

C++ Basics

Arithmetic

Arithmetic

$y = a + b * c - d \rightarrow$ **expression**

$a + b \rightarrow$ **addition operation**

$+ \rightarrow$ **addition operator**

$a, b \rightarrow$ **operands**

$a++ \rightarrow$ **postfix increment operation**

$++ \rightarrow$ **postfix operator**

$a \rightarrow$ **operand**

Rules of Operator Precedence

C++ applies the operators in arithmetic expressions in a precise order determined by the following **rules of operator precedence**:

1. Operators in expressions contained within pairs of parentheses are applied first. Parentheses are said to be at the “*highest level of precedence*.”

In cases of nested, or embedded, parentheses, such as

(a * (b + c))

the operators in the ***innermost*** pair of parentheses (b+c) *are applied first*.

Rules of Operator Precedence

C++ applies the operators in arithmetic expressions in a precise order determined by the following **rules of operator precedence**:

2. *Multiplication(*), division(/) and remainder(%)* operations are evaluated next. If an expression contains several multiplication, division and remainder operations, operators are applied from *left to right*.

Rules of Operator Precedence

C++ applies the operators in arithmetic expressions in a precise order determined by the following **rules of operator precedence**:

3. *Addition(+)* and *subtraction(-)* operations are applied last. If an expression contains several addition and subtraction operations, operators are applied from *left to right*. Addition and subtraction also have the *same level* of precedence.

So, the rules of operator precedence define the order in which C++ applies operators.

Rules of Operator Precedence

$y = a * x * x + b * x + c;$



Step 1. $y = 2 * 5 * 5 + 3 * 5 + 7;$ (Leftmost multiplication)

$2 * 5$ is 10

Step 2. $y = 10 * 5 + 3 * 5 + 7;$ (Leftmost multiplication)

$10 * 5$ is 50

Step 3. $y = 50 + 3 * 5 + 7;$ (Multiplication before addition)

$3 * 5$ is 15

Step 4. $y = 50 + 15 + 7;$ (Leftmost addition)

$50 + 15$ is 65

Step 5. $y = 65 + 7;$ (Last addition)

$65 + 7$ is 72

Step 6. $y = 72$ (Low-precedence assignment—place 72 in y)

Operator precedence and associativity

| Operator | Type | Associativity |
|------------------------------------------------------------------------------------------------------------------------------------|----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|---------------|
| :: :: () | binary scope resolution unary scope resolution grouping parentheses <i>[See caution in Fig. 2.10 regarding grouping parentheses.]</i> | left to right |
| () [] . -> ++ -- typeid dynamic_cast<type> static_cast<type> reinterpret_cast<type> const_cast<type> | function call array subscript member selection via object member selection via pointer unary postfix increment unary postfix decrement runtime type information runtime type-checked cast compile-time type-checked cast cast for nonstandard conversions cast away const-ness | left to right |
| ++ -- + - ! ~ sizeof & * new new[] delete delete[] | unary prefix increment unary prefix decrement unary plus unary minus unary logical negation unary bitwise complement determine size in bytes address dereference dynamic memory allocation dynamic array allocation dynamic memory deallocation dynamic array deallocation | right to left |
| (type) | C-style unary cast | right to left |

Operator precedence and associativity

| Operator | Type | Associativity |
|-----------------------------------------------------------------|------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|---------------|
| . [*] -> [*] | pointer to member via object pointer to member via pointer | left to right |
| * / % | multiplication division remainder | left to right |
| + - | addition subtraction | left to right |
| << >> | bitwise left shift bitwise right shift | left to right |
| < <= > >= | relational less than relational less than or equal to relational greater than relational greater than or equal to | left to right |
| == != | relational is equal to relational is not equal to | left to right |
| & | bitwise AND | left to right |
| ^ | bitwise exclusive OR | left to right |
| | bitwise inclusive OR | left to right |
| && | logical AND | left to right |
| | logical OR | left to right |
| ?: | ternary conditional | right to left |
| = += -= *= /= %= &= ^= = <<= >>= | assignment addition assignment subtraction assignment multiplication assignment division assignment remainder assignment bitwise AND assignment bitwise exclusive OR assignment bitwise inclusive OR assignment bitwise left-shift assignment bitwise right-shift assignment | right to left |
| , | comma | left to right |

Associativity

When we say that certain operators are applied from *left to right* or *right to left*, we are referring to the **associativity** of the operators.

For example, the addition operators (+) in the expression

$$a + b + c + d$$

associate from *left to right* and is parsed as

$$((a + b) + c) + d$$

Associativity

The associativity of **assignment(=)** operator is *right to left*:

$$a = b = c = d$$

is parsed as

$$a = (b = (c = d))$$

not

$$((a = b) = c) = d$$

Order of evaluation

- Order of evaluation of any part of any expression, including order of evaluation of ***function arguments*** is **unspecified** (with some exceptions).
 - The compiler can evaluate operands and other subexpressions in any order and may choose another order when the same expression is evaluated again.
- There is **NO** concept of *left-to-right* or *right-to-left* evaluation in C++. This is not to be confused with left-to-right and right-to-left associativity of operators.

