

Vardaan

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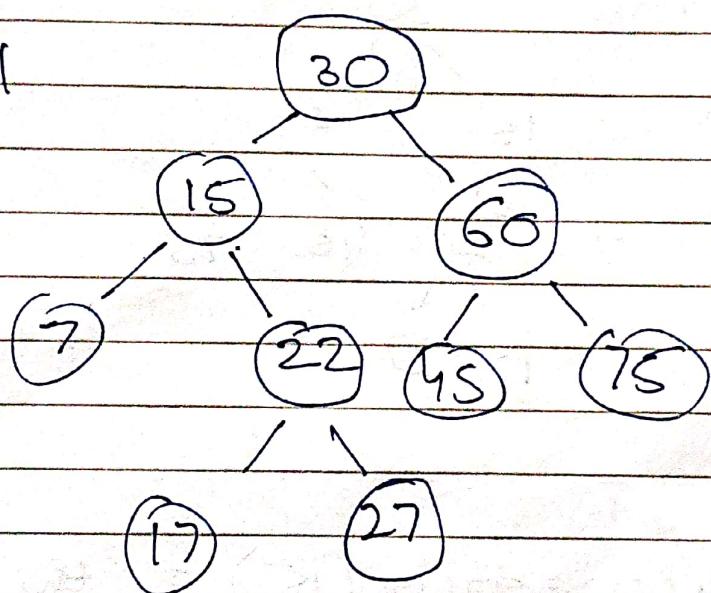
①

### Assignment 3A

i)

30, 15, 60, 7, 22, 45, 75, 17, 27

1.1



1.2

in order (L, Root, R)

→ 7, 15, 17, 22, 27,  
30, 45, 60, 75

Pre order → 30, 15, 7, 22, 17,  
(Root, L, R) 27, 60, 45, 75

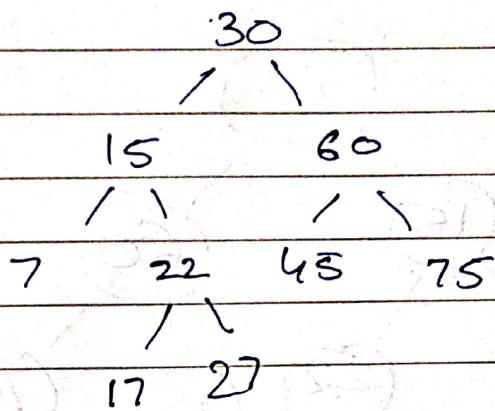
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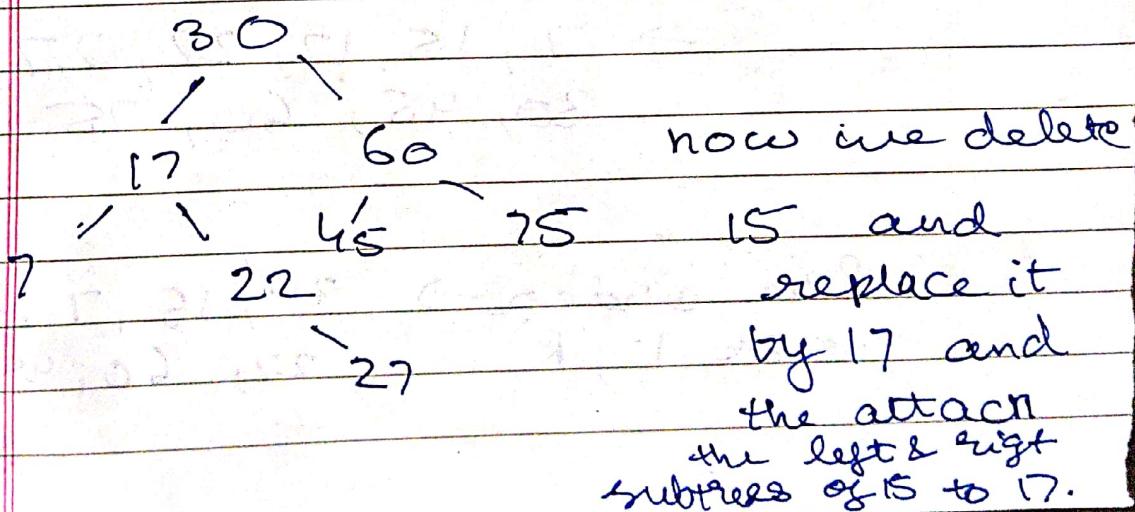
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Post order  $\rightarrow$  7, 17, 27, 22, 15,  
 (L, R, Root) 45, 75, 60, 30

1. 3.

