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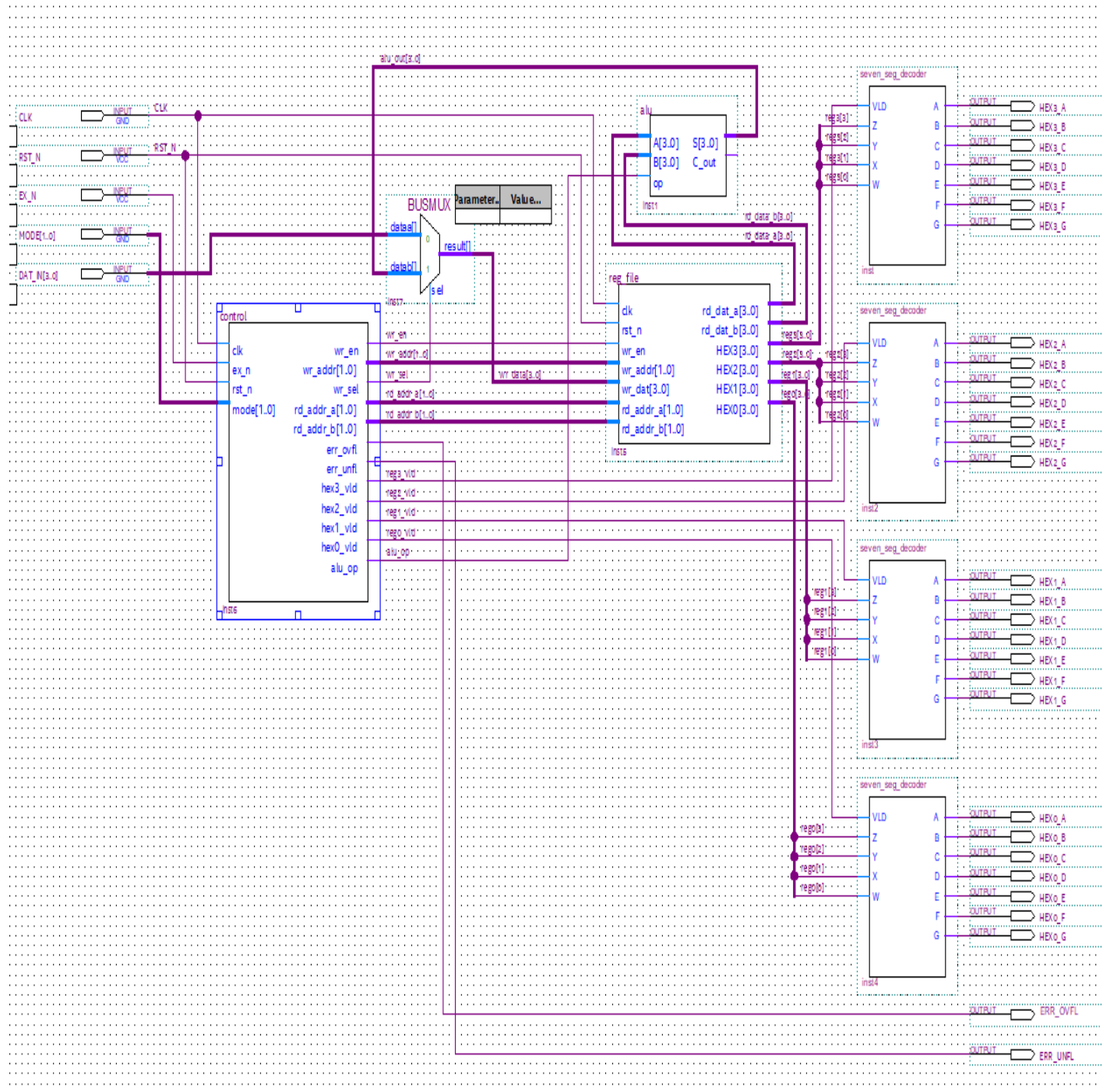
Final Project Report

The final project that I chose was the project on Stack Arithmetic (Project #2). The project report contains all the components that I used to put together this project, which includes, Verilog code for the Finite State Machine logic and the seven-segment display, schematic files for the ALU and Register file, and all the components used to put these top-level components together.

Basic Idea:

The basic idea for the implementation of this project was to use a Control Unit that would act like a Finite State Machine to evaluate the next state of the stack at any given instance, but also to evaluate other components essential to the stack operations such as the error overflow, the error underflow, and the validity of the 4 HEX panels used for the seven-segment display. With respect to loading and manipulating values on the stack, the idea was to use a 4-bit register file with 2 read ports and 1 write port that would not only load values inputted to it, but also send the read data to the Arithmetic Logic Unit (ALU), and the ALU carries out either the addition or subtraction operation, as selected by the user. Furthermore, a BUS MUX would be used to load values into the register either from the ALU or the user input, that is activated by a write select line coded in the control logic and would be sent to the write port of the register file. The register file then sends the data to the seven-segment display that displays the Stack onto the Altera Board.

Top Level Diagram:



Control Unit:

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q
Input				State	Output											
Current Stack Counter	Action	Description		Next Stack Counter	Write Enable	Write Address	Write Select	Read Address A	Read Address B	ALU Op	Error Overflow	Error Underflow	HEX3 VALID	HEX2 VALID	HEX1 VALID	HEX0 VALID
000	00	Push		001	1	00	0	XX	XX	X	0	0	0	0	0	0
001	00	Push		010	1	01	0	XX	XX	X	0	0	1	0	0	0
010	00	Push		011	1	10	0	XX	XX	X	0	0	1	1	0	0
011	00	Push		100	1	11	0	XX	XX	X	0	0	1	1	1	0
100	00	Push		100	0	XX	X	XX	XX	X	1	0	1	1	1	1
000	01	Pop		000	0	XX	X	XX	XX	X	0	1	0	0	0	0
001	01	Pop		000	0	XX	X	XX	XX	X	0	0	1	0	0	0
010	01	Pop		001	0	XX	X	XX	XX	X	0	0	1	1	0	0
011	01	Pop		010	0	XX	X	XX	XX	X	0	0	1	1	1	0
100	01	Pop		011	0	XX	X	XX	XX	X	0	0	1	1	1	1
000	10	Pop with Add		000	0	XX	X	XX	XX	X	0	1	0	0	0	0
001	10	Pop with Add		001	0	XX	X	XX	XX	X	0	1	1	0	0	0
010	10	Pop with Add		001	1	00	1	01	00	0	0	0	1	1	0	0
011	10	Pop with Add		010	1	01	1	10	01	0	0	0	1	1	1	0
100	10	Pop with Add		011	1	10	1	11	10	0	0	0	1	1	1	1
000	11	Pop with Subtract		000	0	XX	X	XX	XX	X	0	1	0	0	0	0
001	11	Pop with Subtract		001	0	XX	X	XX	XX	X	0	1	1	0	0	0
010	11	Pop with Subtract		001	1	00	1	01	00	1	0	0	1	1	0	0
011	11	Pop with Subtract		010	1	01	1	10	01	1	0	0	1	1	1	0
100	11	Pop with Subtract		011	1	10	1	11	10	1	0	0	1	1	1	1

The control unit takes care of the state of stack at any given instance, based on the 4 modes available to the user:

1. Push: Mode 00
2. Pop: Mode 01
3. Pop with Add: Mode 10
4. Pop with Subtract: Mode 11

Next State Counter: Next State Counter derives the next state expression for the stack (positional only) based on the mode of operation selected by the user. There are 5 different possibilities for the number of values on the stack, based on the stack depth of 4 provided in the project description.

K MAP:

Counter Next Counter[2] Next Counter[1] Next Counter[0]

Mode = 2'b00 Counter[1:0] Counter[1:0] Counter[1:0]

Counter[2]	00 01 11 10	00 01 11 10	00 01 11 10
0	0 0 1 0	0 0 1 0	0 0 1 0
1	1 X X X	1 0 X X X	1 0 X X X

Mode = 2'b01 Counter[1:0] Counter[1:0] Counter[1:0]

Counter[2]	00 01 11 10	00 01 11 10	00 01 11 10
0	0 0 0 0	0 0 0 1	0 0 0 1
1	0 X X X	1 1 X X X	1 1 X X X

Mode = 2'b10 or 2'b11 Counter[1:0] Counter[1:0] Counter[1:0]

Counter[2]	00 01 11 10	00 01 11 10	00 01 11 10
0	0 0 0 0	0 0 0 1	0 0 0 1
1	0 X X X	1 1 X X X	1 1 X X X

$next_counter[2] = (counter[2])$
 $next_counter[1] = (\sim mode[0] \& \sim mode[1]) (counter[1] \wedge counter[0]) \mid (mode[0] \& counter[1]) \mid (mode[1] \& counter[0])$
 $next_counter[0] = (counter[1] \wedge counter[0]) \mid (mode[0] \& counter[1]) \mid (mode[1] \& counter[0])$

Write Enable: The write enable activates the write port of the register file, and is only active if we are in Mode 00, 10, 11, and there is no error (overflow or underflow).