The expression

$$a() + b() + c()$$

is parsed as

$$(a() + b()) + c()$$

due to left-to-right associativity of operator+, **but c() may be evaluated first, last, or between a() or b() at run time:**

Order of evaluation

```
#include <cstdio>

int a() { return std::puts("a"); }
int b() { return std::puts("b"); }
int c() { return std::puts("c"); }

void z(int, int, int) {}

int main()
{
    z(a(), b(), c()); // all 6 permutations of output are allowed
    return a() + b() + c(); // all 6 permutations of output are allowed
}
```

Possible output:

```
b
c
a
c
a
b
```

Please refer to this for more details:

https://en.cppreference.com/w/cpp/language/eval_order

Order of evaluation

- Order of evaluation of any part of any expression, including order of evaluation of *function arguments* is **unspecified** (with some exceptions).

Example of some exceptions:

- **Logical AND operator (&&)** guarantees *left-to-right* evaluation: the second operand is not evaluated if the first operand is **false**.

a() && b() // b() will be evaluated if a() is true

- **Logical OR operator (||)** guarantees *left-to-right* evaluation; moreover, the second operand is not evaluated if the first operand evaluates to **true**

a() || b() // b() will NOT be evaluated if a() is true

Examples

```
int x[2] = {0,10};  
int* xPtr = x;
```

- `cout << ++*xPtr;`
- `cout << *++xPtr;`
- `cout << *xPtr++;`
- `cout << (x[0] ? 5 : x[1] ? 555 : 10);`
- `cout << x[0] ? 10 : 1000;`

Examples

```
int x[2] = {0,10};  
int* xPtr = x;
```

- `cout << ++*xPtr; // ++(*xPtr), 1`
- `cout << *++xPtr;`
- `cout << *xPtr++;`
- `cout << (x[0] ? 5 : x[1] ? 555 : 10);`
- `cout << x[0] ? 10 : 1000;`

Examples

```
int x[2] = {0,10};  
int* xPtr = x;
```

- `cout << ++*xPtr; // ++(*xPtr), 1`
- `cout << *++xPtr; // *(++xPtr), 10`
- `cout << *xPtr++;`
- `cout << (x[0] ? 5 : x[1] ? 555 : 10);`
- `cout << x[0] ? 10 : 1000;`

Examples

```
int x[2] = {0,10};  
int* xPtr = x;
```

- `cout << ++*xPtr; // ++(*xPtr), 1`
- `cout << *++xPtr; // *(++xPtr), 10`
- `cout << *xPtr++; // *(xPtr++), 0`
- `cout << (x[0] ? 5 : x[1] ? 555 : 10);`
- `cout << x[0] ? 10 : 1000;`

Examples

```
int x[2] = {0,10};  
int* xPtr = x;
```

- `cout << ++*xPtr; // ++(*xPtr), 1`
- `cout << *++xPtr; // *(++xPtr), 10`
- `cout << *xPtr++; // *(xPtr++), 0`
- `cout << (x[0] ? 5 : (x[1] ? 555 : 10)); // 555`
- `cout << x[0] ? 10 : 1000;`

Examples

```
int x[2] = {0,10};  
int* xPtr = x;
```

- `cout << ++*xPtr; // ++(*xPtr), 1`
- `cout << *++xPtr; // *(++xPtr), 10`
- `cout << *xPtr++; // *(xPtr++), 0`
- `cout << (x[0] ? 5 : (x[1] ? 555 : 10)); // 555`
- `(cout << x[0]) ? 10 : 1000; // 0`

C++ Basics

If...else statement

How this will be parsed by compiler?

```
int x = 10, y = 4;  
if (x > 5)  
    if (y > 5)  
        cout << "x and y are > 5";  
else  
    cout << "x is <= 5";
```

If...else statement

How this will be parsed by compiler?

```
int x = 10, y = 4;  
if (x > 5)  
    if (y > 5)  
        cout << "x and y are > 5";  
else  
    cout << "x is <= 5";
```

Result: x is <= 5

This is called **dangling-else** problem

If...else statement

```
int x = 10, y = 4;  
if (x > 5) {  
    if (y > 5) {  
        cout << "x and y are > 5";  
    } else {  
        cout << "x is <= 5";  
    }  
}
```

- C++ compilers *always* associate an **else** with the immediately preceding **if** unless told to do otherwise by the placement of braces ({ and }).

