

## Defining Haskell Values

- Functions definitions are written as equations
- $\text{double } x = x + x$   
 $\text{quadruple } x = \text{double } (\text{double } x)$
- compute the length of a list
  - length, if empty, is zero  
 $\text{length } [] = 0$
  - length if not empty is one plus the length of its tail  
 $\text{length } (x:xs) = 1 + \text{length } xs$

## Type Polymorphism

- What is the type of length?  
 $> \text{length } [1,2,3]$   
3  
 $> \text{length } ['a','b','c','d']$   
4  
 $> \text{length } [[],[1,2],[3,2,1],[],[6,7,8]]$   
5
- length works for lists of elements of arbitrary type  
 $\text{length} :: [a] \rightarrow \text{Int}$   
Here, 'a' denotes a type variable, so the above reads as "length takes a list of arbitrary type a and returns an Int"

## Laziness

- $\text{from } n = n : (\text{from } (n+1))$   
This recursive definition generates an infinite list of ascending numbers
- take n list - return first n elements of list
- take 10 (from 1)  
Haskell is a lazy language, so values are evaluated only when needed

## Program Compactness

- Sorting the empty list gives the empty list  
 $\text{qsort } [] = []$   
 $\text{qsort } (x:xs)$   
 $= \text{qsort } [y \mid y \leftarrow xs, y \leq x]$   
 $++ [x]$   
 $++ \text{qsort } [z \mid z \leftarrow xs, z > x]$

- Haskell list comprehensions

$[y \mid y \leftarrow xs, y \in x]$  "build list of ys, where y is drawn from xs such that y  $\in$  x"

### Patterns in Mathematics

We characterise things in maths by the laws they obey, laws which often look like patterns:

$$0! = 1$$

$$n! = n * (n - 1)!, n \geq 0$$

$$\text{len}(\langle \rangle) = 0$$

$$\text{len}(L1 \ L2) = \text{len}(L1) + \text{len}(L2)$$

$\langle \rangle$  denotes an empty list and  $L1L2$  is a list concatenation

### Factorial as Patterns

Maths:

$$0! = 1 \quad n! = n * (n - 1)!$$

Haskell (without patterns):

```
factorial_nop n = if n==0 then 1 else n * factorial_nop (n-1)
```

Haskell (with patterns):

$$\text{factorial } 0 = 1 \quad \text{factorial } n = n * \text{factorial } (n-1)$$

Formal argument 0 is shorthand saying check the argument to see if it is zero. If so, do the righthand

Formal argument n says take the argument and refer it to the righthand side as n

**Lists in Haskell** There is a standard approach to constructing lists:

- Empty list using  $[]$
- Given a term x and a list xs we can construct a list consisting of x followed by xs as follows:  $x:xs$
- So the list 1,2,3 can be built as  $1:2:3:[]$ . Brackets show how it is built up:  $1:(2:(3:[]))$
- We can write  $[1,2,3]$  as shorthand for the above list
- Lists can contain characters  $['H', 'E', 'L', 'L', 'O']$  but can be written shorthand as "HELLO".

and  $:$  are list constructors.  $":"$  is pronounced "cons"

### Length with Patterns

Math:

$$\text{len}(\langle \rangle) = 0$$

$$\text{len}(L1L2) = \text{len}(L1) + \text{len}(L2) \quad \text{len}(\langle \rangle) = 1 \quad \text{len}(\langle \rangle L) = 1 + \text{len}(L)$$

Haskell:

$$\text{mylength } [] = 0 \quad \text{mylength } (x:xs) = 1 + \text{mylength } xs$$

The key idea in pattern-matching is that the syntax used to build values can also be used to look at a value, determine how it was built and extract out the individual sub-parts if required.

### **Compact "Truth Tables"**

Patterns can be used to give an elegant expression to certain functions, for instances we can define a function over two boolean arguments like this:

```
myand True True = True myand _ _ = False
```

Here, the underscore pattern "\_" is a wildcard that matches anything

### **Types**

Haskell is strongly typed - every expression/value has a well-defined type:

```
myExpr :: MyType
```

"The value myExpr has type MyType"

Haskell supports type-inference. We don't have to declare types of functions in advance.

### **Atomic Types**

Some Atomic types builtin to Haskell

- () - the unit type has only one value, also written as ().
- Bool - boolean values, of which there are only two: true and false
- Char - character values, representing Unicode characters
- Int - fixed-precision integer type
- Integer - infinite-precision integer type
- Float - floating point number of precision at least that of IEEE single-precision
- Double - floating point of precision at least that of IEEE double-precision

### **Composite Types**

A type built on top of another type

- Functions have a type that indicates the type(s) of its input(s) and the type of its output
- Tuples gather values of other types together in a package
- Algebraic are data types that allow you to define types whose values can have more than one form

### **Tuples**

- We can create pairs, triples, n-tuples of values: (1,2) or (1, 'a', "Hi!")
- The type of the pair (42, 'z') where 42 has type Int, is (Int, Char)  
e.g. (42 'z') :: (Int,Char)

- We can use tuples as patterns in function definitions:

```
sumPair (a,b) = a + b
```

### **Algebraic Data Types**

Haskell lists are an example, with data that can have one of two forms:

Empty [] or non-empty (x:xs)

### **Function Types**

A function type consists of the input type, followed by a right-arrow and then the output type

```
myFun :: MyInputType -> MyOutputType
```

like f: A -> B in maths

```
sumInt :: [Integer] -> Integer
```

adds up a list of integers

Consider a rounding function that converts a floating point number to a fixed-width integer:

```
round :: Double -> Int
```

### **Inferring Function Types**

**FunDef** - Given a function declaration like  $f\ x = e$ , if  $e$  has type  $b$  and (the use of)  $x$  (in  $e$ ) has type  $a$ , then  $f$  must have type  $a \rightarrow b$ .

**FunUse** - Given a function application  $f\ v$ , if  $f$  has type  $a \rightarrow b$ , then  $v$  must have type  $a$ , and  $f\ v$  will have type  $b$ .