K MAP:

Write Enable

(1) Mode = 2'b00 Counter[1:0] Write Enable

Counter[2]	00 01 11 10	Write Enable
0	1 1 1 1	$(\sim Mode[0] \& \sim Mode[1]) \& \sim Counter[2]$
1	0 X X X	$(Mode[0] \& Counter[2]) \mid (Mode[1] \& Counter[1])$

(2) Mode = 2'b01 Counter[1:0]

Counter[2]	00 01 11 10
0	0 0 0 0
1	0 X X X

(3) Mode = 2'b10 or 2'b11 Counter[1:0]

Counter[2]	00 01 11 10
0	0 0 1 1
1	1 1 X X

Write Address: The write address complements the write enable, and is the 2 bit-address that writes to one of the 4 registers in the register file.

K MAP:

Write Address

(1) Mode = 2'b00

Counter [1:0]	
00	01 11 10
0	0 0 1 1
1	X X X X

Counter [2]

Counter [1:0]	
00	01 11 10
0	0 1 1 0
1	X X X X

Counter [2]

(2) Mode = 2'b01

Counter [1:0]	
00	01 11 10
0	X X X X
1	X X X X

Counter [2]

Counter [1:0]	
00	01 11 10
0	X X X X
1	X X X X

Counter [2]

(3) Mode = 2'b10 or 2'b11

Counter [1:0]	
00	01 11 10
0	X X 0 0
1	1 X X X

Counter [2]

Counter [1:0]	
00	01 11 10
0	X X 1 0
1	0 X X X

Counter [2]

Wt-addr[0] = Counter[0]
Wt-addr[1] = Counter[2] | (Mode[1] & Counter[1])

Write Select: The write select is the select line for the BUS MUX that activates the ALU (in Mode 10 and 11) to carry out some arithmetic operations, and read and write the data back into the register file.

K MAP:

Write Select

Mode = 2'b00

Counter [1:0]	
00	01 11 10
0	0 0 0 0
1	0 X X X

Counter [2]

Mode = 2'b01

Counter [1:0]	
00	01 11 10
0	X X X X
1	X X X X

Counter [2]

Mode = 2'b10 or 2'b11

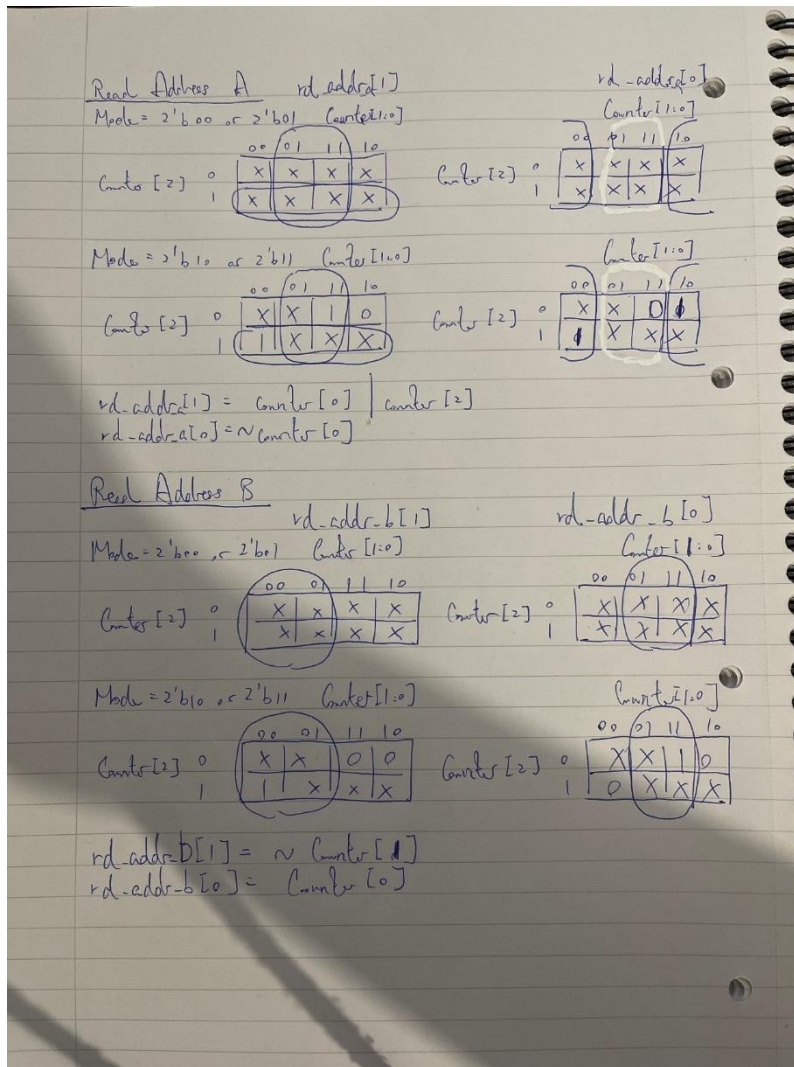
Counter [1:0]	
00	01 11 10
0	X X 1 1
1	1 X X X

Counter [2]

Wt_sel = Mode[1]

Read Address A and B: The read address is a 2-bit address that reads the data to one of the 4 registers in the register file. There are 2 read ports in the register file, that read data from the ALU, and the read address then feeds this data to the respective register in the register files and writes it back into the register file to load the data.

K MAP:



ALU Op: The ALU or the Arithmetic Logic Unit, carries out some arithmetic operations (add/subtract in this case), and feeds it back into the register file with the help of the BUS MUX. In case of subtraction, if the first popped value is less than the second popped value, the 2's complement of the operation is pushed onto the stack. The ALU op in the control unit provides the mode to the ALU to decide whether it must add/subtract.

K MAP:

ALU Op

Mode = 2'b00 Counter[1:0]

Counter[2]	0	00	01	11	10
1	X	X	X	X	X
	X	X	X	X	X

Mode = 2'b01 Counter[1:0]

Counter[2]	0	00	01	11	10
1	X	X	X	X	X
	X	X	X	X	X

Mode = 2'b10 Counter[1:0]

Counter[2]	0	00	01	11	10
1	X	X	0	0	0
	0	X	X	X	X

Mode = 2'b11 Counter[1:0]

Counter[2]	0	00	01	11	10
1	X	X	1	1	1
	1	X	X	X	X

ALU Op = Mode[0]

Error (Overflow and Underflow): The error overflow is the state in which the stack is full, and the user tries to push a value onto the stack, which results in an error displayed (LED illuminates).

The error underflow is the state in which:

1. The stack is empty and user tries to pop
 2. The stack has less than 2 values, and the user tries to pop with add/subtract
- This results in an error displayed (LED illuminates).

K MAP:

Error Overflow

Mode = 2'b00 Counter[1:0]

Counter[2]	0	00	01	11	10
1	0	0	0	0	0
	0	X	X	X	X

Mode = 2'b01 or 2'b10 or 2'b11

Counter[2]	0	00	01	11	10
1	0	X	X	X	X
	0	X	X	X	X

$Err_ovf = \sim mode[0] \& \sim mode[1] \& counter[2] \& \sim counter[1] \& \sim counter[0]$

Error Underflow

Mode = 2'b00 Counter[1:0]

Counter[2]	0	00	01	11	10
1	0	0	0	0	0
	0	X	X	X	X

Mode = 2'b01

Counter[2]	0	00	01	11	10
1	0	1	0	0	0
	0	X	X	X	X

Mode = 2'b10 or 2'b11

Counter[2]	0	00	01	11	10
1	0	1	1	0	0
	0	X	X	X	X

$Err_unf = (\sim mode[1] \& mode[0] \& \sim counter[2] \& \sim counter[1] \& \sim counter[0]) \mid (mode[1] \& \sim counter[2] \& \sim counter[1])$

Display HEX (HEX3, HEX2, HEX1, HEX0): The control unit calculates the validity of the each of the 3 HEX displays at any state of the stack. Based on all the possible operations that the user could carry out, the HEX display that was on was tabulated and the next state logic was derived. The validity of each HEX is inputted to the seven-segment display.