If...else statement

- Always use braces {} in control statements, even for single-statement bodies

```
if (condition) {  
    single-statement or multi-statement body  
}
```

- **Indent** both body statements (or groups of statements) of an if...else statement.

```
if (a > 0) {  
    cout << "a > 0";  
} else if (a < 0) {  
    cout << "a < 0";  
} else {  
    if (b > 0) {  
        cout << "b > 0";  
        if (c > 0) {  
            cout << "c > 0";  
        }  
    }  
}
```


for statement

```
for (initialization; LoopContinuationCondition; increment)  
{  
    statement  
}
```

1. *Initialization*
2. *LoopContinuationCondition*
3. *Statement*
4. *Increment*

while statement

```
initialization;  
while (LoopContinuationCondition) {  
    statement  
    increment;  
}
```

1. *Initialization*
2. *LoopContinuationCondition*
3. *Statement*
4. *Increment*

for & while statements

Can you convert every **for** statement to **while**?

```
int n = 10;
for (int i = 0; i < n; ++i) {
    if (5 == i) {
        cout << "5";
    }
}
```

for & while statements

Can you convert every **for** statement to **while**?

```
int n = 10;
for (int i = 0; i < n; ++i) {
    if (5 == i) {
        cout << "5";
    }
}
```

```
int n = 10;
{
    int i = 0;
    while (i < n) {
        if (5 == i) {
            cout << "5";
        }
        ++i;
    }
}
```

for & while statements

```
int n = 10;
for (int i = 0; i < n; ++i) {
    if (5 == i) {
        continue;
    }
    ...
}
```

```
int n = 10;
{
    int i = 0;
    while (i < n) {
        if (5 == i) {
            continue;
        }
        ...
        ++i;
    }
}
```

for & while statements

```
int n = 10;
for (int i = 0; i < n; ++i) {
    if (5 == i) {
        continue;
    }
    ...
}
```

```
int n = 10;
{
    int i = 0;
    while (i < n) {
        if (5 == i) {
            ++i;
            continue;
        }
        ...
        ++i;
    }
}
```

Confusing the == and = operators

What is the difference between these two statements?

```
if (payCode == 4) { // good
    cout << "You get a bonus!" << endl;
}
```

```
if (payCode = 4) { // bad
    cout << "You get a bonus!" << endl;
}
```

if (paycode = 4) will be parsed as

1. paycode =4
2. If (paycode)

How to avoid from this kind of errors?

lvalues and rvalues

- Variable names are said to be **lvalues** (for “*left values*”) because they can be used on an *assignment operator’s left side*.
- **Literals** are said to be **rvalues** (for “*right values*”) because they can be used on **only** an *assignment operator’s right side*.

`x = 25; // x -> lvalue, 25 -> rvalue`

- **Lvalues** can also be used as **rvalues** on the *right side* of an assignment, **but not vice versa**:

`y = x; // OK, y -> lvalue, x -> lvalue (used as rvalue)`
`25 = x; // ERROR!`

Confusing the == and = operators

- ✓ Programmers normally write conditions such as `x == 7` with the variable name (an lvalue) on the left and the literal (an rvalue) on the right. Placing the literal on the left, as in `7 == x`, enables the compiler to issue an error if you accidentally replace the `==` operator with `=`.

```
if (4 == payCode) { // good
    cout << "You get a bonus!" << endl;
}
```

```
if (4 = payCode) { // ERROR!
    cout << "You get a bonus!" << endl;
}
```

string literals

What is the difference between these two declarations?

```
const char* literal = "Hello";  
literal[1] = 'a';
```

```
char str[] = {"hello"};  
str[1] = 'a';
```

string literals

What is the difference between these two declarations?

```
const char* literal = "Hello";  
literal[1] = 'a'; // compile ERROR!
```

```
char str[] = {"hello"};  
str[1] = 'a'; // OK
```

string literals

What is the difference between these two declarations?

```
const char* literal = "Hello";  
char* p = const_cast<char*>(literal);  
p[1] = 'a';
```

```
char str[] = {"hello"};  
str[1] = 'a'; // OK
```

string literals

What is the difference between these two declarations?

```
const char* literal = "Hello"; // stored in read-only memory  
char* p = const_cast<char*>(literal);  
p[1] = 'a'; // undefined behavior
```

```
-----  
char str[] = {"hello"}; // stored in stack  
str[1] = 'a'; // OK
```

- Attempting to modify a string literal results in **undefined behavior**: they may be stored in **read-only storage** or combined with other string literals:

Scope Rules

The portion of a program where an identifier can be used is known as its **scope**.

- Block Scope

```
{  
    int a;  
}
```

- Global Namespace Scope - An identifier declared **outside** any *function* or *class* has *global namespace scope*.

- Function Scope

- Namespace Scope

```
namespace MyNamespace {  
    int a;  
}  
...  
MyNamespace::a = 5;
```

Scope Rules

```
#include <iostream>
using namespace std;

void useLocal(); // function prototype
void useStaticLocal(); // function prototype
void useGlobal(); // function prototype

int x{1}; // global variable

int main() {
    cout << "global x in main is " << x << endl;

    int x{5}; // local variable to main

    cout << "local x in main's outer scope is " << x << endl;

    { // block starts a new scope
        int x{7}; // hides both x in outer scope and global x

        cout << "local x in main's inner scope is " << x << endl;
        cout << "global x in global namespace scope is " << ::x << endl;
    }

    cout << "local x in main's outer scope is " << x << endl;

    useLocal(); // useLocal has local x
    useStaticLocal(); // useStaticLocal has static local x
    useGlobal(); // useGlobal uses global x
    useLocal(); // useLocal reinitializes its local x
    useStaticLocal(); // static local x retains its prior value
    useGlobal(); // global x also retains its prior value

    cout << "\nlocal x in main is " << x << endl;
}
```

