge, to (at least) these two vertices. If v_j incident with 3 red edges then wlog take $\{a, b\}$ as pair.

But only 3 options for this pair a, b, c

So at least two v_j with same pair, by pigeonhole principle. Hence



Q2(a) Andrew Kaploun



wlog a has ≥ 4 blue edges

WLOG a has ≥ 4 blue edges

Of these, $\leq 1, 2, 3, 4$ are red

If ≥ 2 of $1, 2, 3, 4$ are red

then we have

blue $K_{2,2}$. So wlog

b_1, b_2, b_3 red.

-1

Now consider edges c_1, c_2, c_3 .

If ≥ 2 blue then blue $K_{2,2}$ with b

If ≥ 2 red then red $K_{2,2}$ with b

+ By pigeonhole principle at least one of these cases must hold

missed

missed

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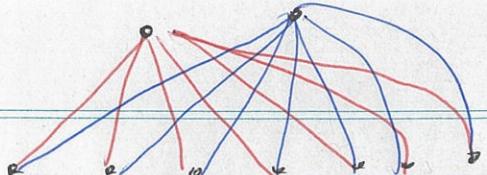
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$K_{2,7}$:

see facing for alternative argument for ③

(b)

[2]



red graph \cong blue graph

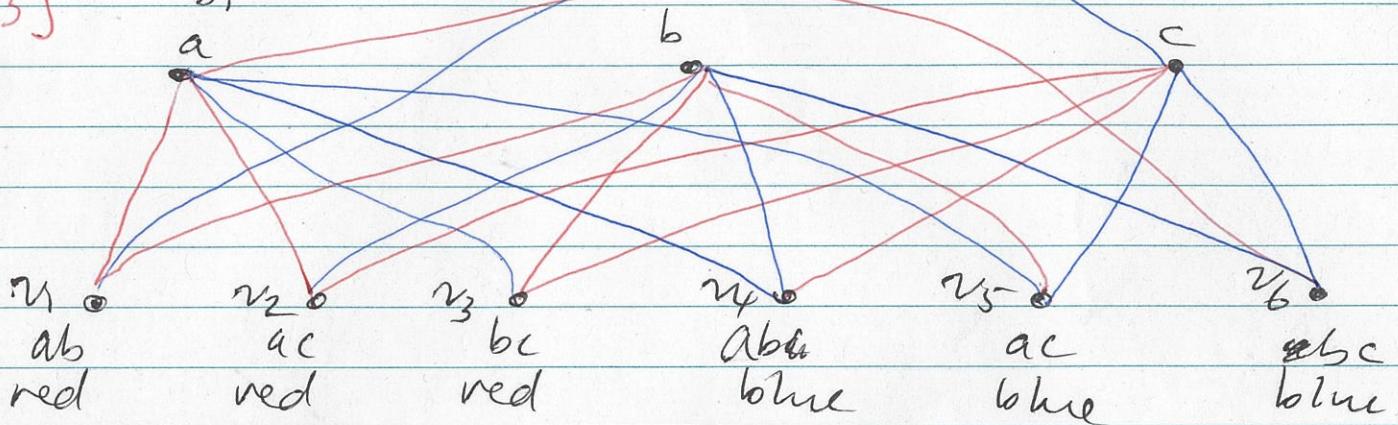
$\cong K_{1,7}$, a tree.

\Rightarrow no $K_{2,2}$.

(c)

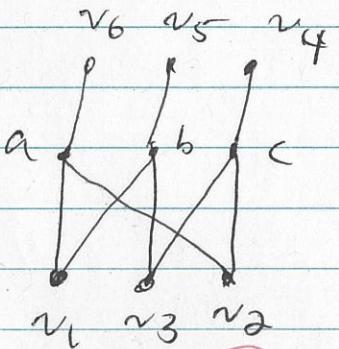
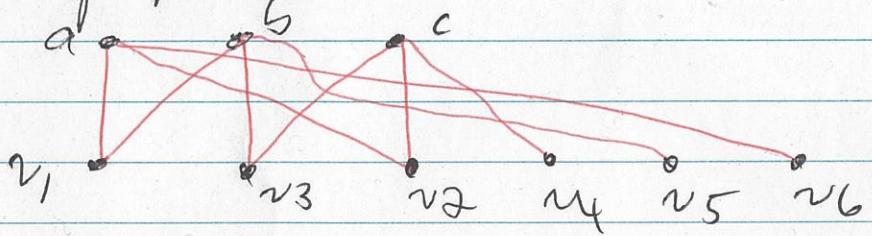
[3]

