

Some operators are "nice"

Some operators have nice properties like having unit values e.g. $0 + a = a = a + 0$

We can code a simplifier for these as follows:

```
uopSimp cons u (Val v) e — v == u = e
uopSimp cons u e (Val v) — v == u = e
uopSimp cons u e1 e2 = cons e1 e2
```

We can deduce that cons is:

```
cons :: Expr -> Expr -> Expr
```

So we use Add and Mul directly:

```
simp (Add e1 e2) = uopSimp Add 0.0 e1 e2
simp (Mul e1 e2) = uopSimp Mul 1.0 e1 e2
```

The data constructors of Expr are in fact functions, whose types are as follows:

```
Val :: Double -> Expr
Var :: Id -> Expr
Add :: Expr -> Expr -> Expr
Mul :: Expr -> Expr -> Expr
Sub :: Expr -> Expr -> Expr
Div :: Expr -> Expr -> Expr
Def :: Id -> Expr -> Expr -> Expr
```

So, cons has to have type `Expr -> Expr -> Expr`, which is why Add and Mul are suitable arguments to pass into `uopSimp`

Given declaration:

```
data MyType = ... — MyCons T1 T2 ... Tn — ....
```

then data constructor MyCons is a function of type:

```
MyCons :: T1 -> T2 -> ... -> Tn -> MyType
```

Partial applications of MyCons are also valid

```
(MyCons x1 x2) :: T3 -> ... -> Tn -> MyType
```

Data constructors are the only functions that can occur in patterns.

Abstractions

A lot of the higher-order functions in the Prelude are examples of abstraction of common programming shapes encountered in functional programs (e.g. map and folds).

Turning Common Shapes Into Functions

Remember these?

```
sum [] = 0 sum (n:ns) = n + sum ns
length [] = 0 length (_,xs) = 1 + length xs
prod [] = 0 prod (n:ns) = n * prod ns
```

They have a common pattern which is typically referred to as folding. Can we abstract this?

Commonalities

- They all have the empty list as a base case

- They all have a non-empty list as the recursive case
- The base case returns a fixed "unit" value, which we will call **u**
 $\text{<abs-fold> } [] = u$
 $\text{<abs-fold> } (a:as) = \dots \text{<abs-fold> } as$