Scope Rules

```
// useLocal reinitializes local variable x during each call
void useLocal() {
    int x{25}; // initialized each time useLocal is called

    cout << "\nlocal x is " << x << " on entering useLocal" << endl;
    ++x;
    cout << "local x is " << x << " on exiting useLocal" << endl;
}

// useStaticLocal initializes static local variable x only the
// first time the function is called; value of x is saved
// between calls to this function
void useStaticLocal() {
    static int x{50}; // initialized first time useStaticLocal is called

    cout << "\nlocal static x is " << x << " on entering useStaticLocal"
        << endl;
    ++x;
    cout << "local static x is " << x << " on exiting useStaticLocal"
        << endl;
}

// useGlobal modifies global variable x during each call
void useGlobal() {
    cout << "\nglobal x is " << x << " on entering useGlobal" << endl;
    x *= 10;
    cout << "global x is " << x << " on exiting useGlobal" << endl;
}
```


Scope Rules

```
global x in main is 1  
local x in main's outer scope is 5  
local x in main's inner scope is 7  
global x in global namespace scope is 1  
local x in main's outer scope is 5
```

```
local x is 25 on entering useLocal  
local x is 26 on exiting useLocal
```

```
local static x is 50 on entering useStaticLocal  
local static x is 51 on exiting useStaticLocal
```

```
global x is 1 on entering useGlobal  
global x is 10 on exiting useGlobal
```

```
local x is 25 on entering useLocal  
local x is 26 on exiting useLocal
```

```
local static x is 51 on entering useStaticLocal  
local static x is 52 on exiting useStaticLocal
```

```
global x is 10 on entering useGlobal  
global x is 100 on exiting useGlobal
```

```
local x in main is 5
```

Function Signature

The portion of a function prototype *that includes the **name of the function** and the **types of its arguments*** is called the **function signature** or simply the **signature**.

- The function's *return type* is **NOT part of the function signature**.

Signature includes the following underlined members:

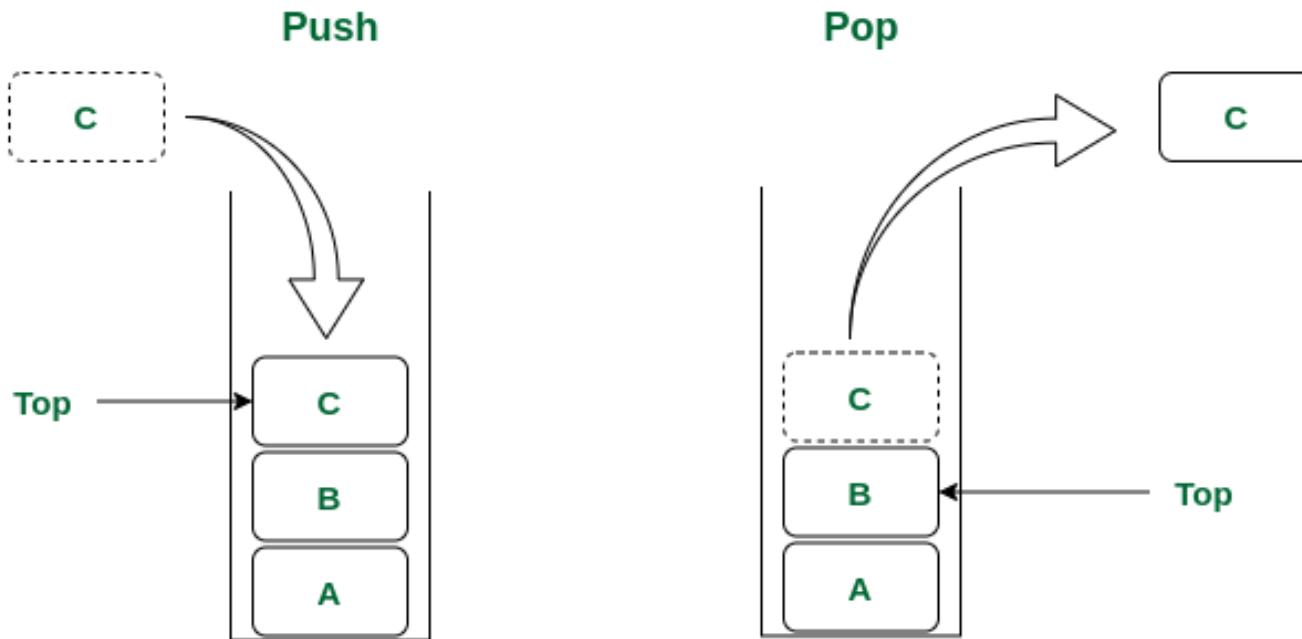
```
retType function(ParamType1 ParamName1, ParamType1 ParamName1)
```

```
...  
void maximum(int, int, int);  
int maximum(int, int, int);  
...
```

Compile Error!!! Functions that differ only in their return type cannot be overloaded.