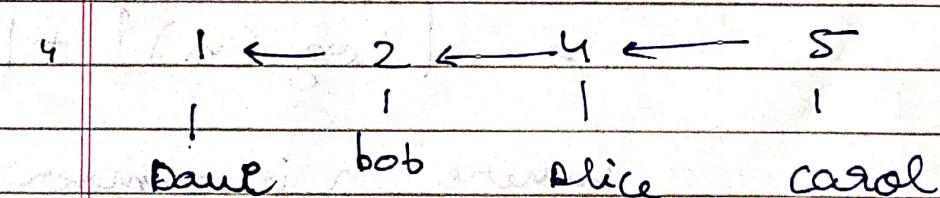
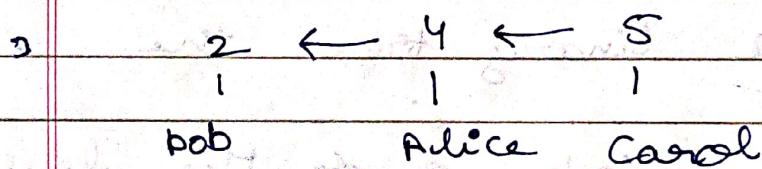
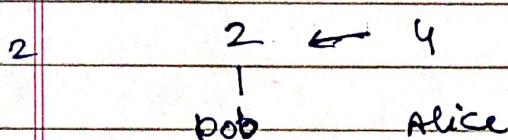
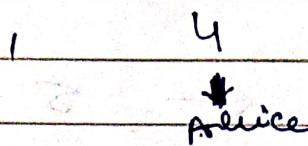


Successor of (15) = the leftmost  
 in the  
 right subtree  
 of 15

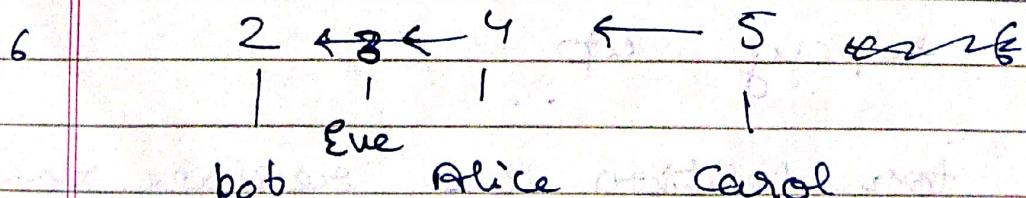
 $\Rightarrow 17$ 

(2)

(2)



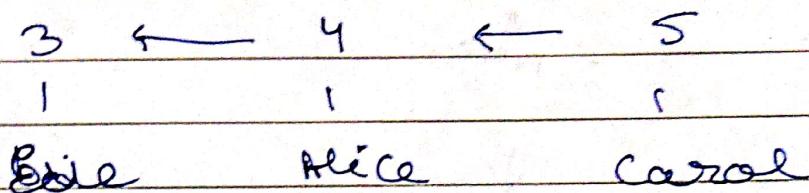
5      1 Dane



7      2 - bob

(1)

2.1



2.2.

In a binary tree the  
max height of the leaf  
from the root

$$= \lceil \log(n) \rceil + 1$$

where  $n$  is number of  
elements

for insertion we insert  
at the bottom and then  
heapify up.

for deletion we replace the  
node to be deleted by the leaf  
and then heapify down  
by this if we delete the root

(5)

Date \_\_\_\_\_  
 Page No. \_\_\_\_\_

then the max time it will take is  $O(\log n)$

Q3. 3.1  $\rightarrow$  code submitted  
separately

- 3.2 .

in usual BST

Right > root > left

but now

Right < root < left.

3.3.

in this approach  
the time complexity  
 $O(n)$

& Space complexity  
 $O(\log n)$   $O(n)$

Space complexity is  $O(h)$

where  $h$  is the height of the tree which is  $\log n \leq h \leq n$



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(6)

Date / /  
Page No.

Q4.

4.2 → In this approach

our time complexity =  $n \log K$  $K \rightarrow$  window size \* the operations $n \rightarrow$  number of elements

but if we use our better approach

↓

 $P_1 = \underset{i=0}{\text{Null}}, P_2 = \text{Null}, P_3 = \text{Null}$ for ( $i < n$ ,  $i++$ ). { $P_1 = i, P_2 = i - 1, P_3 = i - 2$ if ( $P_1, P_2, P_3 \geq 0$ ) {return  $\max(\text{dis}[P_1], \text{lis}[P_2], \text{lis}[P_3])$ 

{

close function

in this approach the time complexity =  $O(n)$