$K_{3,6}$:

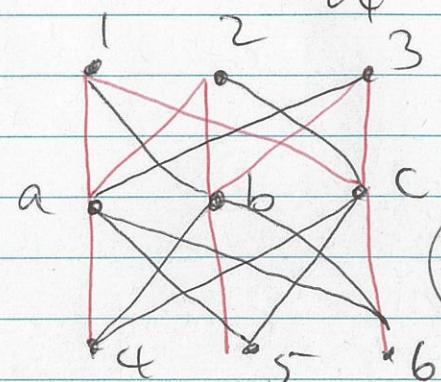
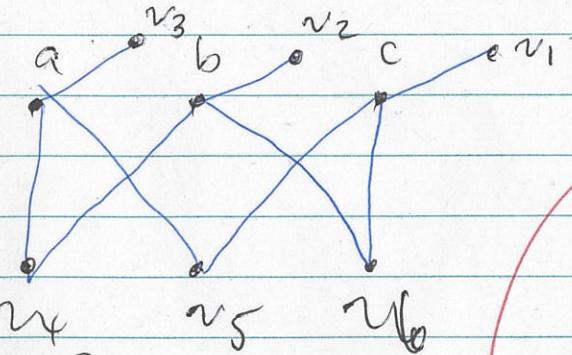


To avoid structures from (a), must be as above.

Red graph



Blue graph



Victoria
Guy

Both are a 6-cycle with 3 pendant edges,
no 4-cycle ($K_{2,2}$) in sight.

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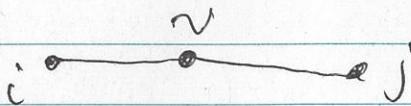
[Q3] $G(n, p)$, $X = X(n, p) = \# \text{ unordered pairs}$
 $\text{of distinct vertices}$
 $\text{with no corner nb.}$

(a) [4]

write $X = \sum_{i < j \leq n} X_{ij}$ where -1

$$X_{ij} = \begin{cases} 1 & \text{if } N(i) \cap N(j) = \emptyset, \\ 0 & \text{otherwise.} \end{cases}$$

Now



$$\Pr(v \in N(i) \cap N(j))$$

$$= p^2, \quad \text{using independence,}$$

for all $v \notin \{i, j\}$.

So, using independence again,

$$\mathbb{E} X_{ij} = \Pr(X_{ij} = 1) = (1-p^2)^{n-2} \quad \text{using independence again,}$$

and

Then by linearity
 of expectation,
as we have $n-2$ choices of v .

$$\mathbb{E} X = \binom{n}{2} (1-p^2)^{n-2} \quad \text{-1}$$

as there are $\binom{n}{2}$ choices for i, j .

(b) Hence, since $1-p \leq e^{-p}$ and $\binom{n}{2} \leq n^2$, -1

[2]

$\mathbb{E} X \leq n^2 e^{-p(n-2)} = \exp(2 \frac{\ln n}{n} - p^2(n-2))$,
 as required.

(c) well $2 \ln n \leq p^2(n-2)$ iff

[2]

$$\frac{2 \ln n}{n-2} \leq p^2, \quad \text{so } p \geq \sqrt{\frac{2 \ln n}{n-2}}.$$

We want strict inequality so take eg

$$p = p(n) = \sqrt{\frac{3 \ln n}{n-2}} \quad -\perp$$

Note for
 $p \in (0, 1)$ for
 a suff (large)
 $n \geq 5$ or so
 (no TE: MARK
 $p \in (0, 1)$)

Then $\mathbb{E}X \leq \exp\left(2 \ln n - \frac{3 \ln n (n-2)}{n-2}\right)$
 with this function P,
~~as~~
 $= \exp(2 \ln n - 3 \ln n) = \frac{1}{n} \rightarrow 0 \text{ as } n \rightarrow \infty$

(d) In $G(n, p)$, with this p , $\Pr(X > 0) \leq \mathbb{E}X < 0$.
 [2]

Hence aas ~~$X \neq 0$~~ , $X=0$, so aas:

given any $i < j$, $i \neq j$ have a common nb
 in $G \in G(n, p(n))$.