Function Call Stack

Stacks are known as **last-in, first-out (LIFO)** data structures—the last item **pushed (inserted)** on the stack is the *first* item **popped (removed)** from the stack.



Stack Data Structure

Function Call Stack

Each time a function calls another function, an *entry* is **pushed** onto the **stack**. This entry is called a **stack frame** or an *activation record*.

Stack frame contains

- Return address to the calling function
 - Non-static Local variables.
- If a function is called, the stack frame for the function call is simply **pushed** onto the call stack.
- If the called function returns instead of calling another function before returning, the stack frame for the function call is **popped**, and *control transfers to the return address in the popped stack frame*.

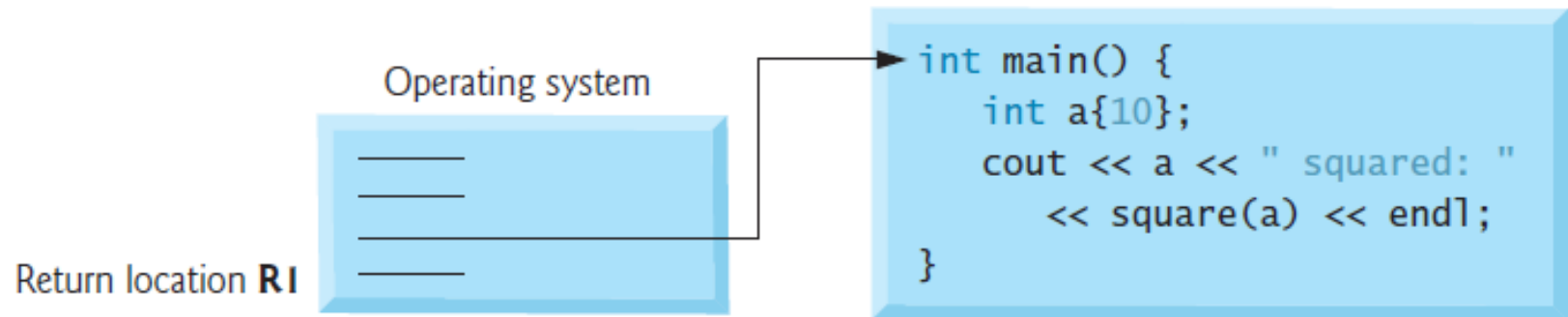
Function Call Stack

```
4  #include <iostream>
5  using namespace std;
6
7  int square(int); // prototype for function square
8
9  int main() {
10     int a{10}; // value to square (local variable in main)
11
12     cout << a << " squared: " << square(a) << endl; // display a squared
13 }
14
15 // returns the square of an integer
16 int square(int x) { // x is a local variable
17     return x * x; // calculate square and return result
18 }
```

