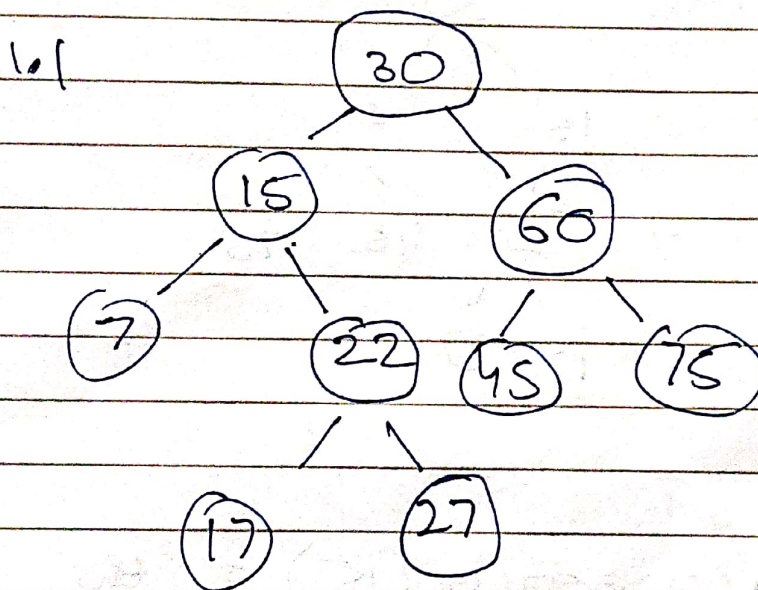


①

## Assignment 3A

i) 30, 15, 60, 7, 22, 45, 75, 17, 27



1.2

in order (L, Root, R)

→ 7, 15, 17, 22, 27,  
30, 45, 60, 75

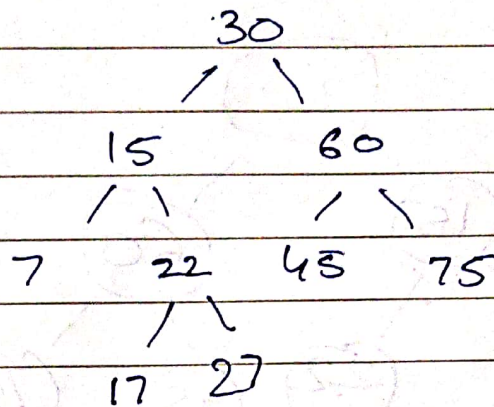
Pre order → 30, 15, 7, 22, 17,  
(root, L, R) 27, 60, 45, 75



②

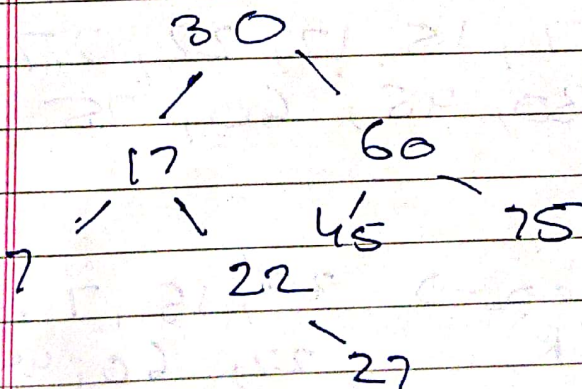
Post order  $\rightarrow$  7, 17, 27, 22, 15,  
(L, R, Root) 45, 75, 60, 30