K MAP:

hex3-vld

	cnt[1:0]			
	00	01	11	10
cnt[2] 0	0	1	1	1
1	1	x	x	x

hex3-vld = cnt[0] | cnt[1] | cnt[2]

hex2-vld

	cnt[1:0]			
	00	01	11	10
cnt[2] 0	0	0	1	1
1	1	x	x	x

hex2-vld = cnt[1] | cnt[2]

hex1-vld

	cnt[1:0]				
	00	01	10	11	
cnt[2] 0	0	0	0	1	
1	1	x	x	x	

hex1-vld = cnt[0] cnt[1] | cnt[2]

hex0-vld

	cnt[1:0]				
	00	01	10	11	
cnt[2] 0	0	0	0	0	
1	1	x	x	x	

hex0-vld = cnt[2]

Control Unit Verilog Code:

```
1  module control(  
2      clk,  
3      rst_n,  
4      ex_n,  
5      mode,  
6      wr_en,  
7      err_ovfl,  
8      err_unfl,  
9      hex3_vld,  
10     hex2_vld,  
11     hex1_vld,  
12     hex0_vld,  
13     alu_op,  
14     wr_sel,  
15     rd_addr_a,  
16     rd_addr_b,  
17     wr_addr  
18 );  
19  
20  
21 input wire  clk;  
22 input wire  rst_n;  
23 input wire  ex_n;  
24 input wire  [1:0] mode;  
25 output wire wr_en;  
26 output wire err_ovfl;  
27 output wire err_unfl;  
28 output wire hex3_vld;  
29 output wire hex2_vld;  
30 output wire hex1_vld;  
31 output wire hex0_vld;  
32 output wire alu_op;  
33 output wire wr_sel;  
34 output wire [1:0] rd_addr_a;  
35 output wire [1:0] rd_addr_b;  
36 output wire [1:0] wr_addr;  
37  
38 wire  cnt_en;  
39 reg   [2:0] cnt;  
40 wire [2:0] nxt_cnt;  
41 reg   ex_dly;  
42  
43 assign  alu_op = mode[0];  
44  
45 assign  hex3_vld = cnt[0] | cnt[1] | cnt[2];  
46 assign  hex2_vld = cnt[1] | cnt[2];  
47 assign  hex1_vld = cnt[0] & cnt[1] | cnt[2];  
48 assign  hex0_vld = cnt[2];  
49
```

```

assign wr_en = cnt_en & (
    (~mode[0] & ~mode[1] & ~cnt[2]) |
    (mode[1] & (cnt[1] | cnt[2]))
);

assign wr_sel = mode[1];

assign wr_addr[1] = cnt[2] | (~mode[1] & cnt[1]);
assign wr_addr[0] = cnt[0];

assign rd_addr_a[1] = cnt[0] | cnt[2];
assign rd_addr_a[0] = ~cnt[0];

assign rd_addr_b[1] = ~cnt[1];
assign rd_addr_b[0] = cnt[0];

assign err_ovfl = ~mode[0] & ~mode[1] & cnt[2];
assign err_unfl = (mode[0] & ~mode[1] & ~cnt[2] & ~cnt[1] & ~cnt[0]) |
    (mode[1] & ~cnt[2] & ~cnt[1]);

// Falling Edge Detector
always @(posedge clk or negedge rst_n) begin
    if (rst_n == 1'b0) begin
        ex_dly <= 1'b1;
    end else begin
        ex_dly <= ex_n;
    end
end
assign cnt_en = ex_dly & ~ex_n;

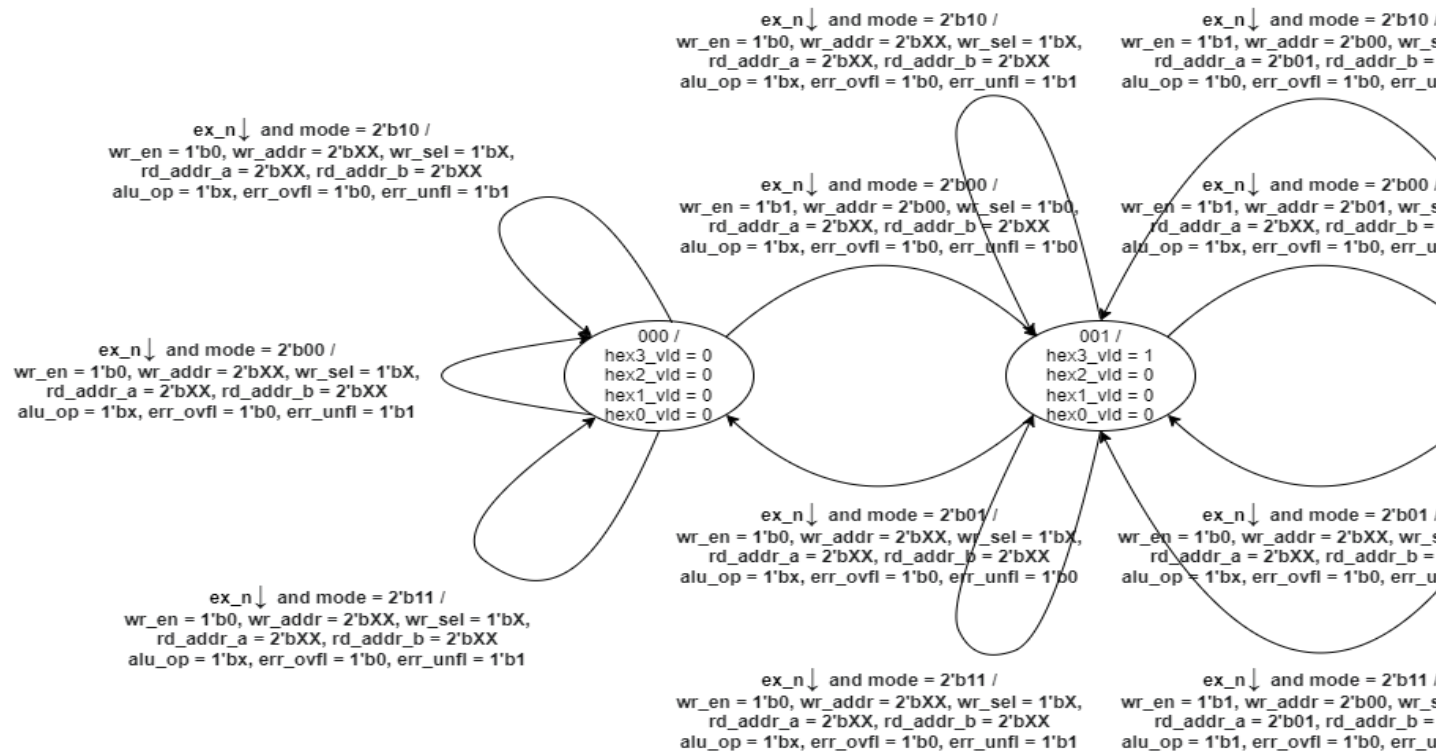
assign nxt_cnt[2] = (~mode[0] & ~mode[1] & (cnt[2] | (cnt[1] & cnt[0])));
assign nxt_cnt[1] = (~mode[0] & ~mode[1] & (cnt[0] ^ cnt[1])) |
    ((mode[0] | mode[1]) & (cnt[2] | (cnt[1] & cnt[0])));
assign nxt_cnt[0] = (~mode[0] & ~mode[1] & ~cnt[2] & ~cnt[0]) |
    ((mode[0] | mode[1]) & cnt[2]) |
    (mode[1] & ~cnt[1] & cnt[0]) |
    (cnt[1] & ~cnt[0]);

always @(posedge clk or negedge rst_n) begin
    if (rst_n == 1'b0) begin
        cnt <= 3'b000;
    end else if (cnt_en) begin
        cnt <= nxt_cnt;
    end
end

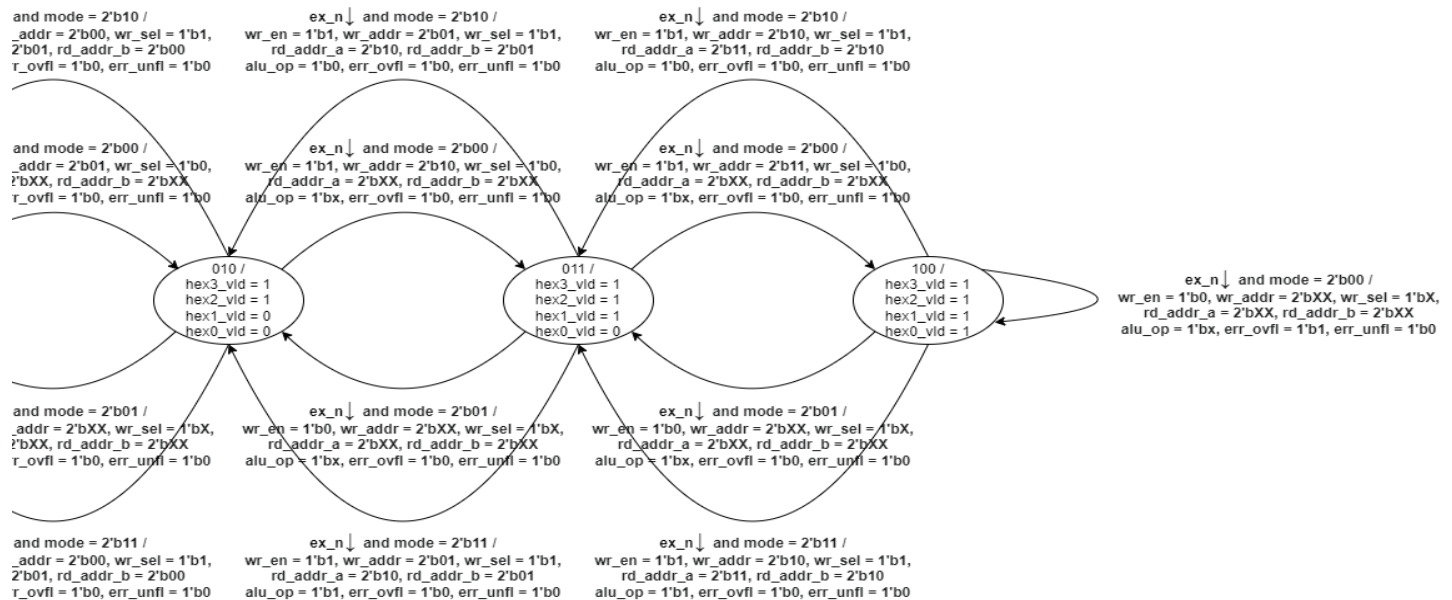
endmodule

