

52) Consider the following ANSI C code segment:

```
z = x + 3 + y → f1 + y → f2;
```

```
for (i = 0; i < 200; i = i + 2) {
```

```
    if (z > i) {
```

```
        p = p + x + 3;
```

```
        q = q + y → f1;
```

```
    } else {
```

```
        p = p + y → f2;
```

```
        q = q + x + 3;
```

```
    }  
}
```

Assume that variable  $y$  points to a struct (allocated on the heap) containing two fields  $f1$  and  $f2$  and the local variables  $x, y, z, p, q$  and  $i$  are allotted registers. Common sub-expression elimination (CSE) optimization is applied on the code. The number of addition and dereference operations (of the form  $y \rightarrow f1$  or  $y \rightarrow f2$ ) in the optimized code, respectively are:

a) 403 & 102

b) 203 & 2

c) 303 & 102

d) 303 & 2

Sol

$t1 = x + 3;$  → addition

$t2 = y \rightarrow f1;$  → dereference

$t3 = y \rightarrow f2;$  → dereference

$z = t1 + t2 + t3;$  → 2 additions

for ( ) {  
 }  
} ⇒ 2 additions, 100 times. and 2 dereference operations.

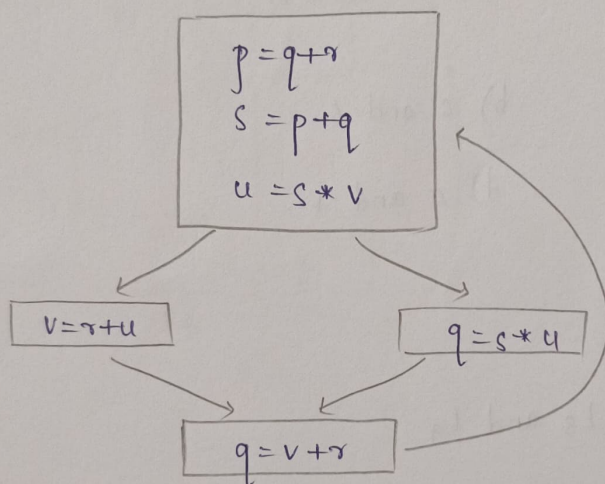
$$100 * 2 + 100 + 3 = 303$$



53) A variable  $x$  is said to be live at a statement  $S_i$  in a program if the following

1. three conditions hold simultaneously.
2. there exists a statement  $S_j$  that uses  $x$ .
3. there is a path from  $S_i$  to  $S_j$  in the flow graph corresponding to the program.

The path has no intervening assignment to  $x$  including at  $S_i$  &  $S_j$



The variables which are live both at the statement in basic block 2 and the statement in basic block 3 of above control flow graph are :

- a)  $p, s, u$       b)  $r, s, u$       ~~c)  $r, u$~~       d)  $q, v$

Sol Variables live at Basic block 2:  $\{r, u\}$

Variables live at Basic block 3:  $\{r, s, u\}$

Variable live at Both basic blocks 2 & 3:  $\{r, u\}$

54) Which of the following options correctly specify the number of basic blocks and the number of instructions in the largest basic



block, respectively? Consider the following pseudo-code

L<sub>1</sub>: t<sub>1</sub> = 1

L<sub>2</sub>: t<sub>2</sub> = 0

L<sub>3</sub>: t<sub>3</sub> = 0

L<sub>4</sub>: t<sub>4</sub> = 4 \* t<sub>3</sub>

L<sub>5</sub>: t<sub>5</sub> = 4 \* t<sub>2</sub>

L<sub>6</sub>: t<sub>6</sub> = t<sub>5</sub> \* M

L<sub>7</sub>: t<sub>7</sub> = t<sub>4</sub> + t<sub>6</sub>

L<sub>8</sub>: t<sub>8</sub> = a[t<sub>7</sub>]

