.var flag 2 ⇒ array of 2 integers

.var turn →

.var count →

.main

lea flag, %fx ⇒ flag → fx (reg)
addr

mov %bx, %cx ⇒ cx = bx ↗ self

neg %cx ⇒ cx = - self

add $1, %cx cx = 1+(-self)

.acquire

mov $1, 0(%fx, %bx, 4) → flag[self] = 1

disp → 0

base → %fx mem addr = %fx + (%bx x 4)

index → %bx

scale → 4

mov %cx, turn → turn = 1 - self

.spin

mov 0(%fx, %cx, 4), %ax → ax = flag[1-self]

test $1, %ax

jne .fini , if flag[1-self] != 1, skip past loop to .fini

meaning go out of spin, exec critical section

.spin 2

mov turn, %ax

test %cx, %ax → compare turn and 1-self

je .spin1 → if turn == 1-self, back to spinning

.fini

mov count, %ax ax = count

add $1, %ax ax = 1+ax

mov %ax, count count = ax

} critical section

.release

mov $0, 0(%fx, %bx, 4) flag[self] = 0

mov %cx, turn turn = 1-self

halt

Let's assume thread 0 acquires lock 1st,

meaning self = 0

Dry run (thread 0)

cx = 1, bx = 0, flag[0] = 1, turn = 1

ax = flag[1], 0 != 1 jne .fini

ax = 1 (count)

flag[0] = 0

turn = 1

Dry run (thread 1)

bx = 1 , cx = 0

flag[1] = 1 , turn = 0

ax = flag[0]

1 == flag[0] → this is 1. since thread 0 is running with lock acquired

Now thread 1 was spinning when thread 0 releases the lock

Above now goes to critical section because the condition, 1 != 0 → flag[0], lock released

Or it was executing .spin 2 section, now

01 = 1 → turn, updated by thread 0.