```

First Half:



Second Half:

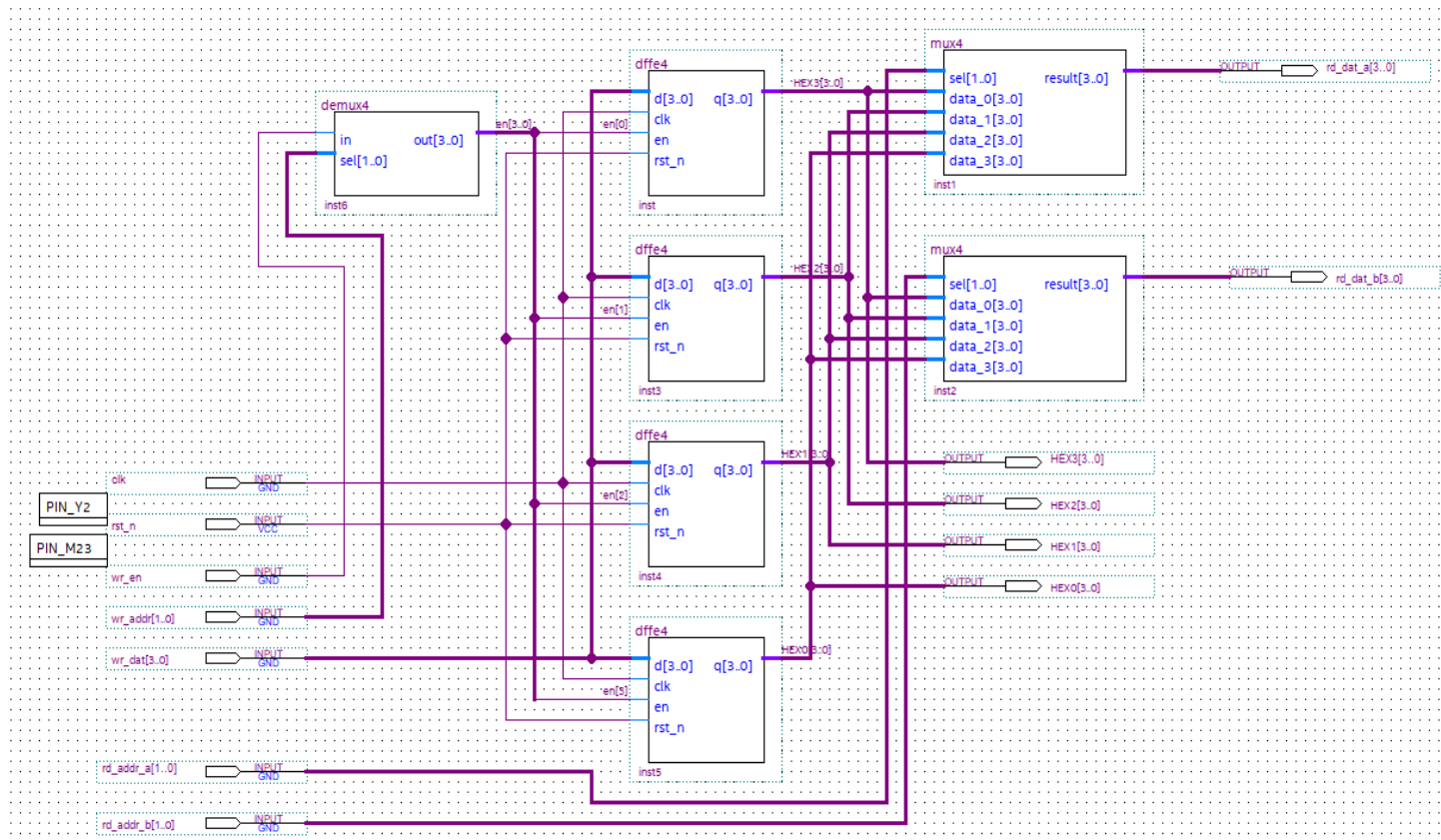


Register File:

The register file is a 4-bit 4 register, register file with 2 read ports and 1 write port. The register file consists of a write enable that enables one of the 4 registers based on their 2-bit write address. The register file has a reset signal and like the control logic, operates based on the positive edge of the system clock (50 MHz) or the negative edge of the reset signal.

Data is loaded onto the register file based on a 4-bit user input, or the arithmetic operations conducted by the ALU, with the help of a BUS MUX of width 4. The add/subtract value from the ALU is sent to the BUS MUX, and if the user selected a mode of operation that would activate the write select (mode 10 or 11), the write select may send the output of the ALU to the write port of the register file.