L<sub>9</sub>: if t<sub>8</sub> ← max goto L<sub>11</sub>

L<sub>10</sub>: t<sub>1</sub> = t<sub>8</sub>

L<sub>11</sub>: t<sub>3</sub> = t<sub>3</sub> + 1

L<sub>12</sub>: if t<sub>3</sub> < M goto L<sub>4</sub>

L<sub>13</sub>: t<sub>2</sub> = t<sub>2</sub> + 1

L<sub>14</sub>: if t<sub>2</sub> < N goto L<sub>3</sub>

L<sub>15</sub>: max = t<sub>1</sub>

a) 7 and 7

b) 6 and 6

☒ c) 7 and 6

d) 6 and 7

Sol

Block 1 = L<sub>1</sub> and L<sub>2</sub>

Block 2 = L<sub>3</sub>

Block 3 = L<sub>4</sub>, L<sub>5</sub>, L<sub>6</sub>, L<sub>7</sub>, L<sub>8</sub> and L<sub>9</sub>

Block 4 = L<sub>10</sub>

Block 5 = L<sub>11</sub> and L<sub>12</sub>

Block 6 = L<sub>13</sub> and L<sub>14</sub>

Block 7 = L<sub>15</sub>

⇒ No of instructions in the largest basic block = 6.

55) Consider the control flow graph shown.

which one of the following choices list the set of live variable at the exit point of each basic block?

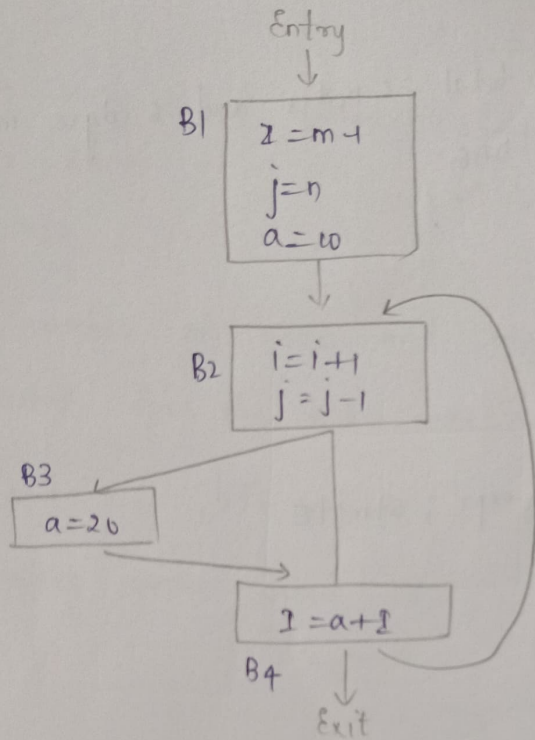
a) B<sub>1</sub>: { } , B<sub>2</sub>: {a} , B<sub>3</sub>: {a} , B<sub>4</sub>: {a}

b) B<sub>1</sub>: {i, j} , B<sub>2</sub>: {a} , B<sub>3</sub>: {a} , B<sub>4</sub>: {i}



c)  $B1: \{a, i, j\}$  ,  $B2: \{a, i, j\}$  ,  $B3: \{a, i\}$  ,  $B4: \{a\}$

d)  $B1: \{a, i, j\}$  ,  $B2: \{a, j\}$  ,  $B3: \{a, j\}$  ,  $B4: \{a, i, j\}$  .



Sol Option D

Live variables at the end of each basic block are:

$B1: \{a, i, j\}$

$B2: \{a, j\}$

$B3: \{a, j\}$

$B4: \{a, i, j\}$  .

56) Consider following C code segment:

$a = b + e;$

$e = a + i;$

$d = b + c;$

$f = d + l;$

$g = e + f;$

In a compiler, this code segment is represented internally as a directed acyclic graph (DAG).