Then $d(i, j) \leq 2$ so $\text{diam}(G(n, p)) \leq 2$. \perp

(ii) But $\text{diam} = 1$ iff $G(n, p) \cong K_n$, an event
 with probability

~~P~~ $P^{\binom{n}{2}} \rightarrow 0$ \perp
 really fast!
 (since $p \rightarrow 0$ alone)

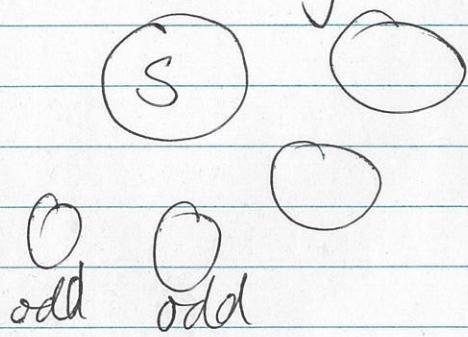
So aas $X=0$ and $G(n, p) \not\cong K_n$
 that is,
 aas $\text{diam}(G) = 2$.

(Q4) Parts (a) + (b) are unrelated.

(a) $G = (V, E)$ in vertex, M a max match.

(5) Fix $S \subseteq V$ + let $M' = \text{edges in } M$ which are disjoint from S .

(i) M' can be thought of as a matching in $G-S$. -1



The maximum number of vertices which can be matched by M' is

$$|G-S| - q(G-S) \quad \text{since there must be at least one unmatched vertex in every odd component of } G. \quad \text{---}$$

But $|G-S| = n - |S|$ so this gives the desired expression.

(ii) The maximum number of edges in $M-M'$ is $|S|$, matching all elements of S to a vertex in $G-S$.

Otherwise, ~~Each edge in $M-M'$ matches at least one vertex of S , so $|M-M'| \leq |S|$.~~

if any edge of $M-M'$ is completely within S then it uses two vertices of S : -1

if number of such edges is i then

$$|M-M'| \leq i + (|S|-i) = |S|-i, \quad \text{maximised when } i=0. \quad \text{---}$$

so

$$|M| \geq |M'| + |M-M'| \leq \frac{1}{2} (n - |S| - q(G-S)) + |S|$$

as claimed. □

$$= \frac{1}{2} (n + |S| - q(G-S)) \quad \text{---}$$

(b) Let $c: E(K_n) \rightarrow \{1, 2, 3\}$ with $c(uv)$

[5] assigned var, & independently
for each edge $uv \in K_n$, -1

Let $X = \# \text{ subsets } U \subseteq C_n \text{ of size } k \text{ st } K_n[U] \text{ is monochromatic.}$

Write $X = \sum_{\substack{U \subseteq [n] \\ |U|=k}} X_U$ where X_U is the
indicator variable
for the event -1
" $K_n[U]$ is monochromatic".

$$\text{Then } \mathbb{E}X_U = \Pr(X_U=1) = 3 \cdot \frac{1}{3^{\binom{k}{2}}} = 3^{1-\binom{k}{2}} \quad -1$$

Since there are 3 choices of colour
and then all $\binom{k}{2}$ edges must choose
this colour. This uses independence.

Hence, by linearity of expectation,

$$\mathbb{E}X = \binom{n}{k} \cdot 3^{1-\binom{k}{2}} \quad -1 \text{ as there are } \binom{n}{k} \text{ choices for } U.$$

By assumption, this gives

$$\mathbb{P}(X>0) \leq \mathbb{E}X = \binom{n}{k} 3^{1-\binom{k}{2}} < 1. \quad \text{why}$$

So there is a particular colouring $c: E(K_n) \rightarrow \{1, 2, 3\}$
of the edges of K_n with no
monochromatic K_k . \square

Or:

Q(a)(i)-2

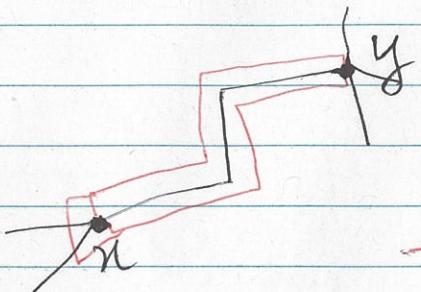
(ii)-62

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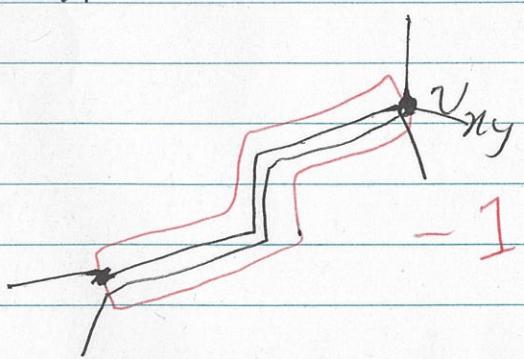
(b) -6

-1

[3]

- Q5 (a) If G is planar and e an edge of G
 Then take a plane embedding $\overset{-1}{\rightarrow}$ of G : (which we also call G)

 Take a thin "ribbon" of width $\epsilon > 0$ around the edge
 from x to y , and around x with ϵ so small that
 this ribbon is disjoint from the rest of G .

Then we can embed G/e in the plane by, say, replacing $\cdot y$ by $\cdot \overset{my}{y}$
 and rerouting all edges incident with x (which are not already covered by $y_w \mapsto \overset{my}{y}_w$, eg do it better if common w)
 to stay within this ϵ -ribbon:

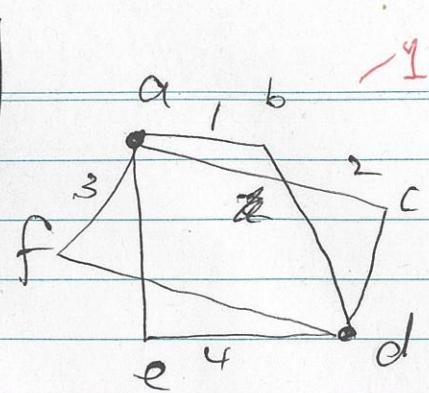
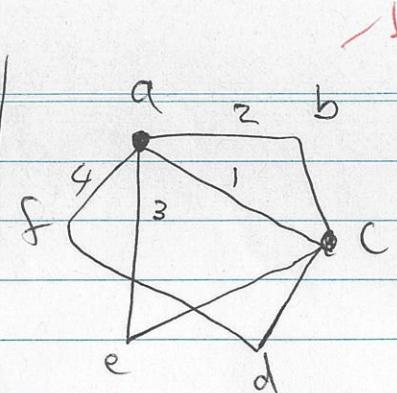
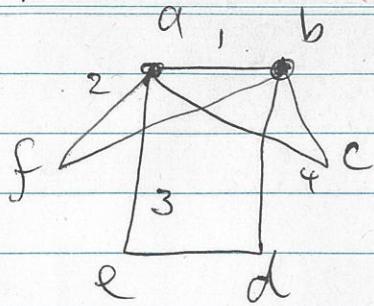


This is a plane embedding of G/e .
 So G/e is planar.

- (i) [5] Rtp for all $u, v \in V(G)$, there are $\overset{-1}{\rightarrow}$ at least 4 indep paths from u to v .
 Then Global Mindeg $\Rightarrow G$ is 4-conn.

By symm, suff to check just for ab, ac, ad (rotate G to ^{cover} all other edges)
 \Rightarrow G is 4-conn.

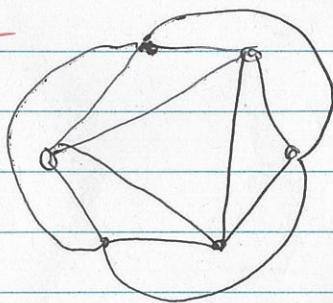
Then



as required.

(ii) Note, G_i is planar:

[2] So if such a sequence



$$G = G_0, G_1, \dots, G_r = H$$

existed, with $G_0 = K_5$, then (a) implies that all G_i are planar, hence $G_0 = K_5$ is planar.

This is a contradiction.

(or): $|H| = 6$ so sequence would have $G_1 = H$.

Contract edge ab, for example, to give graph where c + f are still not adjacent. (an anti-matching)

This always happens: every vertex has a nonneighbor. Of the 4 non-contracted vertices in G_0 , only 2 will have two non-~~edges~~ neighbors contracted so still two nonneighbors. So $H \not\cong K_5$