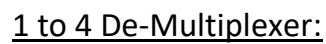
Block Diagram:



reg_file

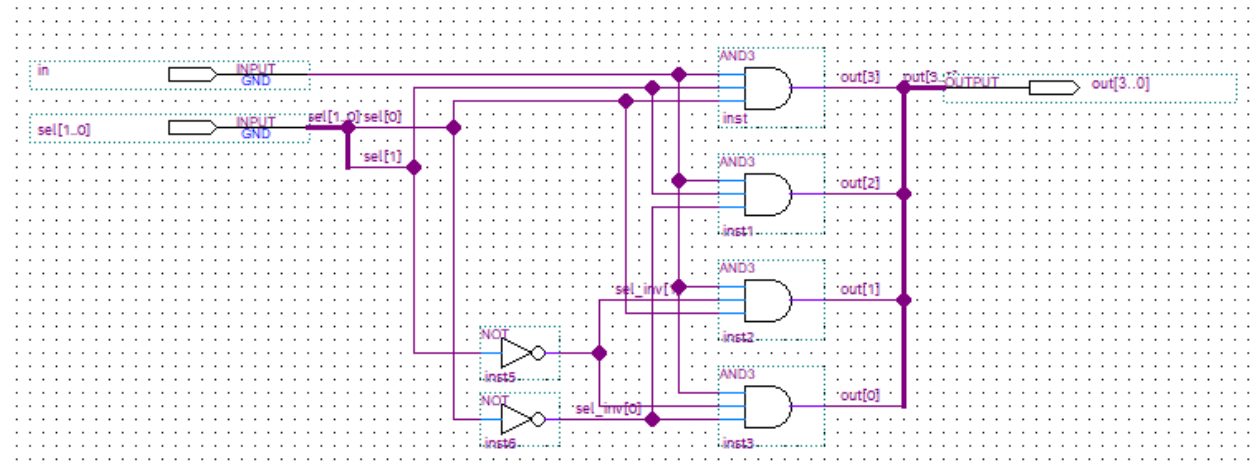
clk	rd_dat_a[3..0]
rst_n	rd_dat_b[3..0]
wr_en	HEX3[3..0]
wr_addr[1..0]	HEX2[3..0]
wr_dat[3..0]	HEX1[3..0]
rd_addr_a[1..0]	HEX0[3..0]
rd_addr_b[1..0]	

inst

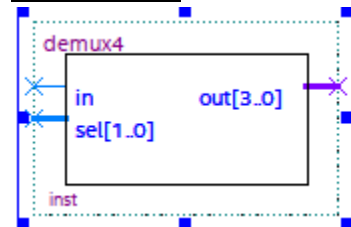


A 1 to 4 De-Multiplexer was used which takes the write enable as an input, and the 2-bit write address act as select lines s0 and s1. The output of the de-multiplexer act as the load for the 4-bit D-flip flops, each representing a single register.

Block Diagram:



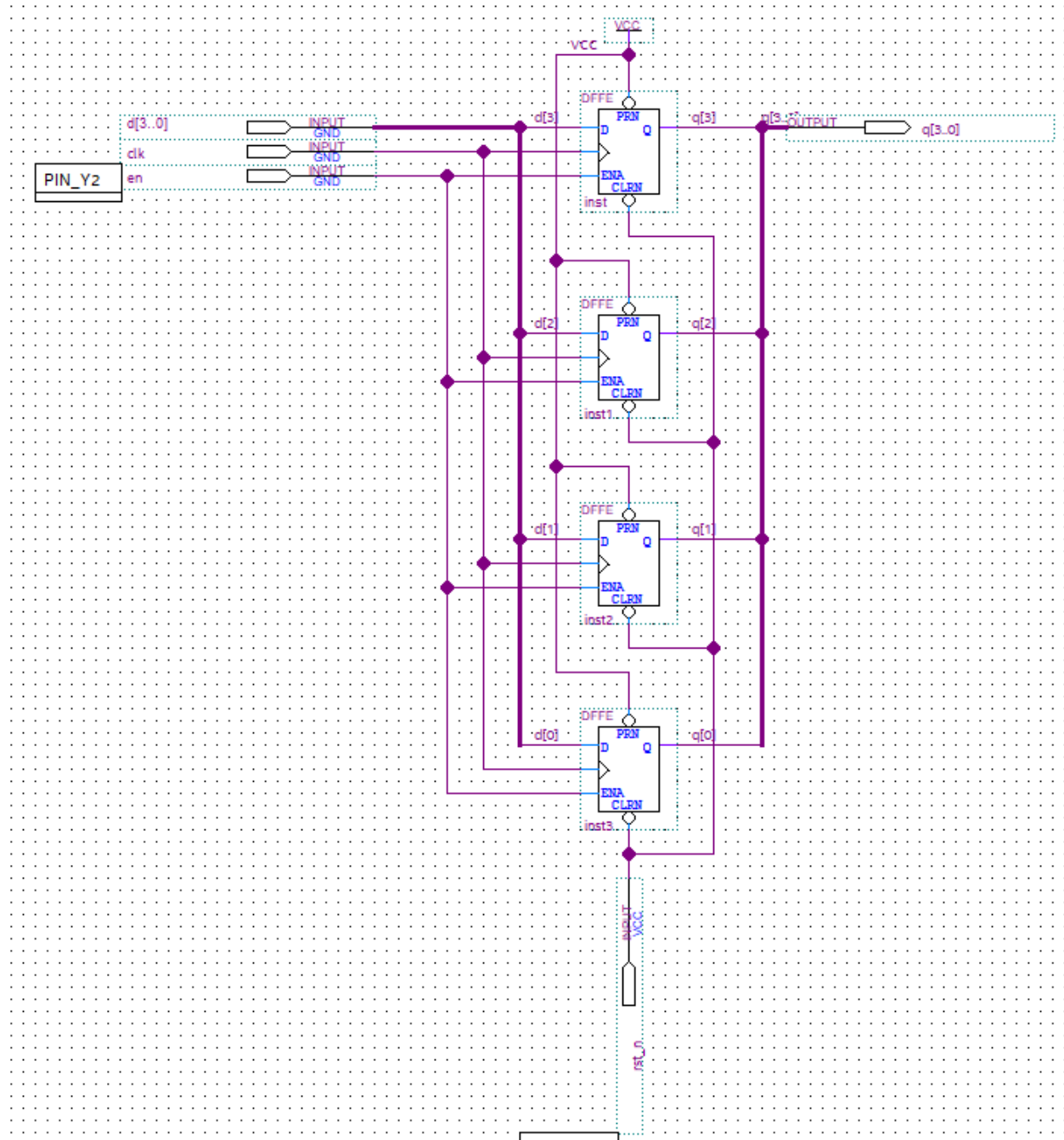
Symbol File:



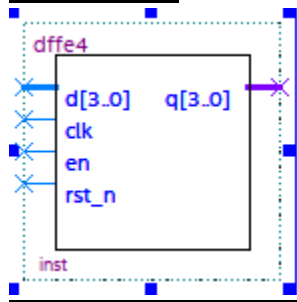
4-bit D-flip flop:

The 4-bit D-flip flop is an arrangement of 4 D-flip flops, each of which represents a single register. The D flip flops act as a memory limit for the register, and load the values onto the register based on the write address, at the positive edge or negative reset of the system clock (50 MHz). The D-flip flop also has a reset signal to clear the values on the registers.

Block Diagram:



Symbol File:



4 to 1 Multiplexer:

The two 4 to 1 Multiplexer's each represent a read port for the register file that has a 2-bit read address as the select lines, which reads the data to one of the 4 registers on the register file. The data from 4 registers is sent to the 4 to 1 MUX, which then outputs the read data based on the read address of the registers in the register file.

Verilog:

```
module mux4(
    data_0,
    data_1,
    data_2,
    data_3,
    sel,
    result
);

input wire [3:0] data_0;
input wire [3:0] data_1;
input wire [3:0] data_2;
input wire [3:0] data_3;
input wire [1:0] sel;
output wire [3:0] result;

assign result[0] = ( ~sel[1] & ~sel[0] & data_0[0] ) |
                  ( ~sel[1] & sel[0] & data_1[0] ) |
                  ( sel[1] & ~sel[0] & data_2[0] ) |
                  ( sel[1] & sel[0] & data_3[0] );

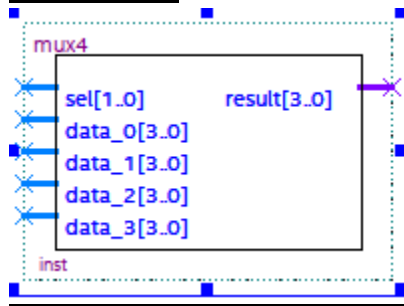
assign result[1] = ( ~sel[1] & ~sel[0] & data_0[1] ) |
                  ( ~sel[1] & sel[0] & data_1[1] ) |
                  ( sel[1] & ~sel[0] & data_2[1] ) |
                  ( sel[1] & sel[0] & data_3[1] );

assign result[2] = ( ~sel[1] & ~sel[0] & data_0[2] ) |
                  ( ~sel[1] & sel[0] & data_1[2] ) |
                  ( sel[1] & ~sel[0] & data_2[2] ) |
                  ( sel[1] & sel[0] & data_3[2] );

assign result[3] = ( ~sel[1] & ~sel[0] & data_0[3] ) |
                  ( ~sel[1] & sel[0] & data_1[3] ) |
                  ( sel[1] & ~sel[0] & data_2[3] ) |
                  ( sel[1] & sel[0] & data_3[3] );

endmodule
```