The no of nodes in the DAG is 6

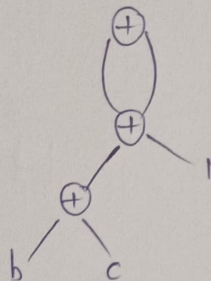


Sol

$$g = e + 1;$$

$$\Rightarrow (a+1) + (d+1) \Rightarrow ((b+c)+1) + ((b+c)+1)$$

Corresponding DAG is:



$\Rightarrow$  total 6 nodes and 6 edges in the DAG

51) Consider the following grammar:

stmt  $\rightarrow$  if expr then expr else expr ; stmt |  $\epsilon$

expr  $\rightarrow$  term relop term | term

term  $\rightarrow$  id | number

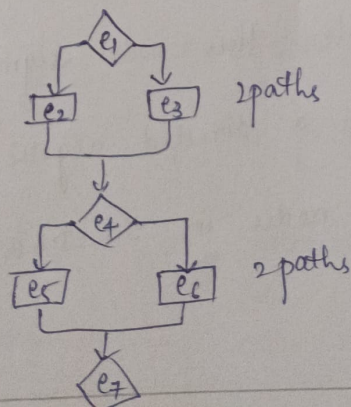
id  $\rightarrow$  a|b|c

number  $\rightarrow$  [0-9]

where relop is a relational operator (eg.  $<$ ,  $>$ ,  $=$ ),  $\epsilon$  refers to the empty statement and if, then, else are terminals. Consider a program P following the above grammar containing ten if terminals. The number of control flow paths in P is 1024. if  $e_1$  then  $e_2$  else  $e_3$  has 2 control flow paths,  $e_1 \rightarrow e_2$  and  $e_1 \rightarrow e_3$ .

Sol

if  $e_1$  then  $e_2$   
else  $e_3$



$$2 \times 2 \times \dots 10 \text{ times} = 2^{10}$$

$$= 1024 \text{ paths.}$$



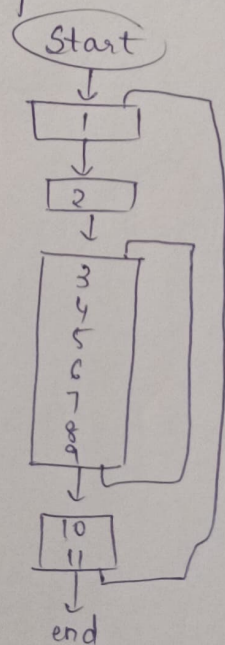
58) Consider the following intermediate code given below:

- |                    |                           |
|--------------------|---------------------------|
| 1. $i = 1$         | 7. $a[t_4] = 1$           |
| 2. $j = 1$         | 8. $j = j + 1$            |
| 3. $t_1 = 5 * i$   | 9. if $j \leq 5$ goto (3) |
| 4. $t_2 = t_1 + j$ | 10. $i = i + 1$           |
| 5. $t_3 = 4 * t_2$ | 11. if $i < 5$ goto (2)   |
| 6. $t_4 = t_3$     |                           |

The number of nodes and edges in the control-flow-graph constructed for the above code, respectively are

- a) 5 and 7    ~~b) 6 and 7~~    c) 5 and 5    d) 7 and 8

Sol Corresponding control graph is



Total, 6 nodes and 7 edges,  
So option B is correct.

59) In the context of abstract syntax tree (AST) and control-flow-graph (CFG) which one of the following is true?

- a) In both AST and CFG, let node  $N_2$  be the successor of node  $N_1$ . In the input program, the code corresponding to  $N_2$  is present



after the code corresponding to  $N_1$ .

b) For any input program, neither AST nor CFG will contain a cycle.

✓ c) The maximum number of successors of a node in an AST and a CFG depends on the input program.

d) Each node in AST and CFG corresponds to at most one statement in the input program.

Sol Option C is correct.

~~E~~ (a)  $\rightarrow$  false. If loop is present then control can move from any statement to any statement.

(b)  $\rightarrow$  false. CFG can have a cycle.

(c)  $\rightarrow$  true.

(d)  $\rightarrow$  false.