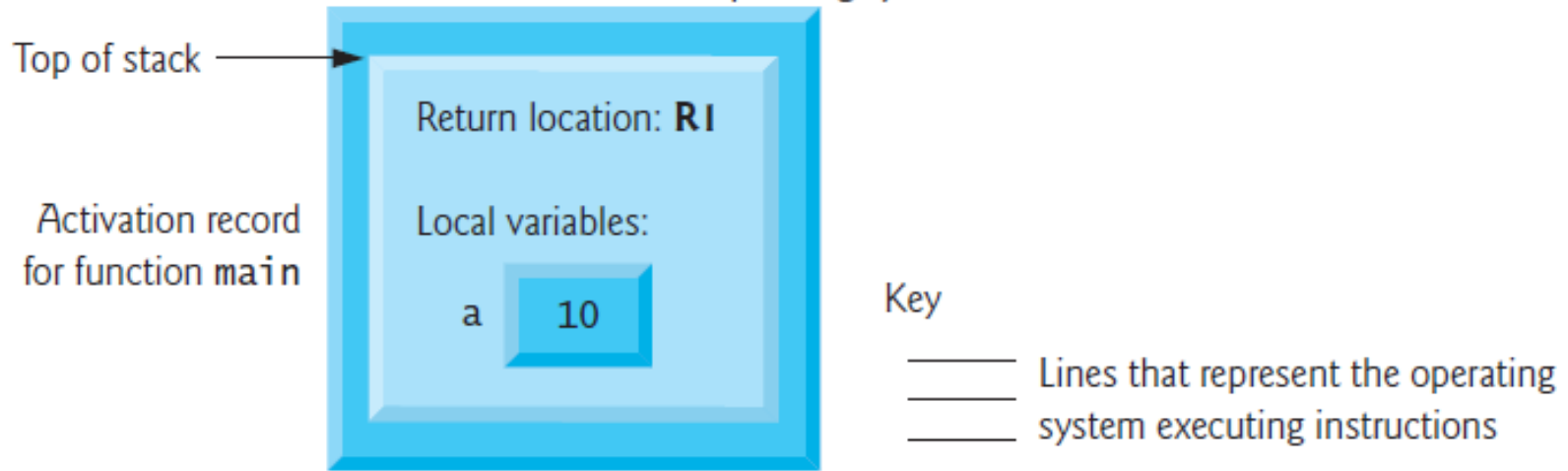
10 squared: 100

Function Call Stack

Step 1: Operating system calls `main` to execute application

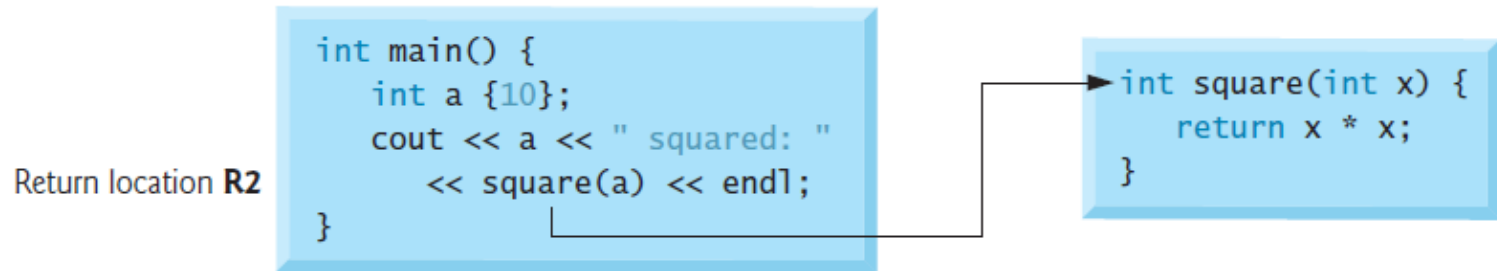


Function call stack after operating system calls `main`

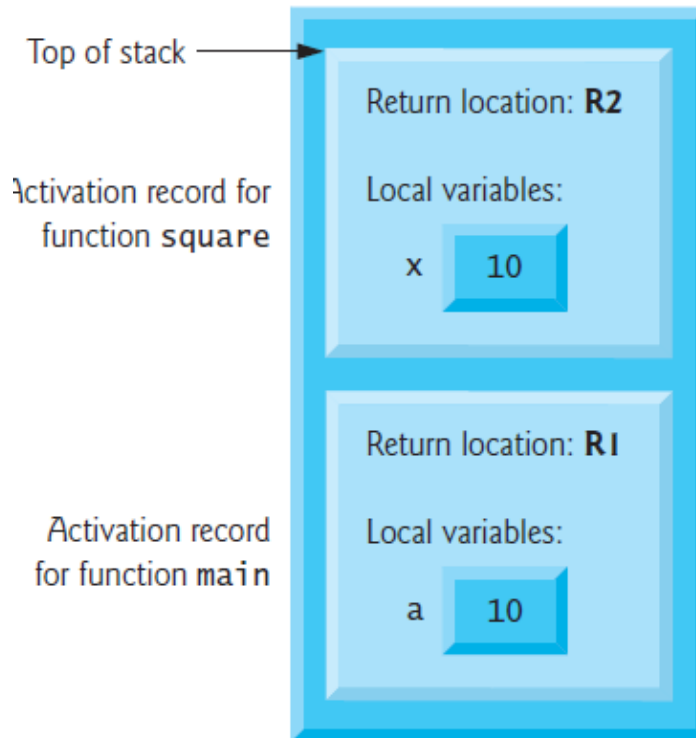


Function Call Stack

Step 2: `main` calls function `square` to perform calculation

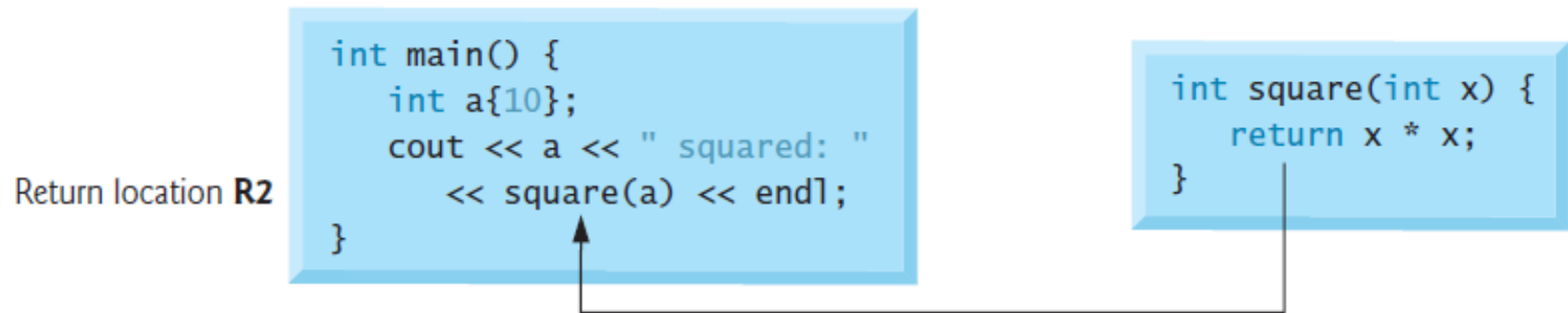


Function call stack after `main` calls `square`

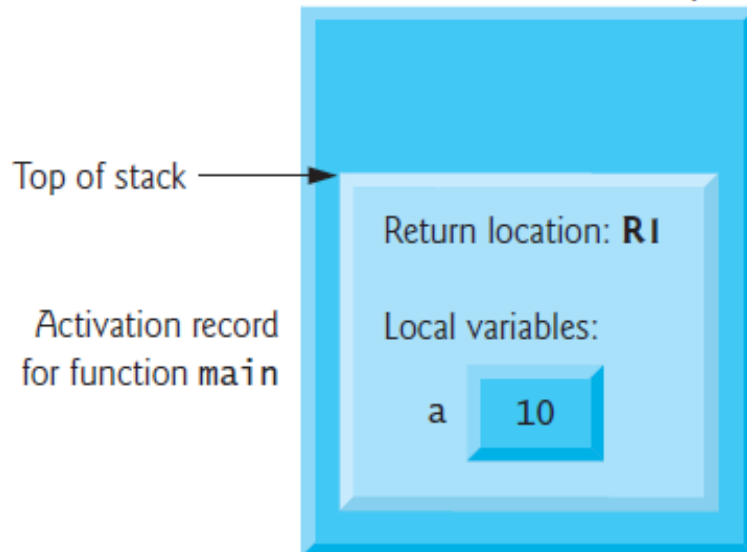


Function Call Stack

Step 3: `square` returns its result to `main`



Function call stack after `square` returns its result to `main`



Function Call Stack

What is **stack overflow**?

- *The amount of memory* in a computer is **finite**, so only a certain amount of memory can be used to store activation records on the function-call stack.
- If more function calls occur than can have their activation records stored on the function-call stack, a fatal error known as **stack overflow** occurs.
- ❖ This is how the website **stackoverflow.com** got its name. This is a great website for getting answers to your programming questions.

Inline Functions

- Function calls involve *execution-time overhead*.
 - C++ provides **inline functions** to help reduce function-call overhead.
 - Placing the qualifier inline before a function's return type in the function definition advises the compiler **to generate a copy of the function's body code in every place where the function is called** (when appropriate) to avoid a function call.
 - This often makes the program larger.
 - The compiler can ignore the inline qualifier and generally does so for all but the *smallest* functions.
 - Reusable inline functions are typically placed in *headers*, so that their definitions can be included in each source file that uses them.
-
- ❖ If you **change** the definition of an inline function, you must **recompile** all of that function's clients.
 - ❖ Compilers can inline code for which you have not explicitly used the inline keyword. Today's optimizing compilers are so sophisticated that it's best **to leave inlining decisions to the compiler**.

Inline Functions

```
3  #include <iostream>
4  using namespace std;
5
6  // Definition of inline function cube. Definition of function appears