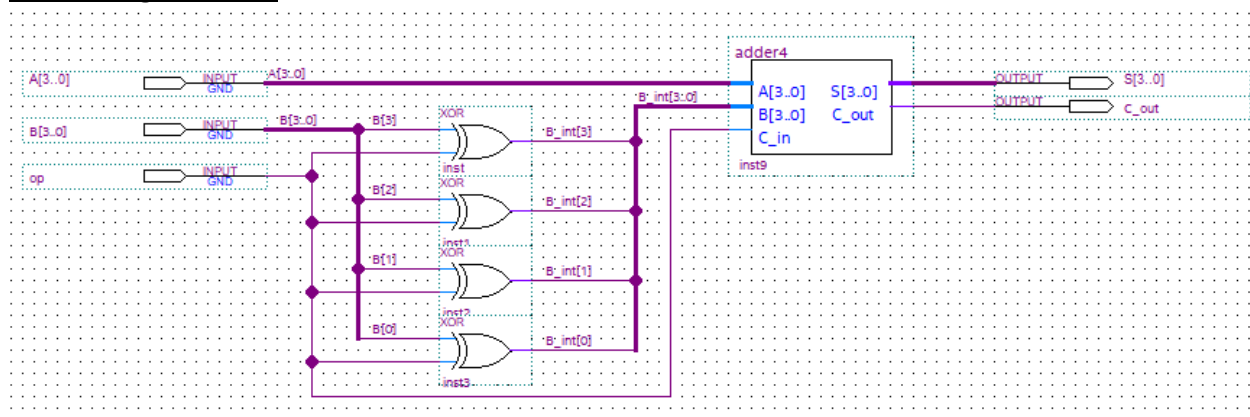
Symbol File:



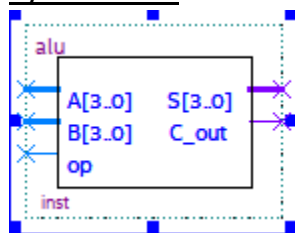
ALU (Arithmetic Logic Unit):

The ALU is used for the operation mode pop with add or pop with subtract, which adds or subtracts the top 2 popped values from the stack, and pushes it back onto the stack. The output of the ALU is sent to the BUS MUX that writes the data to the register file if the write select is active (= 1). The ALU gets the 2 popped values from the read data of the register file, that reads the data from the register corresponding to the read address, and sends it to the ALU to carry out the add/subtract operations. The ALU uses a 4-bit adder and XOR gates to carry out both the add and subtract operations.

Block Diagram File:



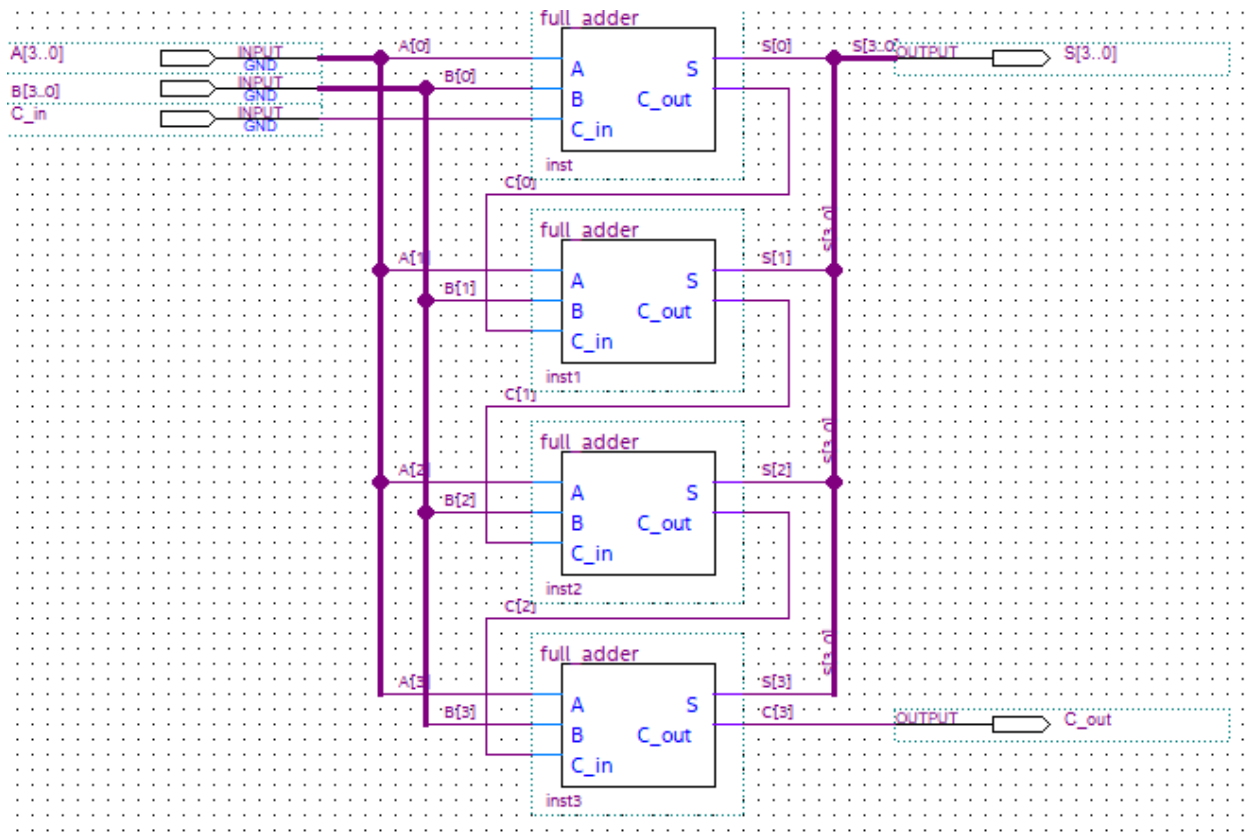
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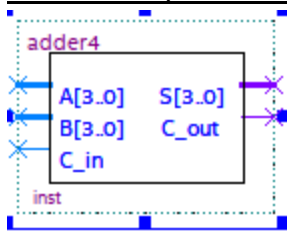
4-bit Adder/Subtractor:

The ALU implements a 4-bit adder circuit to carry out both addition and subtraction arithmetic operations based on the mode of operation (ALU_op) selected by the user. The 4-bit adder circuit implements 4 full adder circuits, each of which have two 4-bit inputs and a carry input which gets the value of ALU_op, and produce a 3-bit sum output and a carry output.

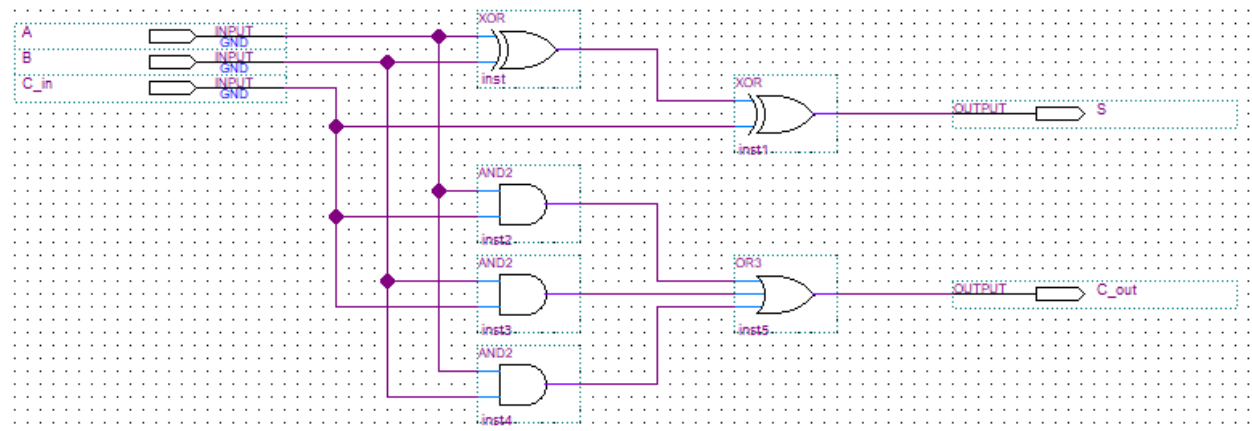
4-bit Adder/Subtractor Block Diagram:



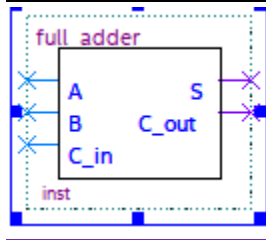
4-bit Adder/Subtractor Symbol File:



Full Adder Block Diagram:



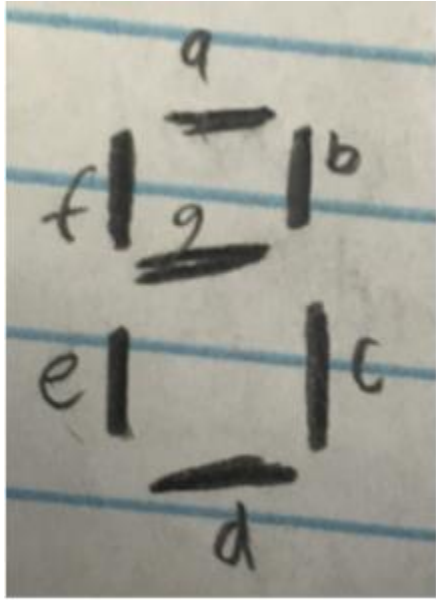
Full Adder Symbol File:



Seven-Segment Display:

The seven-segment display is a key component of the top-level diagram, as it displays the values on the stack. The seven-segment display considers seven different segments(a-g) of the HEX display and displays only the segments relevant to the output. Additionally, the design for my seven-segment display has an additional input VLD, which checks if that HEX is valid, and should be in use. The seven-segment display is activated by an enable, and takes 4 inputs Z, Y, X, W.

Seven-Segment Display segments:

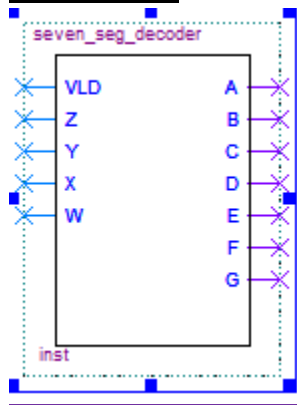


Verilog Code:

```
module seven_seg_decoder(VLD, Z, Y, X, W, A, B, C, D, E, F, G);
    input VLD, Z, Y, X, W;
    output reg A, B, C, D, E, F, G;

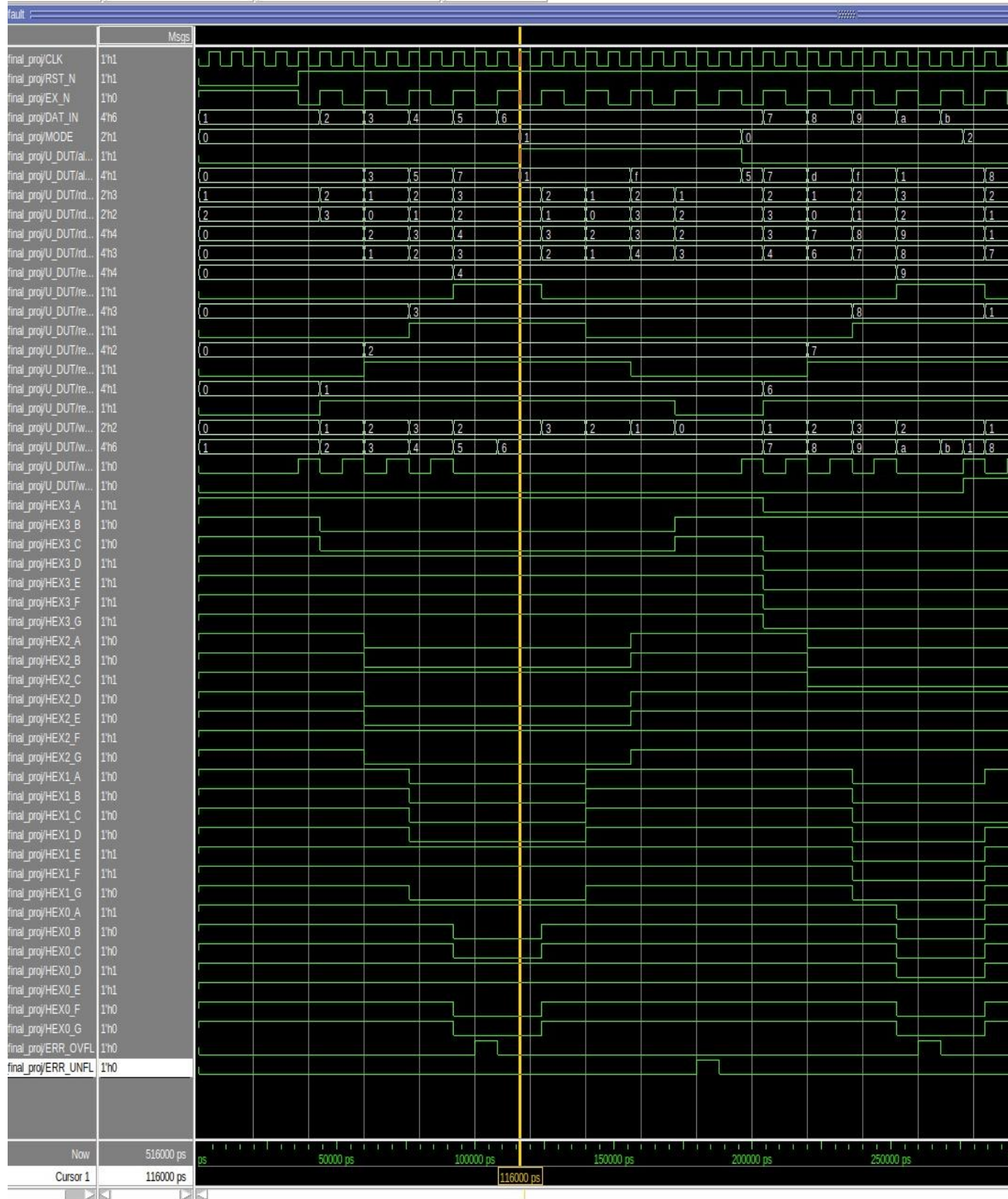
    always @(VLD or Z or Y or X or W)
    begin
        if (VLD == 1'b0) begin
            {A, B, C, D, E, F, G} = 7'b1111111;
        end else begin
            case({Z, Y, X, W})
                //truth tabling
                4'b0000: {A, B, C, D, E, F, G} = 7'b0000001;
                4'b0001: {A, B, C, D, E, F, G} = 7'b1001111;
                4'b0010: {A, B, C, D, E, F, G} = 7'b0010010;
                4'b0011: {A, B, C, D, E, F, G} = 7'b0000110;
                4'b0100: {A, B, C, D, E, F, G} = 7'b1001100;
                4'b0101: {A, B, C, D, E, F, G} = 7'b0100100;
                4'b0110: {A, B, C, D, E, F, G} = 7'b0100000;
                4'b0111: {A, B, C, D, E, F, G} = 7'b0001111;
                4'b1000: {A, B, C, D, E, F, G} = 7'b0000000;
                4'b1001: {A, B, C, D, E, F, G} = 7'b0000100;
                4'b1010: {A, B, C, D, E, F, G} = 7'b0001000;
                4'b1011: {A, B, C, D, E, F, G} = 7'b1100000;
                4'b1100: {A, B, C, D, E, F, G} = 7'b0110001;
                4'b1101: {A, B, C, D, E, F, G} = 7'b1000010;
                4'b1110: {A, B, C, D, E, F, G} = 7'b0110000;
                4'b1111: {A, B, C, D, E, F, G} = 7'b0111000;
            endcase
        end
    end
endmodule
```

