CSE3666

Homework 1 Solutions

1. Understanding instructions.

The answers are already provided.

Explanation:

We could convert everything to bits (not to a decimal number), find the result, and then convert bits to hex digits, but there are faster ways sometimes.

ADD: Do addition in hexadecimal.

AND: We convert each pair of hex digits to bits, AND, and then convert them back to hex digits. For special cases, we can do it quickly with hexadecimal digits. If a hex digit is 0, we know the resulting digit is 0. If a hex digit is F, we know the result is the other digit.

OR: Similar to AND. There are special cases. Note if a hex digit is F, the result will be F.

XOR: Similar to AND. There are also special cases, for example, when a hexadecimal digit is 0.

ADDI: similar to ADD.

ANDI: Convert -16 to hex first. -16 = (-1 << 4) = 0xFFFF FFF0. We then know the left 7 hex digits do not change. The least significant (right-most) hex digit is cleared.

SLLI: This is a special case. Shift hex digits to left by 3 positions.

SRAI: Pay attention to the sign bit because this is SRAI (not SRLI). This question is a special case. Shift hex digits to right by 2 positions. The highest bit is 1 and the instruction is SRAI. The sign is duplicated so the highest two hex digits are 0xFF.

2. Operations on bits.

```
      andi  t0, s1, 0x200         # separate bit 9
      beq   t0, x0, else          # if bit 9 is 0, goto else
      andi  s1, s1, 0x1FF         # set bits 32 to bit 9 to 0
      beq   x0, x0, endif         # skip the else branch
else:
      ori   s1, s1, 0xFFFFFE00    # note the immediate is sign extended
endif:
```

3. Hamming weight.

Solutions:

a. The loop is controlled by the mask. So the number of iterations of the loop does not depend on the bits in s0. The extracted bit decides if s1 is incremented or not. Therefore, the number of executed instructions depends on the number of 1's in s0, but does not depend on the location of 1's.

   The loop is executed for 32 times. In the loop, addi instruction is executed only if the extracted bit is 1. So the number of executed instruction in the loop is 32 * 4 + 16 = 144. Plus two instructions before the loop, the total number of executed instruction is 146.

b. The code is listed below.

```
      addi   s1, x0, 0          # s1 = 0
      add    t0, x0, s0         # make a copy so s0 is not changed
loop:
      bge    t0, x0, skip       # if the bit is 0, do not increment s1
      addi   s1, s1, 1          # increment the counter
skip:
      slli   t0, t0, 1          # shift bits to left by 1
      bne    t0, x0, loop       # if there is any 1 in t0, continue
```

   The loop terminates when t0 is 0. The number of iterations depends on the location of 1's. If t0 is 0 at the beginning, the loop is executed only once. If t0 is 1, the loop is executed 32 times.

   For s0 = 0xFF00FF00, the loop is executed 24 times. In each iteration, the ADDI instruction (in the loop) is executed only when the highest bit in t0 is 1. So the number of executed instruction in the loop is 24 * 3 + 16 = 88. Plus two instructions before the loop, the total number of executed instruction is 90.

   Can you figure out other ways to compute the Hamming weight?

4. Loop and if-else.

Solutions:

```
        addi    s2, x0, 0           # i = 0
        beq     x0, x0, condi

loopi:
        # Test the if condition
        andi    t0, s2, 0x91        # j = 0
        beq     t0, x0, else

        #------------ if branch
        slli    t1, s2, 8
        xor     s3, s3, t1
        beq     x0, x0, endif

        #------------ else branch
else:
        srai    t1, s2, 4
        add     s3, s3, t1

endif:
        addi    s2, s2, 1           # i ++
condi:
        blt     s2, s1, loopi       # i < a?
```