7  // before function is called, so a function prototype is not required.
8  // First line of function definition acts as the prototype.
9  inline double cube(const double side) {
10     return side * side * side; // calculate cube
11 }
12
13 int main() {
14     double sideValue; // stores value entered by user
15     cout << "Enter the side length of your cube: ";
16     cin >> sideValue; // read value from user
17
18     // calculate cube of sideValue and display result
19     cout << "Volume of cube with side "
20          << sideValue << " is " << cube(sideValue) << endl;
21 }
```

```
Enter the side length of your cube: 3.5
Volume of cube with side 3.5 is 42.875
```

Inline Functions

```
9 inline double cube(const double side) {  
10     return side * side * side; // calculate cube  
11 }
```

Why double side is const?

- This tells the compiler that the function does not modify variable *side*.
- This ensures that side's value is *not* changed by the function during the calculation.

❖ *The **const** qualifier should be used to enforce the principle of **least privilege**. Using this principle to properly design software can greatly reduce debugging time and improper side effects and can make a program easier to modify and maintain.*

References and Reference Parameters

Two ways to pass arguments to functions in many programming languages are **pass-by-value** and **pass-by-reference**.

- When an argument is **passed by value**, a copy of the argument's value is made and passed (on the function-call stack) to the called function.
- With **pass-by-reference**, the caller gives the called function the ability **to access the caller's data directly, and to modify that data**.
- ❖ *One disadvantage of pass-by-value is that, if a large data item is being passed, copying that data can take a considerable amount of execution time and memory space.*
- ❖ *Pass-by-reference is good for performance reasons, because it can eliminate the pass-by-value overhead of copying large amounts of data.*
- ❖ *Pass-by-reference can weaken security; the called function can corrupt the caller's data.*

References and Reference Parameters

A **reference parameter** is an alias for its corresponding argument in a function call.

```
int a = 1;  
int& aRef = a; // 'aRef' is a reference to 'a'  
aRef++; // a=2, aRef=2
```