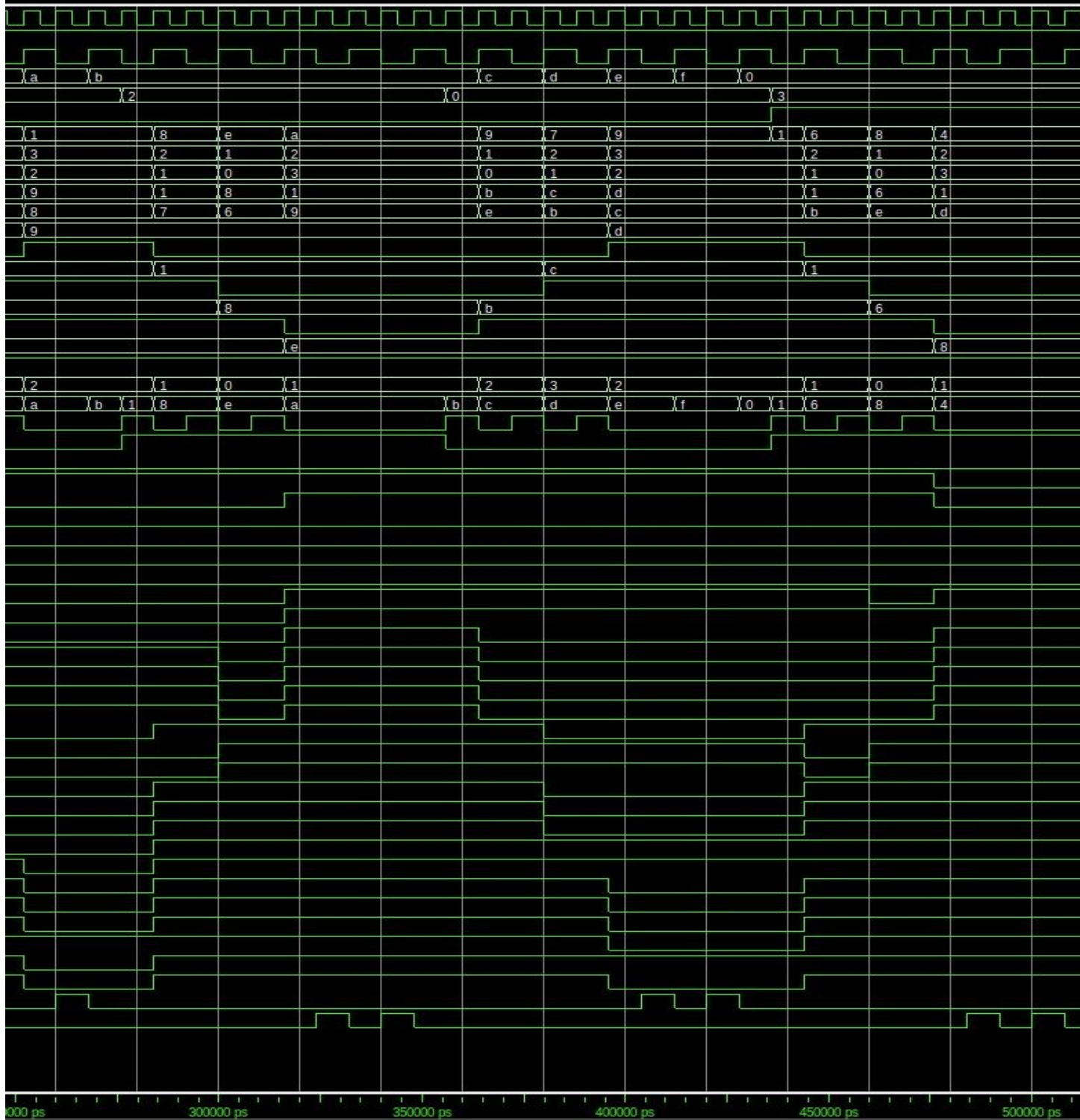
Symbol File:



ModelSim Tests:

I conducted ModelSim tests, for all the top-level components and I've attached the waveform for all the components of this final project below:





Pin Assignments:

Pin Assignment	Description	I/O Direction	Board Signal Name	Pin Assignment	Voltage						
CLK	System Clock	Input	CLOCK_50	PIN_Y2	3.3V		50 MHz	Oscillator			
RST_N	Global Reset (Active Low)	Input	KEY0	PIN_M23	3.3V		Push Button				
EX_N	Execute (Active Low)	Input	KEY1	PIN_M21	3.3V		Push Button				
MODE[1]	Operation Mode	Input	SW1	PIN_AC28	Depending on JP7	2.5V	Slide Switch				
MODE[0]		Input	SW0	PIN_AB28	Depending on JP7	2.5V	Slide Switch				
DAT_IN[3]	Input Data	Input	SW17	PIN_Y23	Depending on JP7	2.5V	Slide Switch				
DAT_IN[2]		Input	SW16	PIN_Y24	Depending on JP7	2.5V	Slide Switch				
DAT_IN[1]		Input	SW15	PIN_AA22	Depending on JP7	2.5V	Slide Switch				
DAT_IN[0]		Input	SW14	PIN_AA23	Depending on JP7	2.5V	Slide Switch				
ERR_OVFL	Overflow Error Indicator	Output	LEDR[0]	PIN_G19	2.5V		LED				
ERR_UNFL	Underflow Error Indicator	Output	LEDR[1]	PIN_F19	2.5V		LED				
HEX0_A		Output	HEX0[0]	PIN_G18	2.5V		7-Segment Displays				
HEX0_B		Output	HEX0[1]	PIN_F22	Depending on JP7	2.5V	7-Segment Displays				
HEX0_C		Output	HEX0[2]	PIN_E17	Depending on JP7	2.5V	7-Segment Displays				
HEX0_D		Output	HEX0[3]	PIN_L26	Depending on JP7	2.5V	7-Segment Displays				
HEX0_E		Output	HEX0[4]	PIN_L25	Depending on JP7	2.5V	7-Segment Displays				
HEX0_F		Output	HEX0[5]	PIN_J22	Depending on JP7	2.5V	7-Segment Displays				
HEX0_G		Output	HEX0[6]	PIN_H22	Depending on JP7	2.5V	7-Segment Displays				
HEX1_A		Output	HEX1[0]	PIN_M24	Depending on JP7	2.5V	7-Segment Displays				
HEX1_B		Output	HEX1[1]	PIN_Y22	Depending on JP7	2.5V	7-Segment Displays				
HEX1_C		Output	HEX1[2]	PIN_W21	Depending on JP7	2.5V	7-Segment Displays				
HEX1_D		Output	HEX1[3]	PIN_W22	Depending on JP7	2.5V	7-Segment Displays				
HEX1_E		Output	HEX1[4]	PIN_W25	Depending on JP7	2.5V	7-Segment Displays				
HEX1_F		Output	HEX1[5]	PIN_U23	Depending on JP7	2.5V	7-Segment Displays				
HEX1_G		Output	HEX1[6]	PIN_U24	Depending on JP7	2.5V	7-Segment Displays				
HEX2_A		Output	HEX2[0]	PIN_AA25	Depending on JP7	2.5V	7-Segment Displays				
HEX2_B		Output	HEX2[1]	PIN_AA26	Depending on JP7	2.5V	7-Segment Displays				
HEX2_C		Output	HEX2[2]	PIN_Y25	Depending on JP7	2.5V	7-Segment Displays				
HEX2_D		Output	HEX2[3]	PIN_W26	Depending on JP7	2.5V	7-Segment Displays				
HEX2_E		Output	HEX2[4]	PIN_Y26	Depending on JP7	2.5V	7-Segment Displays				
HEX2_F		Output	HEX2[5]	PIN_W27	Depending on JP7	2.5V	7-Segment Displays	Jumper			
HEX2_G		Output	HEX2[6]	PIN_W28	Depending on JP7	2.5V	7-Segment Displays	JP7	Short Pins 5 and 6	2.5V	Default
HEX3_A		Output	HEX3[0]	PIN_V21	Depending on JP7	2.5V	7-Segment Displays				
HEX3_B		Output	HEX3[1]	PIN_U21	Depending on JP7	2.5V	7-Segment Displays	Jumper			
HEX3_C		Output	HEX3[2]	PIN_AB20	Depending on JP6	3.3V	7-Segment Displays	JP6	Short Pins 5 and 6	2.5V	
HEX3_D		Output	HEX3[3]	PIN_AA21	Depending on JP6	3.3V	7-Segment Displays		Short Pins 7 and 8	3.3V	Default
HEX3_E		Output	HEX3[4]	PIN_AD24	Depending on JP6	3.3V	7-Segment Displays				
HEX3_F		Output	HEX3[5]	PIN_AF23	Depending on JP6	3.3V	7-Segment Displays				
HEX3_G		Output	HEX3[6]	PIN_Y19	Depending on JP6	3.3V	7-Segment Displays				