## CS 1511 Homework 17

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## 32.)

Let  $L \in BPL$ .

 $\exists$  TM T and  $\exists$  integer k such that:

 $\forall x \forall R \ \mathrm{T}(x,R)$  halts using log space

if  $x \in L$  then  $prob(T(x,R) \text{ accepts}) \ge 3/4$ 

if  $x \notin L$  then  $prob(T(x,R) \text{ accepts}) \le 1/4$ 

LOGSPACE  $\subset$  P. This is because with  $\log n$  space, there are  $2^{\log n}$  possible configurations for R. This simplifies to polynomial time.

Now, there exists some u, where u is a random assignment of Take an input x that has a length of n. Now we will have C be the number of configurations of TM T with input x. We will combine each configuration with another set of the same configurations to form a matrix. We will create this matrix such that each cell will have a 1/2 probability if the second configuration is reachable from the first in one step, and a probability of 0 otherwise. With this created, each cell  $W_t$  with configurations  $c_1$  and  $c_2$  is the probability of reaching configuration  $c_2$  from  $c_1$  in t steps, where  $W_t$  is the matrix created by multiplying W by itself t times. By scaling this up, we can compute the accepting probability of T(x, R) and decide if  $x \in L$ . With this we can see that each probability is a multiple of  $1/2^{poly(n)}$ , which we can represent with a poly number of digits. Therefore,  $L \in P$  and  $BPL \subseteq P$ .