Passing Arguments by Value and by Reference (example)

```
3  #include <iostream>  
4  using namespace std;  
5  
6  int squareByValue(int); // function prototype (value pass)  
7  void squareByReference(int&); // function prototype (reference pass)  
8  
9  int main() {  
10     int x{2}; // value to square using squareByValue  
11     int z{4}; // value to square using squareByReference
```

References and Reference Parameters

```
12
13 // demonstrate squareByValue
14 cout << "x = " << x << " before squareByValue\n";
15 cout << "Value returned by squareByValue: "
16     << squareByValue(x) << endl;
17 cout << "x = " << x << " after squareByValue\n" << endl;
18
19 // demonstrate squareByReference
20 cout << "z = " << z << " before squareByReference" << endl;
21 squareByReference(z);
22 cout << "z = " << z << " after squareByReference" << endl;
23 }
24
25 // squareByValue multiplies number by itself, stores the
26 // result in number and returns the new value of number
27 int squareByValue(int number) {
28     return number *= number; // caller's argument not modified
29 }
30
31 // squareByReference multiplies numberRef by itself and stores the result
32 // in the variable to which numberRef refers in function main
33 void squareByReference(int& numberRef) {
34     numberRef *= numberRef; // caller's argument modified
35 }
```

```
x = 2 before squareByValue
Value returned by squareByValue: 4
x = 2 after squareByValue

z = 4 before squareByReference
z = 16 after squareByReference
```

References and Reference Parameters

Const References

- Const reference specifies that the reference is not be allowed to modify the corresponding argument.

What will you modify here (class set/get methods)?

```
void setName(std::string accountName);  
std::string getName() const;
```


References and Reference Parameters

Const References

- Const reference specifies that the reference is not be allowed to modify the corresponding argument.

What will you modify here?

```
void setName(std::string accountName);  
std::string getName() const;
```

```
void setName(const std::string& accountName);  
const std::string& getName() const;
```

References and Reference Parameters

- ❖ *When returning a reference to a **local variable**—unless that variable is declared **static**—the reference refers to a variable that’s discarded when the function terminates. An attempt to access such a variable yields undefined behavior. References to undefined variables are called **dangling references**.*

```
int& getLocalVarReference() {  
    int x = 5;  
    return x;  
}
```

```
int& danglingRef = getLocalVarReference();  
danglingRef++;  
std::cout << danglingRef << std::endl;
```

Result: Application crashed!

- ❖ *Returning a reference to a local variable in a called function is a logic error for which compilers typically issue a **warning**. Compilation warnings indicate potential problems, so most software-engineering teams have policies **requiring code to compile without warnings**.*

Default Arguments

- It's common for a program to invoke a function repeatedly with the same argument value for a particular parameter.
- In such cases, you can specify that such a parameter has a **default argument**, i.e., a default value to be passed to that parameter.
- When a program **omits** an argument for a *parameter with a default argument* in a function call, **the compiler rewrites the function call and inserts the default value of that argument**.
- Default arguments *must* be the rightmost (trailing) arguments in a function's parameter list.
- ❖ *Using default arguments can simplify writing function calls. However, some programmers feel that explicitly specifying all arguments is clearer.*

```

3  #include <iostream>
4  using namespace std;
5
6  // function prototype that specifies default arguments
7  unsigned int boxVolume(unsigned int length = 1, unsigned int width = 1,
8      unsigned int height = 1);
9
10 int main() {
11     // no arguments--use default values for all dimensions
12     cout << "The default box volume is: " << boxVolume();
13
14     // specify length; default width and height
15     cout << "\n\nThe volume of a box with length 10,\n"
16         << "width 1 and height 1 is: " << boxVolume(10);
17
18     // specify length and width; default height
19     cout << "\n\nThe volume of a box with length 10,\n"
20         << "width 5 and height 1 is: " << boxVolume(10, 5);
21
22     // specify all arguments
23     cout << "\n\nThe volume of a box with length 10,\n"
24         << "width 5 and height 2 is: " << boxVolume(10, 5, 2)
25         << endl;
26 }
27
28 // function boxVolume calculates the volume of a box
29 unsigned int boxVolume(unsigned int length, unsigned int width,
30     unsigned int height) {
31     return length * width * height;
32 }

```

The default box volume is: 1

The volume of a box with length 10,
width 1 and height 1 is: 10

The volume of a box with length 10,
width 5 and height 1 is: 50

The volume of a box with length 10,
width 5 and height 2 is: 100

Unary Scope Resolution Operator

C++ provides the **unary scope resolution operator (::)** to access a global variable when a local variable of the same name is in scope.

```
3  #include <iostream>
4  using namespace std;
5
6  int number{7}; // global variable named number
7
8  int main() {
9      double number{10.5}; // local variable named number
10
11     // display values of local and global variables
12     cout << "Local double value of number = " << number
13          << "\nGlobal int value of number = " << ::number << endl;
14 }
```

```
Local double value of number = 10.5
Global int value of number = 7
```

- ❖ **Always using** the unary scope resolution operator (::) to refer to global variables (even if there is no collision with a local-variable name) makes it clear that you're intending to access a global variable rather than a local variable.