1.3-



Successor of (15) = the leftmost  
in the  
right subtree  
of 15

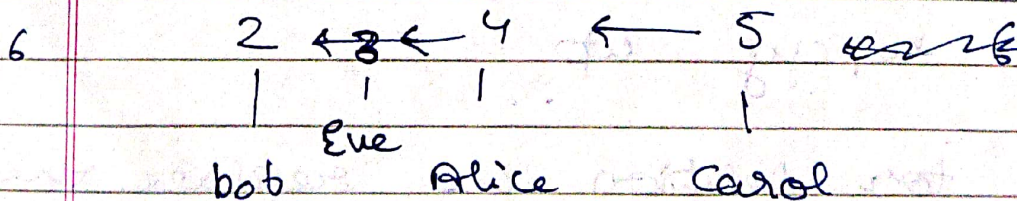
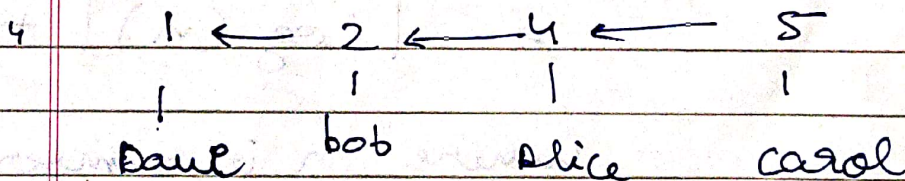
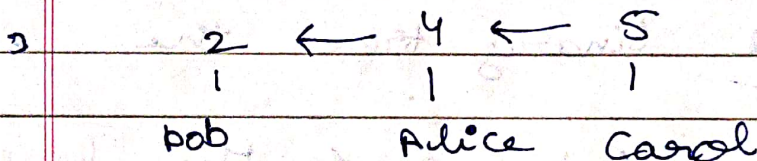
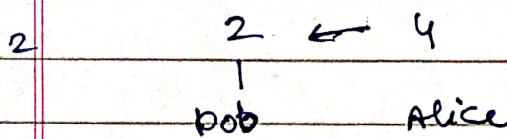
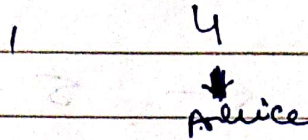
$\Rightarrow$  17



now we delete  
15 and  
replace it  
by 17 and  
the attach  
the left & right  
subtrees of 15 to 17.

3

2)



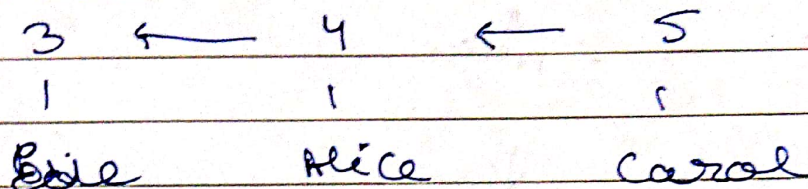


4

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2.1



2.2.

In a binary tree the

max height of the leaf  
from the root

$$= \lfloor \log(n) \rfloor + 1$$

where  $n$  is number of  
elements

for insertion we insert  
at the bottom and then  
heapify up.

for deletion we replace the  
node to be deleted by the leaf  
and then heapify down  
by this if we delete the root



5

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then the max time it will take is  $O(\log cn)$

Q3. 3.1  $\rightarrow$  code submitted separately

- 3.2 .

in usual BST

Right > root > left

but now

Right < root < left.

3.3.

in this approach  
the time complexity  
 $O(n)$

& space complexity  
 $O(\log n)$   $O(n)$

Space complexity is  $O(h)$

where  $h$  is the height of the tree which is  $\log n \leq h \leq n$



6

Q4.

4.2  $\rightarrow$  In this approach

our time complexity =  $n \log K$

$K \rightarrow$  window ~~of the operations~~

$n \rightarrow$  number of elements

~~but if we use an~~  
better approach

$\downarrow$

$P_1 = 0, P_2 = \text{Null}, P_3 = \text{Null}$   
 $i = 0$

for ( $i < n, i++$ ) {  
 $P_1 = i, P_2 = i-1, P_3 = i-2$   
if ( $P_1, P_2, P_3 > 0$ ) {

~~return~~  $\max(\text{lis}[P_1], \text{lis}[P_2], \text{lis}[P_3])$   
}

close continue

in this approach the time  
complexity =  $O(n)$